CSCI 132: Basic Data Structures and Algorithms

Lessons Learned so far + Intro to Stacks

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Fall 2023

https://www.cs.montana.edu/pearsall/classes/fall2023/132/main.html



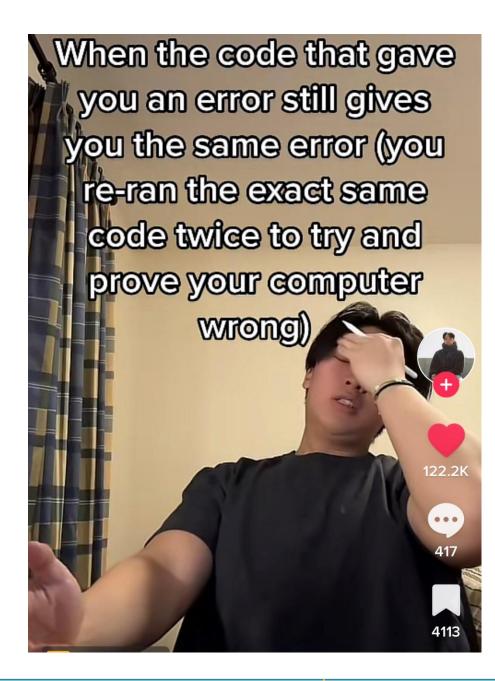
Announcements

Come get your midterm exam

- Exam average was in the 80s
- Don't stress if you didn't do well
- Make sure I calculated your score correctly

Lab 7 and Program 3 will be posted very soon

Class Registration \rightarrow CSCI 232 \rightarrow CS 145 \rightarrow CSCI 215 \rightarrow CSCI 246





Big-O

Big-O notation is a way to describe the running-time/time complexity of an algorithm regarding the number of operations that are executed in the algorithm (in relation to some input n)

• Focus on worst-case scenario, and how the algorithm scales as n gets really big



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Big-O

Big-O notation is a way to describe the running-time/time complexity of an algorithm regarding the number of operations that are executed in the algorithm (in relation to some input n)

• Focus on worst-case scenario, and how the algorithm scales as n gets really big

A very powerful computer and a very weak computer running the same algorithm will both execute the same number of operations (the speed at which they execute these operations will be different)

Takeaway: the asymptotic running time (the big-o running time) will be the same for each computer



Big-O notation is a way to describe the running-time/time complexity of an algorithm regarding the number of operations that are executed in the algorithm (in relation to some input n)

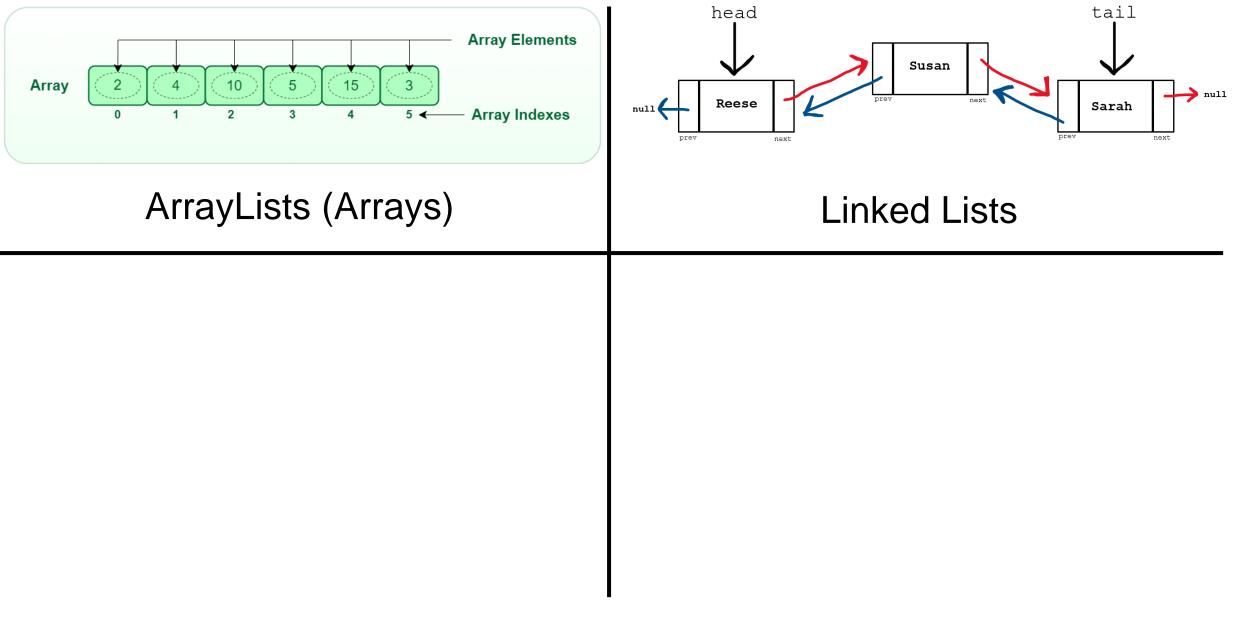
• Focus on worst-case scenario, and how the algorithm scales as n gets really big

To find the total running time of an algorithm, we calculate the runningtime of each operation in the algorithm and then add everything together

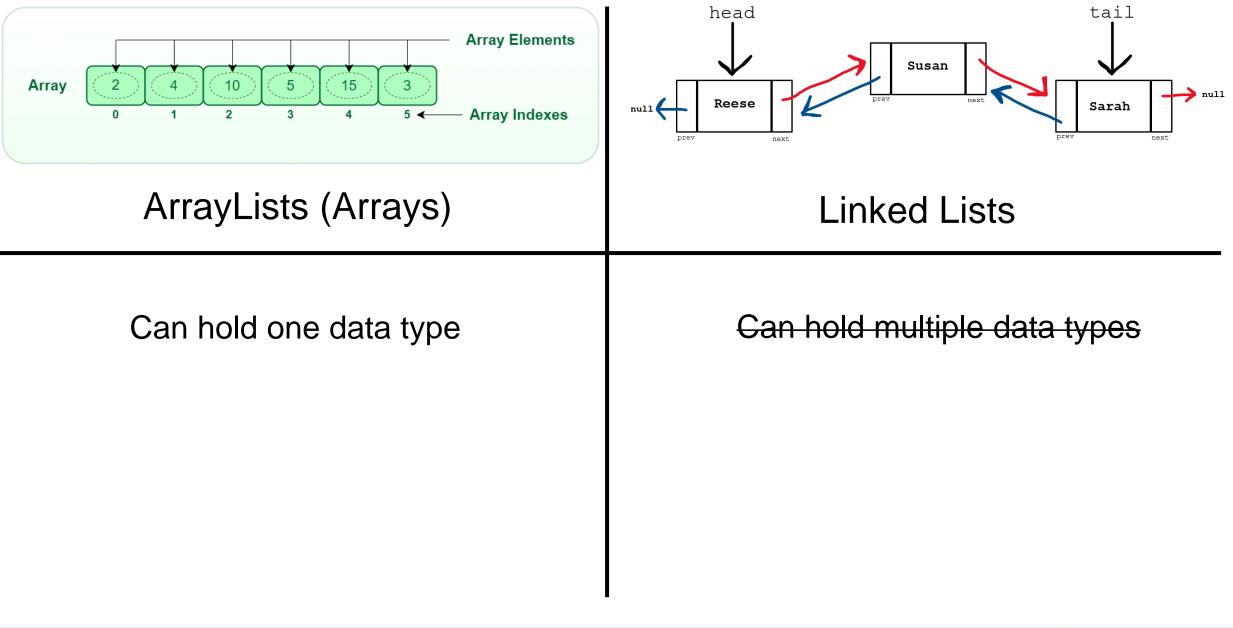
In Big-O, we can drop non-dominant factors and multiplicative constants (coefficients)

