CS418—-Operating Systems

Lecture 5

Memory Management—recent systems (cont.)

Textbook: Operating Systems
by William Stallings
1. Demand Paging (continued)

- How do we swap a page out of memory?
- FIFO (First In First Out).
  FIFO removes the page that has been in the memory the longest.

- LRU (Least Recently Used).
  LRU removes the page that shows the least sign of recent usage.

- MRU (Most Recently Used).
  MRU removes the page that shows the strongest sign of recent usage.
• LFU (Least Frequently Used). LFU removes the page that shows the least amount of recent usage, over certain period of time.

• How do we make use of the PMT (Page Map Table)?

• How to improve the performance of demand paging?
  – **Working set**: a set of pages in memory which do not need to be swapped out back and forth.
  – However, identifying working set is not easy.

• Summary
  – 1. Virtual memory is introduced.
  – 2. Utilizes memory more efficiently.
  – 3. Overhead is heavy.
2. Segmented Memory Allocation

- Both of the paging algorithms divide a job into physically equal-sized pages, which might cause serious problems in reality.

- The idea of segmented memory allocation algorithm is to divide job into logical segments.

- Memory is consequently divided into page frames with different sizes → external fragmentation reappears.

- For each job we associate it with a Segment Map Table (SMT).
• Similar to paging we need to maintain the following data structures: Job Table, Segment Map Table and Memory Map Table.
  
  – 1. Job Table lists every job in process.
  – 2. Segment Map Table lists details about each segment.
  – 3. Memory Map Table monitors the allocation of main memory.

• How to access a specific instruction? You still need to locate SEGMENT NUMBER and DISPLACEMENT.
3. Segmented/Demand Paged Memory Allocation

- IDEA: Divide each segment further into pages of equal size. Hence we need the following 4 data structures:
  - 1. Job Table lists every job in process.
  - 2. Segment Map Table (for each job) lists details about each segment.
  - 3. Page Map Table (for each segment) monitors the pages associated with each segment.
  - 4. Memory Map Table monitors the allocation of main memory.
Now we can move pages at will between main memory and second memory — **Virtual Memory**.

**Advantage of Virtual Memory.**

- 1. Job size has almost nothing to do with size of memory.
- 2. Memory is used more efficiently.
- 3. External fragmentation is eliminated and internal fragmentation is minimized.
- 4. Sharing of code/data is possible.
- 5. Dynamic linking of program segments is facilitated.

**Disadvantage of Virtual Memory.**

- 1. Hardware cost is increased.
- 2. Overhead (for paging interrupts) is increased significantly.
- 3. High cost for preventing thrashing.