

CS 201 Dynamic Data Structures (2)

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Link List Operation (Searching)

Find First occurrence of target in the list.

```

1. /*
2.  * Searches a list for a specified target value. Returns a pointer to
3.  * the first node containing target if found. Otherwise returns NULL.
4.  */
5. list_node_t *
6. search(list_node_t *headp, /* input - pointer to head of list */
7.        int target) /* input - value to search for */
8. {
9.     list_node_t *cur_nodep; /* pointer to node currently being checked */
10.
11.    for (cur_nodep = headp;
12.         cur_nodep != NULL && cur_nodep->digit != target;
13.         cur_nodep = cur_nodep->restp) {}
14.
15.    return (cur_nodep);
16. }
    
```

1. What if I put `cur_nodep++` instead of `cur_nodep -> restp`? Could that work? When?
2. What if the order of the following tests are reversed? `(cur_nodep != NULL) && (cur_nodep -> digit != target)`

Overview

- Linked list basics
- List Searching
- Insertion
- Deletion
- Stack
- Queue

Link List Operation (Insertion at Head)

- Insert at list's head (i.e. at the front of the list)

```

list_node_t * insertH (list_node_t *pHead, int v) {
    list_node_t *newp;
    newp = (list_node_t *) malloc(sizeof(list_node_t));
    newp->digit = v;
    newp->restp = pHead;
    return newp; /* return pointer to the new head of the list */
}
    
```

```

int main(void){
    list_node_t * pHead = NULL;
    pHead = insertH(pHead, 3);
    pHead = insertH(pHead, 5);
    pHead = insertH(pHead, 7);
    pHead = insertH(pHead, 9);
}
    
```

```

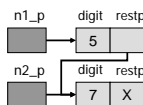
typedef struct list_node_s {
    int digit;
    struct list_node_s *restp;
} list_node_t;
    
```

Linked List basics

- A linked list is a sequence of nodes in which each node but the last contains the address of the next node.

```

typedef struct list_node_s {
    int digit;
    struct list_node_s *restp;
} list_node_t;
list_node_t *n1_p, *n2_p;
n1_p = (list_node_t *) malloc (sizeof(list_node_t));
n2_p = (list_node_t *) malloc (sizeof(list_node_t));
n1_p -> digit = 5;
n2_p -> digit = 7;
n1_p -> restp = n2_p;
n2_p -> restp = NULL;
    
```



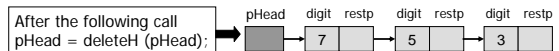
Link List Operation (Deletion at Head)

- After the insertion at last slide, your list now looks like



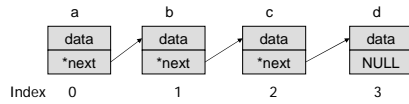
```

list_node_t * deleteH (list_node_t * pHead){
    list_node_t *newp = pHead -> restp;
    /* newp is now pointing to the 2nd element of the list */
    free(pHead);
    return newp;
    /* return pointer to the new head of the list */
}
    
```



Lists are arrays

- Quite often the lists are treated as arrays, that can change their size dynamically.

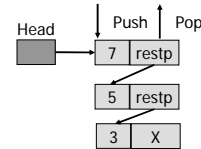


- Assumptions

- Indexing of list starts at 0 (as in arrays).
- Every index value is unique.
- Indices are in growing order (incremented by 1).

Representing a stack with a linked list

- Like a push down stack of books.
- Push (insert) at top (pointed by Head)
- Pop (remove) from the top (pointed by Head)
- Last in first out (LIFO) architecture.



Link List Operation (Deletion at Index)

- Delete element at some index of the list

```
list_node_t * deleteIndex (list_node_t * pHead, int index) {
    int i;
    list_node_t * newwp, *p, * tmpp;
    if (index == 0) { // deleting head
        newwp = pHead -> restp;
        free(pHead);
        return newwp;
    }
}
```

Push on Stack

```
typedef struct list_node_s {
    int digit;
    struct list_node_s *restp;
} list_node_t;
```

Return pointer to the new head of the stack

Pointer to the current head of the stack

```
list_node_t * push (list_node_t * sHead, int v) {
    list_node_t * p = (list_node_t *)
        malloc(sizeof(list_node_t));
    p -> digit = v;
    p -> restp = sHead;
    return p;
}
```

```
int main(void) {
    list_node_t * sHead = NULL;
    /* Function call for push */
    sHead = push(sHead, 3);
    sHead = push(sHead, 5);
    return 0;
}
```

