

CSCI 476: Computer Security

Lecture 2: Computer Systems Review

Reese Pearsall
Fall 2022

Announcements

TA

- **Karishma Rahman**
- karishma.rahman.bd@gmail.com
- Office Hours: Tuesdays 1:00 pm to 3:00 pm
- Location: Barnard 259

Notetaker needed for this class

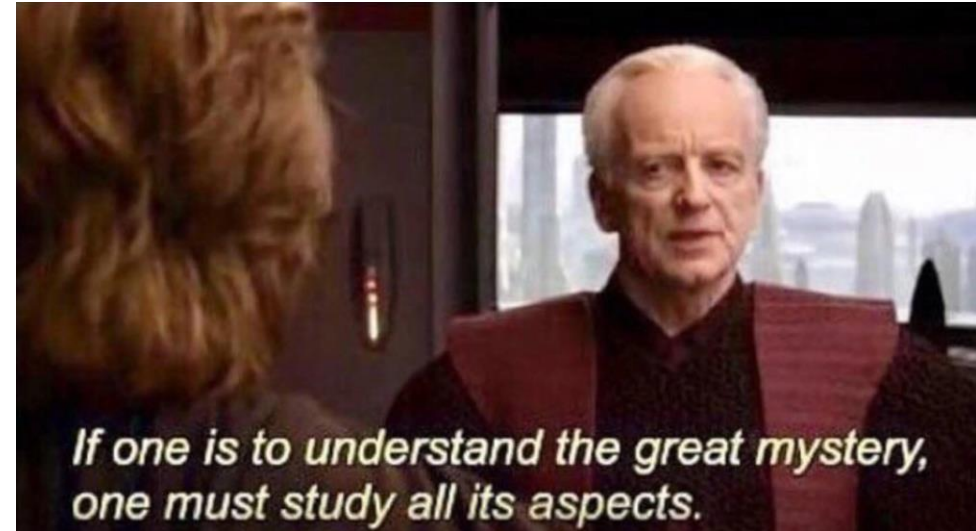
- Contact Office of Disability Services, notetaking@montana.edu, or 406-994-2824
- Paid position

Lab 0 Report Format → Due Tuesday

Research Project Instructions will hopefully be released at the end of next week

Computer Systems Review

To understand the technical aspects of security, we must have a good understanding of how computers work



What is a computer?



What is a computer?

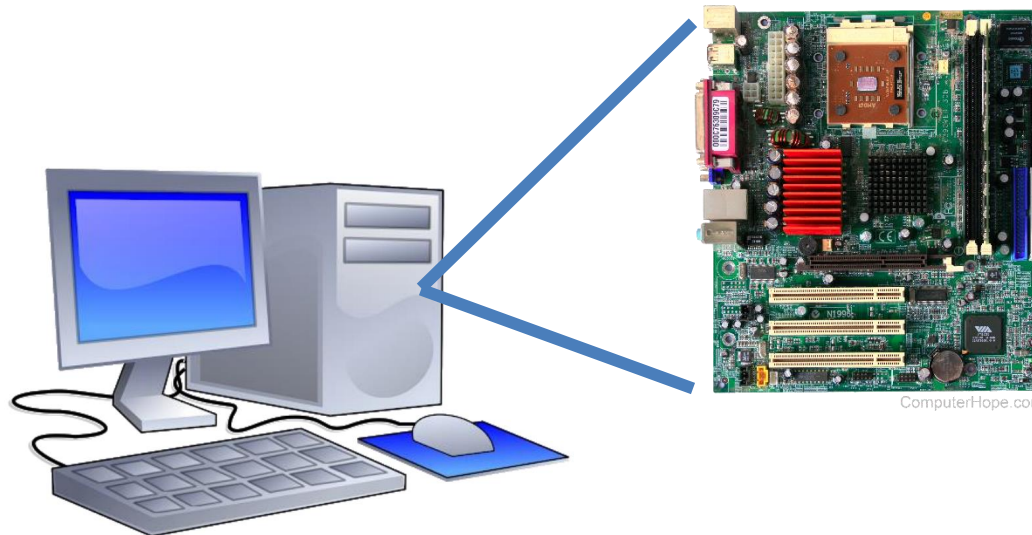


A magical box that does stuff



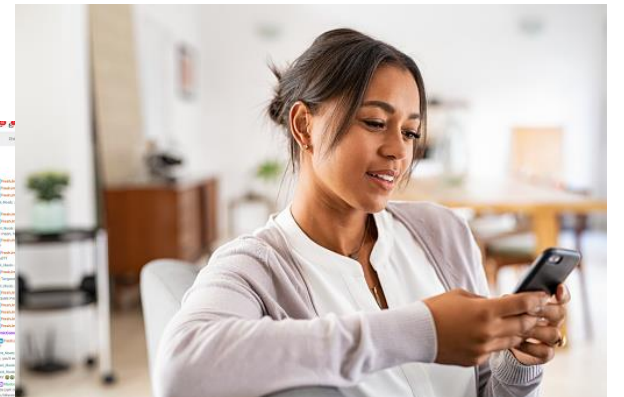
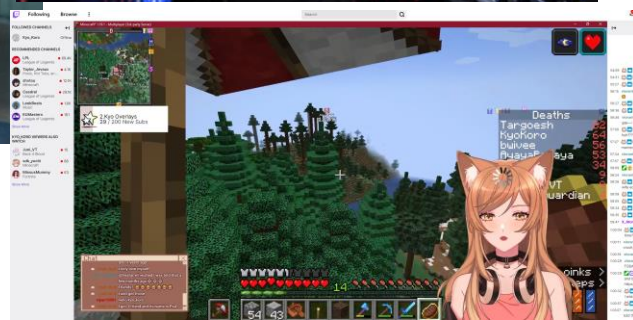
What is a computer?

A **semi**-magical box that ~~does stuff~~ **executes instructions**



What is so magical about a computer?

We use computers every day for many different things



What is so magical about a computer?

Big Idea

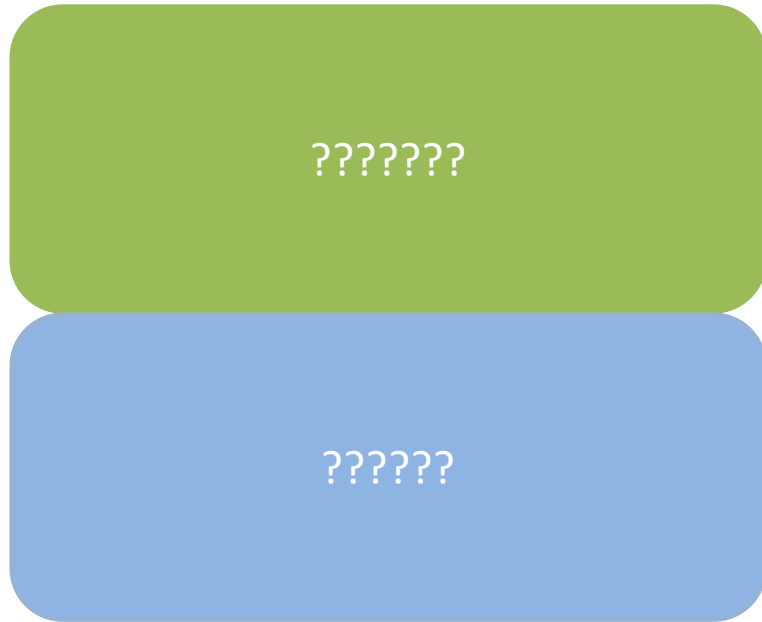
Computers only understand instructions in the form of 0s and 1s (binary)

~~Welcome to CSCI 476~~



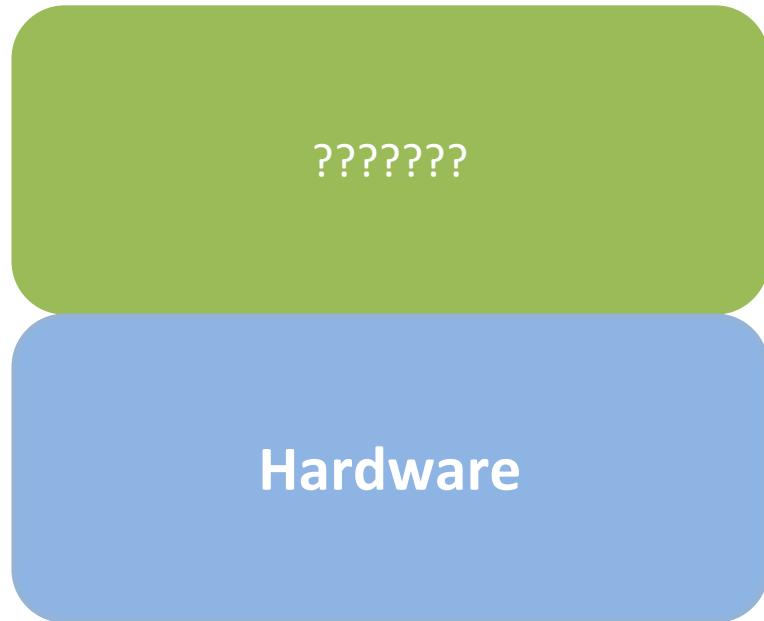
```
01010111 01100101 01101100 01100011 01101111
01101101 01100101 00100000 01110100 01101111
00100000 01000011 01010011 01000011 01001001
00100000 00110100 00110111 00110110
```


How does this happen?



From a high level, we will divide a computer system into two parts

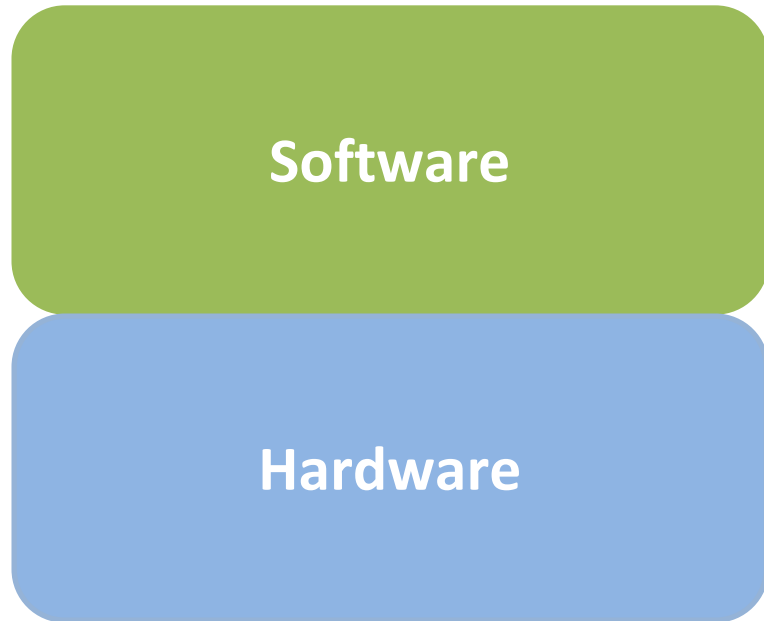
How does this happen?



From a high level, we will divide a computer system into two parts

I. Hardware

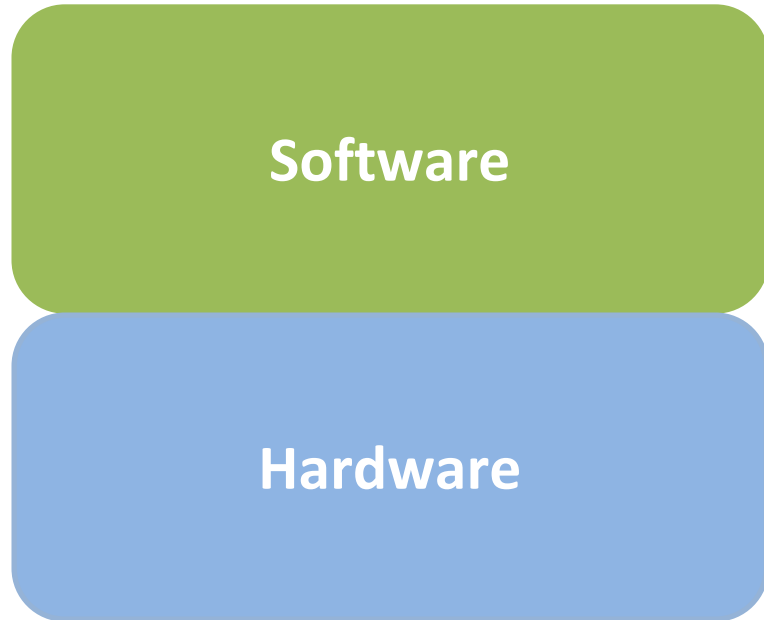
How does this happen?



From a high level, we will divide a computer system into two parts

- I. Hardware**
- II. Software**

How does this happen?



From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**

Symbiotic relationship

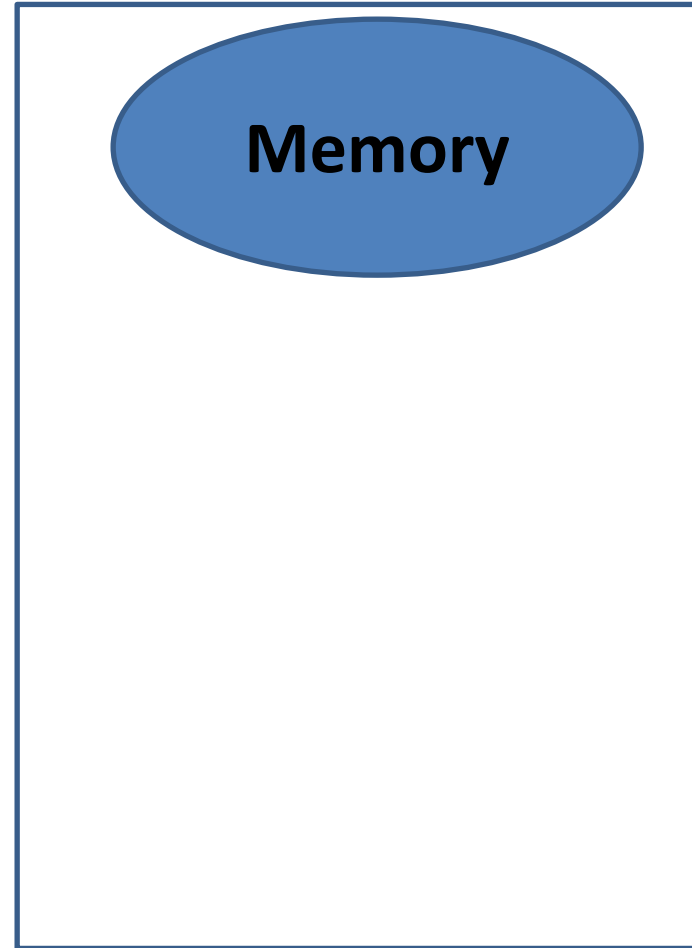
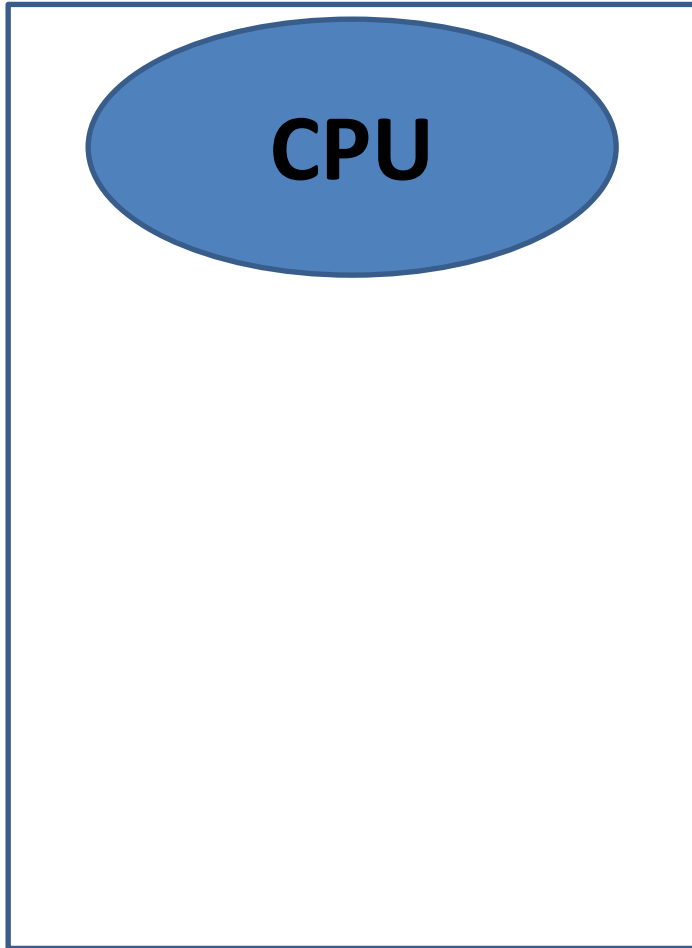


I. Hardware

The **physical** parts of a computer

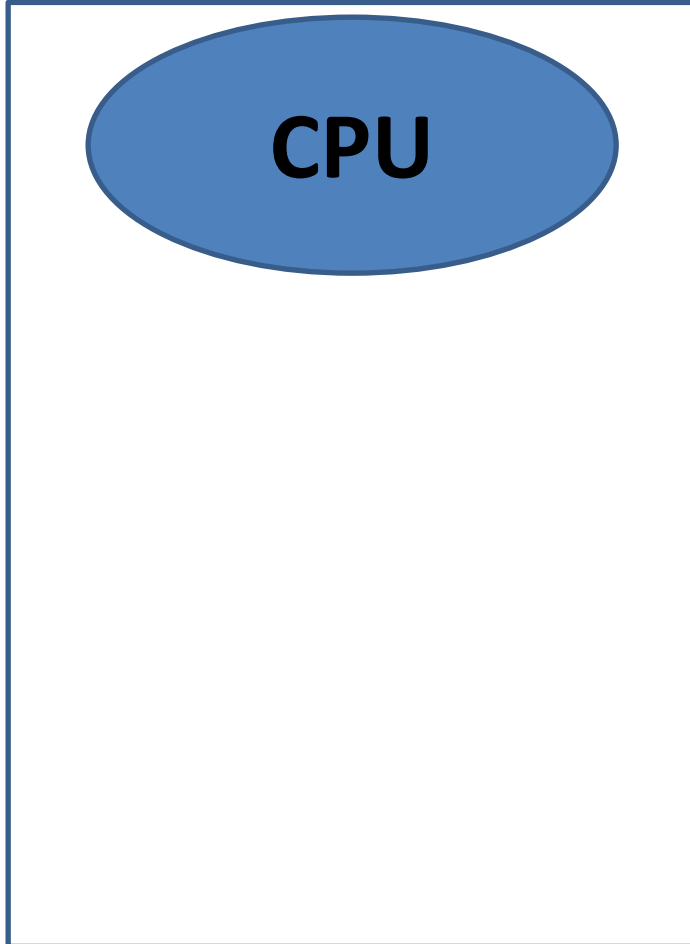


I. Hardware

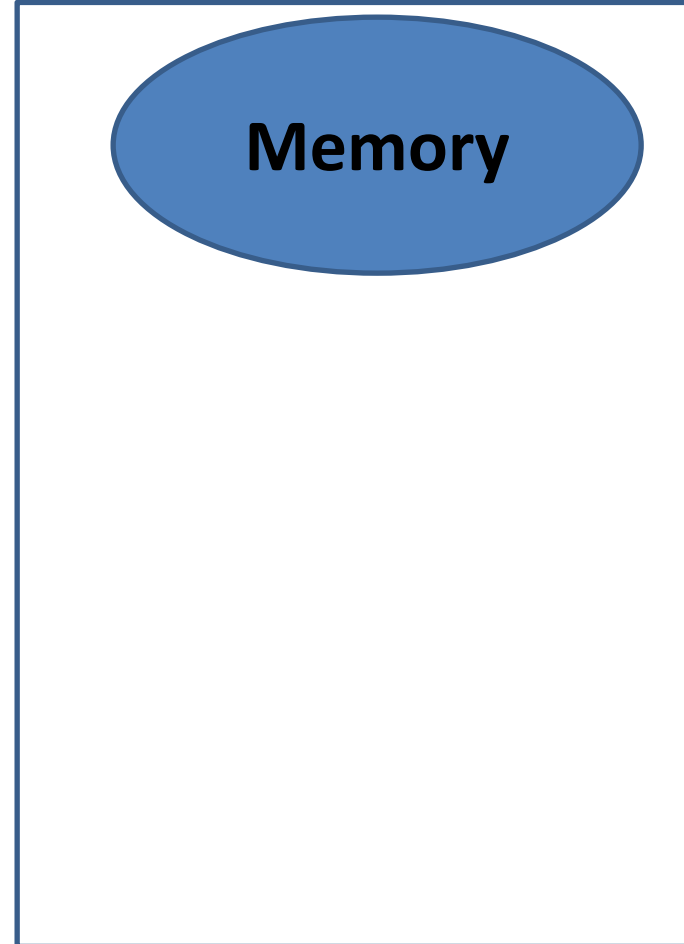


I. Hardware

Brain with no short-term memory

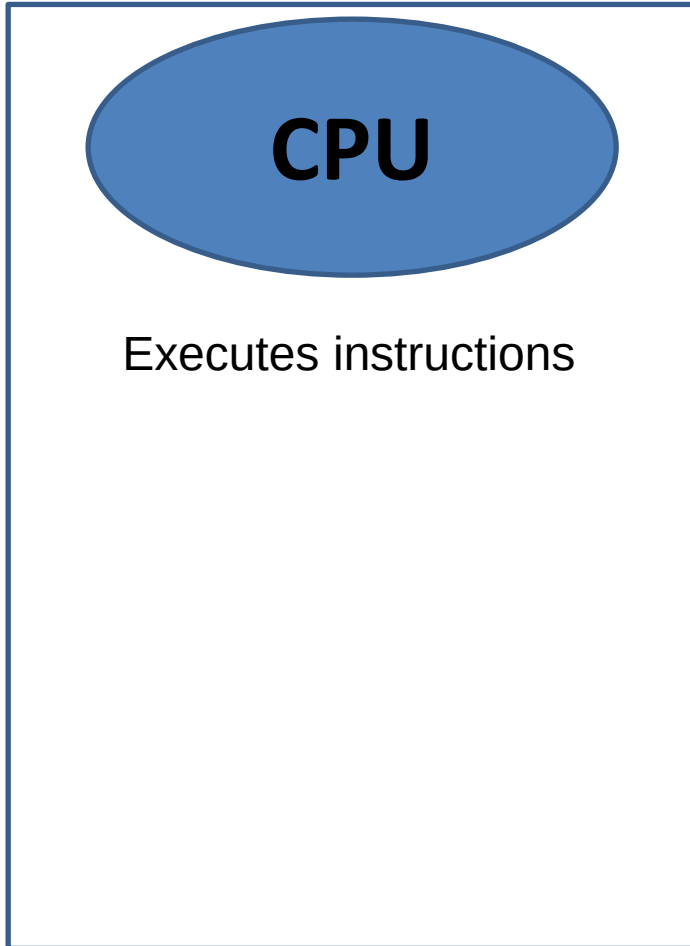


Scratch Pad

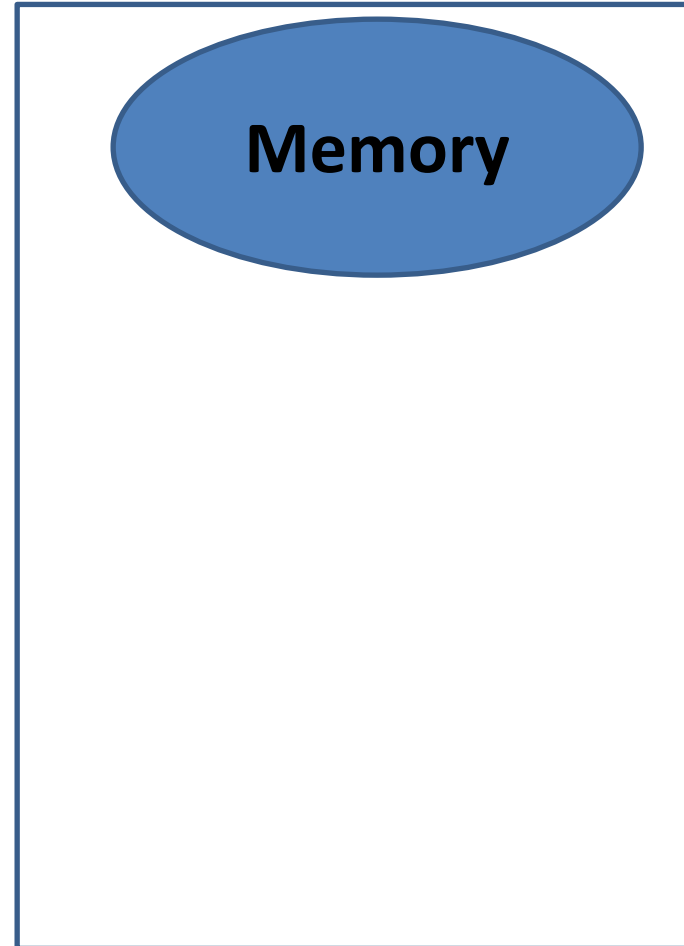


I. Hardware

Brain with no short-term memory



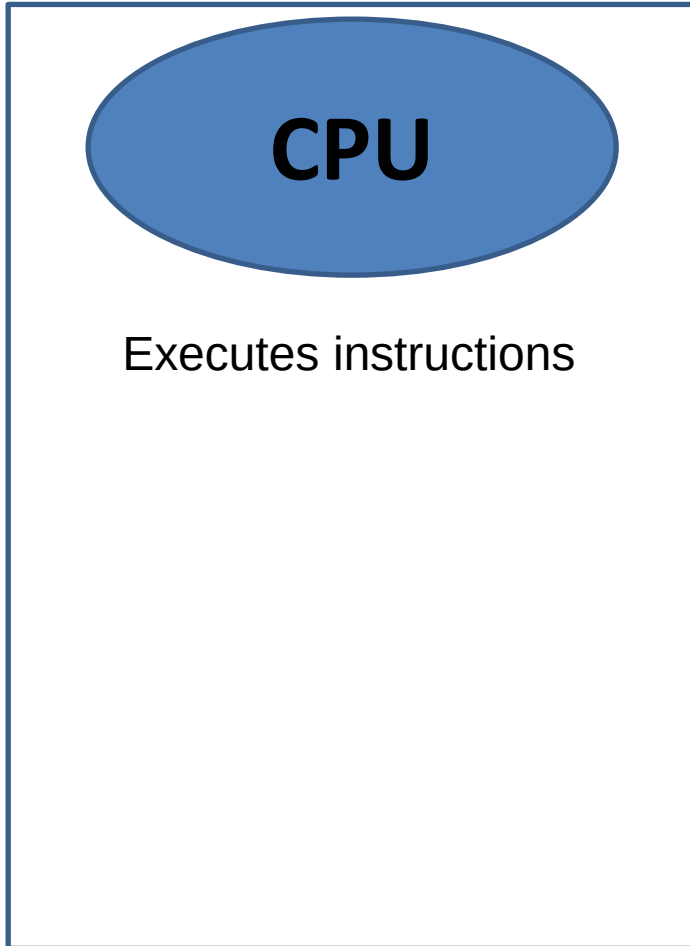
Scratch Pad



I. Hardware



Brain with no short-term memory



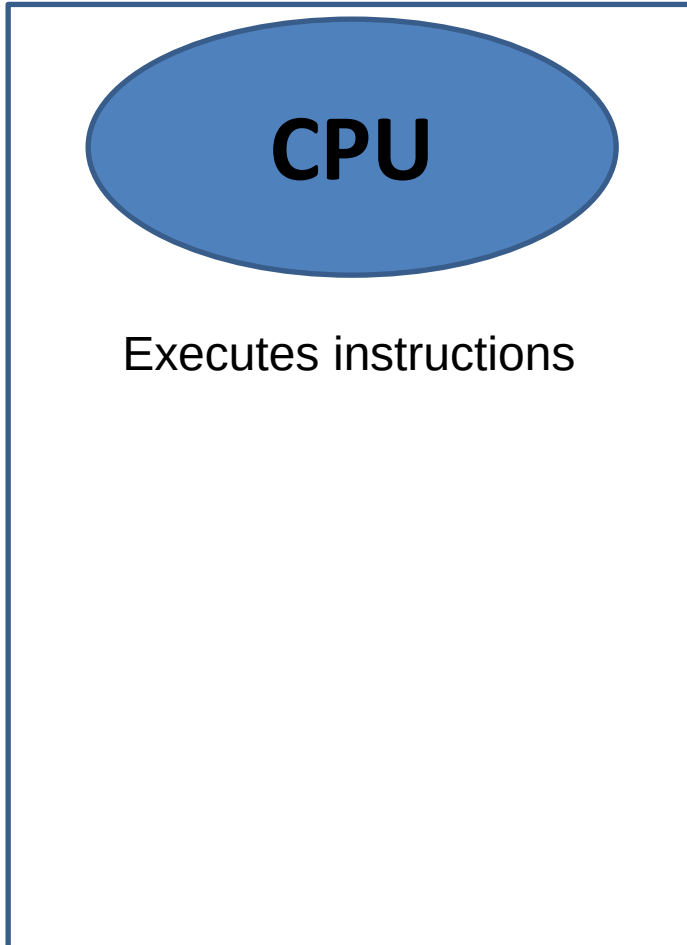
How does it “execute” instructions?

