

1. A. Sweep the tape input from left to right
if there are no more 0s or 1s, ~~REJECT~~
- B. Sweep the tape input from right to left
mark ~~last~~ two 0s with an X and ~~the rest~~ ^{one} with an X
if this was possible, go to A
Otherwise ACCEPT

2. Problem 3.15 (d)

for any ~~decidable~~ language L ,

let M be the Tm that ~~decides~~ it.

Construct Tm M' to decide the complement of L as follows:

"On input w :

1. Run M on w . If M accepts, then REJECT.

If M rejects, then ACCEPT."

3. Problem 4.2

$\text{EQ}_{\text{DFA}, \text{REG}} = \{ \langle A, R \rangle \mid A \text{ is a DFA, } R \text{ is a regular expression and } L(A) = L(R) \}$

$\text{EQ}_{\text{DFA}, \text{REG}}$ is decidable

\exists "On input $\langle A, R \rangle$ where A is a DFA and R is a regular expression"

- 1. Convert regular expression R to an equivalent NFA N using the procedure in Theorem 1.54
- 2. Convert NFA N to an equivalent DFA B using the procedure in Theorem 1.39
- 3. Run TM F from Theorem 4.5 on $\langle A, B \rangle$
- 4. If TF accepts, ACCEPT.
If TF rejects, REJECT."

Construction

4. Problem 4.3

$\text{All}_{\text{DFA}} = \{A \mid A \text{ is a DFA and } L(A) = \Sigma^*\}$ is decidable

Construction:

" On input $\langle A \rangle$ where A is a DFA

1. Construct DFA B as follows

$$Q = \{q_1\}$$

$$\Sigma$$

$$\delta(q_1, a) = q_1, \quad \forall a \in \Sigma$$

$$q_{\text{start}} = q_1$$

$$q_{\text{accept}} = q_1$$

always
accepts

2. Run TM F from Theorem 4.5 on $\langle A, B \rangle$

3. If F accepts, ACCEPT.

If F rejects, REJECT."

5. Problem 4.7

Suppose B is countable. We could then create a table such as

\mathbb{N}	B								
1	0	1	1	0	1	1	0	...	
2	1	0	0	0	1	1	1	...	
3	1	1	0	1	0	1	1	...	
4	0	0	0	0	1	0	1	...	
...									

However, we can always find a number that is not in the table using diagonalization,

e.g. 1 0 0 0 ...

Contradiction! Therefore B is uncountable.

b. Problem 4.8

(Many answers are possible.)

N	i	j	k	
1	1	1	1	$\left. \begin{array}{l} \\ \\ \end{array} \right\} i+j+k=3$
2	1	1	2	$\left. \begin{array}{l} \\ \\ \end{array} \right\} i+j+k=4$
3	1	2	1	$\left. \begin{array}{l} \\ \\ \end{array} \right\} i+j+k=4$
4	2	1	1	
5	1	1	3	
6	1	2	2	$i+j+k=5$
7	1	3	1	(go in
8	2	1	2	ascending order over i,j,k)
9	2	2	1	
10	3	1	1	
	i	j	k	