

CSCI 476: Computer Security

Lecture 3: Operating Systems Review

Reese Pearsall
Fall 2022

Announcements

TA

- **Karishma Rahman**
- karishma.rahman.bd@gmail.com
- Office Hours: Tuesdays 1:00 pm to 3:00 pm
- Location: Barnard 259

Lab 0 due **tonight** at 11:59 PM

Lectures are all recorded

Raytheon cyber security talk tonight @ 5PM in Barnard Hall 347

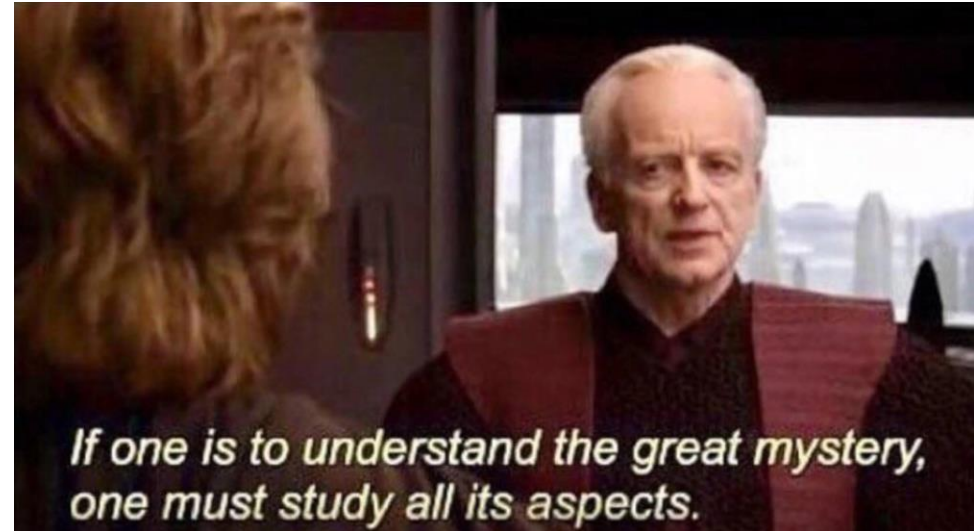
Announcements

[illegible]

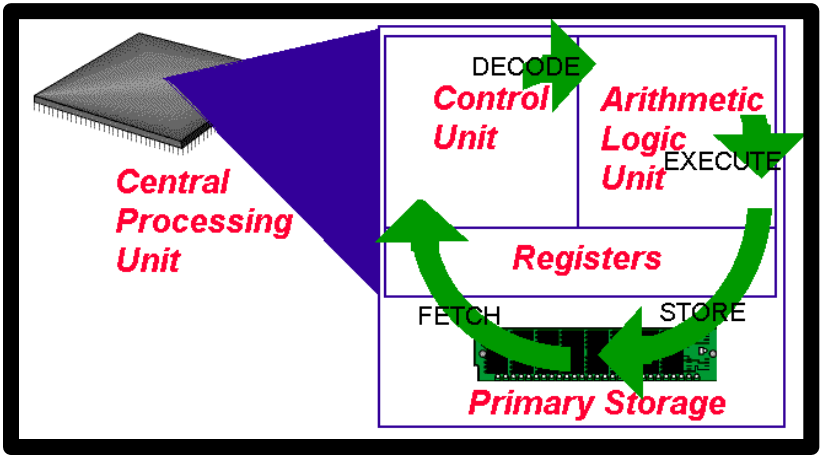
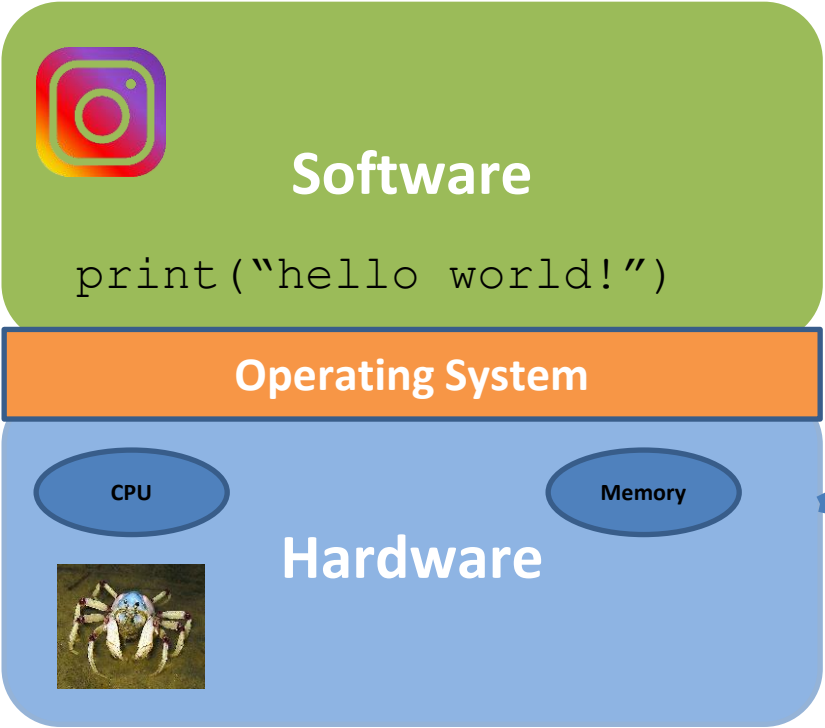
Operating Systems Review

To understand the technical aspects of security, we must have a good understanding of how computers work

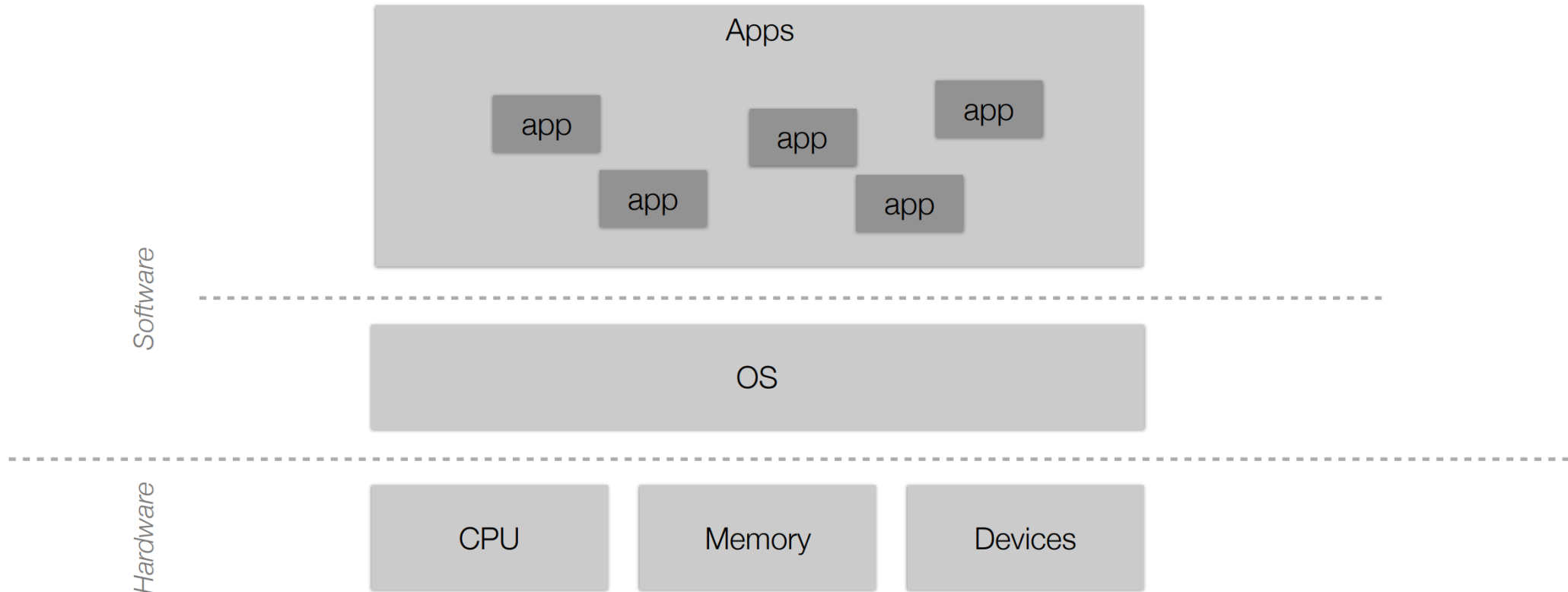
operating systems



Operating Systems Review



Operating Systems Review

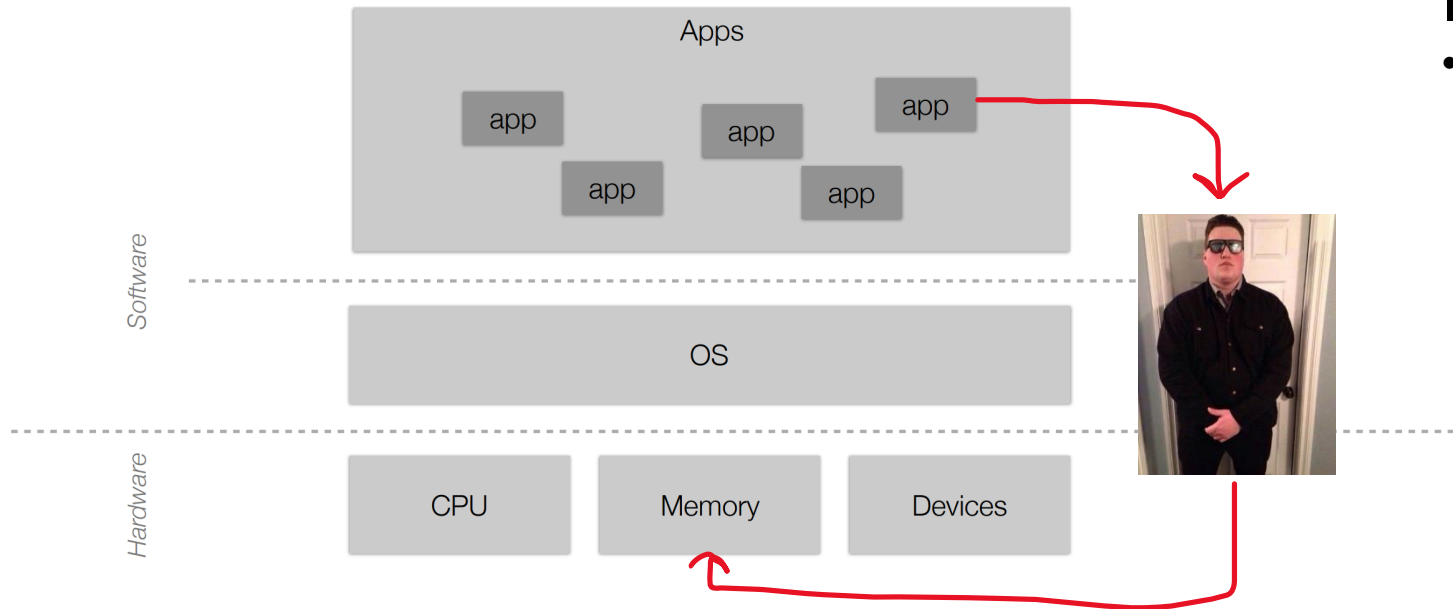


Operating Systems Review

Responsibilities of the OS?

Interface Manager

- Manages communication between apps and hardware



Operating Systems Review

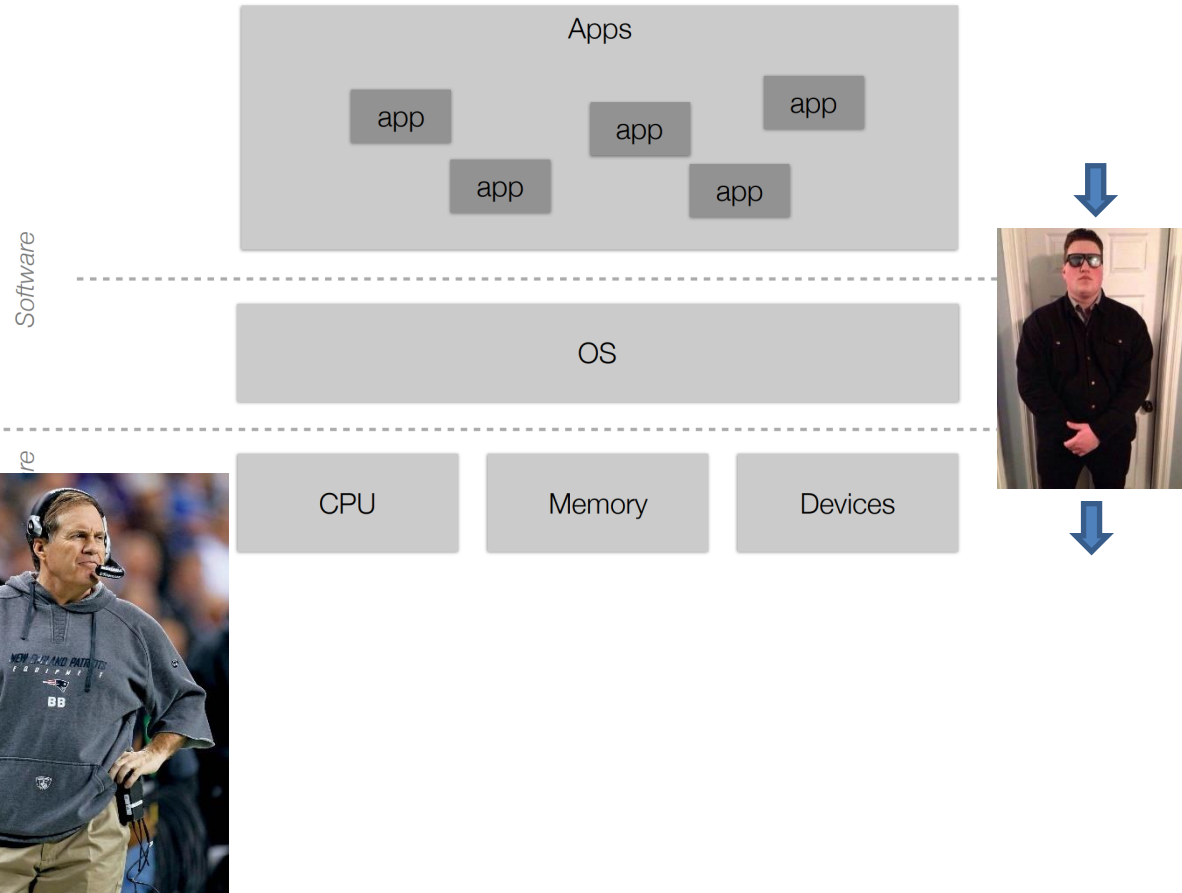
Responsibilities of the OS?

Interface Manager

- Manages communication between apps and hardware

Process Manager

- Manages how processes are structured and how to handle many processes running at once



Operating Systems Review

Responsibilities of the OS?

Interface Manager

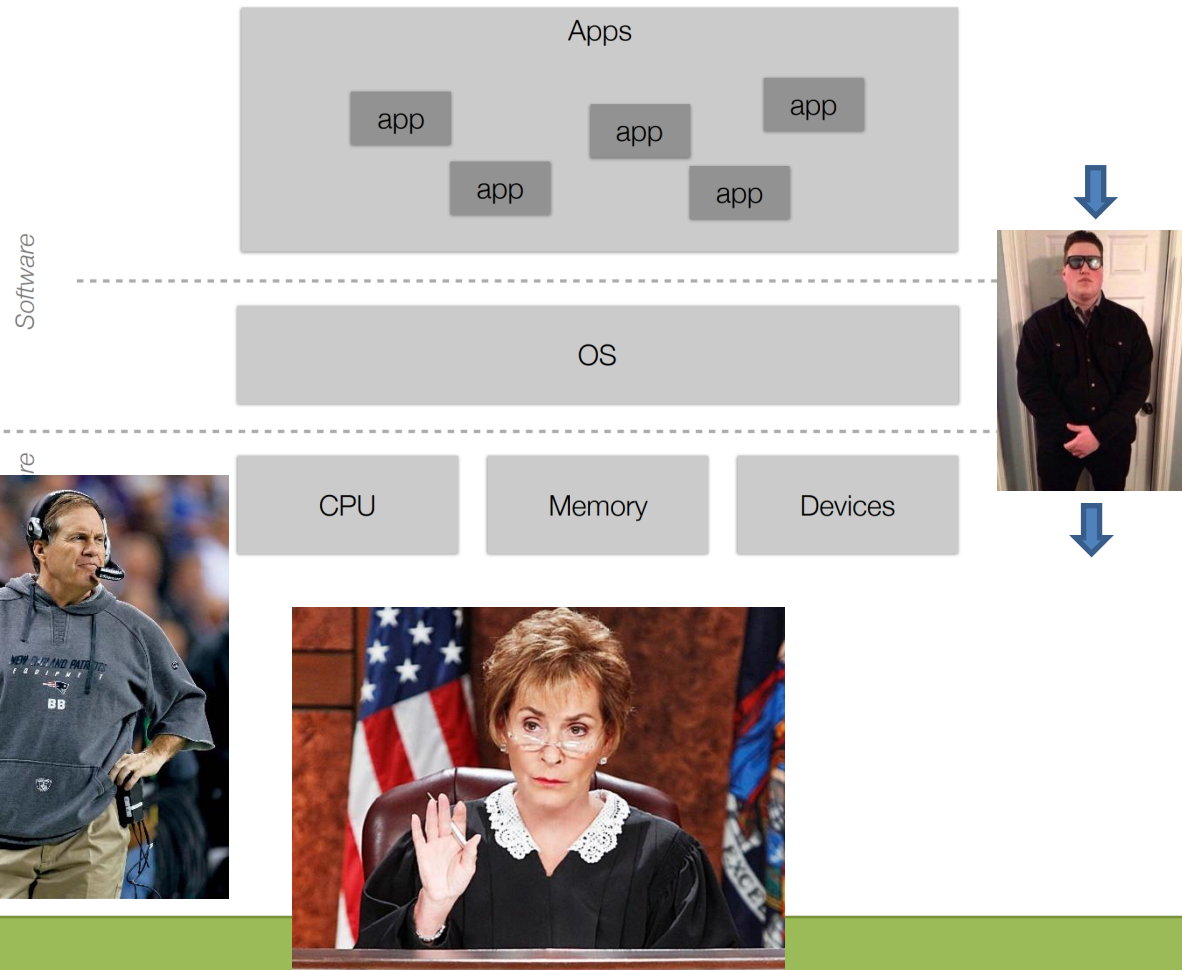
- Manages communication between apps and hardware

Process Manager

- Manages how processes are structured and how to handle many processes running at once

Traffic Manager

- Manages which programs should be executed by the CPU



Operating Systems Review

Responsibilities of the OS?

Interface Manager

- Manages communication between apps and hardware

Process Manager

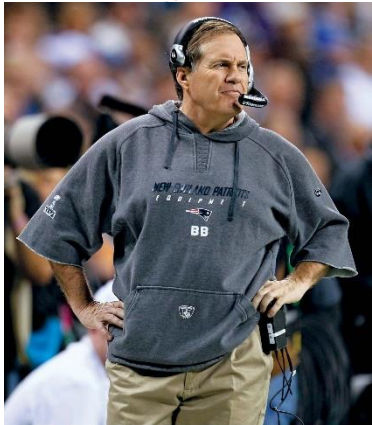
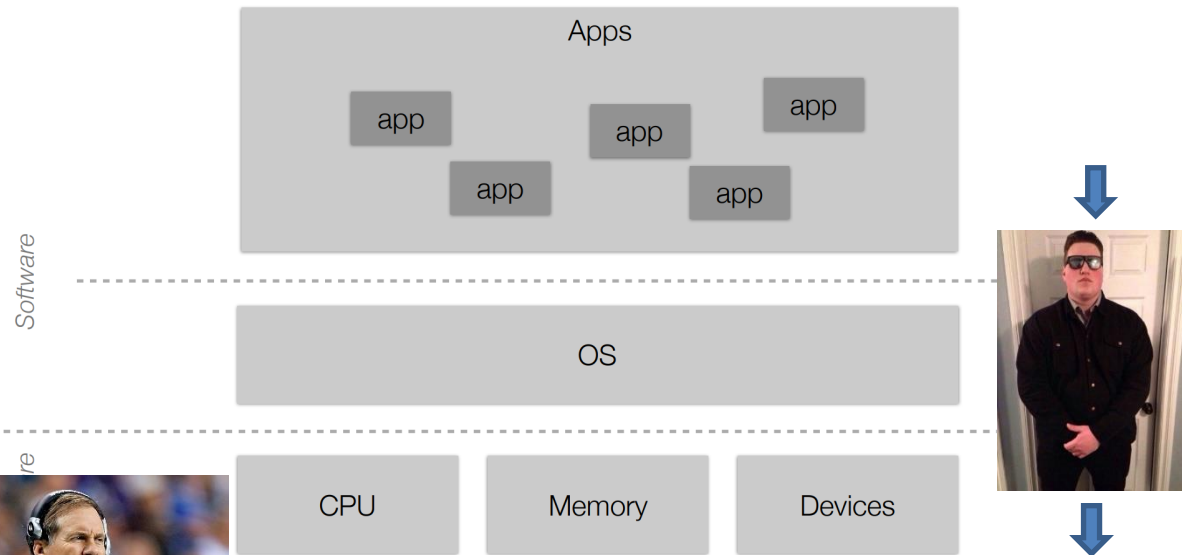
- Manages how processes are structured and how to handle many processes running at once

Traffic Manager

- Manages which programs should be executed by the CPU

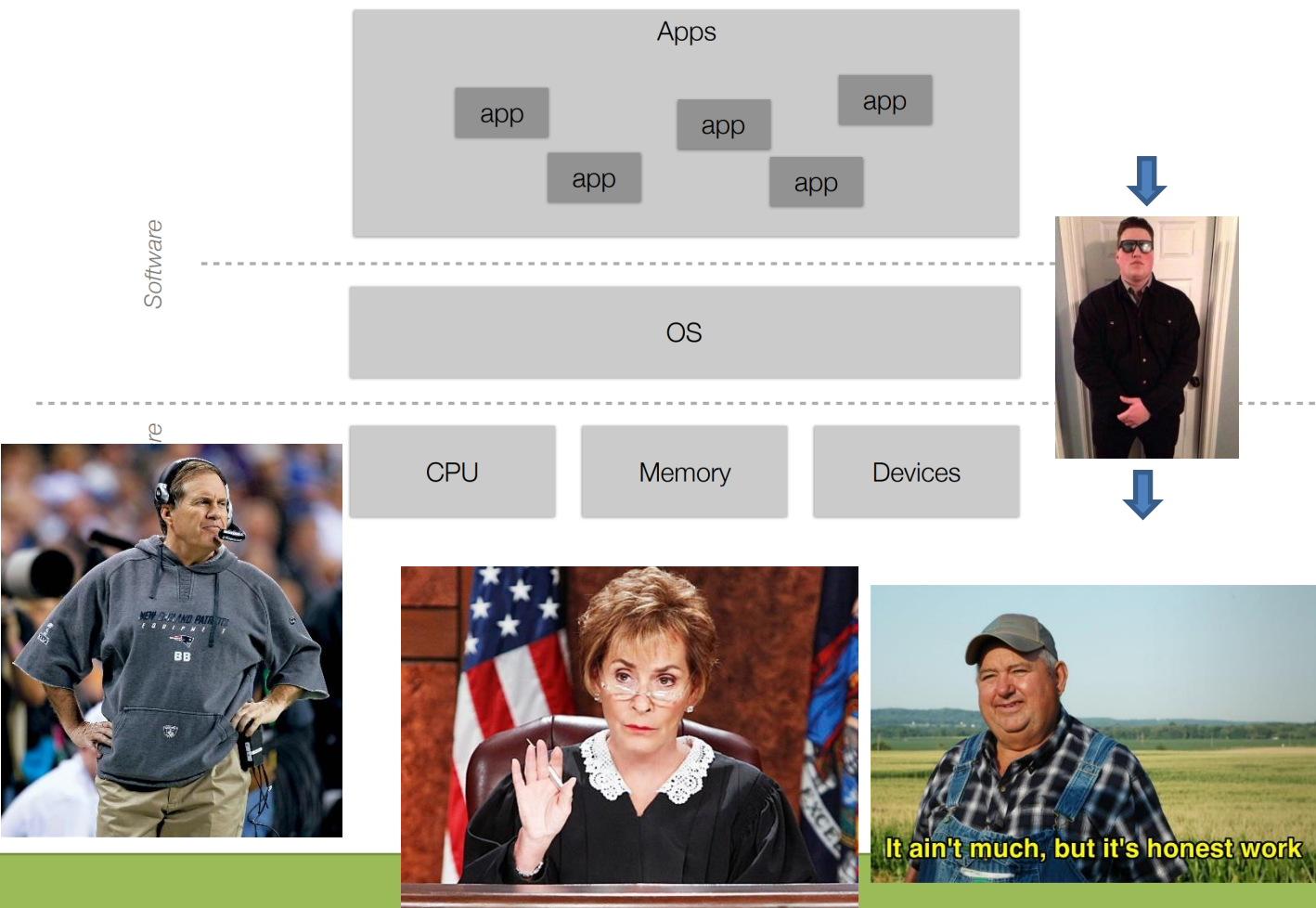
Memory Manager

- Manages how physical memory is utilized



Operating Systems Review

Responsibilities of the OS?



