CSCI 132: Basic Data Structures and Algorithms

Recursion (Part 3)

Reese Pearsall

Fall 2023

https://www.cs.montana.edu/pearsall/classes/fall2023/132/main.html



Announcements

Lab 10 due tomorrow

Program 4 due next Friday

No class on Friday (Veterans Day)



2

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?



3

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

D = [1, 5, 10, 25] K = 37



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D = [1, 5, 10, 25] K = 37

Answer = 4

(Quarter, dime, two pennies)



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(Quarter, dime, two pennies)

Use as many quarters as possible, then as many dimes as possible, ...



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

D = [1, 5, 10, 25] K = 37

Answer = 4

(Quarter, dime, two pennies)

Use as many quarters as possible, then as many dimes as possible, ...

This is known as the **greedy** approach



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 25]

K = 37

Use as many quarters as possible, then as many dimes as possible, ...



9

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 18, 25] Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ... K = 37

What if there were also an 18-cent coin?



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 18, 25] Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ... K = 37

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?



K = 37

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 18, 25] Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ...

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?

Real Answer = 18, 18, 1 (3 coins)



K = 37

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

```
D = [1, 5, 10, 18, 25]
Use as many quarters as possible, then as many 18
cent pieces as possible, then dimes , ...
```

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?

```
Real Answer = 18, 18, 1 (3 coins)
```

Lesson Learned: The Greedy approach works for the United States denominations, but not for a general set of denominations



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

25 + 25 + 10 + 1 + 1 + 1 = 63



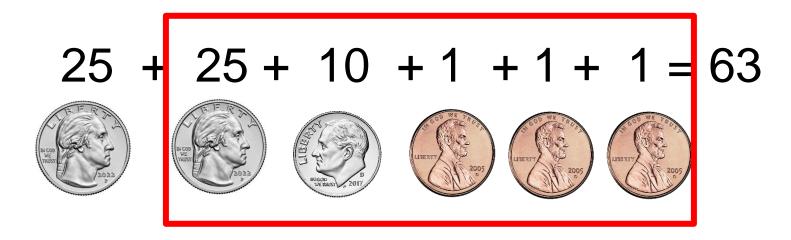
What can you conclude?

Does this provide an answer to any other change making problems?



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

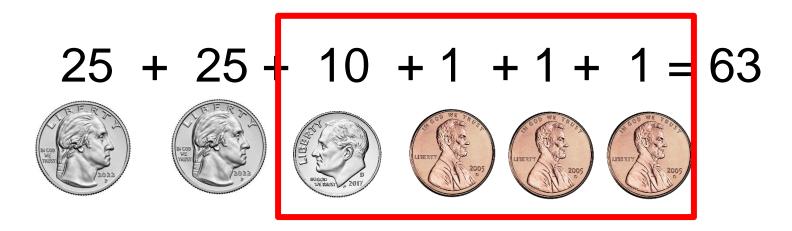


This is the minimum coins needed to make 38 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

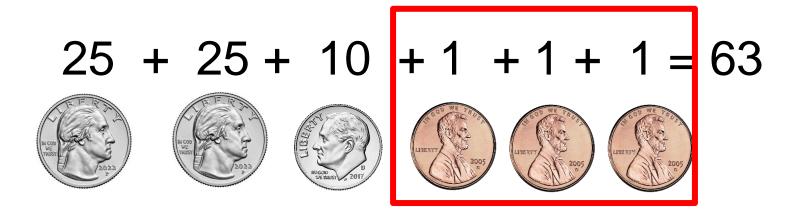


This is the minimum coins needed to make 13 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

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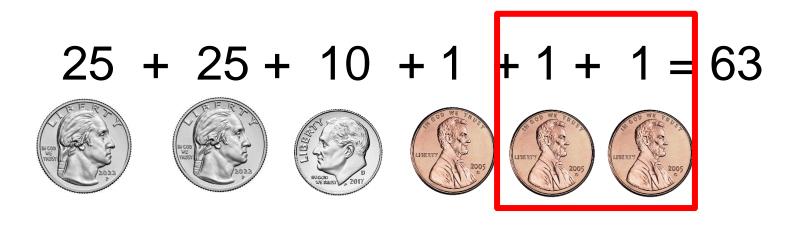


This is the minimum coins needed to make 3 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)



This is the minimum coins needed to make 2 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)



This is the minimum coins needed to make 1 cent



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

25 + 25 + 10 + 1 + 1 + 1 = 63



The solution to the change making problems consists of solutions to smaller change making problems

We can use **recursion** to solve this problem



In general, suppose a country has coins with denominations:

 $1 = d_1 < d_2 < \dots < d_k$ (US coins: $d_1 = 1, d_2 = 5, d_3 = 10, d_4 = 25$)

Algorithm: To make change for p cents, we are going to figure out change for every value x < p. We will build solution for p out of smaller solutions.



C(p) – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.



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$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$
Now find the minimum number of coins needed to make 12 cents



C(p) – minimum number of coins to make p cents.

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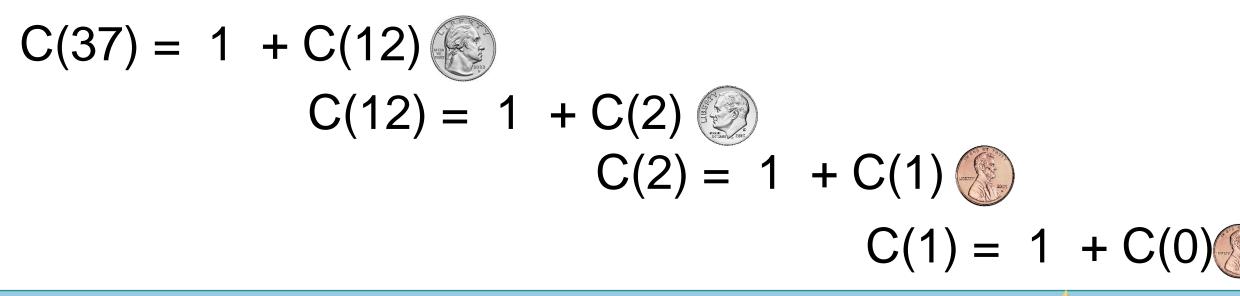
$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$
We used one dime
$$C(12) = 1 + C(2)$$

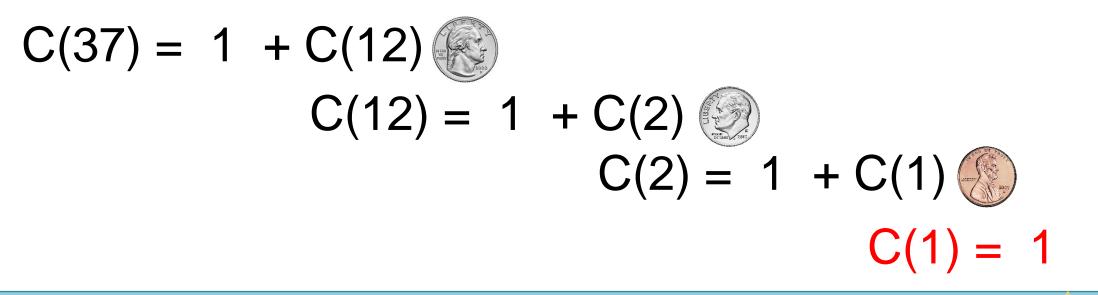
Now find the minimum number of coins needed to make 2 cents



C(p) – minimum number of coins to make p cents.

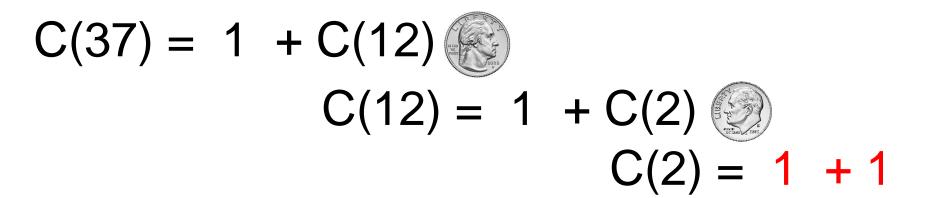


C(p) – minimum number of coins to make p cents.



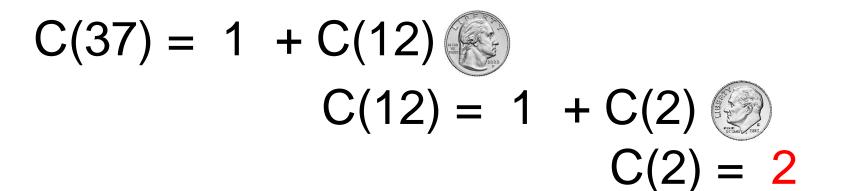


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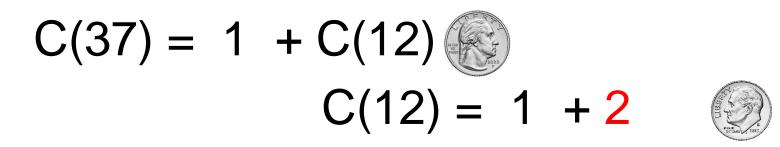


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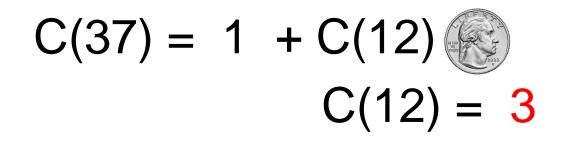


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C(37) = 1 + 3





C(p) – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution. C(p) = 1 + C(p - x).

C(37) = 4

The minimum number of coins needed to make 37 cents is 4



In general, suppose a country has coins with denominations:

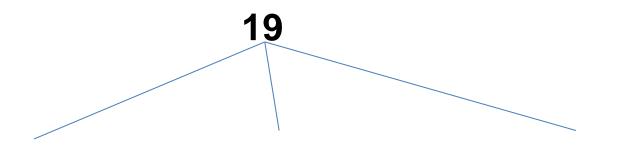
 $1 = d_1 < d_2 < \dots < d_k \qquad (\text{US coins: } d_1 = 1, d_2 = 5, d_3 = 10, d_4 = 25)$

(This algorithm must work for ALL denominations)

Algorithm: To make change for p cents, we are going to figure out change for every value x < p. We will build solution for p out of smaller solutions.

