

CSCI 476: Computer Security

Lecture 5: Set-UID and Environment Variables

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Fall 2022

Announcements

Lab 1 Due **FRIDAY** 9/16 @ 11:59 PM

- Shouldn't be too bad

Note taker still needed

How would you protect your computer and its resources?

who can do what to whom?



users/groups

what is their identity?



Objects

Usually things on a filesystem



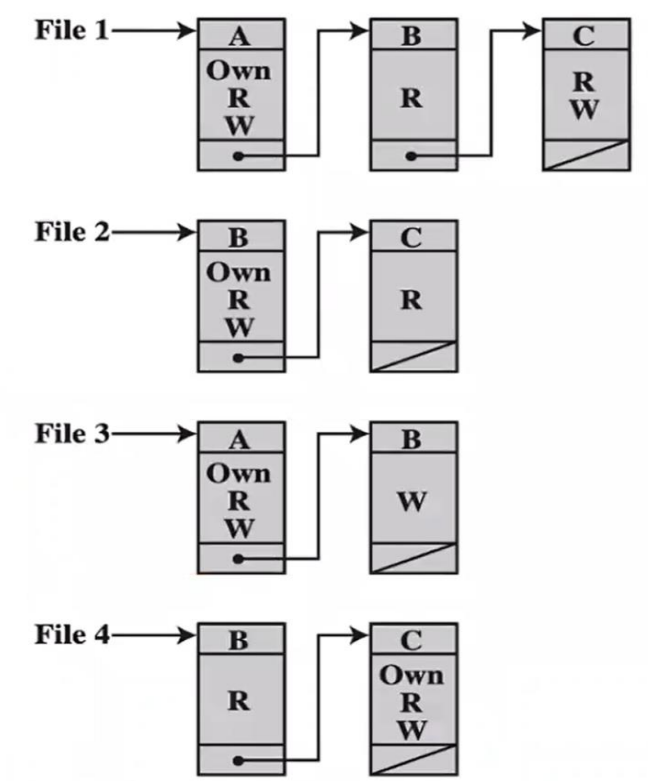
permissions (read/write/execute)

Ok, I know the who– what are *you* permitted to do?

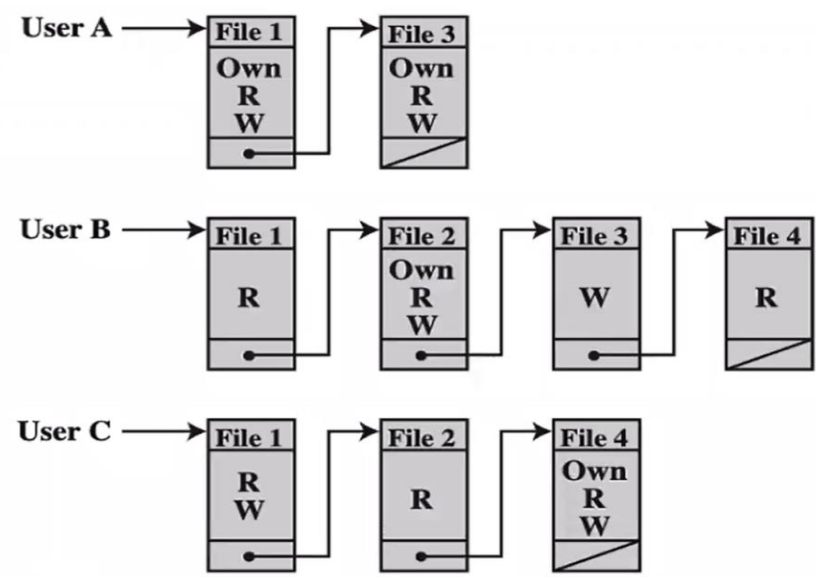
		OBJECTS			
		File 1	File 2	File 3	File 4
SUBJECTS	User A	Own Read Write		Own Read Write	
	User B	Read	Own Read Write	Write	Read
	User C	Read Write	Read		Own Read Write

Access Control Matrix

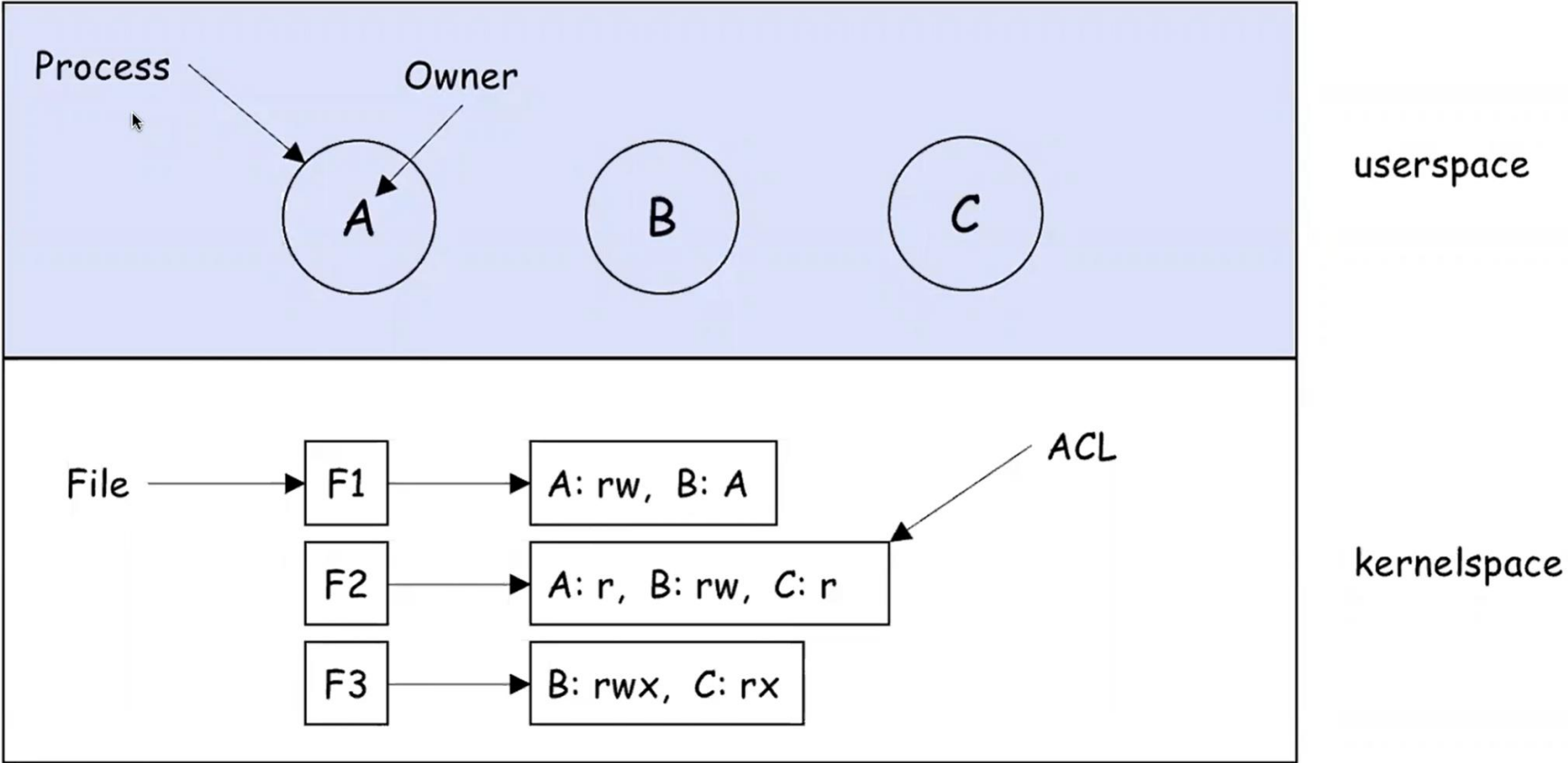
What are some issues with this?



Access Control list (ACL)



Wont take up as much memory!