O(n) + O(n) + O(n): Total running time = $O(3n) \in O(n)$

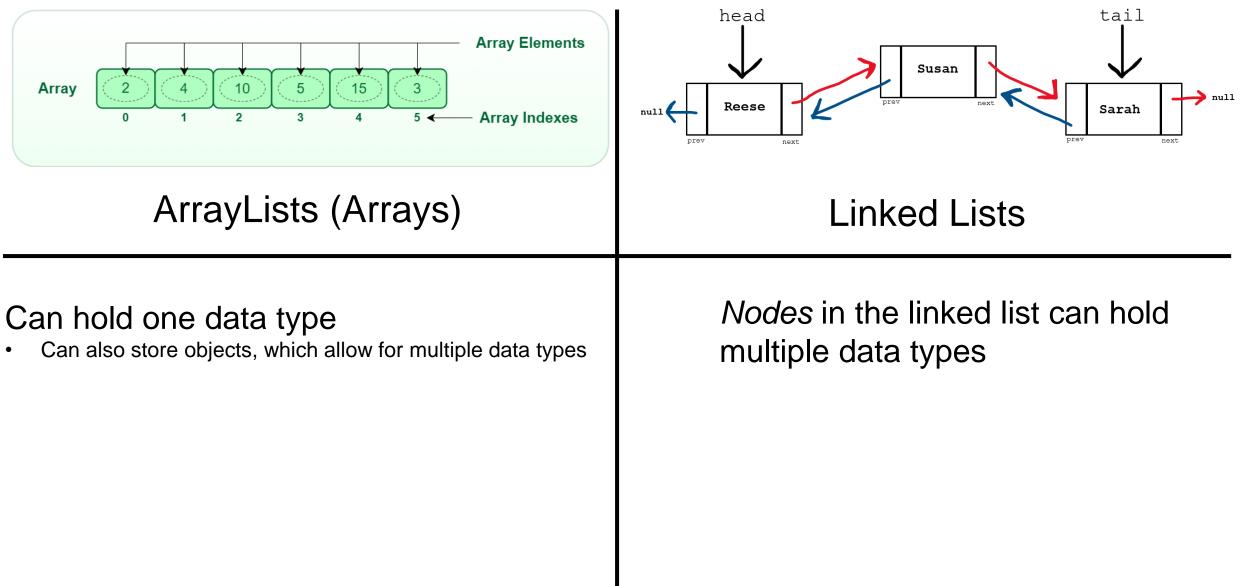




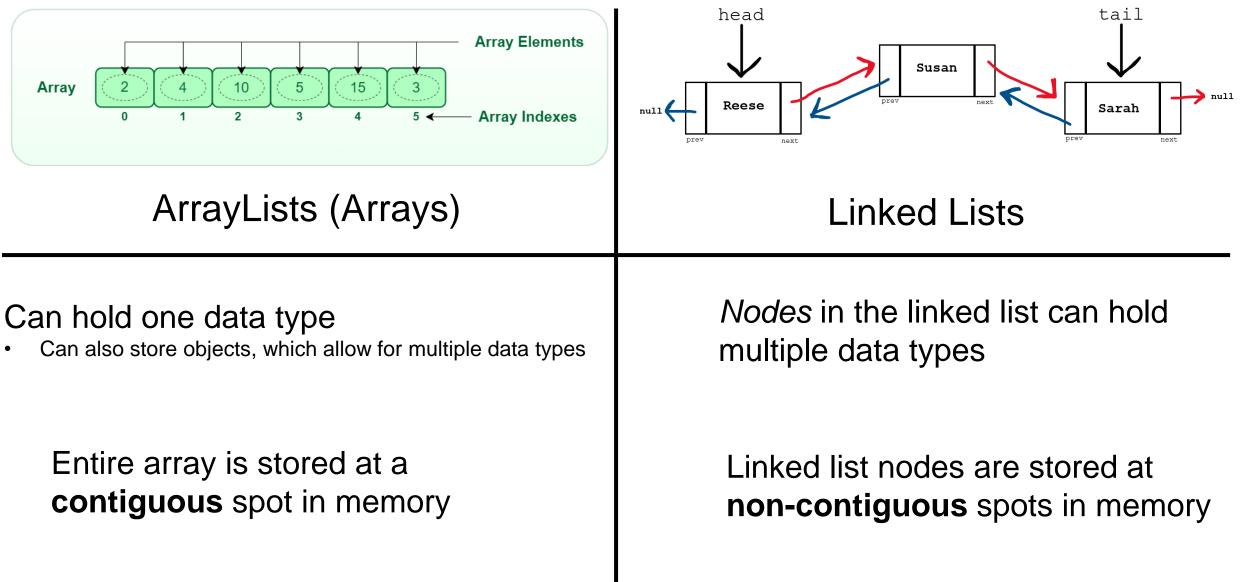




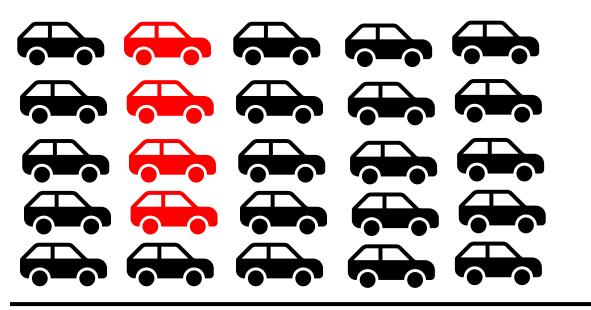














Can hold one data type

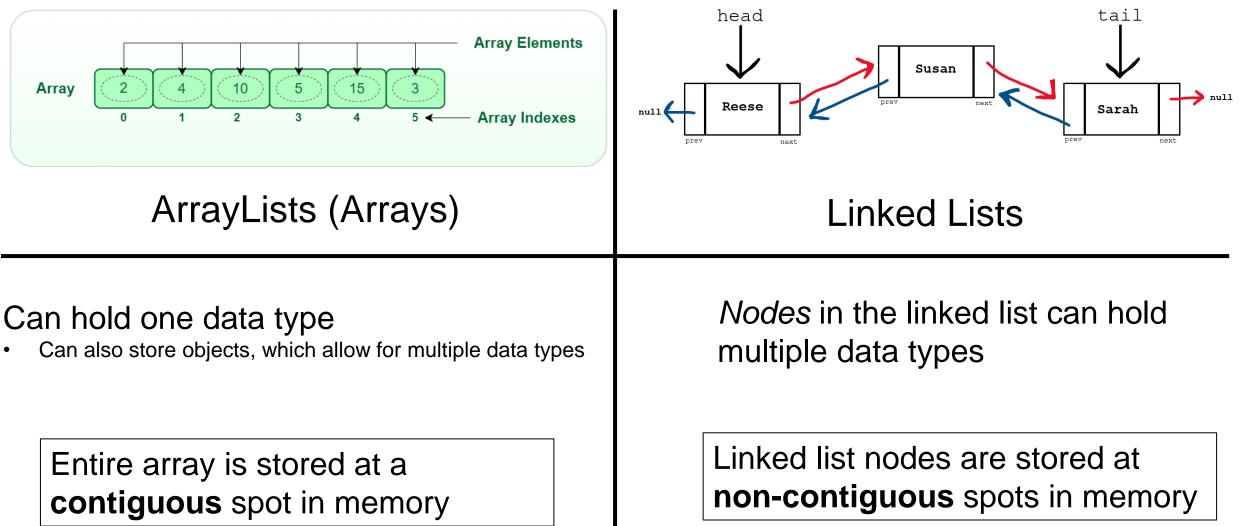
• Can also store objects, which allow for multiple data types

Entire array is stored at a **contiguous** spot in memory

Nodes in the linked list can hold multiple data types

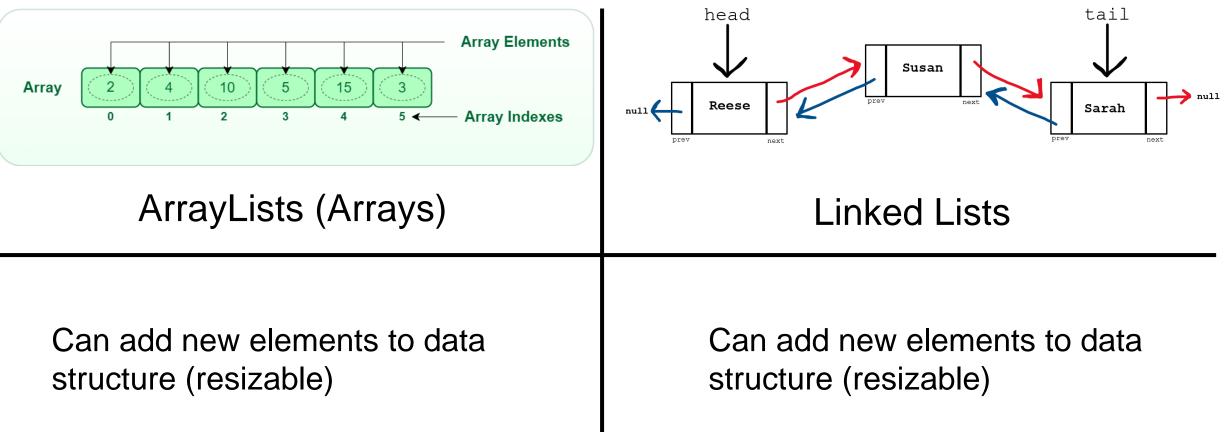
Linked list nodes are stored at **non-contiguous** spots in memory





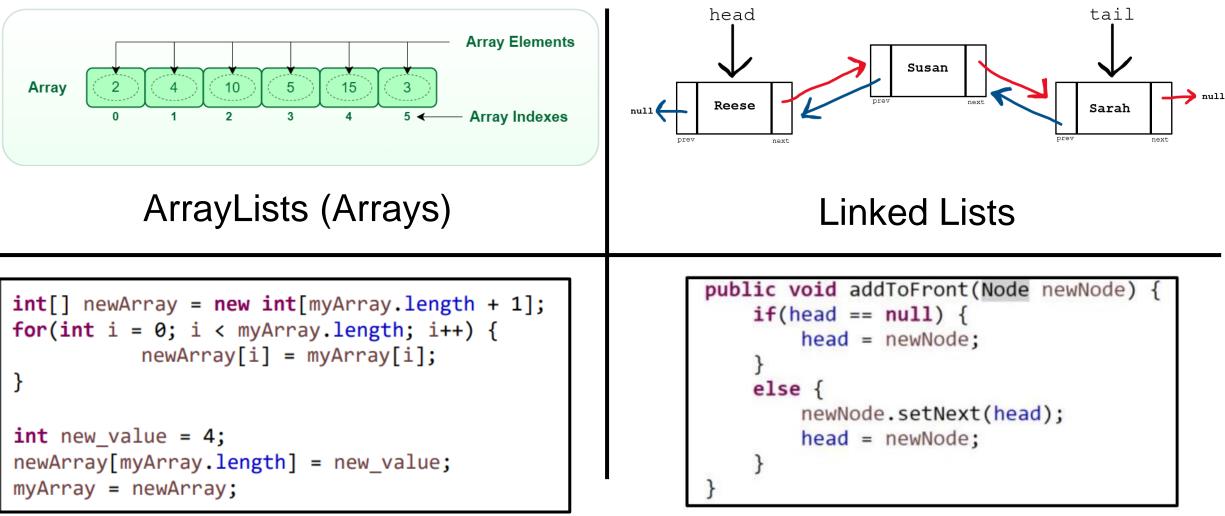
Traversing a linked list requires more work than traversing an array





Both data structures can grow dynamically, and new elements can be added, but they way they add new elements is drastically different

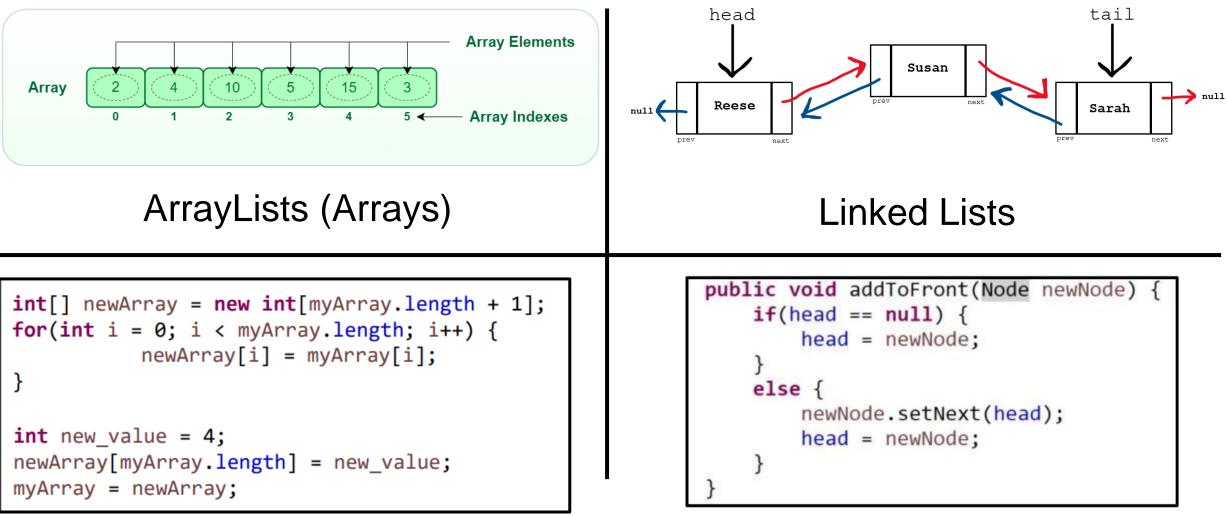




Create a brand-new array, copy everything over from old array

Update pointers



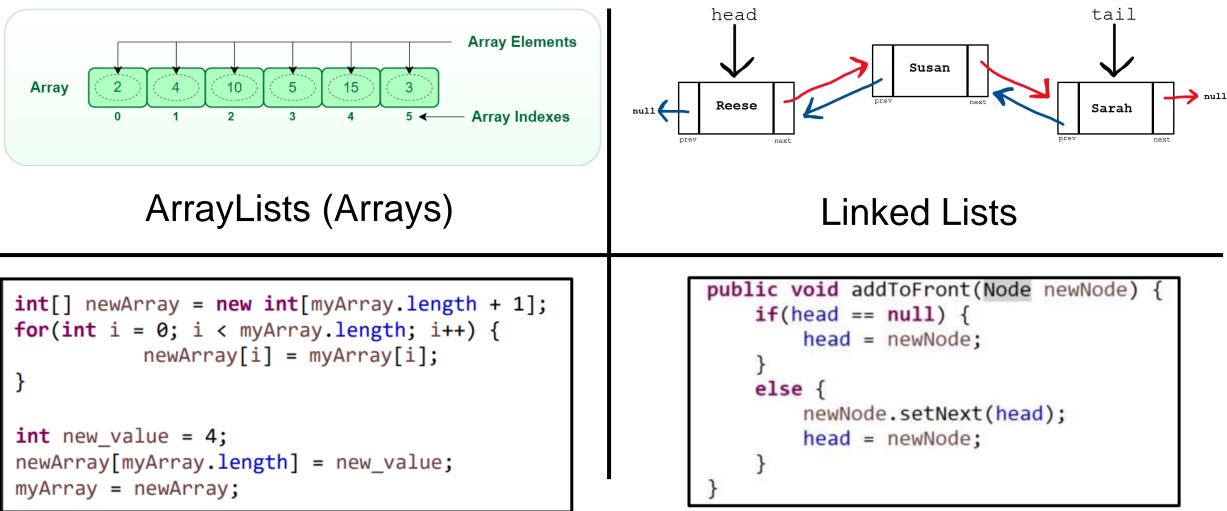


Create a brand-new array, copy everything over from old array O(n)