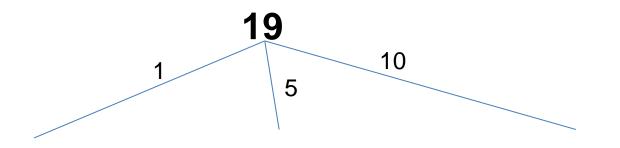


Make \$0.19 with \$0.01, \$0.05, \$0.10



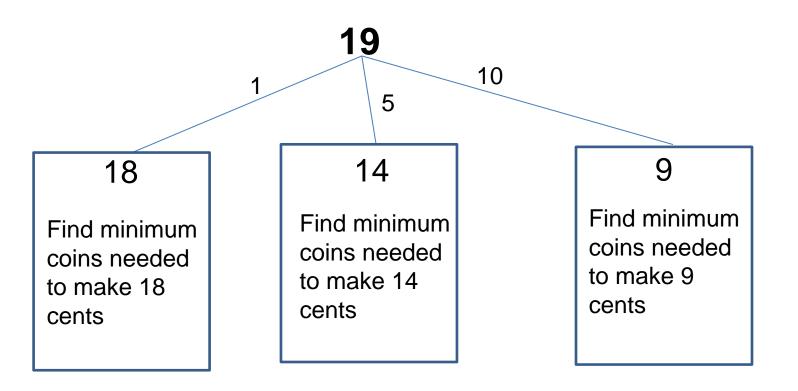


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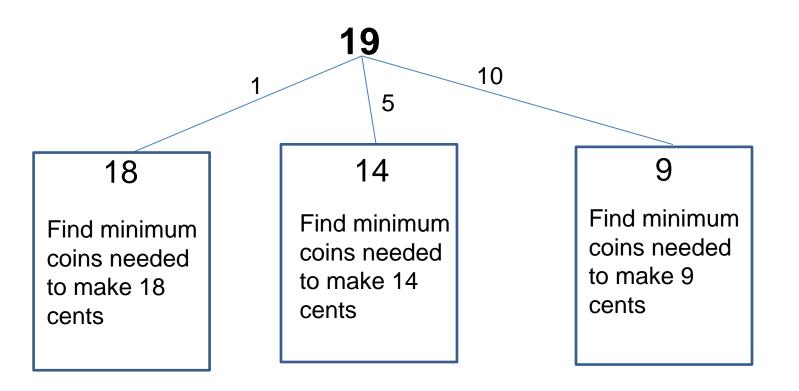
Make \$0.19 with \$0.01, \$0.05, \$0.10



To find the minimum number of coins needed to create 19 cents, we generate **k** subproblems



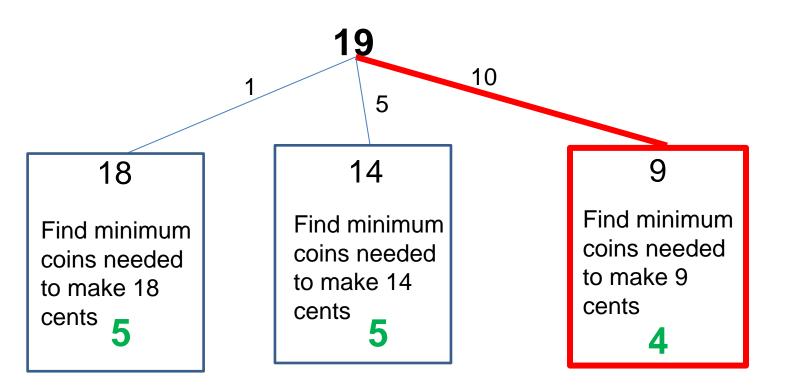
Make \$0.19 with \$0.01, \$0.05, \$0.10



We want to select the **minimum** solution of these three subproblems



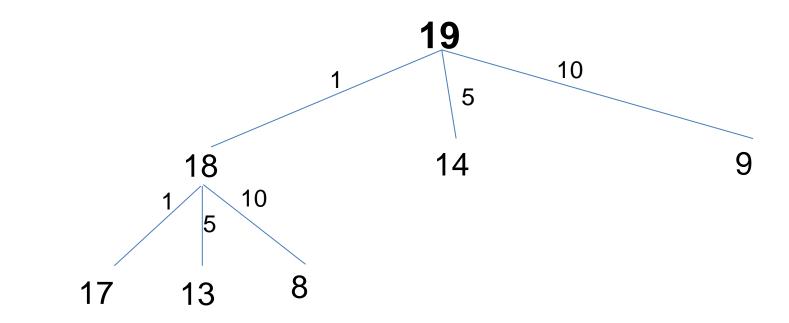
Make \$0.19 with \$0.01, \$0.05, \$0.10 k = # denominations



For the solution of our original problem (19), we want to select this branch (one dime used)



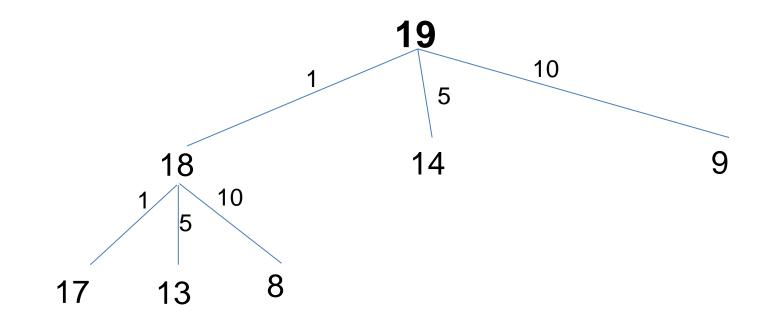
Make \$0.19 with \$0.01, \$0.05, \$0.10



| Find minimum | Find minimum | Find minimum |
|--------------|--------------|--------------|
| coins needed | coins needed | coins needed |
| to make 17 | to make 13 | to make 8 |
| cents | cents | cents |

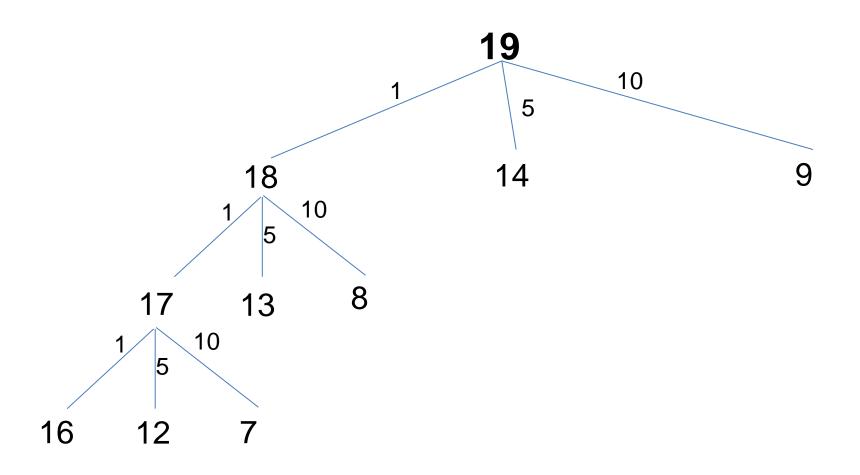


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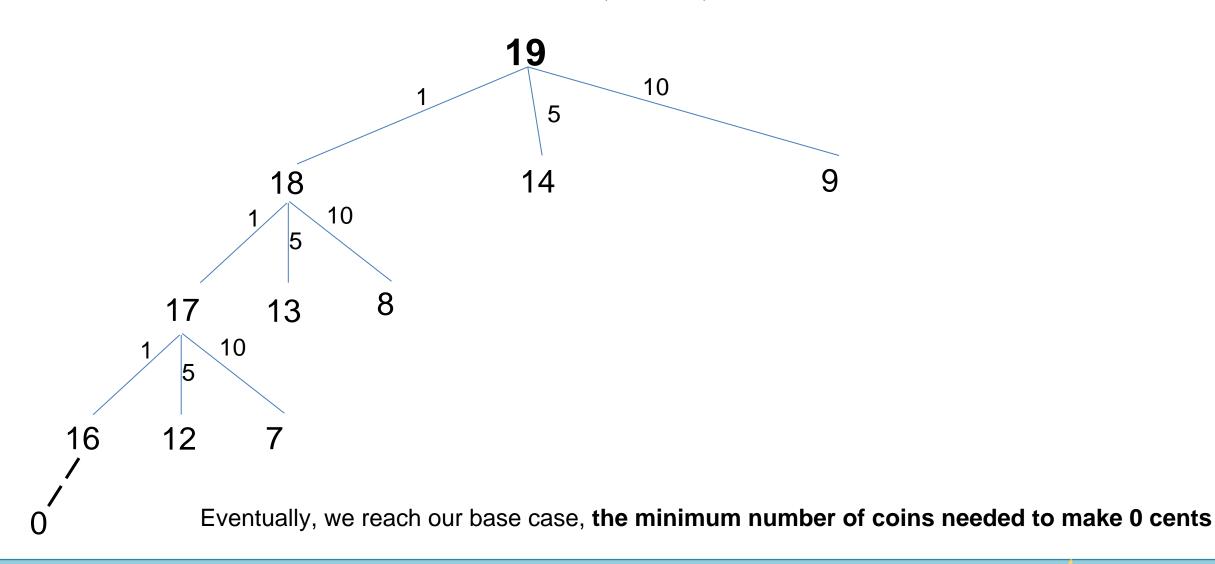


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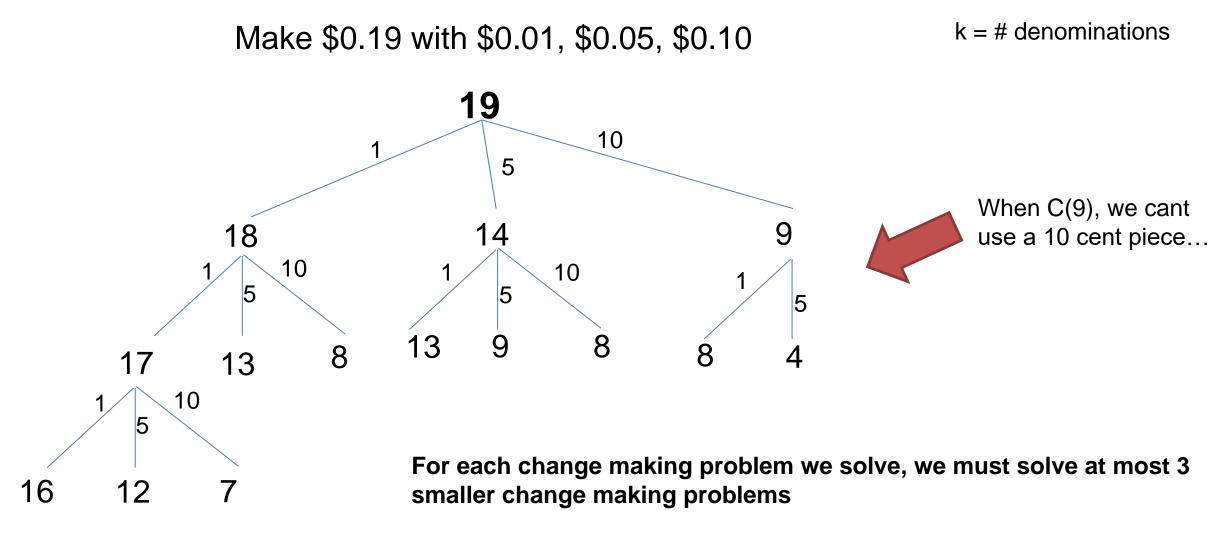




