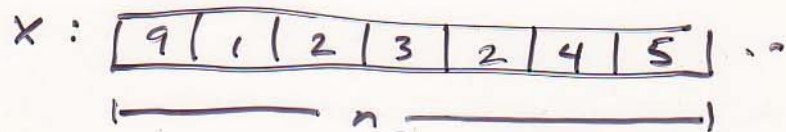


CS 223 Lec 4

Statistics on Arrays



e.g. find mean? ✓

min?

max?

median?

mode?

Finding Minimum

⇒ $n-1$ comparisons

Finding Maximum

⇒ $n-1$ comparisons

Find Both the Min + Max

x_1	x_2	x_3	...	x_{10}
6	1	5	3	9 .. 2

$$x_1 \leq x_5 ?$$

⋮

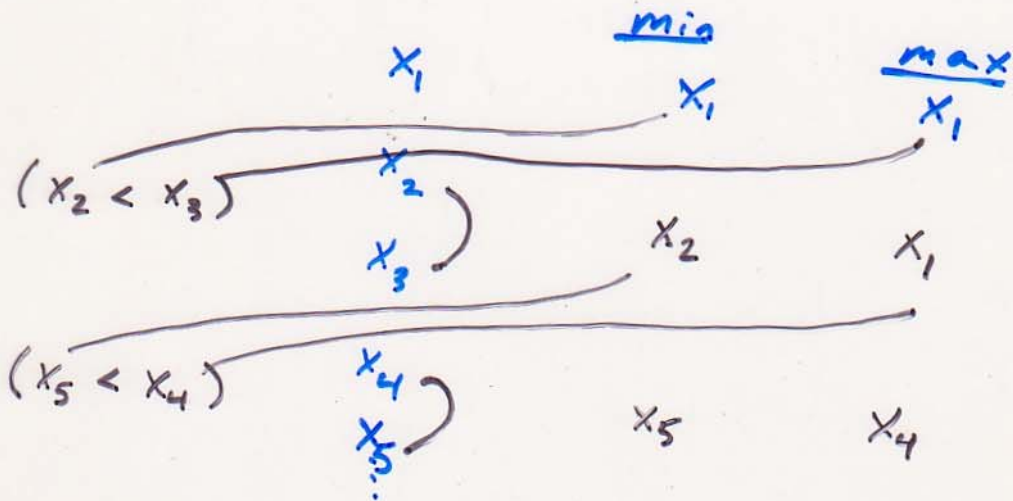
x_2 is largest

x_5 is smallest

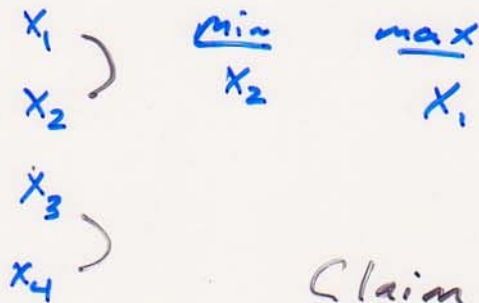
17, 15, 13, 14, 13

Algorithm for finding

Min / Max ...



x_n
 (n odd vs. even)