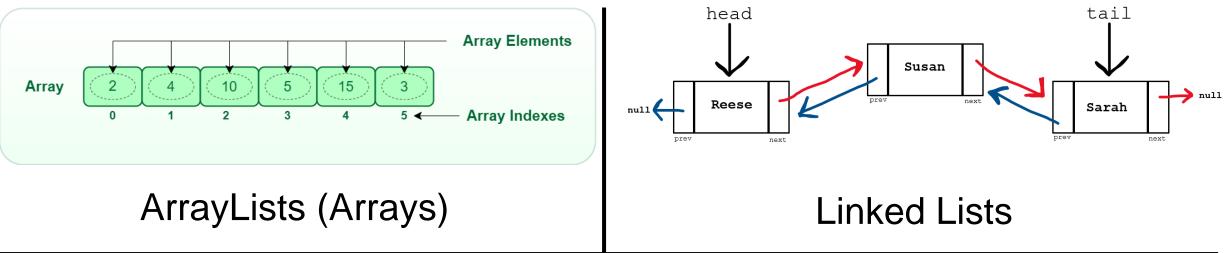
Update pointers



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Takeaway: Adding a new element to an ArrayList requires much more work than adding a new element to a Linked List

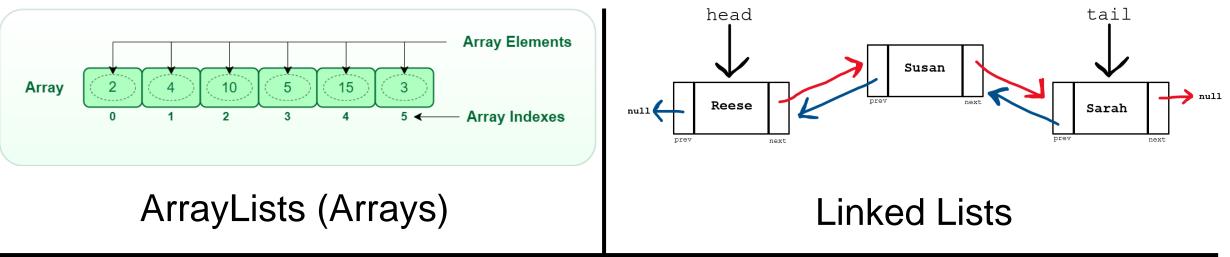


Arrays are generally much easier to sort than Nodes in a Linked List

If you are constantly needing to add new elements to the data structure, using a Linked List requires much less work in the long run

Arrays are more memory efficient (adding is not very memory efficient though)





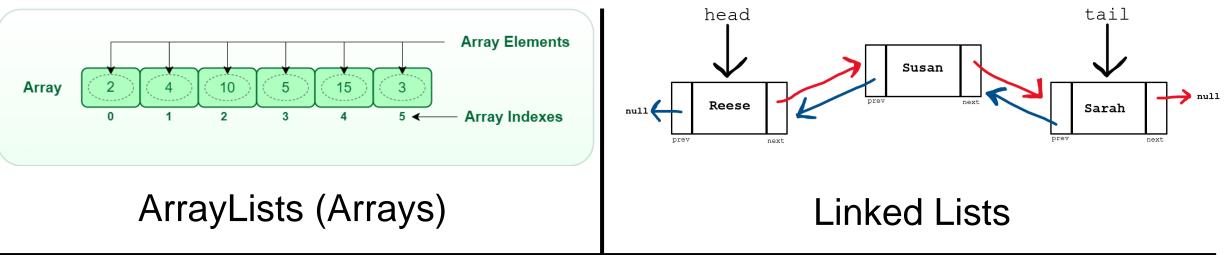
When to use each data structure?

It depends on how you are using your data and if you know how much data you have

If you don't know how much data you need to store, or if you are constantly needing to add new elements to the data structure \rightarrow Linked Lists

If you know how much data you need to store, and if you can add all your data at once \rightarrow Arrays/ArrayLists





These two data structures are <u>implementations</u> of a **List** Abstract Data Type (ADT)

ADT is a class whose behavior is defined by a set of operations and how a user interacts with it.

A list data type must be able to **get** an element, **add** an element, **remove** an element, etc \rightarrow How they do these operations is up to the subclass (LL and AL)

As programmers, we use handy methods that were written by other people that allows us to use these data structures



We will no longer be writing our own Linked List class, instead we will now import the Java-provided Linked List Class

import java.util.LinkedList;



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```
import java.util.LinkedList;
```

```
LinkedList<String> names = new LinkedList<String>();
The data type the
linked list will be
holding
```



The Linked List Class

https://docs.oracle.com/javase/7/docs/api/java/util/LinkedList.html

public class LinkedList<E>
extends AbstractSequentialList<E>
implements List<E>, Deque<E>, Cloneable, Serializable

Doubly-linked list implementation of the List and Deque interfaces. Implements all optional list operations, and permits all elements (including null).

All of the operations perform as could be expected for a doubly-linked list. Operations that index into the list will traverse the list from the beginning or the end, whichever is closer to the specified index.

The **documentation** describe how the LinkedList class was implemented, and all the methods/operations we can do with the Linked List class

Methods	
Modifier and Type	Method and Description
boolean	add(E_e) Appends the specified element to the end of this list.
void	add(int index, E element) Inserts the specified element at the specified position in this list.
boolean	addAll(Collection extends E c) Appends all of the elements in the specified collection to the end of this list, in the order that they are returned by the sp
boolean	addAll(int index, Collection extends E c) Inserts all of the elements in the specified collection into this list, starting at the specified position.
void	addFirst(E e) Inserts the specified element at the beginning of this list.
void	addLast(E e) Appends the specified element to the end of this list.
void	clear() Removes all of the elements from this list.
Object	clone() Returns a shallow copy of this LinkedList.
boolean	contains(Object o) Returns true if this list contains the specified element.
Iterator <e></e>	descendingIterator() Returns an iterator over the elements in this deque in reverse sequential order.
E	element() Retrieves, but does not remove, the head (first element) of this list.
E	get(int_index) Returns the element at the specified position in this list.
E	getFirst() Raturns the first element in this list

when you start coding in a new language without reading the documentation:

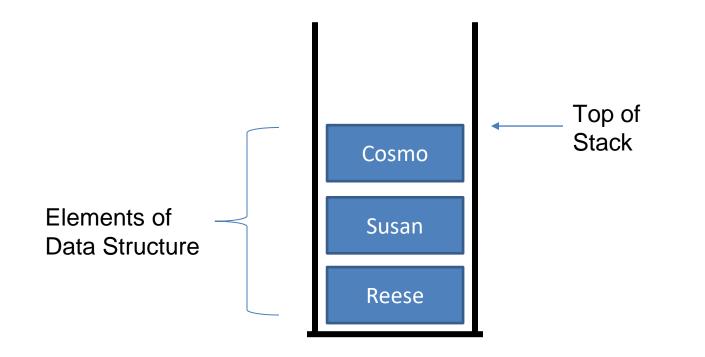




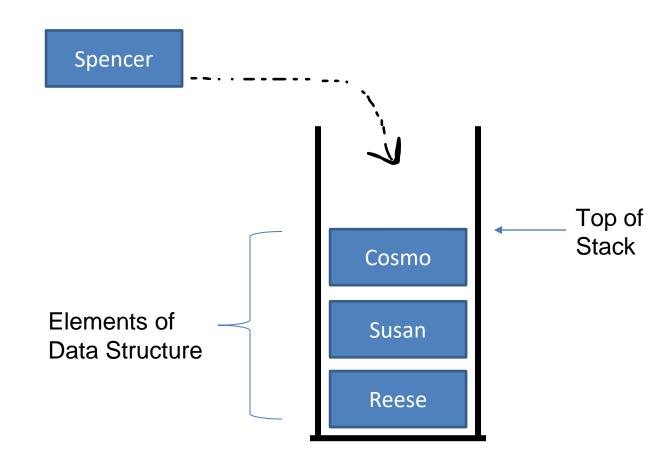
The Linked List Class

