## Erratum: The Approximability of the Exemplar Breakpoint Distance Problem

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Abstract. The paper "The Approximability of the Exemplar Breakpoint Distance Problem" [1], which appeared in AAIM 2006, contained several negative results and one positive result — a claimed  $O(\log n)$ factor greedy approximation for the One-sided Exemplar Breakpoint Distance Problem. Here, we show that the analysis was incorrect and the approximation factor of the greedy algorithm could be  $\Theta(n)$ , where n is the size of the alphabet.

In Section 5 of [1], a greedy algorithm is presented for the One-sided Exemplar Breakpoint Distance Problem. The claimed approximation factor is  $O(\log n)$ . We show that the factor could be  $\Theta(n)$  with an example. In our example, G is exemplar, so it satisfies the k-span condition. We start with a small n = 9.

 $G = \langle 1, 2, 3, 4, 5, 6, 7, 8, 9 \rangle$ , and  $H = \langle 9, 8, 7, 6, 1, 6, 2, 7, 3, 8, 4, 9, 5, 1, 2, 3, 4 \rangle$ .

The optimal solution is to have  $H^* = \langle \underline{6}, 7, 8, 9, 5, \underline{1}, 2, 3, 4 \rangle$ . In other words, we will have two breakpoints between G and  $H^*$ .

The greedy algorithm would first select the NB-interval in  $H: \langle 1, 2, 3, 4, 5 \rangle$ . So the greedy algorithm would have a solution  $H' = \langle 9, 8, 7, 6, 1, 2, 3, 4, 5 \rangle$ . In other words, we will have four breakpoints between G and H'. (We thank Minghui Jiang for the idea regarding this example.)

By generalizing the alphabet to be n = 2m + 1, i.e., |G| = 2m + 1 and |H| = 4m + 1, the greedy algorithm would generate *m* breakpoints while the optimal solution only introduces two breakpoints. So the approximation factor of the greedy algorithm is  $m/2 = \Theta(n)$ .

It is an open question whether the One-Sided Exemplar Breakpoint Distance Problem admits a polynomial time o(n)-factor approximation. The only known negative result is the APX-hardness of the problem.

## References

 Z. Chen, B. Fu and B. Zhu. The approximability of the exemplar breakpoint distance problem. In: Cheng, S.-W., Poon, C.K. (eds.) AAIM 2006. LNCS, vol. 4041, pp. 291-302. Springer, Heidelberg (2006)