CSCI 132: Basic Data Structures and Algorithms

Searching (Binary Search)

Reese Pearsall Fall 2023

Announcements

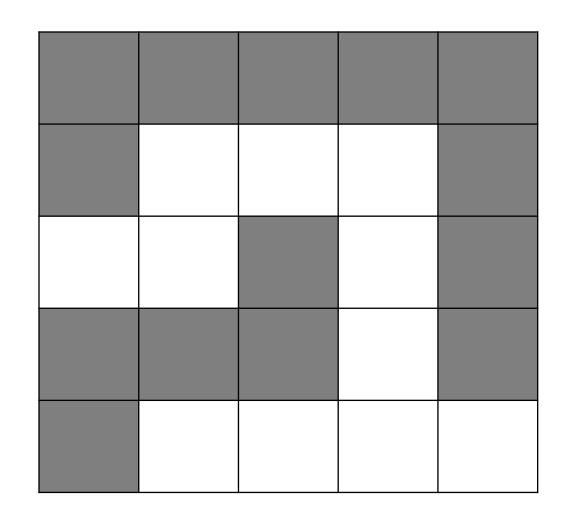
Lab 12 due tomorrow @11:59 PM

Program 5 due Sunday December 10th

Gradebook



char[][] maze [[#, #, #, # ,#], [#,.,.,#], [.,.,#,.,#], [#, #, #, . , #], [#,.,.,.],



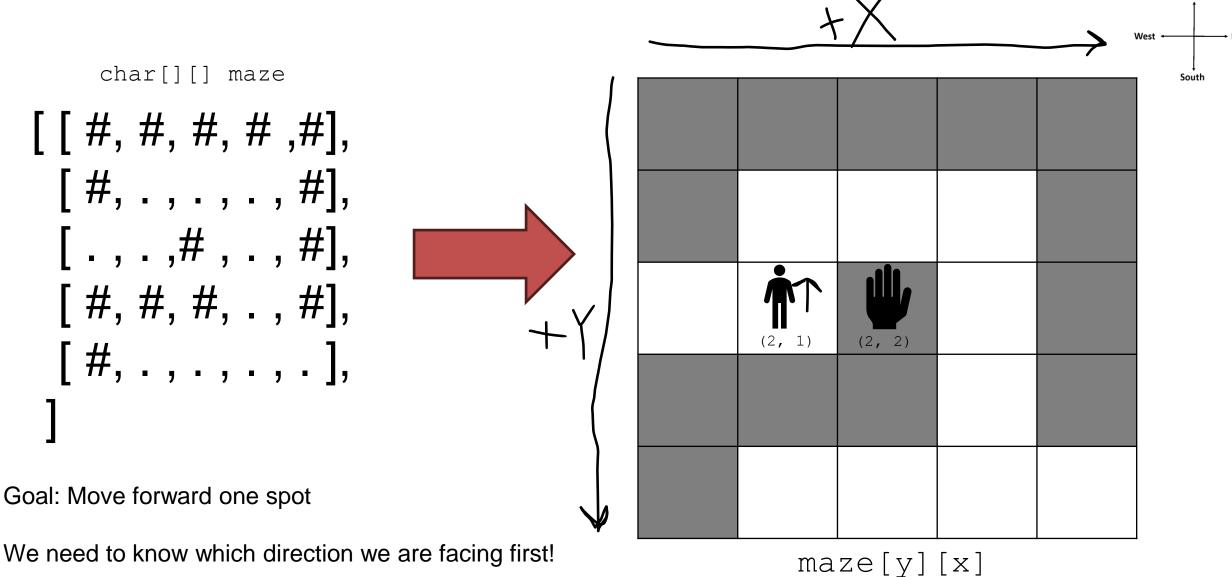
char[][] maze [[#, #, #, # ,#], [#,.,.,#], [.,.,#,.,#], [#, #, #, ., #], [#,.,.,.], maze[0]

char[][] maze [[#, #, #, #, #], [#, ., ., ., #],[.,.,#,.,#], [#, #, #, . , #], [#,.,.,.], maze[1]

char[][] maze [[#, #, #, #, #], [#, ., ., ., #],[.,.,#,.,#], [#, #, #, ., #], [#,.,.,.], maze[1][0]