```
import java.util.LinkedList;
public class march20demo {
    public static void main(String[] args) {
        LinkedList<String> names = new LinkedList<String>();
        names.add("Reese");
        names.add("Spencer");
        names.add("Susan");
        System.out.println(names);
    }
```



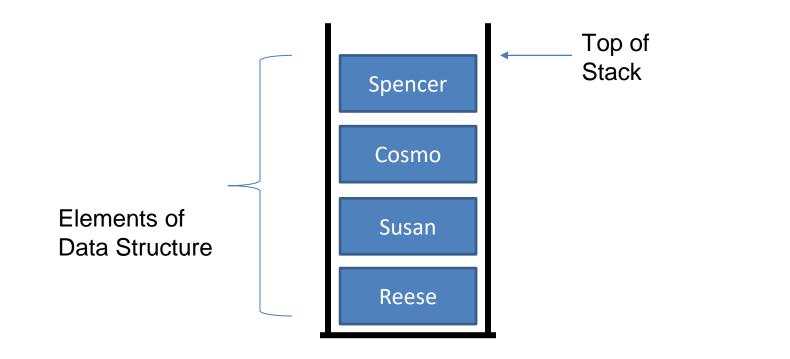






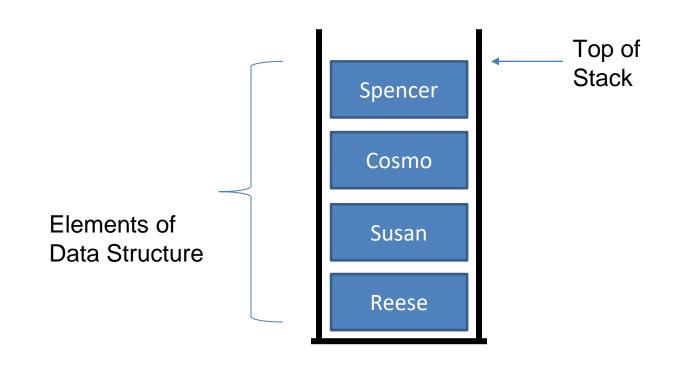
When only interact with the top of the stack.

If we want to add a new element, we must put it on the top of the stack



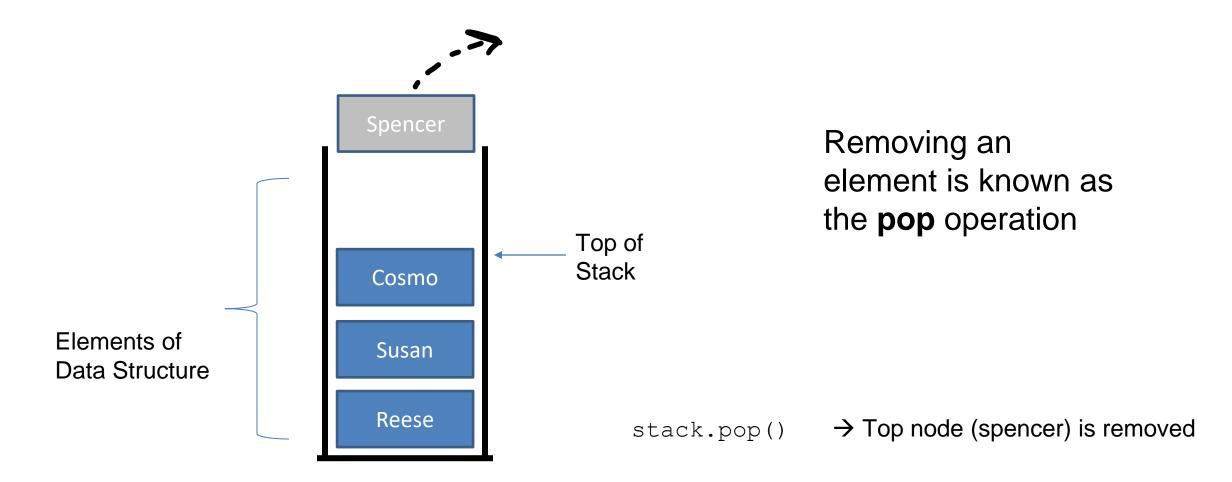
Adding something to a stack is known as the **push** operation





If we want to remove something, we must always remove the element on the top of the stack

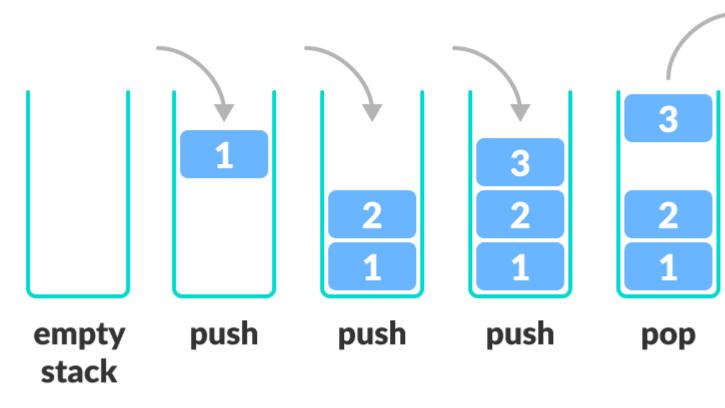






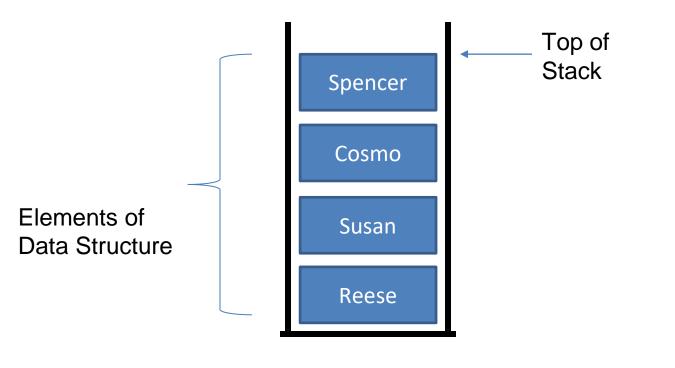
We can:

- Add an element to the top of the stack (push)
- Remove the top element (pop)



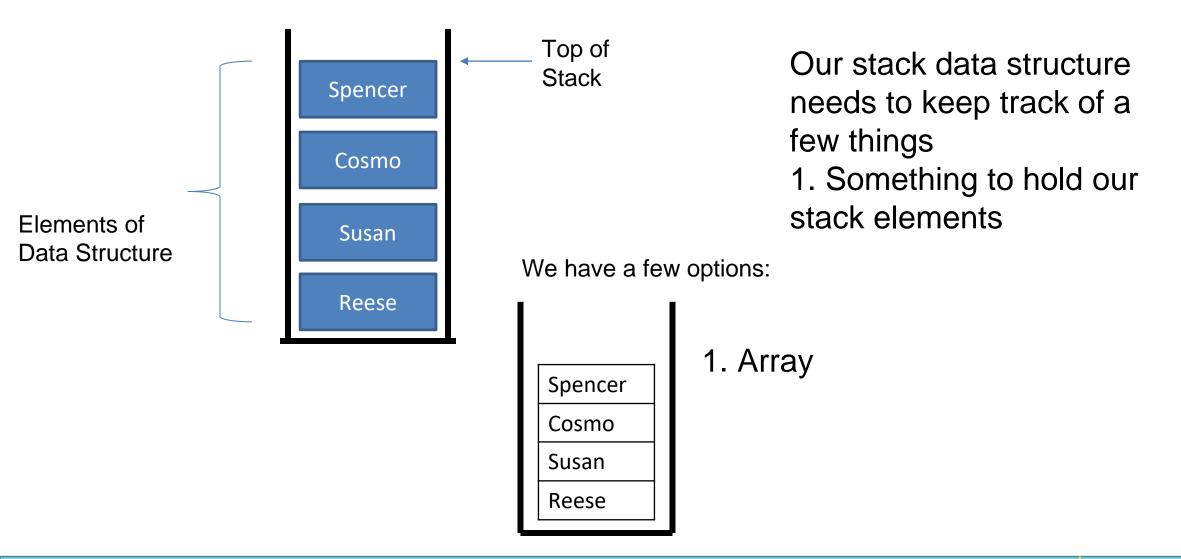




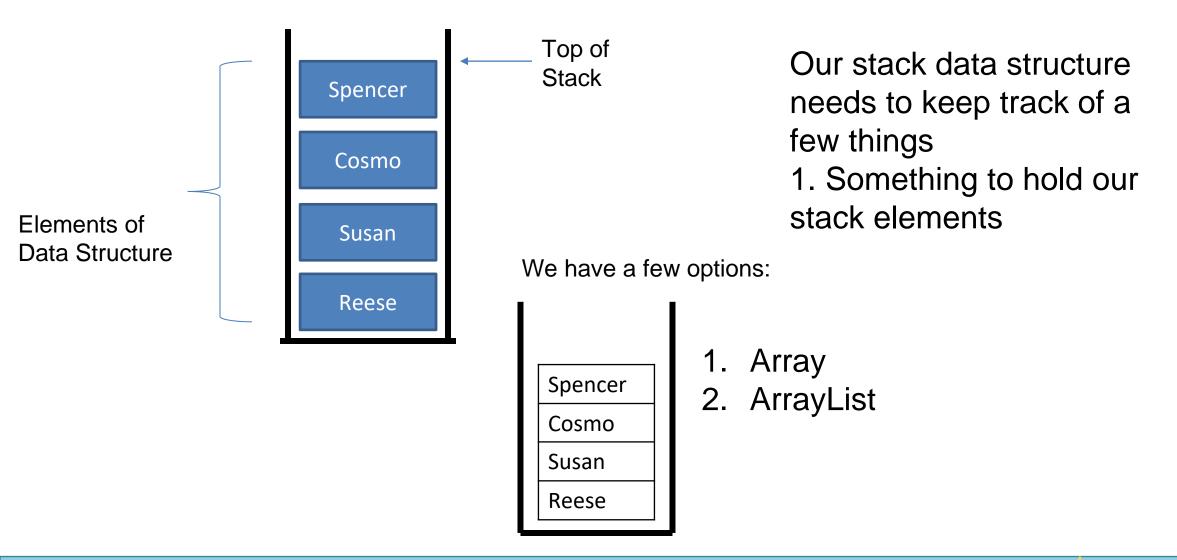