Interfacing

- Managing hardware

Providing

- Managing memory

Managing

- Managing execution

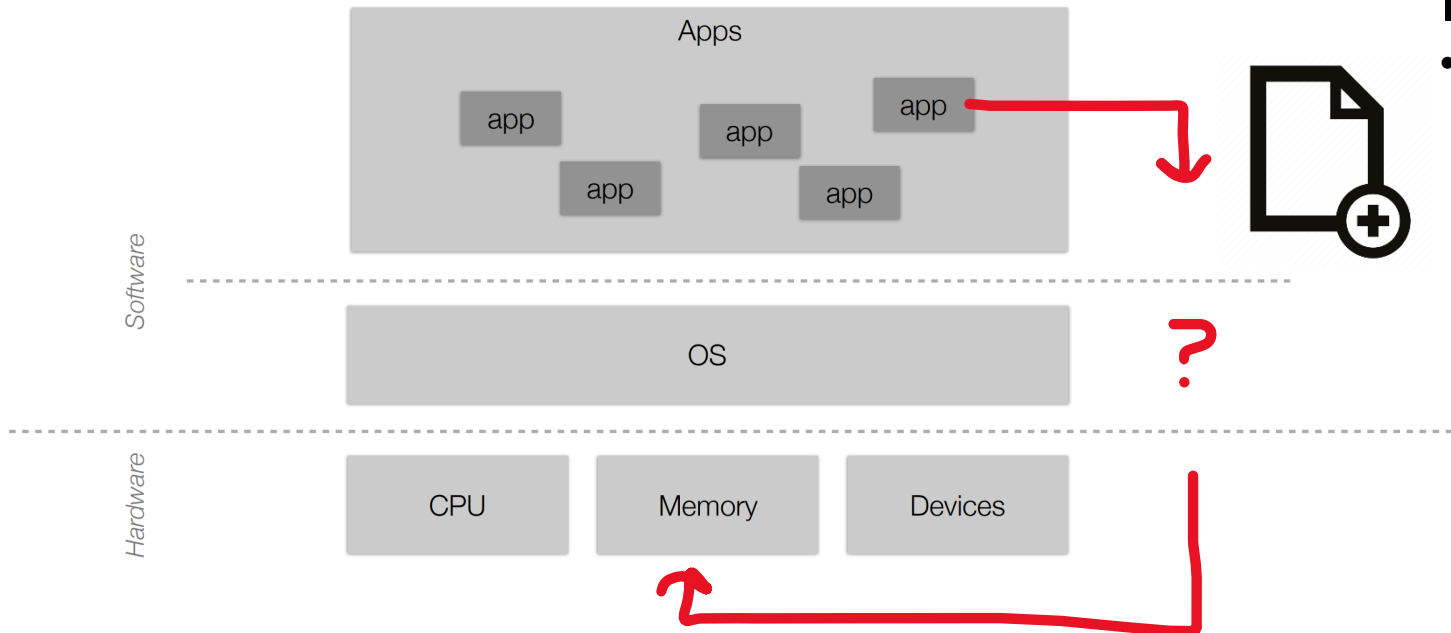
Managing

- Managing memory



Operating Systems Review

Responsibilities of the OS?



Interface Manager

- Manages communication between apps and hardware

How does an application get access to a computer's resources?



Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
int main(void)
{
    printf("Hello, World!\n");

    return 0;
}
```


Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
int main(void)
{
    printf("Hello, World!\n");

    return 0;
}
```



```
int main(void)
{
    write(1, "Hello, World!\n", 14);

    return 0;
}
```



```
int main(void)
{
    syscall(SYS_write, 1, "Hello, World!\n", 14);

    return 0;
}
```

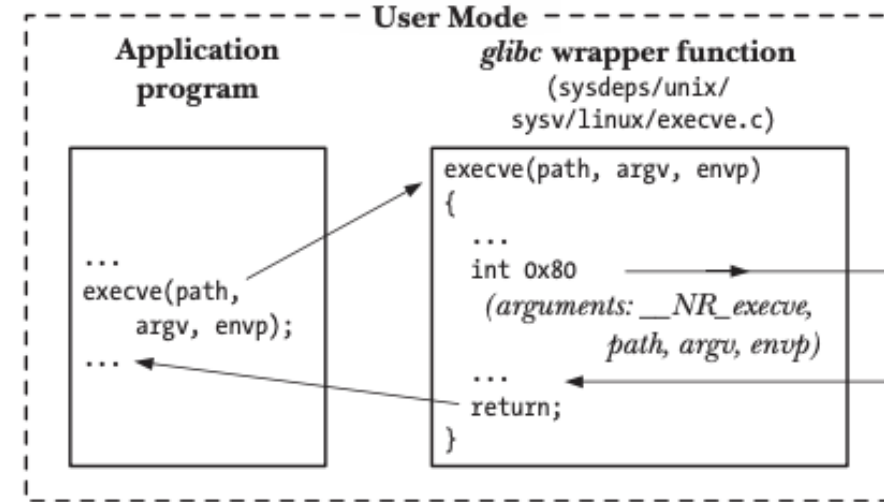
Number	Name	Description
1	exit	terminate process execution
2	fork	fork a child process
3	read	read data from a file or socket
4	write	write data to a file or socket
5	open	open a file or socket
6	close	close a file or socket
37	kill	send a kill signal
90	old_mmap	map memory
91	munmap	unmap memory
301	socket	create a socket
303	connect	connect a socket

Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0], name, NULL);
    return 0;
} syscall
```

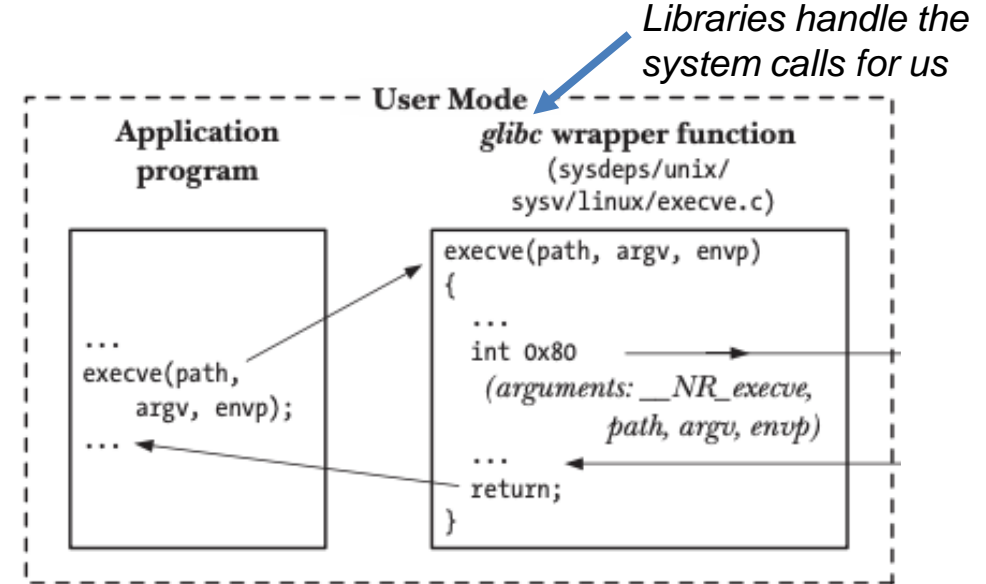


Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0], name, NULL);
    return 0;
} syscall
```

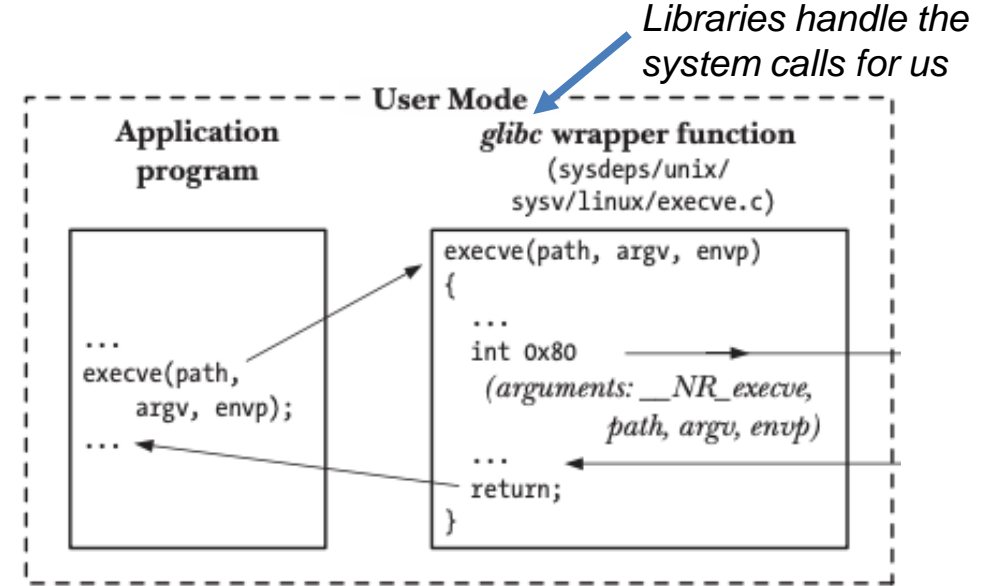


Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0], name, NULL);
    return 0;
} syscall
```



EAX

EBX

ECX

EDX

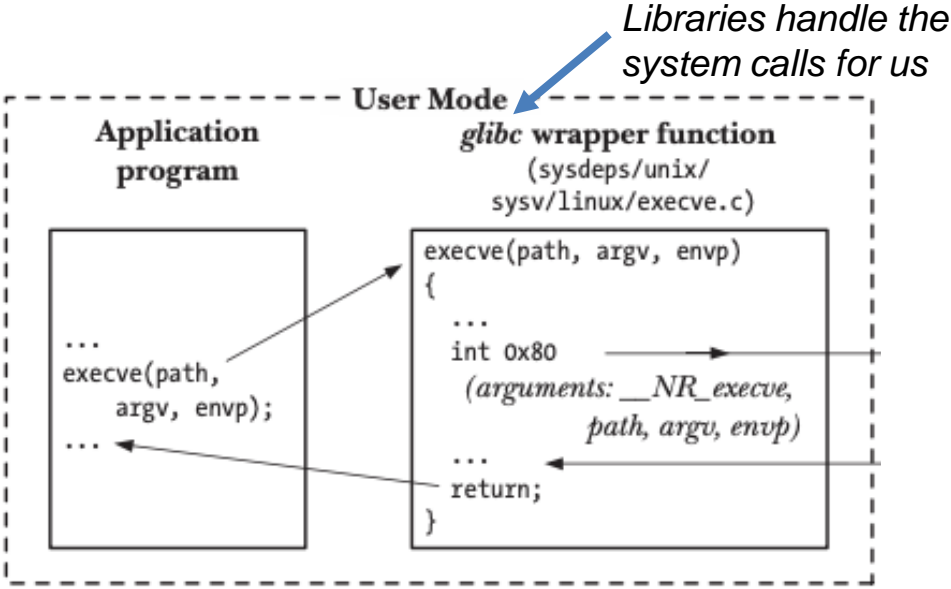
The operating system have hundreds of different syscalls, and different syscalls have different parameters, we need a way to distinguish them

Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0], name, NULL);
    return 0;
} syscall
```



EAX	System Call Number	
EBX	Address of “/bin/bc”	
ECX	0 or 1	Environment variables
EDX	INT 0x80	send trap to kernel and invoke the syscall

The operating system have hundreds of different syscalls, and different syscalls have different parameters, we need a way to distinguish them

The OS will look at the values at certain registers!

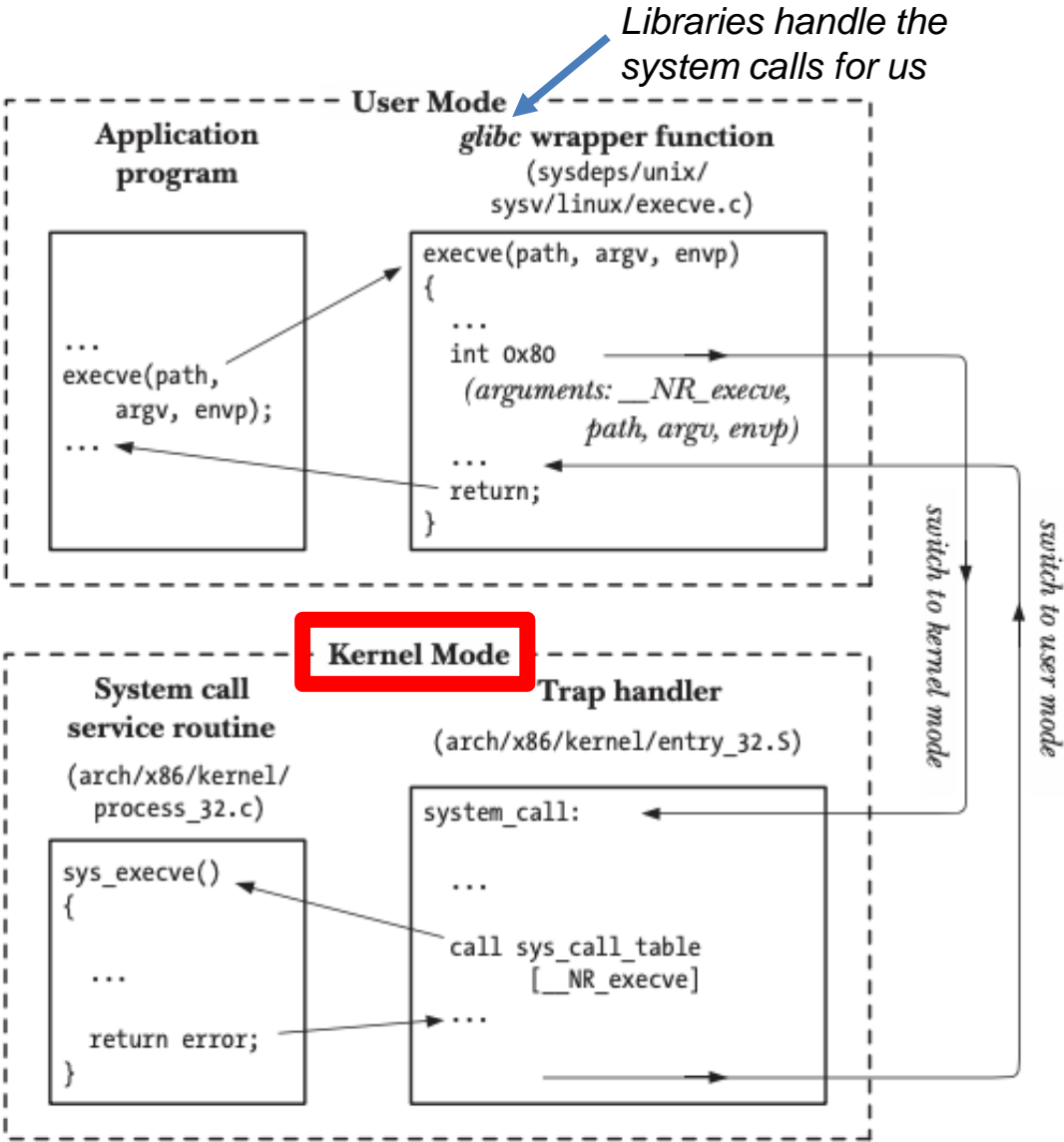
Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

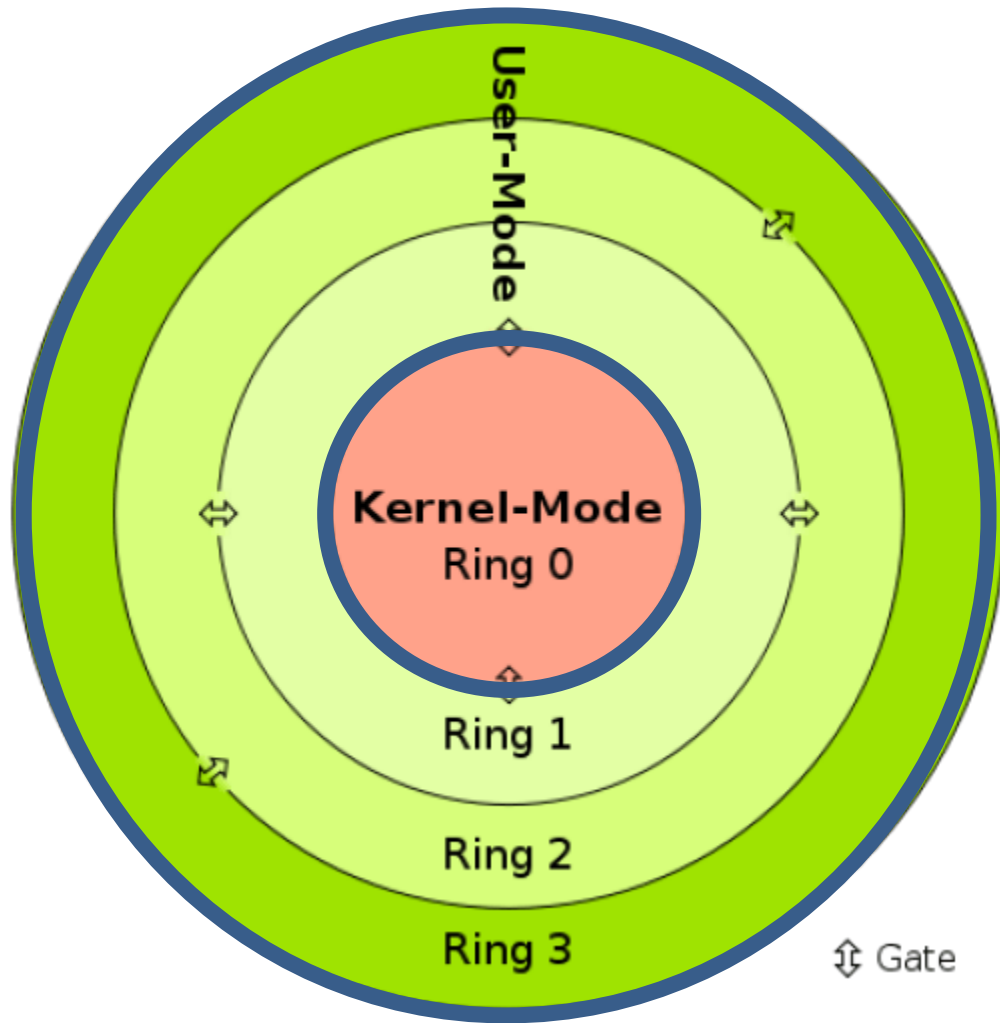
```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0], name, NULL);
    return 0;
} syscall
```

- EAX** System Call Number
- EBX** Address of “/bin/bc”
- ECX** 0 or 1 Environment variables
- EDX** INT 0x80 send trap to kernel and invoke the syscall



Syscalls

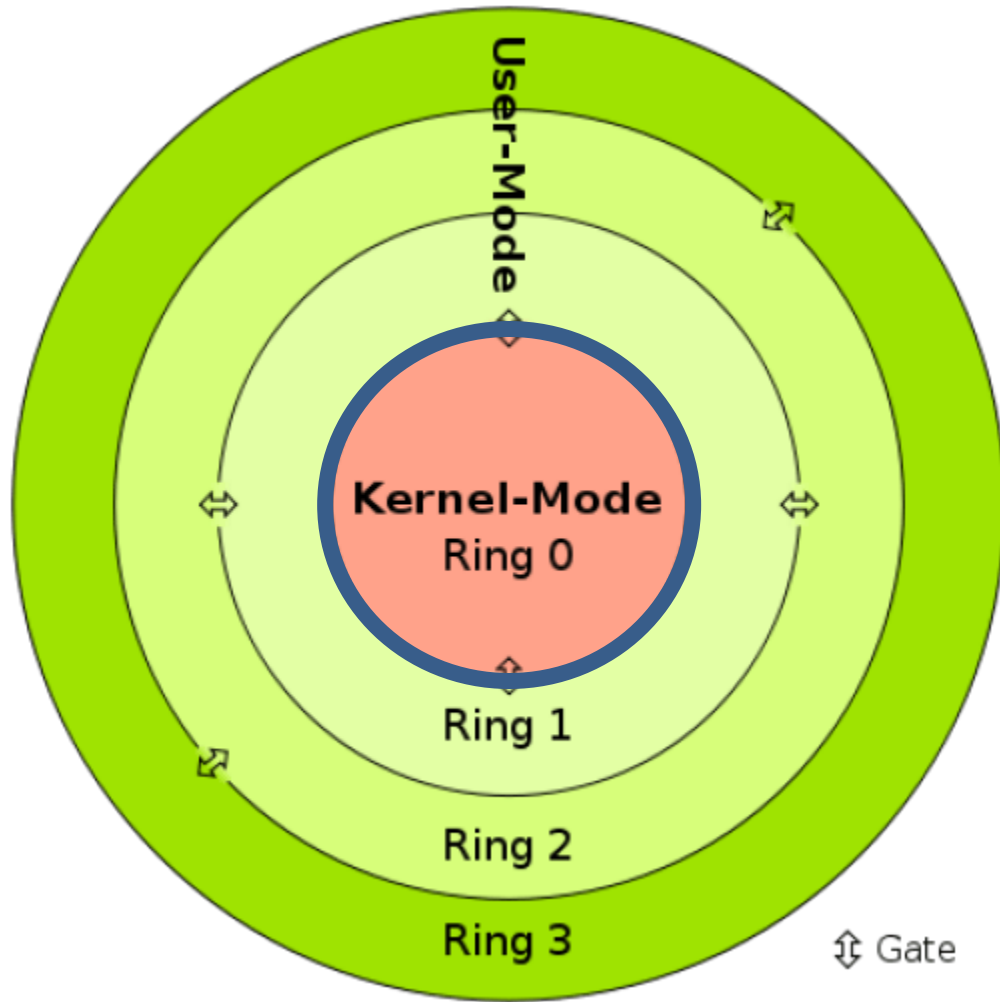


All applications run in user mode.

