Def. Let $M$ be a TM and $w$ an input string. An accepting computation history for $M$ on $w$ is a sequence $C_1, C_2, \ldots, C_e$, where

1. $C_1$ is the starting configuration
2. $C_e$ is the accepting configuration
3. $C_i$ follows from $C_{i-1}$

- If $C_e$ is a rejecting configuration, then $C_1, \ldots, C_e$ is a rejecting history.

Q: What if $M$ doesn't halt on $w$?