

# CS418—Operating Systems

## Lecture 6

### Memory Management

Textbook: Operating Systems  
by William Stallings

## 1. Why buddy system?

- In real systems, a page is usually of size  $2^a$  — a power of 2.
- A combination of dynamic partition and paging.
- Allocation algorithm
  - 1. Take a free block  $B_i$ .
  - 2. If the job requests more than 50% of  $B_i$ , allocate  $B_i$  to the job. Otherwise break  $B_i$  into two blocks  $B_{i1}, B_{i2}$  with equal size and proceed recursively on  $B_{i1}$ .
- Deallocation algorithm
  - 1. Take a block  $B_i$  to be released.
  - 2. If the buddy of  $B_i$  (i.e., adjacent to  $B_i$  and has the same size as  $B_i$ ),  $B_j$ , is free, then combine  $B_i, B_j$  into a free block  $B_k$  with twice of the size.
  - 3. Recurse on  $B_k$ .

## 2. Hash Tables

- How to put  $N$  items in a table of size  $M$  ( $N \leq M$ )?
  - 1. Sequential search —  $O(M)$  time in the worst case.
  - 2. Associative search —  $O(1)$  time, but needs special hardware.
  - 3. Binary search —  $O(\log M)$  time, but adding new items is hard.
- Hashing: An element with key  $k$  is mapped to slot  $h(k)$  in the table.
- How do we handle the situation when two elements are mapped to the same slot (collision)?
  - 1. Using an overflow area.
  - 2. Without using an overflow area.

### 3. Hashing with chaining

- Insert: If collision occurs, insert the element in an overflow area.
- Search: If slot  $h(k)$  is empty, report that no element exists. Otherwise, search slot  $h(k)$  then possibly the corresponding slots in the overflow area.
- Deletion: Just mark 'NIL'.
- Can you think of some good hash function?

## 4. Hashing without overflow area

- Insert: To perform insertion using this method, we successively examine (or probe) the hash table until we find an empty slot for the element.
- Search: Similar to Insert.
- Deletion: You can't just mark 'NIL', mark 'DELETED' instead. Why?
- Linear hashing
  
- Double hashing
  
- **End of Memory Management.** □