The code has no ability to directly access hardware

Code running in user mode must use
API/syscalls to access hardware and memory

Syscalls



All applications run in user mode.

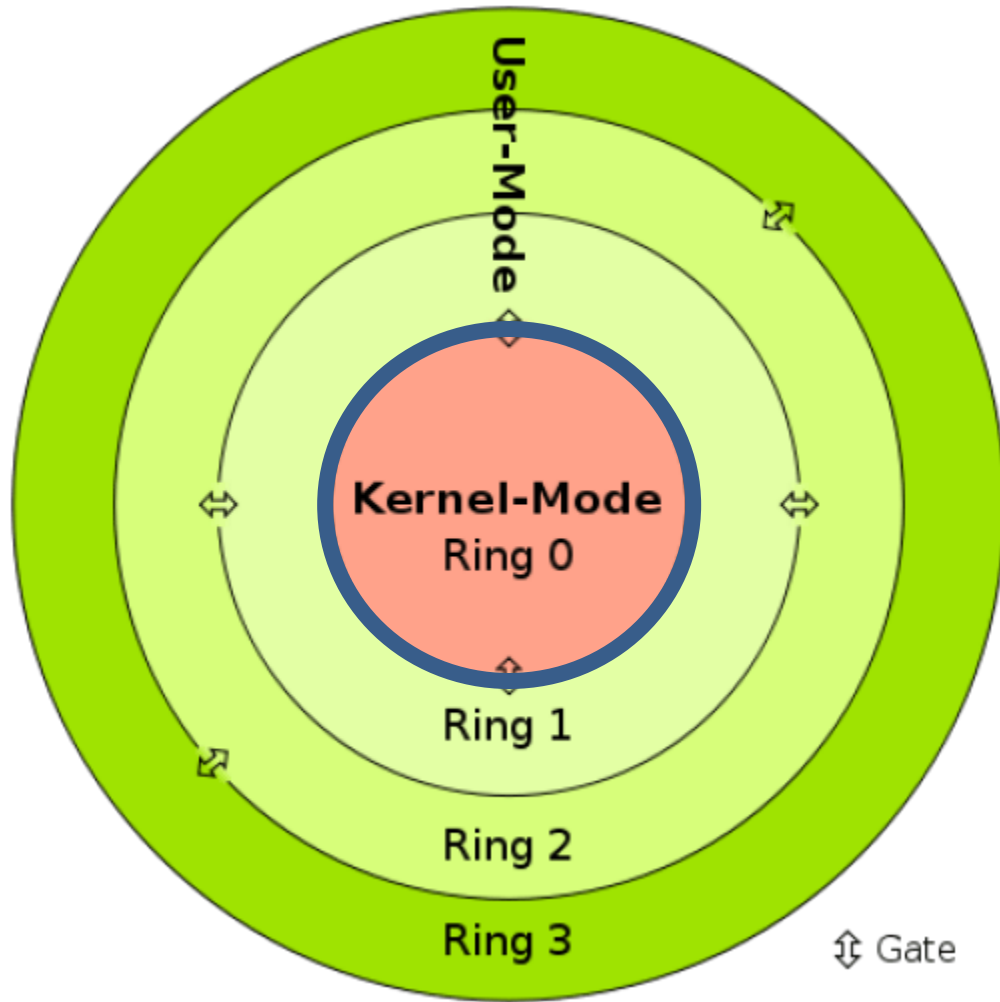
The code has no ability to directly access hardware

Code running in user mode must use
API/syscalls to access hardware and memory

Code running in kernel-mode has complete, unrestricted access to computer resources

Reserved for the lowest-level trusted functions
of the operating system

Syscalls



The collective functionality and services of the OS that manages the computer and its resources is called the **kernel**

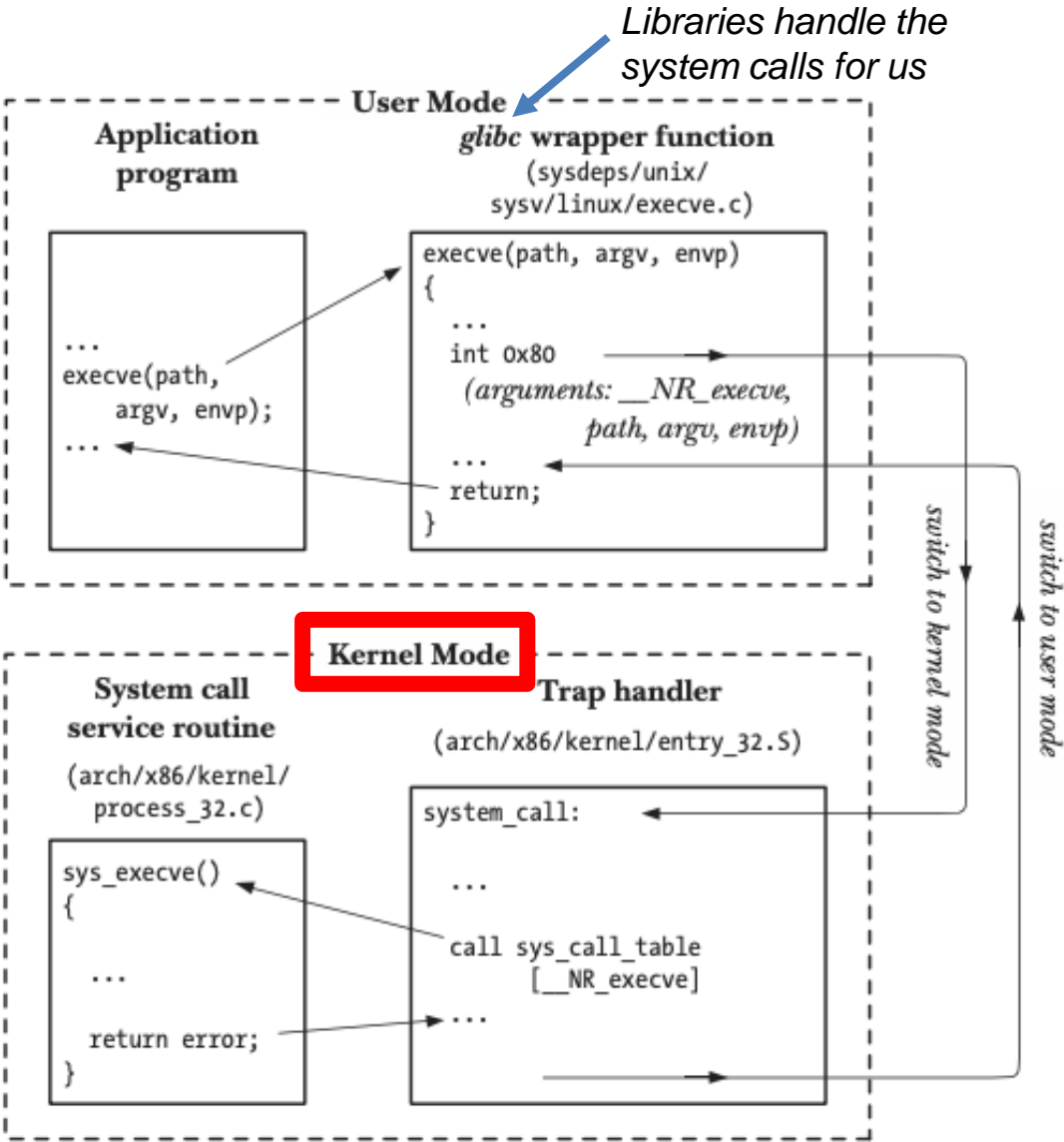
Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0],name, NULL);
    return 0;
} syscall
```

- EAX** System Call Number
- EBX** Address of “/bin/bc”
- ECX** 0 or 1 Environment variables
- EDX** INT 0x80 send trap to kernel and invoke the syscall



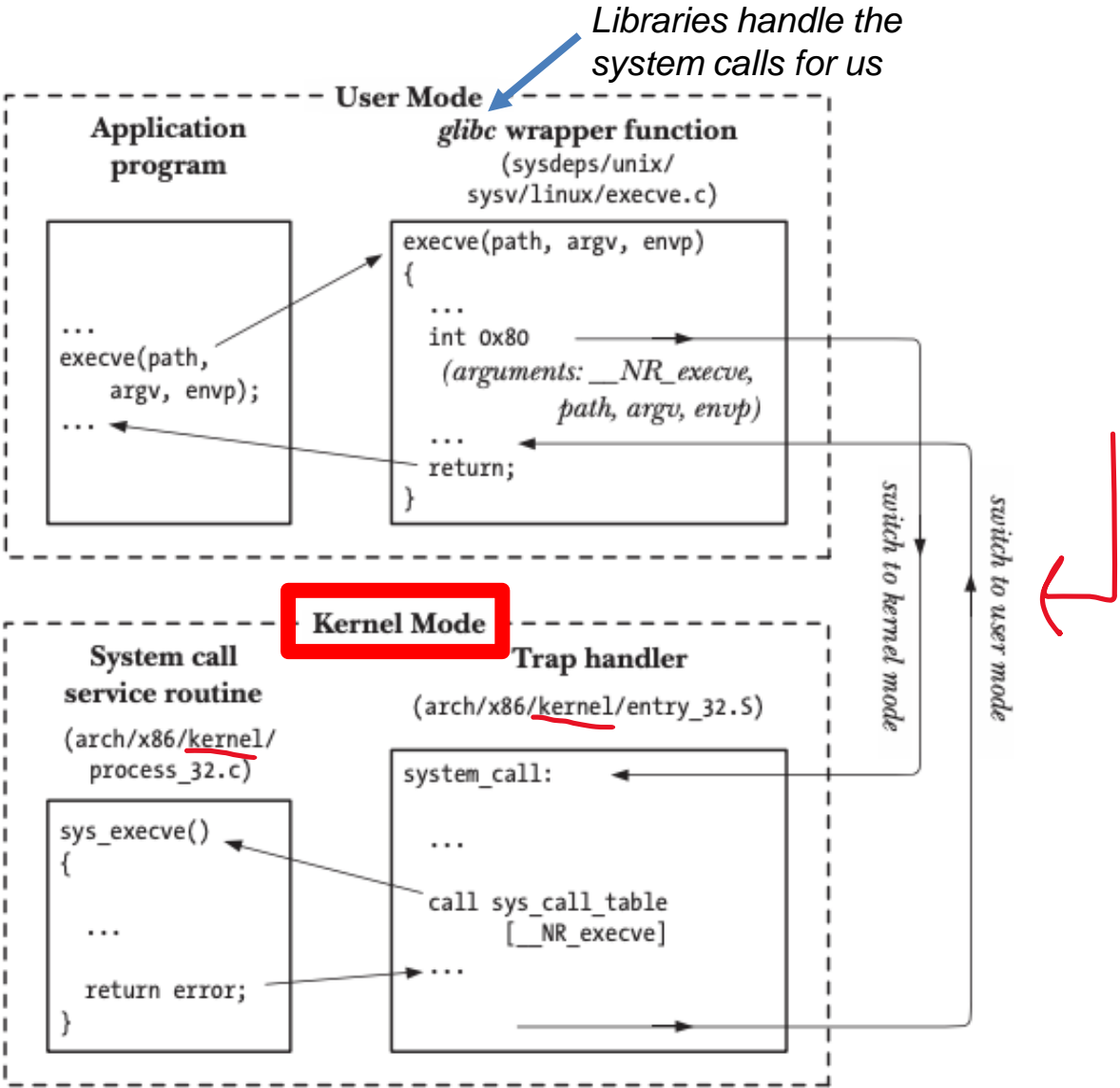
Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

int main()
{
    char *name[2];
    name[0] = "/bin/bc";
    name[1] = NULL;
    execve(name[0], name, NULL);
    return 0;
} syscall
```

- EAX** System Call Number
- EBX** Address of “/bin/bc”
- ECX** 0 or 1 Environment variables
- EDX** INT 0x80 send trap to kernel and invoke the syscall



Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

NR	syscall name	references	%eax	arg0 (%ebx)	arg1 (%ecx)	arg2 (%edx)	arg3 (%esi)	arg4 (%edi)	arg5 (%ebp)
0	restart_syscall	man/ cs/	0x00	-	-	-	-	-	-
1	exit	man/ cs/	0x01	int error_code	-	-	-	-	-
2	fork	man/ cs/	0x02	-	-	-	-	-	-
3	read	man/ cs/	0x03	unsigned int fd	char *buf	size_t count	-	-	-
4	write	man/ cs/	0x04	unsigned int fd	const char *buf	size_t count	-	-	-
5	open	man/ cs/	0x05	const char *filename	int flags	umode_t mode	-	-	-
6	close	man/ cs/	0x06	unsigned int fd	-	-	-	-	-
7	waitpid	man/ cs/	0x07	pid_t pid	int *stat_addr	int options	-	-	-
8	creat	man/ cs/	0x08	const char *pathname	umode_t mode	-	-	-	-
9	link	man/ cs/	0x09	const char *oldname	const char *newname	-	-	-	-
10	unlink	man/ cs/	0x0a	const char *pathname	-	-	-	-	-
11	execve	man/ cs/	0x0b	const char *filename	const char *const *argv	const char *const *envp	-	-	-
12	chdir	man/ cs/	0x0c	const char *filename	-	-	-	-	-

EDX

INT 0x80

send trap to kernel and invoke the syscall

Syscalls

Applications evoke operating system defined functions, or **system calls (syscalls)**, to access computing resources

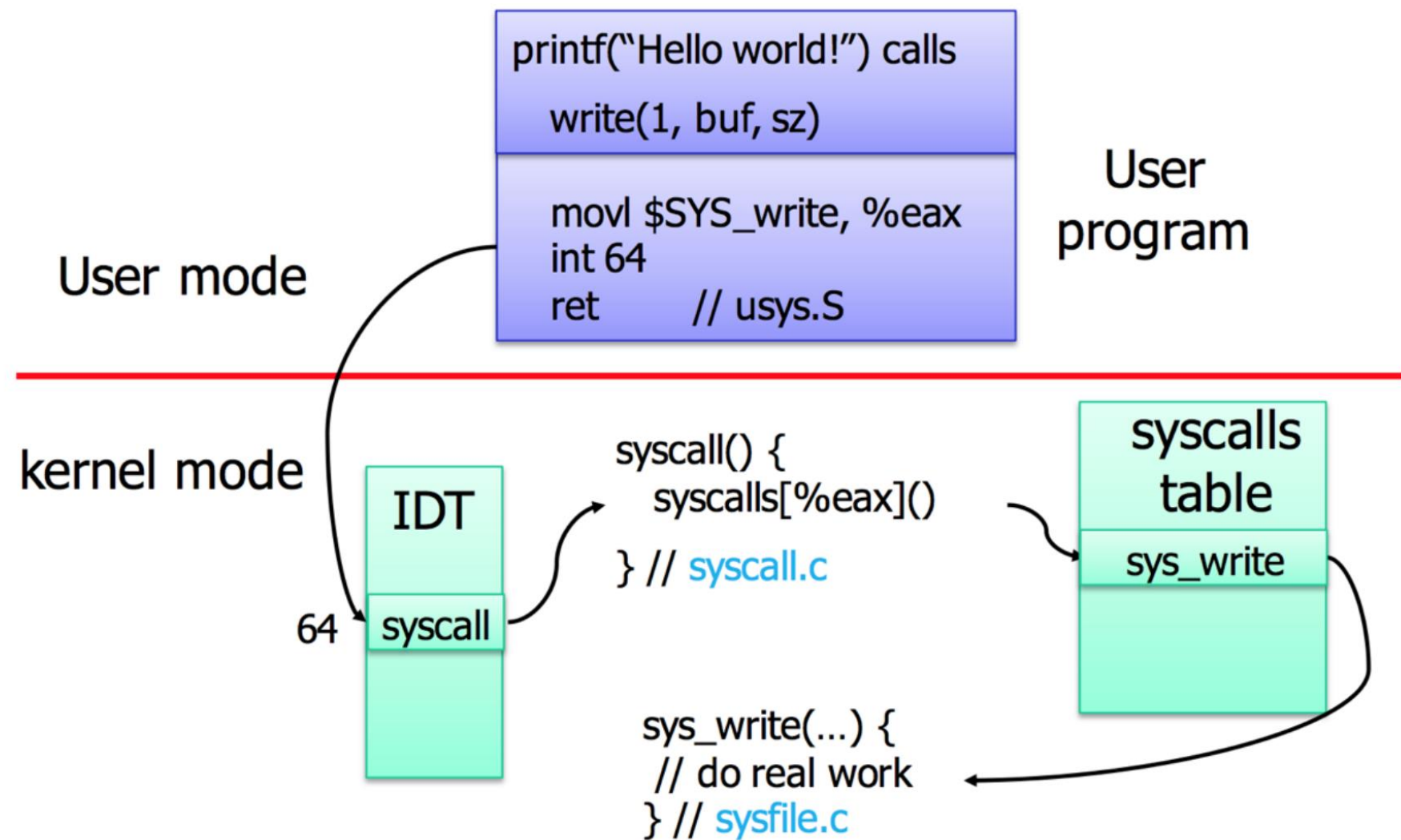
NR	syscall name	references	%eax	arg0 (%ebx)	arg1 (%ecx)	arg2 (%edx)	arg3 (%esi)	arg4 (%edi)	arg5 (%ebp)
0	restart_syscall	man/ cs/	0x00	-	-	-	-	-	-
1	exit	man/ cs/	0x01	int error_code	-	-	-	-	-
2	fork	man/ cs/	0x02	-	-	-	-	-	-
3	read	man/ cs/	0x03	unsigned int fd	char *buf	size_t count	-	-	-
4	https://chromium.googlesource.com/chromiumos/docs/+/master/constants/syscalls.md#x86-32_bit								
5									
6									
7									
8									
9									
10	unlink	man/ cs/	0x0a	const char *pathname	-	-	-	-	-
11	execve	man/ cs/	0x0b	const char *filename	const char *const *argv	const char *const *envp	-	-	-
12	chdir	man/ cs/	0x0c	const char *filename	-	-	-	-	-

EDX

INT 0x80

send trap to kernel and invoke the syscall

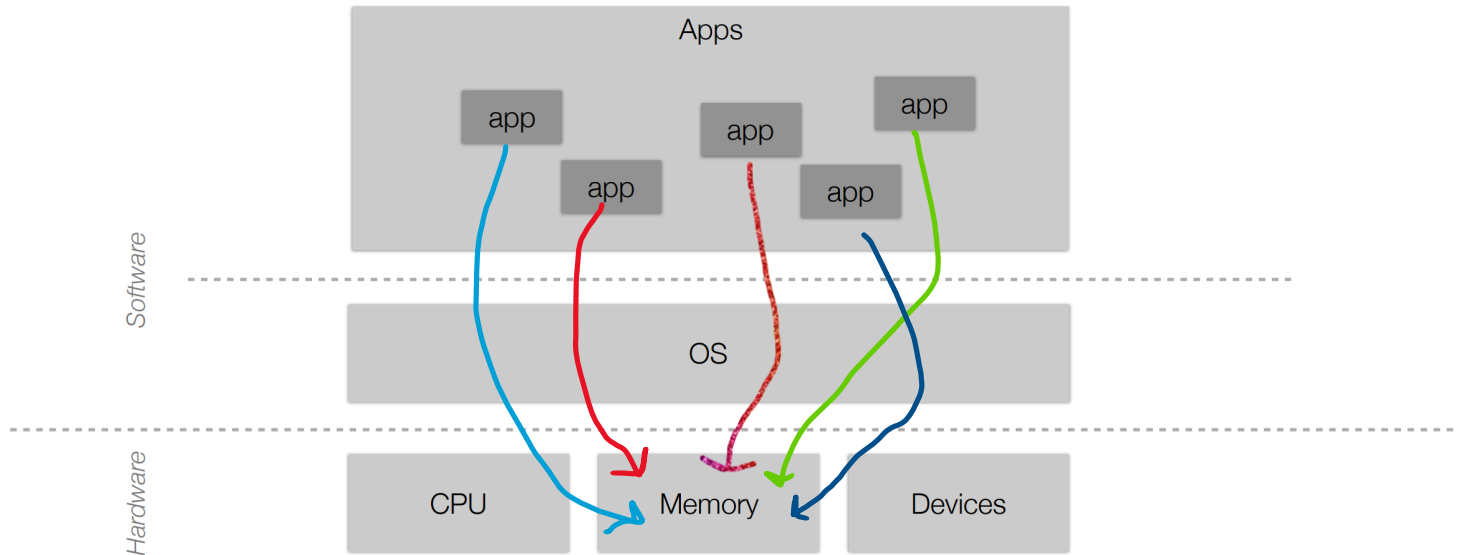
Syscalls



Applications Layout in Memory

Process Manager

- Manages how processes are structured and how to handle many processes running at once

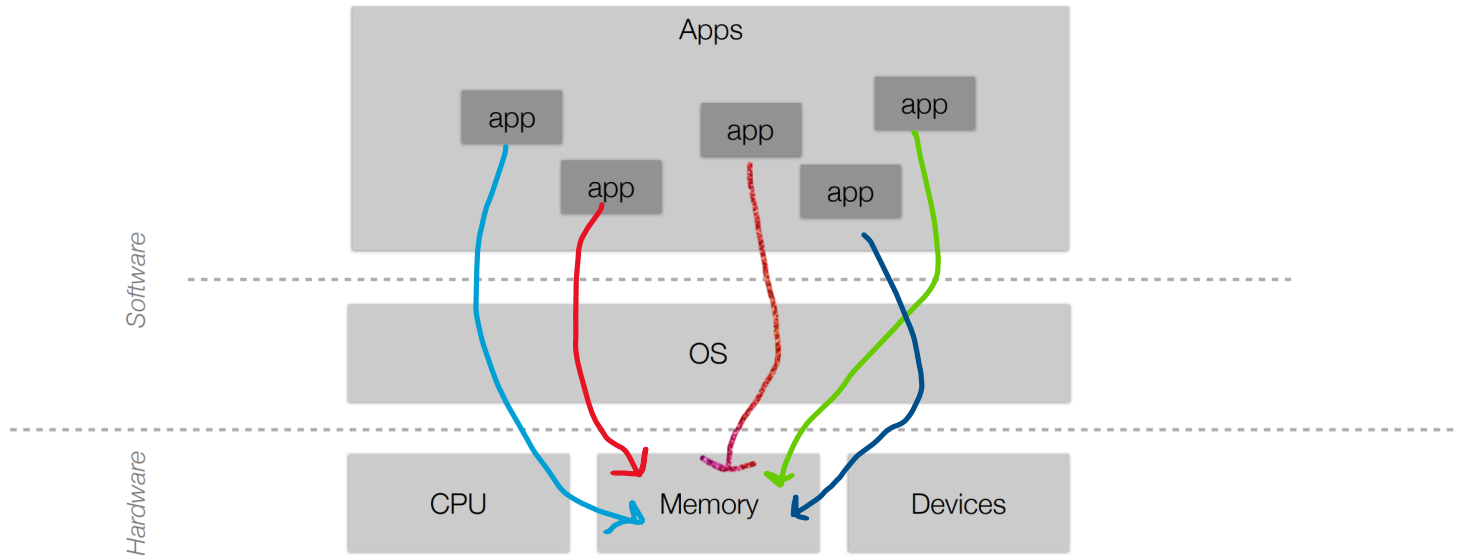


How does a **program** get loaded into memory?