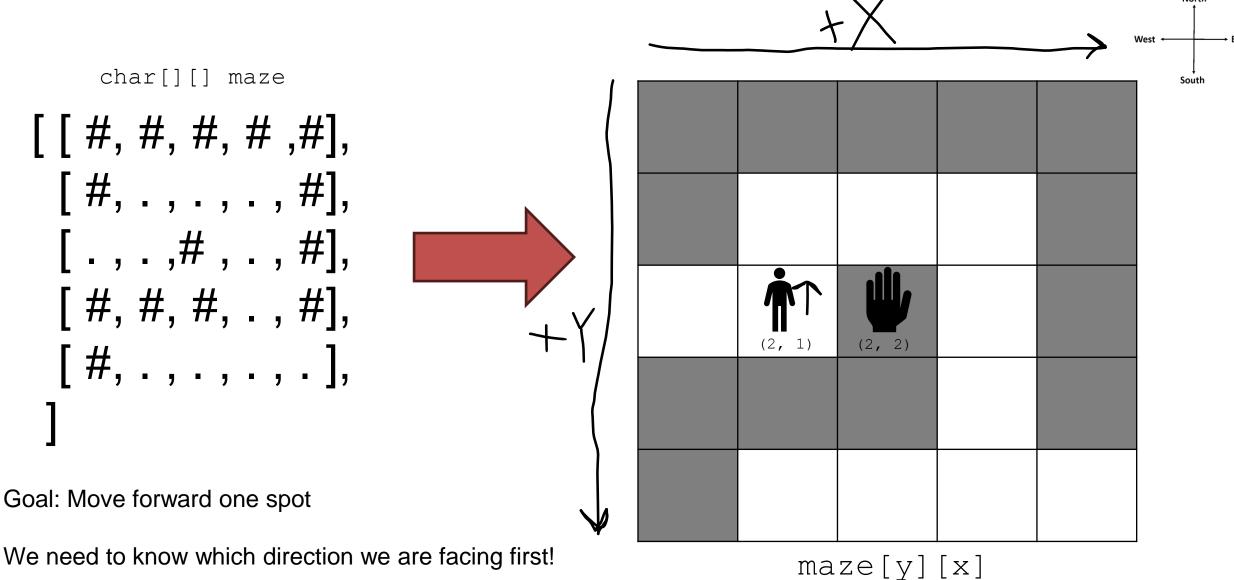
```
char[][] maze
[[#, #, #, #, #],
 [#,.,.,#]
 [.,.,#,.,#],
 [ #, #, #, . , #],
 [#,.,.,.],
   maze[1][2]
```

char[][] maze [[#, #, #, # ,#], [#,.,.,#], [.,.,#,.,#], [#, #, #, . , #], [#,.,.,.], maze[y][x]