It is sent instructions from another part of the computer

I. Hardware



Brain with no short-term memory



How does it “execute” instructions?

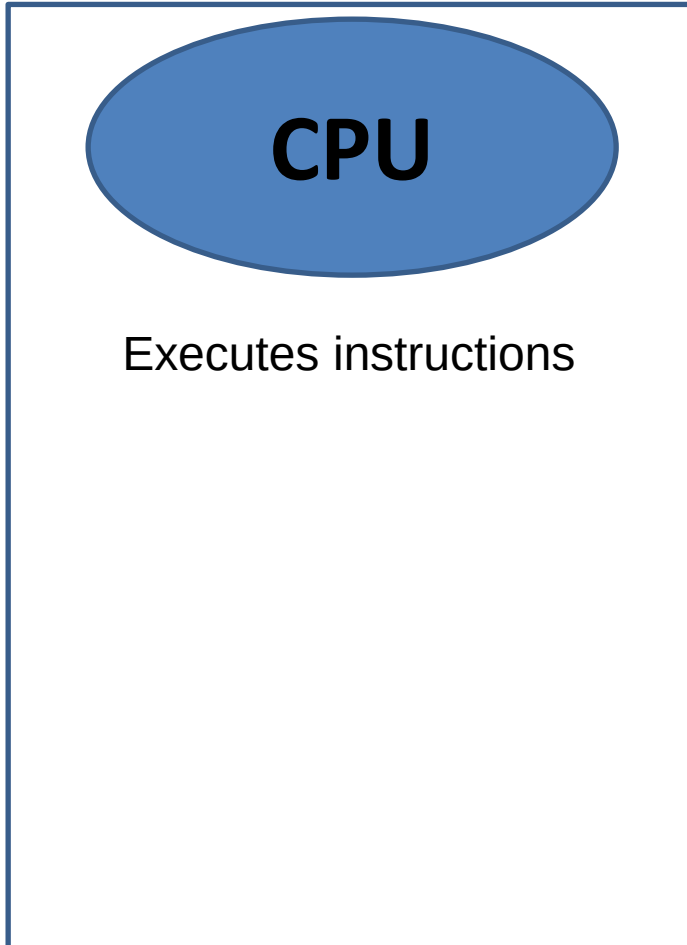
It is sent instructions from another part of the computer

01001100000000110100011100001010

I. Hardware



Brain with no short-term memory



How does it “execute” instructions?

It is sent instructions from another part of the computer

00000000101000010001100000100000 → 00 A1 18 20

00 A1 18 20

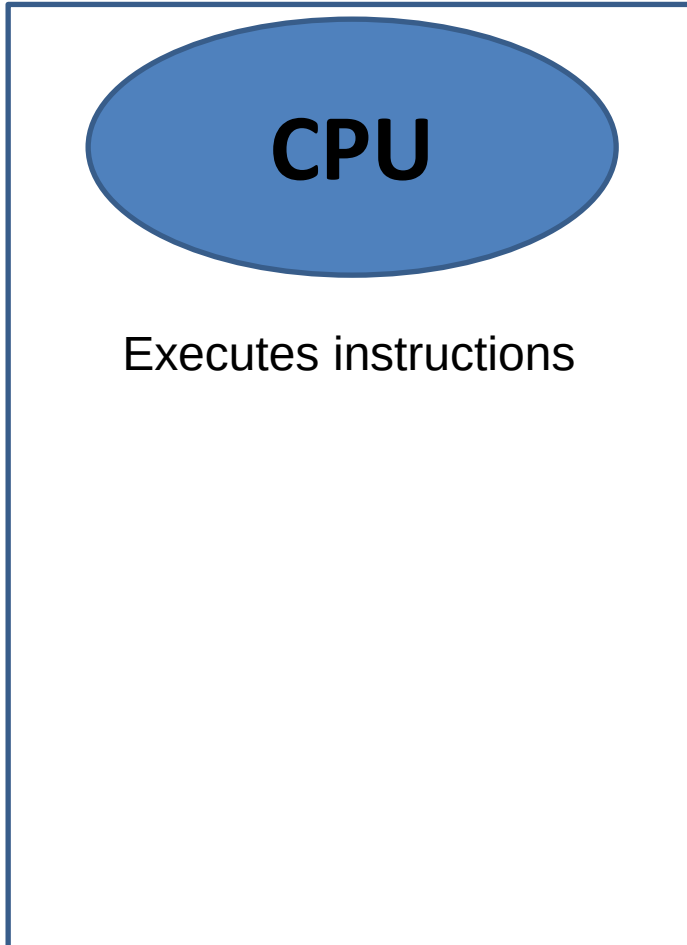
0000	0	1000	8
0001	1	1001	9
0010	2	1010	A
0011	3	1011	B
0100	4	1100	C
0101	5	1101	D
0110	6	1110	E
0111	7	1111	F

Hex_(hexadecimal) is a common representation for binary

I. Hardware



Brain with no short-term memory



How does it “execute” instructions?

It is sent instructions from another part of the computer

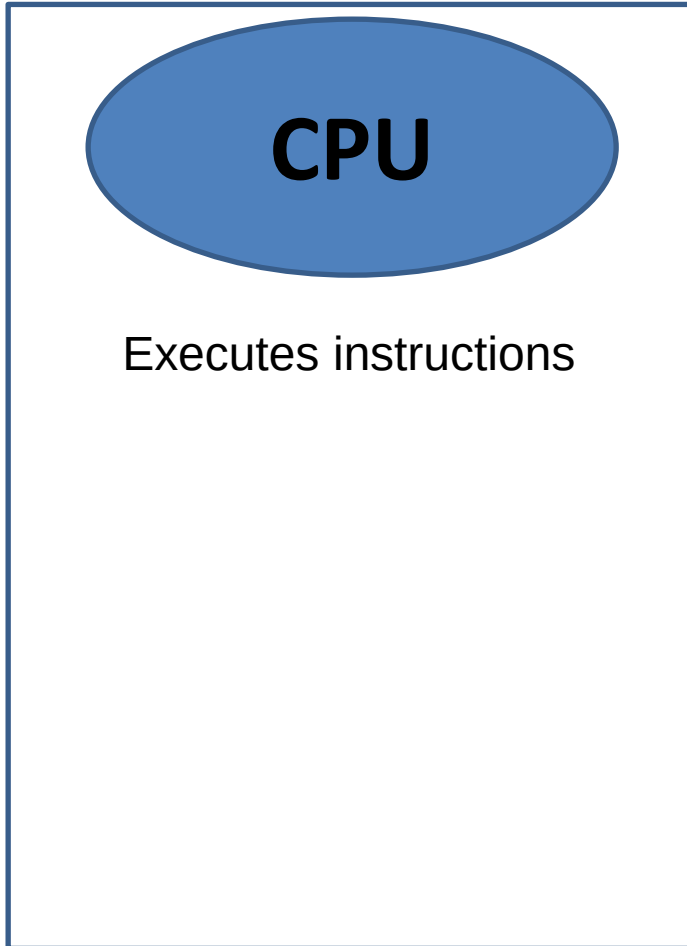
00000000101000010001100000100000

What does this instruction do?????

I. Hardware



Brain with no short-term memory



How does it “execute” instructions?

It is sent instructions from another part of the computer

00000000010100001000110000010000

Opcode

Common MIPS instructions

Notes: *op*, *funct*, *rd*, *rs*, *rt*, *imm*, *address*, *shamt* refer to fields in the instruction format

PC is assumed to point to the next instruction, **Mem** is the byte addressed main memory

Assembly Instruction	Instr Format	op op/funct	Meaning	Comments
<u>add \$rd, \$rs, \$rt</u>	R	0/32	$\$rd = \$rs + \$rt$	Add contents of two registers
sub \$rd, \$rs, \$rt	R	0/34	$\$rd = \$rs - \$rt$	Subtract contents of two registers
addi \$rt, \$rs, imm	I	8	$\$rt = \$rs + imm$	Add signed constant
addu \$rd, \$rs, \$rt	R	0/33	$\$rd = \$rs + \$rt$	Unsigned, no overflow
subu \$rd, \$rs, \$rt	R	0/35	$\$rd = \$rs - \$rt$	Unsigned, no overflow
addiu \$rt, \$rs, imm	I	9	$\$rt = \$rs + imm$	Unsigned, no overflow
mfc0 \$rt, \$rd	R	16	$\$rt = \rd	<i>rd</i> = coprocessor register (e.g. epc, cause, status)
mult \$rs, \$rt	R	0/24	Hi, Lo = $\$rs * \rt	64 bit signed product in Hi and Lo
multu \$rs, \$rt	R	0/25	Hi, Lo = $\$rs * \rt	64 bit unsigned product in Hi and Lo
div \$rs, \$rt	R	0/26	Lo = $\$rs / \rt , Hi = $\$rs \bmod \rt	
divu \$rs, \$rt	R	0/27	Lo = $\$rs / \rt , Hi = $\$rs \bmod \rt (unsigned)	
mfhi \$rd	R	0/16	$\$rd = \text{Hi}$	Get value of Hi
mflo \$rd	R	0/18	$\$rd = \text{Lo}$	Get value of Lo
and \$rd, \$rs, \$rt	R	0/36	$\$rd = \$rs \& \$rt$	Logical AND
or \$rd, \$rs, \$rt	R	0/37	$\$rd = \$rs \$rt$	Logical OR
andi \$rt, \$rs, imm	I	12	$\$rt = \$rs \& imm$	Logical AND, unsigned constant
ori \$rt, \$rs, imm	I	13	$\$rt = \$rs imm$	Logical OR, unsigned constant
sll \$rd, \$rs, shamt	R	0/0	$\$rd = \$rs \ll shamt$	Shift left logical (shift in zeros)
srl \$rd, \$rs, shamt	R	0/2	$\$rd = \$rs \gg shamt$	Shift right logical (shift in zeros)



Common MIPS instructions

Notes: *op*, *funct*, *rd*, *rs*, *rt*, *imm*, *address*, *shamt* refer to fields in the instruction format

PC is assumed to point to the next instruction, **Mem** is the byte addressed main memory

Assembly Instruction	Instr Format	op op/funct	Meaning	Comments
add <i>\$rd</i> , <i>\$rs</i> , <i>\$rt</i>	R	0/32	$\$rd = \$rs + \$rt$	Add contents of two registers
sub <i>\$rd</i> , <i>\$rs</i> , <i>\$rt</i>	R	0/34	$\$rd = \$rs - \$rt$	Subtract contents of two registers
addi <i>\$rt</i> , <i>\$rs</i> , <i>imm</i>	I	8	$\$rt = \$rs + imm$	Add signed constant

MIPS Instruction formats

Format	Bits 31-26	Bits 25-21	Bits 20-16	Bits 15-11	Bits 10-6	Bits 5-0
R	op	rs	rt	rd	shamt	funct
I	op	rs	rt	imm		
J	op	address				

Assembly Instruction	Instr Format	op/funct	Meaning	Comments
div <i>\$rs</i> , <i>\$rt</i>	R	0/26	$Lo = \$rs / \rt , $Hi = \$rs \bmod \rt	Get value of Hi
divu <i>\$rs</i> , <i>\$rt</i>	R	0/27	$Lo = \$rs / \rt , $Hi = \$rs \bmod \rt (unsigned)	Get value of Lo
mfhi <i>\$rd</i>	R	0/16	$\$rd = Hi$	Get value of Hi
mflo <i>\$rd</i>	R	0/18	$\$rd = Lo$	Get value of Lo
and <i>\$rd</i> , <i>\$rs</i> , <i>\$rt</i>	R	0/36	$\$rd = \$rs \& \$rt$	Logical AND
or <i>\$rd</i> , <i>\$rs</i> , <i>\$rt</i>	R	0/37	$\$rd = \$rs \$rt$	Logical OR
andi <i>\$rt</i> , <i>\$rs</i> , <i>imm</i>	I	12	$\$rt = \$rs \& imm$	Logical AND, unsigned constant
ori <i>\$rt</i> , <i>\$rs</i> , <i>imm</i>	I	13	$\$rt = \$rs imm$	Logical OR, unsigned constant
sll <i>\$rd</i> , <i>\$rs</i> , <i>shamt</i>	R	0/0	$\$rd = \$rs \ll shamt$	Shift left logical (shift in zeros)
srl <i>\$rd</i> , <i>\$rs</i> , <i>shamt</i>	R	0/2	$\$rd = \$rs \gg shamt$	Shift right logical (shift in zeros)

I. Hardware



Brain with no short-term memory

CPU

Executes instructions



Must decipher
what instruction
to execute

How does it “execute” instructions?

It is sent instructions from another part of the computer

00000000001010000100011000000100000

Opcode

I. Hardware



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0000000000101000010001100000100000

Opcode \$rs

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000000000101000010001100000100000

Opcode

\$rs

\$rt

I. Hardware



Brain with no short-term memory

CPU

Executes instructions

Field	Bits	Function
opcode	6	Operation code
rs	5	Source register
rt	5	Target register
rd	5	Destination register
shamt	5	Shift amount
funct	6	Function code

Must decipher
what instruction
to execute

How does it “execute” instructions?

It is sent instructions from another part of the computer

0000000001010000100011000000100000

Opcode

\$rs

\$rt

\$rd

I. Hardware



Brain with no short-term memory

CPU

Executes instructions

Field	Bits	Description
opcode	6	Operation code
rs	5	Source register
rt	5	Target register
rd	5	Destination register
shamt	5	Shift amount
immediate	16	Immediate value

Must decipher
what instruction
to execute

How does it “execute” instructions?

It is sent instructions from another part of the computer

0000000000101000010001100000100000

Opcode

\$rs

\$rt

\$rd

shamt

I. Hardware



Brain with no short-term memory

CPU

Executes instructions

Field	Bits	Value
opcode	6	000000
rs	5	00101
rt	5	10000
rd	5	10001
shamt	5	00000
funct	6	100000

Must decipher
what instruction
to execute

How does it “execute” instructions?

It is sent instructions from another part of the computer

000000000101000010001100000100000

Opcode

\$rs

\$rt

\$rd

shamt

funct

I. Hardware



Brain with no short-term memory

CPU

Executes instructions

Field	Bits	Function
opcode	6	Operation code
rs	5	Source register
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000000000101000010001100000100000

Opcode

\$rs

\$rt

\$rd

shamt

funct

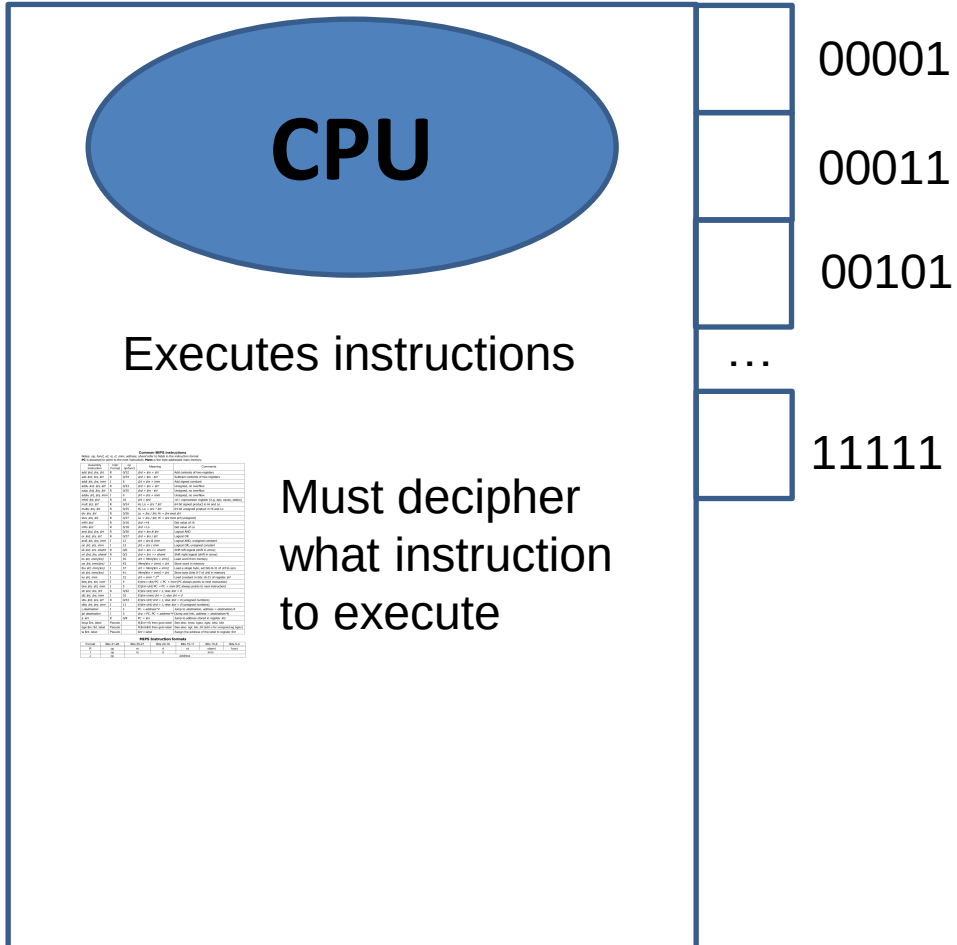
Damn.... I kinda **don't care**



I. Hardware



Brain with no short-term memory



How does it “execute” instructions?

It is sent instructions from another part of the computer

0000000000101000010001100000100000

Opcode

\$rs

\$rt

\$rd

shamt

funct

\$ denotes that it is a **register**

I. Hardware

Brain with no

Execu



Register	Value
EAX	00000000
EBX	00000000
ECX	00000000
EDX	00000000
ESI	00000000
EDI	00000000
ESP	00000000
EBP	00000000
ESI	00000000
EDI	00000000
ESP	00000000
EBP	00000000
ESI	00000000
EDI	00000000
ESP	00000000
EBP	00000000

General-purpose Registers

EAX

AX

AH

AL

EBX

BX

BH

BL

ECX

CX

CH

CL

EDX

DX

DH

DL

ESI

EDI

ESP

(stack pointer)

EBP

(base pointer)

32 bits

16 bits
8 bits 8 bits

Instructions?

part of the computer

000000100000

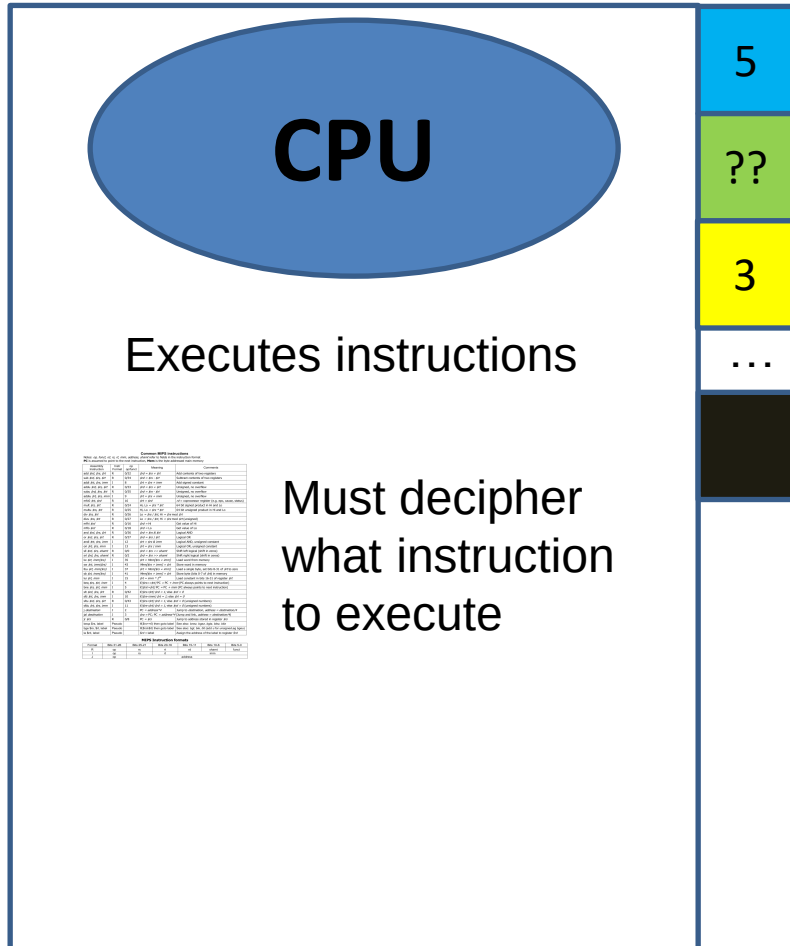
shamt funct

er

I. Hardware



Brain with no short-term memory



How does it “execute” instructions?

It is sent instructions from another part of the computer

0000000000101000010001100000100000

Opcode \$rs \$rt \$rd shamt funct

ADD \$rs, \$rt, \$rd

I. Hardware



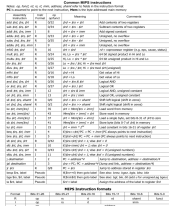
Brain with no short-term memory

Registers

How does it “execute” instructions?

CPU

Executes instructions



Must decipher
what instruction
to execute



The CPU can
add, subtract,
logic, move stuff
around

5	00001
??	00011
3	00101
...	
	11111

It is sent instructions from another part of the computer

0000000000101000010001100000100000

Opcode \$rs \$rt \$rd shamt funct

ADD \$rs, \$rt, \$rd

I. Hardware



Brain with no short-term memory

Registers

How does it “execute” instructions?

CPU

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Must decipher
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5	00001
8	00011
3	00101
...	
	11111

It is sent instructions from another part of the computer

000000**00101****100001****100011**00000100000

Opcode \$rs \$rt \$rd shamt funct

ADD \$rs, \$rt, \$rd

I. Hardware



Brain with no short-term memory

CPU

Executes instructions



Must decipher
what instruction
to execute



The CPU can
add, subtract,
logic, move stuff
around

Register

5

0000

8

0000

3

...

11

CPU uses

Electricity™



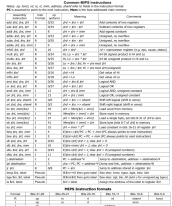
I. Hardware



Brain with no short-term memory

CPU

Executes instructions



Address	Disassembly	Comment
00000000	JMP 00000000	
00000001	JMP 00000000	
00000002	JMP 00000000	
00000003	JMP 00000000	
00000004	JMP 00000000	
00000005	JMP 00000000	
00000006	JMP 00000000	
00000007	JMP 00000000	
00000008	JMP 00000000	
00000009	JMP 00000000	
0000000A	JMP 00000000	
0000000B	JMP 00000000	
0000000C	JMP 00000000	
0000000D	JMP 00000000	
0000000E	JMP 00000000	
0000000F	JMP 00000000	

Must decipher
what instruction
to execute



The CPU can
add, subtract,
logic, move stuff
around

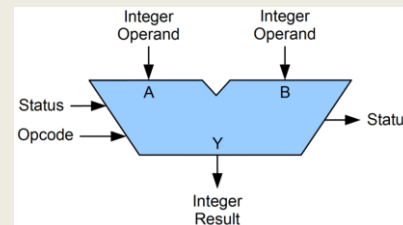
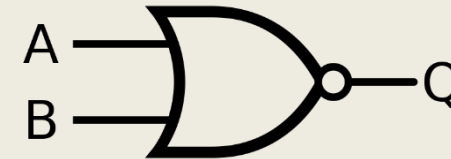
Register

5	0000
8	0000
3	
...	
15	

CPU uses

Electricity™

To decipher and execute instructions



Brain with m

Exec



ions

ISSUE DATE
1 - 11/12/86
2 - 15/12/86
3 - 26/1/87

I. Hardware



Brain with no short-term memory

CPU

Executes instructions



Must decipher
what instruction
to execute



The CPU can
add, subtract,
logic, move stuff
around

Register

5

0000

8

0000

3

...

