

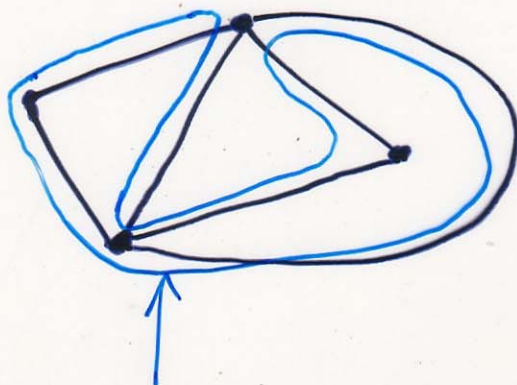
CS 222 Lec 22

Graphs

$$G = (V, E)$$

↑            ↑  
vertices    edges

Eulerian Graphs



Eulerian circuit

Thm : A graph is Eulerian  
iff all vertices have  
even degree.



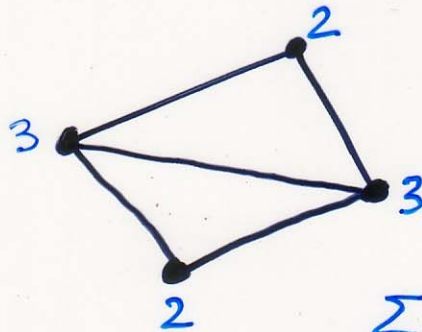
$\deg(v) = 4$

## Properties of Graphs

Thm If  $G = (V, E)$  is undirected

Then

$$\sum_{v \in V} \deg(v) = 2 \cdot |E|$$



$$\begin{aligned} \sum \deg(v) &= 10 \\ &= 2 \cdot |E| \end{aligned}$$

### Proof

Every edge contributes +1 to the degrees of two vertices, so increases  $\sum \deg(v)$  by +2.

## Corollary

In any graph, the number of odd degree vertices must be even.

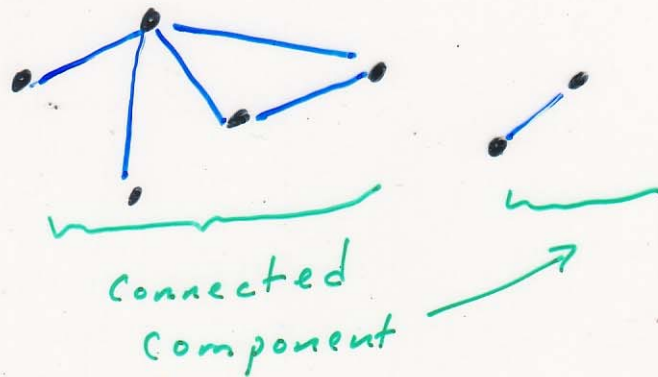
Question is it possible to have a graph with the following degree sequence:

2, 3, 3, 4, 5 ?

↙ ↘  
odd  $\Rightarrow$  so no.

Defn A graph is connected

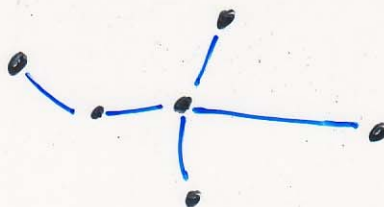
if  $\forall a, b \in V$ , there is a path in  $G$  from  $a$  to  $b$



Defn A graph is a tree

if it is:

- 1) connected
- 2) does not contain any cycles.

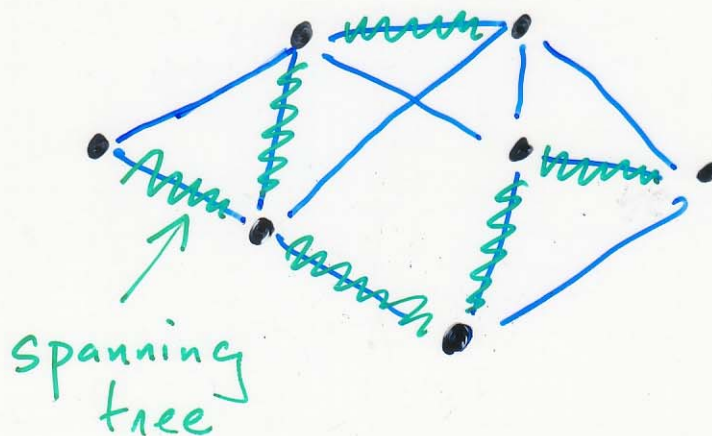


Observe :

Any tree  $T$  with  $n$  vertices has  $n-1$  edges.