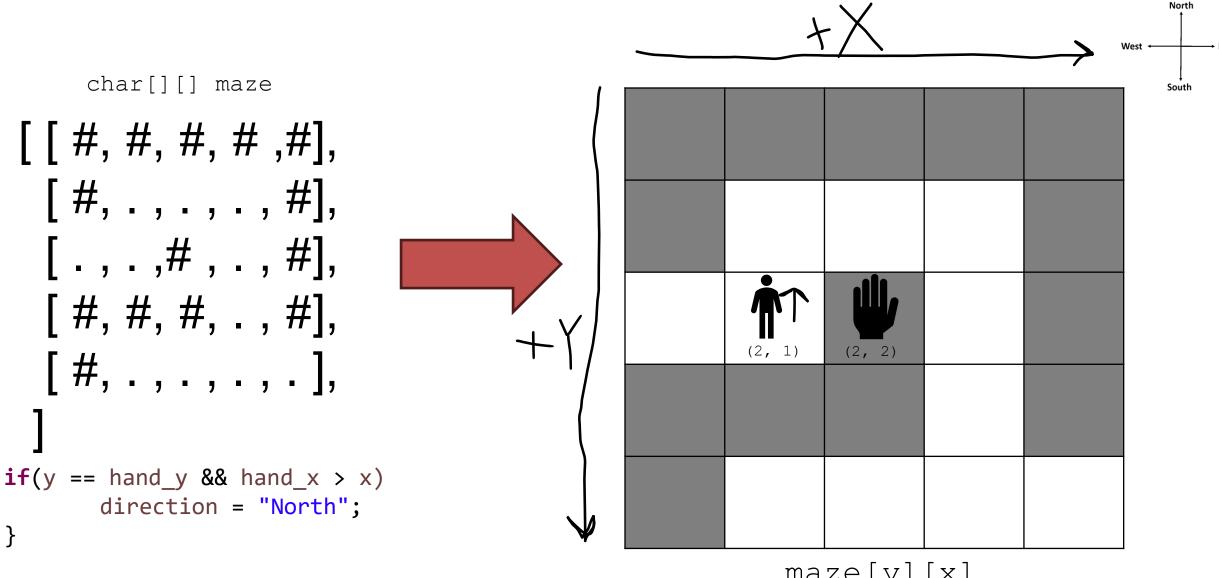


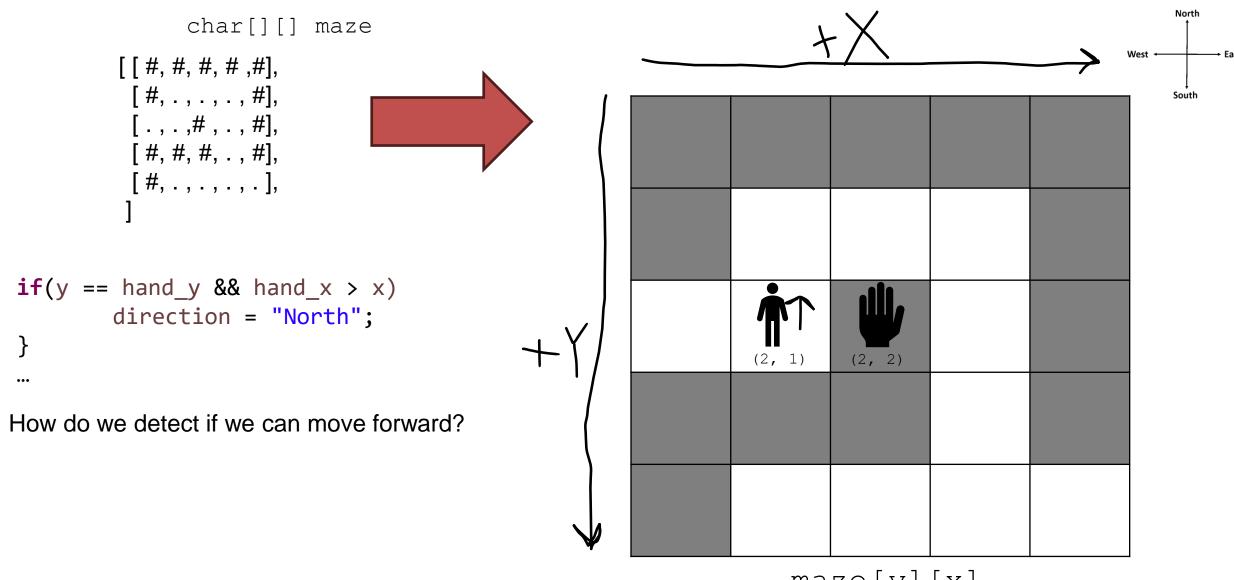
We need to know which direction we are facing first!