Unix File Modes and Permissions

Every Unix file has a set of permissions that determine whether someone can read, write, or run the file

```
ls -l ~
```

```
ls -l /dev
```

```
[09/13/22] seed@VM:~$ ls -l ~
```

```
total 44
```

```
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Desktop
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Documents
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Downloads
drwxrwxr-x 2 seed seed 4096 Sep  1 14:37 lab0
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Music
drwxrwxr-x 2 seed seed 4096 Sep  6 15:23 os-review
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Pictures
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Public
drwxrwxr-x 2 seed seed 4096 Aug 25 13:41 shared
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Templates
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Videos
```


Unix File Modes and Permissions

Every Unix file has a set of permissions that determine whether someone can read, write, or run the file

```
[09/13/22] seed@VM:~$ ls -l ~
total 44
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Desktop
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Documents
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drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Public
drwxrwxr-x 2 seed seed 4096 Aug 25 13:41 shared
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Templates
drwxr-xr-x 2 seed seed 4096 Nov 24 2020 Videos
```

ls -l ~

ls -l /dev

Permissions for the file

Unix File Modes and Permissions

Every Unix file has a set of permissions that determine whether someone can read, write, or run the file

```
ls -l ~
ls -l /dev
```

Permissions for the file

Owner/group information

[09/13/22] seed@VM:~\$ ls -l ~

total 44

drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Desktop
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Documents
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Downloads
drwxrwxr-x	2	seed	seed	4096	Sep	1	14:37	lab0
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Music
drwxrwxr-x	2	seed	seed	4096	Sep	6	15:23	os-review
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Pictures
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Public
drwxrwxr-x	2	seed	seed	4096	Aug	25	13:41	shared
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Templates
drwxr-xr-x	2	seed	seed	4096	Nov	24	2020	Videos

Unix File Modes and Permissions

Every Unix file has a set of permissions that determine whether someone can read, write, or run the file

```
$ ls -l file  
-rw-r--r-- owner group date/time file
```

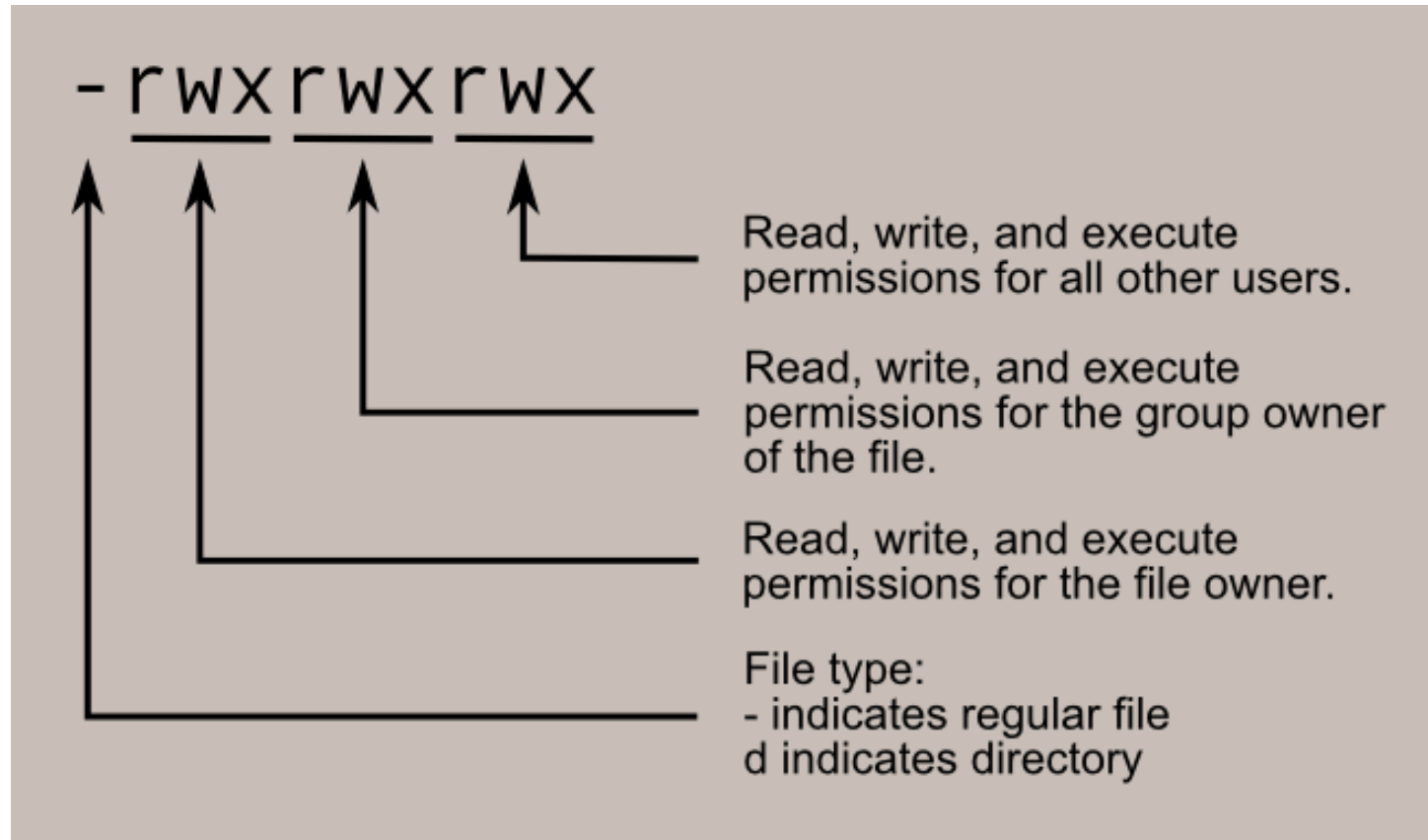
File permissions (4 parts)

- [file type][user][group][other]

Unix File Modes and Permissions

File permissions (4 parts)

- [file type][user][group][other]



Suppose you have the following file:

```
$ ls -l F  
-rw-rw-r-- B G ... F
```

If user **A** asks to perform some operation **O** on a file object **F**, the OS checks:

- Is **A** the owner of **F**?

Suppose you have the following file:

```
$ ls -l F  
-rw-rw-r-- B G ... F
```

If user **A** asks to perform some operation **O** on a file object **F**, the OS checks:

- Is **A** the owner of **F**?

No, B is the owner

Suppose you have the following file:

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$ ls -l F  
-rw-rw-r-- B G ... F
```

If user **A** asks to perform some operation **O** on a file object **F**, the OS checks:

- Is **A** the owner of **F**?
- Is **A** a member of **F**'s group?

Suppose you have the following file:

```
$ ls -l F  
-rw-rw-r-- B G ... F
```

If user **A** asks to perform some operation **O** on a file object **F**, the OS checks:

- Is **A** the owner of **F**?
- Is **A** a member of **F**'s group? Suppose $G = \{B, C, F\}$

A is not in F's group

Suppose you have the following file:

```
$ ls -l F  
-rw-rw-r-- B G ... F
```

If user **A** asks to perform some operation **O** on a file object **F**, the OS checks:

- Is **A** the owner of **F**?
- Is **A** a member of **F**'s group?
- Otherwise, what can they do?

Suppose you have the following file:

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$ ls -l F  
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```

If user **A** asks to perform some operation **O** on a file object **F**, the OS checks:

- Is **A** the owner of **F**?
- Is **A** a member of **F**'s group?
- Otherwise, what can they do?

Everyone can **read** file F

Suppose user C asks to execute a file object F2. Will they be able to do so?

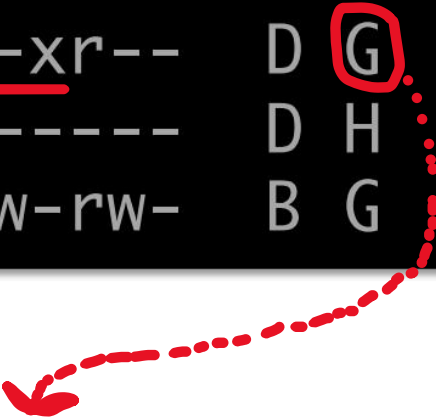
```
$ ls -l F
-rwxrwxrwx  B H  ...  F1
-rwxr-xr--  D G  ...  F2
-rw-r----- D H  ...  F3
-rw-rw-rw-  B G  ...  F4
```

Note:

- Group = G = {A, C, K, M, Q, Z}
- Group = H = {A, B, C, Q}

Suppose user C asks to execute a file object F2. Will they be able to do so?

```
$ ls -l F
-rwxrwxrwx B H ... F1
-rwxr-xr-- D G ... F2
-rw-r----- D H ... F3
-rw-rw-rw- B G ... F4
```



Note:

- Group = G = {A, C, K, M, Q, Z}
- Group = H = {A, B, C, Q}

Limitations of File-Based Access Control

When would a non-privilege user require more power/permissions?

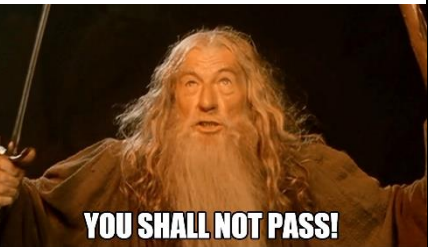
Limitations of File-Based Access Control

When would a non-privilege user require more power/permissions?

Changing password!