11

CPU uses

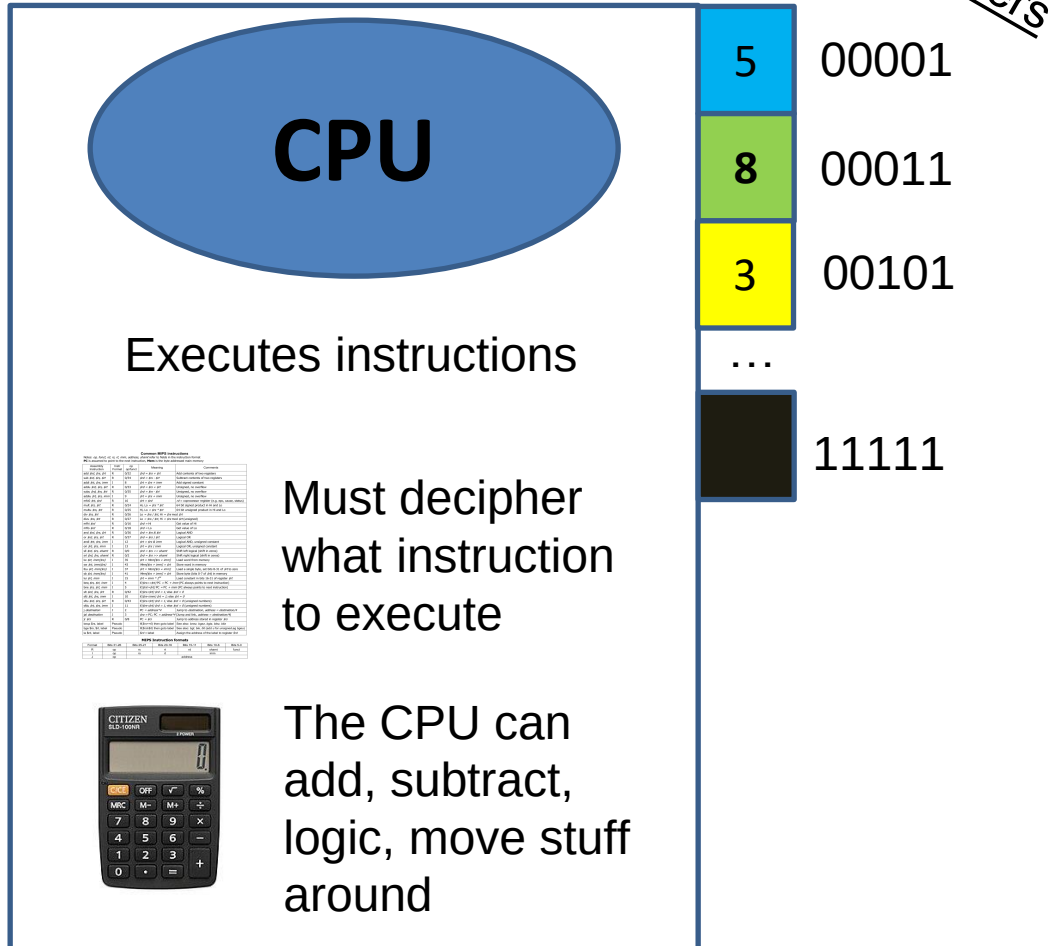
Electricity™

To decipher and execute instructions

*Everything is just **electricity**
on or **electricity** off*

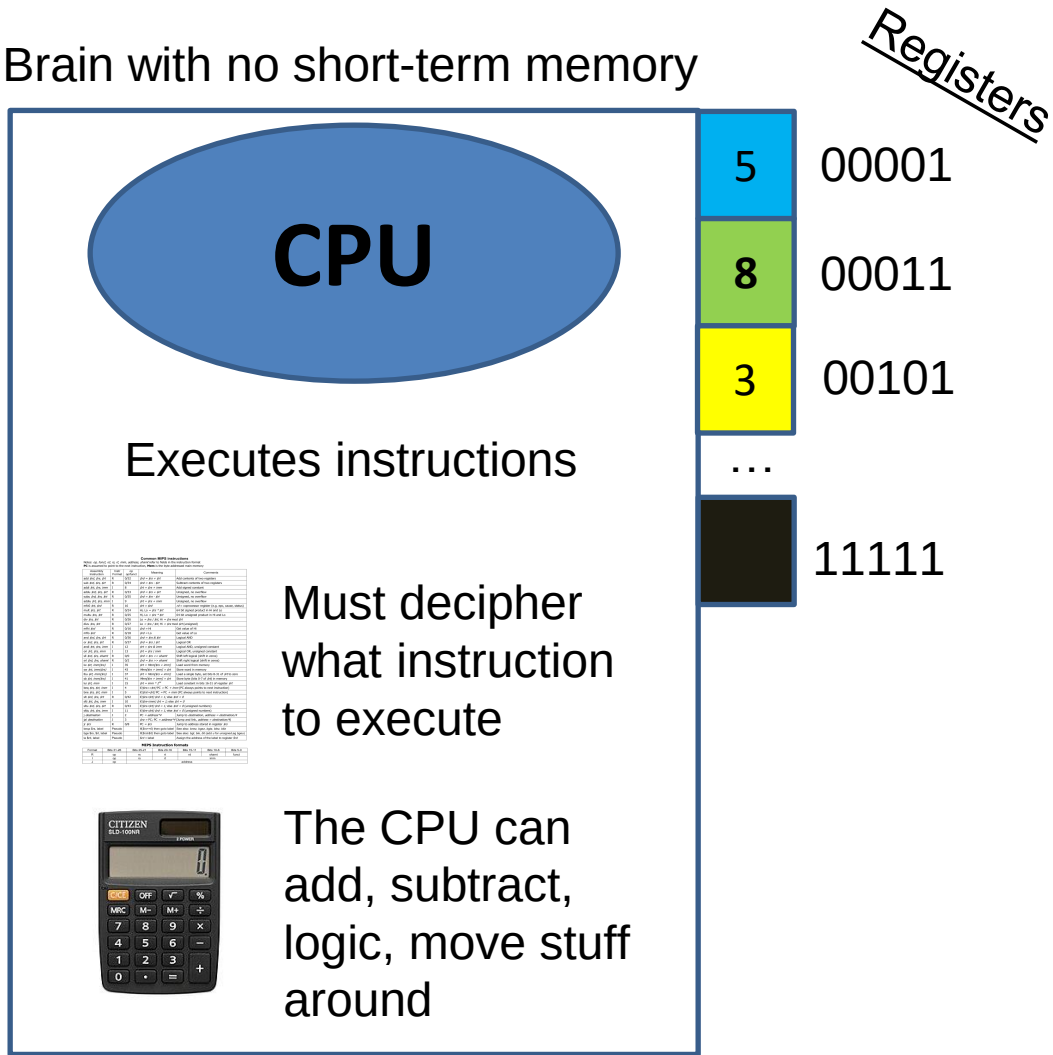
I. Hardware

Brain with no short-term memory



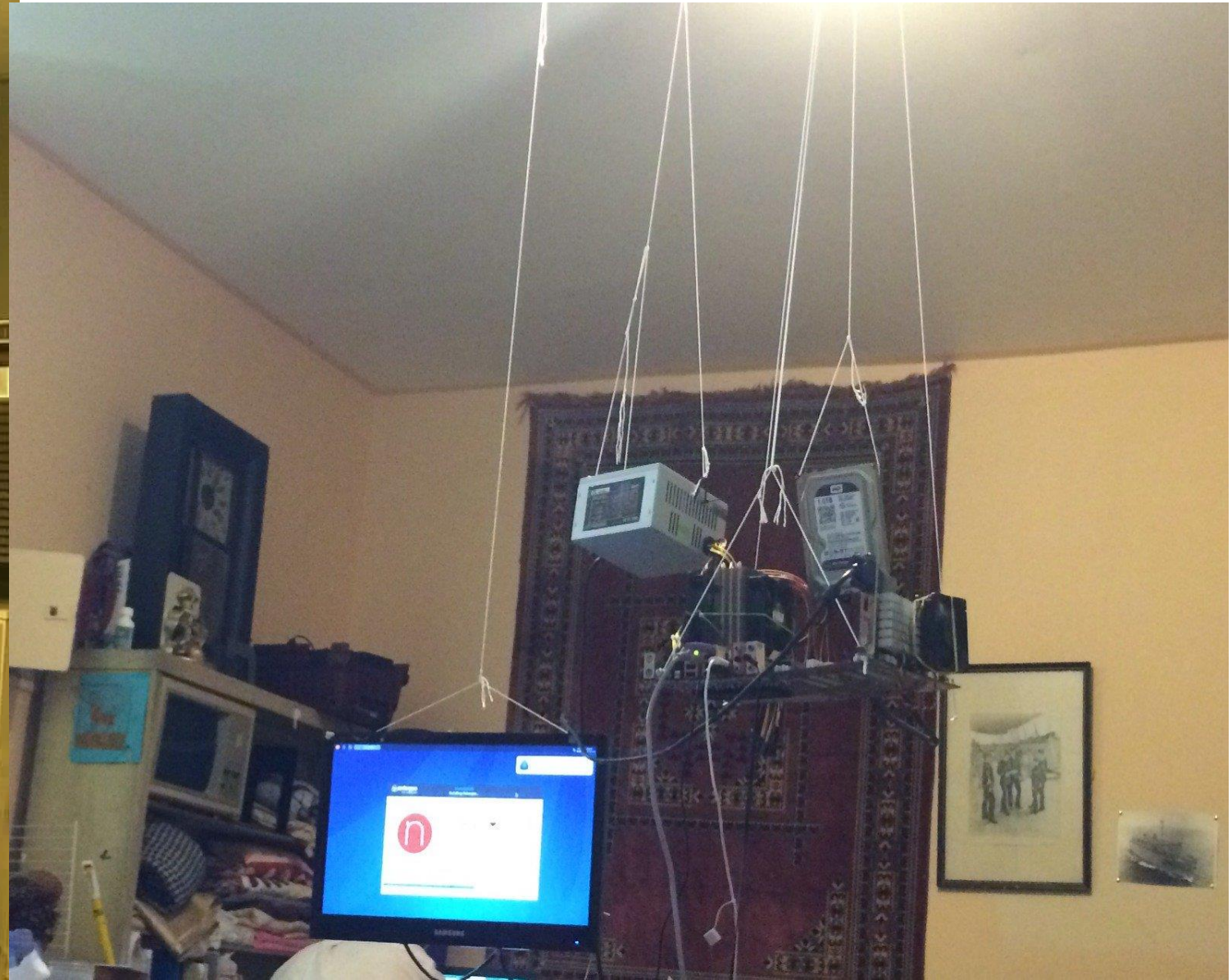
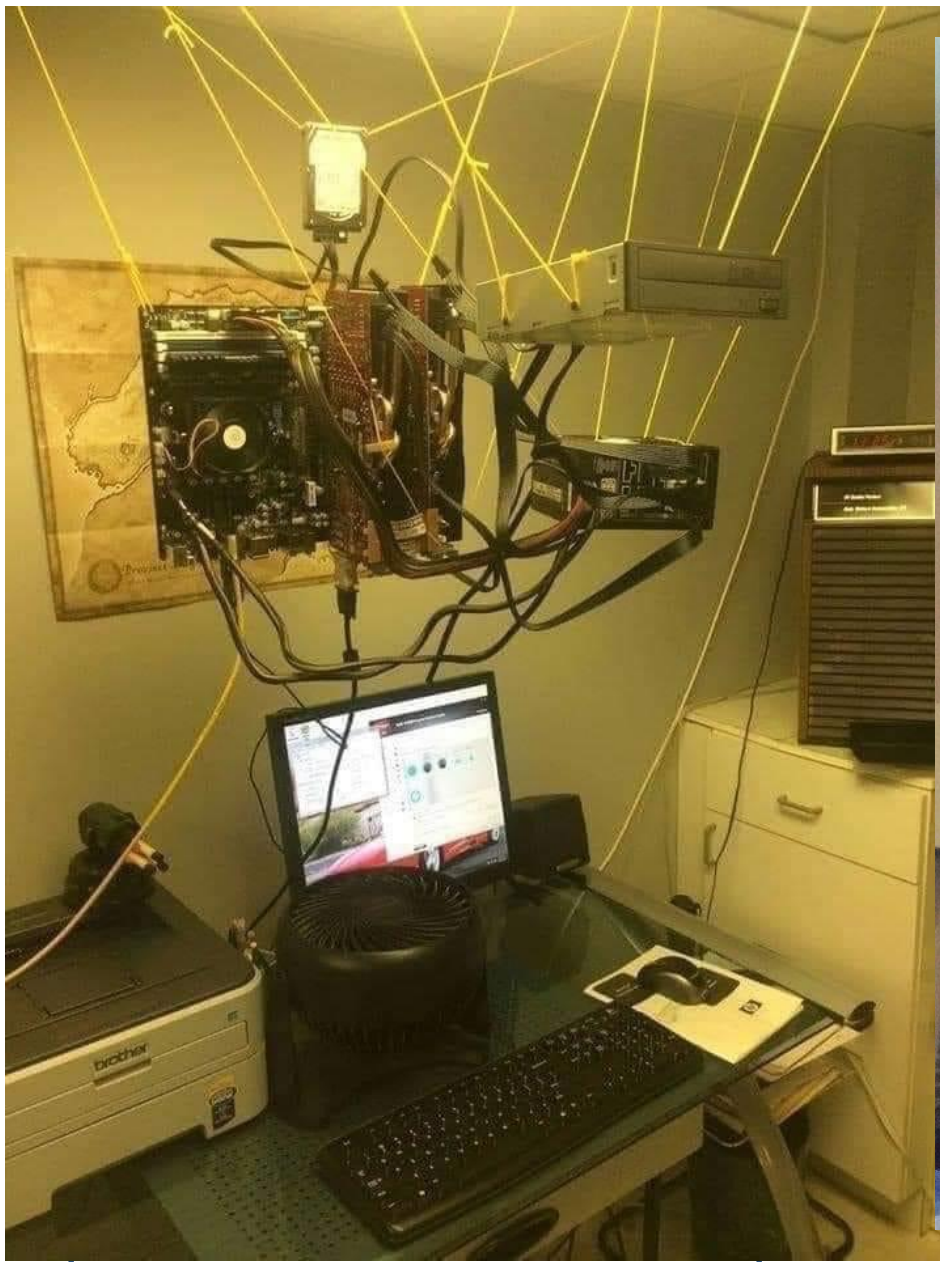
32 bit vs 64 bit?

I. Hardware



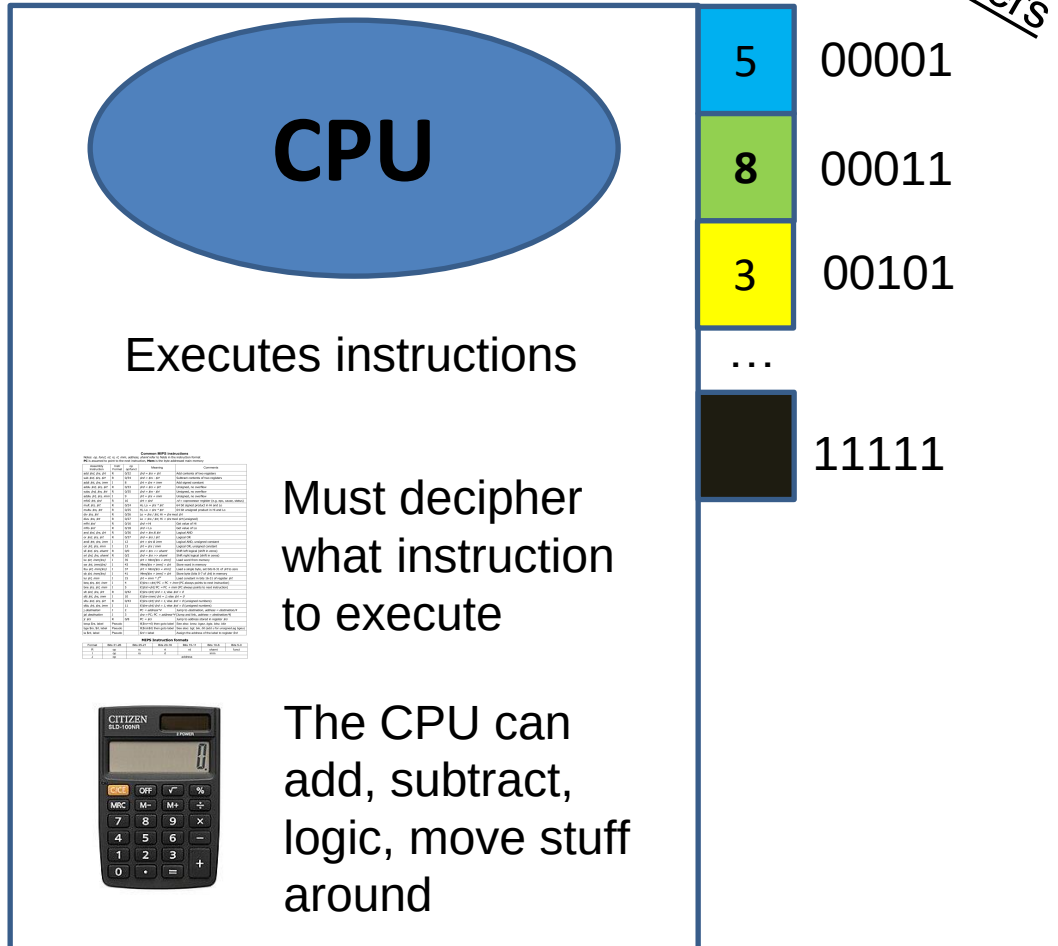
32 bit vs 64 bit?

Parameter	32-bit processors	64-bit processors
Addressable space	It has 4 GB addressable space	64-bit processors have 16 exabytes addressable space
Application support	64-bit applications and programs won't work	32-bit applications and programs will work
OS support	Need a 32-bit operating system.	It can run on 32 and the 64-bit operating system.
Support for multi-tasking	Not an ideal option for stress testing and multi-tasking.	Works best for performing multi-tasking and stress testing.
OS and CPU requirement	32-bit operating systems and applications require 32-bit CPUs	64-bit OS demands 64-bit CPU, and 64-bit applications require 64-bit OS and CPU.
System available	Support Windows 7, 8 Vista, XP, and, Linux.	Windows XP Professional, Windows Vista, Windows 7, Windows 8, Windows 10, Linux, and Mac OS X.

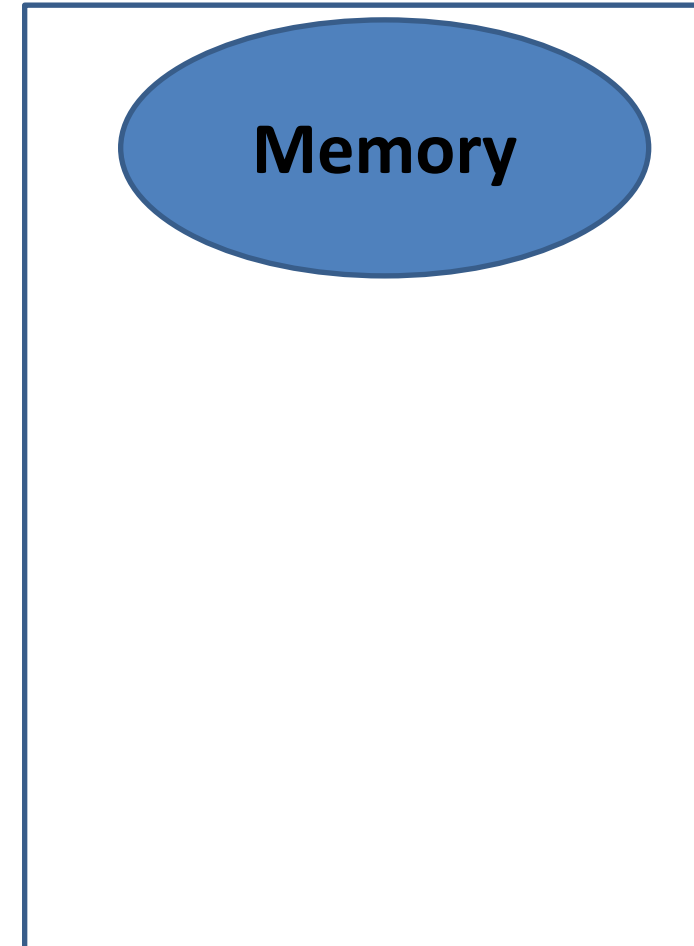


I. Hardware

Brain with no short-term memory

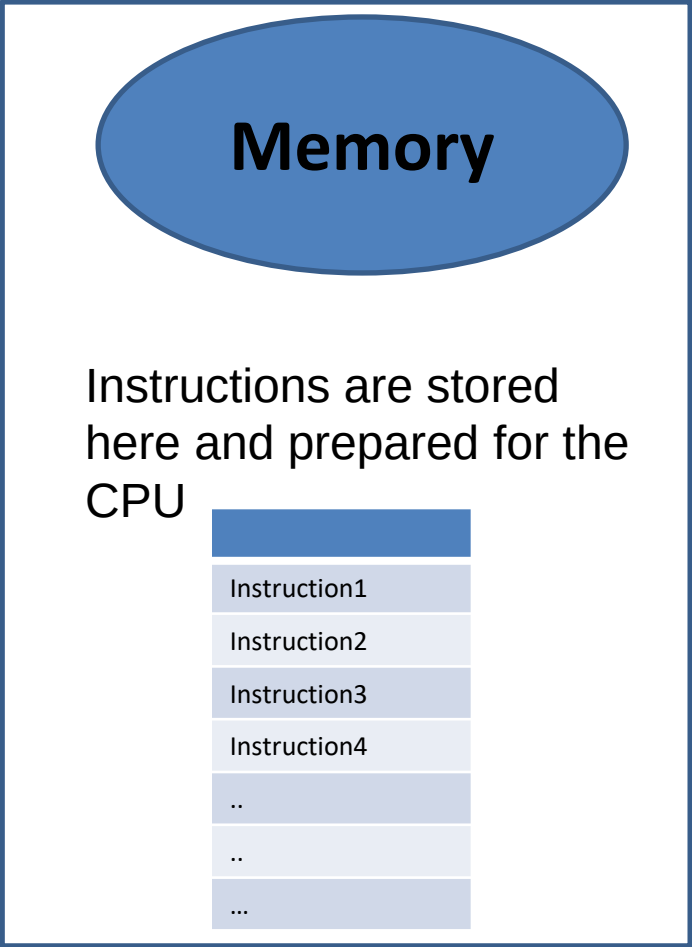


Scratch Pad



I. Hardware

Scratch Pad



Memory

Instructions are stored
here and prepared for the
CPU

Instruction1

Instruction2

Instruction3

Instruction4

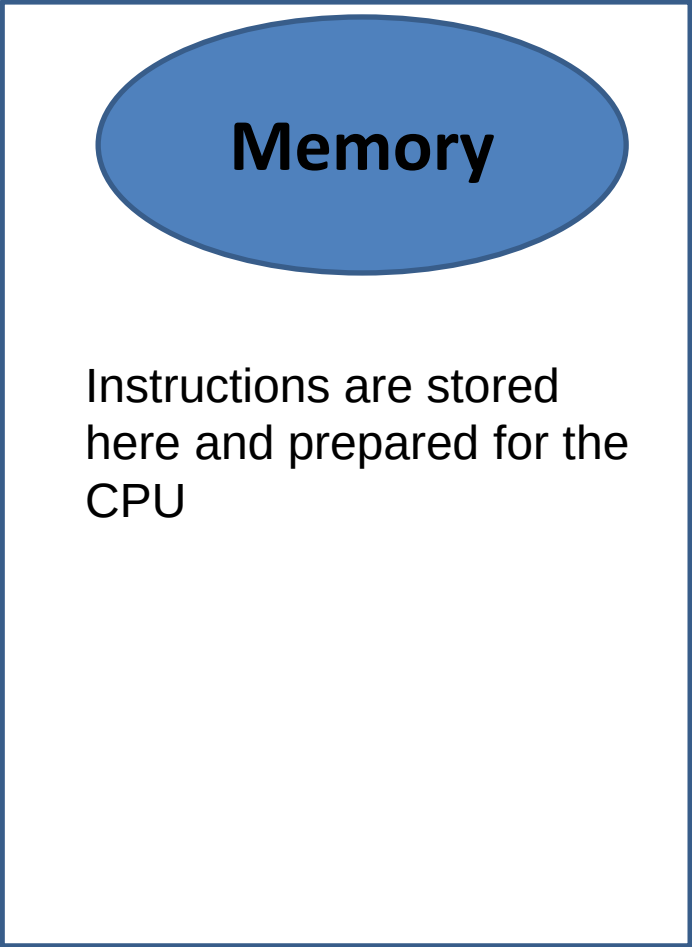
..

..

...

I. Hardware

Scratch Pad



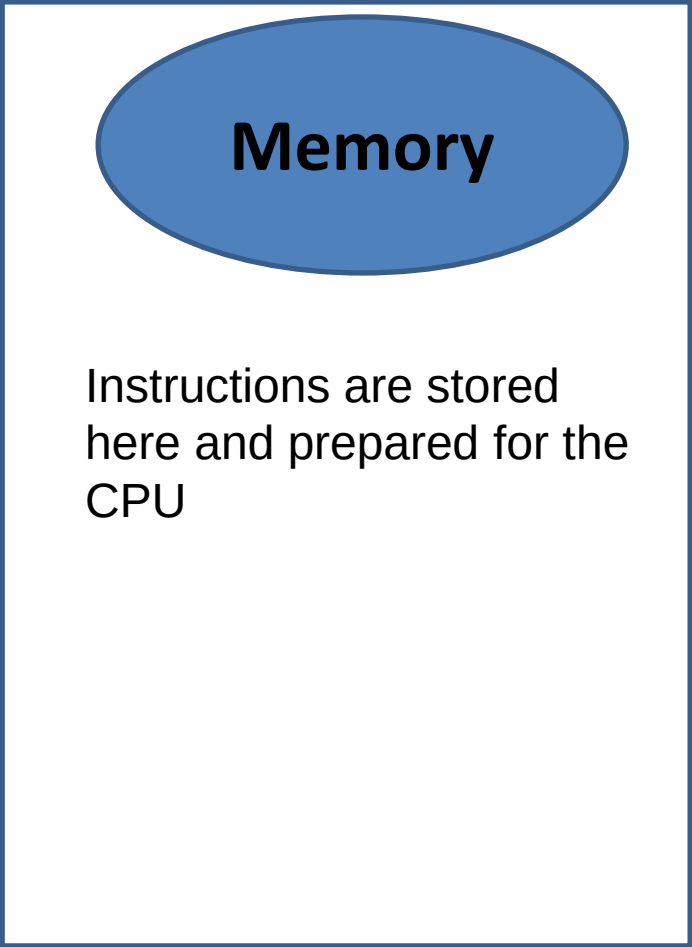
Memory

Instructions are stored
here and prepared for the
CPU

When computer programs are executed, their instructions will eventually get stored in memory

I. Hardware

Scratch Pad



Memory

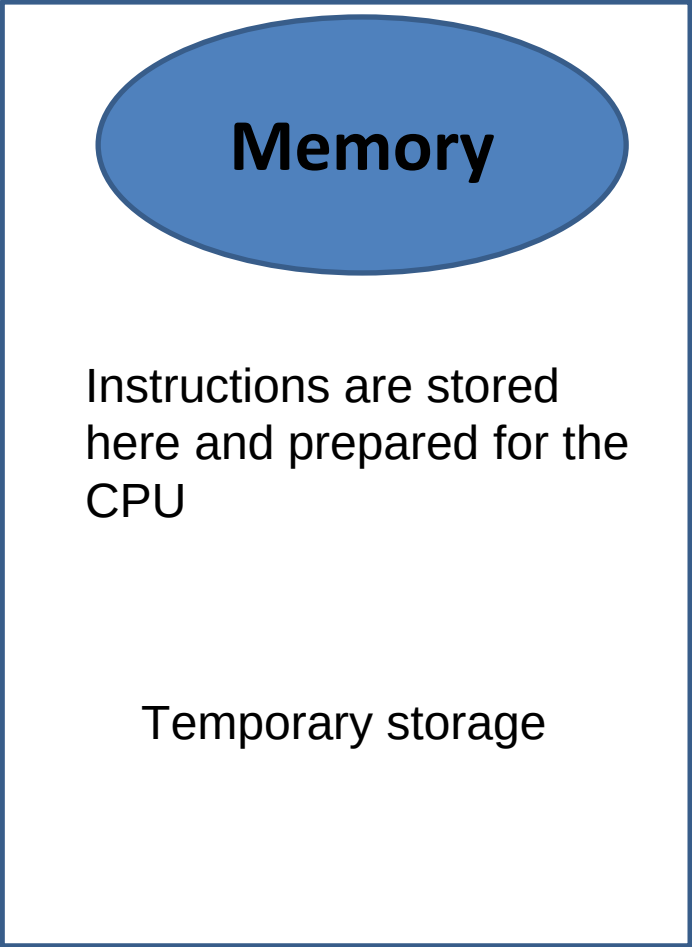
Instructions are stored
here and prepared for the
CPU

When computer programs are executed, their
instructions will eventually get stored in memory

Permanently?!?!?

I. Hardware

Scratch Pad



The diagram shows a large rectangular box with a blue border. Inside this box, at the top, is a blue oval with the word "Memory" in bold black text. Below the oval, the text "Instructions are stored here and prepared for the CPU" is written. At the bottom of the box, the text "Temporary storage" is written.