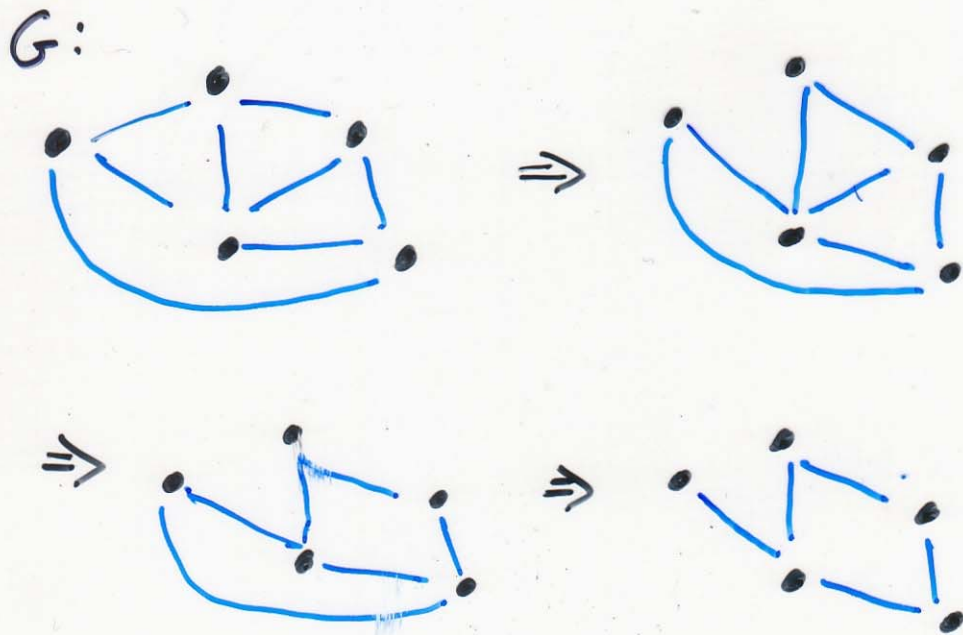
### Spanning Trees

A spanning tree  $T$  for a graph  $G$  is a subset of the edges of  $G$  that forms a tree, and each vertex in  $G$  is connected by  $T$ .



## Algorithm to find Spanning Trees.

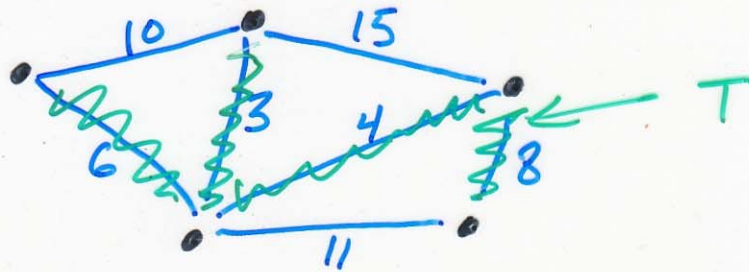
- 1) Let  $G$  be input graph (connected)
- 2) Find a cycle in  $G$ , delete any edge in cycle.
- 3) Repeat step 2. until no cycles remain.





## Minimum Spanning Trees

given: weighted undirected graph



$$\text{cost}(T) = 22.$$

$$\text{cost}(T) = \sum_{e \in T} w(e)$$

$w(e)$  = weight of edge  $e$ .

Problem find a  $\text{MST}_G^{T^*}$  for  $G$

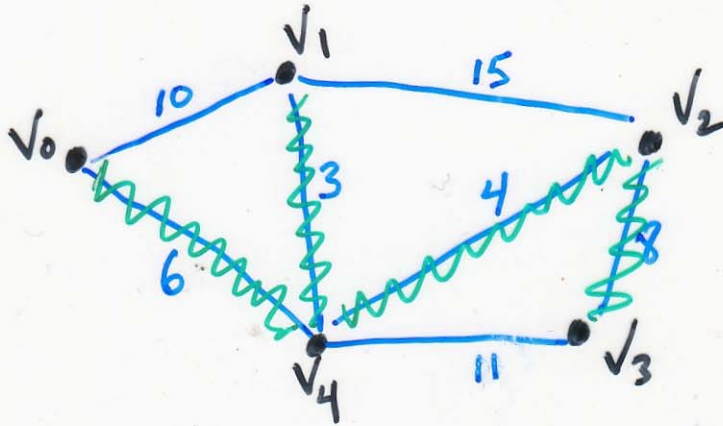
where

$$\text{cost}(T^*) \leq \text{cost}(T)$$

$T$  is any spanning tree.

## Prim's algorithm

"grow a MST"



1) start by picking a random vertex.

$$T = \{V_0\}$$

2) repeat: find the cheapest edge that goes outside of  $T$  and add to  $T$ .

3) repeat until all vertices belong to  $T$ .

Another MST alg:

## Kruskal's algorithm

1. sort all edges by weight.

$$w(e_1) \leq w(e_2) \leq w(e_3) \leq \dots \leq w(e_m)$$

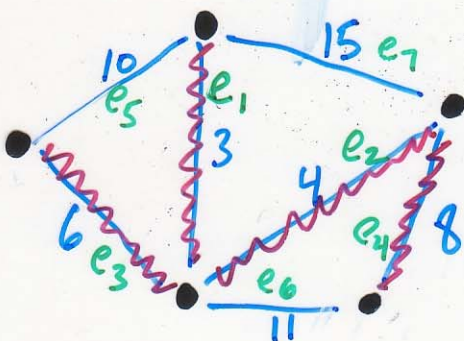
$$T = \{\}$$

2. for  $i = 1$  to  $m$

add  $e_i$  to  $T$

provided doesn't create  
a cycle in  $T$ .

3. return  $T$ .



Question:

Suppose  $G$  has  
all unique edge  
weights? Is  $G$ 's  
MST unique?

yes.