Make \$0.19 with \$0.01, \$0.05, \$0.10

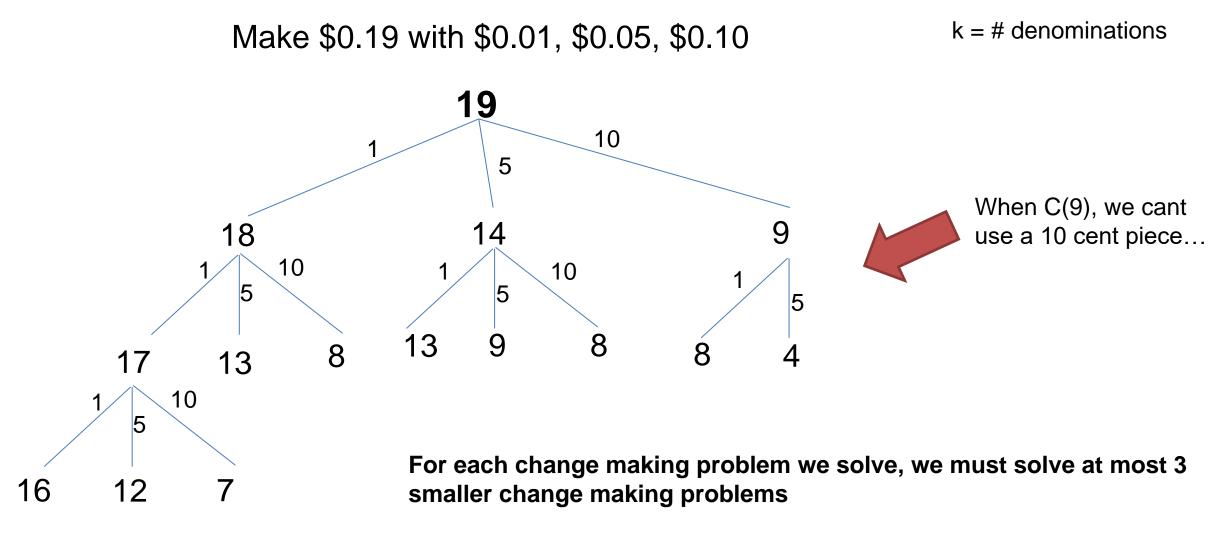






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$$C(p) = \begin{cases} \min_{i:d_i \le p} C(p - d_i) + 1, p > 0\\ 0, p = 0 \end{cases}$$

Least change for 19 cents = minimum of:

- least change for 19-10 = 9 cents
- least change for 19-5 = 14 cents
- least change for 19-1 = 18 cents

For each problem P, we will solve the problem for (P - d), where d represents each possible denomination



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D = array of denominations [1, 5, 10, 18, 25]p = desired change (37)

min_coins(D, p)



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if p == 0
 return 0;

Base Case



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min coins(D, p) if p == 0Base Case return 0; else min = 🕶 int min = Integer.MAX_VALUE; int a = Integer.MAX_VALUE;; a = 🚥 for each d_i in D $if (p - d_i) >= 0$

 $a = min_coins(D, p - d_i)$

Recurse, and find the minimum number of coins needed using each valid denomination



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min coins(D, p) if p == 0Base Case return 0; else min = 🕶 int min = Integer.MAX_VALUE; int a = Integer.MAX_VALUE;; a = ∞ for each d; in D Recurse, and find the minimum number of coins if $(p - d_i) >= 0$ needed using each valid $a = min coins(D, p - d_i)$ denomination if a < min Select the branch that has min = athe minimum value

MONTANA 54

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| min_coins(D, p) | |
|-----------------|---|
| if p == 0 | Base Case |
| return 0; | |
| else | <pre>int min = Integer.MAX_VALUE; int a = Integer.MAX_VALUE;;</pre> |
| min = ∞ | |
| a = ∞ | |

for each
$$d_i$$
 in D
if $(p - d_i) \ge 0$
 $a = \min_{i=1}^{1} coins(D, p - d_i)$
if $a < \min_{i=1}^{1} a$

Recurse, and find the minimum number of coins needed using each valid denomination

Select the branch that has the minimum value

D = array of denominations [1, 5, 10, 18, 25] p = desired change (37)

| min_coins(D, p) | | |
|-----------------|---|--|
| if p == 0 | Base Case | |
| return 0; | | |
| else | <pre>int min = Integer.MAX_VALUE; int a = Integer.MAX_VALUE;;</pre> | |
| min = ∞ | | |
| a = ∞ | | |

Recurse, and find the minimum number of coins needed using each valid denomination

Select the branch that has the minimum value

return 1 + min

Once, our for loop finishes, we should know the branch that had the minimum, so return (1 + min), 1 because one coin was used in the current method call



 $min_coins(D, p)$ if p == 0 return 0; else min = ∞ a = ∞

return 1 + min

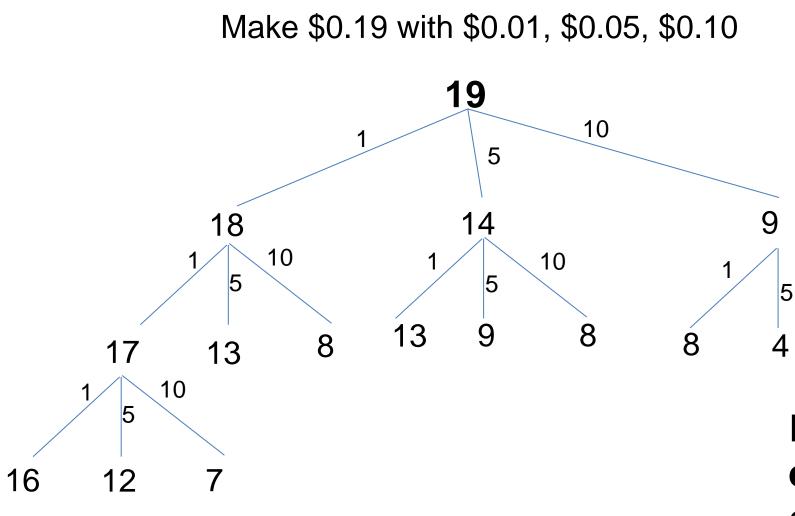


min_coins(D, p) if p == 0return 0; else min = ∞ $a = \infty$

Running time?

```
return 1 + min
```



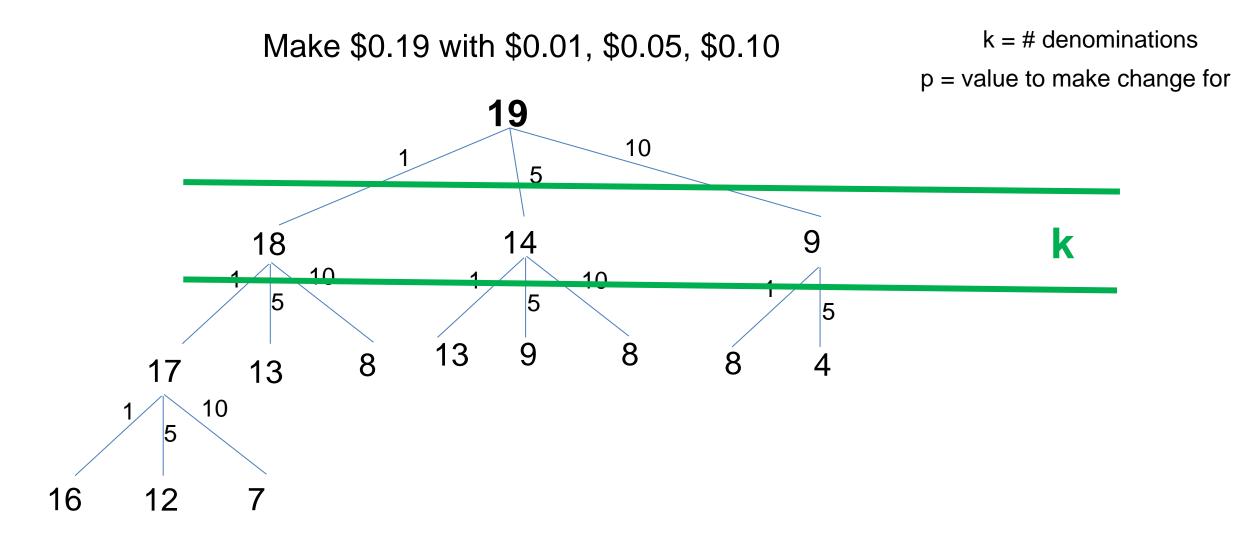


k = # denominations

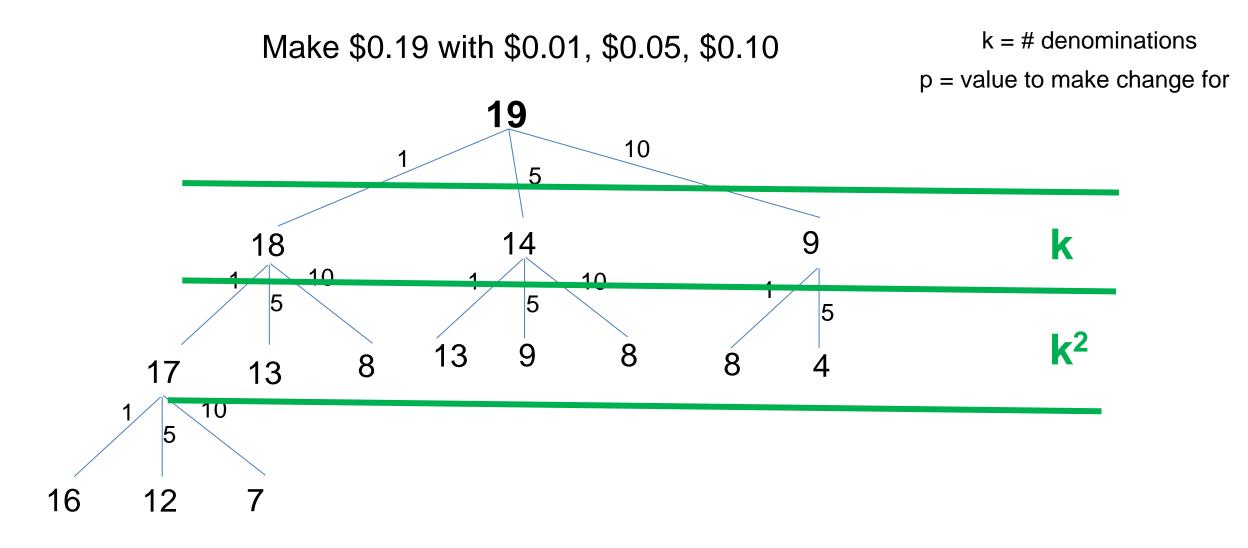
p = value to make change for

For sufficiently large p, every permutation of denominations is included.