Applications Layout in Memory

Process Manager

- Manages how processes are structured and how to handle many processes running at once

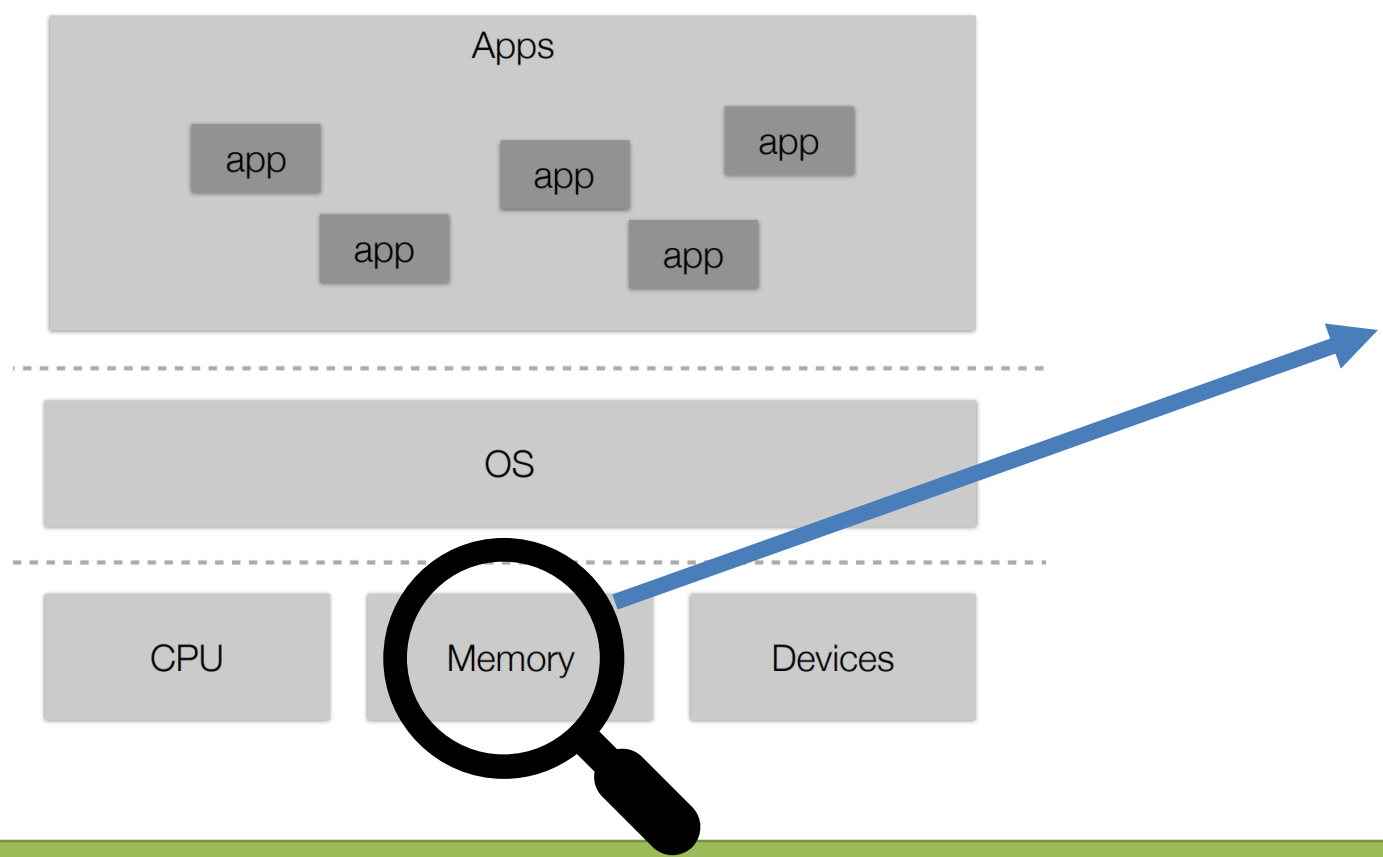


How does a **program** get loaded into memory?

An active program running on a computer is called a **process**

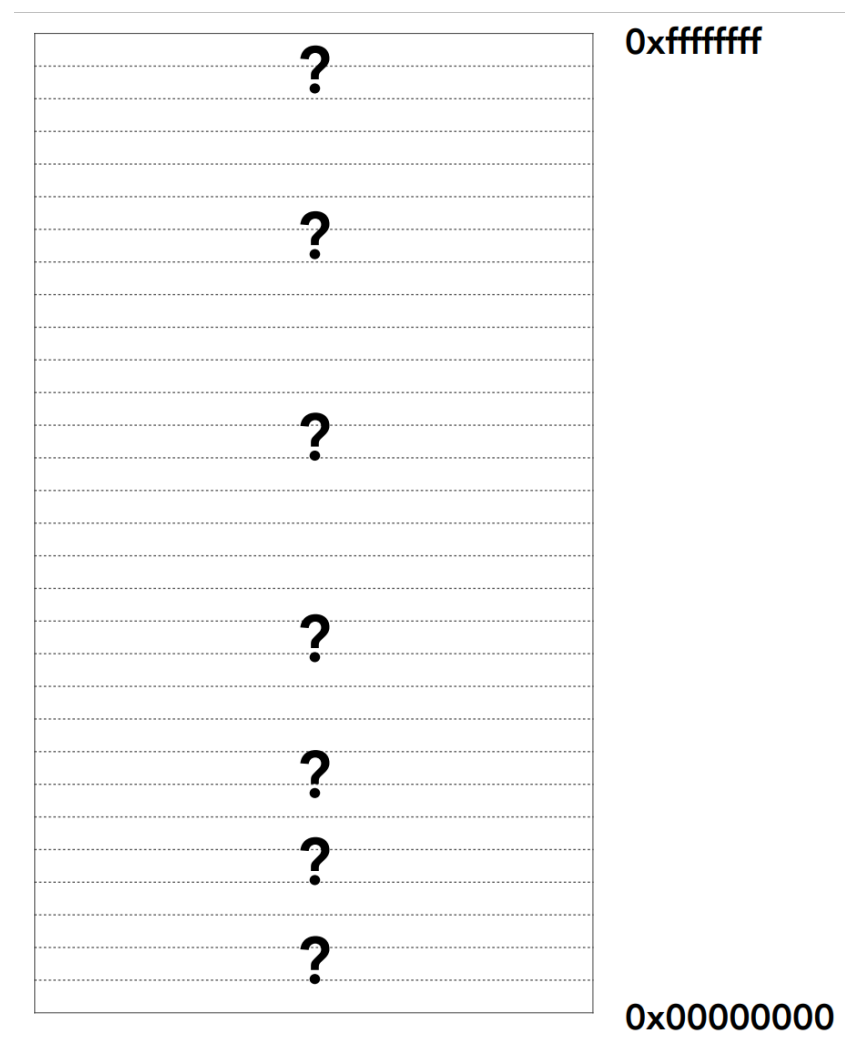
Applications Layout in Memory

What does a program look like in memory?



Process Manager

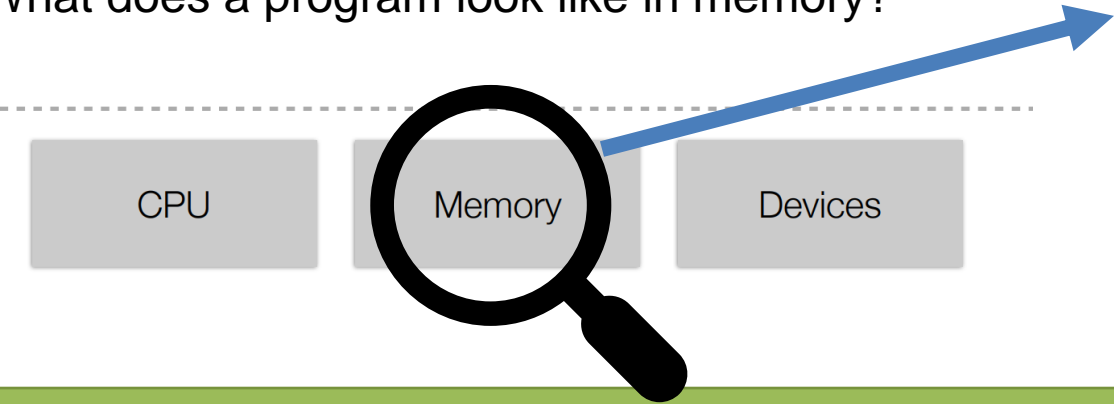
- Manages how processes are structured and how to handle many processes running at once



Applications Layout in Memory

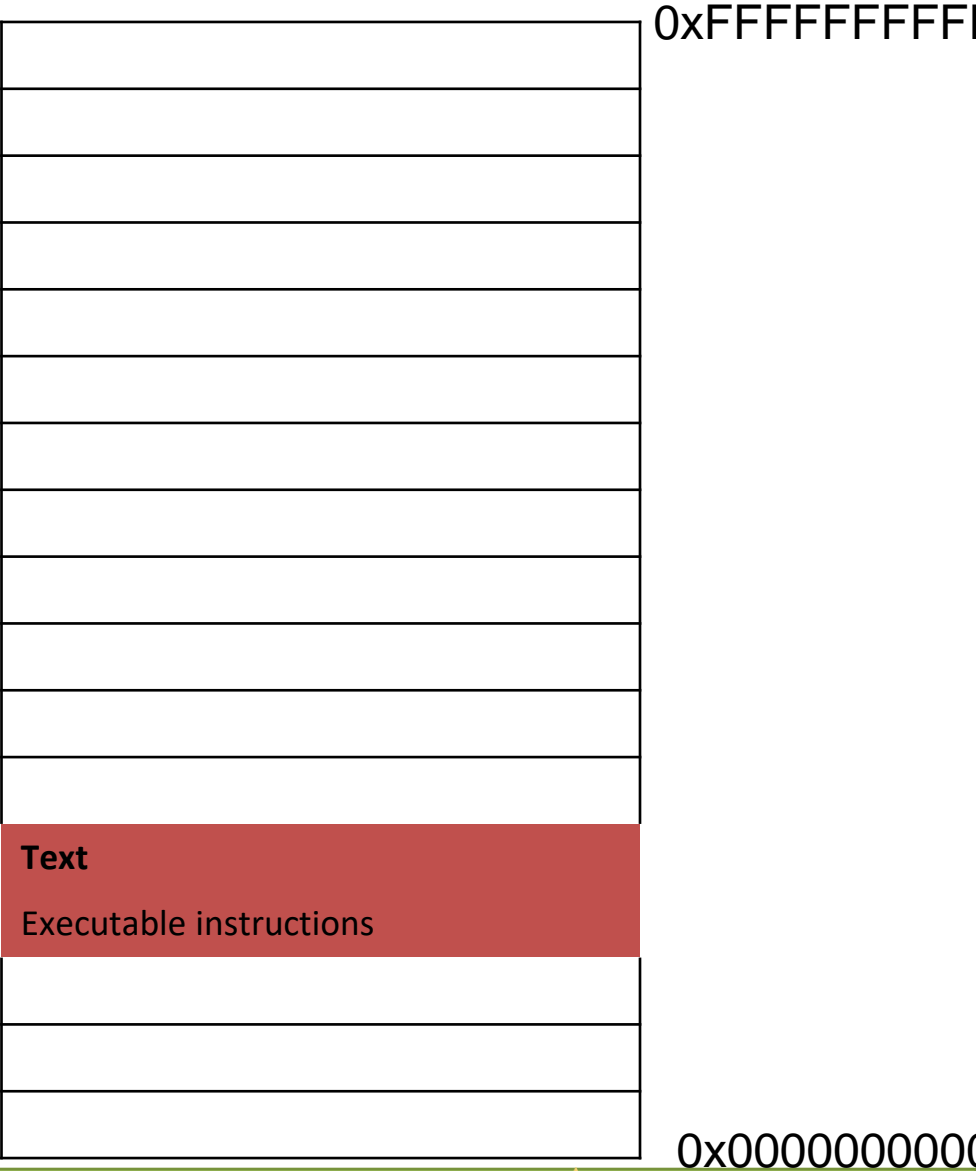
Text Segment- binary executable instructions for the process

What does a program look like in memory?



Process Manager

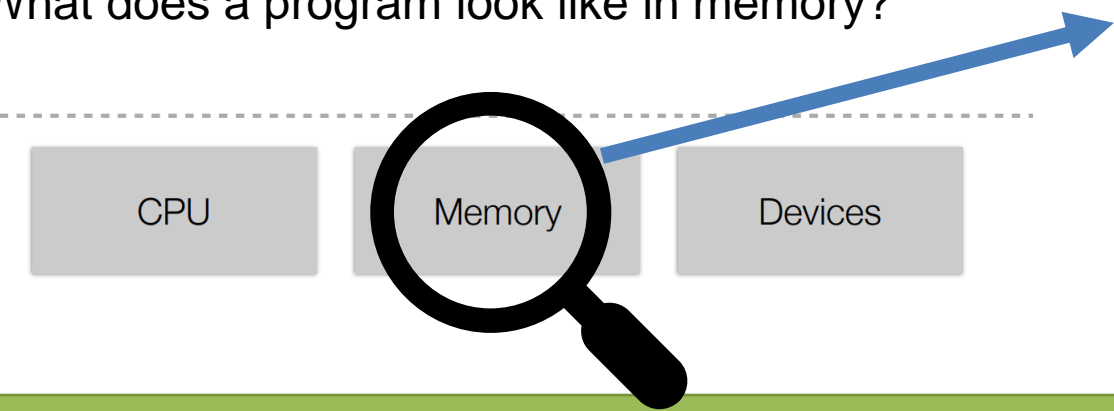
- Manages how processes are structured and how to handle many processes running at once



Applications Layout in Memory

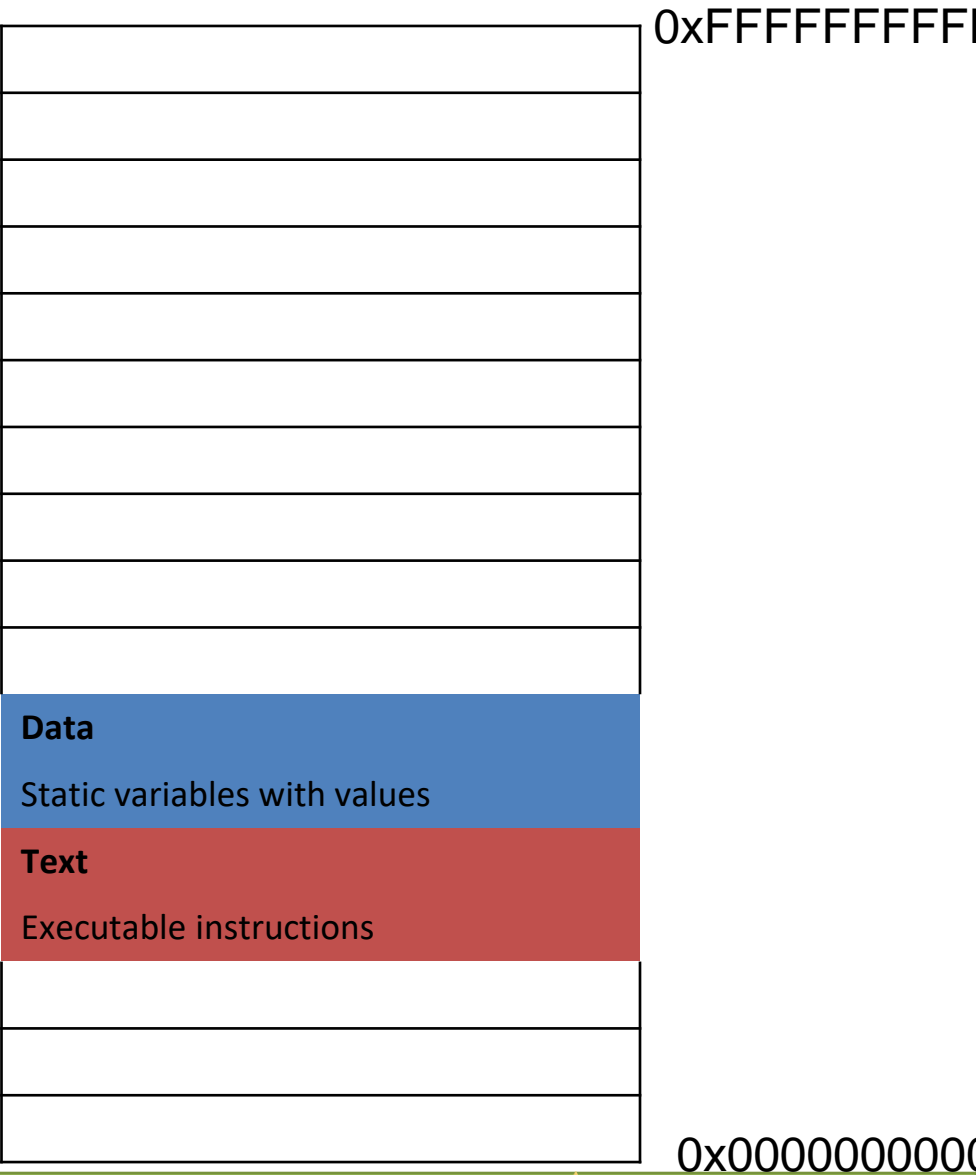
Data Segment- Static variables initialized by the programmer

What does a program look like in memory?