Memory

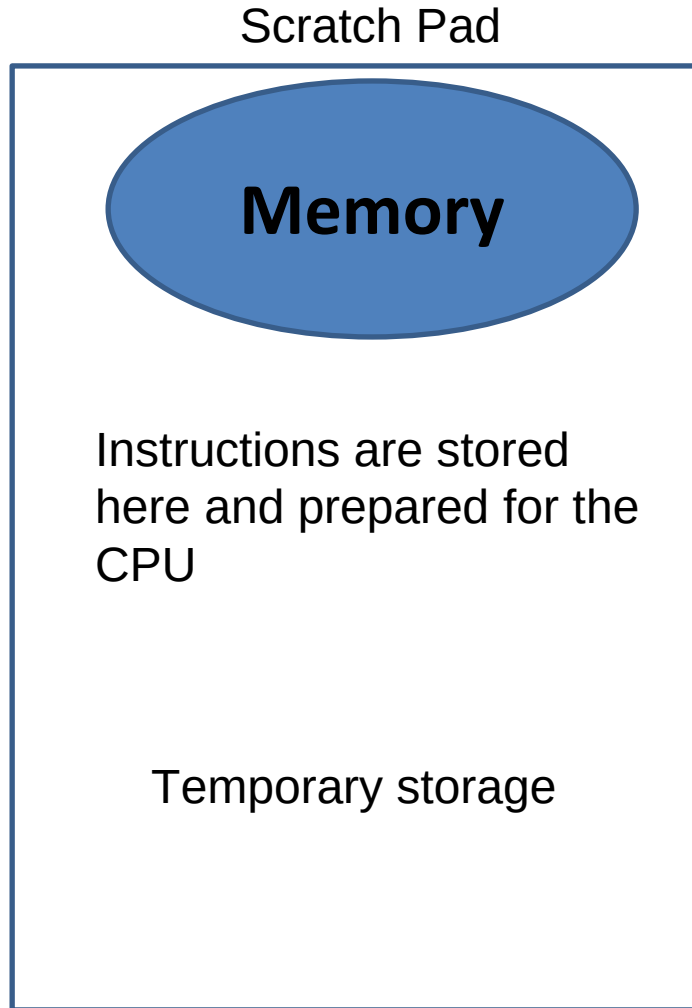
Instructions are stored here and prepared for the CPU

Temporary storage

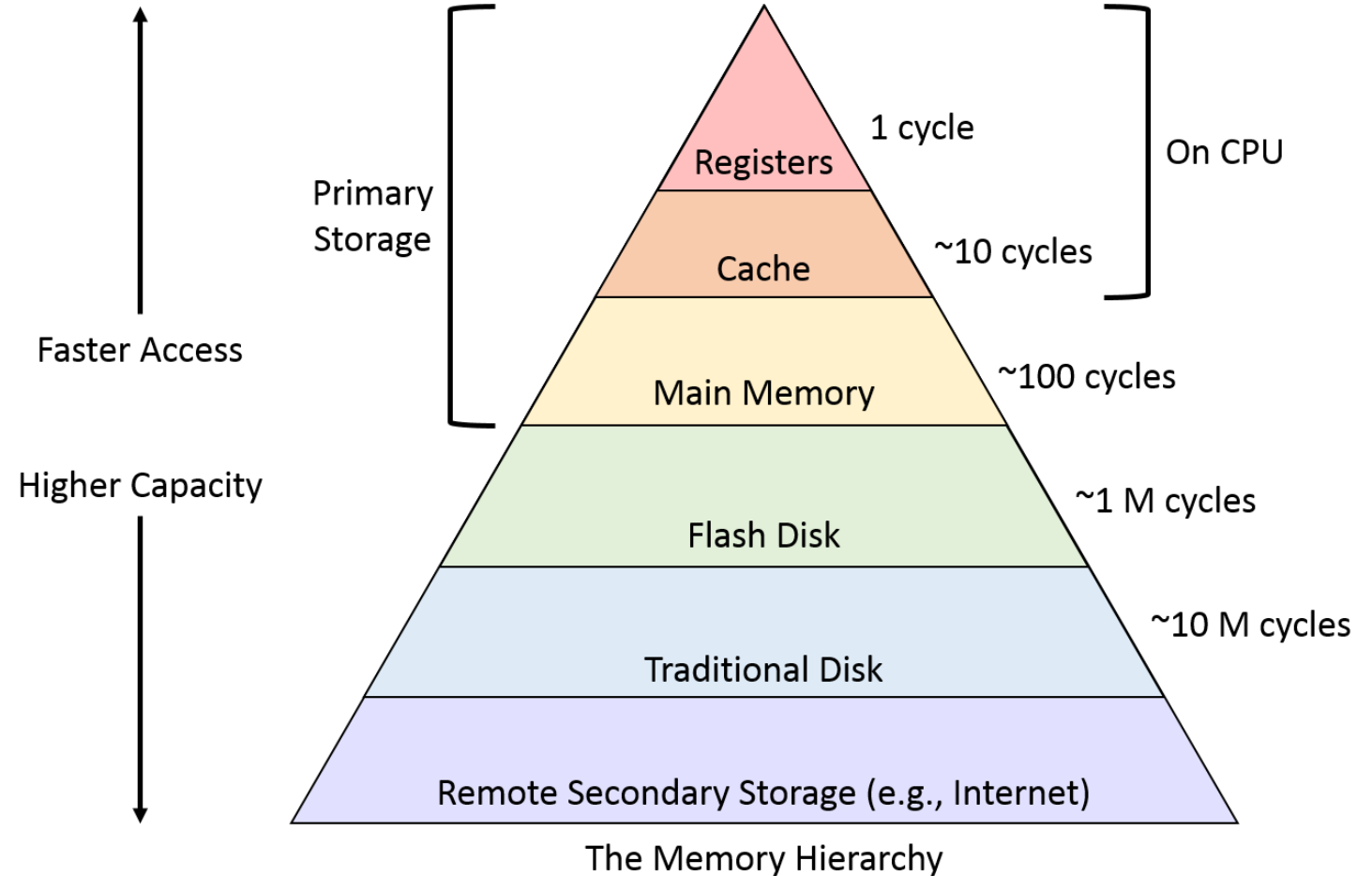
When computer programs are executed, their instructions will eventually get stored in memory

Main memory is **volatile**

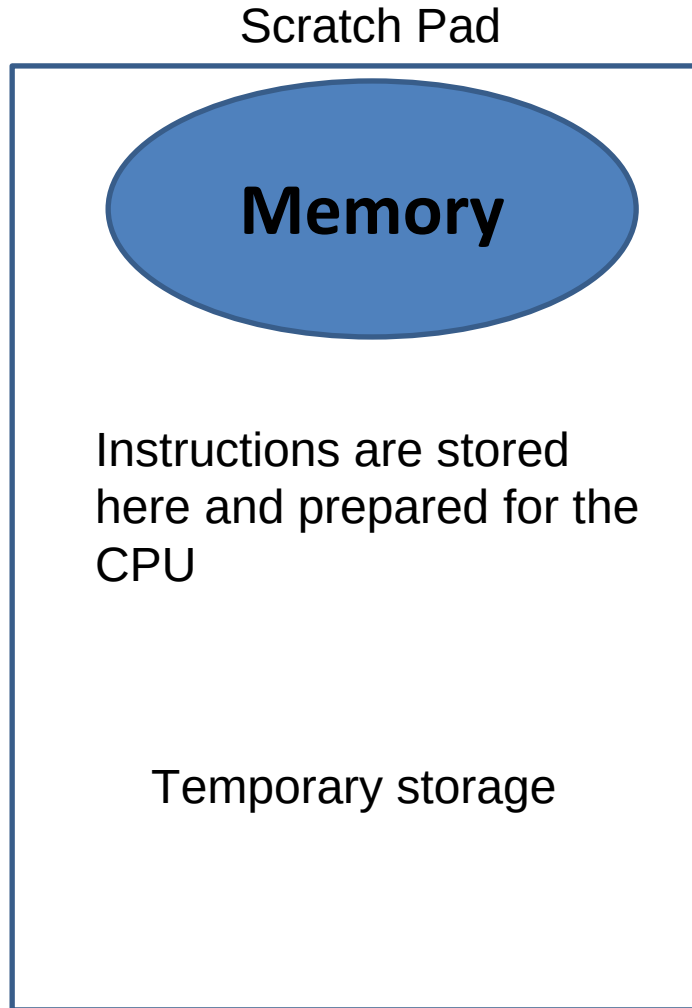
I. Hardware



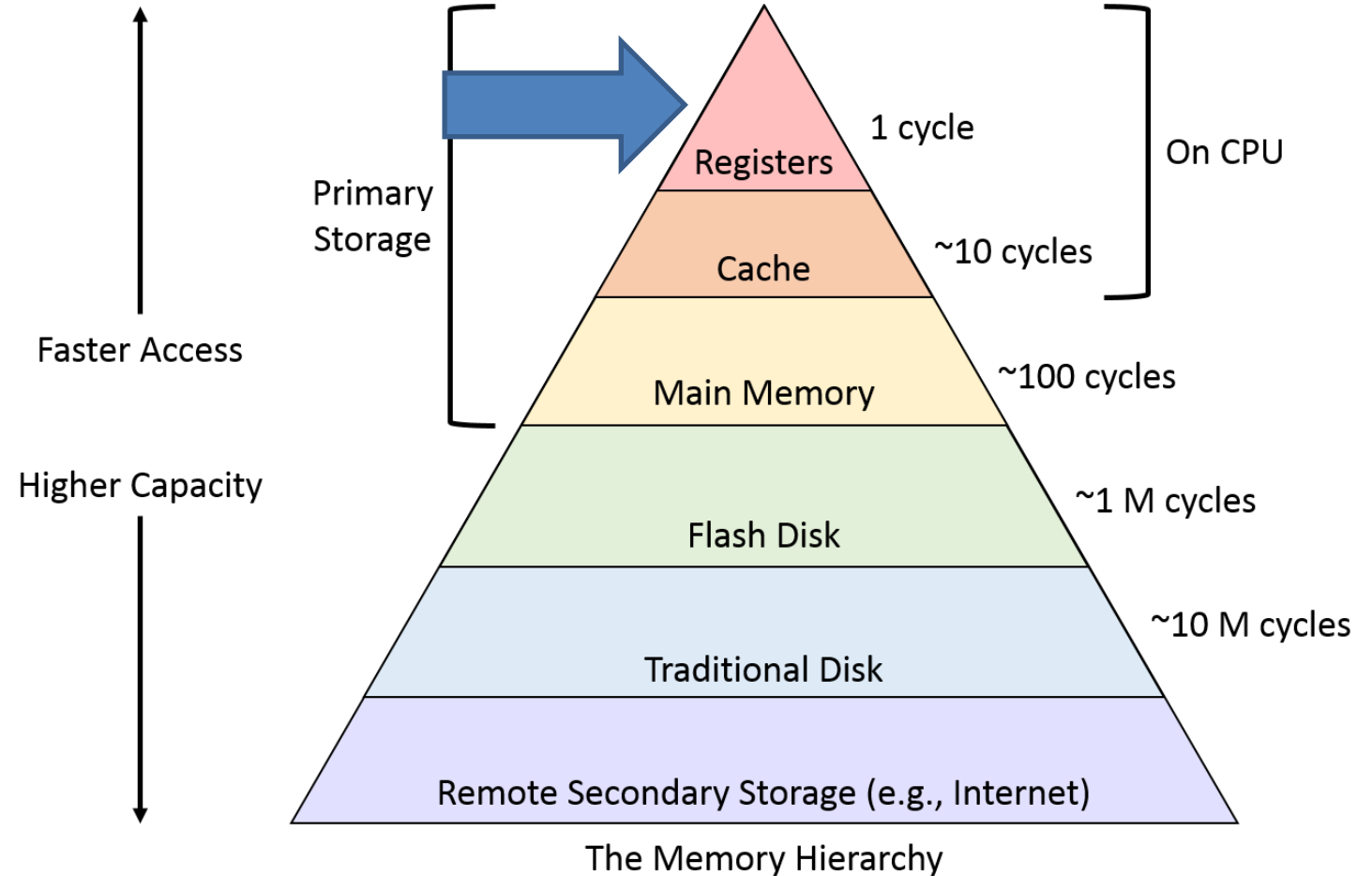
The Memory Hierarchy



I. Hardware



The Memory Hierarchy



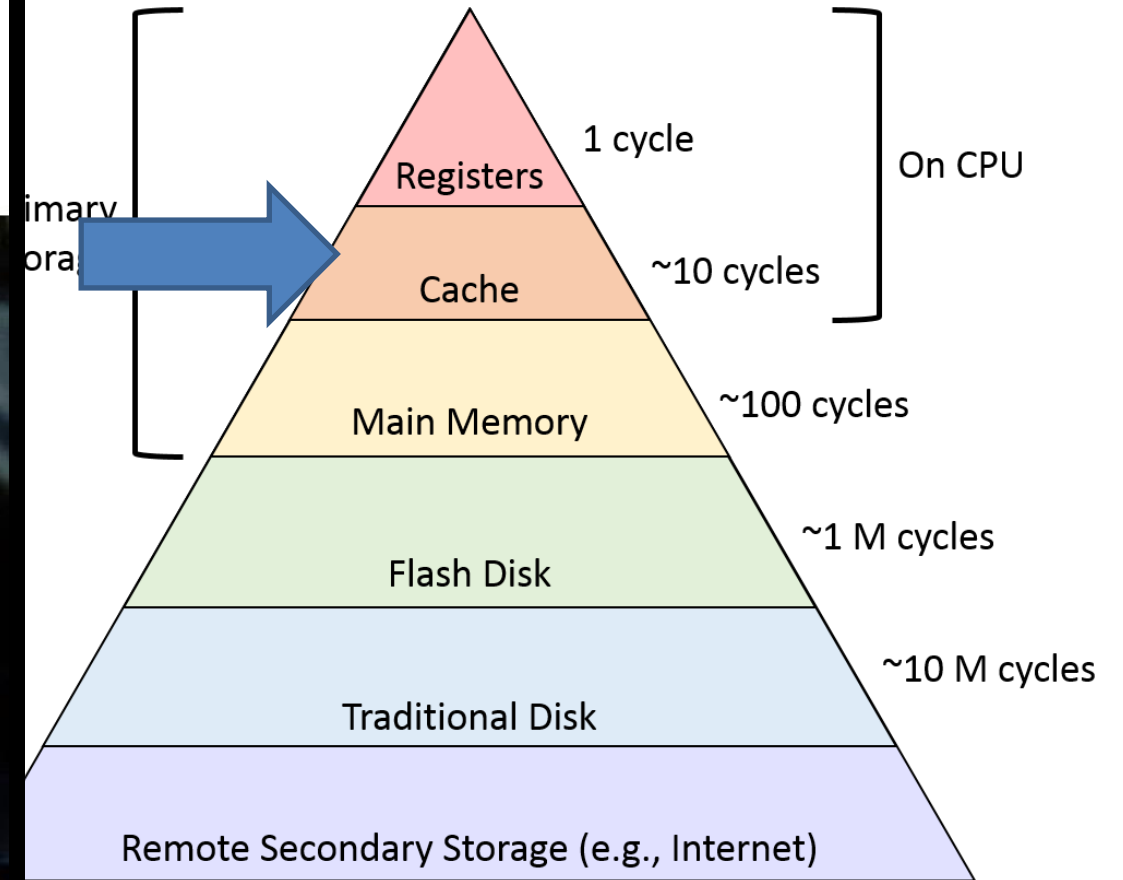
I. Hardware

My CPU when the L1 cache misses



This little maneuver is gonna cost us 3 nanoseconds

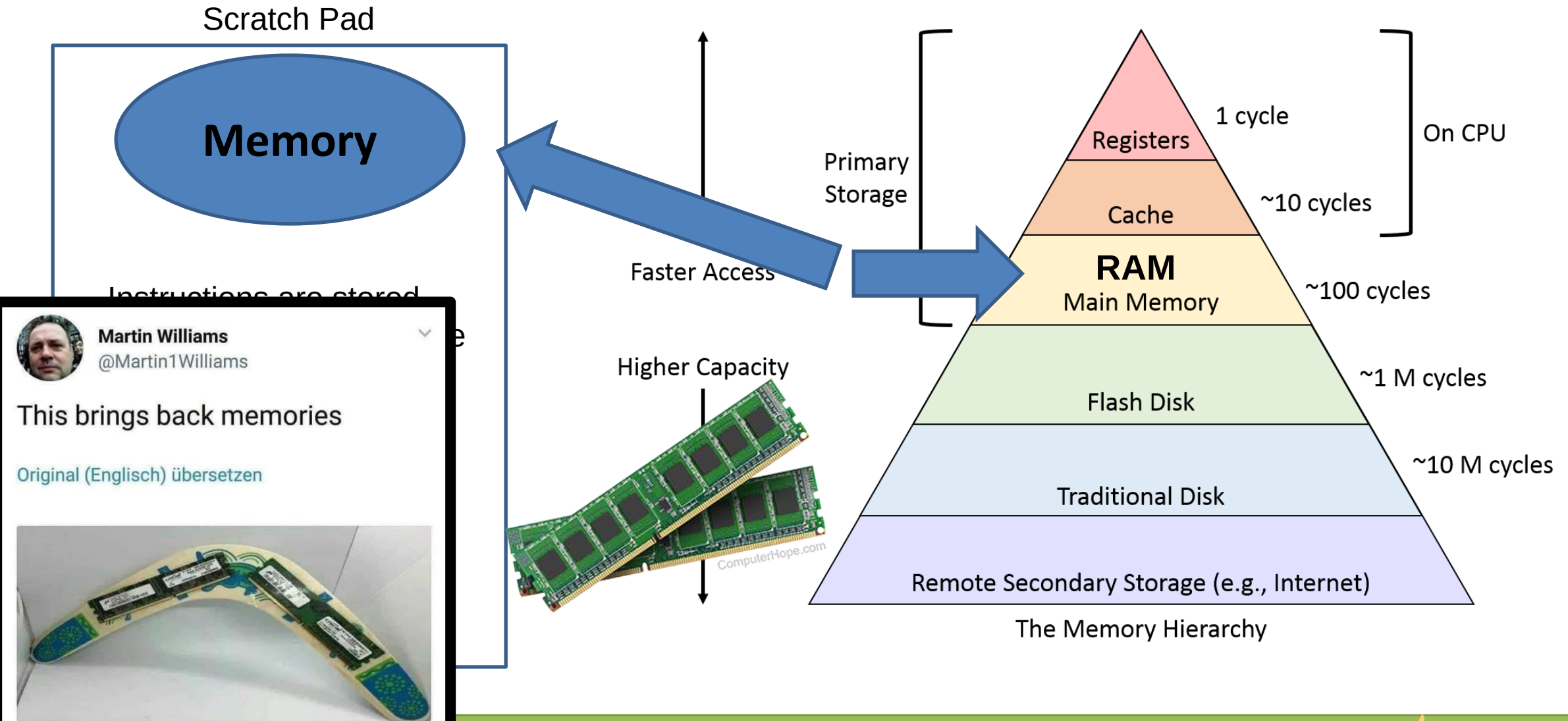
The Memory Hierarchy



The Memory Hierarchy

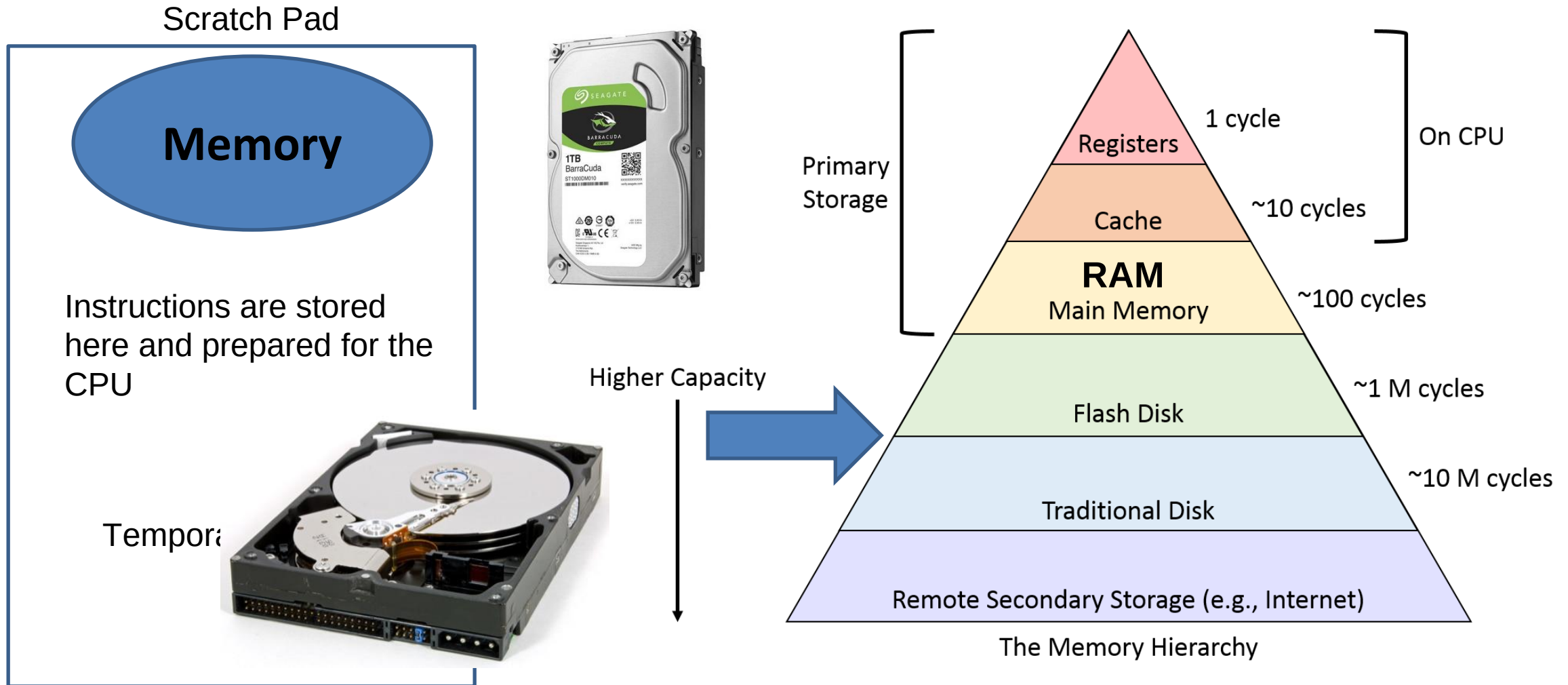
I. Hardware

The Memory Hierarchy

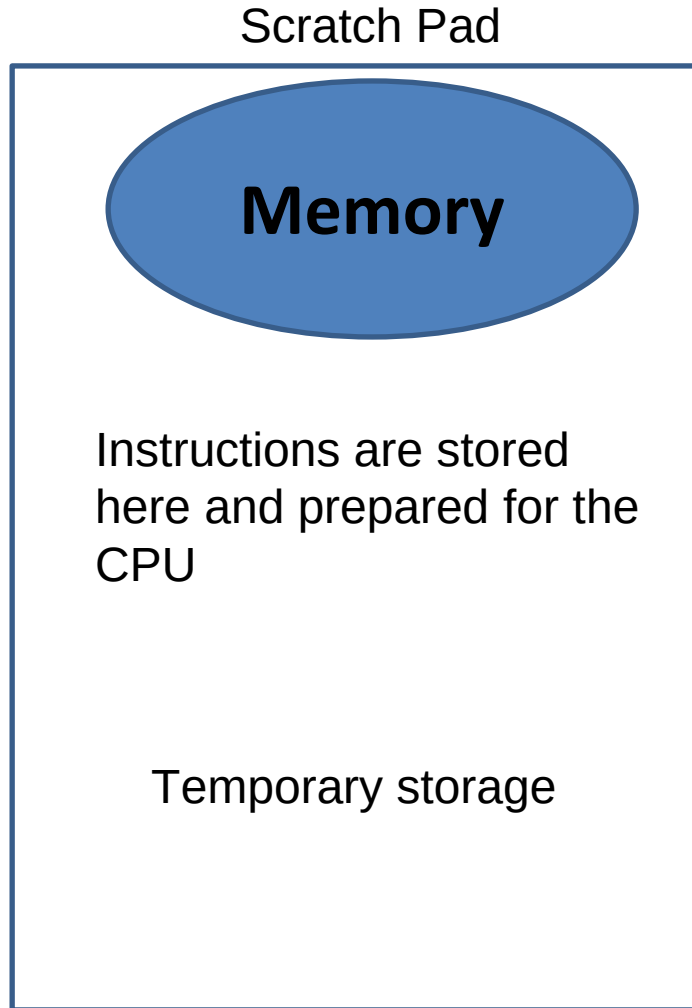


I. Hardware

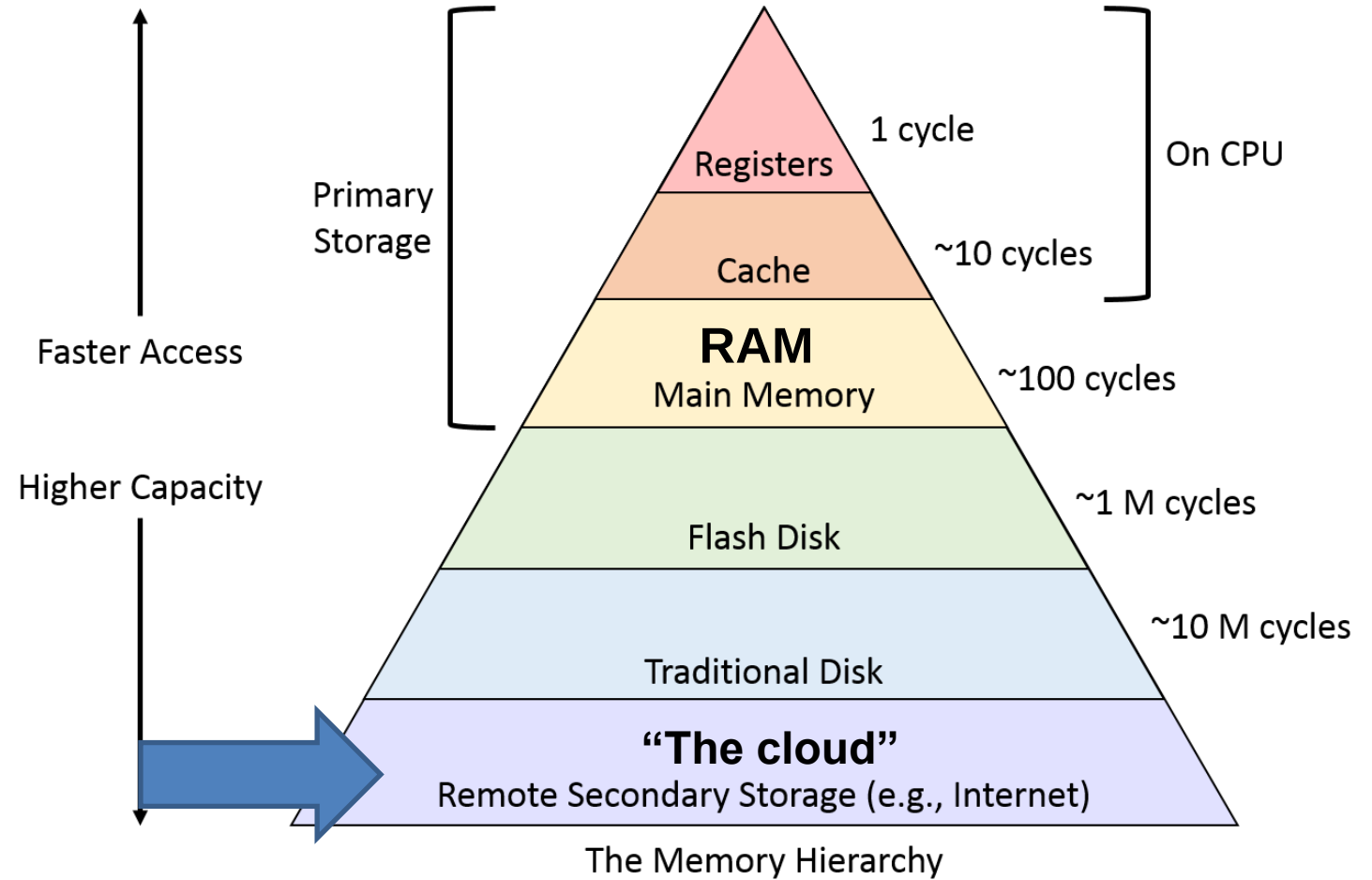
The Memory Hierarchy



I. Hardware



The Memory Hierarchy



I. Hardware



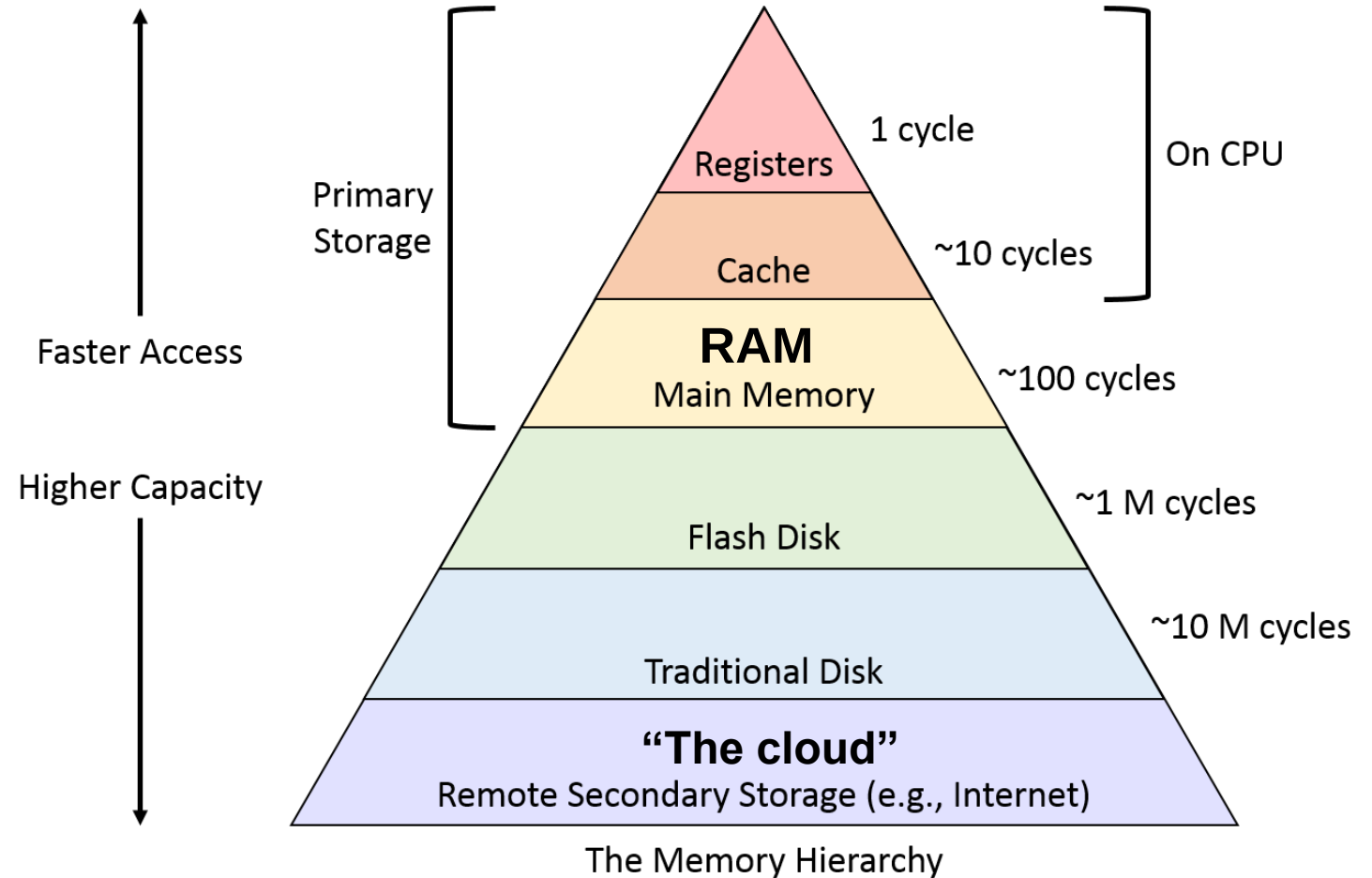
Scratch Pad

Memory

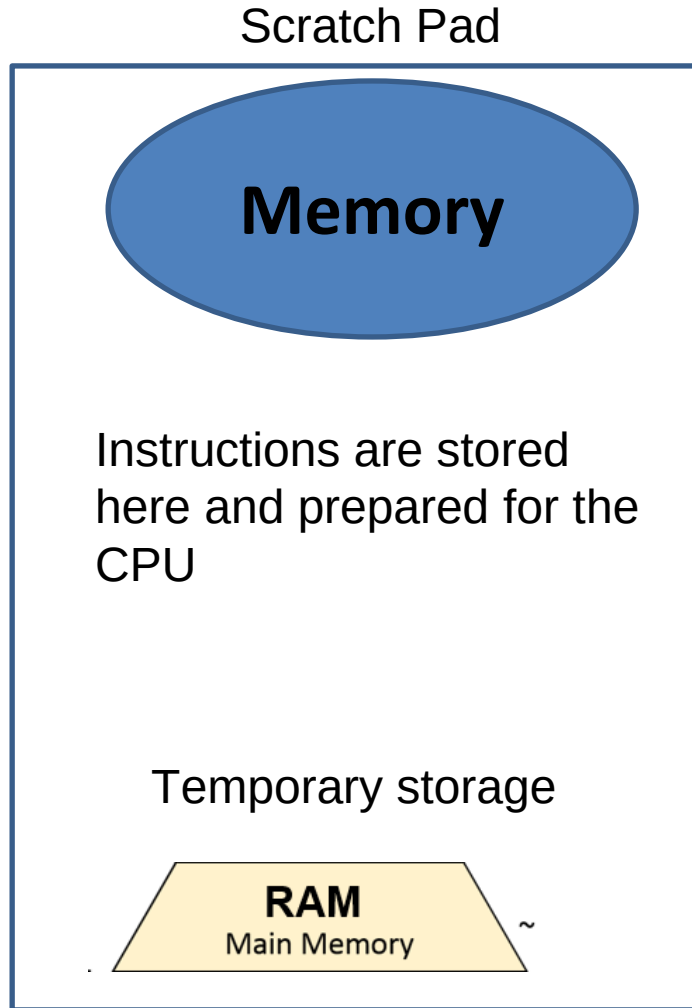
Instructions are stored here and prepared for the CPU