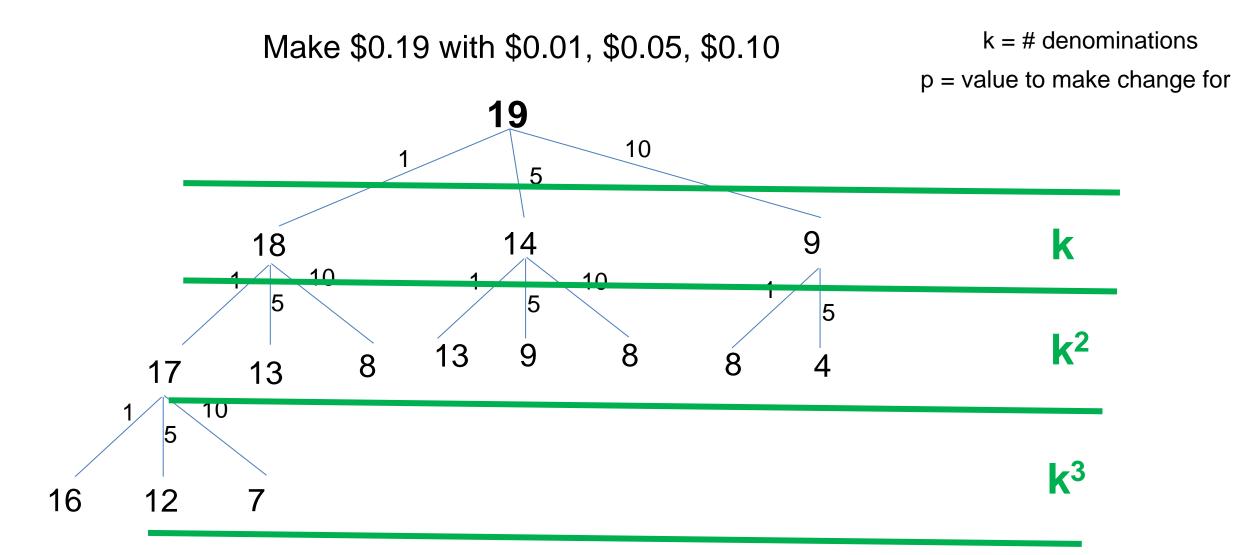




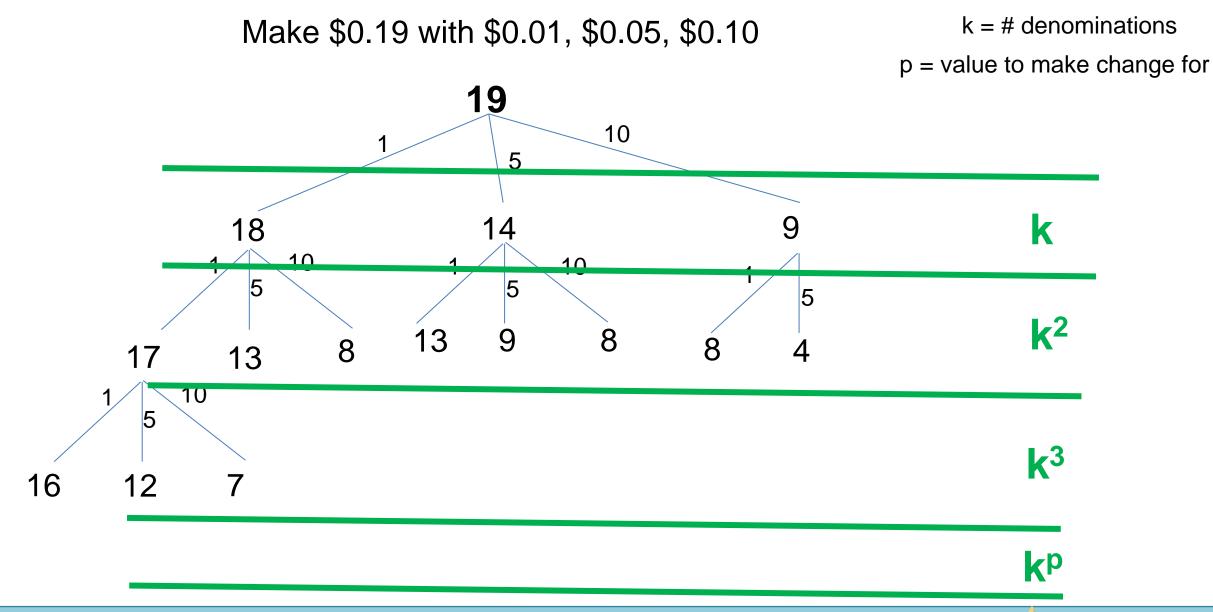














Make \$0.19 with \$0.01, \$0.05, \$0.10

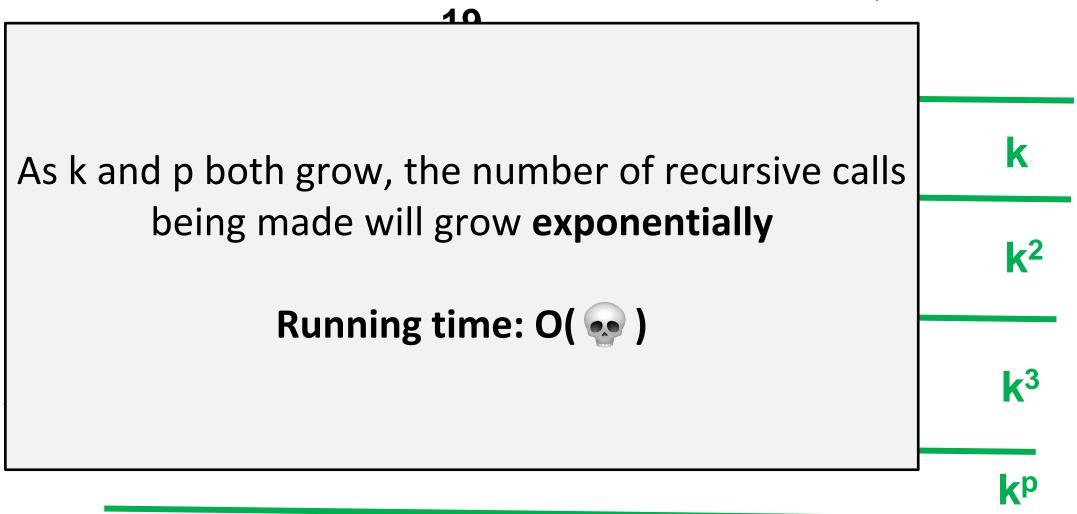
p = value to make change for

10 As k and p both grow, the number of recursive calls k being made will grow **exponentially k**² If we have a lot of coin denominations, we will have a lot of branching **k**³ kp





p = value to make change for

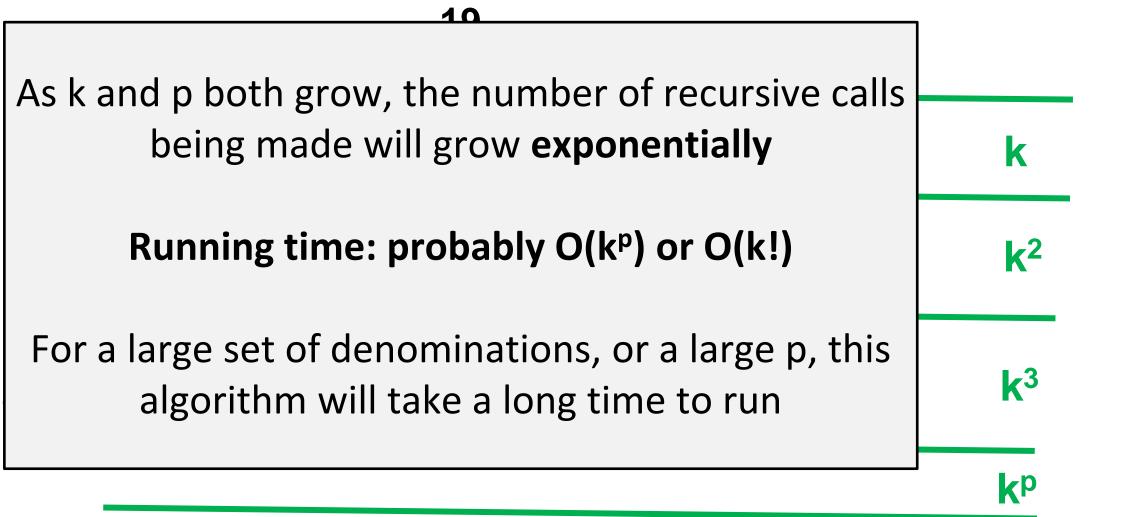




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p = value to make change for



This algorithm returns the minimum number of coins needed (ie 4), but it does not tell us what coins were used in that solution



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D = Array of coin denominations [1, 5, 10, 25]p = value to make change forn = minimum number of coins used to make p cents

Goal: Find an **n-length** combination of coins from **D** that were used to make **p**



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D = Array of coin denominations [1, 5, 10, 25]p = value to make change forn = minimum number of coins used to make p cents

Goal: Find an **n-length** combination of coins from **D** that were used to make **p**

To do this, we will compute **all n-length combinations**, but only return the combinations that add up to be **p** (not very efficient)



Generating Combinations for finding coins

Denominations (D) = [1, 5, 10]

For **n=3**, these are the combinations to be generated:

- [1, 1, 1]
- [1, 1, 5] • [1, 1, 10]
- [1, 1, 10]
 [1, 5, 5]

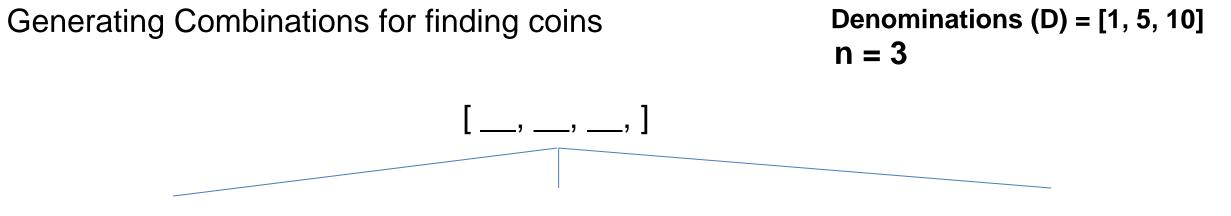
• [1, 5, 10]

- [1, 10, 10]
- [5, 5, 5]
- [5, 5, 10]
 - [5, 10, 10]
 - [10, 10, 10]