Our stack data structure needs to keep track of a few things 1. Something to hold our stack elements

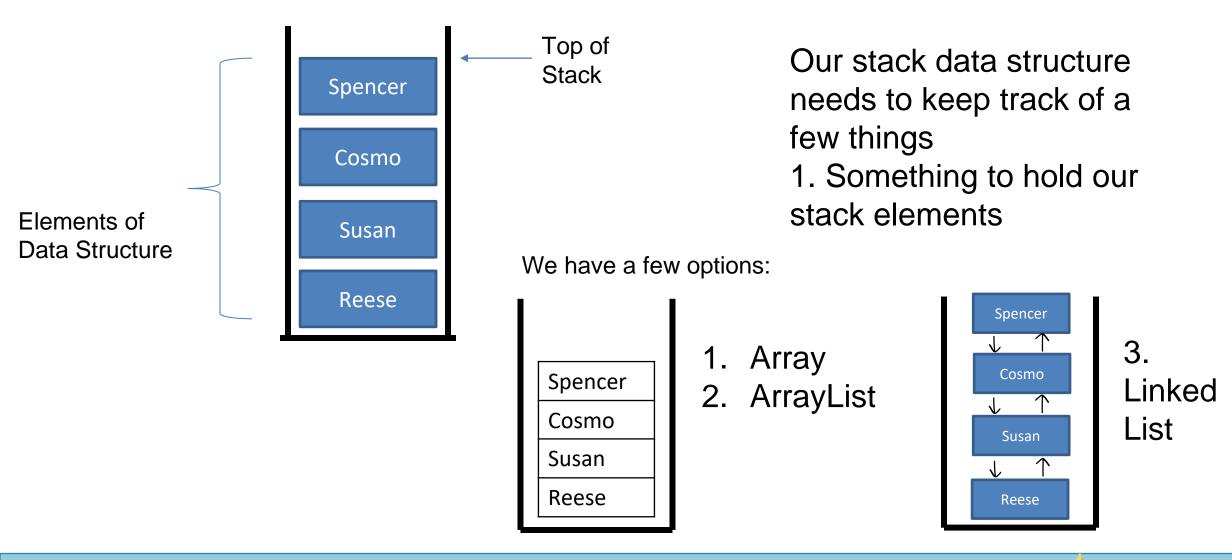




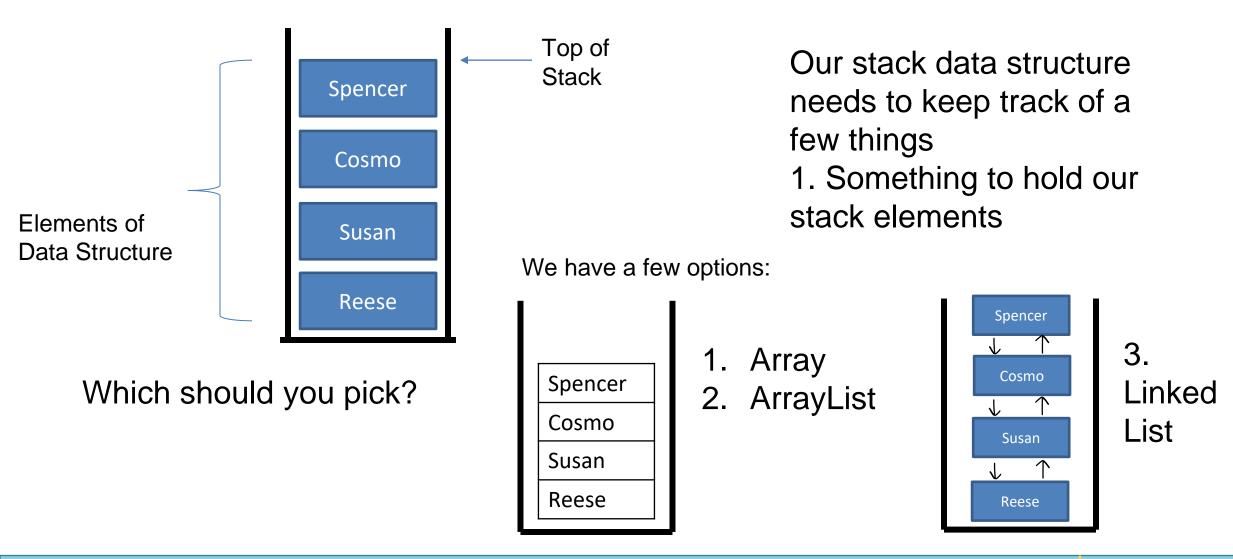




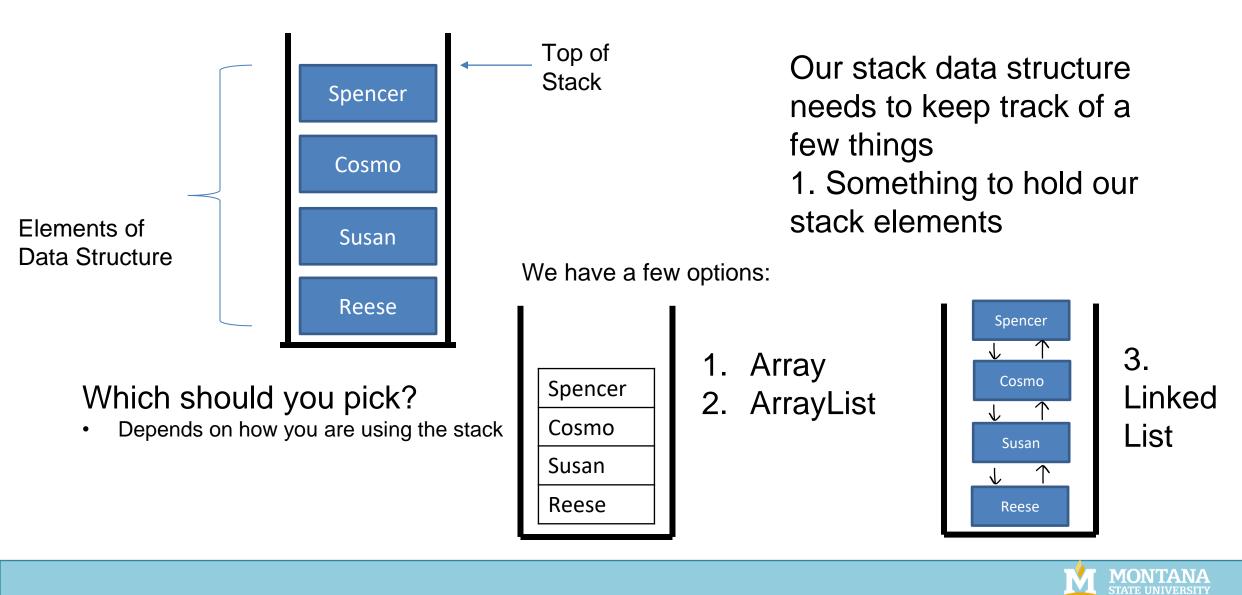


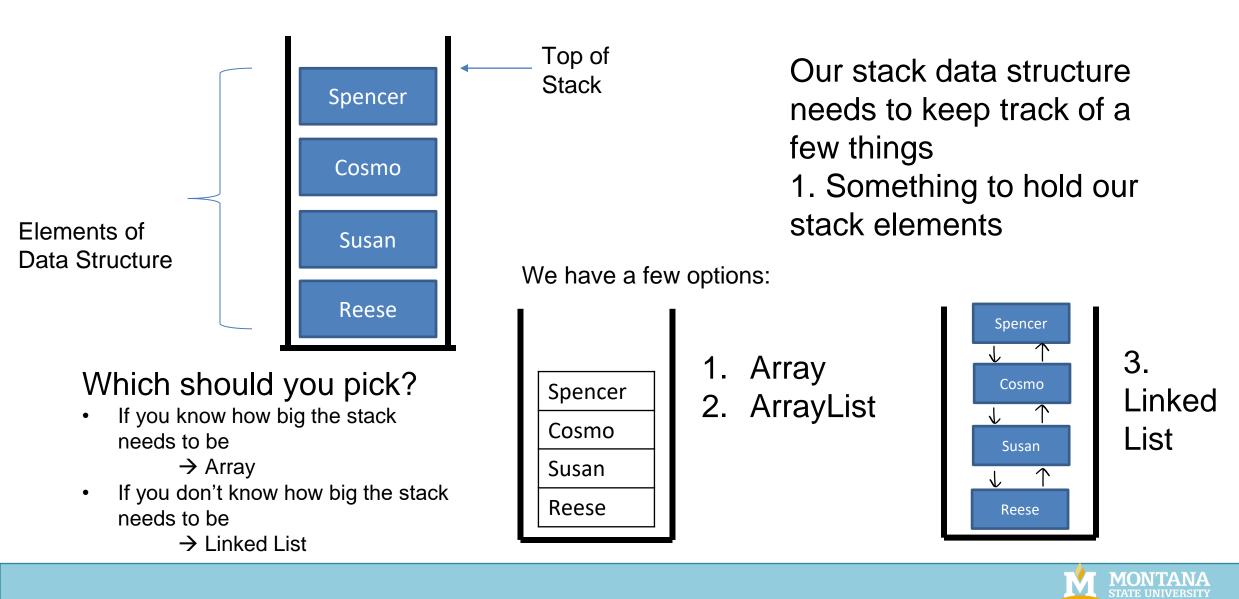


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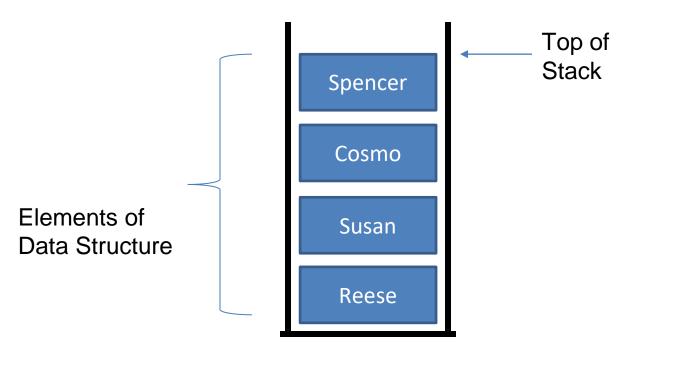


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Our stack data structure needs to keep track of a few things

- Something to hold our stack elements (Array/LinkedList)
- 2. Something that points the current top element of the stack
- 3. The size of the stack



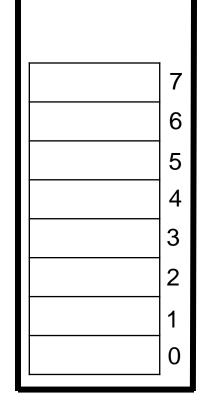
Here, we've created an array of size 8 to hold our stack data

To Do List:

- Push()
- Pop()

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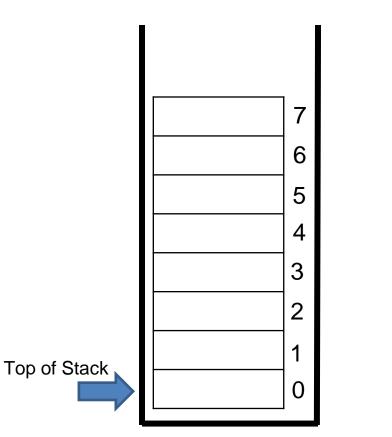
Peek()



IsEmpty()



Here, we've created an array of size 8 to hold our stack data