Process Manager

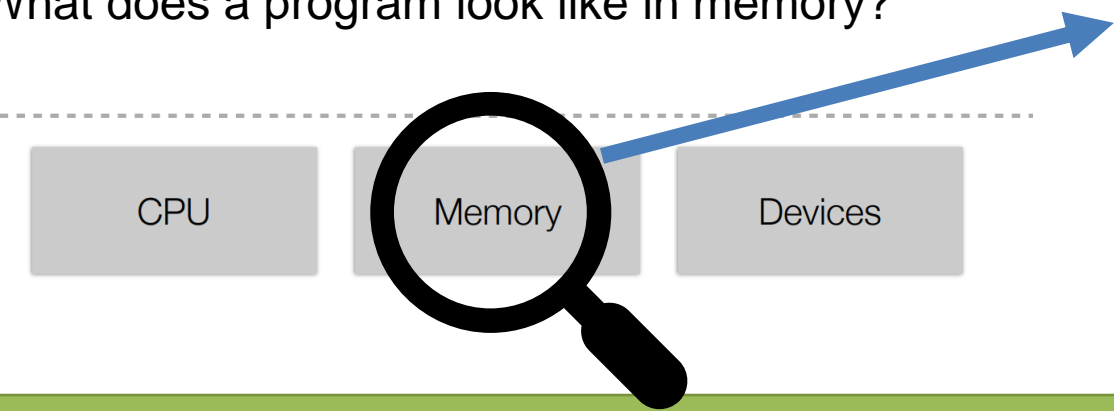
- Manages how processes are structured and how to handle many processes running at once



Applications Layout in Memory

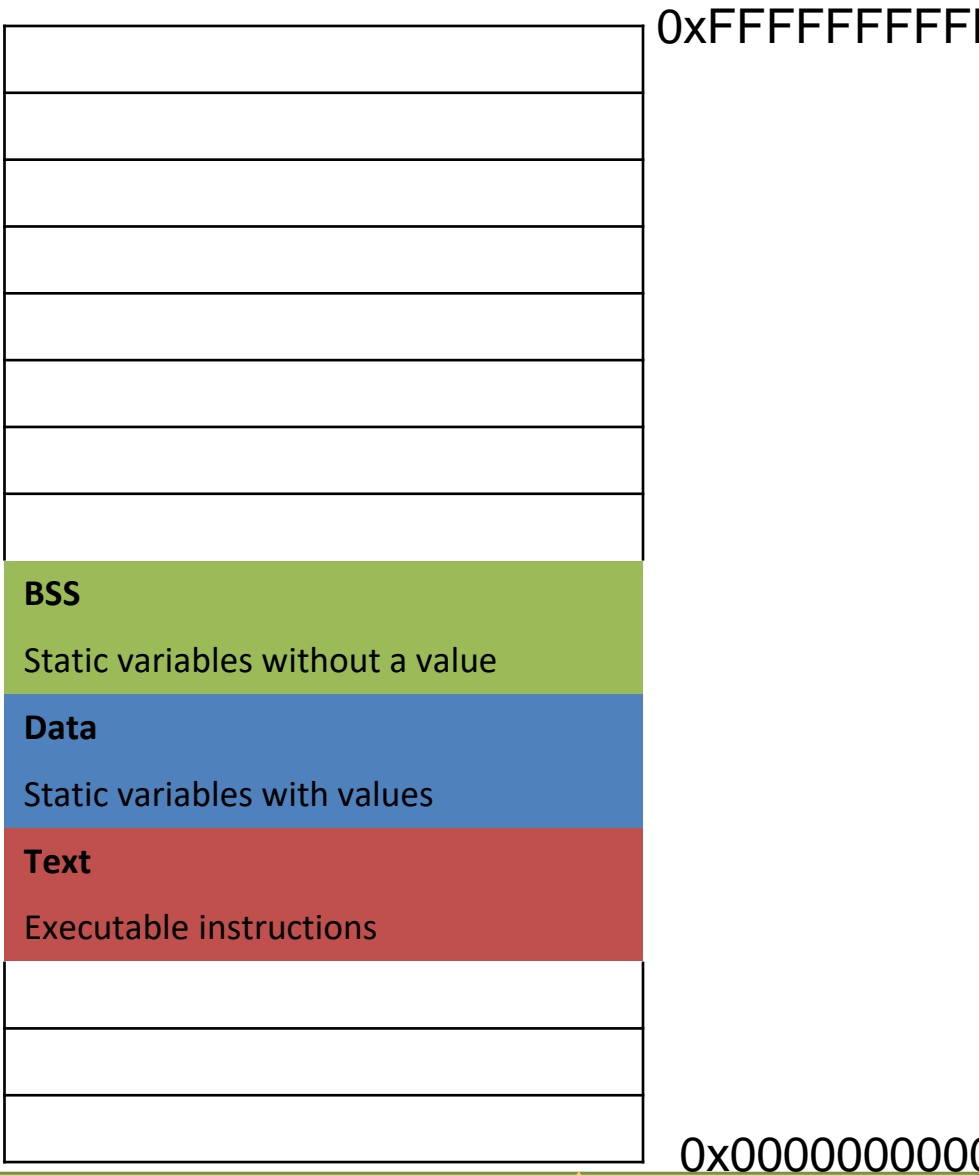
BSS Segment- contains statically allocated variables that are declared, but have not been assigned a value yet

What does a program look like in memory?



Process Manager

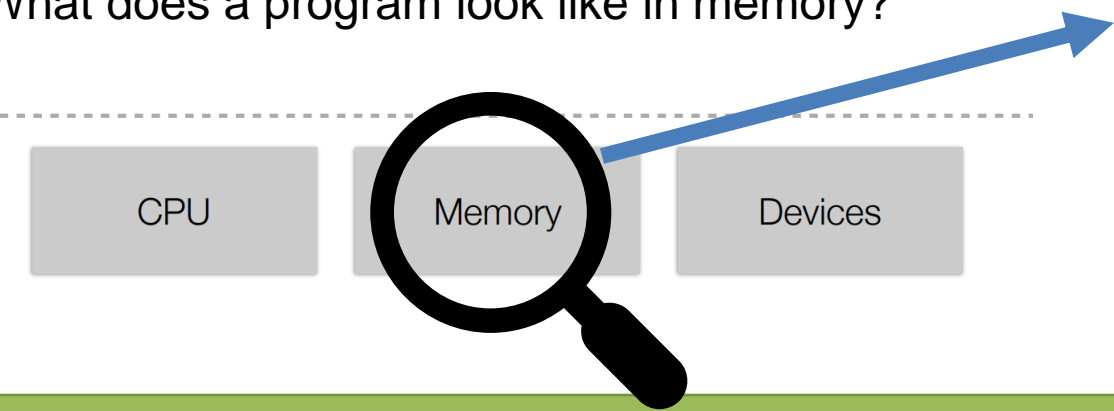
- Manages how processes are structured and how to handle many processes running at once



Applications Layout in Memory

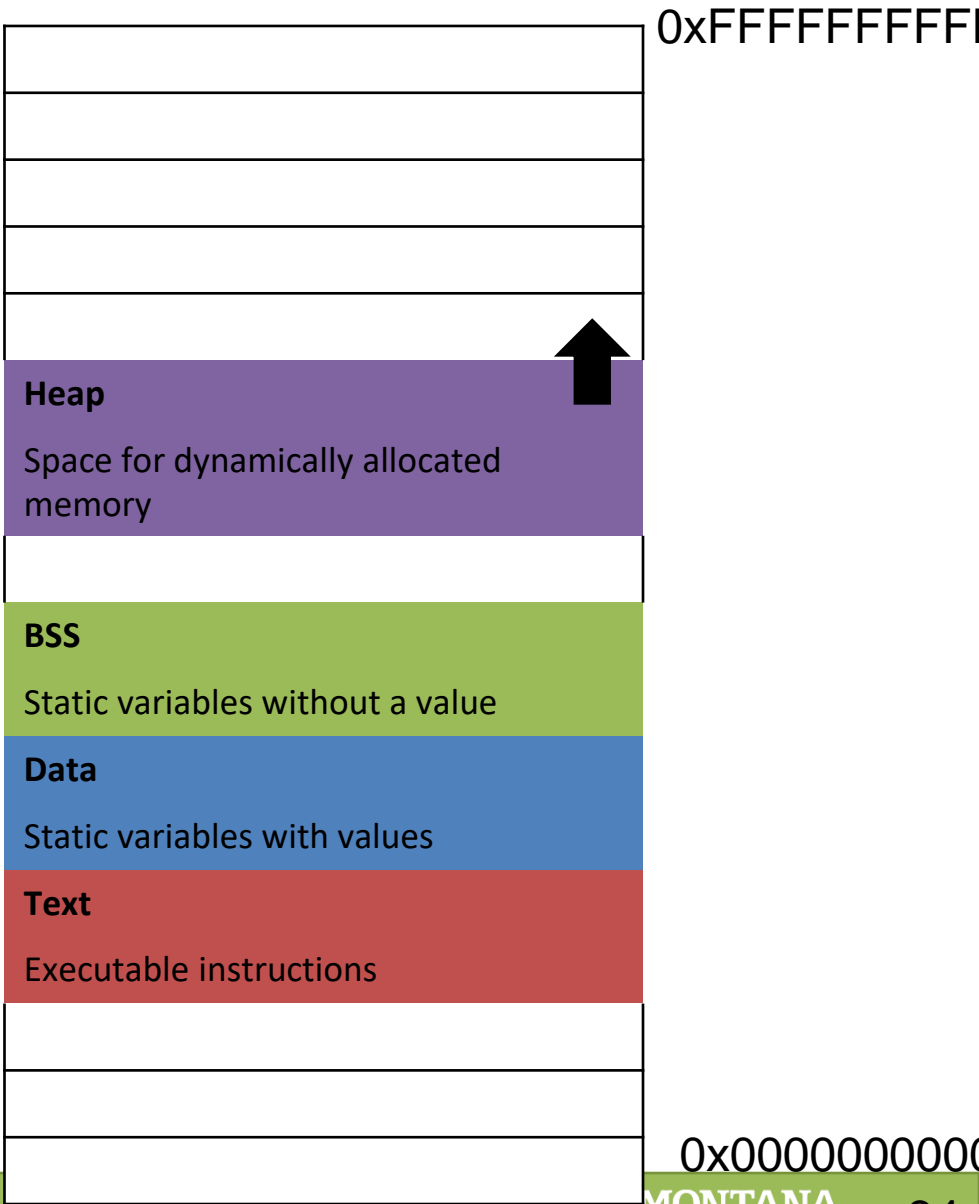
Heap- memory set aside for dynamic allocation (e.g. malloc). Grows “up” as more memory is allocated

What does a program look like in memory?



Process Manager

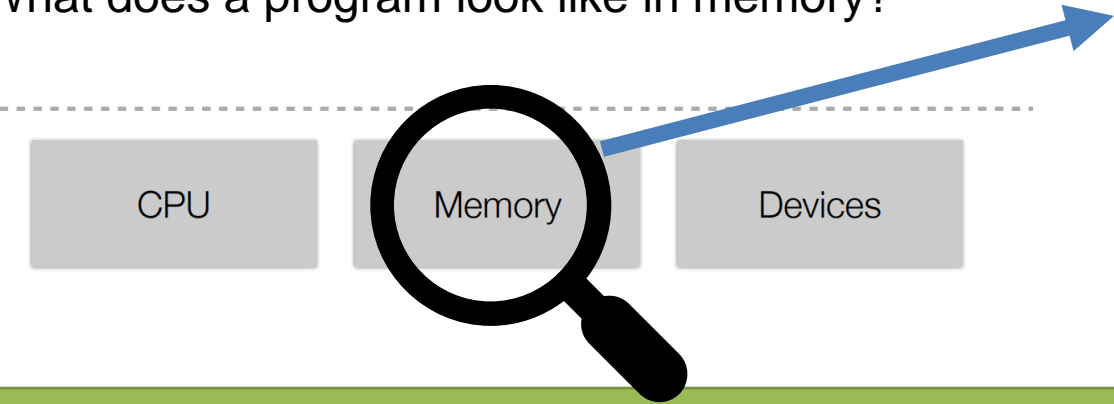
- Manages how processes are structured and how to handle many processes running at once



Applications Layout in Memory

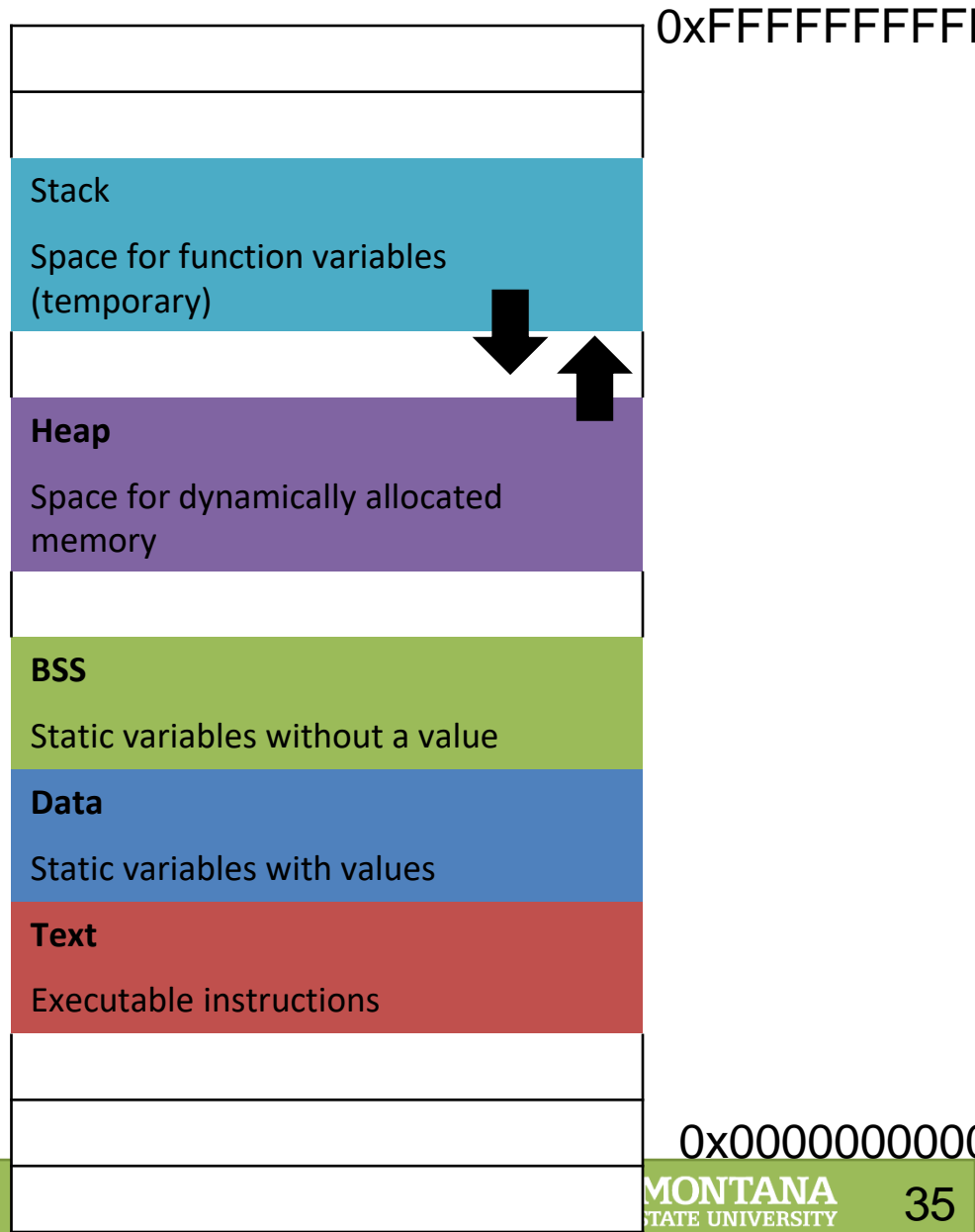
Stack – memory for storing function variables. Grows “down” as additional functions are called

What does a program look like in memory?



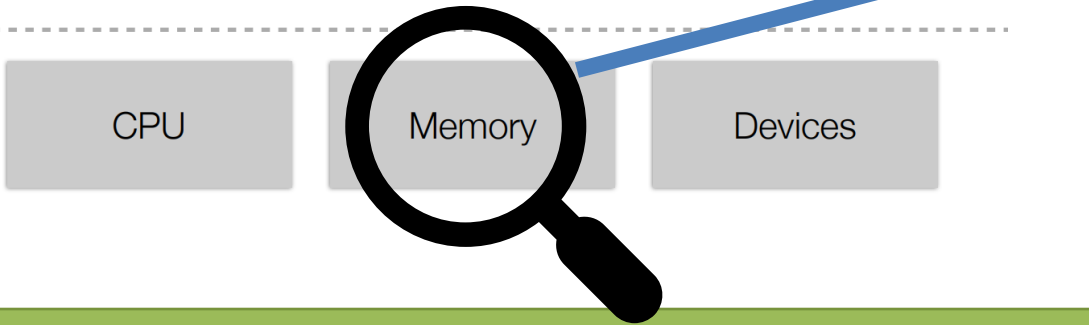
Process Manager

- Manages how processes are structured and how to handle many processes running at once



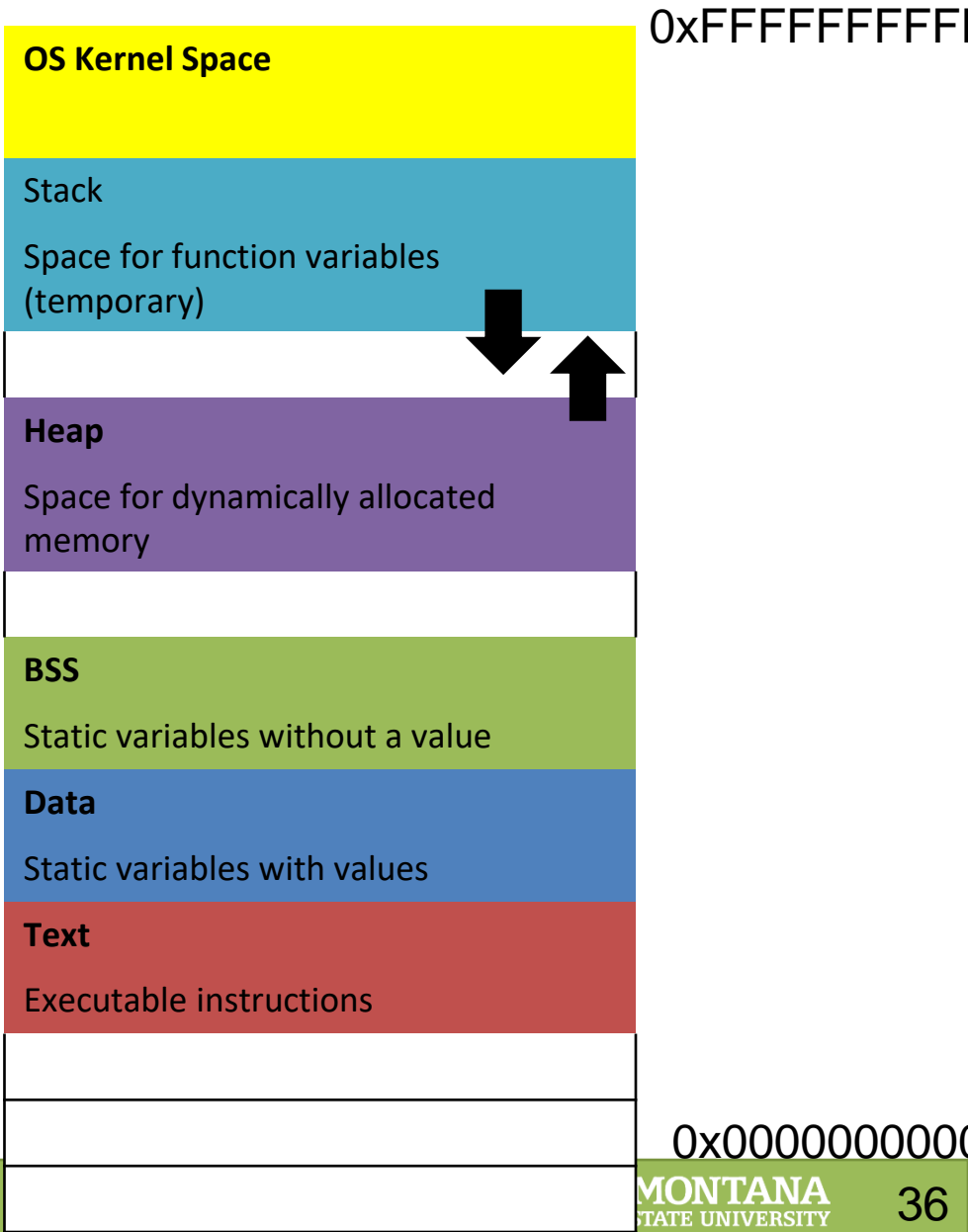
Applications Layout in Memory

What does a program look like in memory?



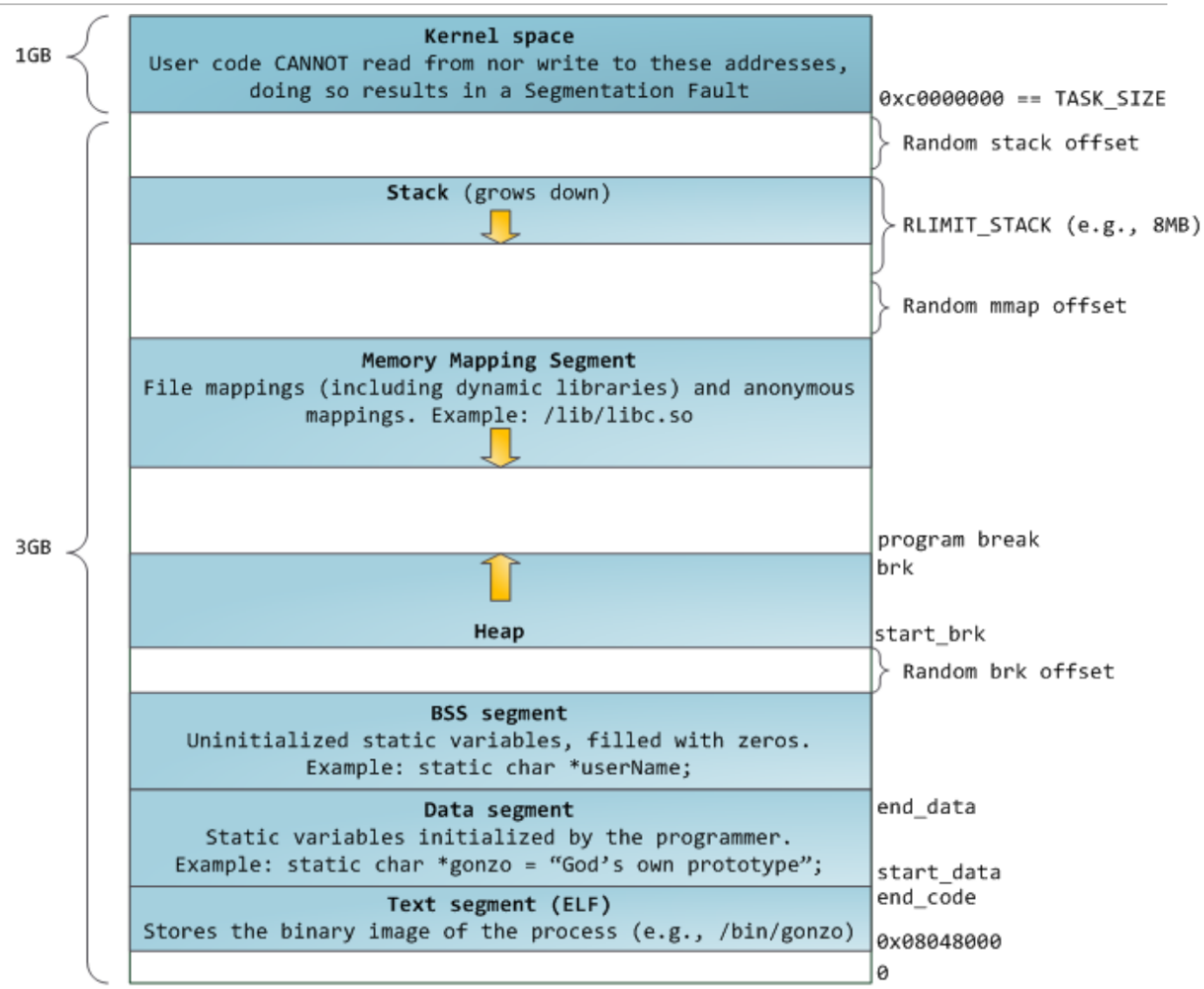
Process Manager

- Manages how processes are structured and how to handle many processes running at once



Applications Layout in Memory

Demo?



