

CS 223 Lec 6

Linear Time Selection



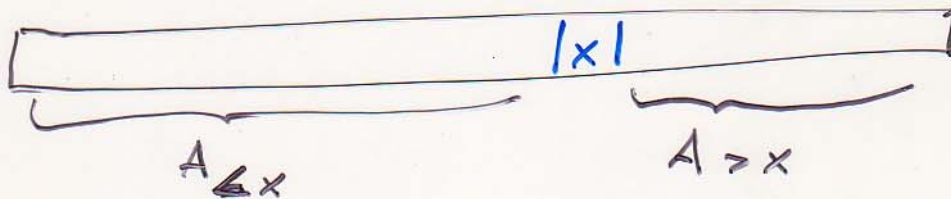
Problem: find the smallest
(largest) value in A?

M



select median x of M \Leftarrow rec. call #1

- use x as a pivot to partition A.



- depending on $i < k$ or $i > k$
we recurse on left hand
or right hand side...

\Leftarrow rec. call #2.

Running Time Analysis

$T(n)$ = time used to select any position on an input array of size n .

Claim :

rec. call #1

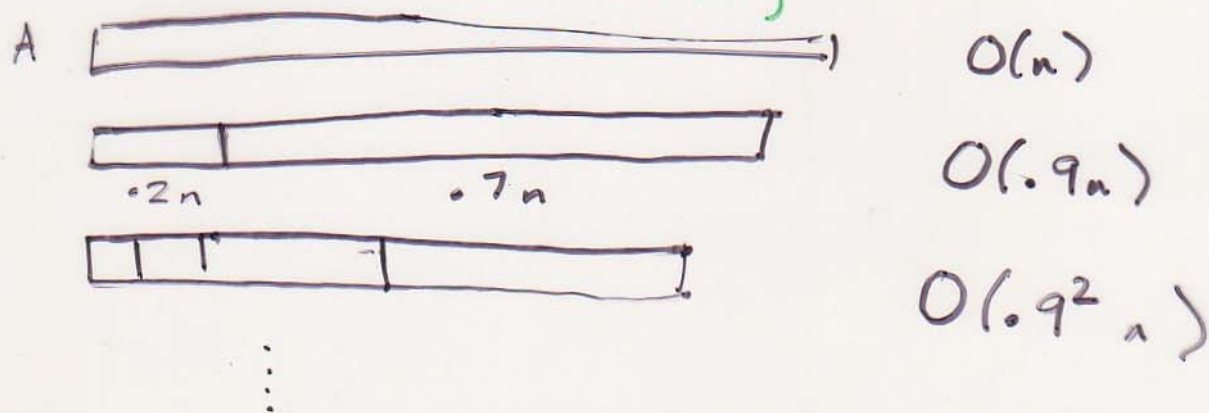
rec. call #2

$$T(n) \leq T\left(\left\lceil \frac{n}{5} \right\rceil\right) + T\left(\frac{7n}{10} + 6\right) + O(n)$$

||
|M|
(finding median of M)

partition around x

call history



$$O(n) + O(.9n) + O(.9^2 n) + \dots$$

$$= O(n) (1 + .9 + (.9)^2 + (.9)^3 + \dots)$$

geometric series

$$1 + x + x^2 + x^3 + \dots$$

$$= O(n) \left(\frac{1}{1 - (.9)} \right)$$

$$= 10 \cdot O(n)$$

$$= O(n)$$

So

$$T(n) = O(n).$$

Hash Tables

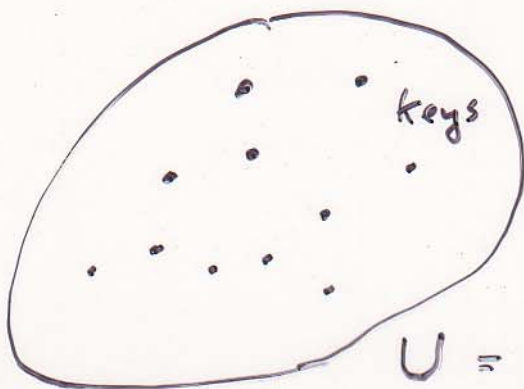
- for supporting dynamic sets of elements
- each element has a key value (not necessarily unique)