Temporary storage

The Memory Hierarchy

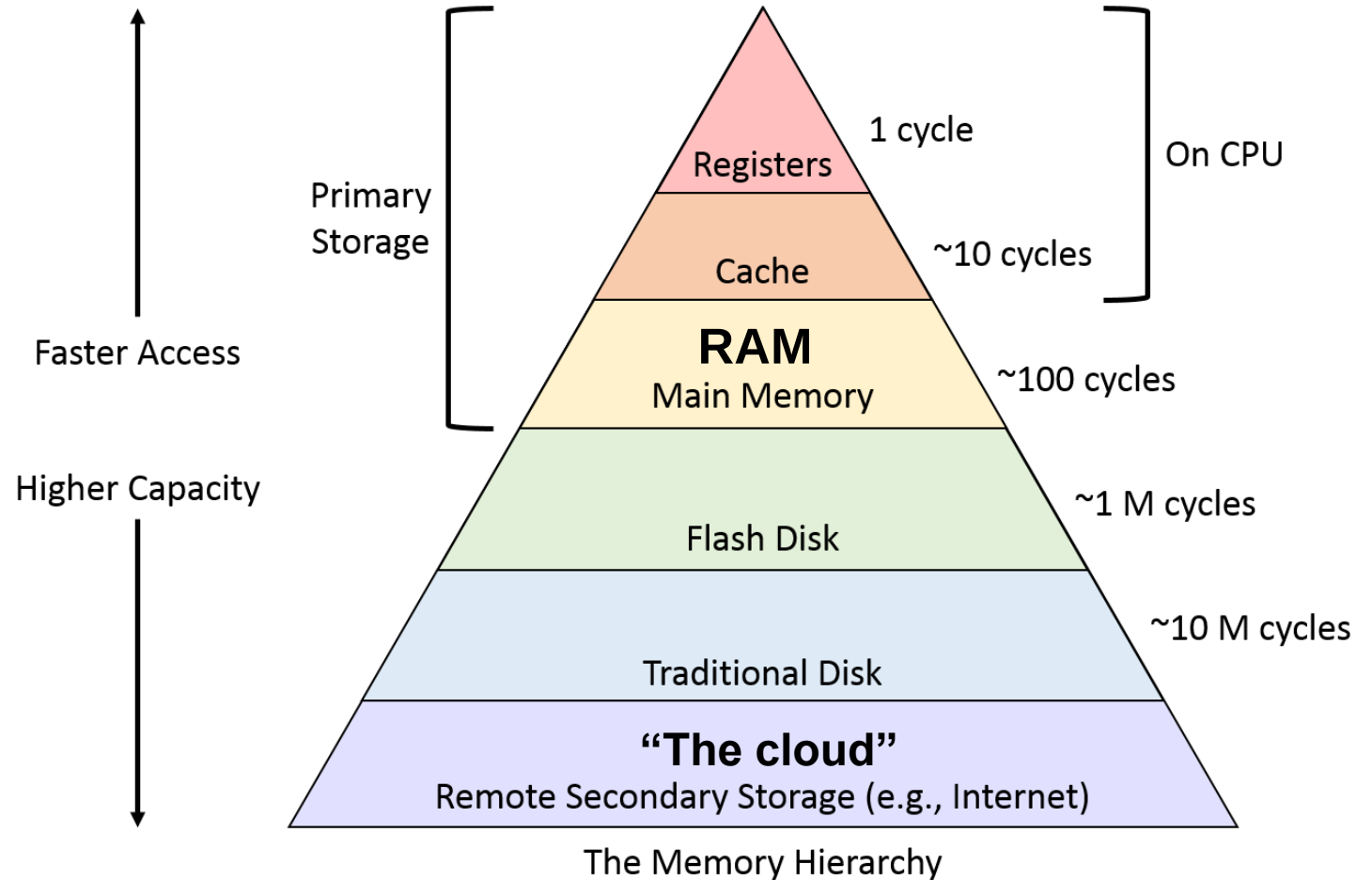


I. Hardware



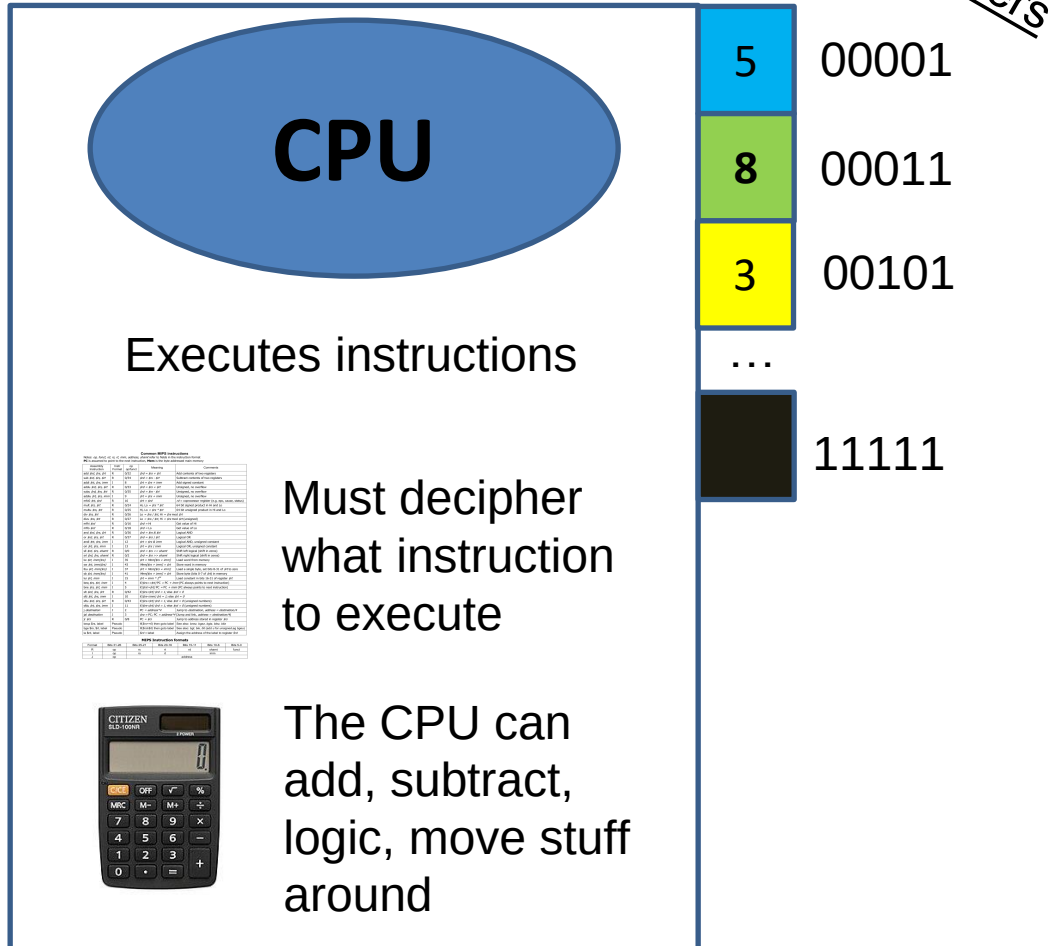
*Why not use
memory/registers
for everything??*

The Memory Hierarchy

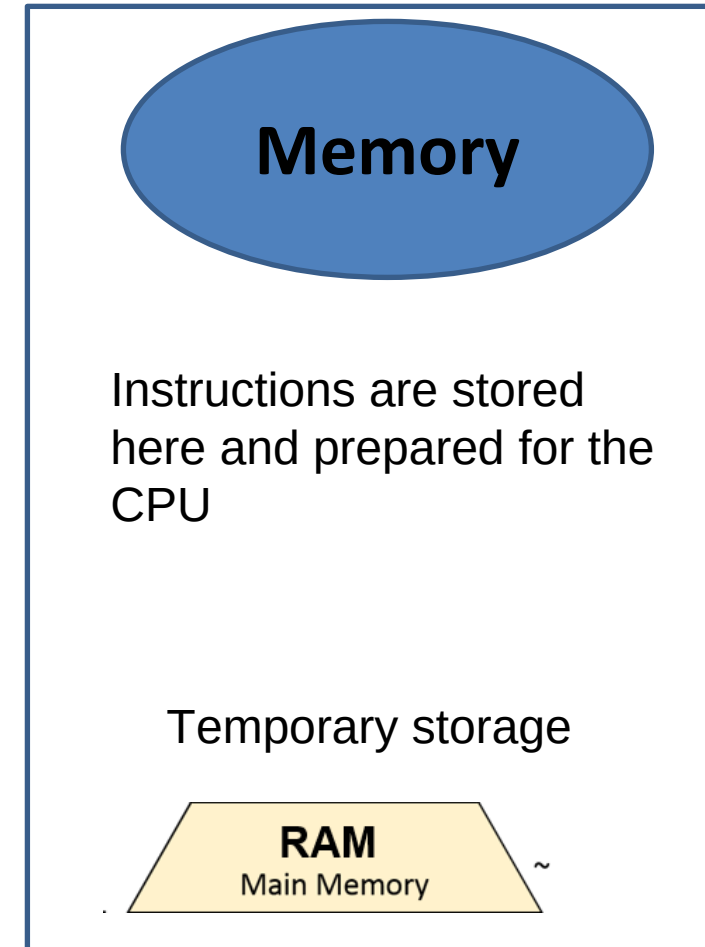


I. Hardware

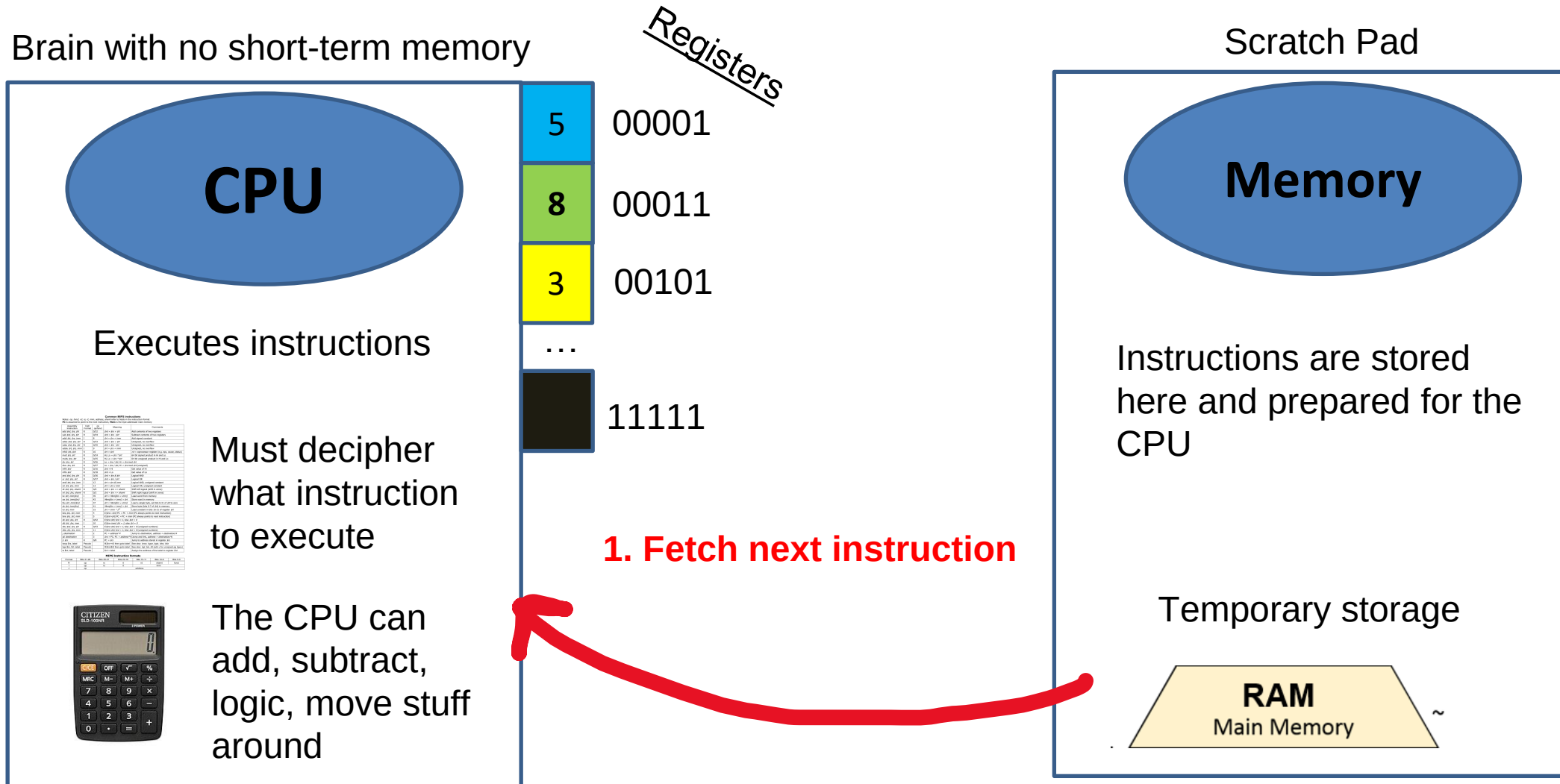
Brain with no short-term memory



Scratch Pad

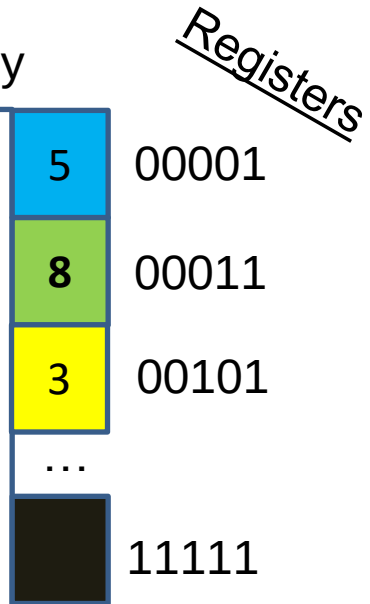
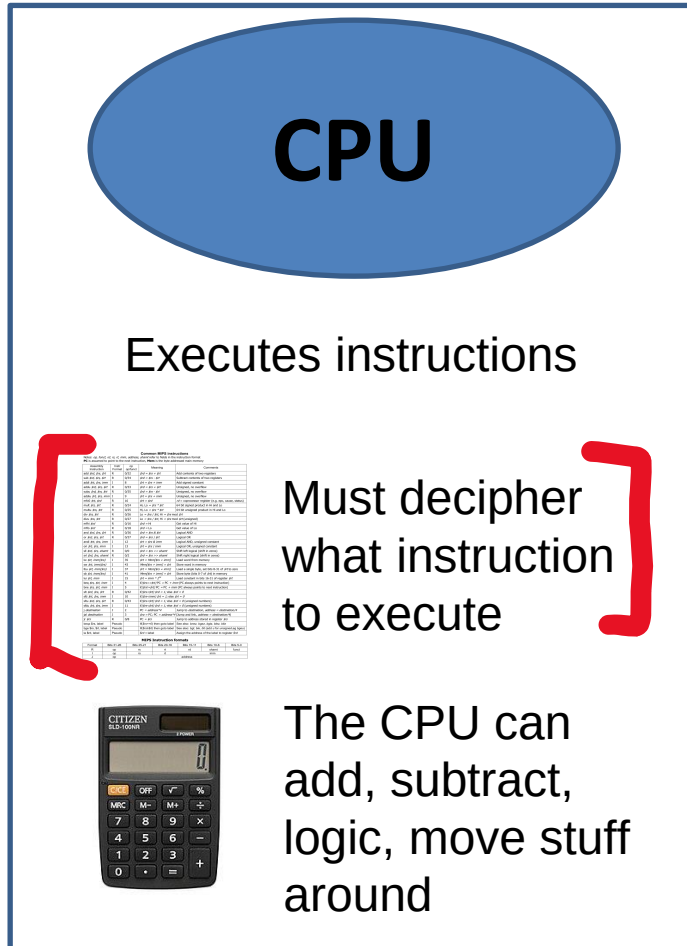


I. Hardware



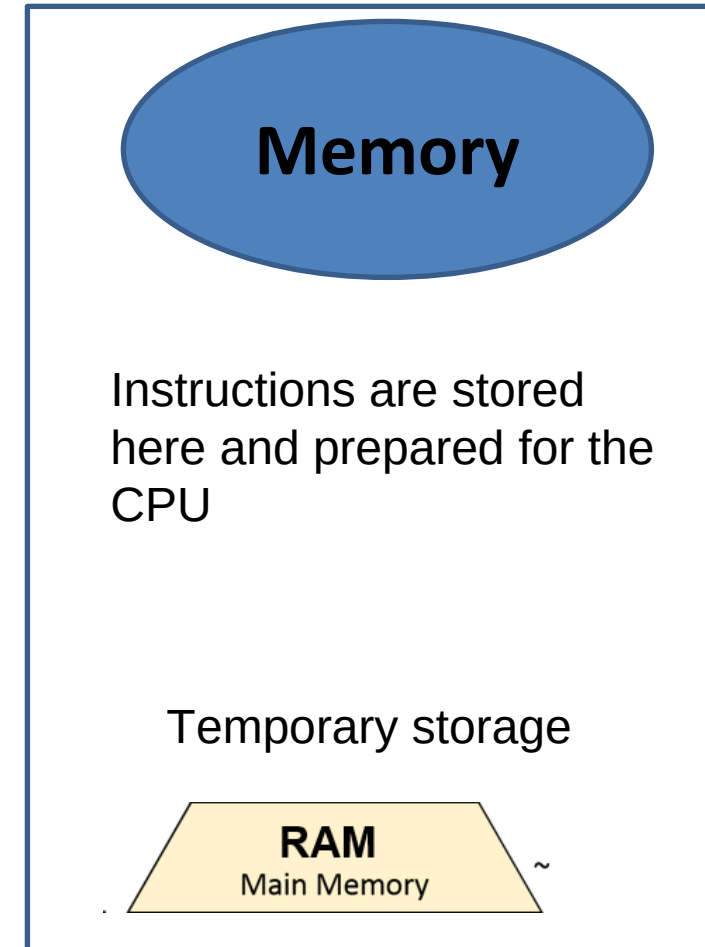
I. Hardware

Brain with no short-term memory



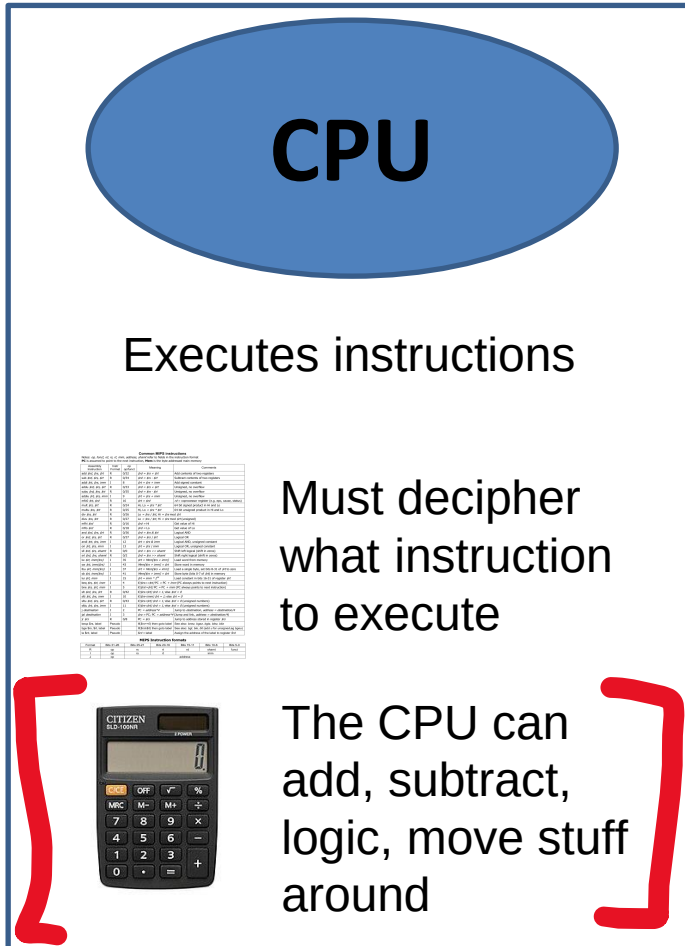
1. Fetch next instruction
2. Decode Instruction

Scratch Pad



I. Hardware

Brain with no short-term memory

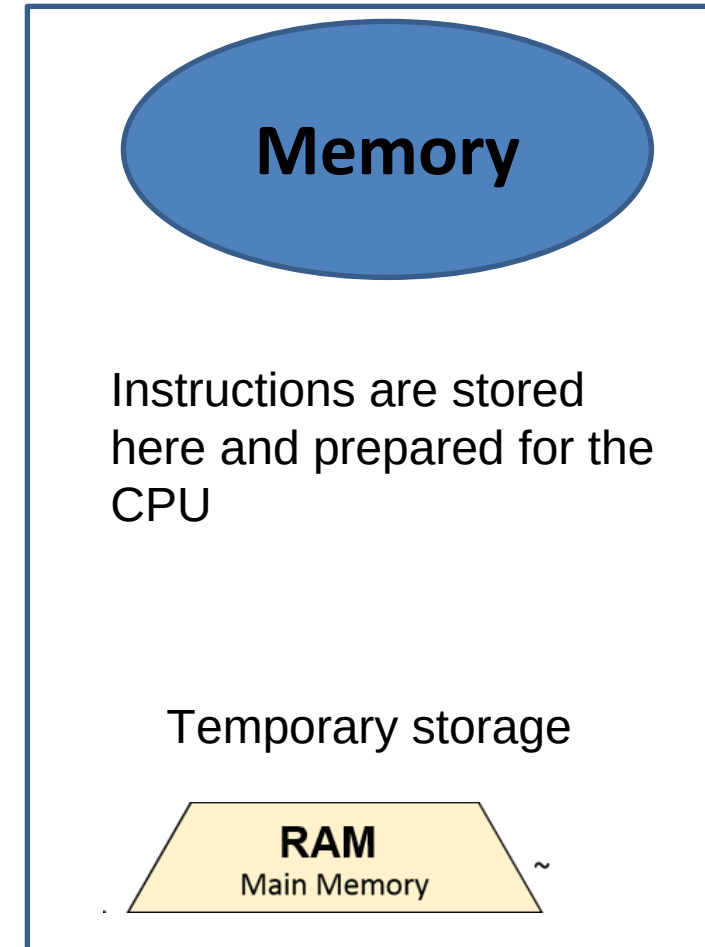


Registers

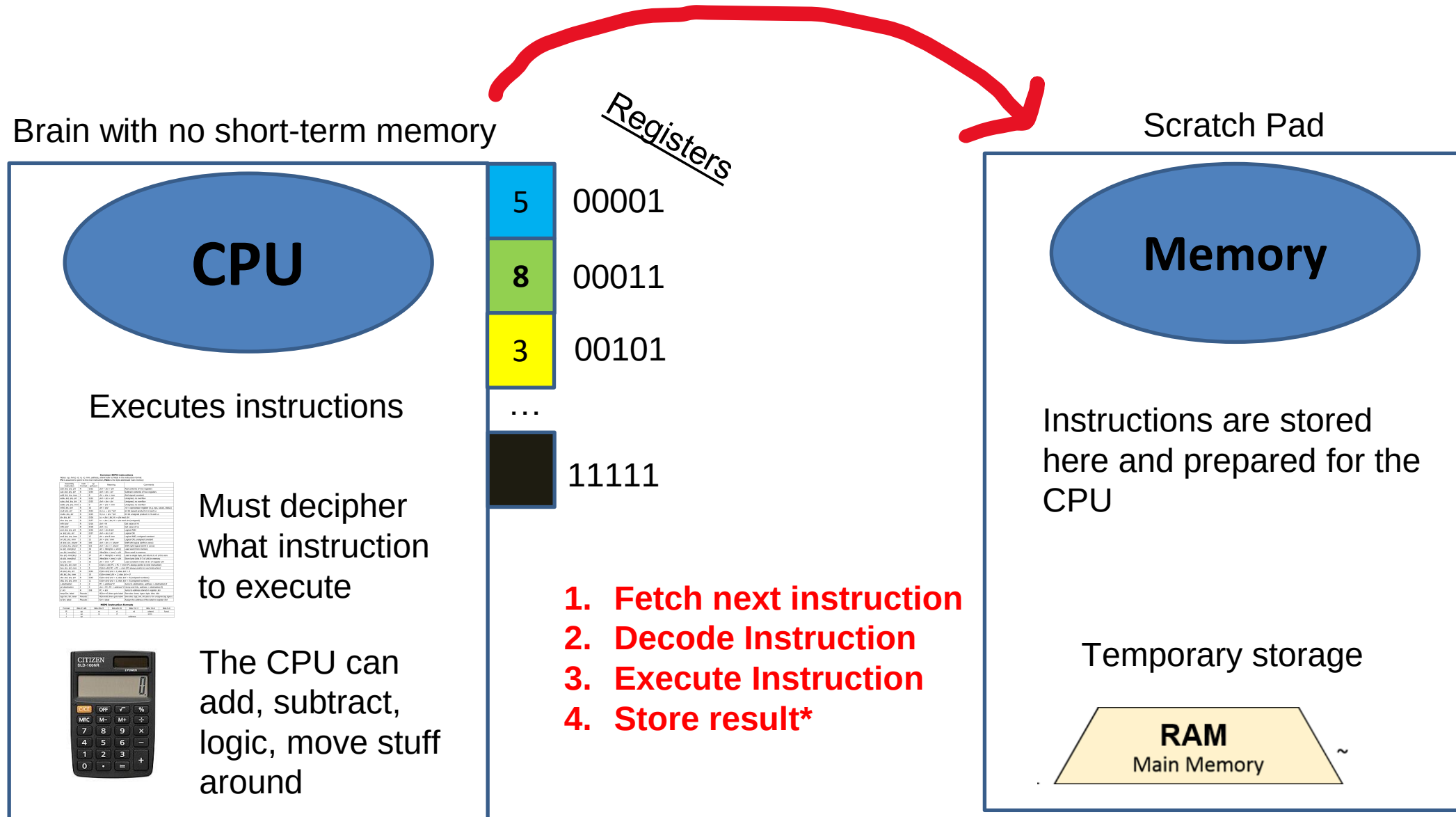
5	00001
8	00011
3	00101
...	
	11111

1. **Fetch next instruction**
2. **Decode Instruction**
3. **Execute Instruction**