 IsEmpty()
 The bottom of the stack will always be at index 0, and grows towards the higher indices String[] data = new String[8]

When the stack is empty, the index of the bottom of the stack, and the index of the top of the stack will be the same

The size of the stack will start at 0 size = 0



To Do List:

• Push()

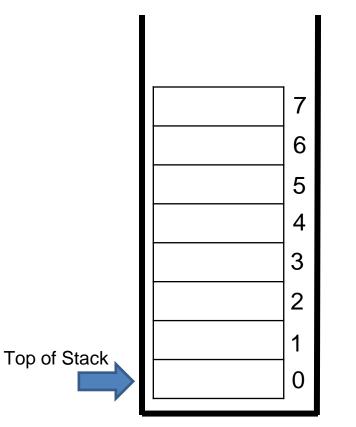
Pop()

Peek()

top_of_stack = 0

Here, we've created an array of size 8 to hold our stack data

}



public void push(newElement){



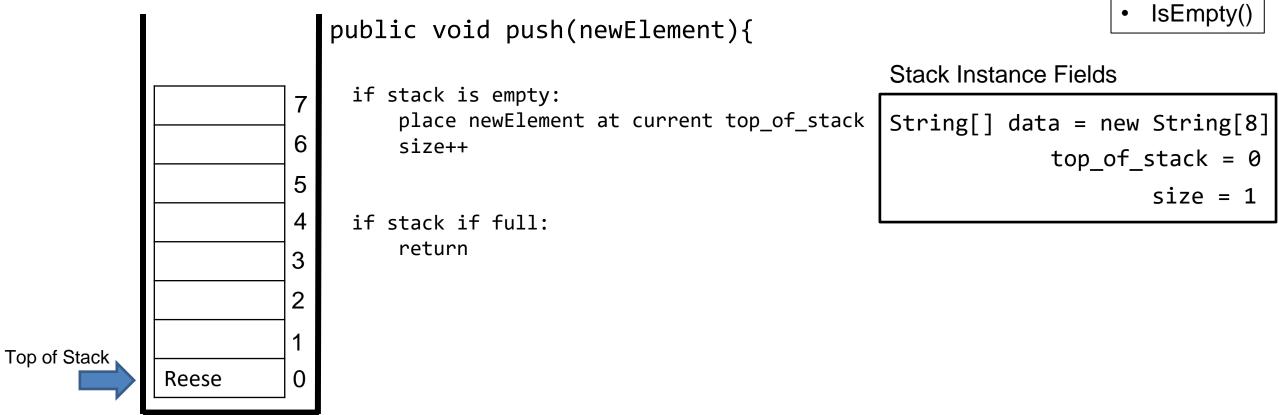
Stack Instance Fields

String[]	data = new String[8]
	top_of_stack = 0)
	size = 0	



Here, we've created an array of size 8 to hold our stack data

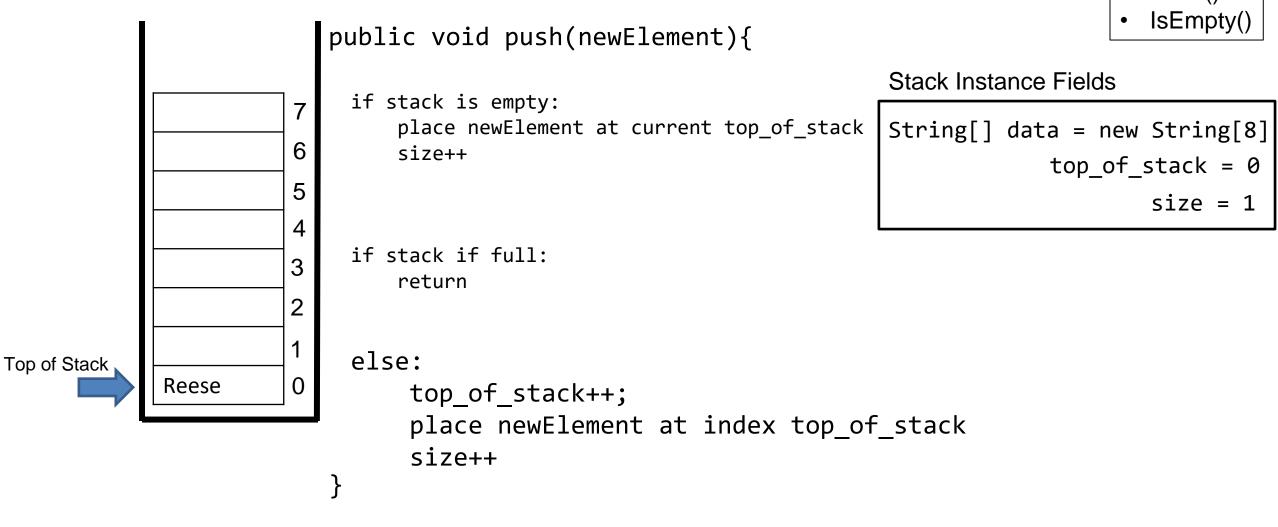
}

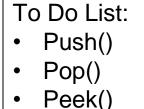


To Do List:Push()Pop()

Peek()

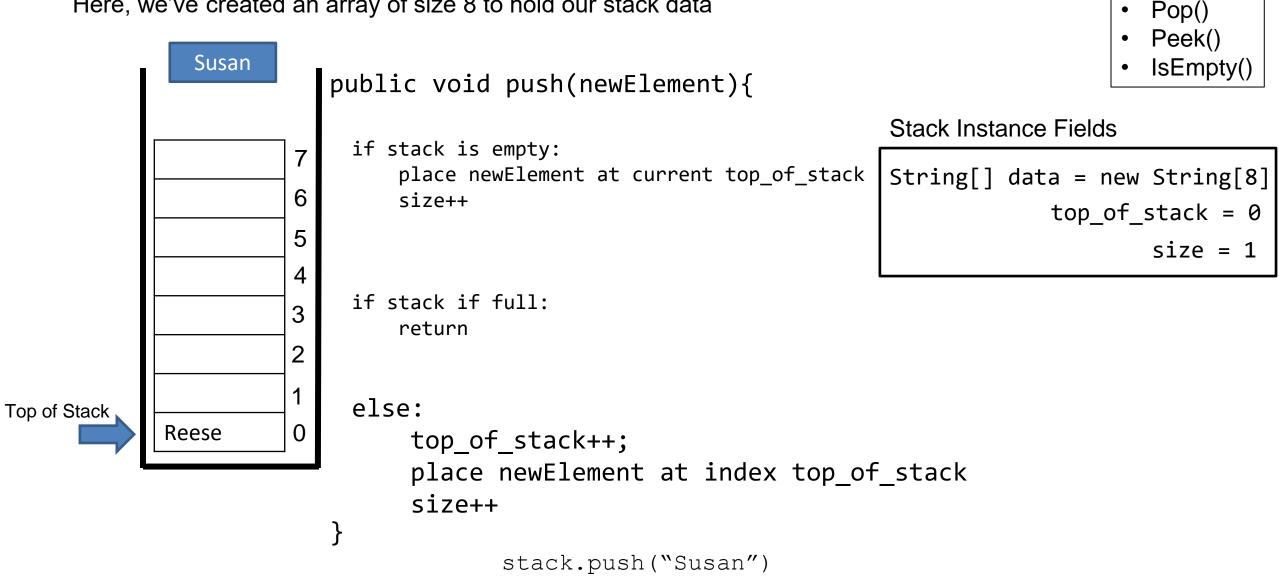
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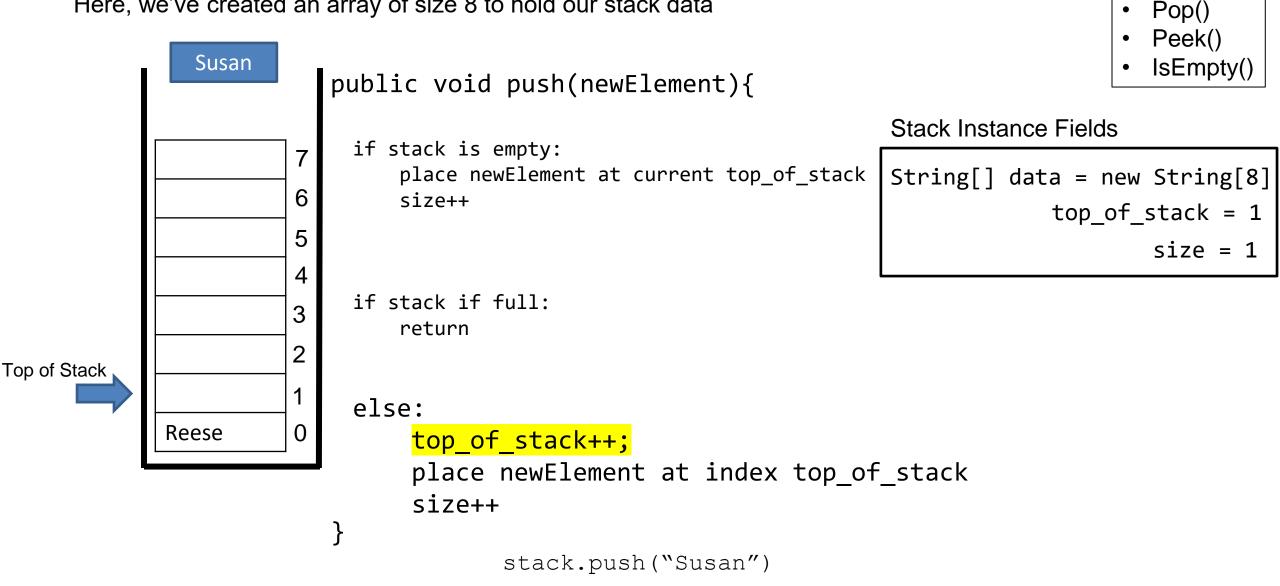