```
[seed@VM][~]$ ls -al /etc/passwd
-rw-r--r-- 1 root root 2886 Nov 24 09:12 /etc/passwd

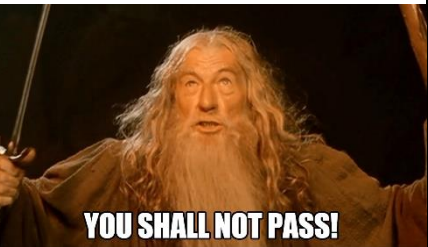
[seed@VM][~]$ ls -al /etc/shadow
-rw-r----- 1 root shadow 1514 Nov 24 09:12 /etc/shadow
```



Limitations of File-Based Access Control

When would a non-privilege user require more power/permissions?

Changing password!



```
[seed@VM][~]$ ls -al /etc/passwd
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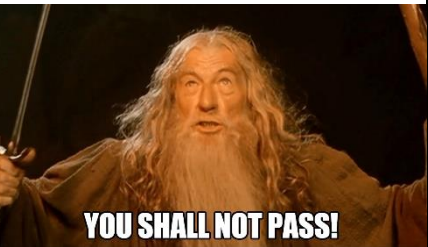
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`/etc/passwd` and `/etc/shadow` hold encrypted passwords for the user, in order to change our password, we will need to have access to those directories

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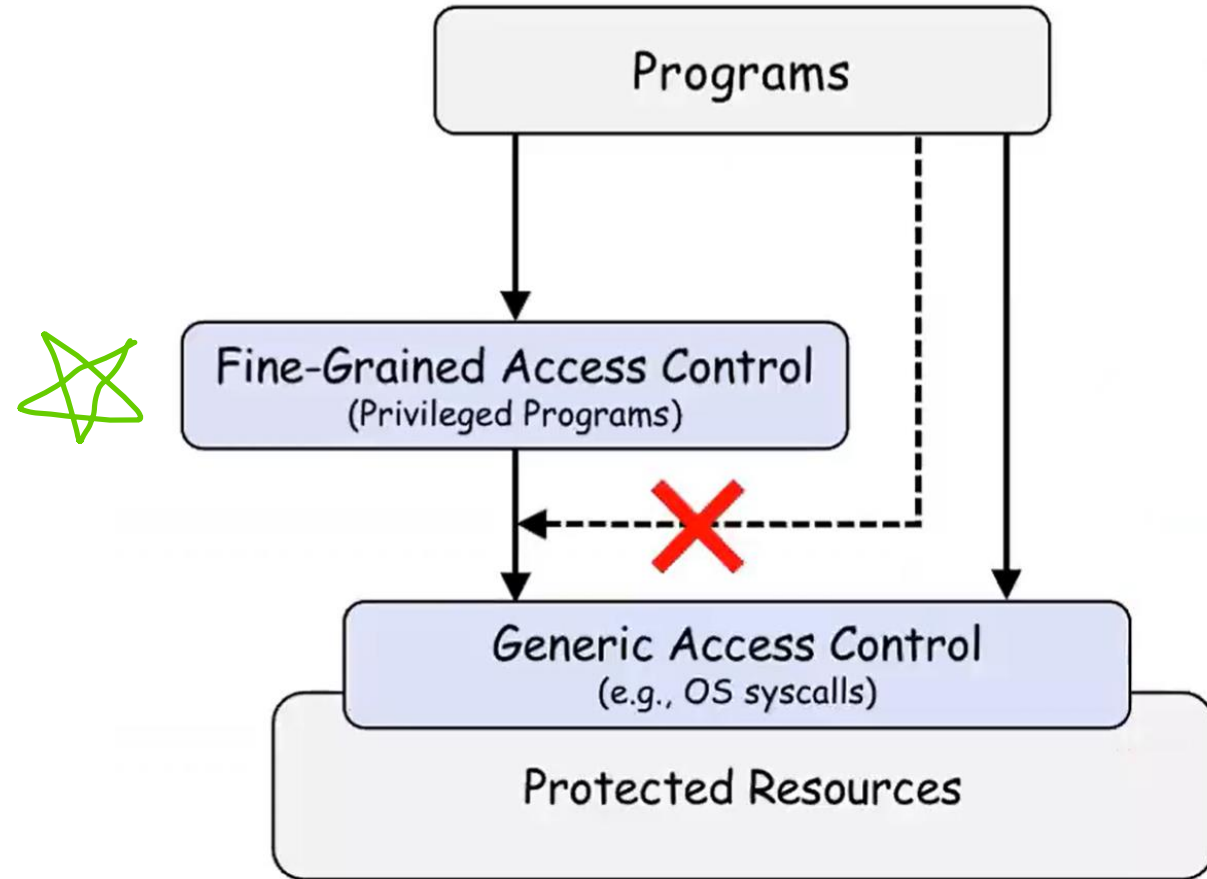
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```

`/etc/passwd` and `/etc/shadow` hold encrypted passwords for the user, in order to change our password, we will need to have access to those directories

root (aka admin) is the only person that has write permissions!

Limitations of File-Based Access Control

Instead of having a user deal with sensitive actions, lets have a privileged program do it for us!



Types of Privileged Programs

- **Daemons**

- Computer program that runs in the background
- Needs to run as root or other privileged users

- **Set-UID Programs**

- Widely used in UNIX systems
- A normal program... but marked with a special bit

The superman story

Superman got tired of saving the city every day

So, he decided to create a “super suit” that would give normal people his powers

Problem: Not all super people are good.....



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Super suit 2.0

Super suit with a dope computer

Programmed to perform a specific task

No way to deviate from the pre-programmed task



The superman story



Task: Stop Bowser

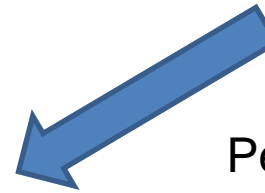
1. Fly North

2. Turn left and move forward

3. Punch



Super suit 2.0



People can hop in, and do the specific task to stop bowser

The superman story



Task: Stop Bowser

1. Fly North
2. Turn left and move forward
3. Punch

The superman story



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This works great! People can only do the predetermined task and don't have control!

The superman story



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This works great! People can only do the predetermined task and don't have control!

Exploitable?

The superman story



Task: Stop Bowser

1. Fly North
2. Turn left and move forward
3. Punch



Suppose I come along,
and I see the power suit

And I decide to flip the suit around

Now what happens???

The superman story



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Task: Stop Bowser

1. Fly North

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Suppose I come along,
and I see the power suit

And I decide to flip the suit around

I still followed the steps, but now we have a totally different outcome

My plan was to rob the bank, and I had friends waiting this whole time!

Set-UID In a Nutshell


Set-UID allows a user to run a program with the program owner's privilege

- User runs a program w/ temporarily elevated privileges

Created to deal with inflexibilities of UNIX access control

Example: The **passwd** program

```
[seed@VM][~]$ ls -al /usr/bin/passwd  
-rwsr-xr-x 1 root root 68208 May 28 2020 /usr/bin/passwd
```



Set-UID In a Nutshell

Set-UID allows a user to run a program with the program owner's privilege

- User runs a program w/ temporarily elevated privileges

Every process has two User IDs

- Real UID (RUID)– Identifies the **owner** of the process
- Effective UID (EUID)– Identifies **current privilege** of the process

When a normal program is executed

- RUID == EUID

When a Set-UID program is executed

- RUID != EUID
- EUID == ID of the program's owner



**If a program owner == root,
The program runs with root privileges**

Set-UID In a Nutshell

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- Real UID (RUID)– Identifies the **owner** of the process
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**If a program owner == root,
The program runs with root privileges**

Set-UID Program Demo


```
[seed@VM][~]$ cp /bin/cat ./mycat  
[seed@VM][~]$ sudo chown root mycat  
[seed@VM][~]$ ls -al mycat  
-rwxr-xr-x 1 root seed 43416 Jan 25 21:15 mycat
```



Change the owner of a file to root

Set-UID Program Demo

```
[seed@VM][~]$ cp /bin/cat ./mycat  
[seed@VM][~]$ sudo chown root mycat  
[seed@VM][~]$ ls -al mycat  
-rwxr-xr-x 1 root seed 43416 Jan 25 21:15 mycat
```



Change the owner of a file to root

```
[seed@VM][~]$ mycat /etc/shadow  
mycat: /etc/shadow: Permission denied
```

Running to program (normally)

Set-UID Program Demo

```
[seed@VM][~]$ cp /bin/cat ./mycat  
[seed@VM][~]$ sudo chown root mycat  
[seed@VM][~]$ ls -al mycat  
-rwxr-xr-x 1 root seed 43416 Jan 25 21:15 mycat
```

Change the owner of a file to root

```
[seed@VM][~]$ mycat /etc/shadow  
mycat: /etc/shadow: Permission denied
```

Running to program (normally)

```
[seed@VM][~]$ sudo chmod 4755 mycat  
[seed@VM][~]$ ls -al mycat  
-rwsr-xr-x 1 root seed 43416 Jan 25 21:15 mycat  
[seed@VM][~]$ mycat /etc/shadow  
root:!:18590:0:99999:7::  
daemon:*:18474:0:99999:7:::
```

Enable the Set-UID bit

We have successfully made a Set-UID program!