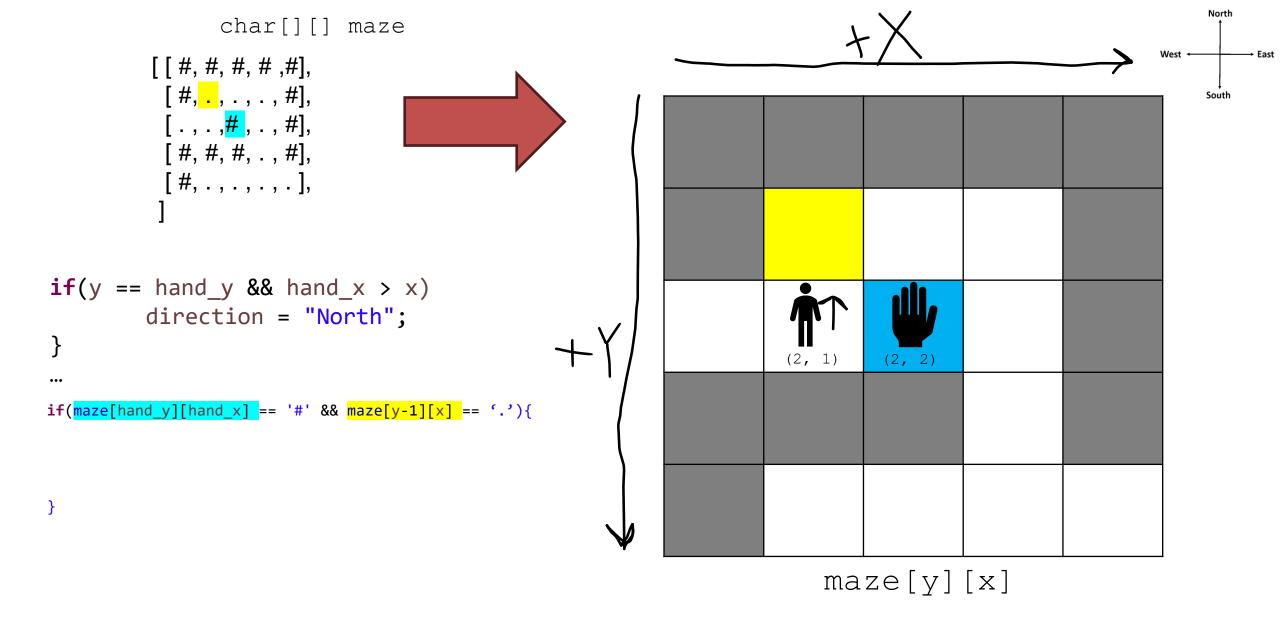
How do we know direction we are facing?