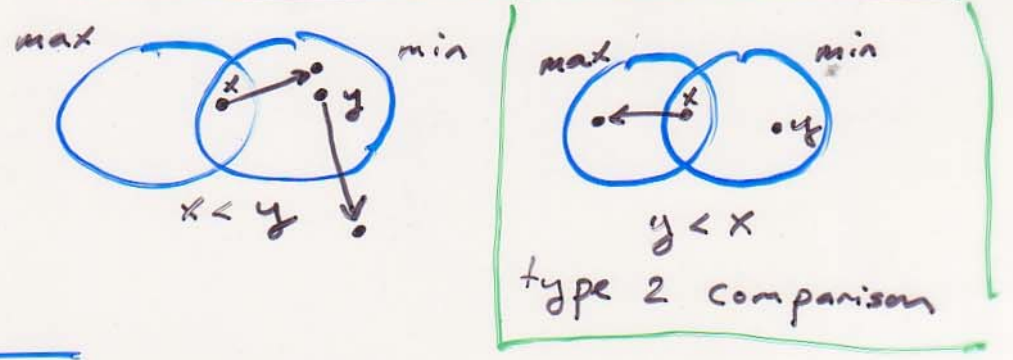
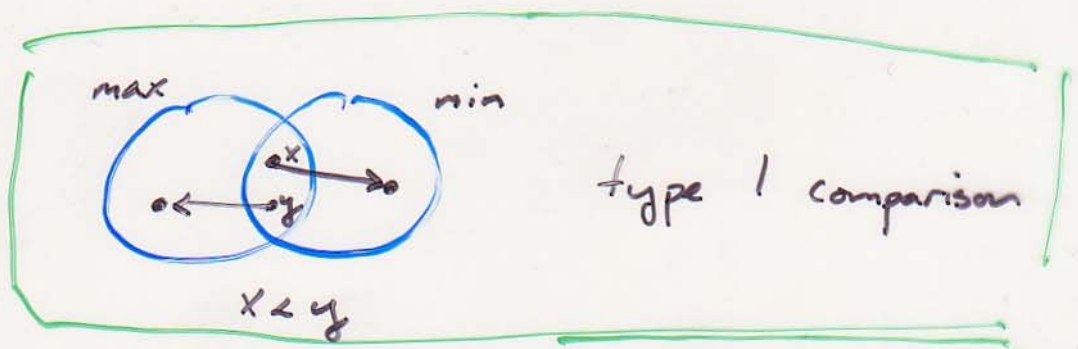
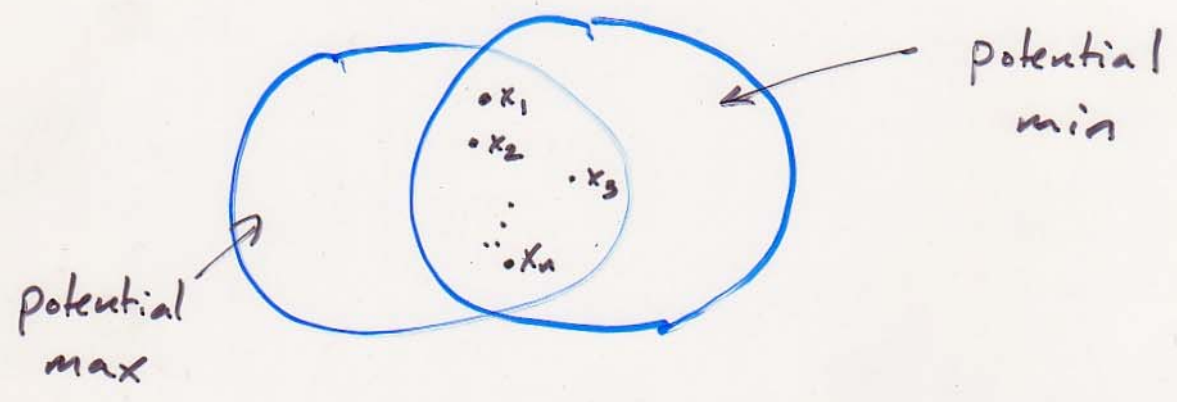


Claim $\lfloor \frac{3n}{2} \rfloor$
 comparisons.

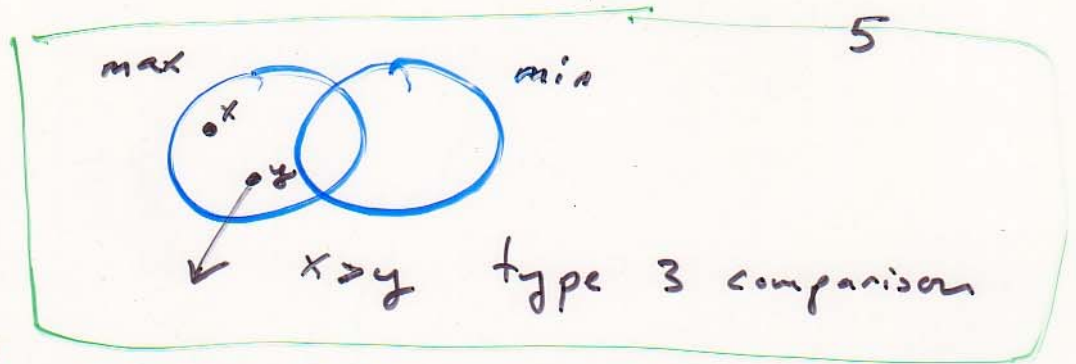
Finding a Lower Bound on
the # of comparisons
needed identify max/min
in an unsorted array



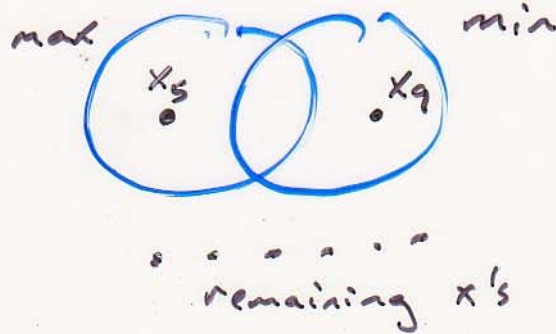
at start:



can
force
only
comparisons



at end:



of comparisons

$$\geq \lceil \frac{n}{2} \rceil + n - 2$$

$$= \lceil \frac{3n}{2} \rceil - 2$$

The Selection Problem

given: an unsorted array

A 1 | 2 | 5 | 4 | 3 | 6 | ...