Announcements

Lab 1 Due **TOMORROW** 9/16 @ 11:59 PM

- Shouldn't be too bad

Lab 2 (SET-UID) is posted. due ???

- You will be able to complete it after today

Set-UID

Demo

A Set-UID program is just like any other program, except that it has a special bit set

```
[09/15/22]seed@VM:~/lab2$ cp /usr/bin/id ./myid
[09/15/22]seed@VM:~/lab2$ chown root myid★
chown: changing ownership of 'myid': Operation not permitted
[09/15/22]seed@VM:~/lab2$ sudo chown root myid
[09/15/22]seed@VM:~/lab2$ ./myid
bash: ./myid: No such file or directory
[09/15/22]seed@VM:~/lab2$ ./myid
uid=1000(seed) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
```

If the set-uid bit is enabled, the EUID is set according to the file owner

```
[09/15/22]seed@VM:~/lab2$ chmod 4755 myid
chmod: changing permissions of 'myid': Operation not permitted
[09/15/22]seed@VM:~/lab2$ sudo chmod 4755 myid★
[09/15/22]seed@VM:~/lab2$ ./myid
uid=1000(seed) gid=1000(seed) euid=0(root) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),120(lpadmin),131(lxd),132(sambashare),136(docker)
```

Steps for creating a set-uid program

1. Change file ownership to root (chown)
2. Enable the Set-uid bit (chmod)

4 = setuid bit

755 = owner r/w/x,
group/others can r/w

Access control decisions made based on EUID, not RUID !

4755

So.... Is Set-UID secure?

- Allows normal users to escalate privileges
 - This is different from directly giving escalated privileges (such as **sudo**)
 - Restricted behavior (think **power suit 2.0**)

Are there any programs that **should not** be Set-UID programs?

So.... Is Set-UID secure?

- Allows normal users to escalate privileges
 - This is different from directly giving escalated privileges (such as **sudo**)
 - Restricted behavior (think **power suft 2.0**)

Are there any programs that **should not** be Set-UID programs?

/bin/sh

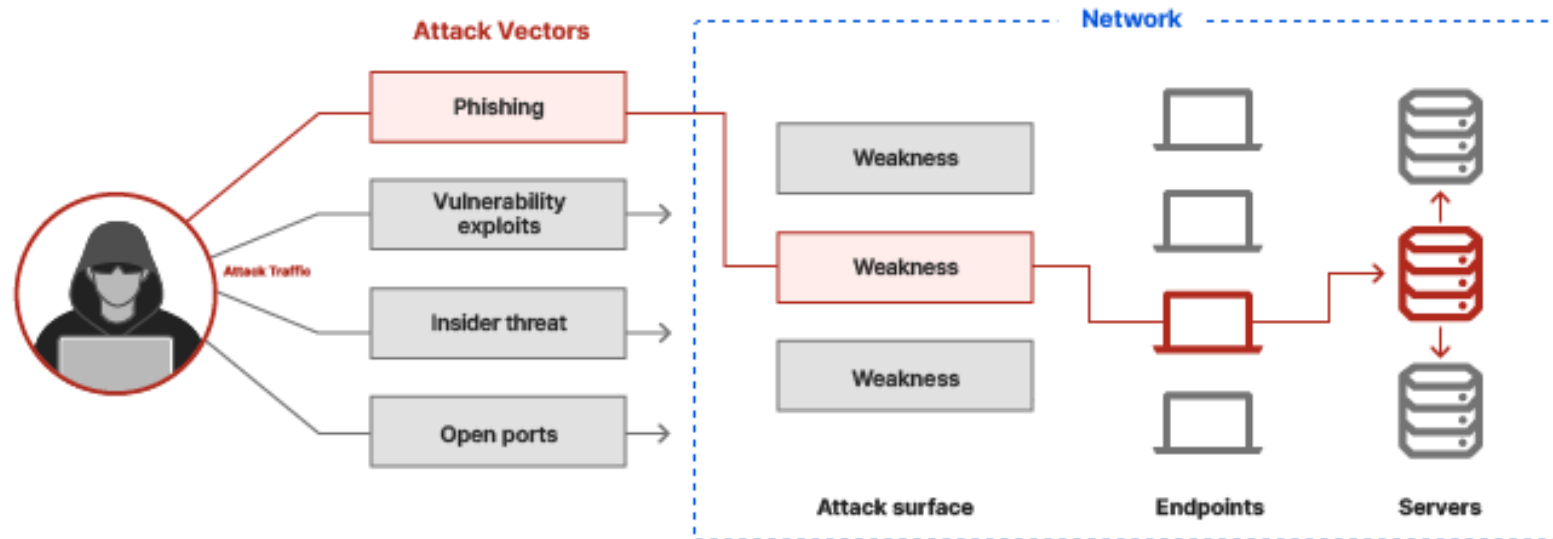
“root shell”



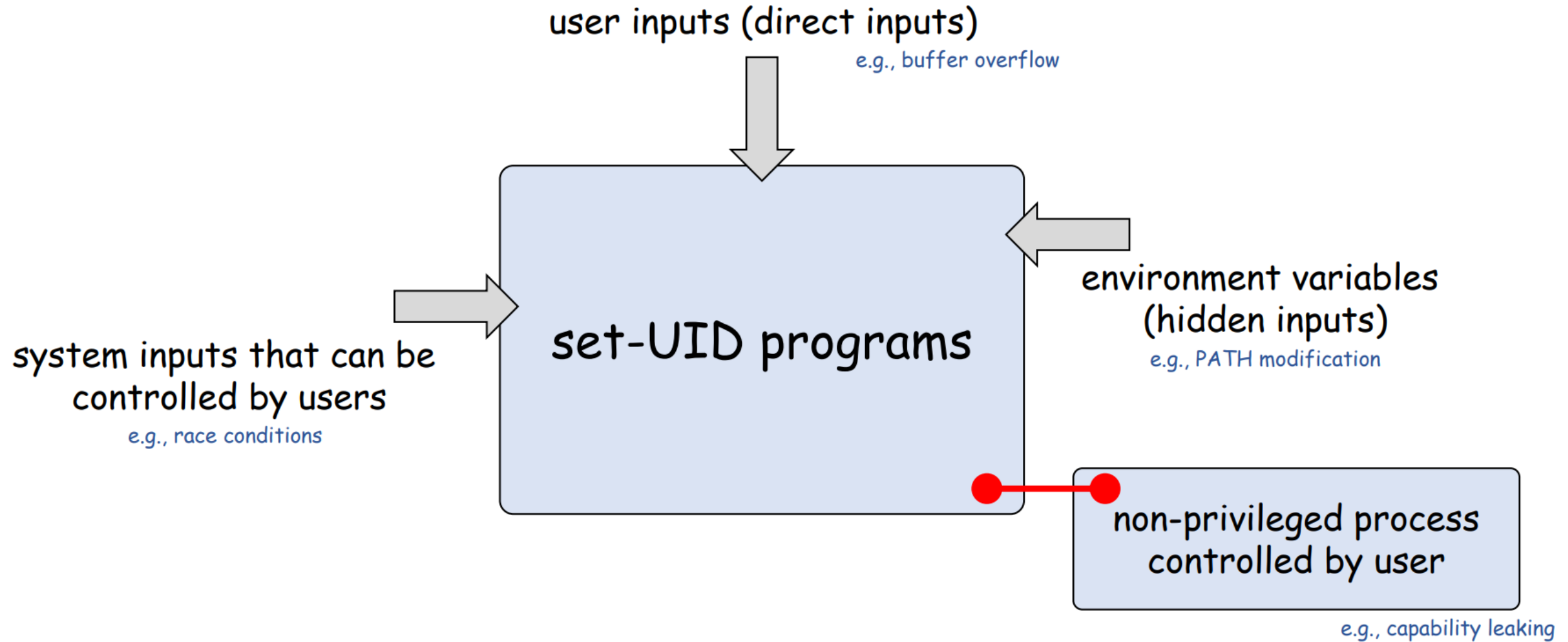
```
[09/15/22]seed@VM:~/lab2$ sudo /bin/sh
# cat /etc/shadow
root:!:18590:0:99999:7:::
daemon:*:18474:0:99999:7:::
bin:*:18474:0:99999:7:::
sys:*:18474:0:99999:7:::
```