Link List Operation (Deletion at Index)

```
else { // deleting element other than the head
    for (p = pHead, i = 1; (i < index) && (p -> restp != NULL); i++)
        /* searching for the element that has a pointer to the one to be
        deleted */
        p = p -> restp;
    tmpp = p -> restp;
    /* tmpp now points at the element to be deleted */
    if (tmpp != NULL)
        p -> restp = tmpp -> restp;
    /* p now points at the element after the one to be deleted */
    else
        /* we are dealing with the last element of the list */
        p -> restp = NULL;
    free(tmpp);
    return pHead;
}
```

Pop from Stack

Return pointer to the new head of the stack

Pointer to the current head of the stack

Popped digit as output parameter

```
list_node_t * pop (list_node_t * sHead, int * v) {
    list_node_t * p;
    *v = sHead -> digit;
    p = sHead -> restp;
    free (sHead);
    return p;
}
```

```
int main(void) {
    list_node_t * sHead = NULL;
    int val;
    sHead = push(sHead, 3);
    sHead = push(sHead, 5);
    /* Function call for pop */
    sHead = pop (sHead, &val);
    printf("Popped : %d", val);
    return 0;
}
```

Stack (Summary)

- In stack only top element can be accessed.
- You could make a stack with an array.
 - Linked list is just one way.
- Common design for function invocation.
- Both push and pop are constant time operations on stack.

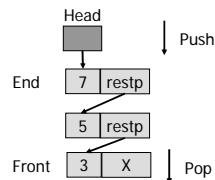
Pop from Queue (from the bottom)

```
list_node_t * pop (list_node_t * qHead, int * v) {
    list_node_t * qEnd, * qFront = NULL;
    if (qHead -> restp = NULL) { // Queue has only one element
        *v = qHead -> digit;
        free (qHead);
        return NULL;
    }
    for (qEnd = qHead; qEnd -> restp != NULL; qEnd = qEnd -> restp)
        qFront = qEnd;
    *v = qEnd -> digit;
    qFront -> restp = NULL;
    free(qEnd);
    return qHead;
}
```

Can we write this more efficiently?

Representing a Queue with a linked list

- Like a queue of people waiting
- Push at the Head (i.e at the end of the list).
- Pop from the Bottom (i.e from the front of the list)
- First In First Out (FIFO)



Queue (Summary)

- You could implement it as an array too.
- You could make a hybrid of stack/queue to access at either end.
- Common design for process scheduling, event processing, buffering, input/output etc.
- In our design push is constant time, but pop is $O(n)$ linear time (where n is the number of elements in the queue).
- If we record two pointers (front and end) instead of only one pointer pointing to the head of the list – both push and pop would have constant time.
 - See the implementation in your textbook.

Push on Queue

Return pointer to the new head of the queue

Push is same as stack (at Head)

```
list_node_t * push (list_node_t * qHead, int v) {
    list_node_t * p = (list_node_t *)
        malloc(sizeof(list_node_t));
    p -> digit = v;
    p -> restp = qHead;
    return p;
}
```

Return the new node as the head of the queue

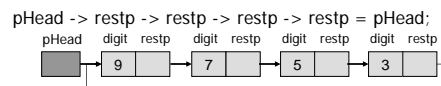
```
typedef struct list_node_s {
    int digit;
    struct list_node_s *restp;
} list_node_t;
```

Pointer to the current head of the queue

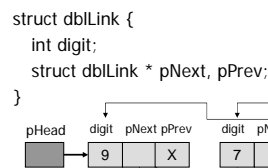
```
int main(void) {
    list_node_t * qHead = NULL;
    /* Function call for push */
    qHead = push(qHead, 3);
    qHead = push(qHead, 5);
    return 0;
}
```

Circular and double linked list

- Circular linked list



- Double linked list



Important !

- Lab 6 grades are posted.
- Remember Exam 2 next Wednesday.
- Covers
 - Array
 - String
 - Recursion
 - Structure & Union
 - Dynamic data storage (upto queue)
- All rules of exam 1 will be applicable.