To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data

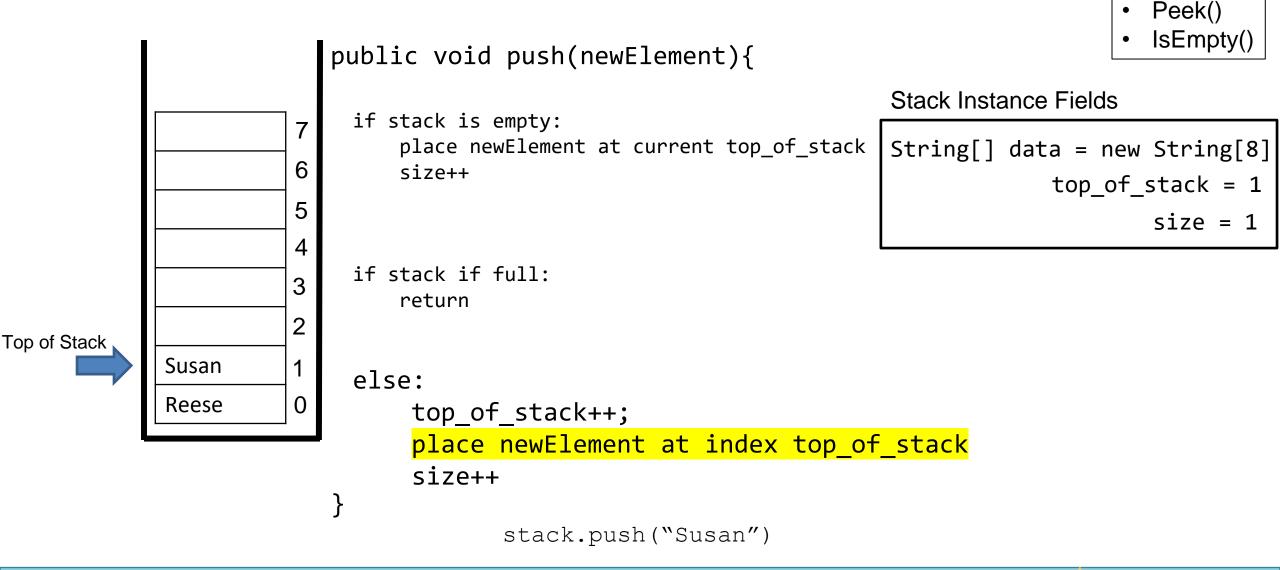




To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data



To Do List:

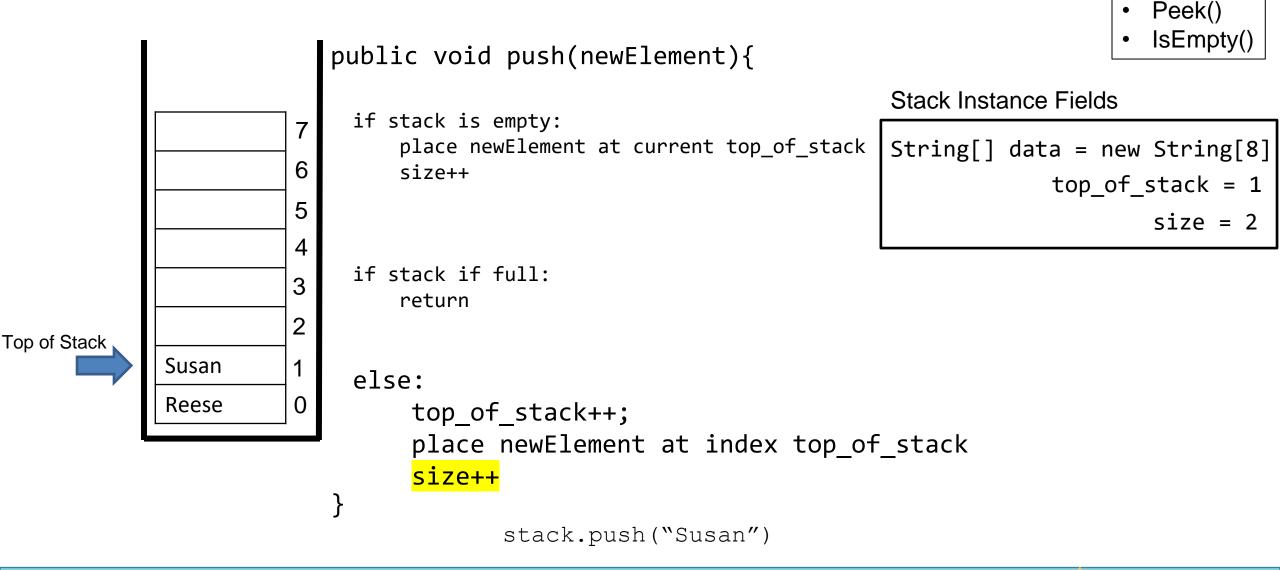
• Push()

• Pop()

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Here, we've created an array of size 8 to hold our stack data

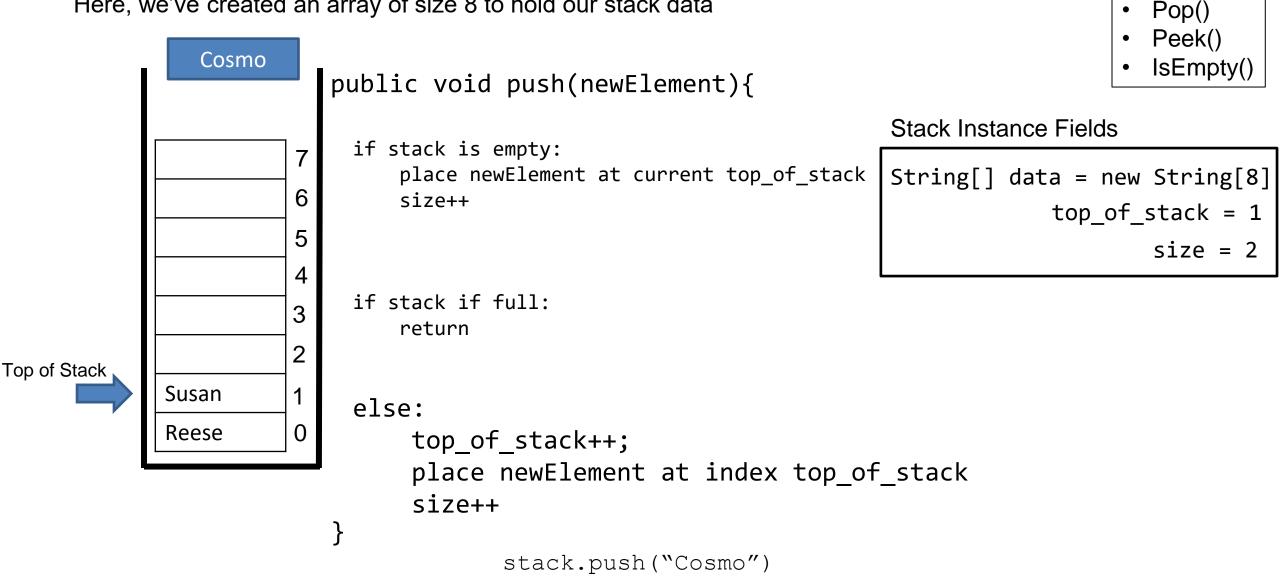


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To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data

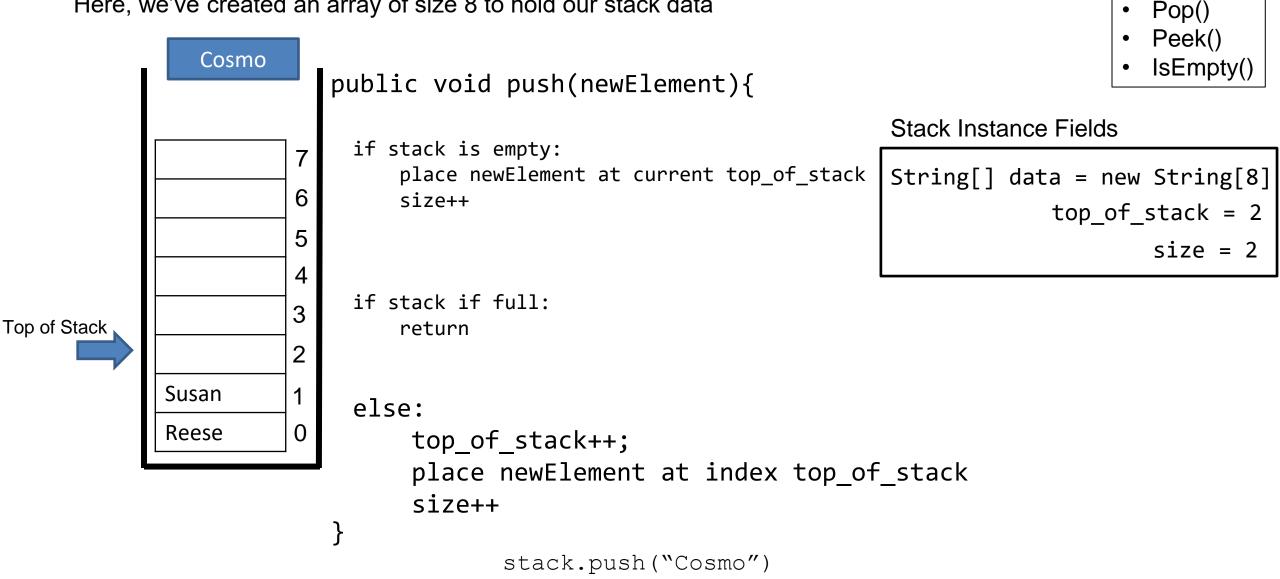




To Do List:

• Push()

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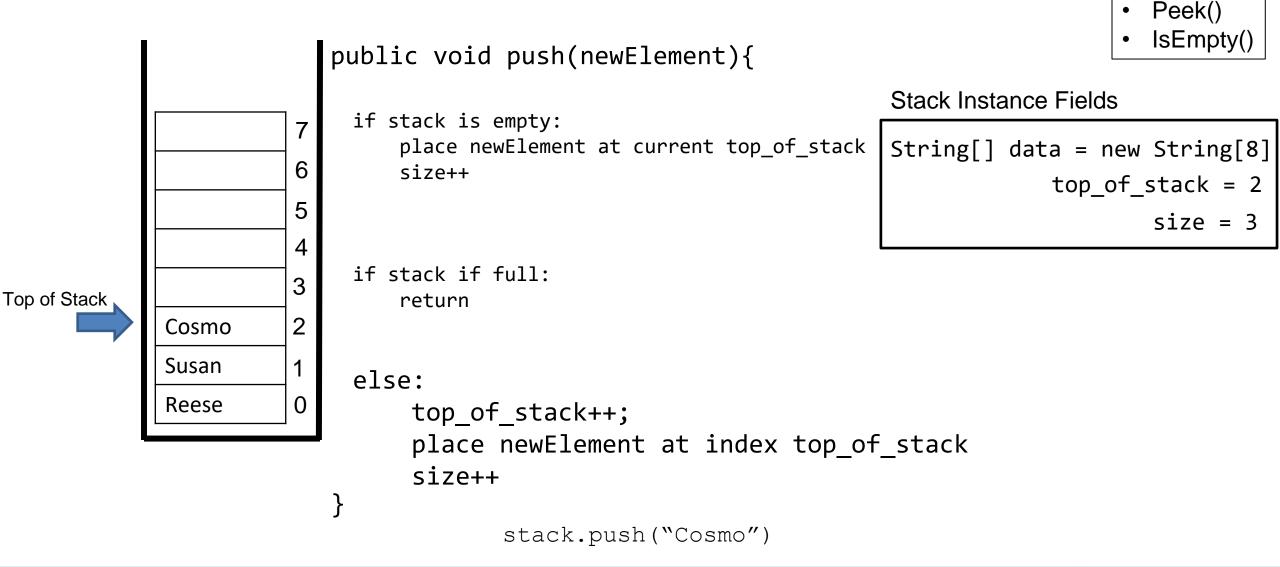




To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data



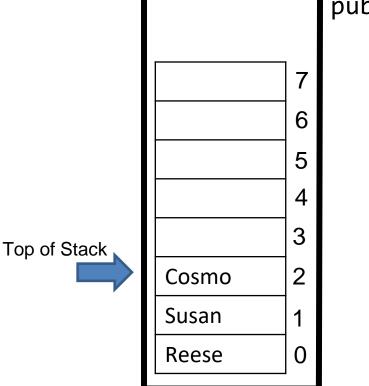


To Do List:

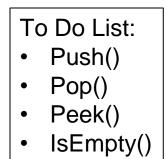
• Push()

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}



public void pop(){



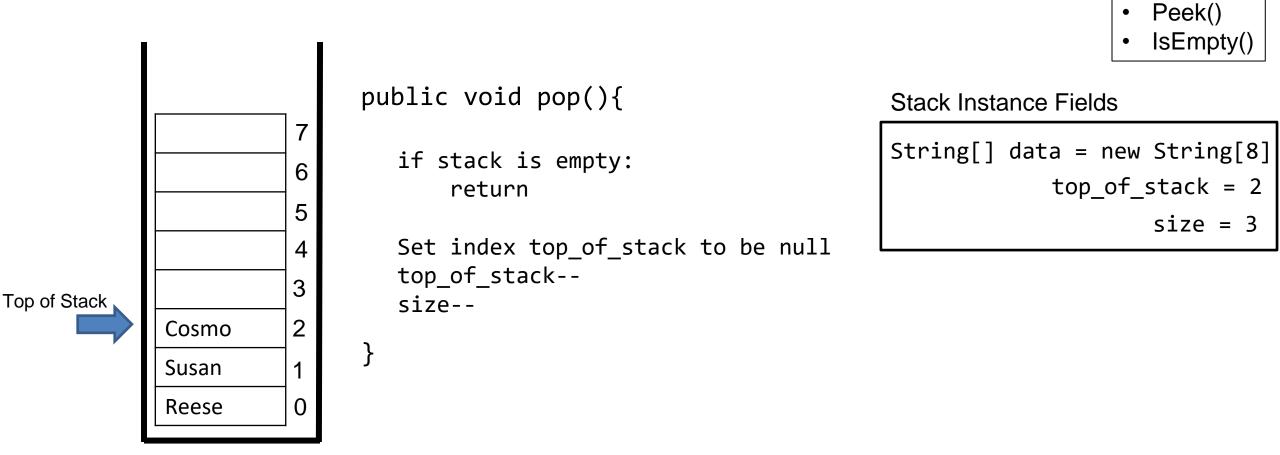
Stack Instance Fields