Attack Surface of (Set-UID) Programs

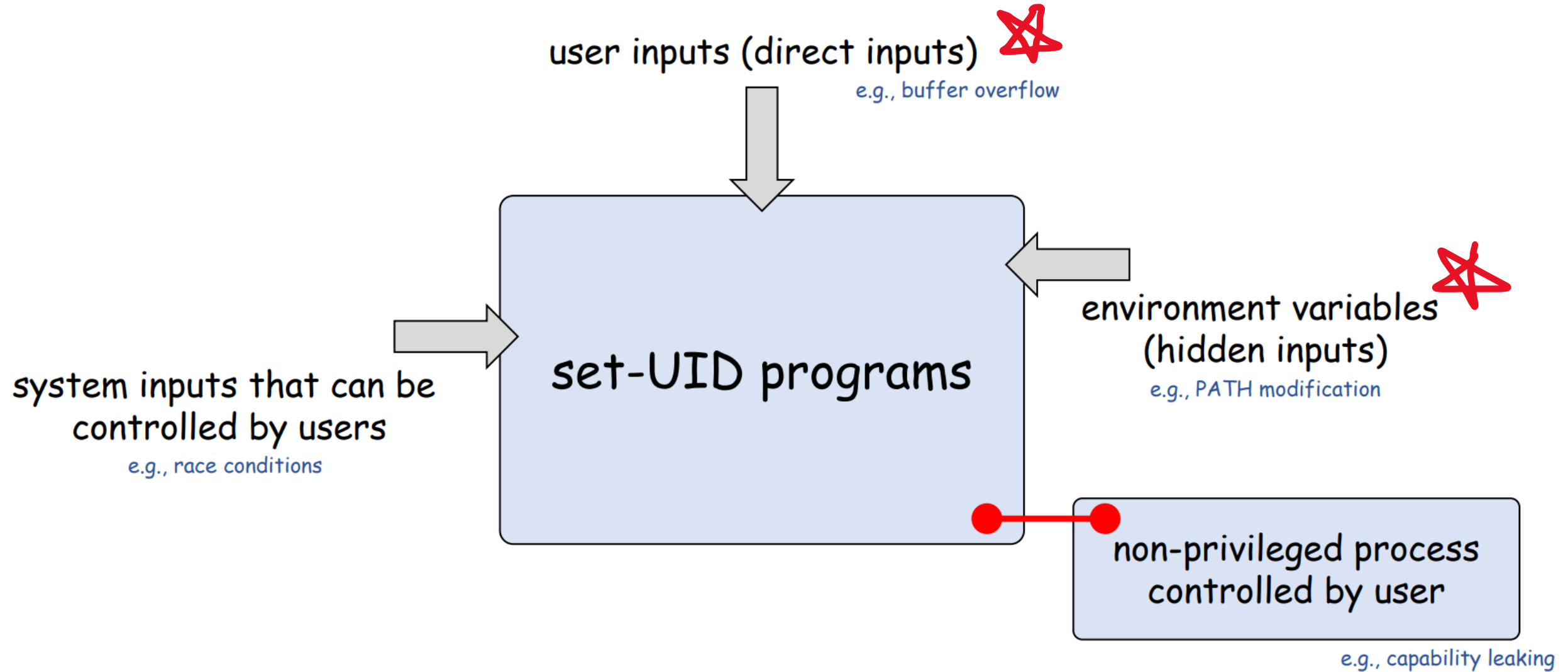
An **attack surface** is the aggregation of all exposed entry points/weaknesses into the system to gain unauthorized access



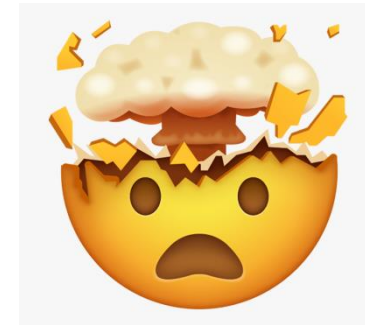
Attack Surface of (Set-UID) Programs



Attack Surface of (Set-UID) Programs




Invoking Programs from within programs



Preliminary setup for attack

/bin/sh is an alias for
/bin/dash. /bin/dash has
countermeasures for some of our
attacks

We will need to run a command to set the unsafe version of shell



```
[seed@VM][~]$ sudo ln -sf /bin/zsh /bin/sh # set shell to zsh (no countermeasure)  
[seed@VM][~]$ sudo ln -sf /bin/dash /bin/sh # set shell to dash (has countermeasure)
```



Invoking Program with a program

We can invoke external commands/programs from INSIDE another program

- `system()`
- `exec()`-family


System()

usage: **system** (*command*)

- Spawns a new process that executes the shell command that is specified in *command*

```
#include <stdlib.h>
#include <stdio.h>

int main()
{
    printf("I am going to start the calculator program! \n");
    system("/bin/bc");
}
```



- Suppose you are preparing for an audit. An auditor may need the access to view certain files.
- We will create a privileged program that will let the auditor view the content some file

```
./audit company_data.csv
```

```
./audit ../lab0/solution.docx
```

```
#include <string.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main(int argc, char *argv[])
{
    char *v[3];

    if (argc < 2) {
        printf("Audit! Please type a file name.\n");
        return 1;
    }

    v[0] = "/bin/cat"; v[1] = argv[1]; v[2] = 0;
    char *command = malloc(strlen(v[0]) + strlen(v[1]) + 2);
    sprintf(command, "%s %s", v[0], v[1]);

    /*
     * Use only one of the following (comment out the other):
     */

    system(command);
    //execve(v[0], v, 0);

    return 0;
}
```

*The command line argument
(file path) is appended to the
string "/bin/cat"*

Spawns a new process that executes

/bin/cat [FILE_PATH]

ex. /bin/cat my_file.txt

- Suppose you are preparing for an audit. An auditor may need the access to view certain files.
- We will create a privileged program that will let the auditor view the content some file

```
./audit company_data.csv
```



Command that
is executed: /bin/cat company_data.csv

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```



/bin/cat ../lab0/solution.docx

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Command that
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/bin/cat ../lab0/solution.docx

We have some control over the behavior of the program

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- We will create a privileged program that will let the auditor view the content some file

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./audit company_data.csv
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Command that
is executed: /bin/cat company_data.csv

```
./audit ../lab0/solution.docx
```



/bin/cat ../lab0/solution.docx

We have some control over the behavior of the program

If this is a Set-UID program.... things could get interesting

```
#include <string.h>
#include <stdio.h>
#include <stdlib.h>
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int main(int argc, char *argv[])
{
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    v[0] = "/bin/cat"; v[1] = argv[1]; v[2] = 0;
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    return 0;
}
```

System is a very unsafe function

What type of input could we provide to exploit this?