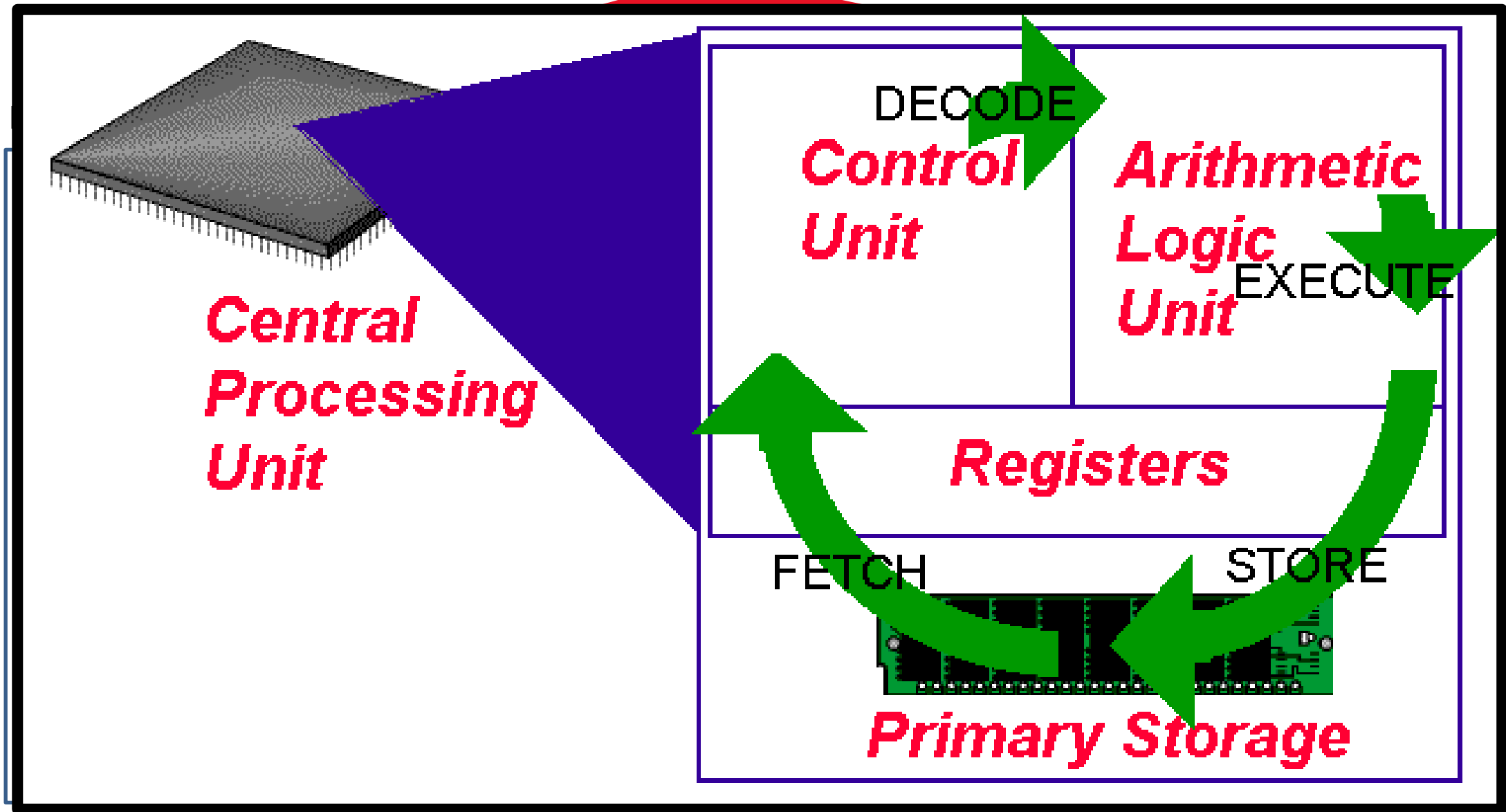
Scratch Pad



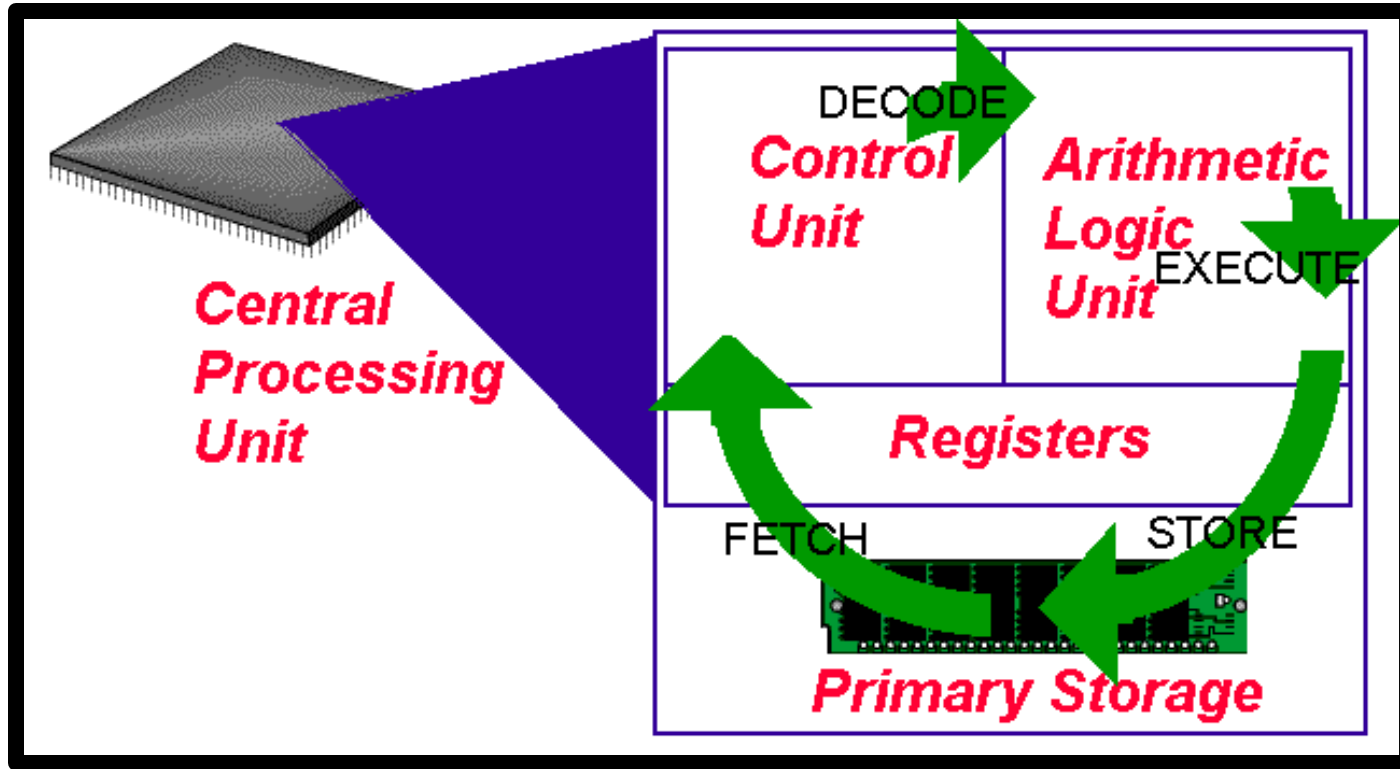
I. Hardware



I. Hardware

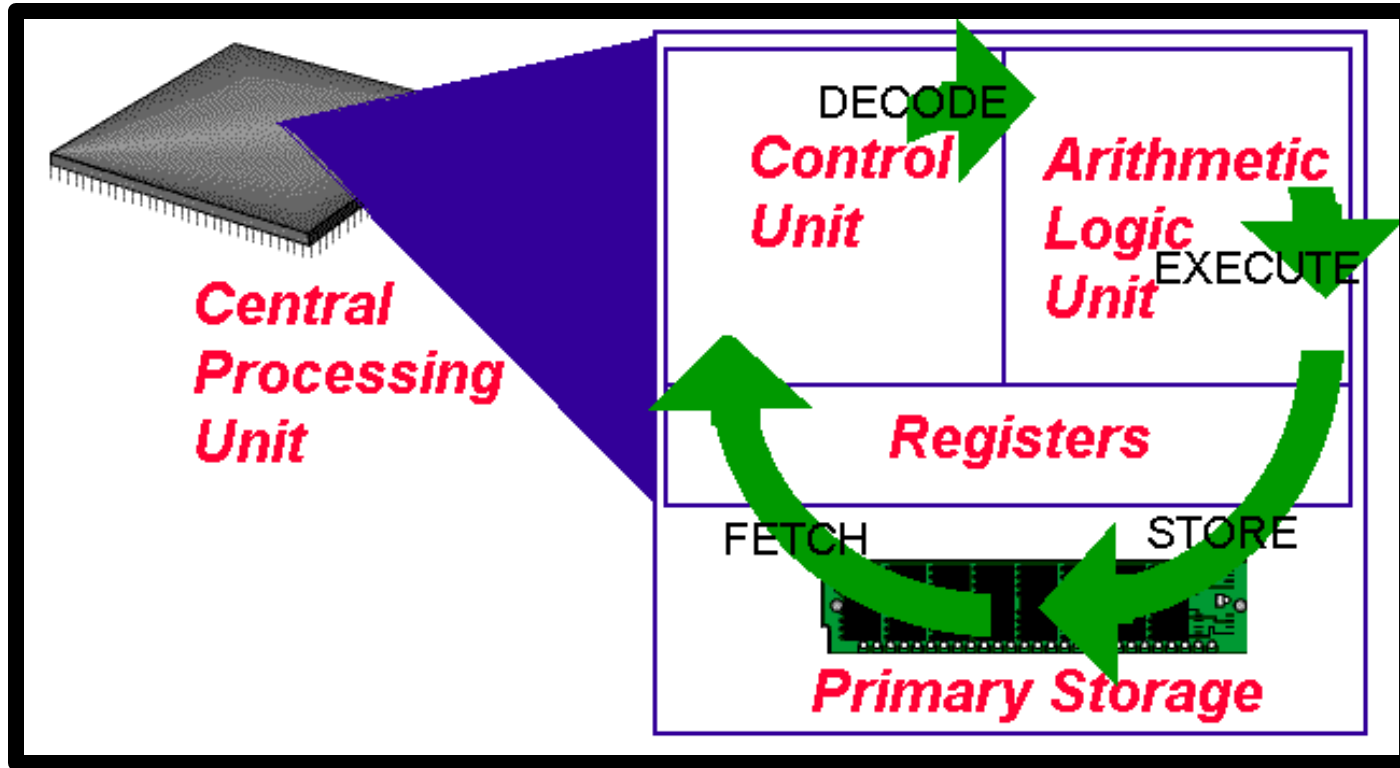


I. Hardware



This process happens really fast

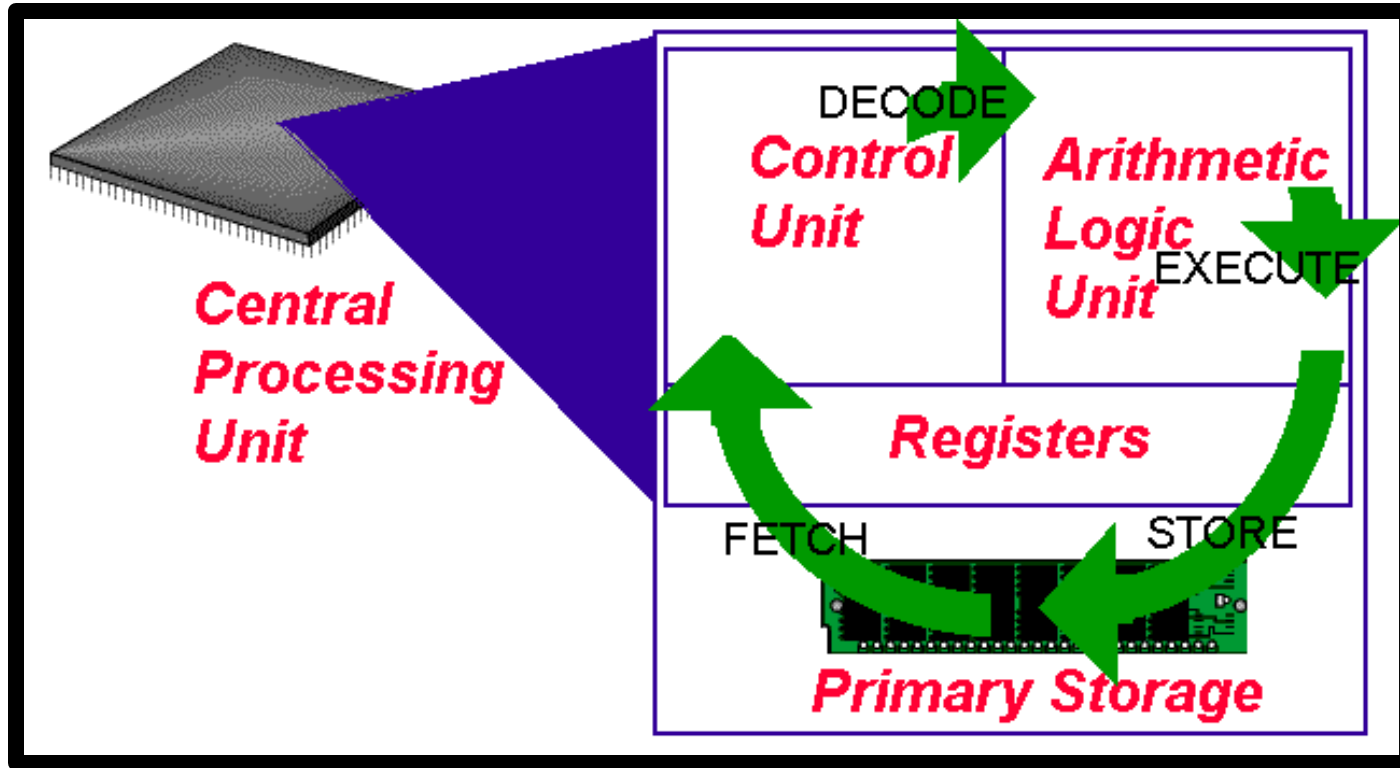
I. Hardware



This process happens really fast

... like *really* fast

I. Hardware

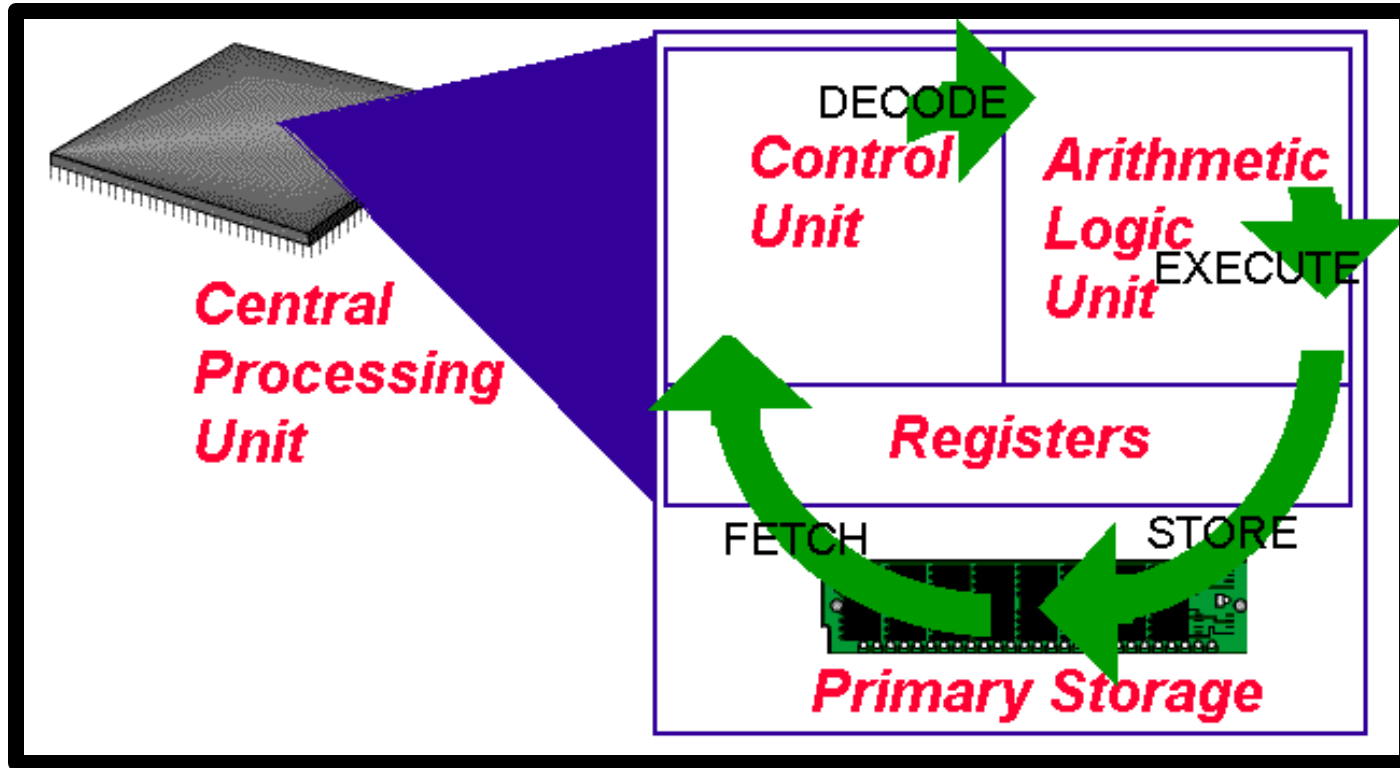


This process happens really fast

... like *really* fast

Computers can execute one or more instructions per clock cycle*

I. Hardware



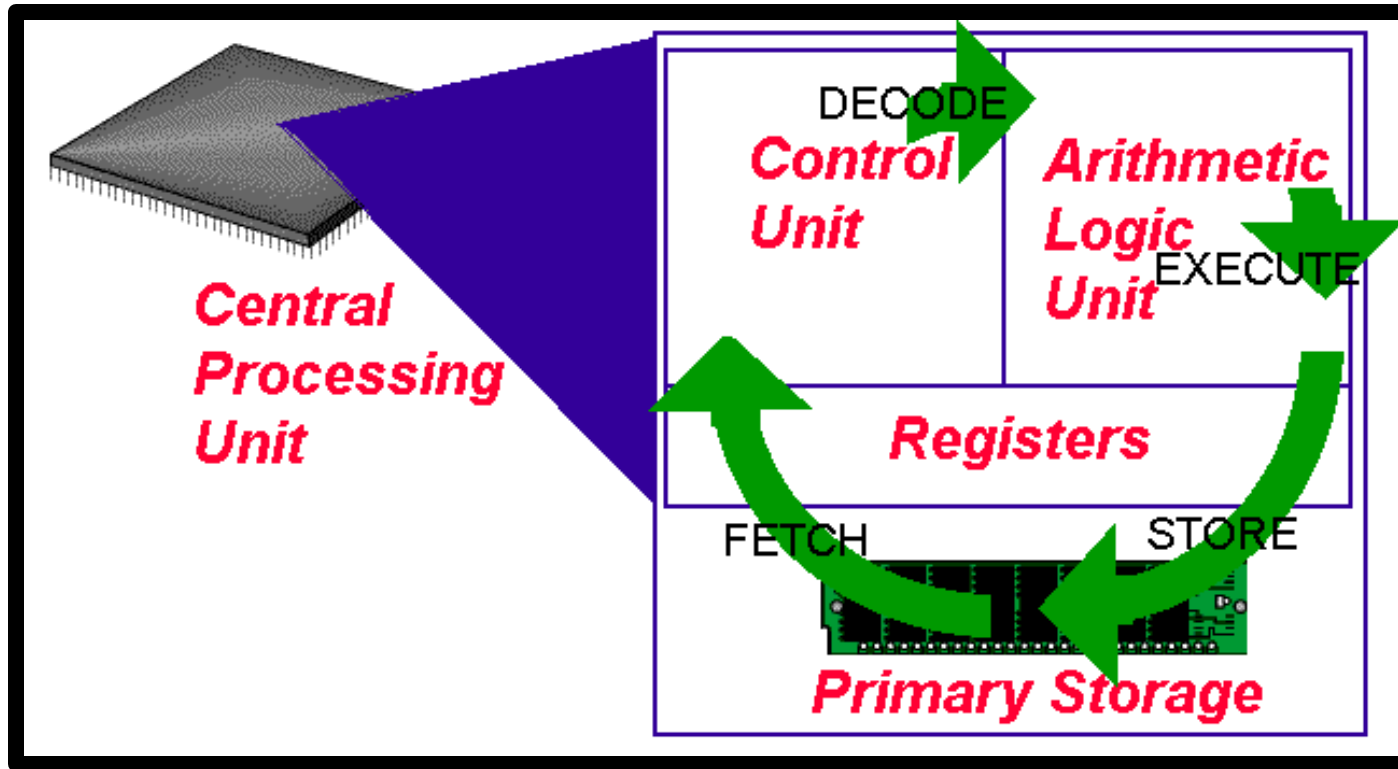
This process happens really fast

... like *really* fast

Computers can execute one or more instructions per clock cycle*

4GHz CPU speed = 4,000,000,000 clock cycles per second

I. Hardware



This process happens really fast

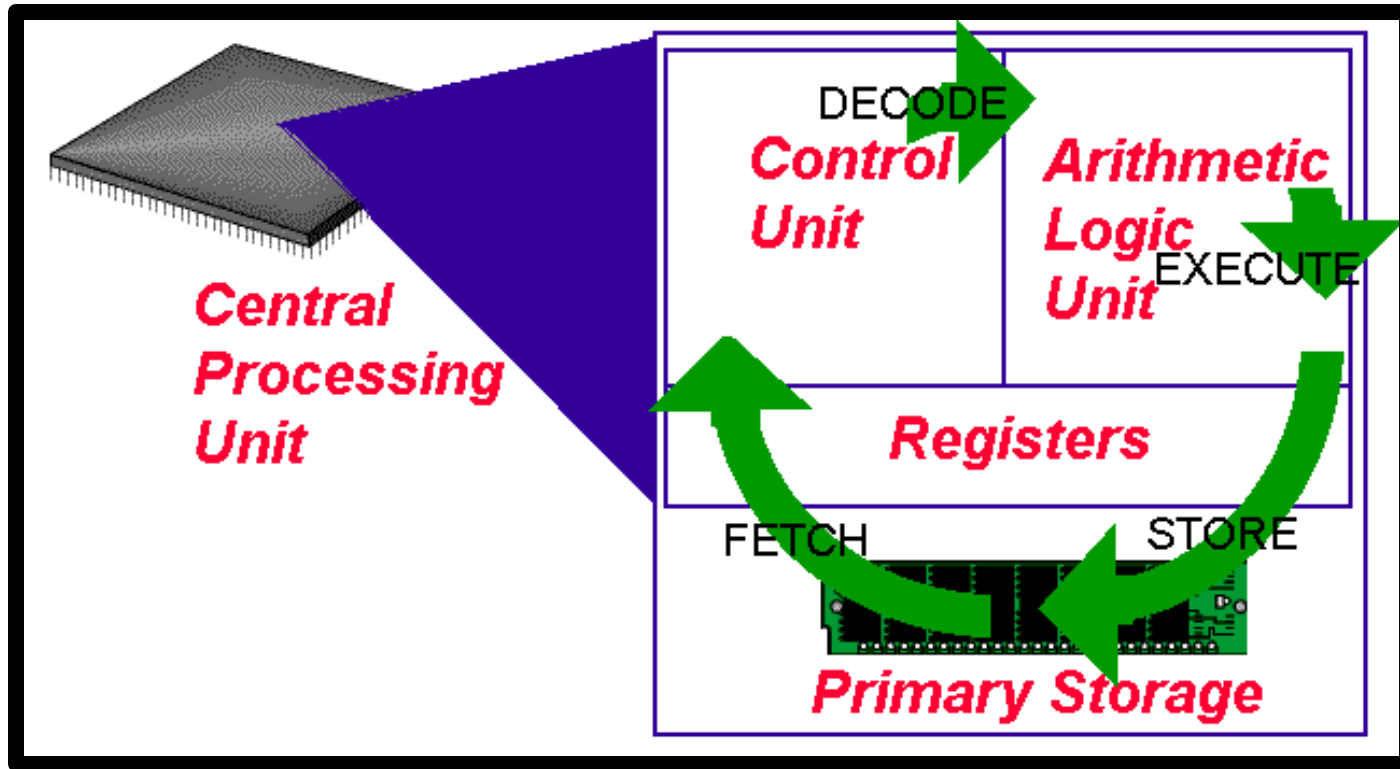
... like *really* fast

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I. Hardware



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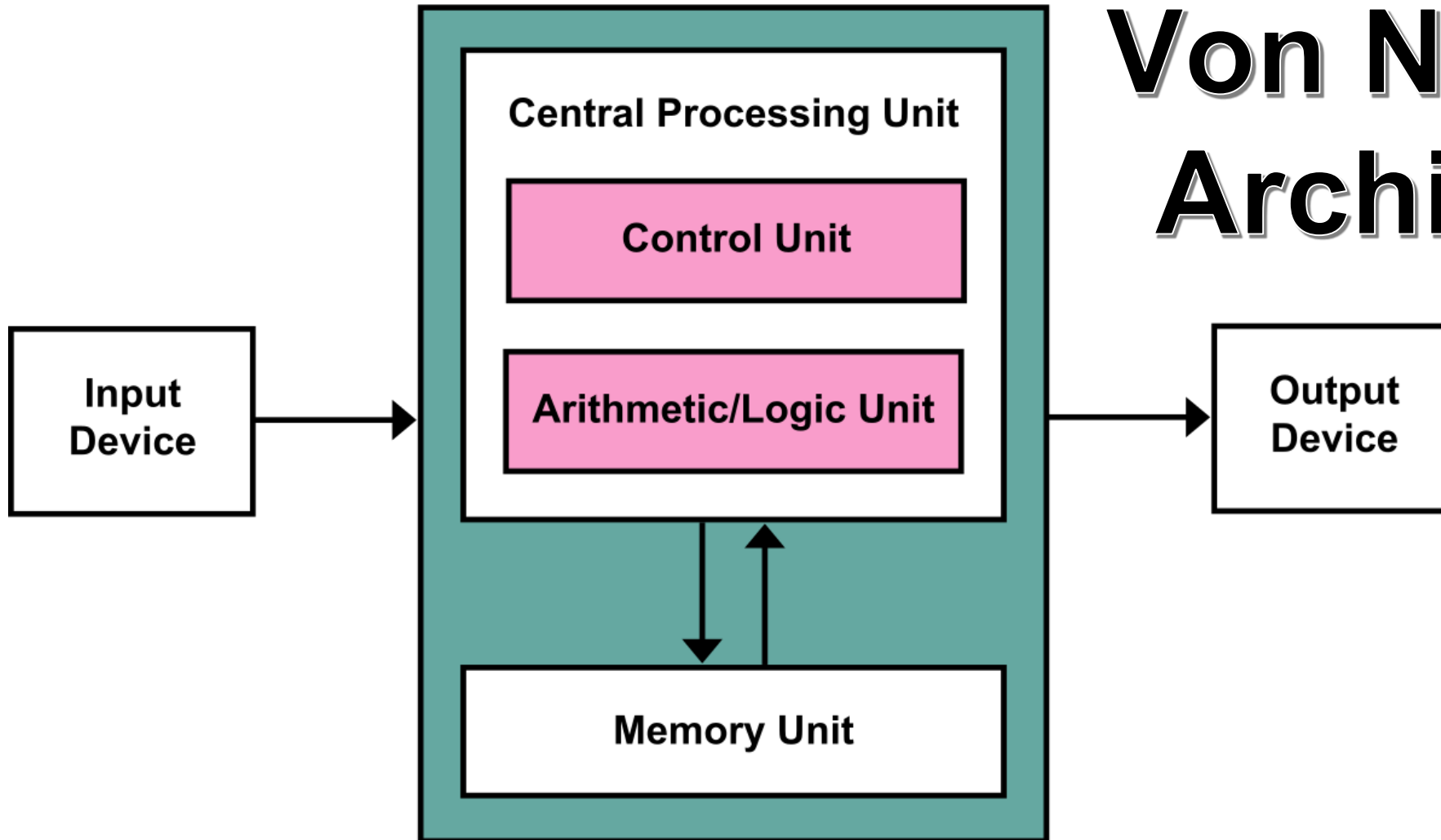
Computers can execute one or more instructions per clock cycle*



4GHz CPU speed = 4,000,000,000 clock cycles per second

Multi-core systems are even faster

I. Hardware



Von Neumann Architecture

I. Hardware

Brain with no short-term memory

CPU

Executes instructions



A screenshot of a text-based assembly code listing, showing various instructions and their corresponding memory addresses.