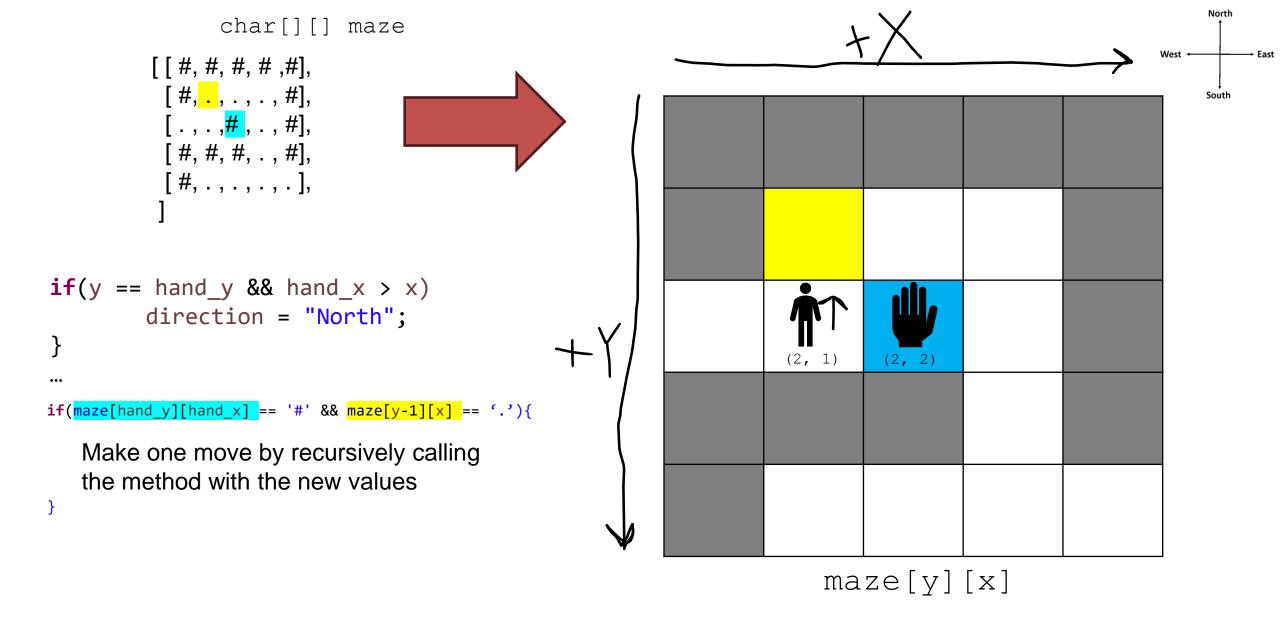


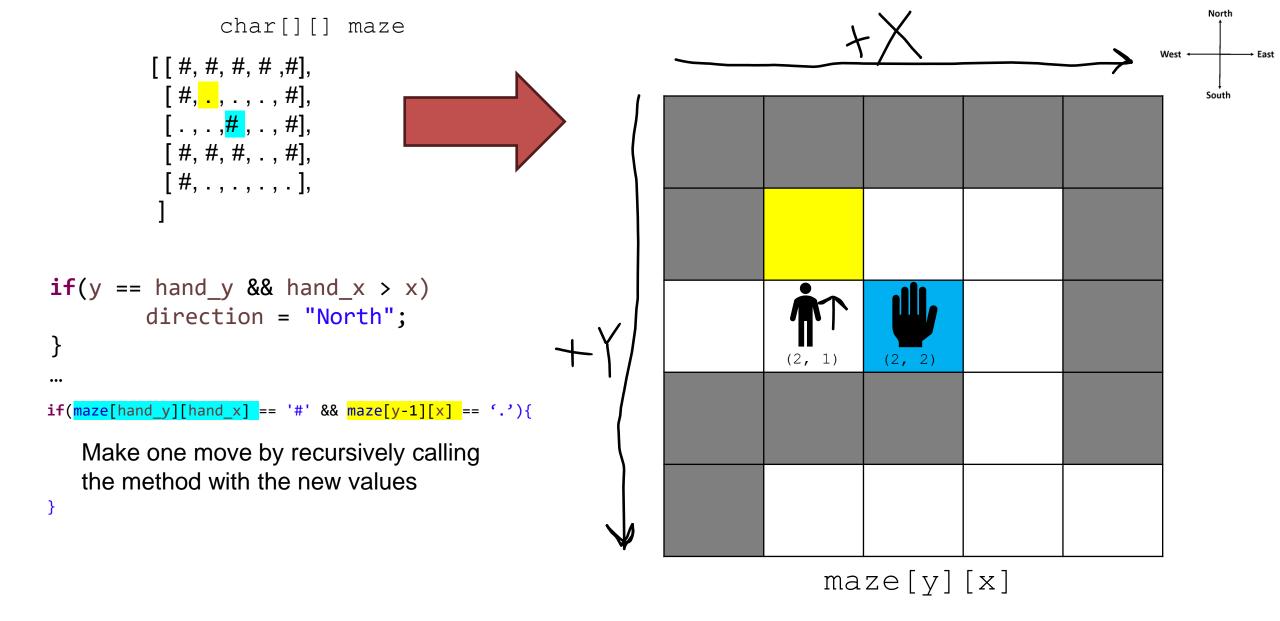
Our character Y value and our hand's Y value is the same, And our character's X value is *less than* our hands' X value