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#include <string.h>
#include <stdio.h>
#include <stdlib.h>
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    }

    v[0] = "/bin/cat"; v[1] = argv[1]; v[2] = 0;
    char *command = malloc(strlen(v[0]) + strlen(v[1]) + 2);
    sprintf(command, "%s %s", v[0], v[1]);

    /*
     * Use only one of the following (comment out the other):
     */

    system(command);
    //execve(v[0], v, 0);

    return 0;
}
```

System is a very unsafe function

What type of input could we provide to exploit this?

hint: the string passed into system can include multiple commands

```
#include <string.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main(int argc, char *argv[])
{
    char *v[3];

    if (argc < 2) {
        printf("Audit! Please type a file name.\n");
        return 1;
    }

    v[0] = "/bin/cat"; v[1] = argv[1]; v[2] = 0;
    char *command = malloc(strlen(v[0]) + strlen(v[1]) + 2);
    sprintf(command, "%s %s", v[0], v[1]);

    /*
     * Use only one of the following (comment out the other):
     */

    system(command);
    //execve(v[0], v, 0);

    return 0;
}
```

System is a very unsafe function

What type of input could we provide to exploit this?

hint: the string passed into system can include multiple commands

./audit "my_info.txt; /bin/sh"

```
./audit "my_info.txt; /bin/sh"
```



```
system(/bin/cat "my_info.txt; /bin/sh")
```

```
./audit "my_info.txt; /bin/sh"
```



```
system(/bin/cat "my_info.txt; /bin/sh")
```

```
[09/15/22]seed@VM:~/lab2$ ./audit "my_info.txt; /bin/sh"  
I have some information  
#
```

System() interprets this as two different commands

```
./audit "my_info.txt; /bin/sh"
```



```
system (/bin/cat my_info.txt; /bin/sh)
```

```
[09/15/22]seed@VM:~/lab2$ ./audit "my_info.txt; /bin/sh"  
I have some information  
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./audit "my_info.txt; /bin/sh"
```



```
system (/bin/cat my_info.txt; /bin/sh)
```

```
[09/15/22]seed@VM:~/lab2$ ./audit "my_info.txt; /bin/sh"
```

```
I have some information
```

```
# whoami
```

```
root
```

```
# cat /etc/shadow
```

```
root:!:18590:0:99999:7:::
```

```
daemon:*:18474:0:99999:7:::
```

```
bin:*:18474:0:99999:7:::
```



Because this is a Set-UID program.

When owner = root, the shell will be run with root permissions

```
./audit "my_info.txt; /bin/sh"
```



```
system (/bin/cat my_info.txt; /bin/sh)
```

```
[09/15/22]seed@VM:~/lab2$ ./audit "my_info.txt; /bin/sh"
```

```
I have some information
```

```
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```
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```
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```

```
root:!:18590:0:99999:7:::
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```

```
bin:*:18474:0:99999:7:::
```

Because this is a Set-UID program.

When owner = root, the shell will be run with root permissions

We have gained access into the system!

A safer way to invoke programs

```
int execve(const char *pathname, char *const argv[], char *const envp[]);
```

`execve()` executes the program referred to by *pathname*. `argv[]` is the command line arguments for the command

A safer way to invoke programs

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`execve()` executes the program referred to by *pathname*. *argv[]* is the command line arguments for the command

Using `execve()` instead of `system()`

```
[09/15/22] seed@VM:~/lab2$ ./audit "aa;/bin/sh"  
/bin/cat: 'aa;/bin/sh': No_such file or directory
```

Fail!

A safer way to invoke programs

```
int execve(const char *pathname, char *const argv[], char *const envp[]);
```

`execve()` executes the program referred to by `pathname`. `argv[]` is the command line arguments for the command

Using `execve()` instead of `system()`

```
[09/15/22] seed@VM:~/lab2$ ./audit "aa;/bin/sh"  
/bin/cat: 'aa;/bin/sh': No such file or directory
```

Fail!

```
execve("/bin/cat", ["aa;/bin/sh"])
```



/bin/cat "aa;/bin/sh"

Treated as an entire argument to the command

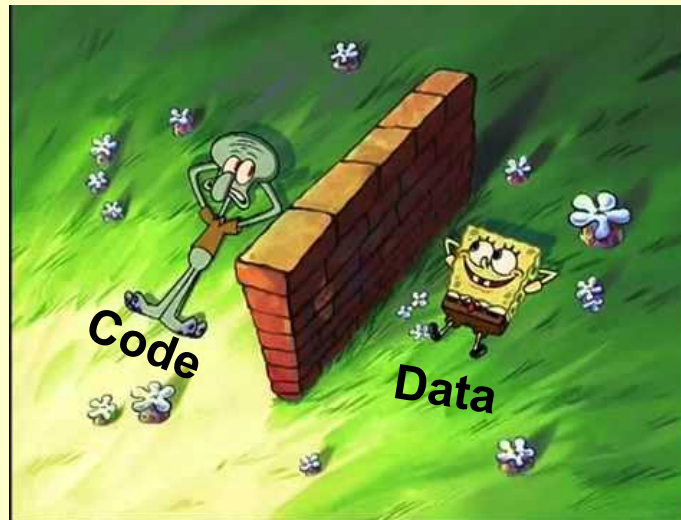
Fail!

What makes this unsafe? Why was this program exploitable?

Principle of Isolation

There needs to be a clear separation of **data** and **code**

If user input needs to be treated as data, NONE of the contents should be treated as code

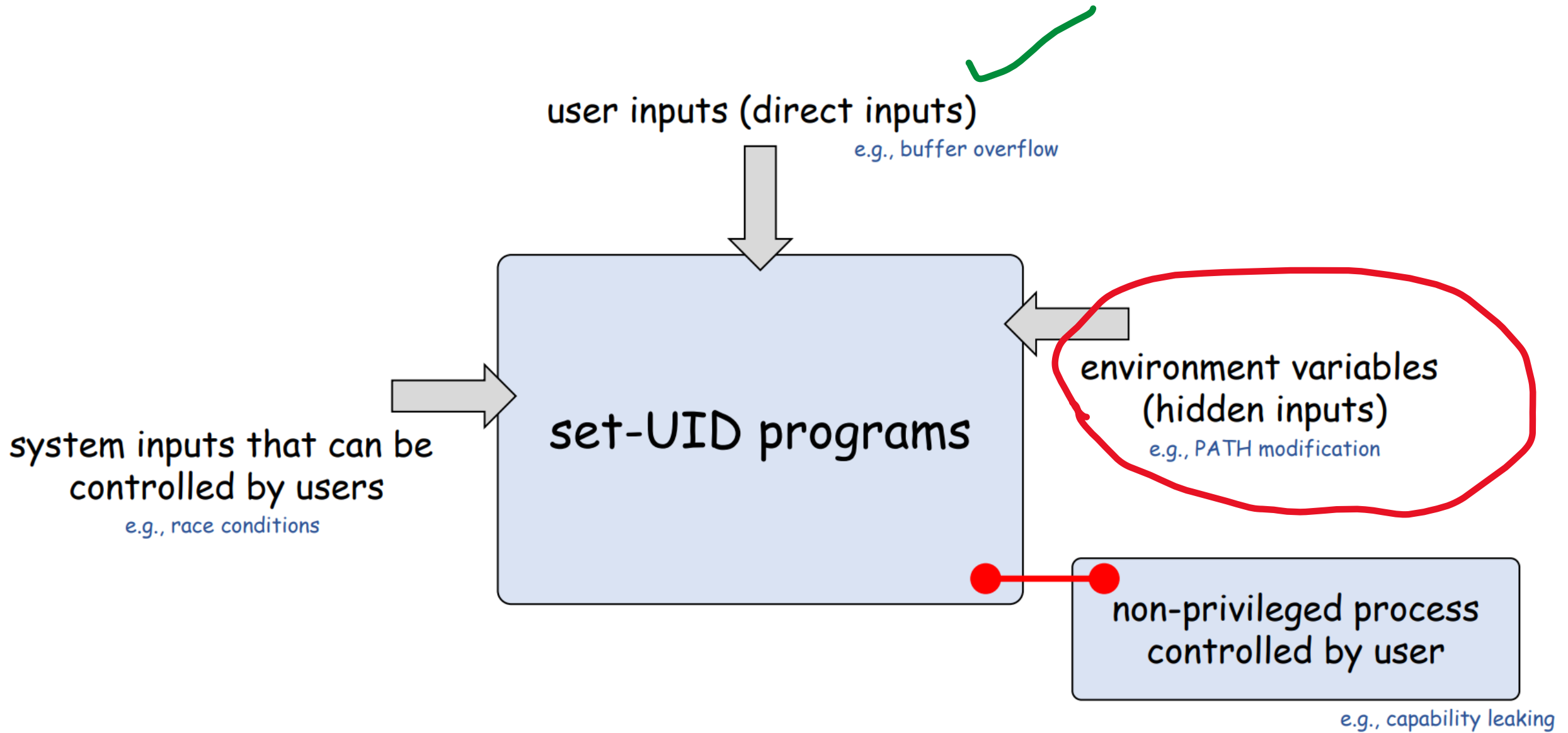


The ability (and risks) of invoking external commands is not limited to C

Python has a `system` call...

Perl has `open()`

PHP has `system`



Environment variable are a set of dynamic named values that affect the way a running process will behave *(key-value pairs)*

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Example: The PATH variable

- We use command such as `ls` and `passwd`

We could be in any directory; how does it know to run `/bin/ls` ?

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- We use command such as `ls` and `passwd`

We could be in any directory; how does it know to run `/bin/ls` ?

If the full path is not provided, the shell process will use the PATH env. Variable to search for it!