```
String[] data = new String[8]
    top_of_stack = 2
    size = 3
```

The pop method will always remove the element on the top of the stack



Here, we've created an array of size 8 to hold our stack data



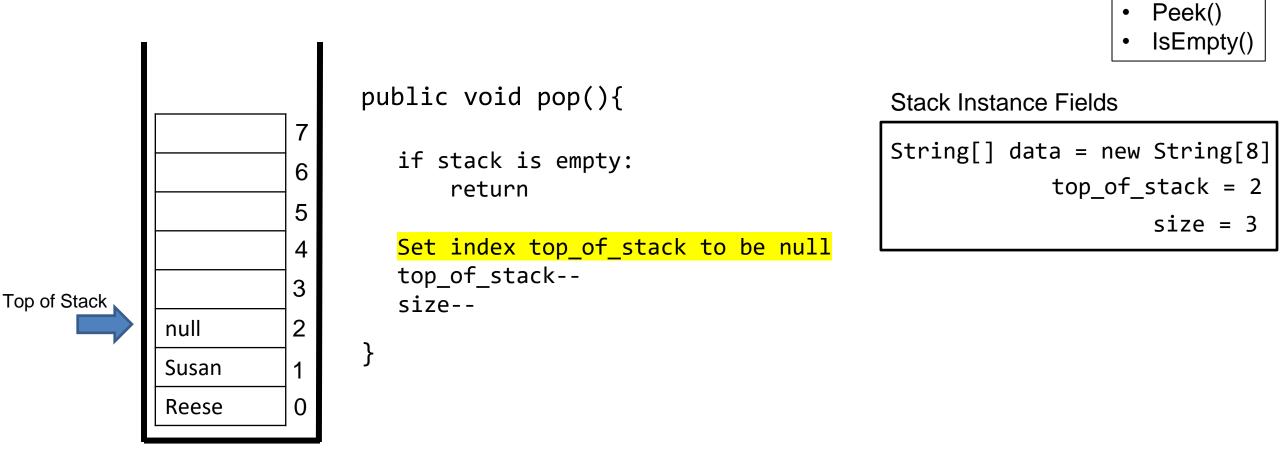
stack.pop()



To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data



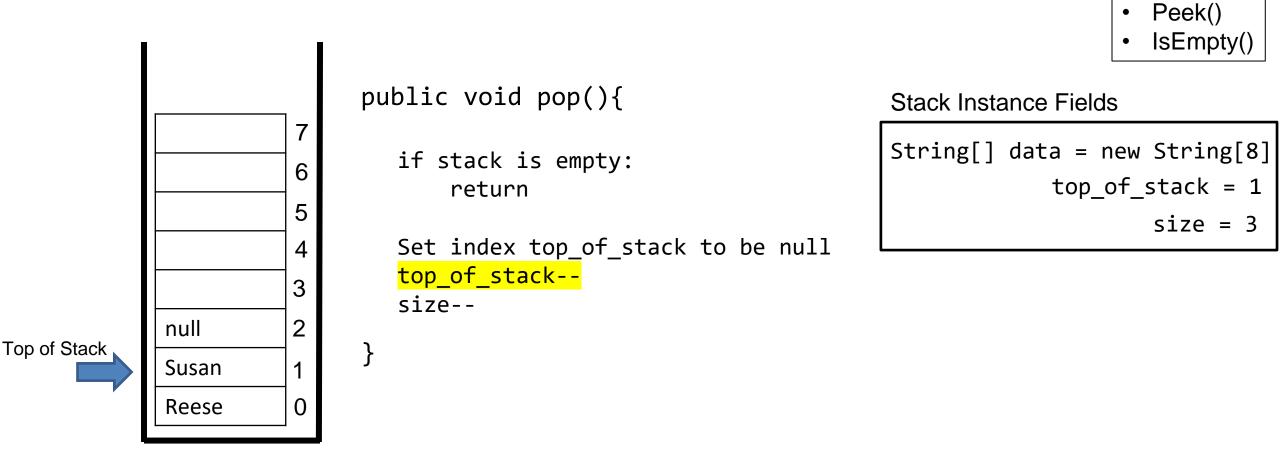
stack.pop()



To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data



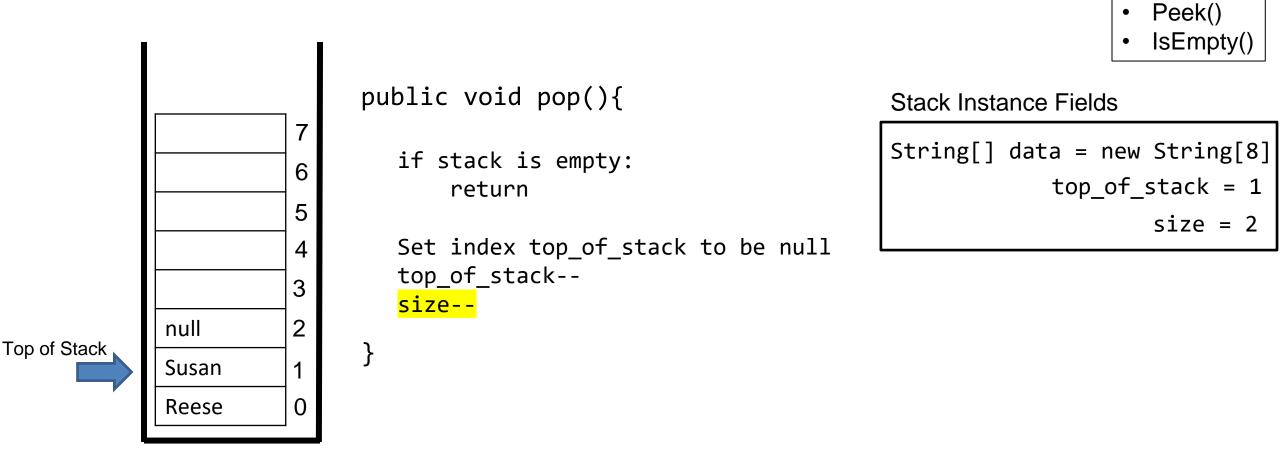
stack.pop()



To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data



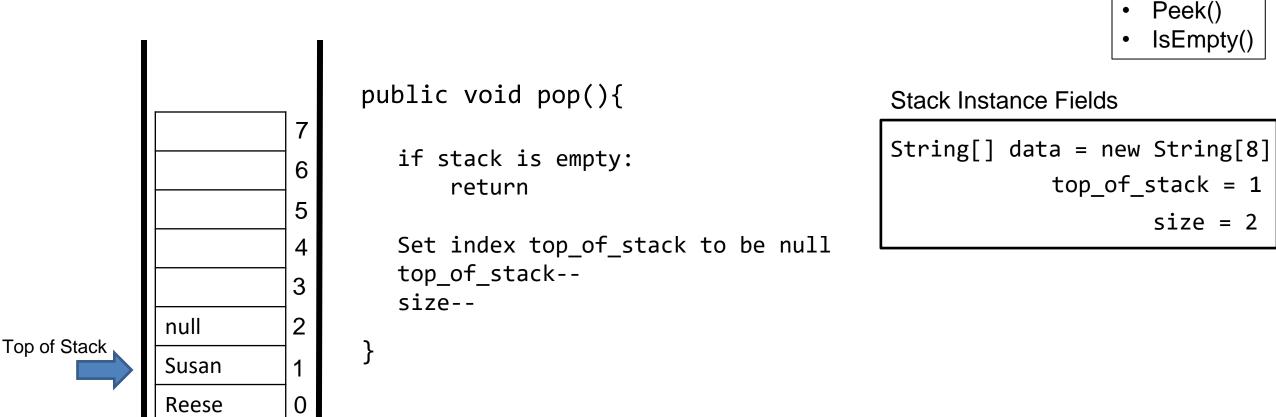
stack.pop()



To Do List:

• Push()

Here, we've created an array of size 8 to hold our stack data



Note: This method does not return the element that was removed, however there may be times where the pop() method returns the element that got removed



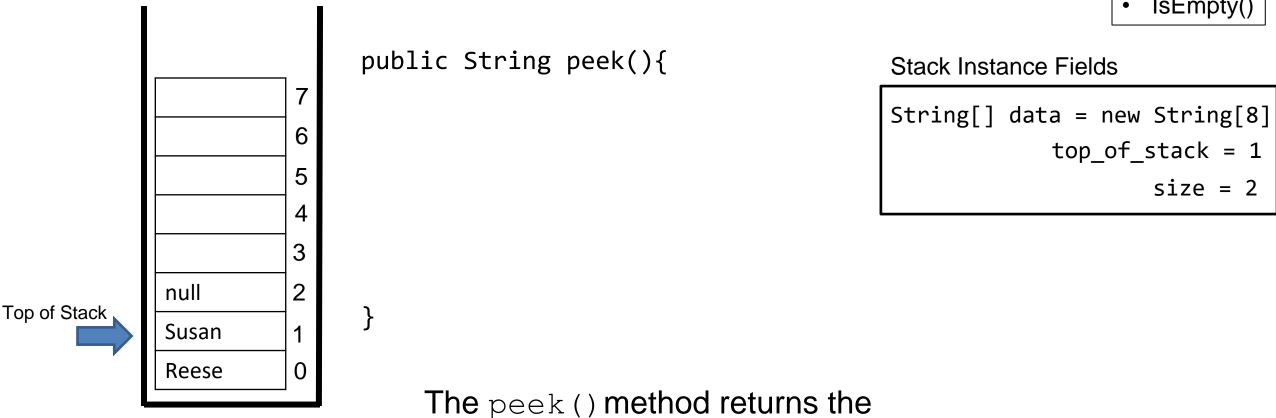
To Do List:

• Push()

Pop()

Here, we've created an array of size 8 to hold our stack data

To Do List:
Push()
Pop()
Peek()
IsEmpty()

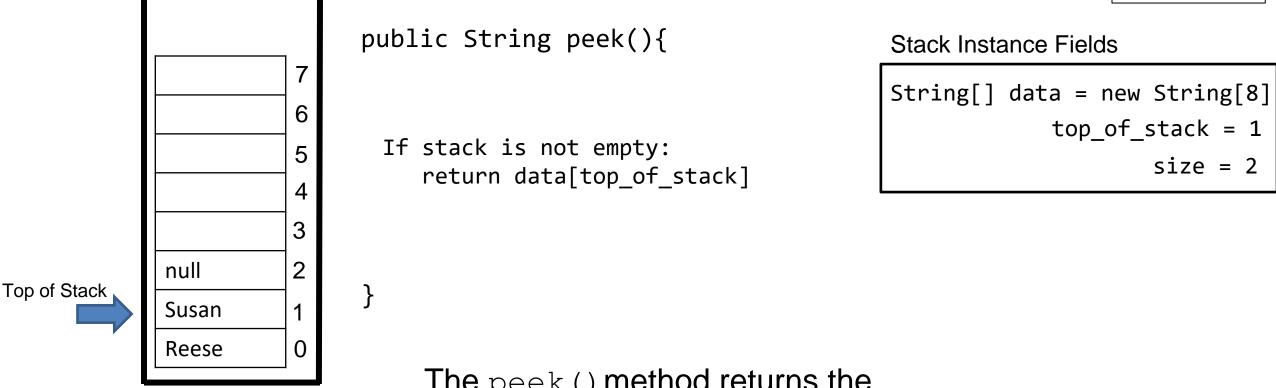


The peek () method returns the element that is currently on the top of the stack



Here, we've created an array of size 8 to hold our stack data

To Do List: • Push() • Pop() • Peek() • IsEmpty()

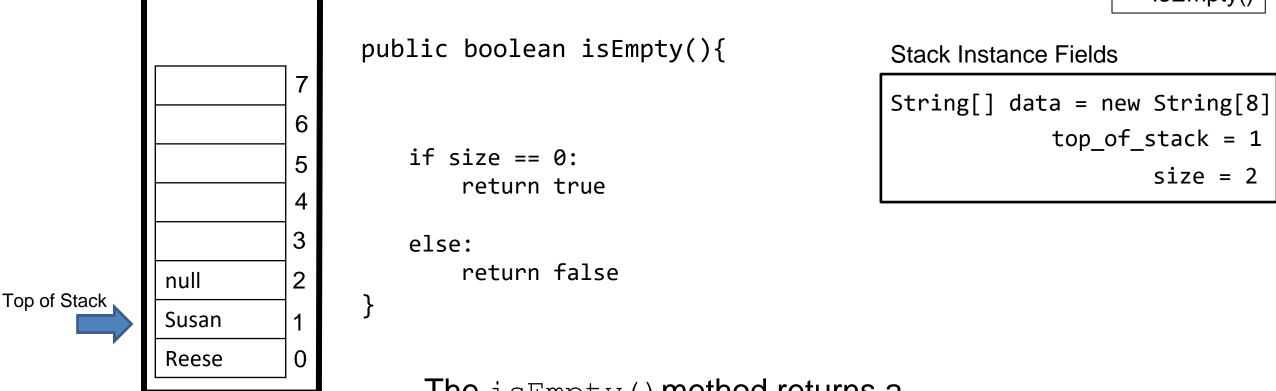


The peek () method returns the element that is currently on the top of the stack



Here, we've created an array of size 8 to hold our stack data

To Do List:
Push()
Pop()
Peek()
IsEmpty()



The isEmpty() method returns a boolean: true if the stack is empty, false if the stack is not empty