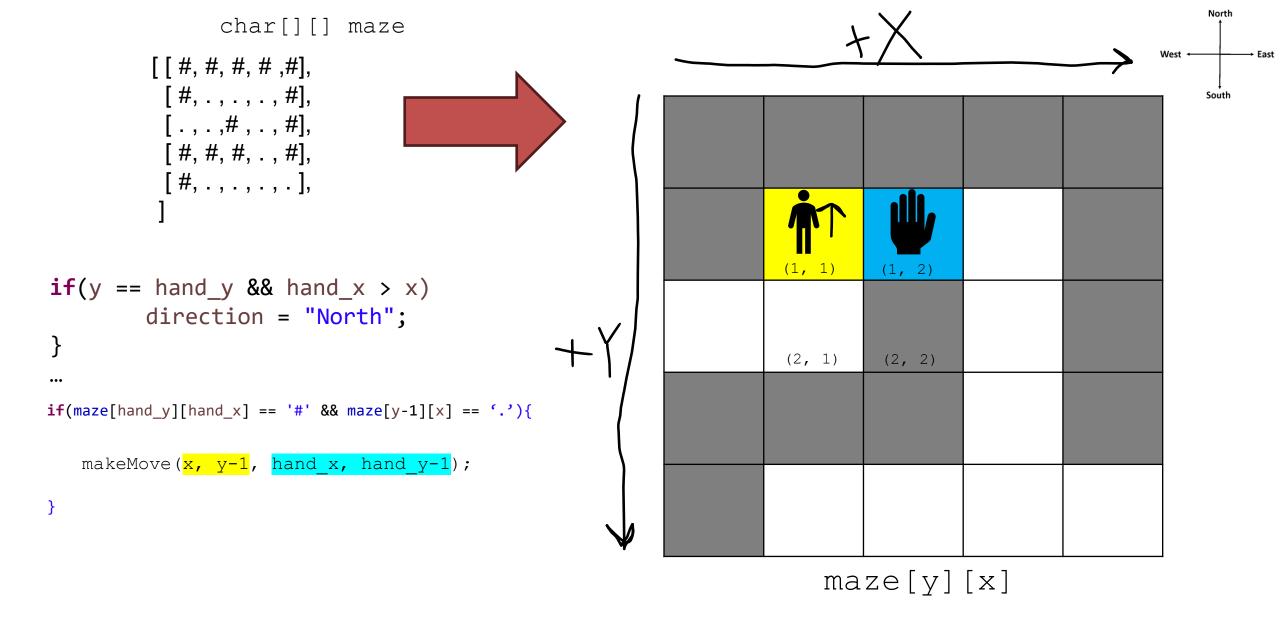




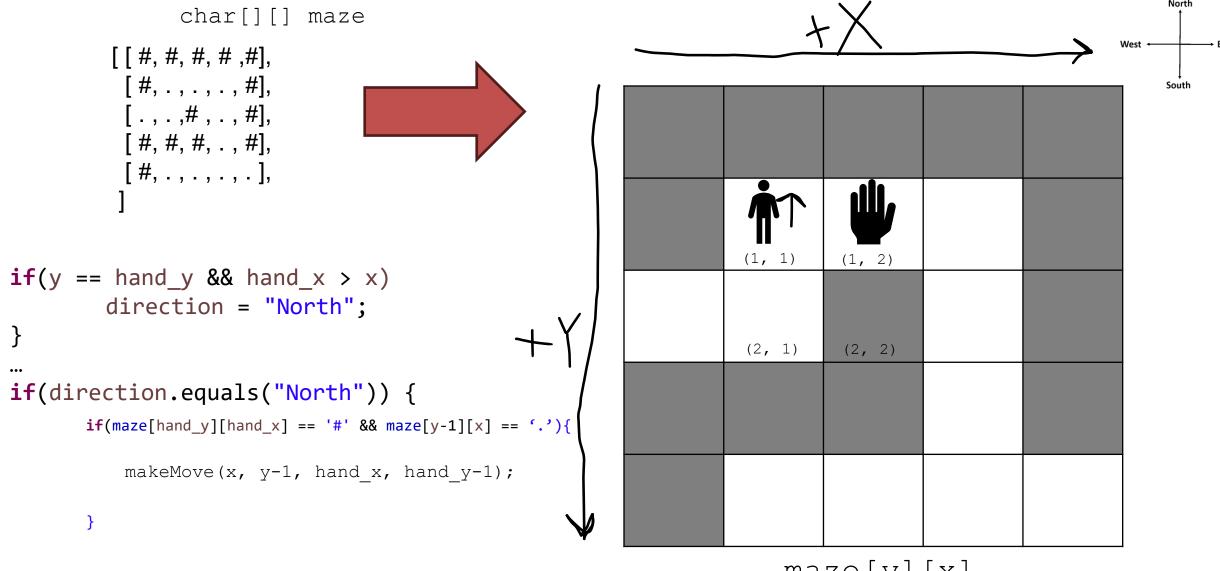








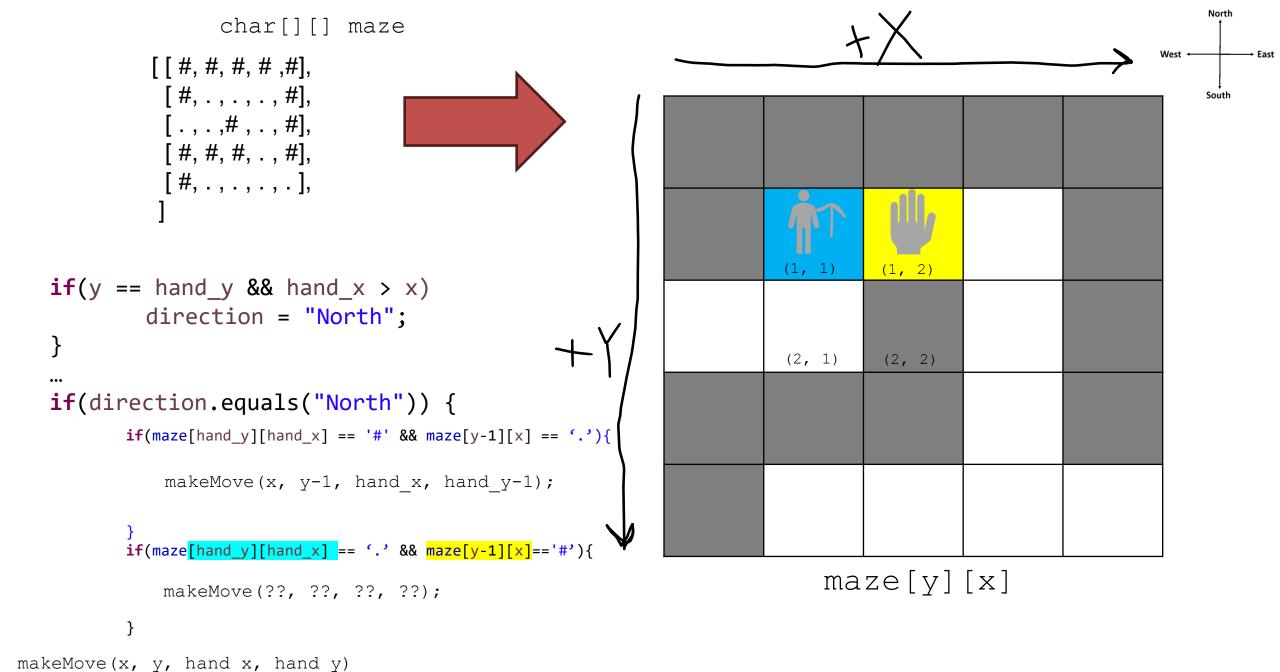
```
North
               char[][] maze
         [[#, #, #, #, #],
          [#,.,.,#],
          [.,.,#,.,#],
          [#, #, #, ., #],
         [#,.,.,.],
                                                                 (1, 1)
                                                                          (1, 2)
if(y == hand_y \&\& hand_x > x)
        direction = "North";
                                                                  (2, 1)
                                                                          (2, 2)
if(direction.equals("North")) {
      if(maze[hand_y][hand_x] == '#' && maze[y-1][x] == '.'){
          makeMove(x, y-1, hand_x, hand_y-1);
                                                                     maze[y][x]
```