```
|PATH=/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr/bin:/sbin:/bin:/usr/games:/usr/local/games:/snap/bin:.
```

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If the full path is not provided, the shell process will use the PATH env. Variable to search for it!

```
|PATH=/usr/local/sbin:/usr/local/bin:/usr/sbin:/usr/bin:/sbin:/bin:/usr/games:/usr/local/games:/snap/bin:.
```

You can run the `env` command to print out all the environment variables

Where do Env Variables come from?

Processes can get environment variables in one of two ways:

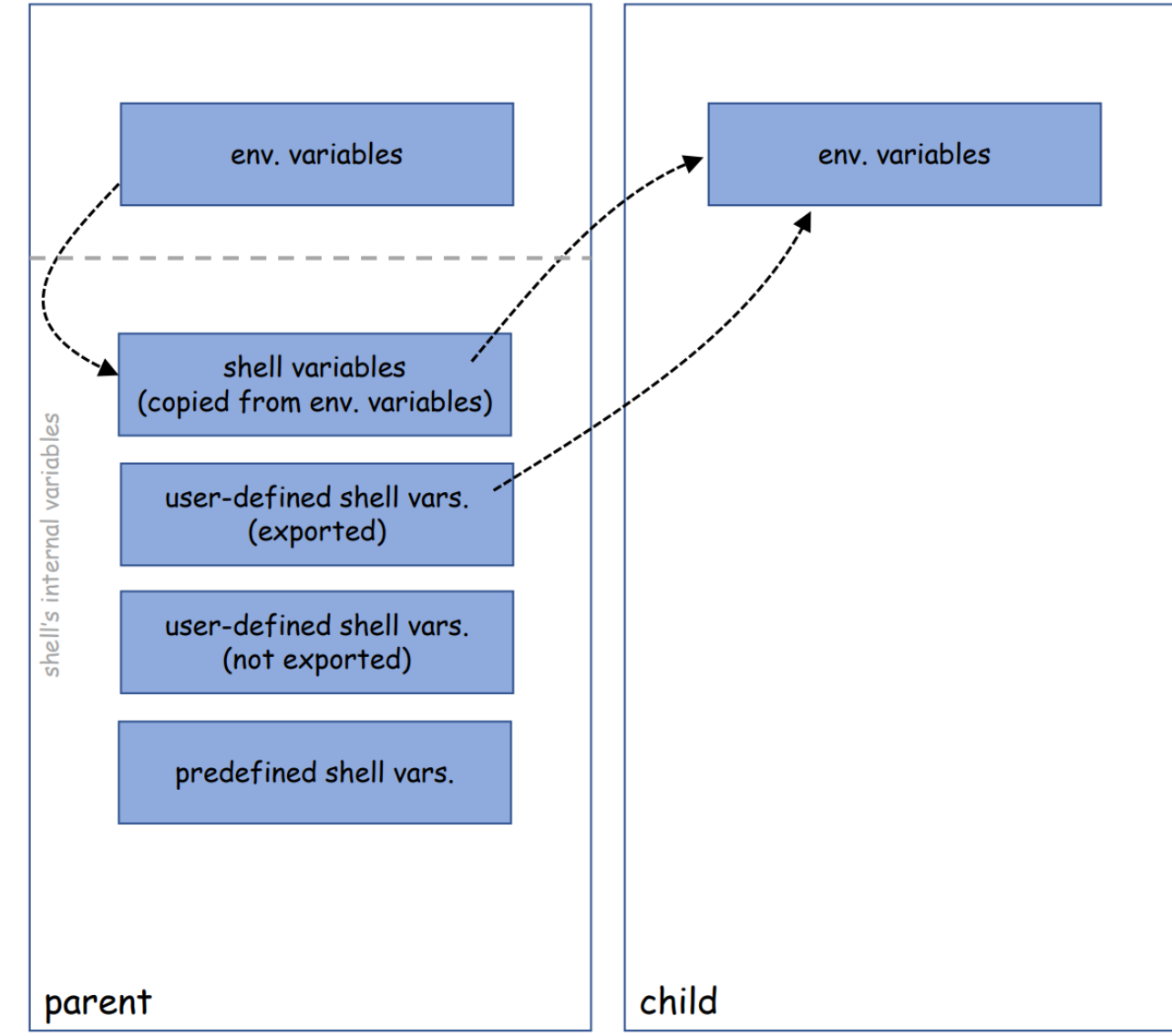
fork() → the child process inherits its parent process's environment variables.

exec() → the memory space is overwritten, and all old environment variables are lost. However, **execve()** can explicitly pass environment variables from one process to another

```
./passenv 1  
./passenv 2  
./passenv 3
```

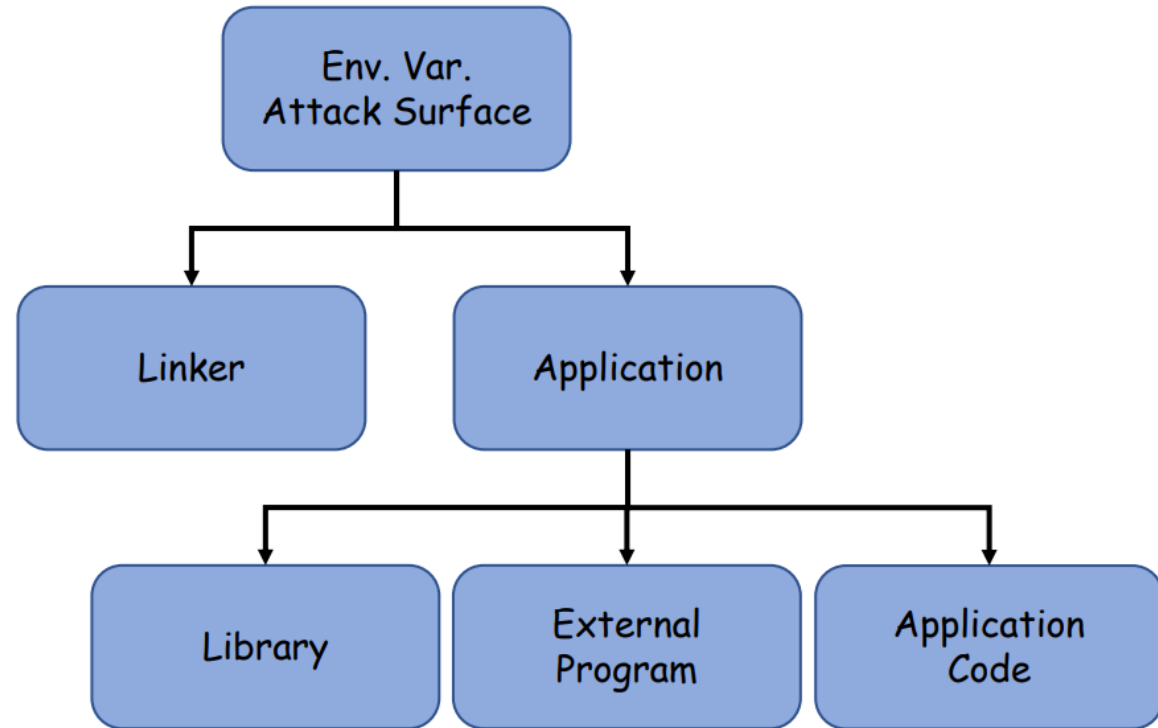
```
1  #include <stdio.h>  
2  #include <unistd.h>  
3  
4  extern char ** environ;  
5  void main(int argc, char* argv[], char* envp[])  
6  {  
7      int i = 0; char* v[2]; char* newenv[3];  
8      if (argc < 2) return;  
9  
10     // Construct the argument array  
11     v[0] = "/usr/bin/env"; v[1] = NULL;  
12  
13     // Construct the environment variable array  
14     newenv[0] = "AAA=aaa"; newenv[1] = "BBB=bbb"; newenv[2] = NULL;  
15  
16     switch(argv[1][0]) {  
17         case '1': // Passing no environment variable.  
18             execve(v[0], v, NULL);  
19         case '2': // Passing a new set of environment variables.  
20             execve(v[0], v, newenv);  
21         case '3': // Passing all the environment variables.  
22             execve(v[0], v, environ);  
23         default:  
24             execve(v[0], v, NULL);  
25     }  
26 }
```

Where do Env Variables come from?



Attacks with Env Variables

- Hidden usage of environment variable is part of what makes them so dangerous
- Users can also modify environment variables...
- If Set-UID programs make use of environment variables, they become part of the attack surface



Attacks with Env Variables

- A program may invoke an external program (e.g., via `system()`) to do some work
- **PATH** contains a list of directories to search for executable programs
- If a program is invoked without using the absolute path (e.g., `system("ls")`), the **PATH** env. variable is used to search for the program
- **PATH** can be set by users....

Any ideas???



Attacks with Env Variables

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- **PATH** can be set by users....

Any ideas???



Task 6 in Lab 2 ☺

Attacks with Env Variables