Applications Layout in Memory

Output of `pmap` (process mapping tool)

[illegible][illegible]

Applications Layout in Memory

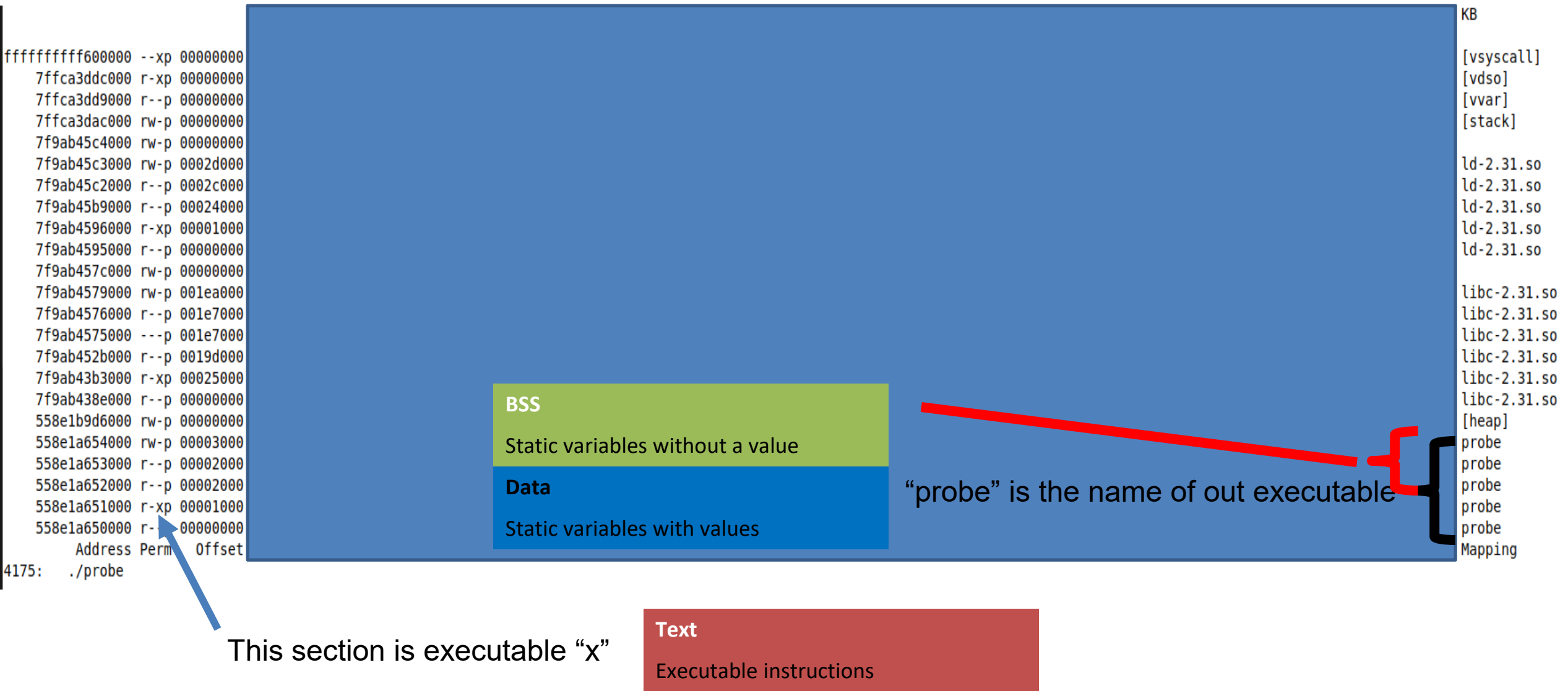
Output of `pmap` (process mapping tool)

			KB
fffffffff600000	--xp 00000000		[vsyscall]
7ffca3ddc000	r-xp 00000000		[vdso]
7ffca3dd9000	r--p 00000000		[vvar]
7ffca3dac000	rw-p 00000000		[stack]
7f9ab45c4000	rw-p 00000000		
7f9ab45c3000	rw-p 0002d000		ld-2.31.so
7f9ab45c2000	r--p 0002c000		ld-2.31.so
7f9ab45b9000	r--p 00024000		ld-2.31.so
7f9ab4596000	r-xp 00001000		ld-2.31.so
7f9ab4595000	r--p 00000000		ld-2.31.so
7f9ab457c000	rw-p 00000000		
7f9ab4579000	rw-p 001ea000		libc-2.31.so
7f9ab4576000	r--p 001e7000		libc-2.31.so
7f9ab4575000	---p 001e7000		libc-2.31.so
7f9ab452b000	r--p 0019d000		libc-2.31.so
7f9ab43b3000	r-xp 00025000		libc-2.31.so
7f9ab438e000	r--p 00000000		libc-2.31.so
558e1b9d6000	rw-p 00000000		[heap]
558e1a654000	rw-p 00003000		probe
558e1a653000	r--p 00002000		probe
558e1a652000	r--p 00002000		probe
558e1a651000	r-xp 00001000		probe
558e1a650000	r--p 00000000		probe
	Address Perm Offset		Mapping
4175:	./probe		

“probe” is the name of out executable

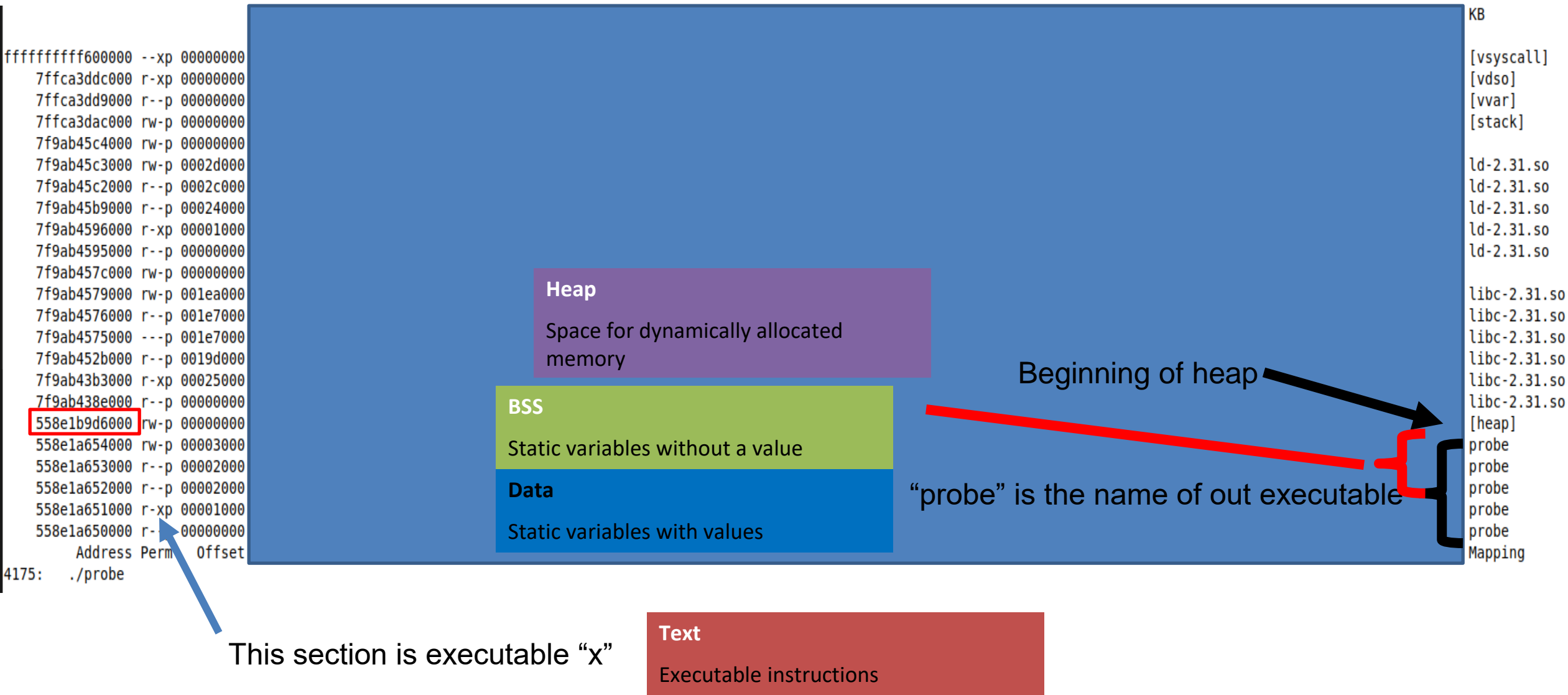
Applications Layout in Memory

Output of `pmap` (process mapping tool)



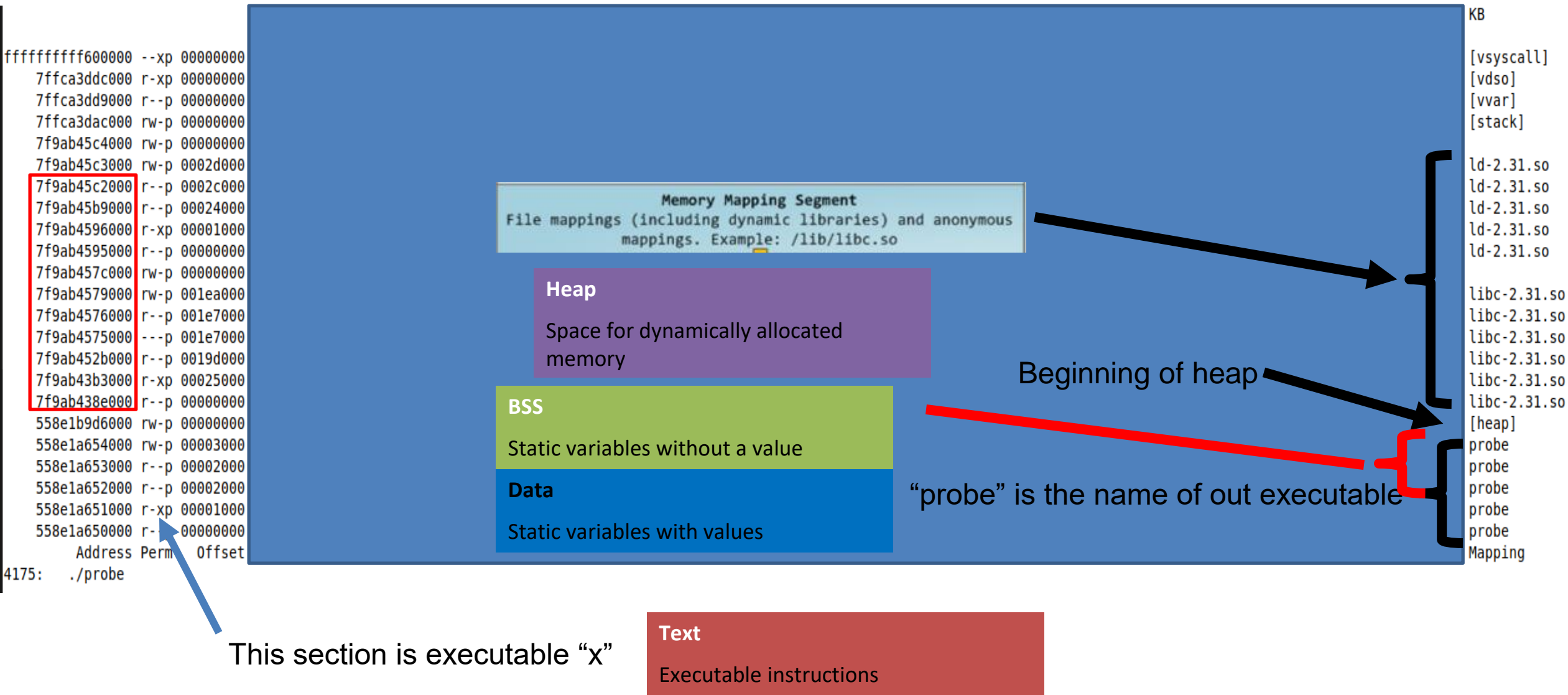
Applications Layout in Memory

Output of `pmap` (process mapping tool)



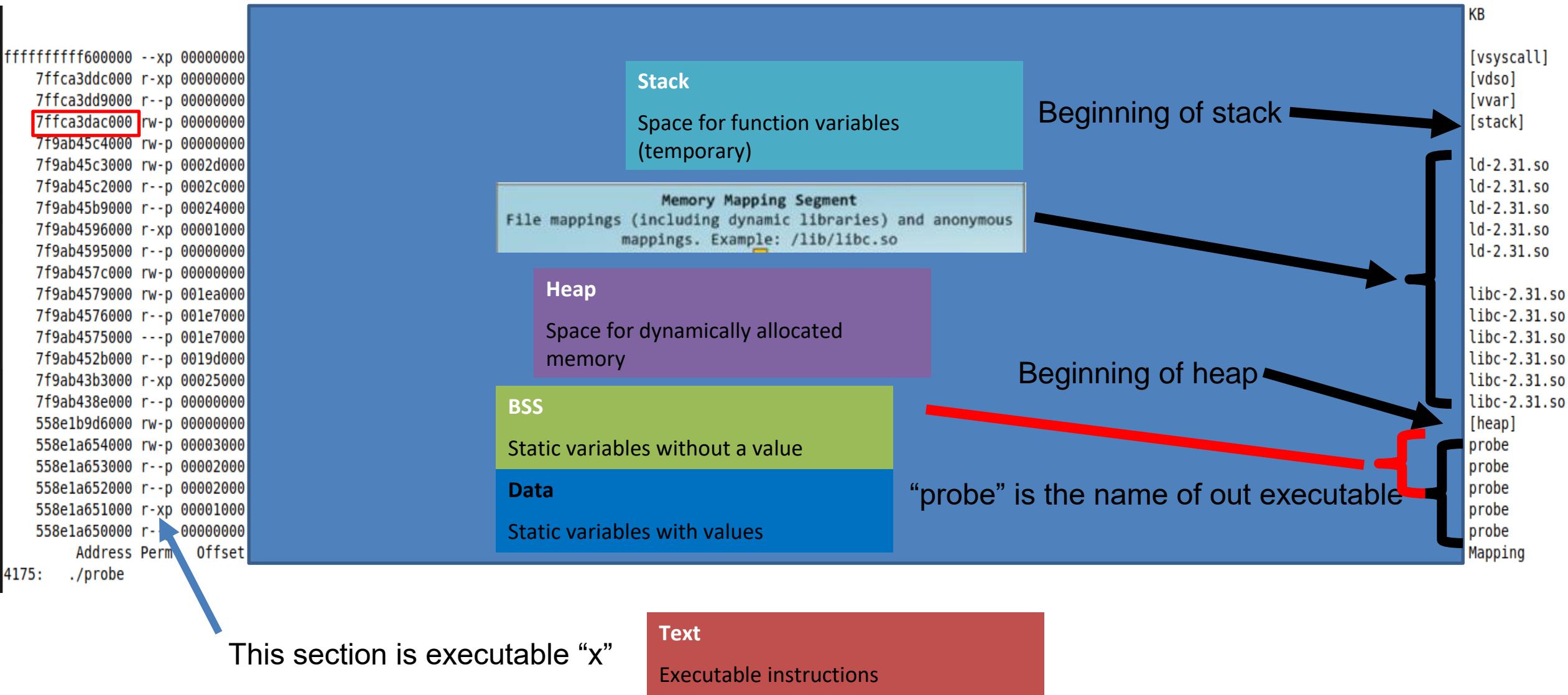
Applications Layout in Memory

Output of `pmap` (process mapping tool)



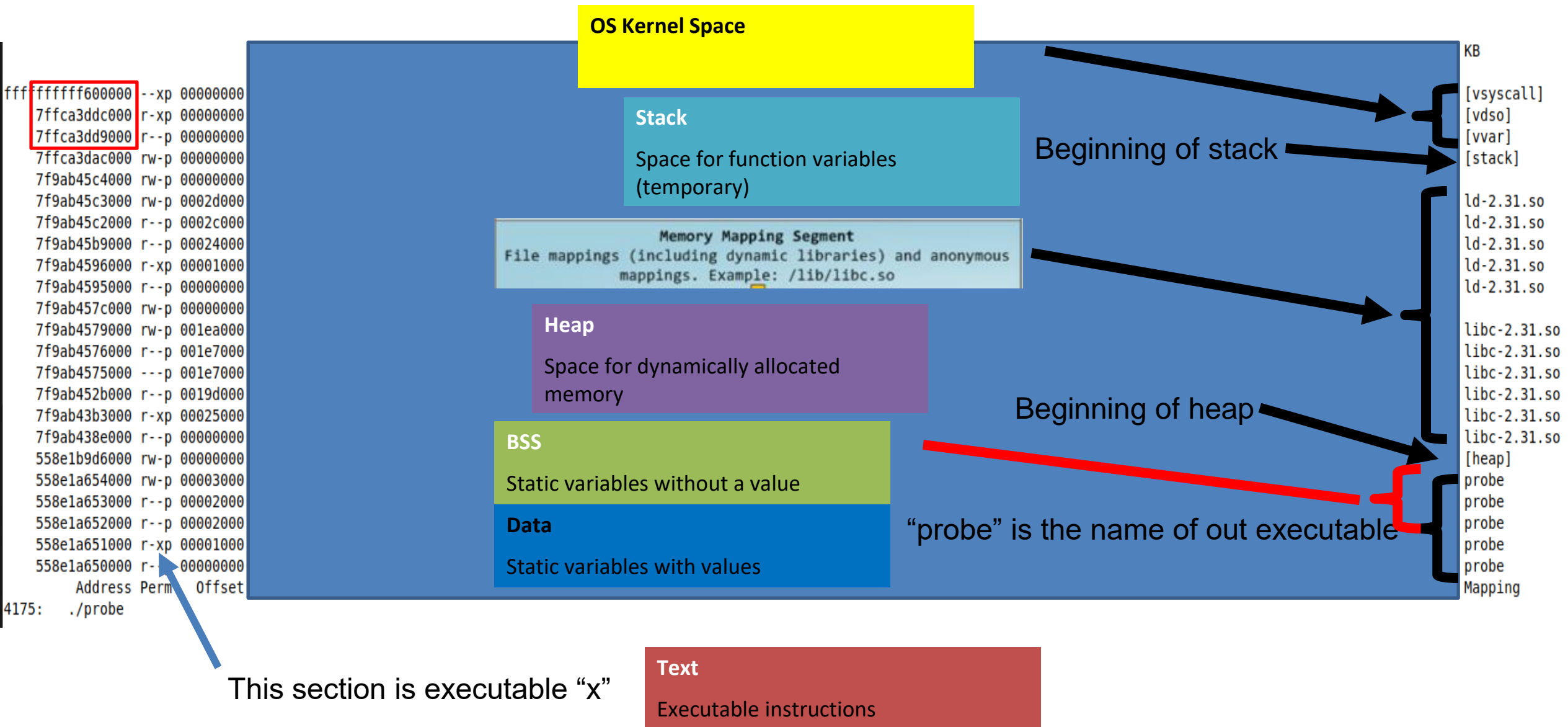
Applications Layout in Memory

Output of `pmap` (process mapping tool)



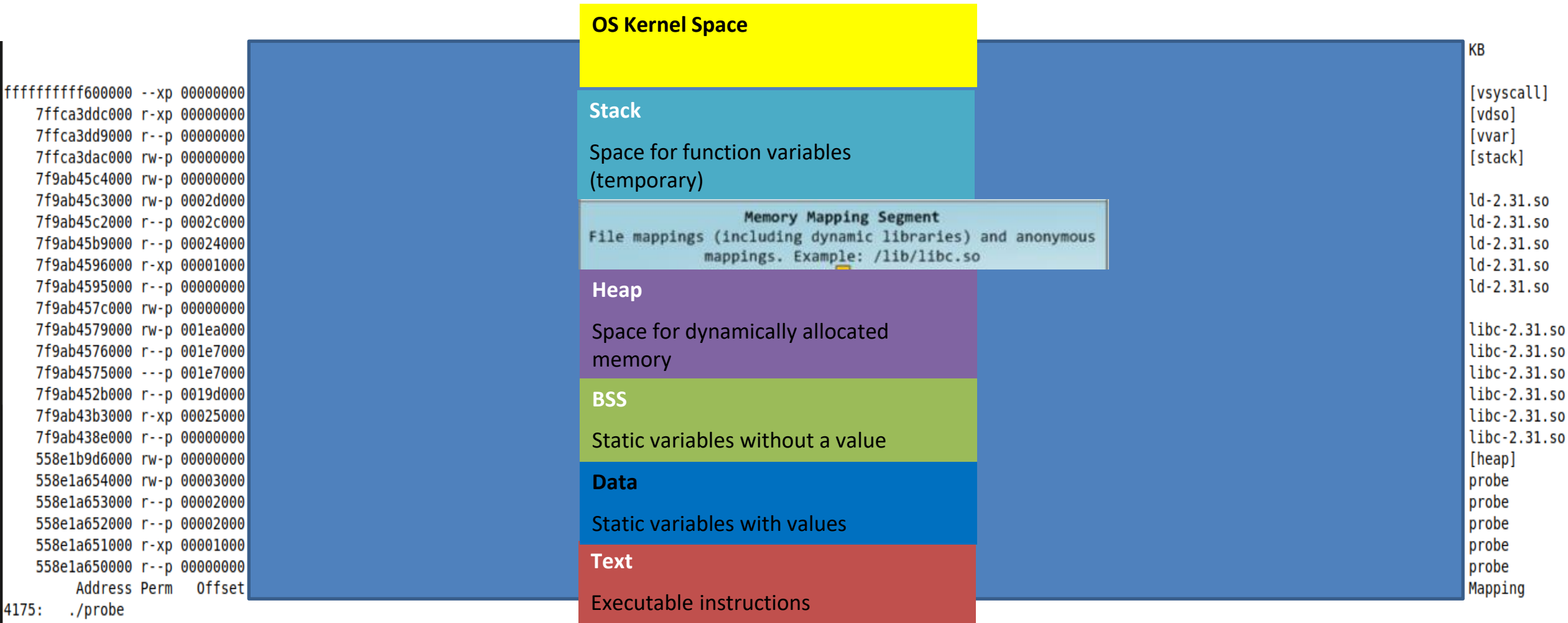
Applications Layout in Memory

Output of `pmap` (process mapping tool)



Applications Layout in Memory

Output of `pmap` (process mapping tool)



Applications Layout in Memory

Output of pmap (process mapping tool)

```
fffffffff600000 --xp 00000000
7ffca3ddc000 r-xp 00000000
7ffca3dd9000 r--p 00000000
7ffca3dac000 rw-p 00000000
7f9ab45c4000 rw-p 00000000
7f9ab45c3000 rw-p 0002d000
7f9ab45c2000 r--p 0002c000
7f9ab45b9000 r--p 00024000
7f9ab4596000 r-xp 00001000
7f9ab4595000 r--p 00000000
7f9ab457c000 rw-p 00000000
7f9ab4579000 rw-p 001ea000
7f9ab4576000 r--p 001e7000
7f9ab4575000 ---p 001e7000
7f9ab452b000 r--p 0019d000
7f9ab43b3000 r-xp 00025000
7f9ab438e000 r--p 00000000
558e1b9d6000 rw-p 00000000
558e1a654000 rw-p 00003000
558e1a653000 r--p 00002000
558e1a652000 r--p 00002000
558e1a651000 r-xp 00001000
558e1a650000 r--p 00000000
Address Perm Offset
4175: ./probe
```

-> the address of main = 0x558e1a651249
-> the address of printf = 0x7f9ab43f2e10
-> the address of getenv = 0x7f9ab43d7020
-> a stack address = 0x7ffca3dcb3b0
-> a global address = 0x558e1a6540c4
-> the argv address = 0x7ffca3dcb4f8
-> argv[0] = 0x7ffca3dcc45f
value is [./probe]
-> the environ address = 0x7ffca3dcb508
-> the envp address = 0x7ffca3dcb508
-> getenv("PWD") = 0x7ffca3dcc5ff
value is [/home/seed/os-review]
-> a heap address = 0x558e1b9d66b0

KB
[vsyscall]
[vdso]
[vvar]
[stack]
ld-2.31.so
ld-2.31.so
ld-2.31.so
ld-2.31.so
ld-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
[heap]
probe
probe
probe
probe
probe
Mapping

Applications Layout in Memory

Output of pmap (process mapping tool)

Where is "main" located in memory?