OPERATIONS

INSERT
DELETE
SEARCH

Direct-access table

use a simple array that is indexed by set of key values



$U =$ universe of
(potential)
key values

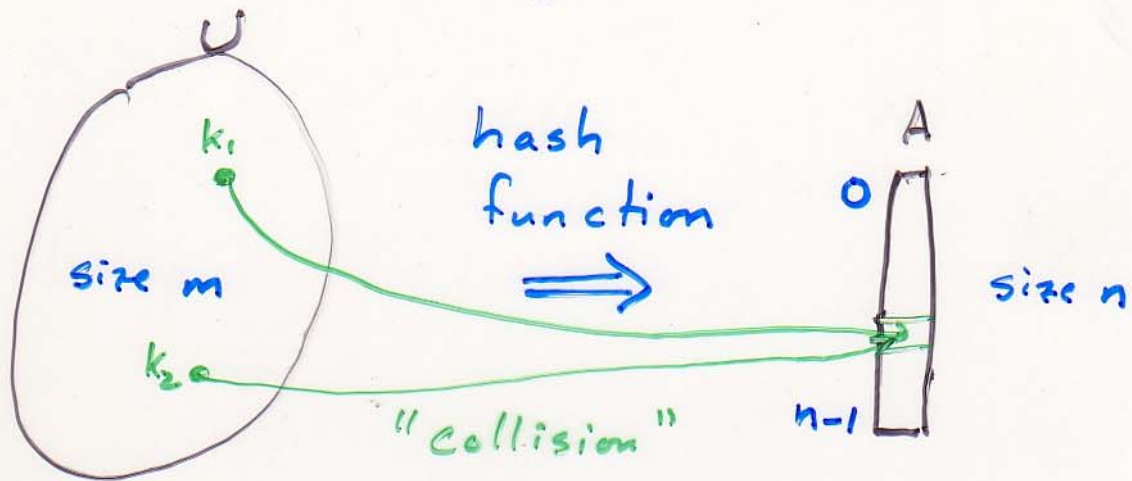
example

$U =$ positive integers
 $\leq 10^9$

$K =$ the actual keys
used in
an application

typically $|K| \ll |U|$

Hash Tables



Example

$$h(x) = x \bmod n$$

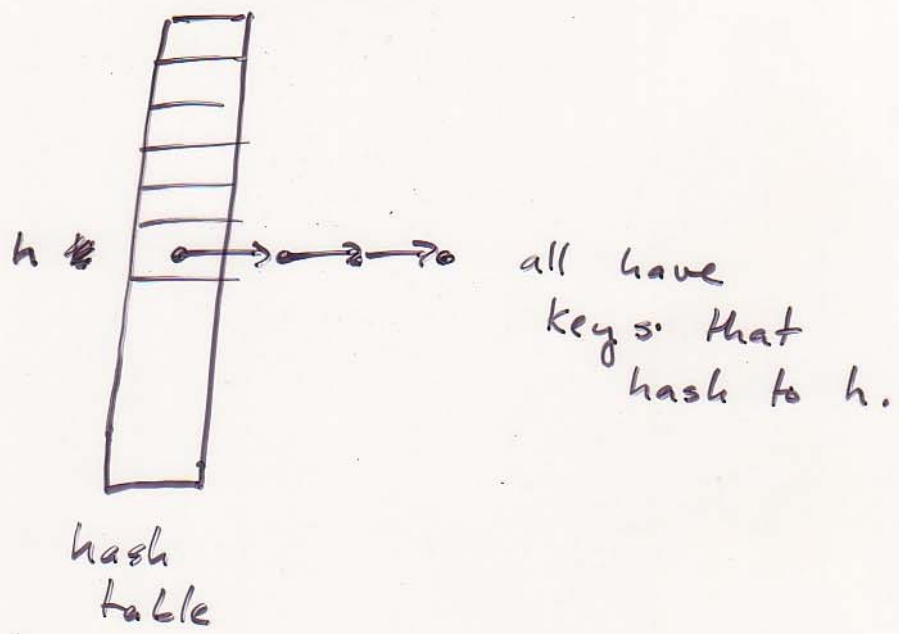
↑
assuming x

is a positive integer.

How to deal with collisions?

"Chaining"

- use a linked list
 at each position in
 the hash table to
 store ~~of~~ multiple elements
 ...



Worst case behavior?

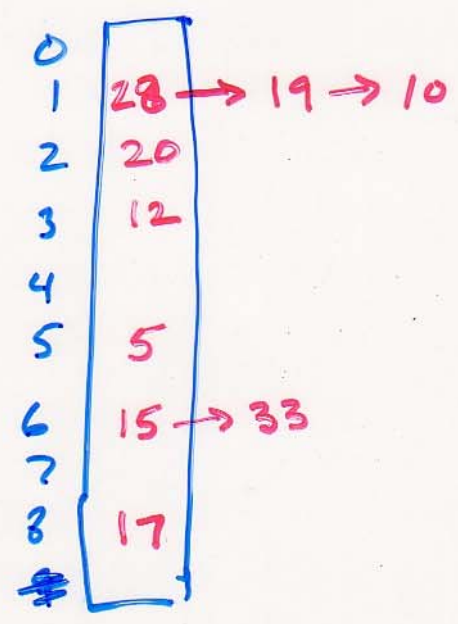
is bad, it occurs ~~if~~
 when all keys have
 the same hash value..

Example

Insert keys 5, 28, 19, 15, 20, 33,
12, 17, 10 into

a hash table with 9 slots

$h(k) = k \text{ mod } 9.$ (use chaining)



Analysis of Hash Tables

§ suppose

1. hash table has size m
2. there are n elements to store.

Defn the load factor
is $\frac{n}{m}$

Assumption

simple uniform hashing

- $P_r \left[\begin{array}{l} \text{a given key} \\ \text{hashes to} \\ \text{location } i \end{array} \right] = \frac{1}{m}$
- hashing is independent.