```
[seed@VM][~]$ sudo ln -sf /bin/zsh /bin/sh # set shell to zsh (no countermeasure)
[seed@VM][~]$ sudo ln -sf /bin/dash /bin/sh # set shell to dash (has countermeasure)
```

Compile and set as Set-UID program

```
[seed@VM][~]$ gcc -o ls_vul ls_vul.c
[seed@VM][~]$ sudo chown root ls_vul
[seed@VM][~]$ sudo chmod 4755 ls_vul
[seed@VM][~]$ ls -al ls_vul
...
```

Q: How to get ls_vul to run attacker code for ls instead of /bin/ls program?!?!?

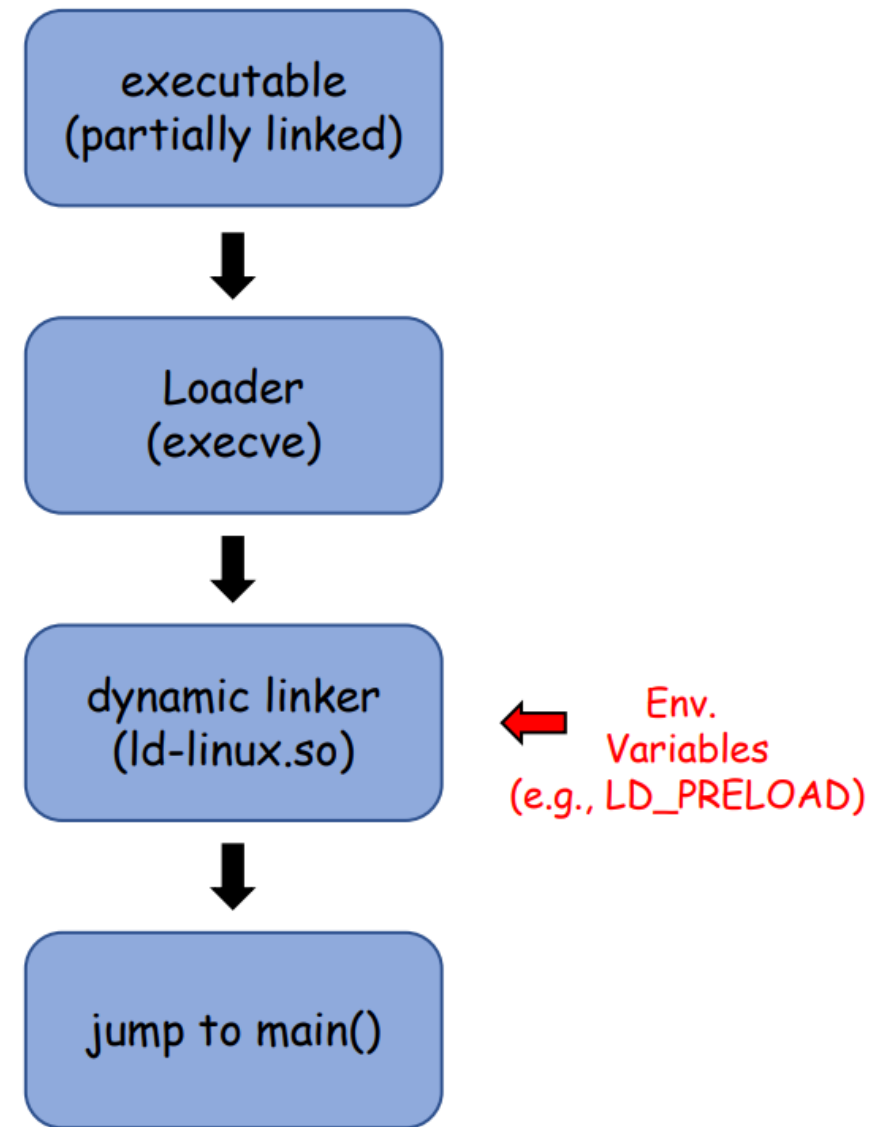
```
[seed@VM][~]$ export PATH=/home/seed:$PATH # set PATH to look in seed's home dir first...
[seed@VM][~]$ echo $PATH
...
```

...and now...



export is used to define new variables

- **Linking** finds the external library code referenced in a program
- **Static linking** – linker combines program code/external code into final executable
- **Dynamic Linking**- linker uses env. Variables to locate external dependencies (increase the attack surface)



How does the linker know where to look?

- **LD_PRELOAD** contains a list of shared libraries to search first
- Provides precedent over standard function calls (malloc, free, etc)
- If functions are not found, it will consult the location specified in **LD_LIBRARY_PATH**
- We can set both values, which gives us an opportunity for users to influence linking



Task 7

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```
/* myprog.c */  
#include <unistd.h>  
int main()  
{  
    sleep(1);  
    return 0;  
}
```

This will run the standard libc sleep

mylib.c

```
#include <stdio.h>  
void sleep (int s)  
{  
    /* If this is invoked by a privileged  
    printf("I am not sleeping!\n");  
}
```

Write our own sleep
function within our mylib.c
program

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Add function to shared library

```
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Set variable to look at our library, not default one

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./myprog.c ???


```
#include <fcntl.h>
```

```
int main()
{
    int fd;

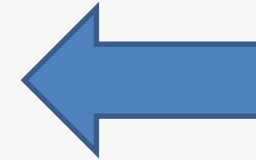
    /*
     * Assume that /etc/zxx is an important system file,
     * and it is owned by root with permission 0644.
     * Before running this program, you should create
     * the file /etc/zxx first.
     */
    fd = open("/etc/zxx", O_RDWR | O_APPEND);
    if (fd == -1) {
        printf("Cannot open /etc/zxx\n");
        exit(0);
    }

    // Simulate the tasks conducted by the program
    sleep(1);
}
```

```
/* After the task, elevated privileges are no longer needed;
 * it is time to relinquish these privileges!
 * NOTE: getuid() returns the real UID (RUID)
 */
setuid(getuid());
```

```
if (fork()) { /* parent process */
    close (fd);
    exit(0);
} else { /* child process */
```

```
/*
 * Now, assume that the child process is compromised, and that
 * malicious attackers have injected the following statements into this process
 */
write (fd, "Malicious Data\n", 15);
close (fd);
}
```



We eventually fork and create a new process

Capability leaking occurs when some process gains escalated privileges, but does not clean up the privileges when downgraded

```
#include <fcntl.h>
```

```
int main()
{
    int fd;

    /*
     * Assume that /etc/zxx is an important system file,
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        exit(0);
    }

    // Simulate the tasks conducted by the program
    sleep(1);
}
```

fd is defined
before the fork()

/etc/zxx is only
writeable by root

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Drop privileges

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if (fork()) { /* parent process */
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```
#include <fcntl.h>
```

```
int main()
{
    int fd;

    /*
     * Assume that /etc/zzz is an important system file,
     * and it is owned by root with permission 0644.
     * Before running this program, you should create
     * the file /etc/zzz first.
     */
    fd = open("/etc/zzz", O_RDWR | O_APPEND);
    if (fd == -1) {
        printf("Cannot open /etc/zzz\n");
        exit(0);
    }

    // Simulate the tasks conducted by the program
    sleep(1);
}
```

fd is defined
before the fork()

/etc/zzz is only
writeable by root

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if (fork()) { /* parent process */
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Close file in
parent process

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 */
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close (fd);
}
```

The file descriptor is still open in
the child process!

```
[09/15/22] seed@VM: ~/lab2$ sudo touch /etc/zxx
[09/15/22] seed@VM: ~/lab2$ ./cap_leak
[09/15/22] seed@VM: ~/lab2$ sudo cat /etc/zxx
Malicious Data
```

Principle of Least Privilege

Subjects and Programs should be given only the privileges needed to complete their task

Disable privileges when they aren't needed



Lab 2