-> the address of main = 0x558e1a651249

-> the address of printf = 0x7ffca3dd2e10

-> the address of getenv = 0x7ffca3dd7020

-> a stack address = 0x7ffca3dcb3b0

-> a global address = 0x558e1a6540c4

-> the argv address = 0x7ffca3dcb4f8

-> argv[0] value is [./probe]

-> the environ address = 0x7ffca3dcb508

-> the envp address = 0x7ffca3dcb508

-> getenv("PWD") value is [/home/seed/os-review]

-> a heap address = 0x558e1b9d66b0

ffffffffff600000 --xp 00000000

7ffca3ddc000 r-xp 00000000

7ffca3dd9000 r--p 00000000

7ffca3dac000 rw-p 00000000

7f9ab45c4000 rw-p 00000000

7f9ab45c3000 rw-p 0002d000

7f9ab45c2000 r--p 0002c000

7f9ab45b9000 r--p 00024000

7f9ab4596000 r-xp 00001000

7f9ab4595000 r--p 00000000

7f9ab457c000 rw-p 00000000

7f9ab4579000 rw-p 001ea000

7f9ab4576000 r--p 001e7000

7f9ab4575000 ---p 001e7000

7f9ab452b000 r--p 0019d000

7f9ab43b3000 r-xp 00025000

7f9ab438e000 r--p 00000000

558e1b9d6000 rw-p 00000000

558e1a654000 rw-p 00003000

558e1a653000 r--p 00002000

558e1a652000 r--p 00002000

558e1a651000 r-xp 00001000

558e1a650000 r--p 00000000

Address Perm Offset

4175: ./probe

KB

[vsyscall]

[vdso]

[vvar]

[stack]

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

[heap]

probe

probe

probe

probe

Mapping

Applications Layout in Memory

Output of `pmap` (process mapping tool)

Where is "main" located in memory?

558e1a651000

r-xp

00001000

558e1a650000

558e1a653000

558e1a654000

558e1b9d6000

558e1b9d6000

558e1b9d6000

Address Perm Offset

4175: ./probe

the address of main

=

0x558e1a651249

the address of printf

=

0x7ff9ab43f2e10

the address of getenv

=

0x7ff9ab43d7020

a stack address

=

0x7ffca3dcb3b0

a global address

=

0x558e1a6540c4

the argv address

=

0x7ffca3dcb4f8

argv[0]

value is

['./probe']

the environ address

=

0x7ffca3dcb508

the envp address

=

0x7ffca3dcb508

getenv("PWD")

=

0x7ffca3dcc5ff

a heap address

=

0x558e1b9d66b0

Mapping

probe

probe

[heap]

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so


[stack]

[vvar]

[vdso]

[vsyscall]

`main` is code in our program, so it goes inside the text segment

 MONTANA
STATE UNIVERSITY

48

Applications Layout in Memory

Output of `pmap` (process mapping tool)

```

ffffffffff600000 --xp 00000000
7ffca3ddc000 r-xp 00000000
7ffca3dd9000 r--p 00000000
7ffca3dac000 rw-p 00000000
7f9ab45c4000 rw-p 00000000
7f9ab45c3000 rw-p 0002d000
7f9ab45c2000 r--p 0002c000
7f9ab45b9000 r--p 00024000
7f9ab4596000 r-xp 00001000
7f9ab4595000 r--p 00000000
7f9ab457c000 rw-p 00000000
7f9ab4579000 rw-p 001ea000
7f9ab4576000 r--p 001e7000
7f9ab4575000 ---p 001e7000
7f9ab452b000 r--p 0019d000
7f9ab43b3000 r-xp 00025000
7f9ab438e000 r--p 00000000
558e1b9d6000 rw-p 00000000
558e1a654000 rw-p 00003000
558e1a653000 r--p 00002000
558e1a652000 r--p 00002000
558e1a651000 r-xp 00001000
558e1a650000 r--p 00000000

```

Address Perm Offset

Where is “printf” located in memory?

```

-> the address of main      = 0x558e1a651249
-> the address of printf    = 0x7f9ab43f2e10
-> the address of getenv    = 0x7f9ab43d7020
-> a stack address         = 0x7ffca3dcb3b0
-> a global address        = 0x558e1a6540c4
-> the argv address        = 0x7ffca3dcb4f8
-> argv[0]                 = 0x7ffca3dcc45f
    value is [./probe]
-> the environ address     = 0x7ffca3dcb508
-> the envp address        = 0x7ffca3dcb508
-> getenv("PWD")           = 0x7ffca3dcc5ff
    value is [/home/seed/os-review]
-> a heap address          = 0x558e1b9d66b0

```

```

4175: ./probe

```

```

[vsyscall]
[vdso]
[vvar]
[stack]

ld-2.31.so
ld-2.31.so
ld-2.31.so
ld-2.31.so
ld-2.31.so

libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
[heap]
probe
probe
probe
probe
Mapping

```

KB

Applications Layout in Memory

Output of pmap (process mapping tool)

Where is "printf" located in memory?

-> the address of main = 0x558e1a651249

-> the address of printf = 0x7f9ab43f2e10

-> the address of getenv = 0x7f9ab43d7020

-> a stack address = 0x7ffca3dcb3b0

-> a global address = 0x558e1a6540c4

-> the argv address = 0x7ffca3dcb4f8

-> argv[0] value is [./probe]

-> the environ address = 0x7ffca3dcb508

-> the envp address = 0x7ffca3dcb508

-> getenv("PWD") value is [/home/seed/os-review]

-> a heap address = 0x558e1b9d66b0

7f9ab43b3000 r-xp 00025000

lib-2.31.so

```
fffffffb600000 --xp 00000000
7ffca3ddc000 r-xp 00000000
7ffca3dd9000 r--p 00000000
7ffca3dac000 rw-p 00000000
7f9ab45c4000 rw-p 00000000
7f9ab45c3000 rw-p 0002d000
7f9ab45c2000 r--p 0002c000
7f9ab45b9000 r--p 00024000
7f9ab4596000 r-xp 00001000
7f9ab4595000 r--p 00000000
7f9ab457c000 rw-p 00000000
7f9ab4579000 rw-p 001ea000
7f9ab4576000 r--p 001e7000
7f9ab4575000 ---p 001e7000
7f9ab452b000 r--p 0010d000
7f9ab43b3000 r-xp 00025000
7f9ab438e000 r--p 00000000
558e1b9d6000 rw-p 00000000
558e1a654000 rw-p 00003000
558e1a653000 r--p 00002000
558e1a652000 r--p 00002000
558e1a651000 r-xp 00001000
558e1a650000 r--p 00000000
Address Perm Offset
4175: ./probe
```

KB

[vsyscall]

[vdso]

[vvar]

[stack]

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

[heap]

probe

probe

probe

probe

probe

Mapping

printf is executable code from a shared library (libc) so we are in the memory mapping segment!

Applications Layout in Memory

Output of `pmap` (process mapping tool)

Where is "argv" located in memory?

-> the address of main = 0x558e1a651249

-> the address of printf = 0x7f9ab43f2e10

-> the address of getenv = 0x7f9ab43d7020

-> a stack address = 0x7ffca3dcb3b0

-> a global address = 0x558e1a6540c4

-> the argv address = 0x7ffca3dcb4f8

argv[0]

value is [./probe]

-> the environ address = 0x7ffca3dcb508

-> the envp address = 0x7ffca3dcb508

-> getenv("PWD") = 0x7ffca3dcc5ff

value is [/home/seed/os-review]

-> a heap address = 0x558e1b9d66b0

4175: ./probe

Address Perm Offset

KB

[vsyscall]

[vdso]

[vvar]

[stack]

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

ld-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

libc-2.31.so

[heap]

probe

probe

probe

probe

Mapping

`argv` is an array that holds the command line parameters passed into this program

Applications Layout in Memory

Output of `pmap` (process mapping tool)

Where is "argv" located in memory?

-> the address of main = 0x558e1a651249
-> the address of printf = 0x7f9ab43f2e10
-> the address of getenv = 0x7f9ab43d7020
-> a stack address = 0x7ffca3dcb3b0
-> a global address = 0x558e1a6540c4
-> the argv address = 0x7ffca3dcb4f8
-> **argv[0]** = 0x7ffca3dcc45f
value is [./probe]
-> the environ address = 0x7ffca3dcb508
-> the envp address = 0x7ffca3dcb508
-> getenv("PWD") = 0x7ffca3dcc5ff
value is [/home/seed/os-review]
-> a heap address = 0x558e1b9d66b0

ffffffffff600000	--xp	00000000
7ffca3ddc000	r-xp	00000000
7ffca3dd0000	r--p	00000000
7ffca3dac000	rw-p	00000000
7f9ab45c4000	rw-p	00000000
7f9ab45c3000	rw-p	0002d000
7f9ab45c2000	r--p	0002c000
7f9ab45b9000	r--p	00024000
7f9ab4596000	r-xp	00001000
7f9ab4595000	r--p	00000000
7f9ab457c000	rw-p	00000000
7f9ab4579000	rw-p	001ea000
7f9ab4576000	r--p	001e7000
7f9ab4575000	---p	001e7000
7f9ab452b000	r--p	0019d000
7f9ab43b3000	r-xp	00025000
7f9ab438e000	r--p	00000000
558e1b9d6000	rw-p	00000000
558e1a654000	rw-p	00003000
558e1a653000	r--p	00002000
558e1a652000	r--p	00002000
558e1a651000	r-xp	00001000
558e1a650000	r--p	00000000
Address Perm Offset		
4175: ./probe		

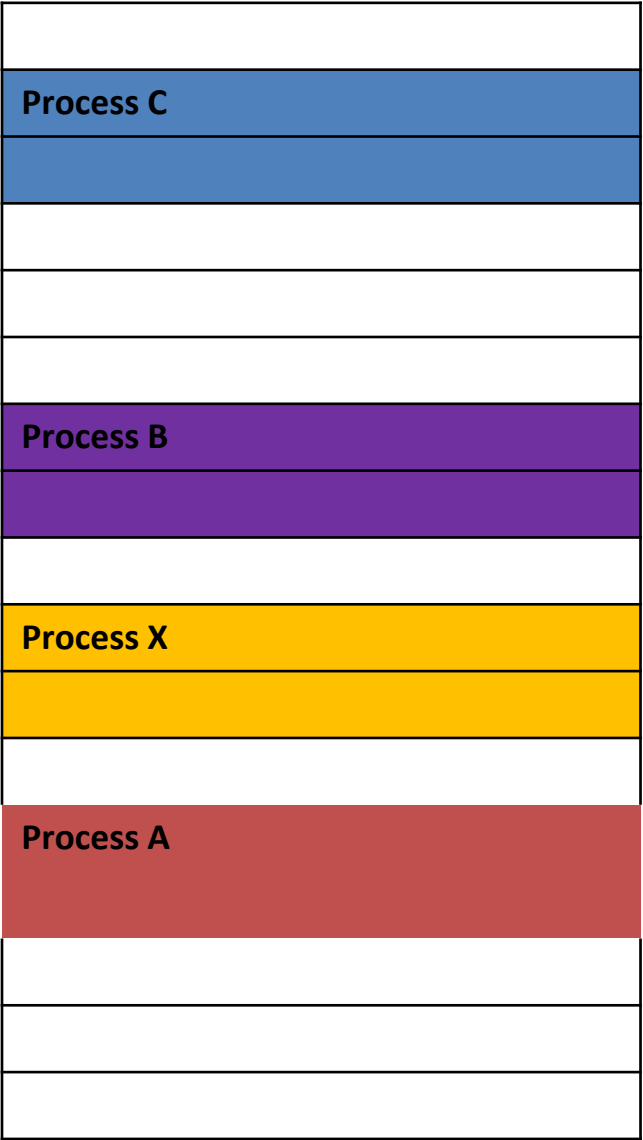
KB

[vsyscall]
[vdso]
[tunctl]
[stack]
ld-2.31.so
ld-2.31.so
ld-2.31.so
ld-2.31.so
ld-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
libc-2.31.so
[heap]
probe
probe
probe
probe
probe
Mapping

argv is the argument to the main function, so we are in the stack!

Applications Layout in Memory

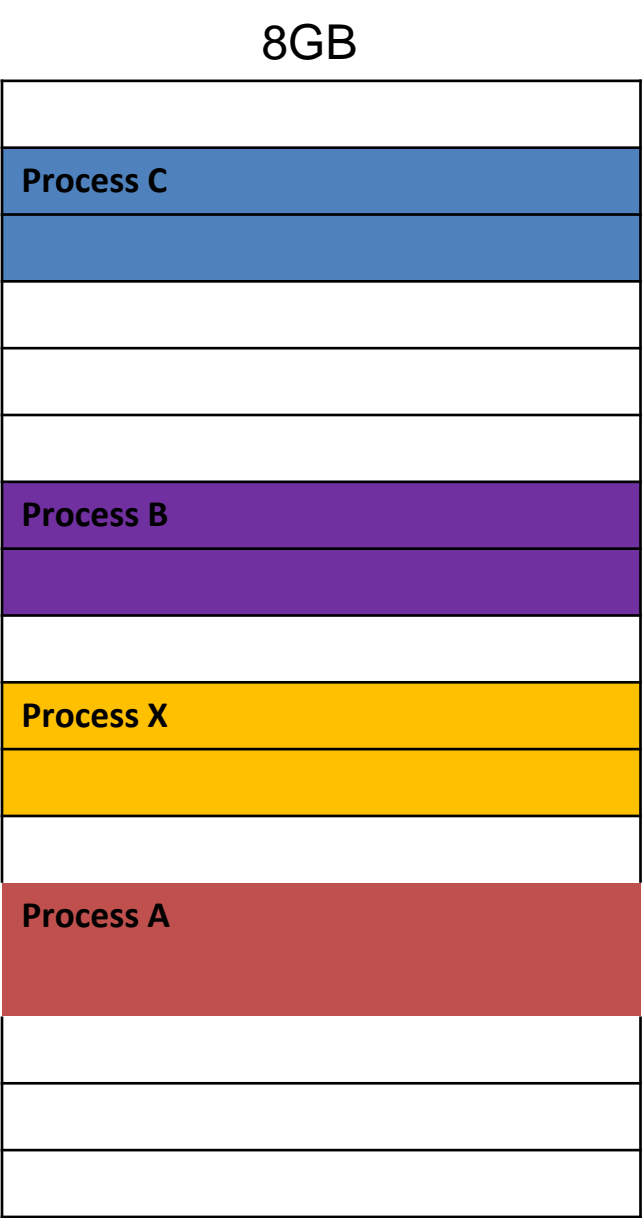
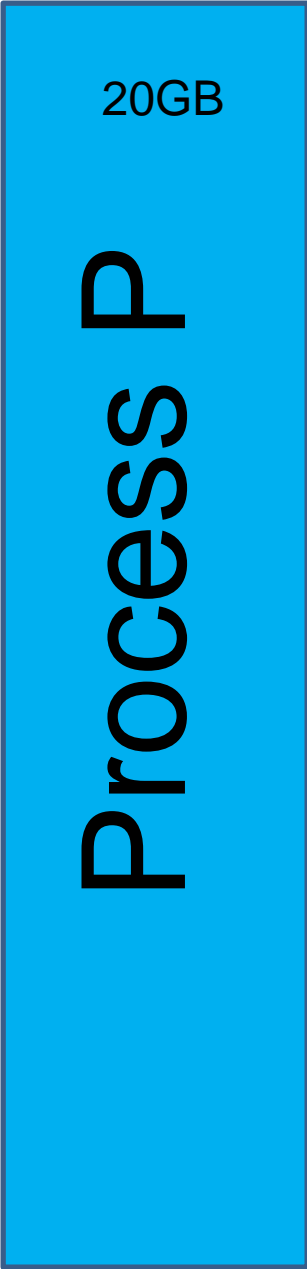
We have many programs that are actively running on our computer



Applications Layout in Memory

We have many programs that are actively running on our computer

What if we have a program that is bigger than our entire main memory?

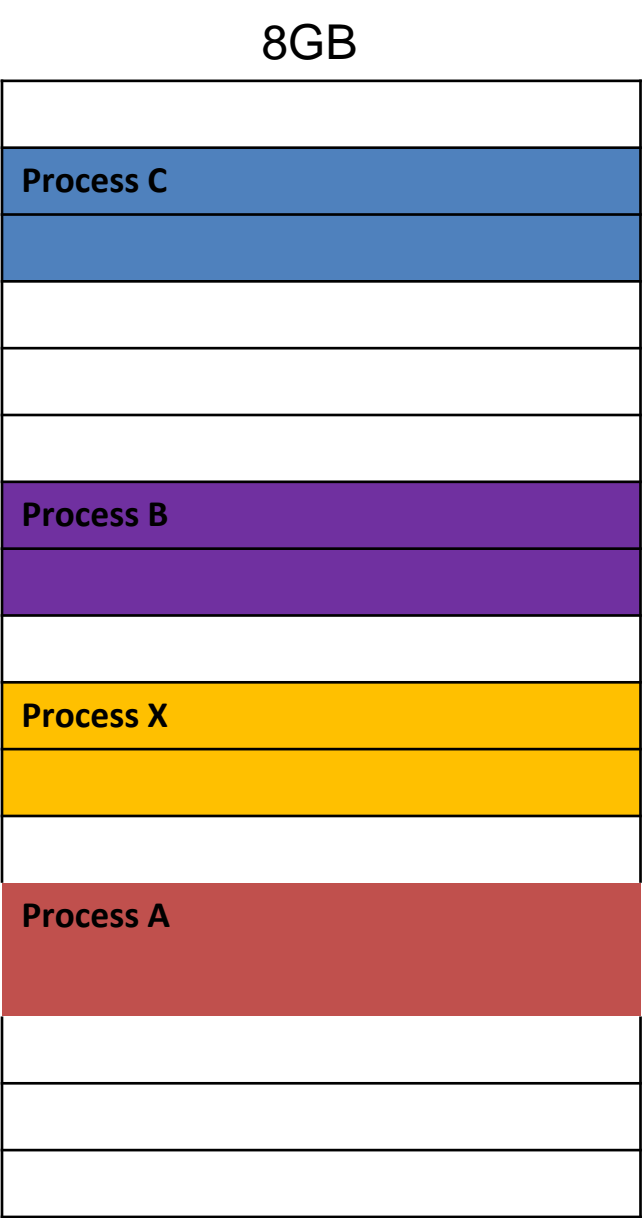
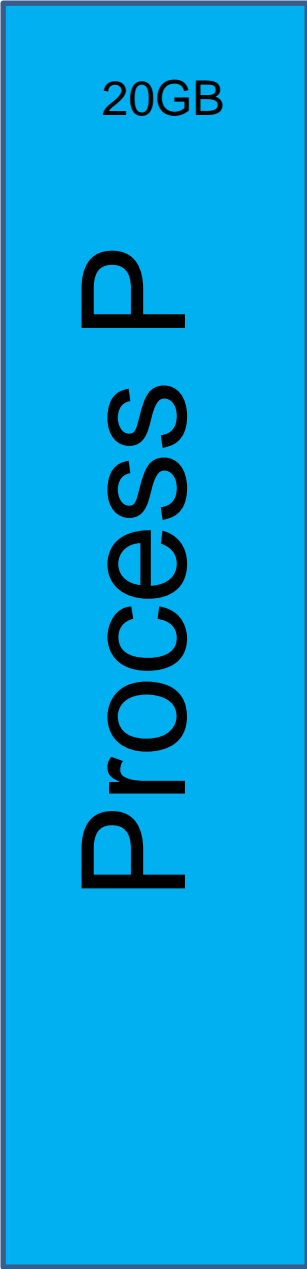


Applications Layout in Memory

We have many programs that are actively running on our computer

What if we have a program that is bigger than our entire main memory?

Does our computer crash?