Must decipher
what instruction
to execute



The CPU can
add, subtract,
logic, move stuff
around

Registers

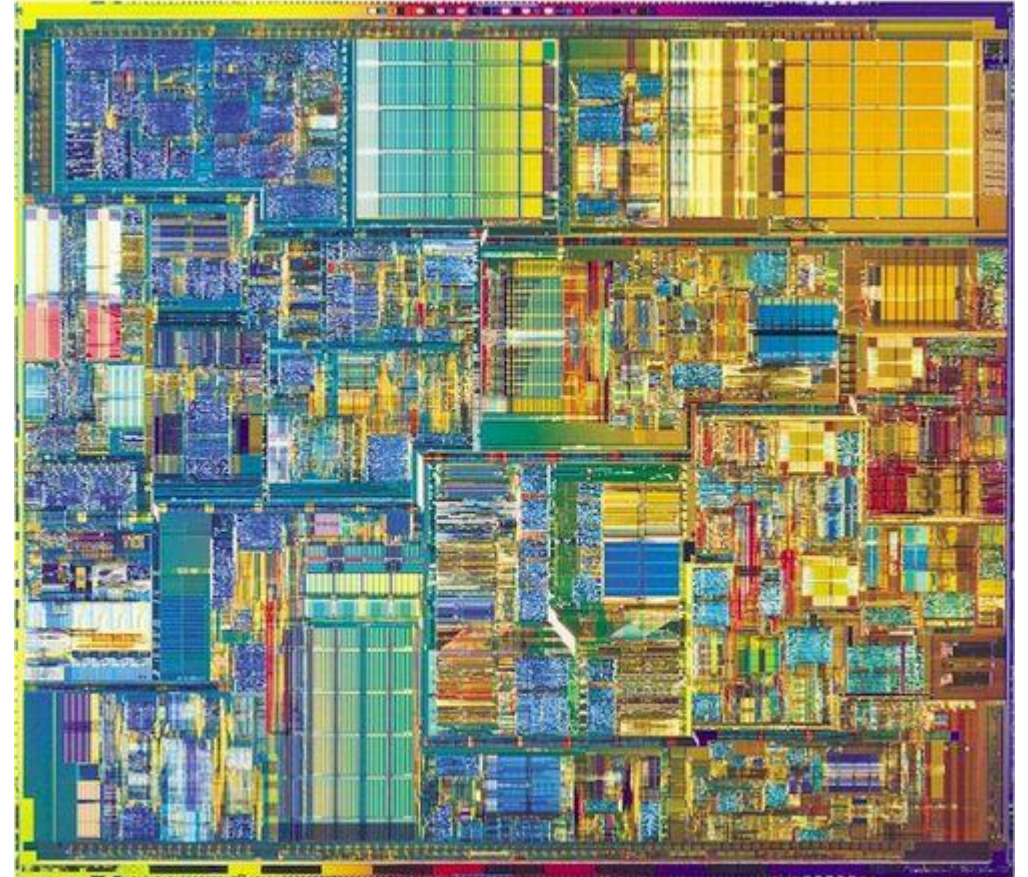
5 00001

8 00011

3 00101

...

11111



I. Hardware

Brain with no short-term memory

CPU

Executes instructions



Instruction	Operation
ADD	Add
SUB	Subtract
MUL	Multiply
DIV	Divide
MOV	Move
LD	Load
ST	Store
HALT	Stop

Must decipher what instruction to execute



The CPU can add, subtract, logic, move stuff around

Registers

5 00001

8 00011

3 00101

...

11111



16-bit ALU in minecraft

9,383,323 views • Sep 28, 2010

People have been able to create CPU components and fully functional, multi-core computers in games such as Minecraft

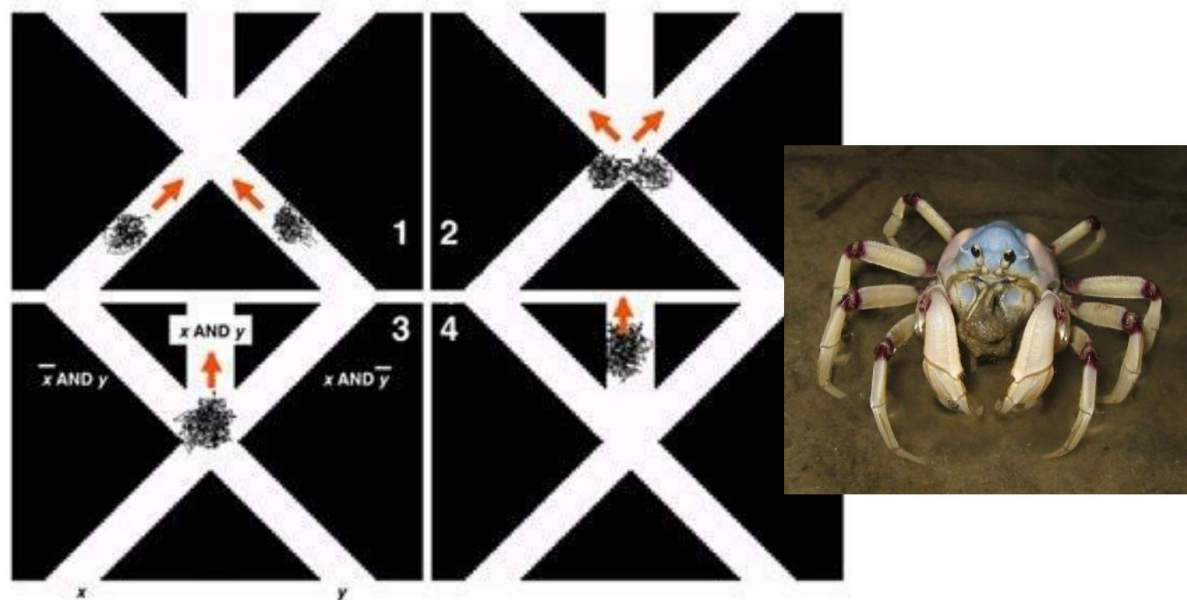
I. Hardware

WIRED STAFF

BUSINESS 04.14.2012 03:28 PM

Computer Built Using Swarms Of Soldier Crabs

Computer scientists at Kobe University in Japan have built a computer that draws inspiration from the swarming behavior of soldier crabs. The computer is based on theories from the early 1980s that studies how it could be possible to build a computer out of billiard balls. Proposed by Edward Fredkin and Tommaso Toffoli, the mechanical [...]



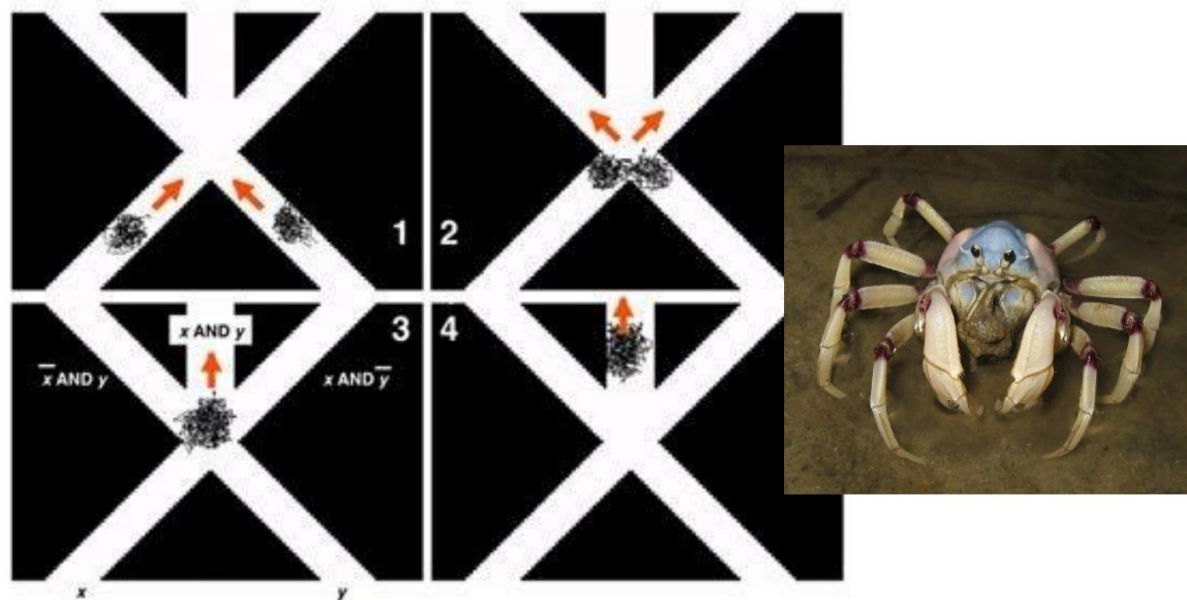
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This is very real

Robust Soldier Crab Ball Gate

Yukio-Pegio Gunji
Yuta Nishiyama
Department of Earth and Planetary Sciences
Kobe University
Kobe 657-8501, Japan

Andrew Adamatzky
Unconventional Computing Centre
University of the West of England
Bristol, United Kingdom

Soldier crabs *Mictyris guinotae* exhibit pronounced swarming behavior. Swarms of the crabs are tolerant of perturbations. In computer models and laboratory experiments we demonstrate that swarms of soldier crabs can implement logical gates when placed in a geometrically constrained environment.

1. Introduction

All natural processes can be interpreted in terms of computations. To implement a logical gate in a chemical, physical, or biological spatially extended medium, Boolean variables must be assigned to disturbances, defects, or localizations traveling in the medium. These traveling patterns collide and the outcome of their collisions are converted

<https://wpmedia.wolfram.com/uploads/sites/13/2018/02/20-2-2.pdf>

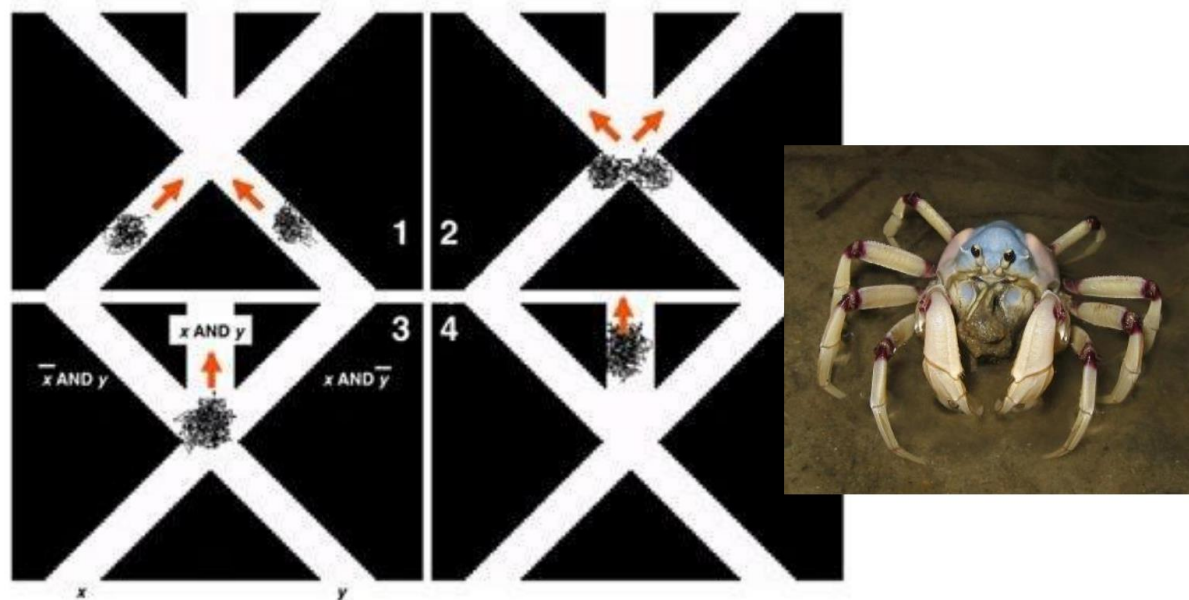
I. Hardware

WIRED STAFF

BUSINESS 04.14.2012 03:28 PM

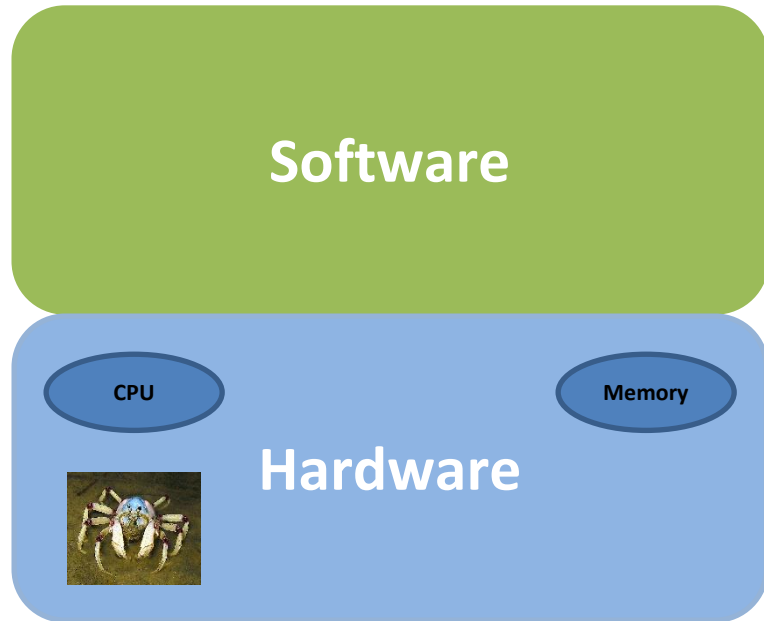
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(In theory) If you wanted to play Doom (1993) using a CPU made from soldier crabs, you would need 22 million crabs

How does this happen?



From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**

II. Software

A sequence of instructions, or **program**,
that tells the computer how to work



II. Software

A sequence of instructions, or **program**,
that tells the computer how to work

Humans write code in binary ?

```
01110011 00111011 01101111 01110101 01101100 01101110 01100111 01101001
01100001 00111011 01110011 01100111 01101110 00111011 01110011 01101011
01100100 01110110 01101110 01100001 00111011 01100100 01100111 01110100
00111011 01001111 01001100 01001011 01010011 01000100 01000110 00111011
01001100 01010011 01000100 01001000 01000111 01000100 01001100 00111011
01000110 01010011 01001011 01000111 01100110 01100100 01101000 01100100
01100110 01110011 01101000 01110011 01100110 01110111 01100101 01110111
01100101 01110010 01110111 01100101 01110010 01100110 01110110 01100111
01110011 01100100 01100010 01100100 01100110 01101000 01100010 01100100
01100110 01100100 01110011 01100100 01100110 01110011 01100100 01100110
01110011 01100100 01100110 01110011 01100100 01101000 01100100 01100110
01101000 01110011 01100100 01100110 01110011 01100100 01100110 01110011
01100100 01100110 01101000 01110100 01110010 01100100 01100110 01101000
01100100 01100110 01100111 01101000 01100110 01100111 01100111 01110100
01110010 01111001 01101000 01110100 01110010 01101000 01110010 01110011
01100110 01101000 01100111 01100110 01101000 01110011 01110010 01110100
01101000 01110100 01110010 01101000 01110011 01110100 01110010 01100100
01100110 01101000 01100100 01100110 01101000 01100100 01100110 01101000
01100100 01100110 01101000 01100101 01110111 01100110 01110011 01100100
01100110 01110010 01100101 01110111 01110100 01100111 01100101 01100111
01100100 01100110 01110011 01100111 01110011 01100110 0110010001110011
00111011 01101111 01110101 01101100 01101110 01100111 01101001 01100001
00111011 01110011 01100111 01101110 00111011 01110011 01101011 01100100
01110110 01101110 01100001 00111011 01100100 01100111 01110100 00111011
01001111 01001100 01001011 01010011 01000100 01000110 00111011 01001100
01010011 01000100 01001000 01000111 01000100 01001100 00111011 01000110
01010011 01001011 01000111 01100110 01100100 01101000 01100100 01100110
01110011 01101000 01110011 01100110 01110111 01100101 01110111 01100101
01110010 01110111 01100101 01110010 01100110 01110110 01100111 01110011
01100100 01100010 01100100 01100110 01101000 01100010 01100100 01100110
01100100 01110011 01100100 01100110 01110011 01100100 01100110 01110011
01100100 01100110 01110011 01100100 01101000 01100100 01100110 01101000
01110011 01100100 01100110 01110011 01100100 01100110 01110011 01100100
01100110 01101000 01110100 01110010 01100100 01100110 01101000 01100100
01100110 01100111 01101000 01100110 01100111 01100111 01110100 01110010
01111001 01101000 01110100 01110010 01101000 01110010 01110011 01100110
01101000 01100111 01100110 01101000 01110011 01110010 01110100 01101000
```

II. Software

A sequence of instructions, or **program**,
that tells the computer how to work

Humans write code in binary ?

NO

01110011 00111011 01101111 01110101 01101100 01101110 01100111 01101001
01100001 00111011 01110011 01100111 01101110 00111011 01110011 01101011
01100100 01110110 01101110 01100001 00111011 01100100 01100111 01110100
00111011 01001111 01001100 01001011 01010011 01000100 01000110 00111011
01001100 01010011 01000100 01001000 01000111 01000100 01001100 00111011
01000110 01010011 01001011 01000111 01100110 01100100 01101000 01100100
01100110 01110011 01101000 01110011 01100110 01110111 01100101 01110111
01100101 01110010 01110111 01100101 01110010 01100110 01110110 01100111
01110011 01100100 01100010 01100100 01100110 01101000 01100010 01100100
01100110 01100100 01110011 01100100 01100110 01110011 01100100 01100110
01110011 01100100 01100110 01110011 01100100 01101000 01100100 01100110
01101000 01110011 01100100 01100110 01110011 01100100 01100110 01110011
01100100 01100110 01101000 01110100 01110010 01100100 01100110 01101000
01100100 01100110 01100111 01101000 01100110 01100111 01100111 01110100
01110010 01111001 01101000 01110100 01110010 01101000 01110010 01110011
01100110 01101000 01100111 01100110 01101000 01110011 01110010 01110100
01101000 01110100 01110010 01101000 01110011 01110100 01110010 01100100
01100110 01101000 01100100 01100110 01101000 01100100 01100110 01101000
01100100 01100110 01101000 01100101 01110111 01100110 01110011 01100100
01100110 01110010 01100101 01110111 01110100 01100111 01100101 01100111
01100100 01100110 01110011 01100111 01110011 01100110 0110010001110011
00111011 01101111 01110101 01101100 01101110 01100111 01101001 01100001
00111011 01110011 01100111 01101110 00111011 01110011 01101011 01100100
01110110 01101110 01100001 00111011 01100100 01100111 01110100 00111011
01001111 01001100 01001011 01010011 01000100 01000110 00111011 01001100
01010011 01000100 01001000 01000111 01000100 01001100 00111011 01000110
01010011 01001011 01000111 01100110 01100100 01101000 01100100 01100110
01110011 01101000 01110011 01100110 01110111 01100101 01110111 01100101
01110010 01110111 01100101 01110010 01100110 01110110 01100111 01110011
01100100 01100010 01100100 01100110 01101000 01100010 01100100 01100110
01100100 01110011 01100100 01100110 01110011 01100100 01100110 01110011
01100100 01100110 01110011 01100100 01101000 01100100 01100110 01101000
01110011 01100100 01100110 01110011 01100100 01100110 01110011 01100100
01100110 01101000 01110100 01110010 01100100 01100110 01101000 01100100
01100110 01100111 01101000 01100110 01100111 01100111 01110100 01110010
01111001 01101000 01110100 01110010 01101000 01110010 01110011 01100110
01101000 01100111 01100110 01101000 01110011 01110010 01110100 01101000

II. Software

We write programs in a readable, higher-level language

```
#include <stdio.h>

int main() {
    printf("Hello WOrld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n",x,y,z);
    return 0;
```

```
class Person():
    #method to initialize name and age attributes.
    def __init__(self,name, age):
        self.name = name
        self.age = age
    #method to demonstrate what a person eats
    def eat(self):
        print(self.name.title() + "eats Matooke and rice")
        print("She is"+ str(self.age) + " years old")
    def drink(self):
        print("Drinks water")
    #instantiating a class.
    my_sister = Person("Haniifa", 30)
    #Accessing the class method through the class object.
    my_sister.eat()
```



II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

int main() {
    printf("Hello WOrld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n",x,y,z);
    return 0;
}
```

```
class Person():
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}
```

II. Software

```
#include <stdio.h>

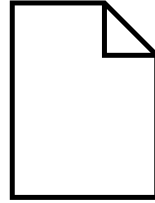
int main() {
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    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

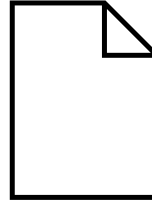
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
 0: f3 0f 1e fa          endbr64
 4: 55                   push    %rbp
 5: 48 89 e5             mov     %rsp,%rbp
 8: 48 83 ec 10          sub     $0x10,%rsp
 c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                 leaveq  %eax
53: c3                 retq
```

- Converted to assembly code
- .s file

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

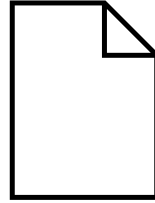
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10          sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                   leaveq  %eax
53: c3                   retq
```

- Converted to assembly code
- .s file

Assembler

```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 00000000000000101
4 11101111 10011110 0000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 0000000000010101
7 11101111 00000010 11111011 0000000000010111
8 11110100 10101101 11011111 0000000000011110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

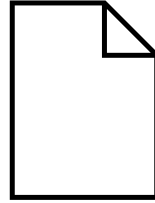
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10          sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                   leaveq  %eax
53: c3                   retq
```

- Converted to assembly code
- .s file

Assembler

We still need to resolve function calls
i.e. printf, sleep, sqrt, etc

Linker

```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 00000000000000101
4 11101111 10011110 0000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 0000000000010101
7 11101111 00000010 11111011 0000000000010111
8 11110100 10101101 11011111 0000000000011110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

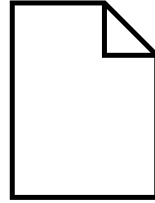
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10          sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                   leaveq  %eax
53: c3                   retq
```

Assembler

- Converted to assembly code
- .s file

Two methods:

Program A

Library 1 string.h

Library 2 stdio.h

Linker

```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 000000000000000101
4 11101111 10011110 00000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 0000000000010101
7 11101111 00000010 11111011 0000000000010111
8 11110100 10101101 11011111 0000000000011110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

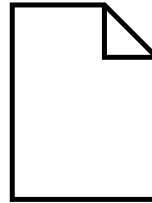
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10          sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                 leaveq  %eax
53: c3                 retq
```

- Converted to assembly code
- .s file

Static Linking- required code and data copied into executable at compile time



Linker

```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 00000000000000101
4 11101111 10011110 0000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 0000000000010101
7 11101111 00000010 11111011 0000000000010111
8 11110100 10101101 11011111 0000000000011110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11101000 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```


II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

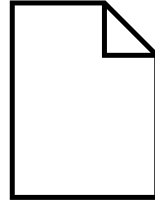
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10          sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                   leaveq  %eax
53: c3                   retq
```

Assembler

- Converted to assembly code
- .s file

Two methods:

Program A

Library 1 string.h

Library 2 stdio.h

Linker

```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 000000000000000101
4 11101111 10011110 00000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 00000000000010101
7 11101111 00000010 11111011 0000000000001011
8 11110100 10101101 11011111 0000000000011110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

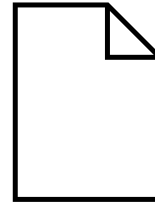
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

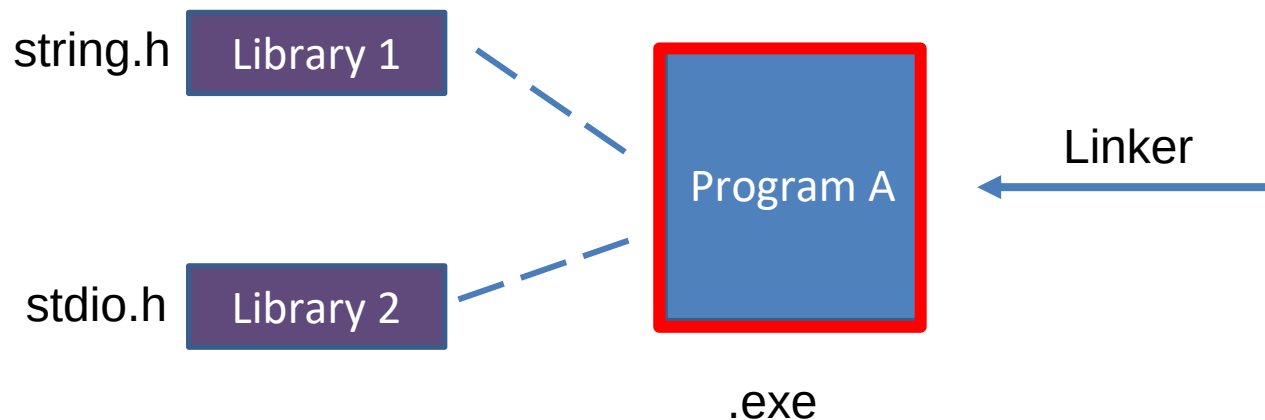
Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10         sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00     callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00     mov     $0x0,%eax
48: e8 00 00 00 00     callq  4d <main+0x4d>
4d: b8 00 00 00 00     mov     $0x0,%eax
52: c9                 leaveq  %eax
53: c3                 retq
```

Assembler

- Converted to assembly code
- .s file

Two methods:



```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 000000000000000101
4 11101111 10011110 00000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 00000000000010101
7 11101111 00000010 11111011 0000000000001011
8 11110100 10101101 11011111 0000000000001110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```

Dynamic Linking - required code and data is linked to executable at runtime

II. Software

We need a way to convert **source** code to **binary**

```
#include <stdio.h>

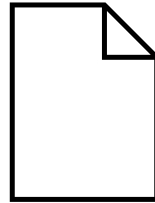
int main() {
    printf("Hello W0rld! \n");

    int x = 0;
    int y = 3;

    int z = x + y;

    printf("%d %d %d \n", x, y, z);
    return 0;
}
```

Preprocessor



- Removal of comments
- Expand Macros

Compiler

```
0000000000000000 <main>:
0: f3 0f 1e fa          endbr64
4: 55                   push    %rbp
5: 48 89 e5             mov     %rsp,%rbp
8: 48 83 ec 10          sub     $0x10,%rsp
c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 13 <main+0x13>
13: e8 00 00 00 00      callq  18 <main+0x18>
18: c7 45 f4 00 00 00 00 movl    $0x0,-0xc(%rbp)
1f: c7 45 f8 03 00 00 00 movl    $0x3,-0x8(%rbp)
26: 8b 55 f4             mov     -0xc(%rbp),%edx
29: 8b 45 f8             mov     -0x8(%rbp),%eax
2c: 01 d0               add     %edx,%eax
2e: 89 45 fc             mov     %eax,-0x4(%rbp)
31: 8b 4d fc             mov     -0x4(%rbp),%ecx
34: 8b 55 f8             mov     -0x8(%rbp),%edx
37: 8b 45 f4             mov     -0xc(%rbp),%eax
3a: 89 c6               mov     %eax,%esi
3c: 48 8d 3d 00 00 00 00 lea     0x0(%rip),%rdi    # 43 <main+0x43>
43: b8 00 00 00 00      mov     $0x0,%eax
48: e8 00 00 00 00      callq  4d <main+0x4d>
4d: b8 00 00 00 00      mov     $0x0,%eax
52: c9                 leaveq  %eax
53: c3                 retq
```

- Converted to assembly code
- .s file

Assembler

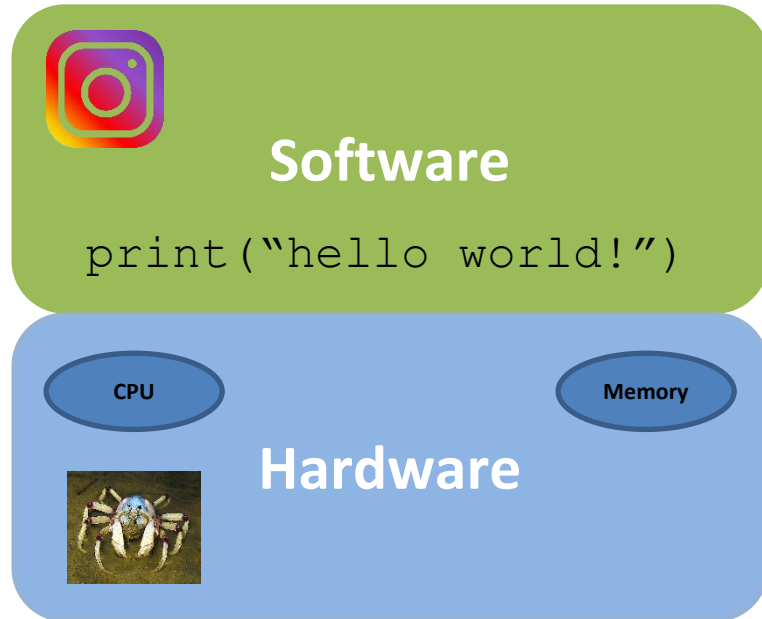
To be continued

Linker