order statistics

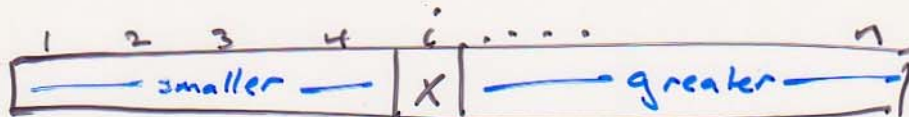
Defn The i th order statistic for A is the i th smallest value in A (x)

in other words there are

$$i-1 \text{ values } \leq x$$

$$n-i \text{ values } \geq x$$

Sort A ...



SELECTION

Given an unsorted array A ,
find the i th order statistic.

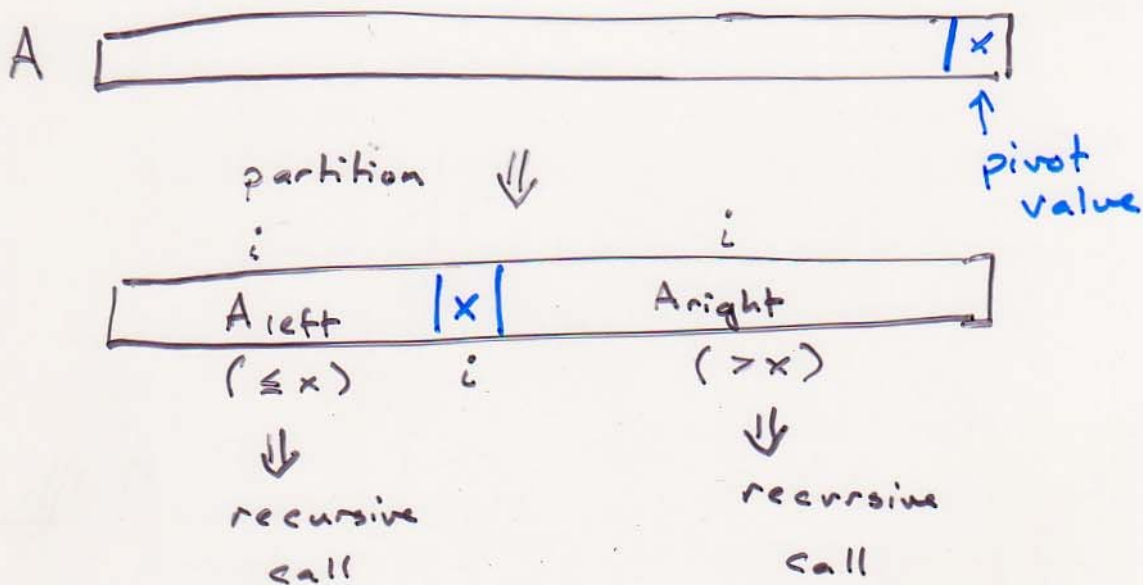
a possible solution

- sort the array

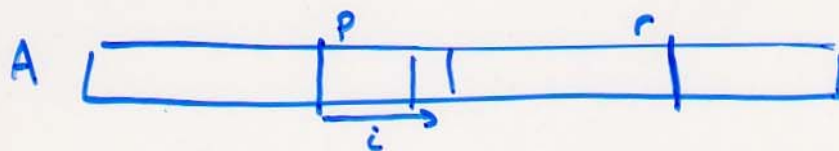
$\Rightarrow O(n \lg n)$ time

it turns out, there is a
linear time ($O(n)$) algorithm
to solve this.

First approach: use
quicksort algorithm
ideas...



$SELECT(A, p, r, i)$



p = start index

r = end index

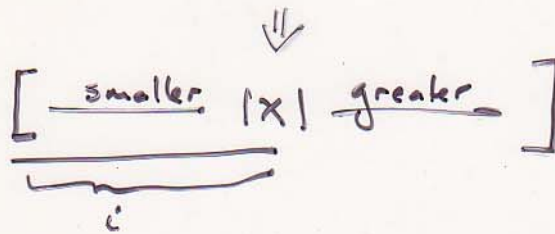
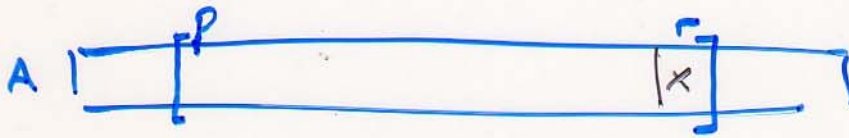
i = order statistic to
find inside $A[p \dots r]$

PARTITION (A, p, r)

A = array

p = start index

r = end index



returns i