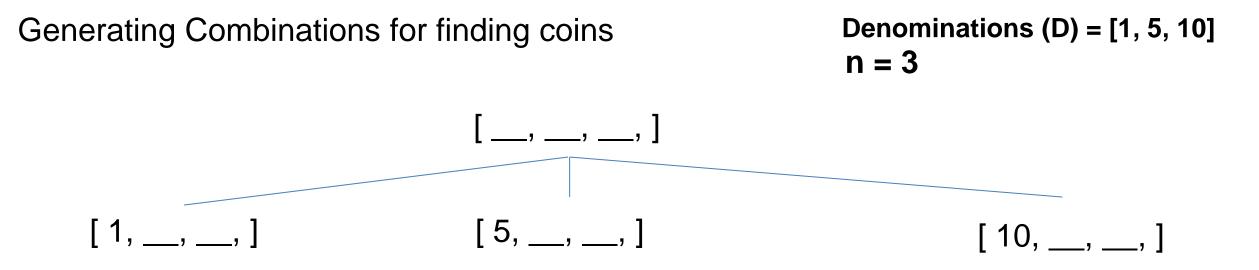
Note:

[5, 1, 1] is not a "valid" combination, because it is the same thing as [1, 1, 5]

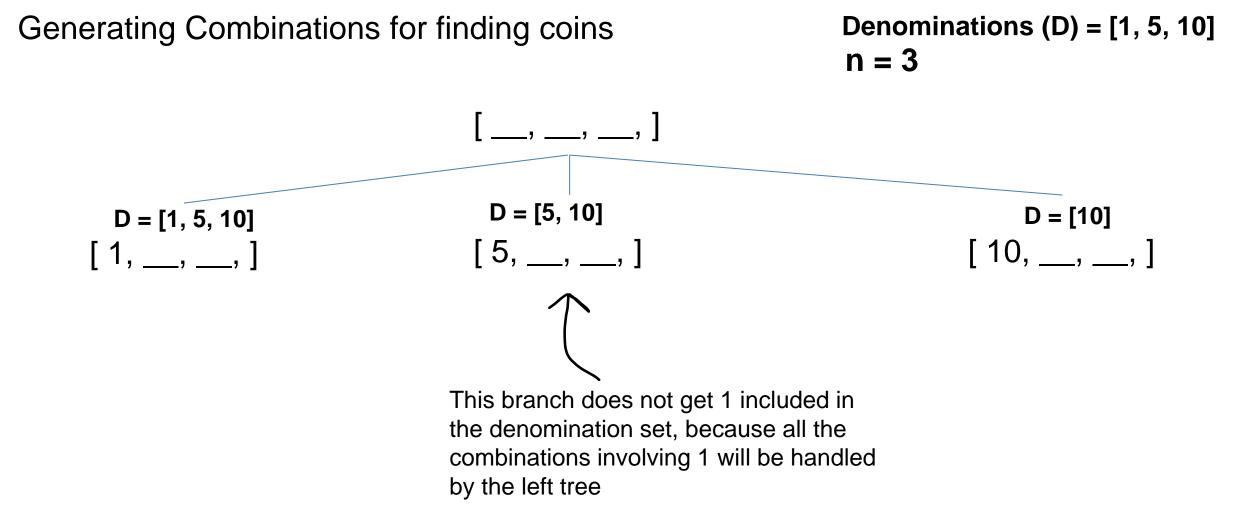




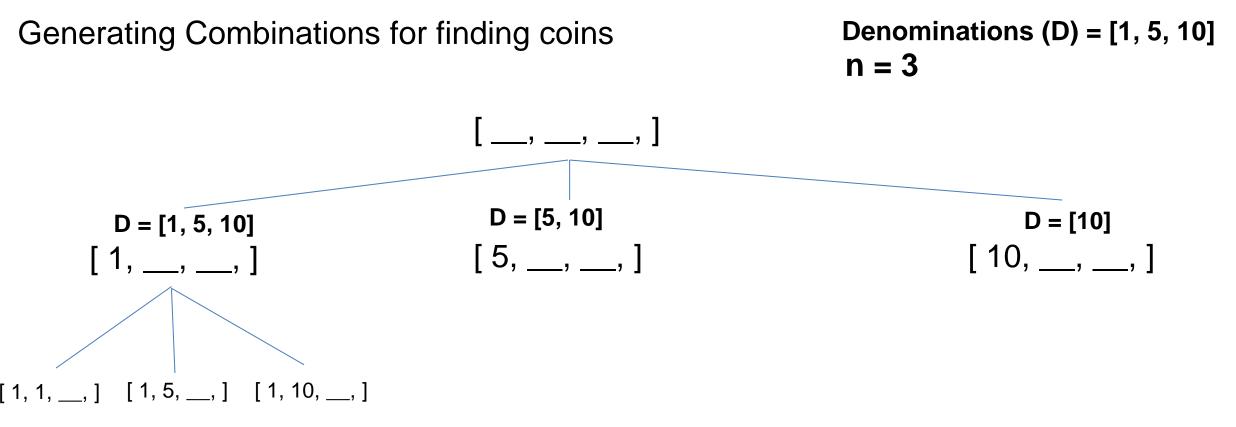




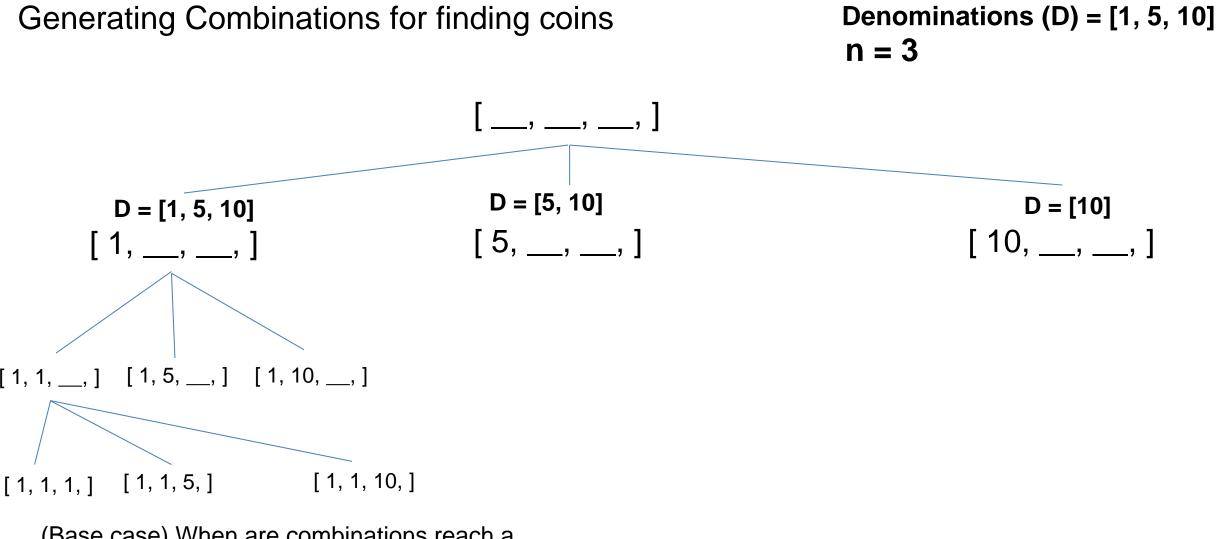








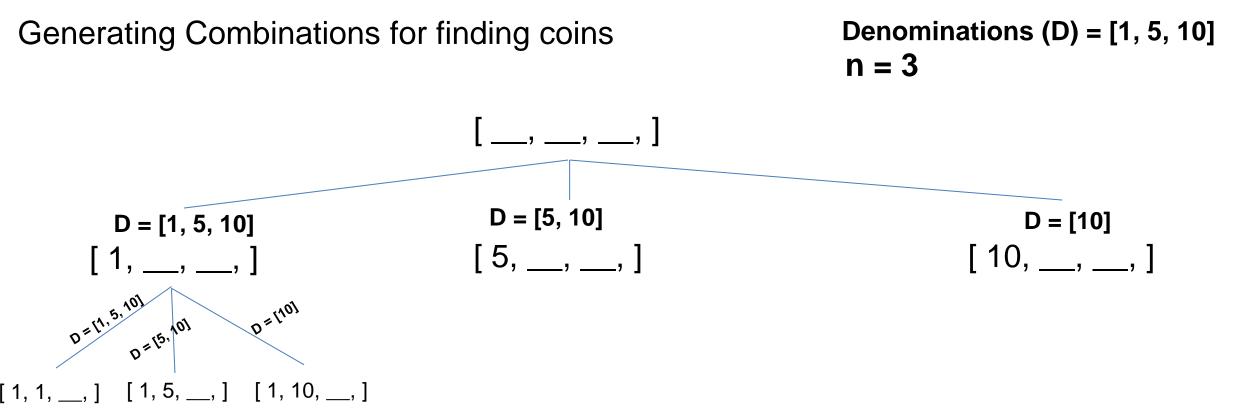




(Base case) When are combinations reach a length of 3, we will stop recursing

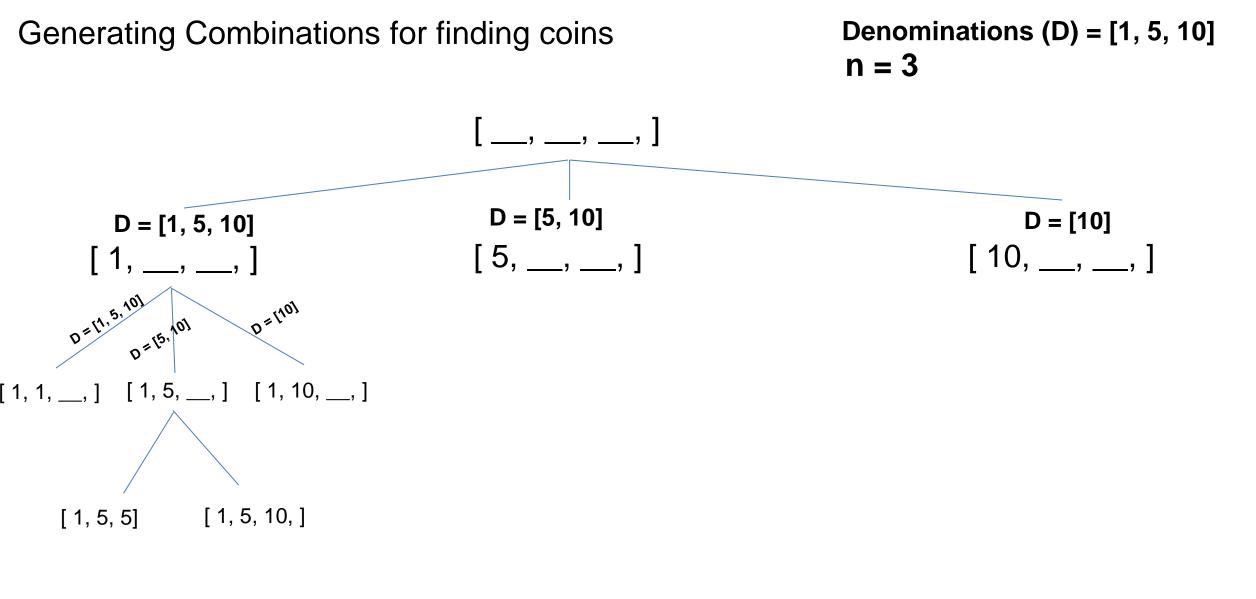
[1, 1, 1] [1, 1, 5] [1, 1, 10]