```
1 00000000 00000100 0000000000000000
2 01011110 00001100 11000010 0000000000000010
3 11101111 00010110 000000000000000101
4 11101111 10011110 00000000000001011
5 11111000 10101101 11011111 0000000000010010
6 01100010 11011111 0000000000010101
7 11101111 00000010 11111011 0000000000010111
8 11110100 10101101 11011111 0000000000011110
9 00000011 10100010 11011111 0000000000100001
10 11101111 00000010 11111011 0000000000100100
11 01111110 11110100 10101101
12 11111000 10101110 11000101 0000000000101011
13 00000110 10100010 11111011 0000000000110001
14 11101111 00000010 11111011 0000000000110100
15 01010000 11010100 0000000000111011
16 00000100 0000000000111101
```

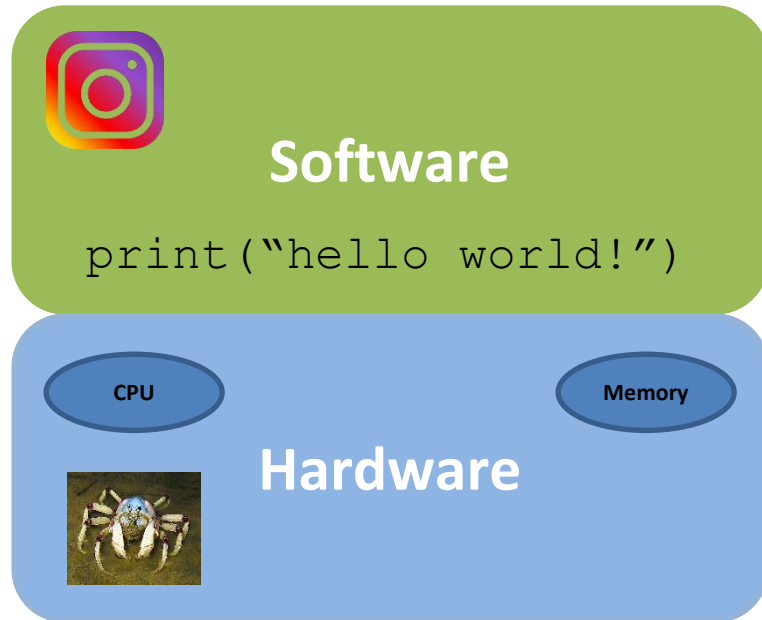
How does this happen?



From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**

How does this happen?

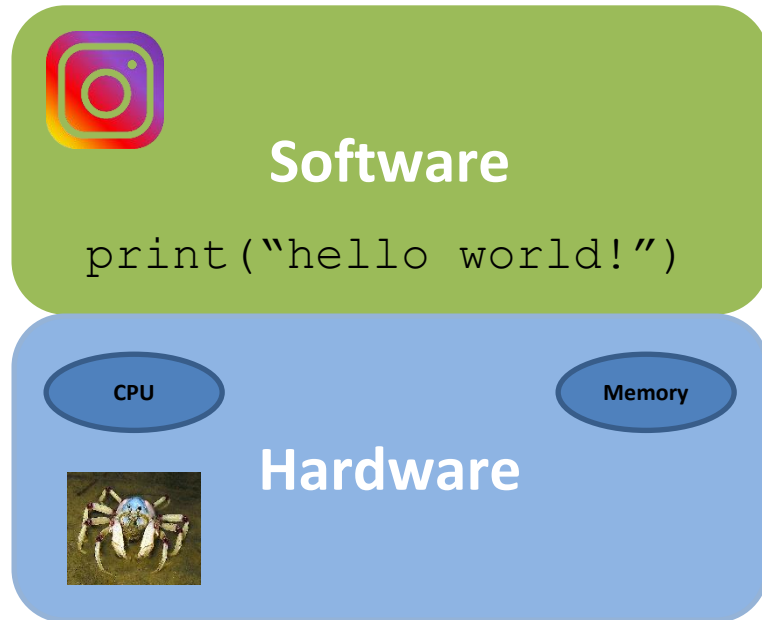


From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**

Software is nothing without hardware

How does this happen?

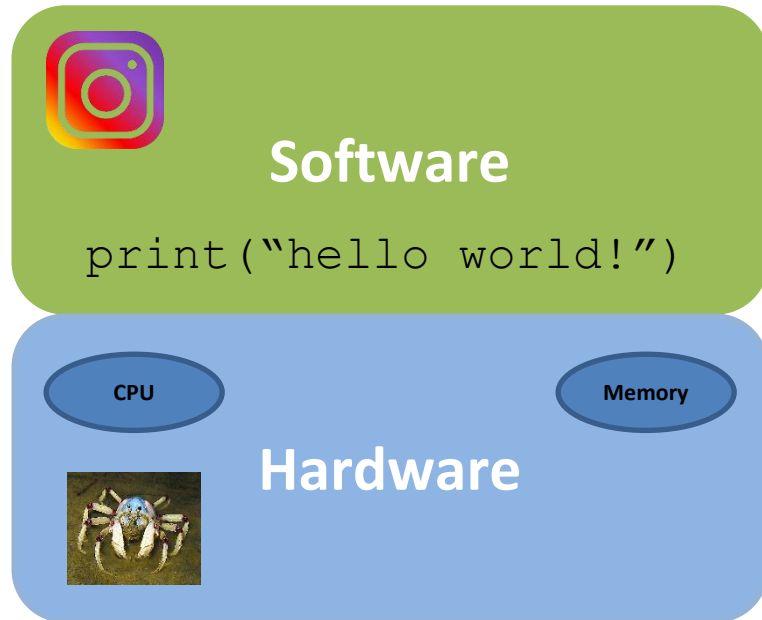


From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**

Hardware is *mostly* nothing without software

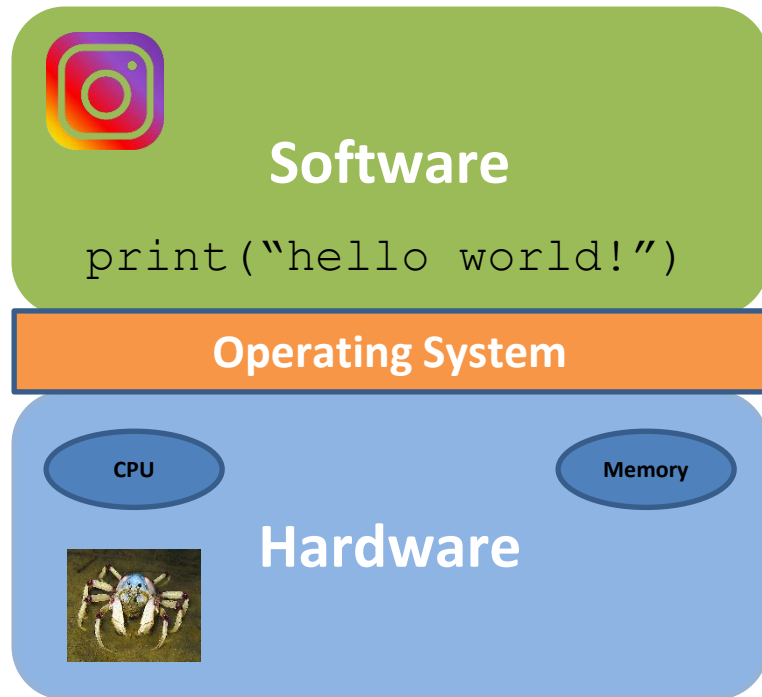
How does this happen?



From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**
- III. **???**

How does this happen?

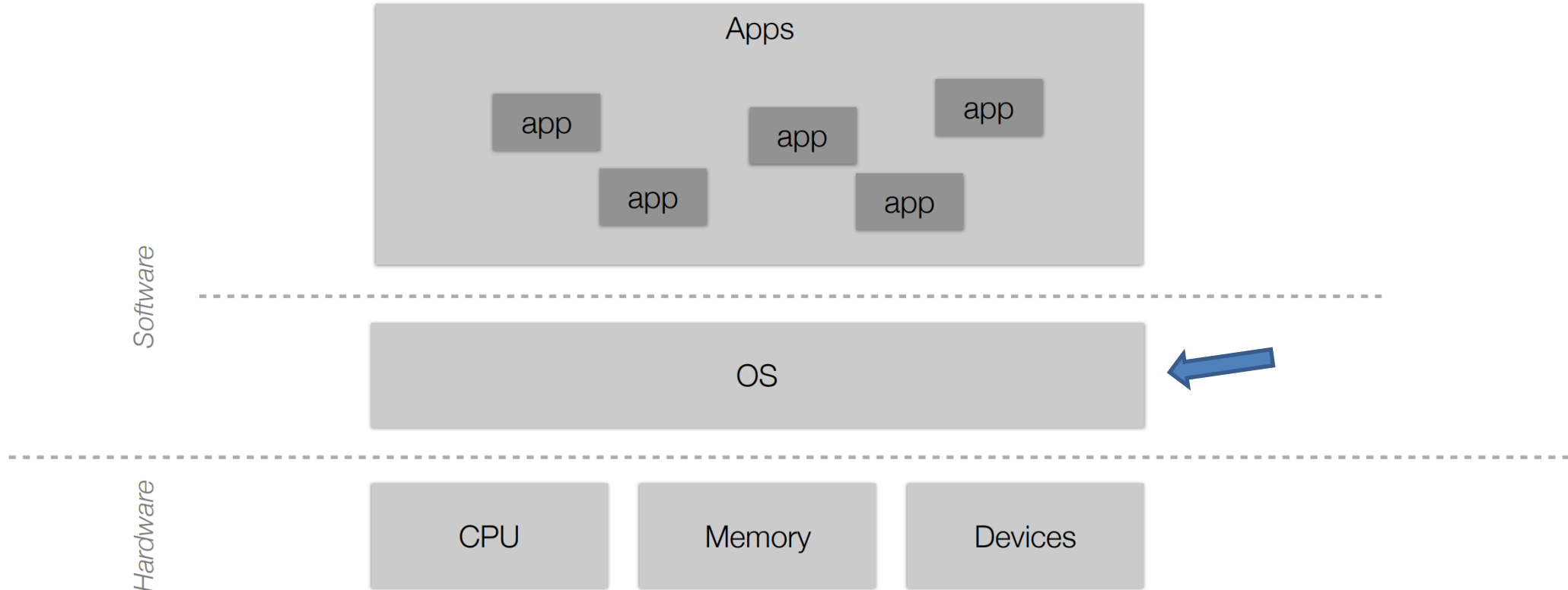


From a high level, we will divide a computer system into two parts

- I. **Hardware**
- II. **Software**
- III. **Operating System**

Typical Layers of a Computer

The **operating system** is a vital component of a computer

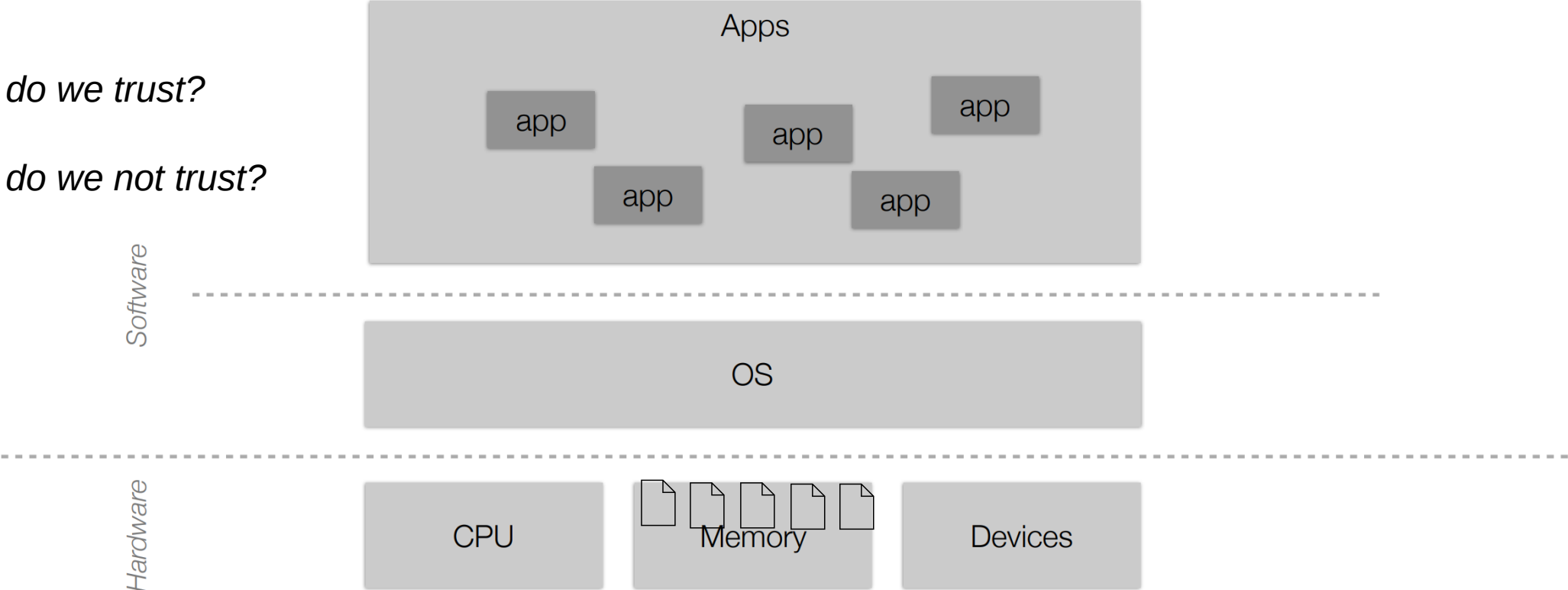


Typical Layers of a Computer

The **operating system** is a vital component of a computer

What do we trust?

What do we not trust?



Typical Layers of a Computer

The **operating system** is a vital component of a computer

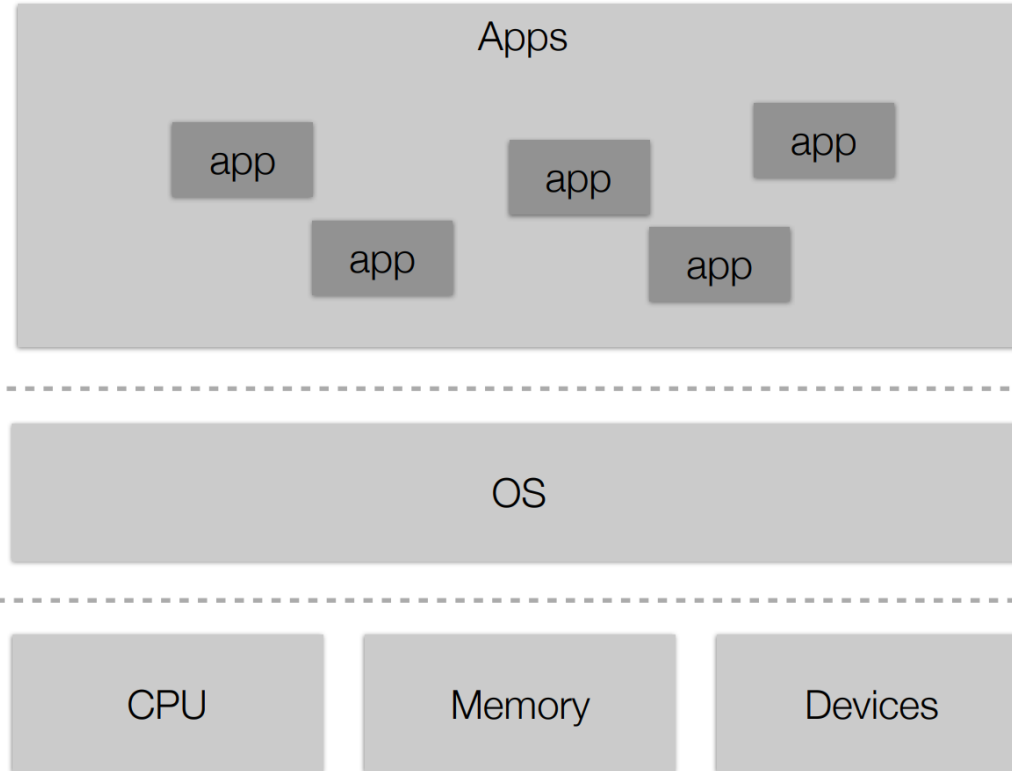
Not privileged

Privileged

Hardware

Software

Hardware



Typical Layers of a Computer

The **operating system** is a vital component of a computer

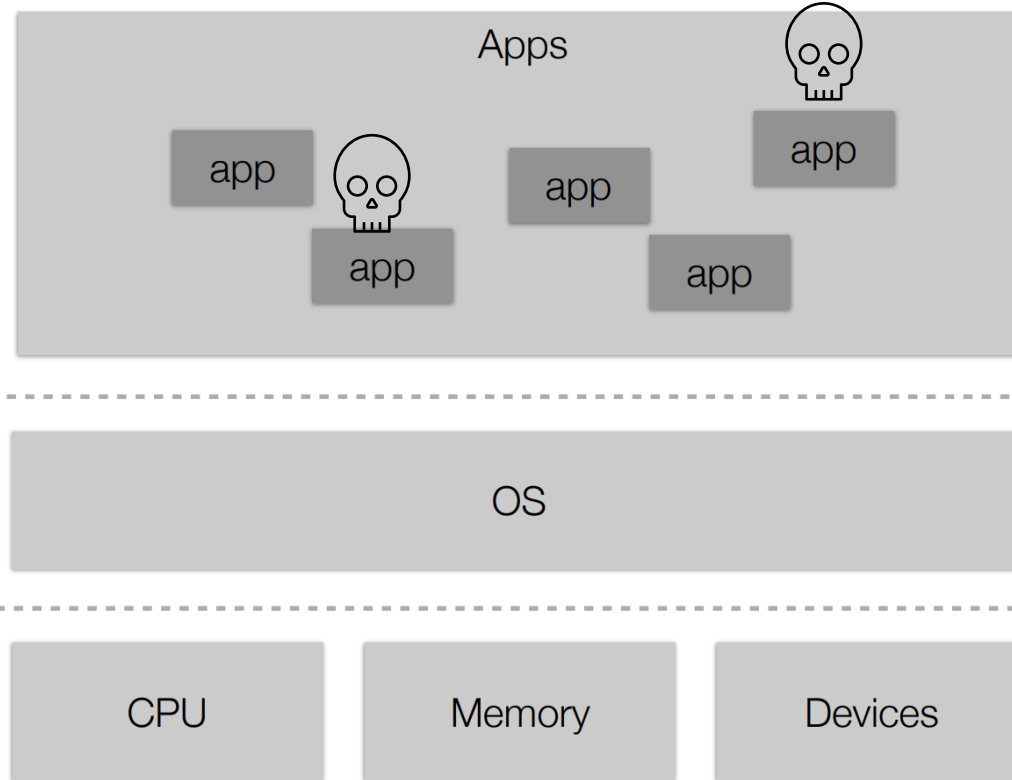
Not privileged

Privileged

Hardware

Software

Hardware



Typical Layers of a Computer

The **operating system** is a vital component of a computer

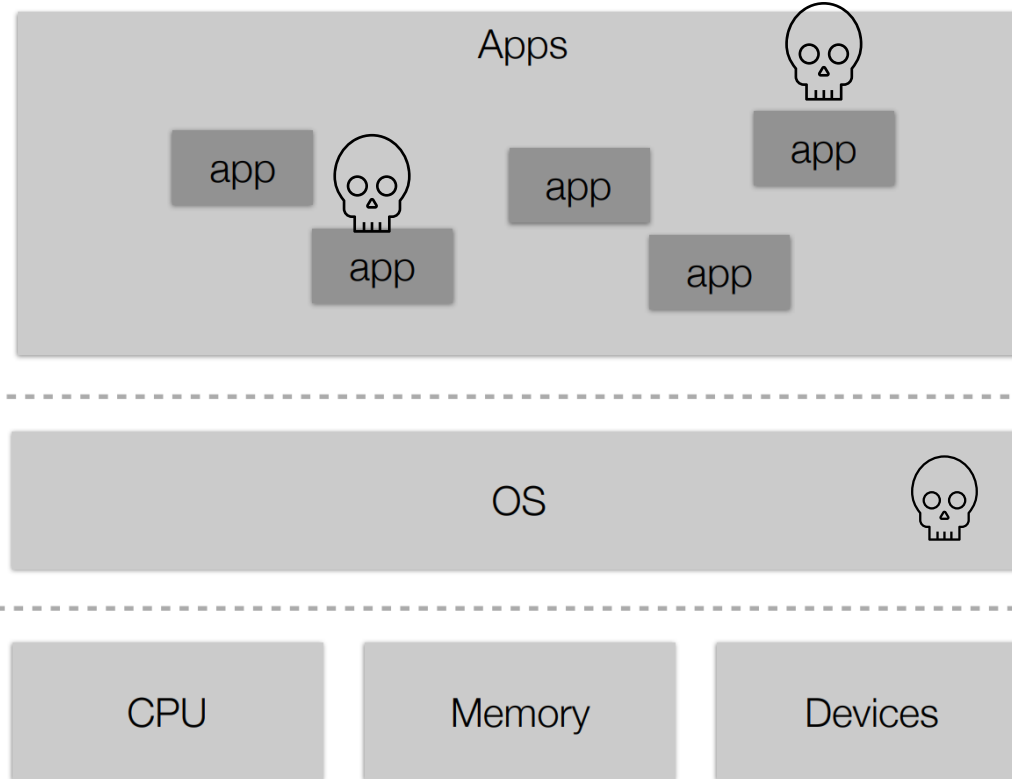
Not privileged

Privileged

Hardware

Software

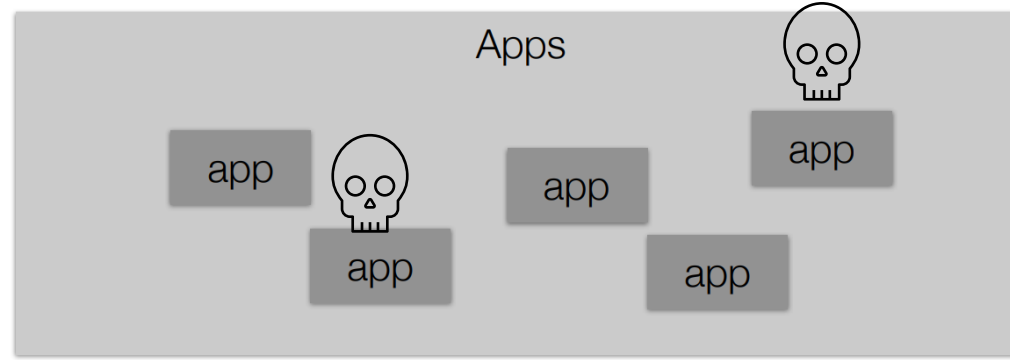
Hardware



Typical Layers of a Computer

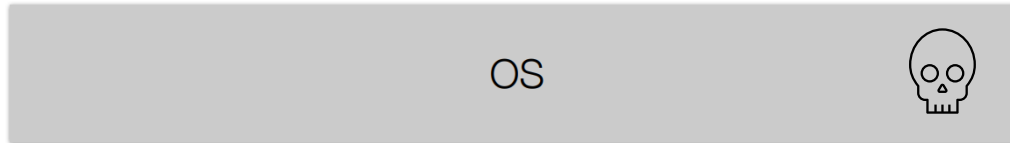
The **operating system** is a vital component of a computer

Not privileged



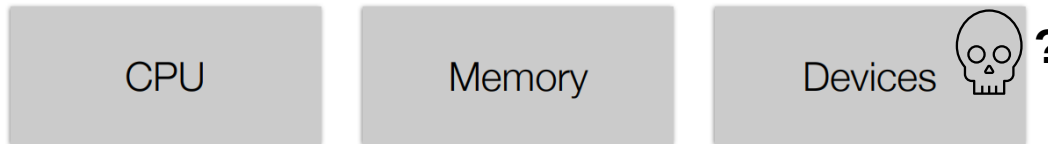
Privileged

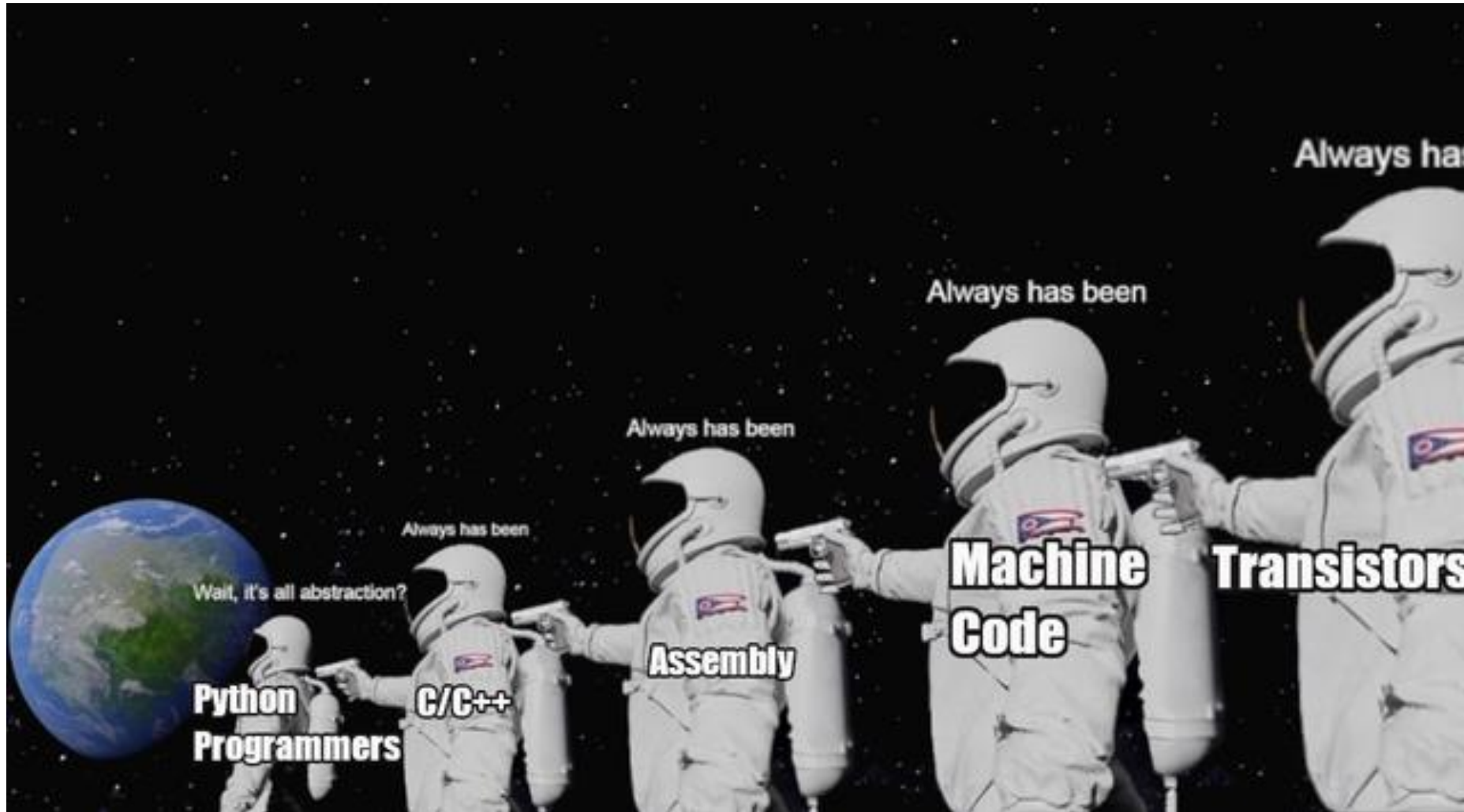
Software



Hardware

Hardware





Meme credit: Carson Gross