Memory management

Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage



Secondary Storage

8GB

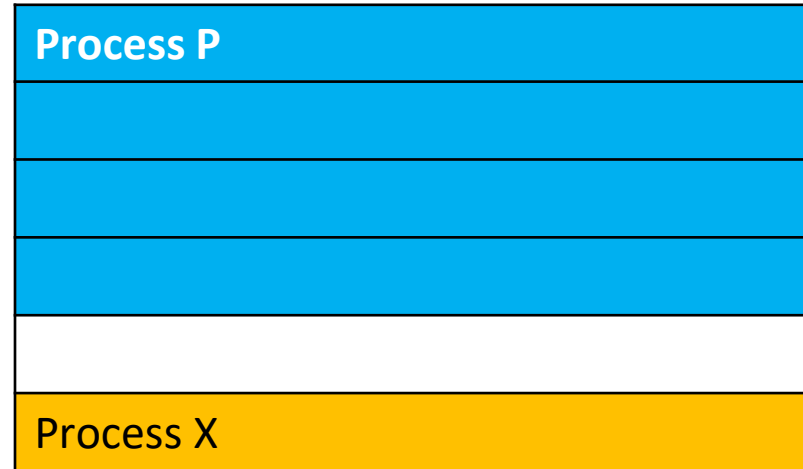
Process C
Process B
Process X
Process A

Memory management

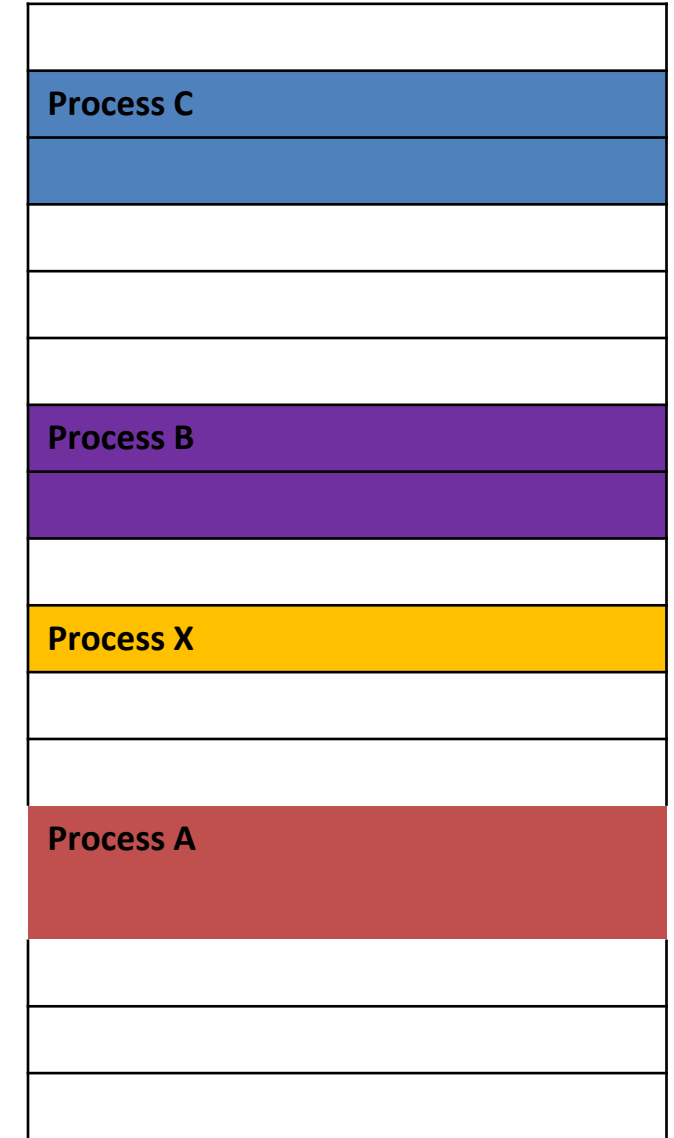
Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage

We split the process into smaller **pages**. Load pages into memory only when needed

Secondary Storage



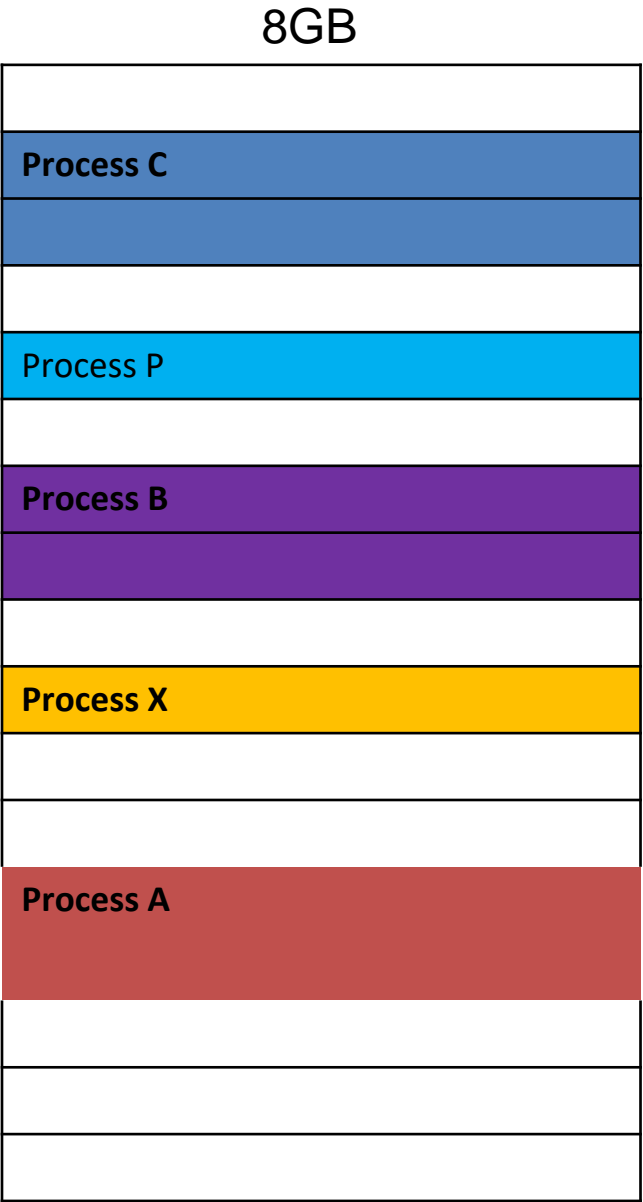
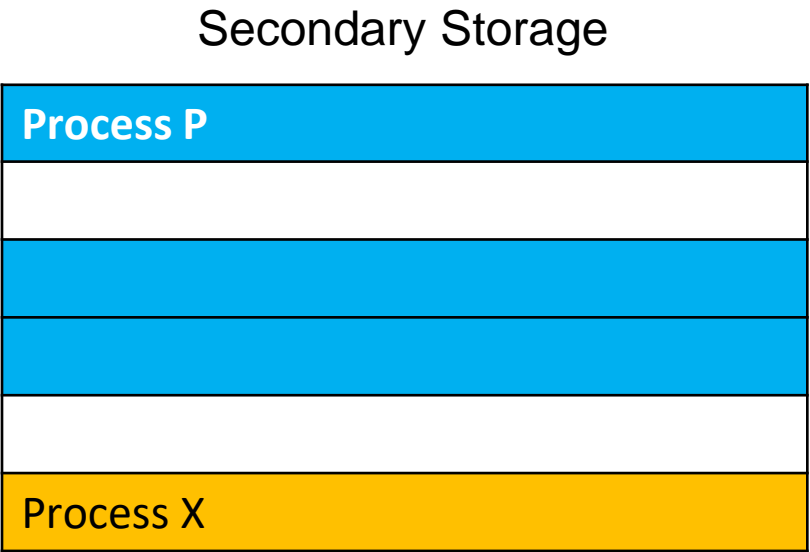
8GB



Memory management

Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage

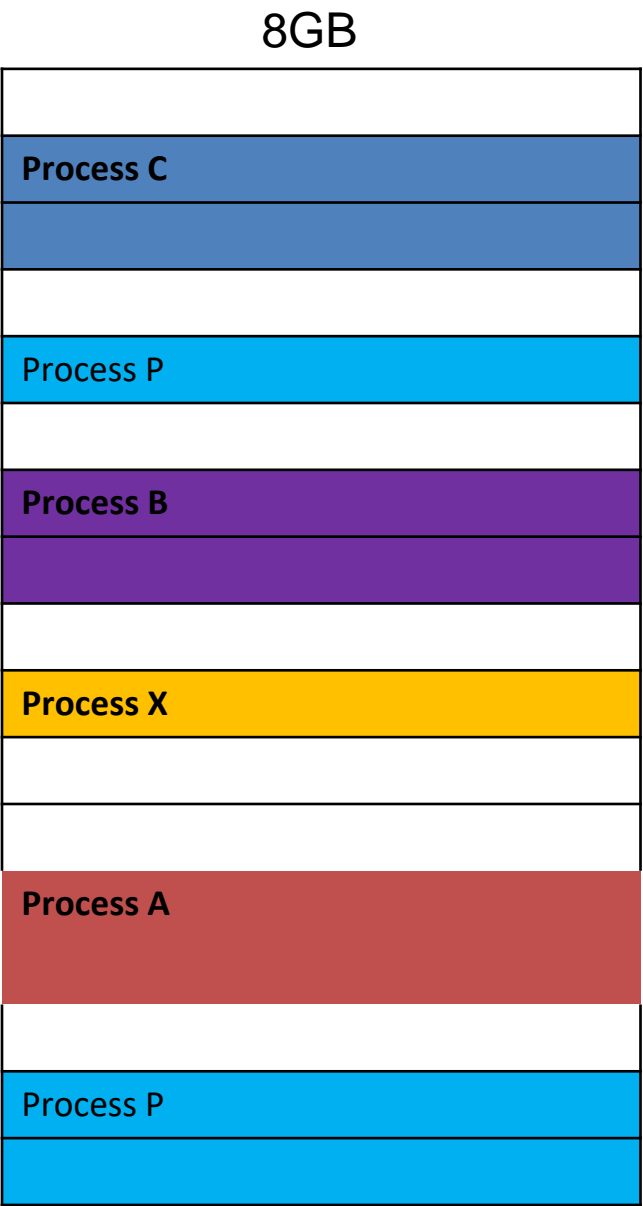
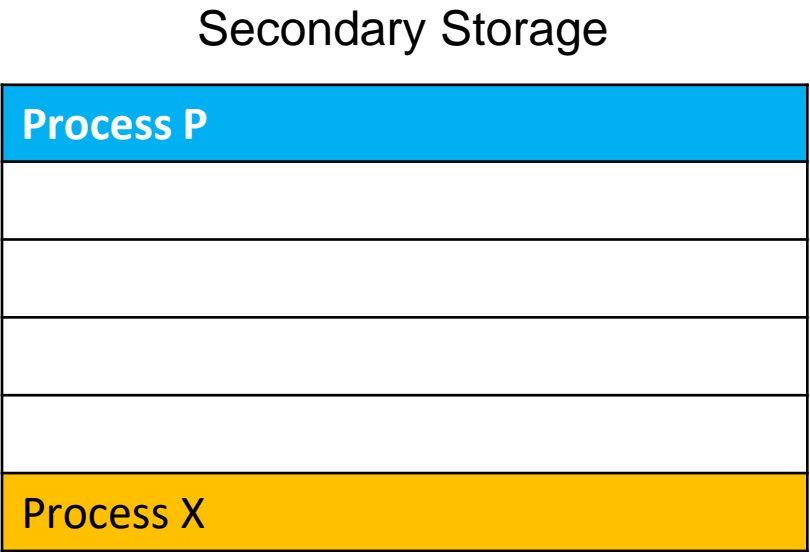
We split the process into smaller **pages**. Load pages into memory only when needed



Memory management

Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage

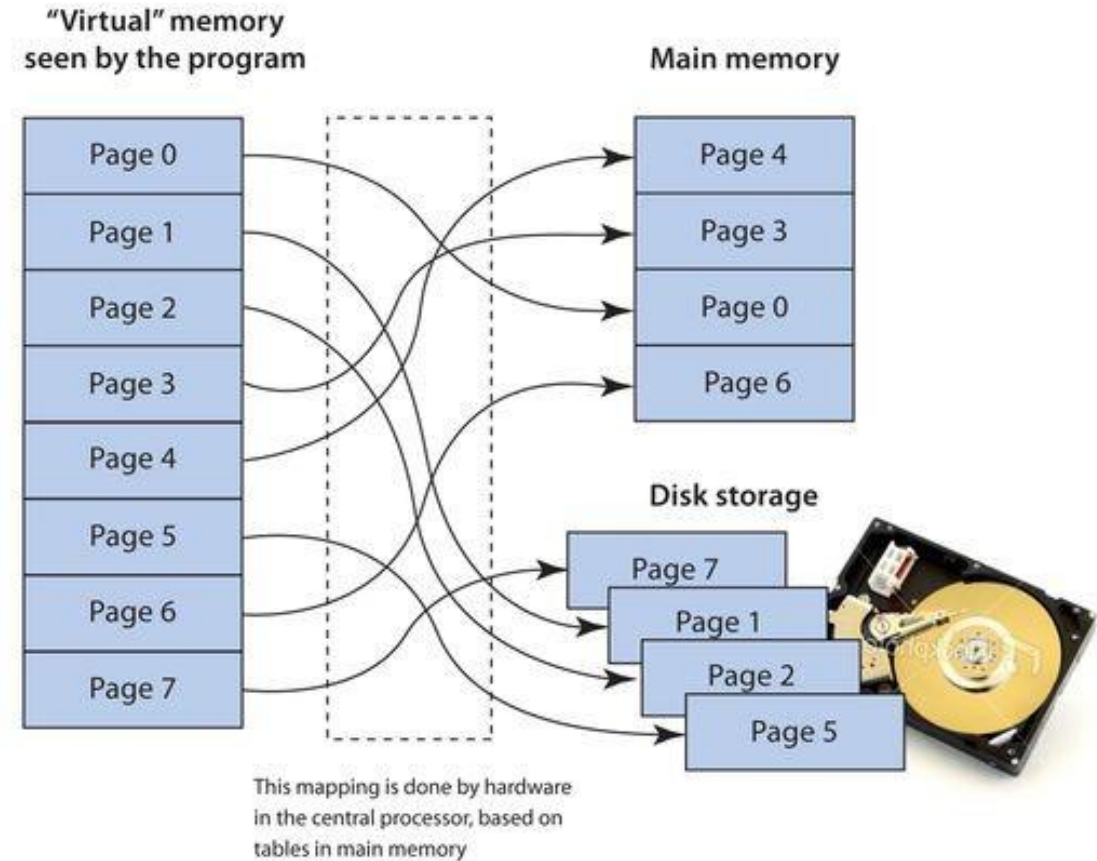
We split the process into smaller **pages**. Load pages into memory only when needed



Memory management

Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage

We split the process into smaller **pages**. Load pages into memory only when needed

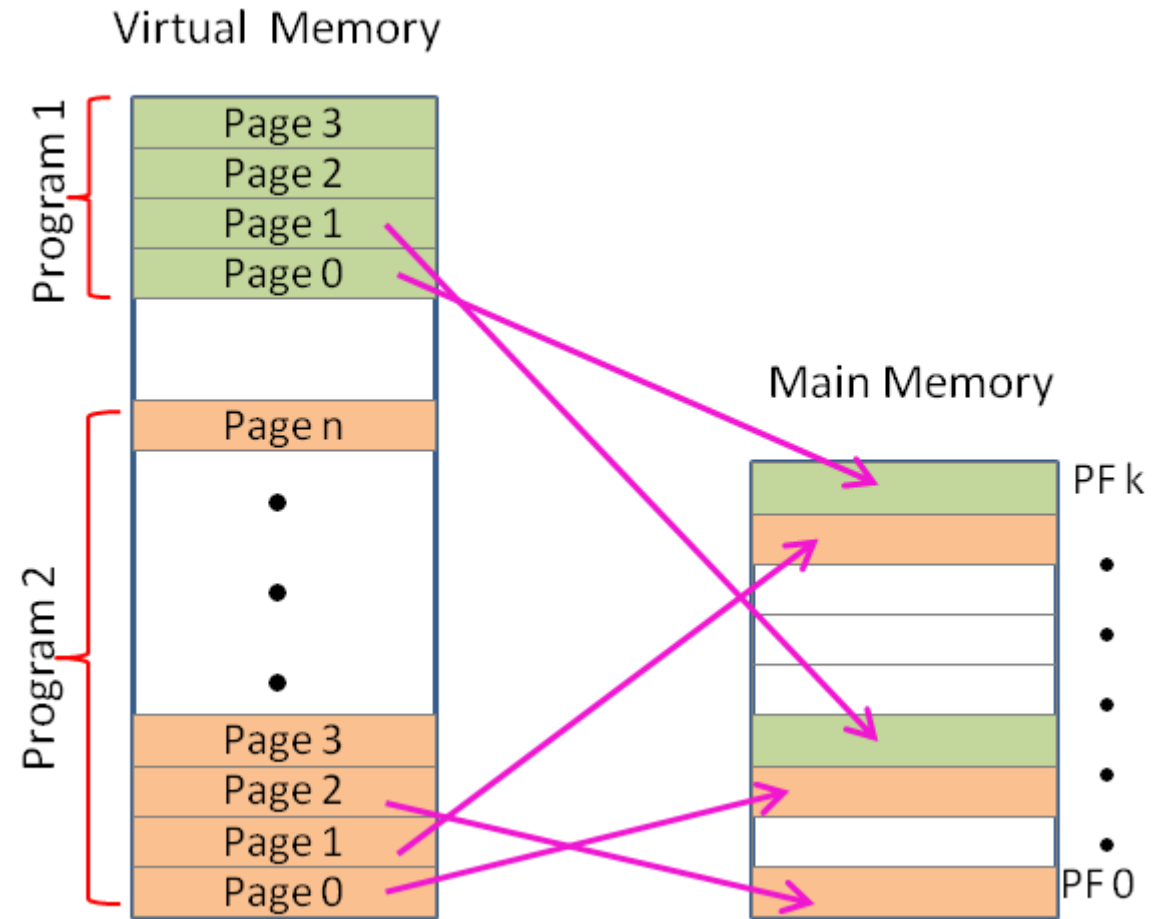


Constantly swapping stuff in and out of main memory

Memory management

Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage

We split the process into smaller **pages**. Load pages into memory only when needed

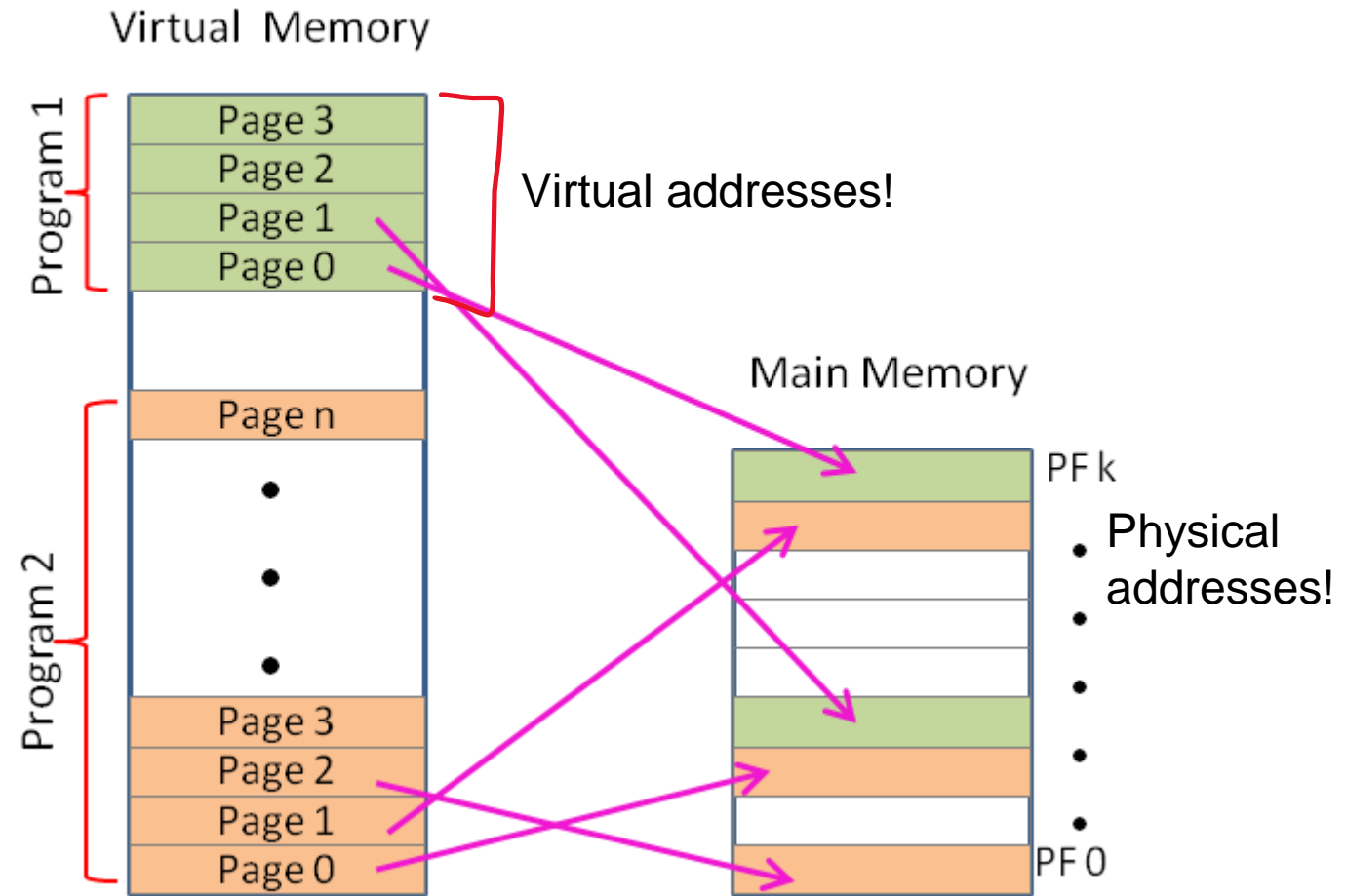


A process in memory is not contiguous

Memory management

Virtual Memory uses secondary storage to give programs the illusion that they have infinite storage

We split the process into smaller **pages**. Load pages into memory only when needed



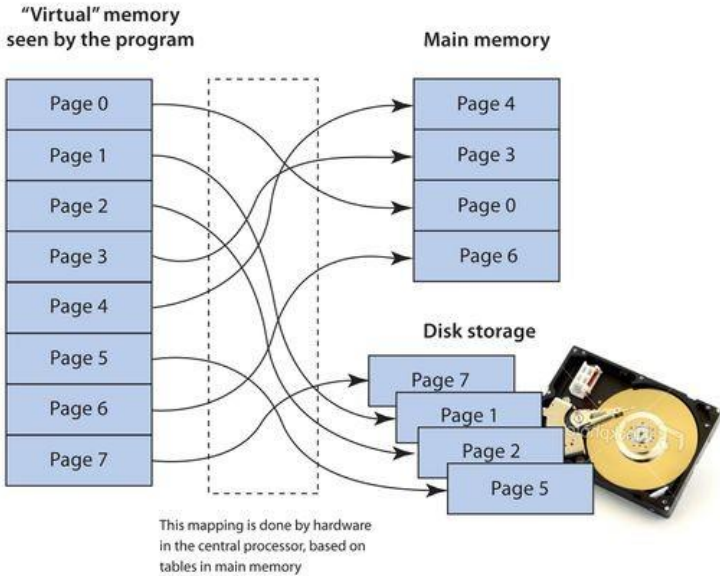
A process in memory is not contiguous

In probe.c, we are seeing virtual addresses!

OS Review

Memory Manager

- Manages how physical memory is utilized



Process Manager

- Manages how processes are structured and how to handle many processes running at once



Interface Manager

- Manages communication between apps and hardware