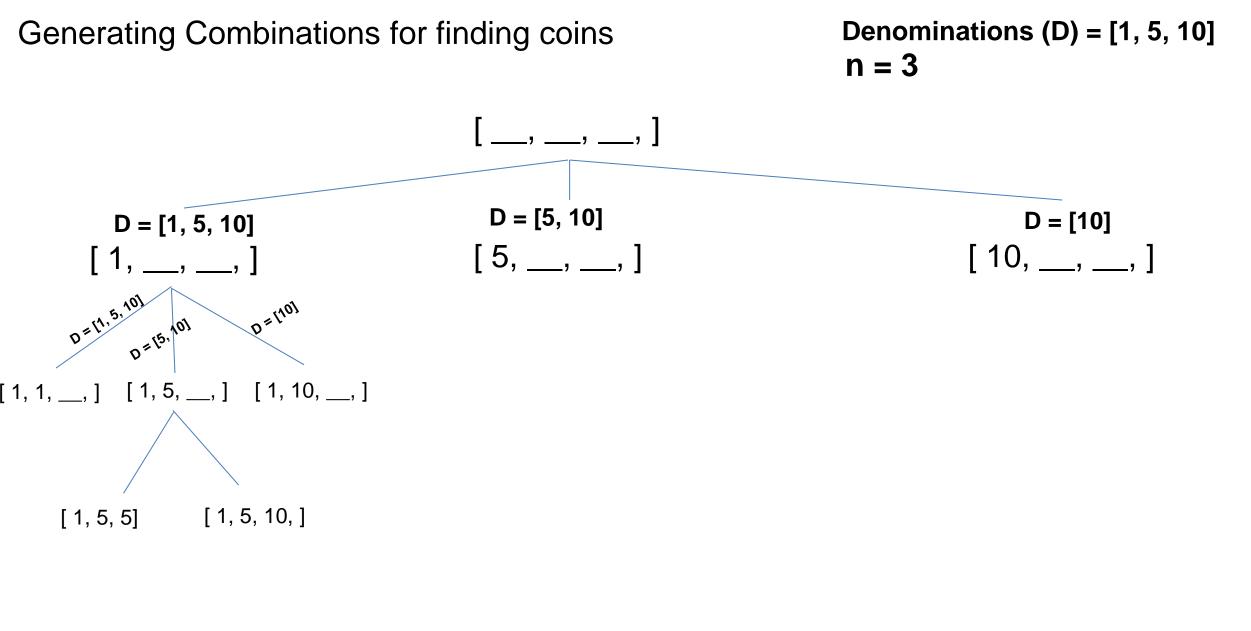
[1, 1, 1] [1, 1, 5] [1, 1, 10]





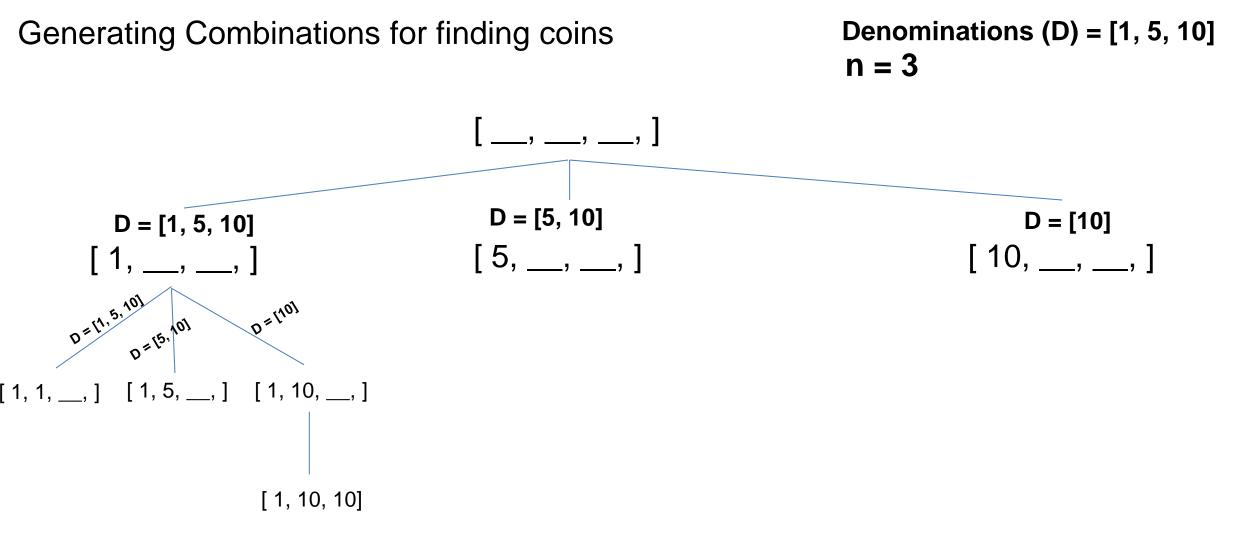
[1, 1, 1] [1, 1, 5] [1, 1, 10]





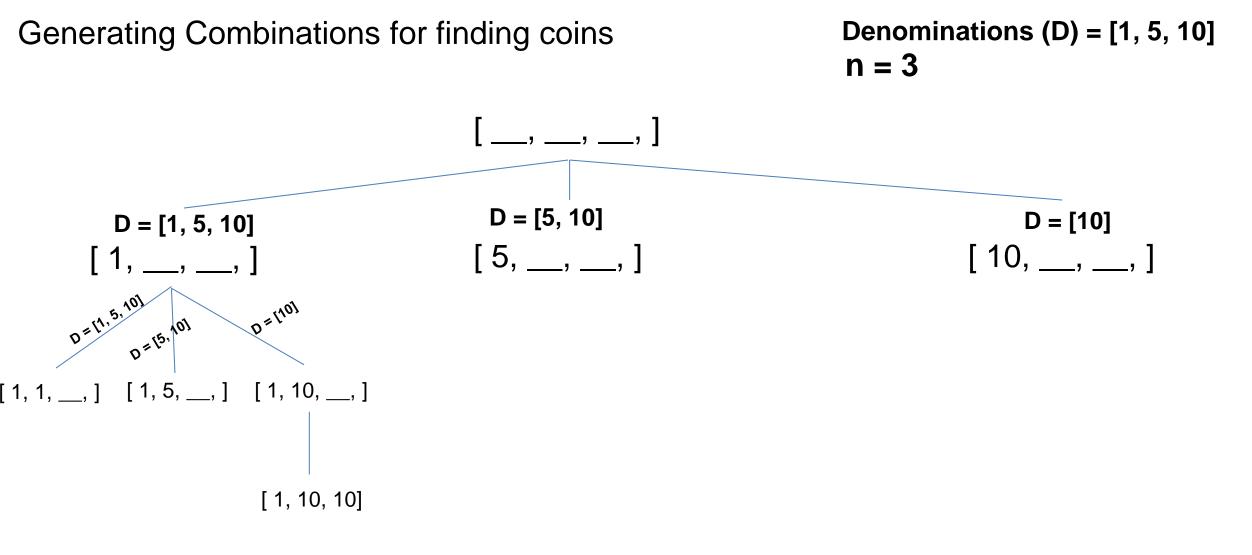
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10]





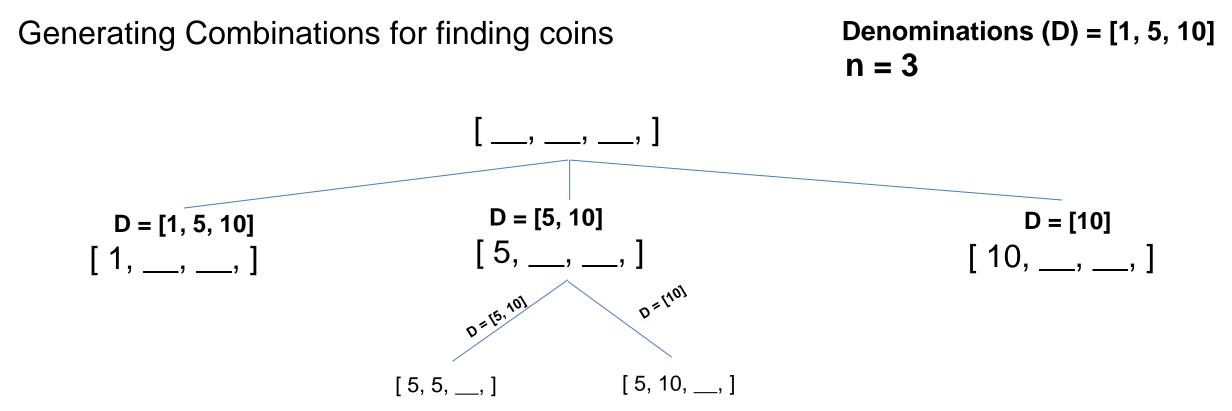
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10]





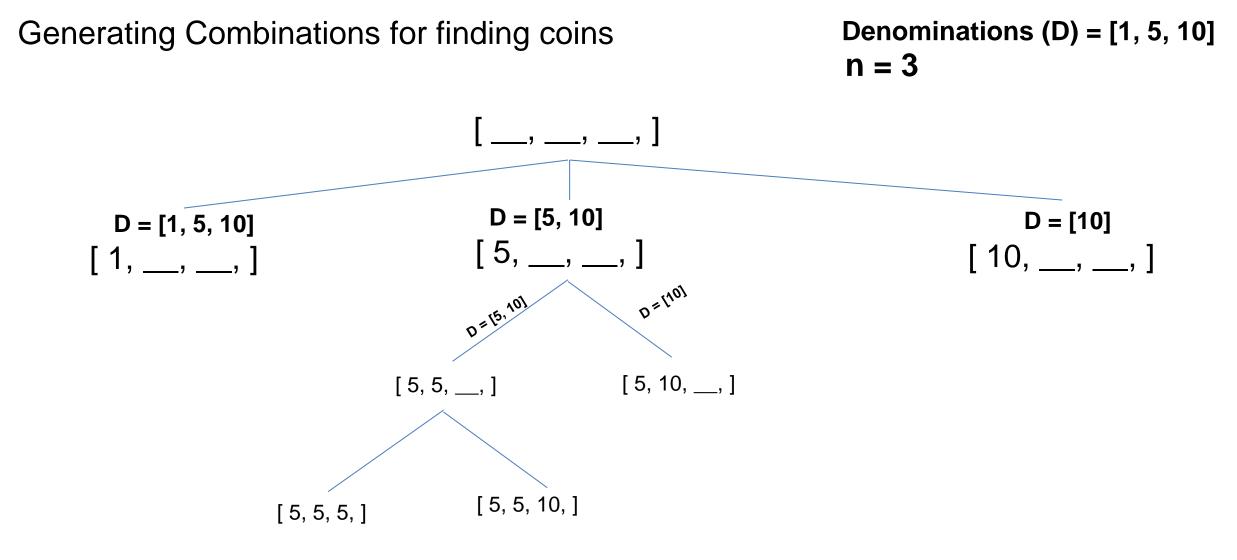
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]





