CSCI 460—Operating Systems

Lecture 5

Memory Management–recent systems (cont.)

Textbook: Operating Systems
by William Stallings

1. Demand Paging (continued)

- How do we swap a page out of memory?
- FIFO (First In First Out).
  FIFO removes the page that has been in the memory the longest.

- LRU (Least Recently Used).
  LRU removes the page that shows the least sign of recent usage.

- MRU (Most Recently Used).
  MRU removes the page that shows the strongest sign of recent usage.

- LFU (Least Frequently Used).
  LFU removes the page that shows the least amount of recent usage, over certain period of time.

- How do we make use of the PMT (Page Map Table)?

- How to improve the performance of demand paging?
  - `Working set`: a set of pages in memory which do not need to be swapped out back and forth.
  - However, identifying working set is not easy.
- Summary
  - 1. Virtual memory is introduced.
  - 2. Utilizes memory more efficiently.
  - 3. Overhead is heavy.

2. Segmented Memory Allocation

- Both of the paging algorithms divide a job into physically equal-sized pages, which might cause serious problems in reality.

- The idea of segmented memory allocation algorithm is to divide job into logical segments.

- Memory is consequently divided into page frames with different sizes $\rightarrow$ external fragmentation reappears.

- For each job we associate it with a Segment Map Table (SMT).

- Similar to paging we need to maintain the following data structures: Job Table, Segment Map Table and Memory Map Table.
  - 1. Job Table lists every job in process.
  - 2. Segment Map Table lists details about each segment.
  - 3. Memory Map Table monitors the allocation of main memory.
- How to access a specific instruction? You still need to locate SEGMENT NUMBER and DISPLACEMENT.

3. Segmented/Demand Paged Memory Allocation

- IDEA: Divide each segment further into pages of equal size. Hence we need the following 4 data structures:
  - 1. Job Table lists every job in process.
  - 2. Segment Map Table (for each job) lists details about each segment.
  - 3. Page Map Table (for each segment) monitors the pages associated with each segment.
  - 4. Memory Map Table monitors the allocation of main memory.

- Now we can move pages at will between main memory and second memory — **Virtual Memory**.

- Advantage of Virtual Memory.
  - 1. Job size has almost nothing to do with size of memory.
  - 2. Memory is used more efficiently.
  - 3. External fragmentation is eliminated and internal fragmentation is minimized.
  - 4. Sharing of code/data is possible.
  - 5. Dynamic linking of program segments is facilitated.
- Disadvantage of Virtual Memory.
  - 1. Hardware cost is increased.
  - 2. Overhead (for paging interrupts) is increased significantly.
  - 3. High cost for preventing thrashing.
- **End of Memory Management.** Before we say that, how does UNIX memory management work?

4. Why buddy system?

- In real systems, a page is usually of size 2^a — a power of 2.
- A combination of dynamic partition and paging.
- Allocation algorithm
  - 1. Take a free block B_i.
  - 2. If the job requests more than 50% of B_i, allocate B_i to the job. Otherwise break B_i into two blocks B_{i1}, B_{i2} with equal size and proceed recursively on B_{i1}.
- Deallocation algorithm
  - 1. Take a block B_i to be released.
  - 2. If the buddy of B_i (i.e., shares the same parent and has the same size as B_i), B_j, is free, then combine B_i, B_j into a free block B_k with twice of the size.
  - 3. Recurse on B_k.