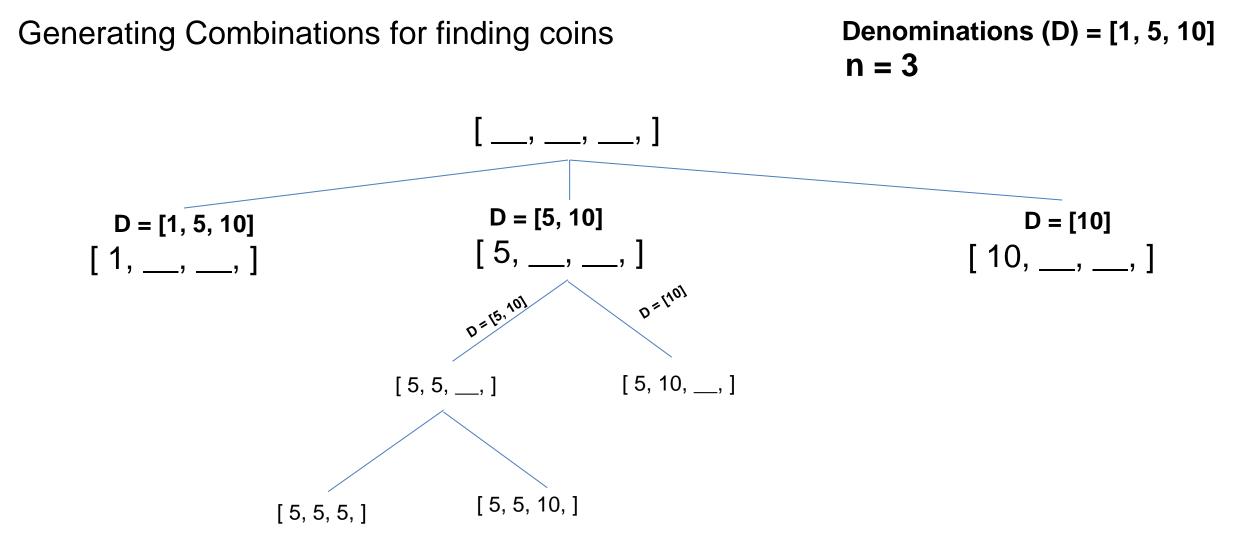
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]





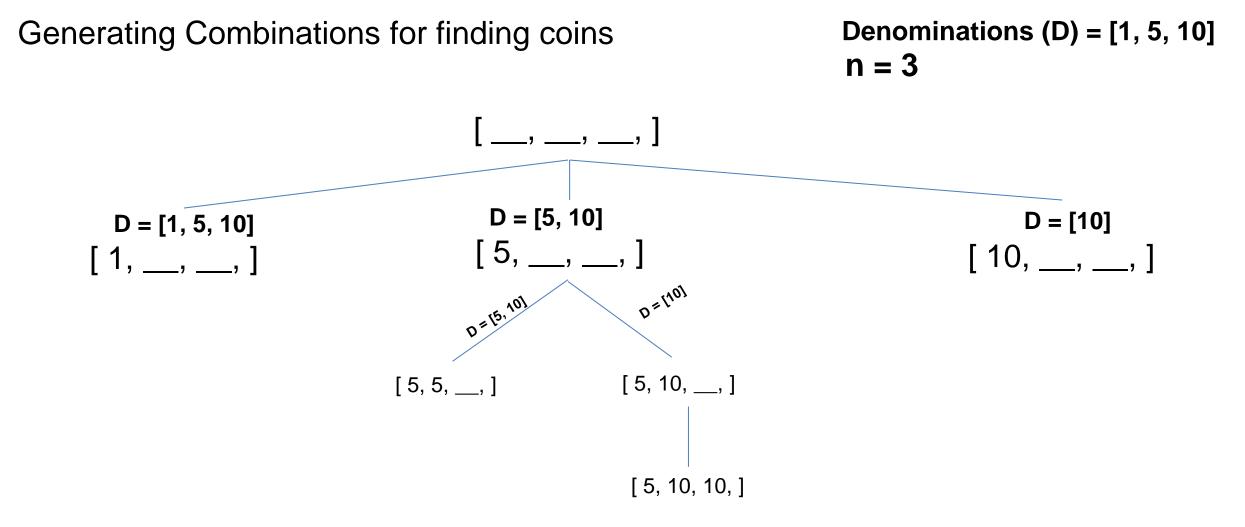
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]





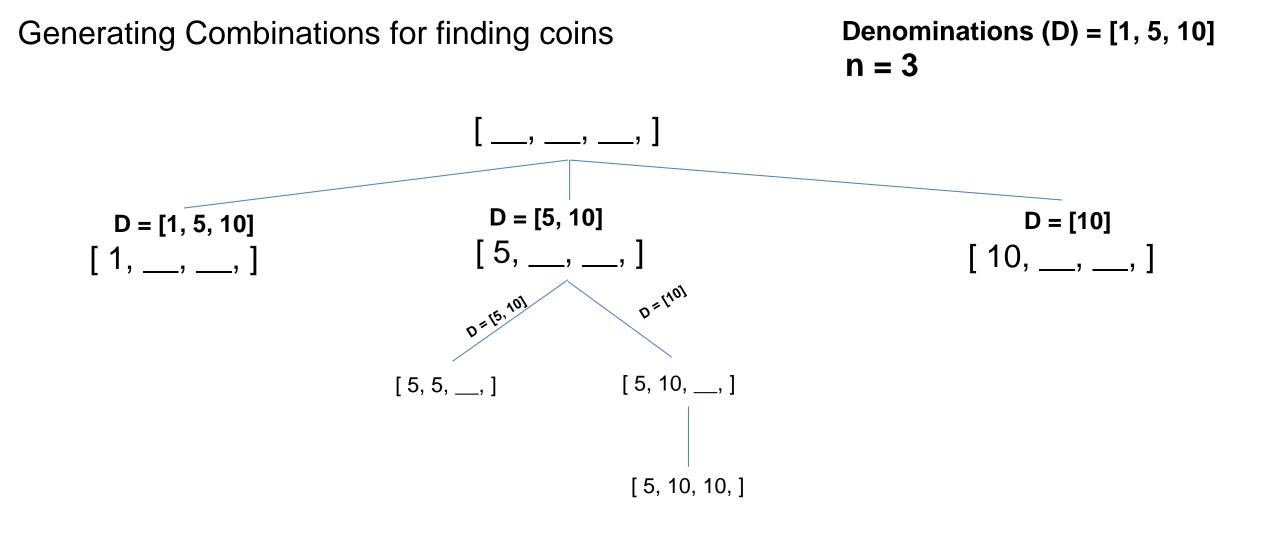
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10]





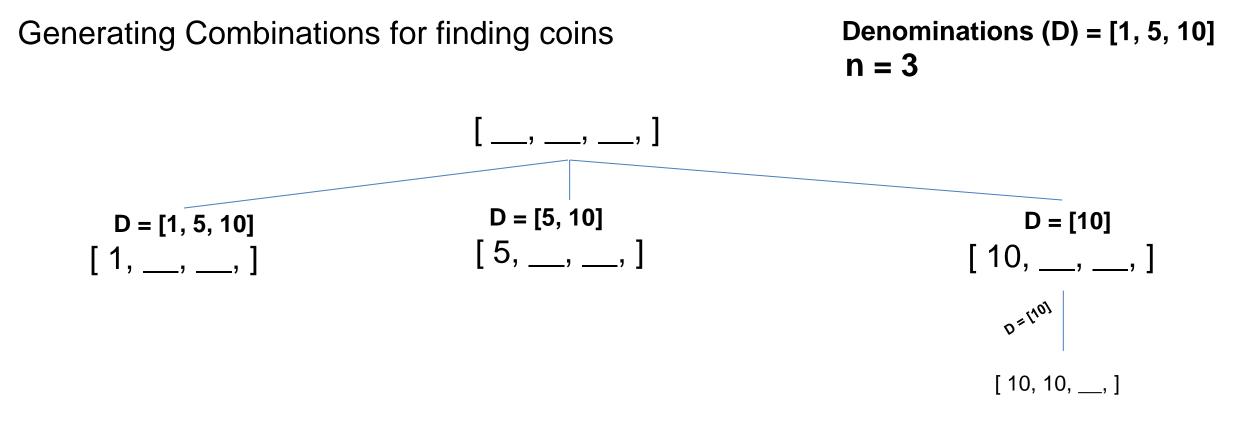
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10]





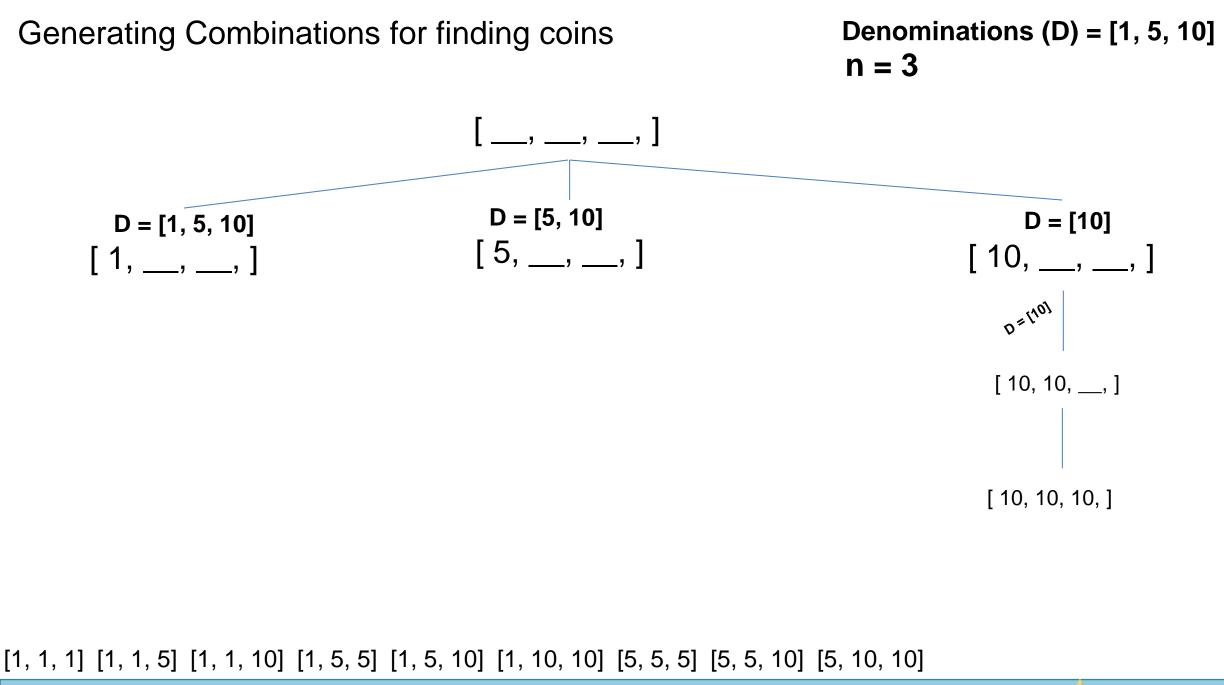
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10] [5, 10, 10]



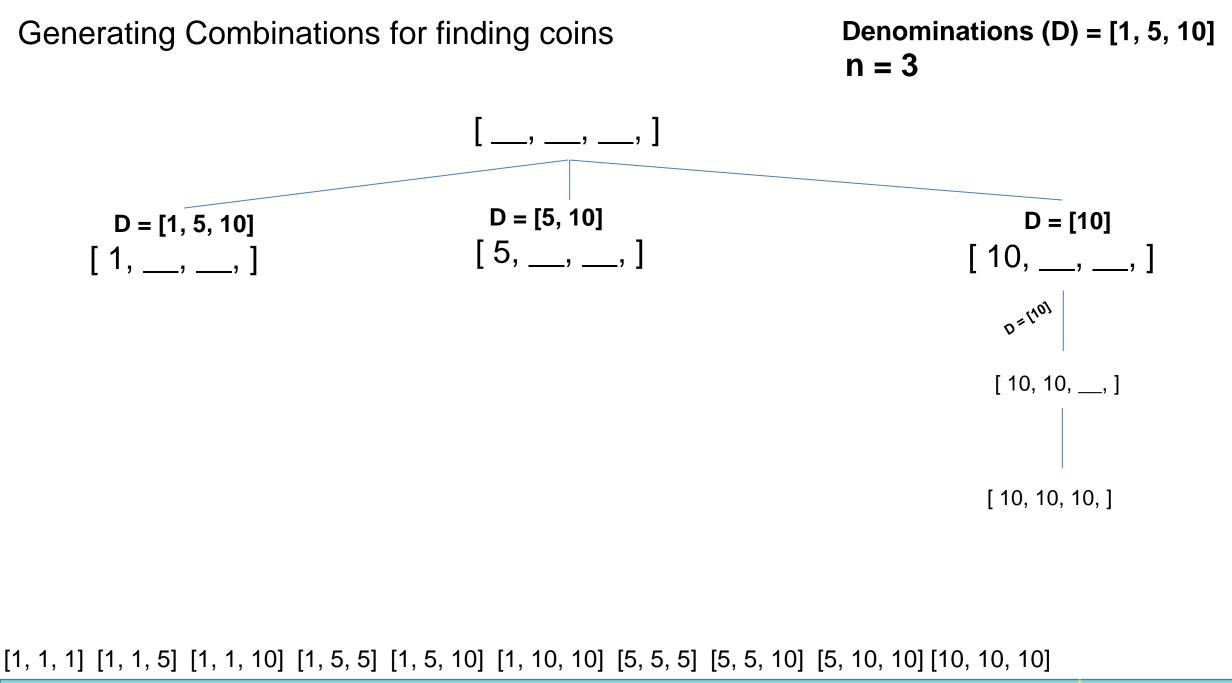


[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10] [5, 10, 10]





MONTANA 87



MONTANA

Denominations (D) = [1, 5, 10] n = 3

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]



We've generated all combinations of length 3



Denominations (D) = [1, 5, 10] n = 3

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]

We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to ${\bf K}$



Denominations (D) = [1, 5, 10] n = 3 K = 16

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]



We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to ${\bf K}$

Suppose K = 16 (a minimum of 3 coins is needed to make 16 cents)



Denominations (D) = [1, 5, 10] n = 3 K = 16

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]



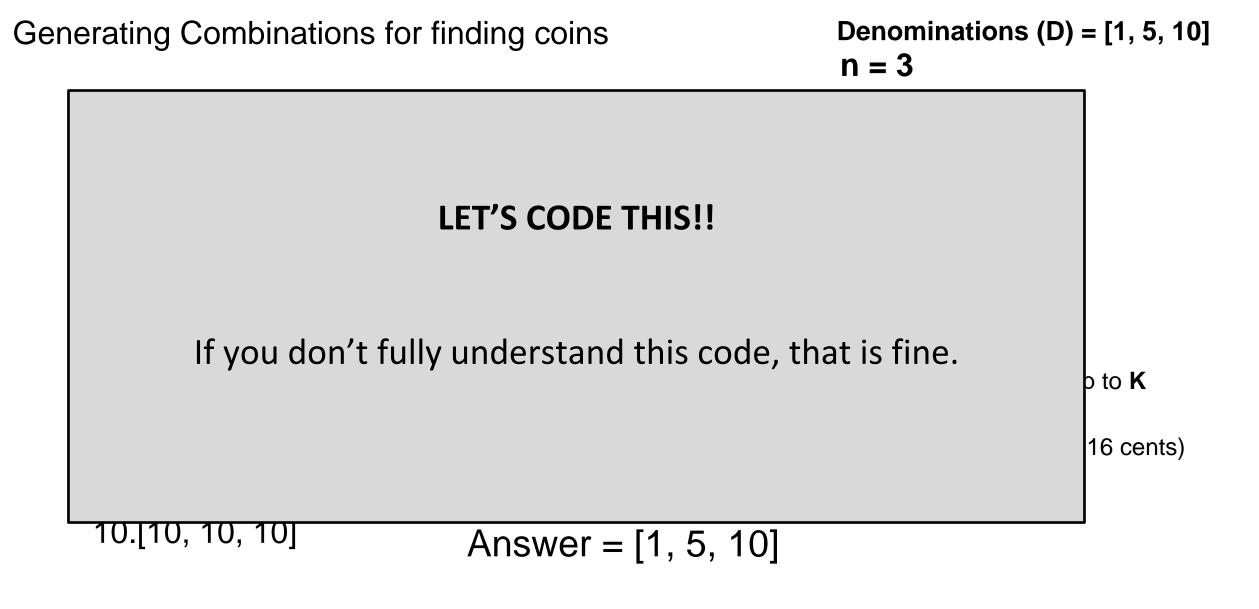
We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to K

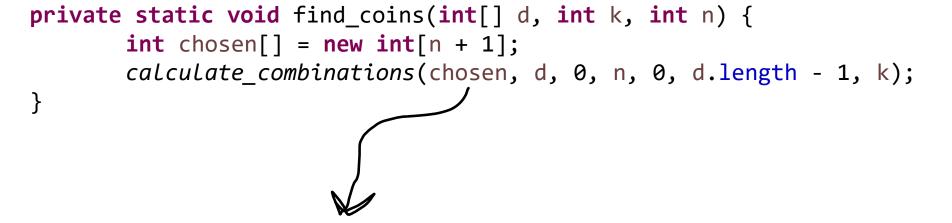
Suppose K = 16 (a minimum of 3 coins is needed to make 16 cents)

Answer = [1, 5, 10]









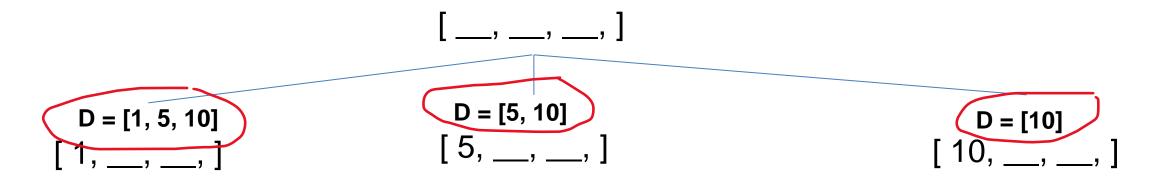
Array that we build up over time. Holds **indices** of currently selected denominations for some combination

[1,___, ___] [1,1, ___] [1,1, 5]



```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```

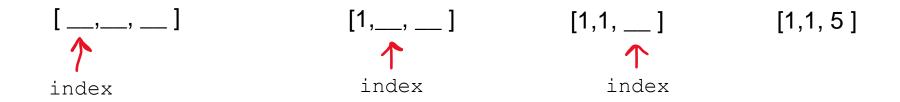
Array of denominations we pass for each recursive call





```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```

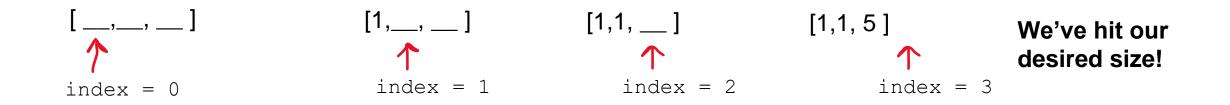
The next index that we need to insert at for chosen array





```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```

The desired size of the combination. When index == r, we have reached the desired combination size





```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
       if (index == r) {
               int counter = 0;
               ArrayList<Integer> coins = new ArrayList<Integer>();
               for (int i = 0; i < r; i++) {</pre>
                       counter += arr[chosen[i]];
                       coins.add(arr[chosen[i]]);
               if(counter == target) {
                       System.out.println(coins);
               return;
       for (int i = start; i <= end; i++) {</pre>
               chosen[index] = i;
               calculate combinations(chosen, arr, index + 1, r, i, end, target);
        }
       return;
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
       if (index == r) {
                                                                             If we hit our base
                int counter = 0;
                                                                             case, we know we
                ArrayList<Integer> coins = new ArrayList<Integer>();
                                                                             have N things, so
  Base case
                for (int i = 0; i < r; i++) {</pre>
                                                                             put them in an
                        counter += arr[chosen[i]];
                                                                             ArrayList and add
                        coins.add(arr[chosen[i]]);
                                                                             them up
                                                       Only print out the combination if it
                if(counter == target) { 
                                                       adds up to target
                        System.out.println(coins);
                return;
        for (int i = start; i <= end; i++) {</pre>
                chosen[index] = i;
                calculate combinations(chosen, arr, index + 1, r, i, end, target);
        return;
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
        if (index == r) {
                int counter = 0;
                ArrayList<Integer> coins = new ArrayList<Integer>();
                for (int i = 0; i < r; i++) {</pre>
                         counter += arr[chosen[i]];
                         coins.add(arr[chosen[i]]);
                if(counter == target) {
                         System.out.println(coins);
                return;
Recursive Case
        for (int i = start; i <= end; i++) {
    chosen[index] = i;</pre>
                calculate_combinations(chosen, arr, index + 1, r, i, end, target);
        return:
}
```

Otherwise, insert selected coin into the chosen array

create (end-start) branches, and give it a smaller section of D