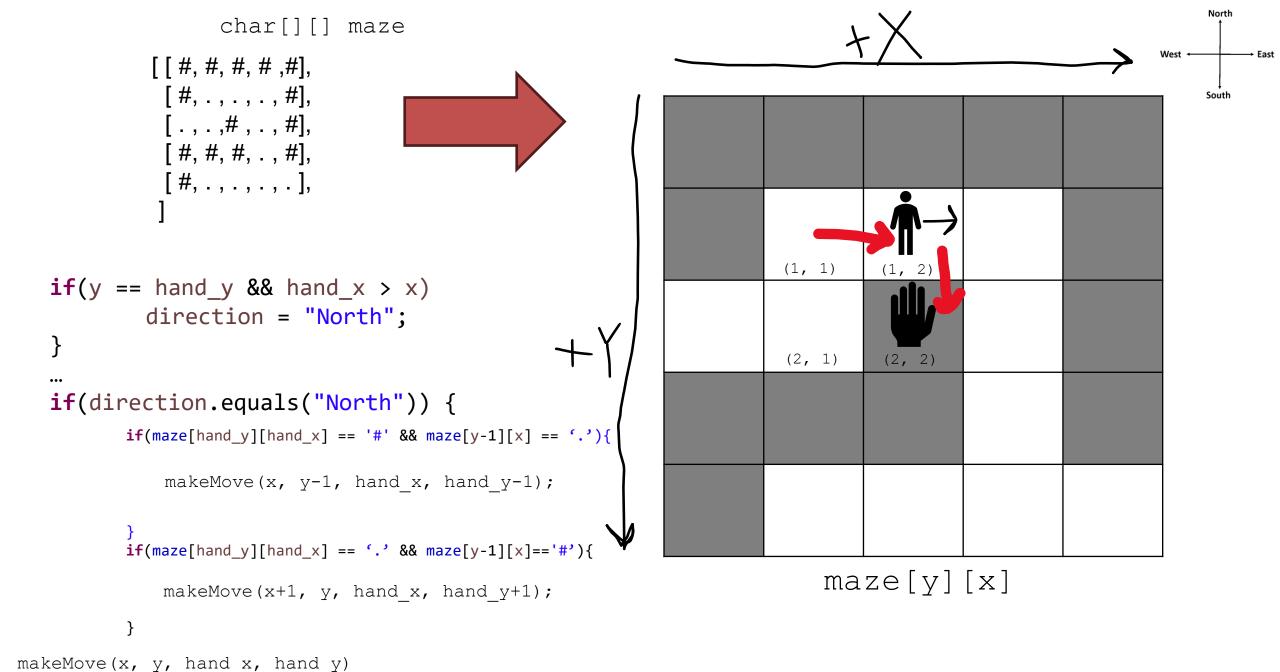


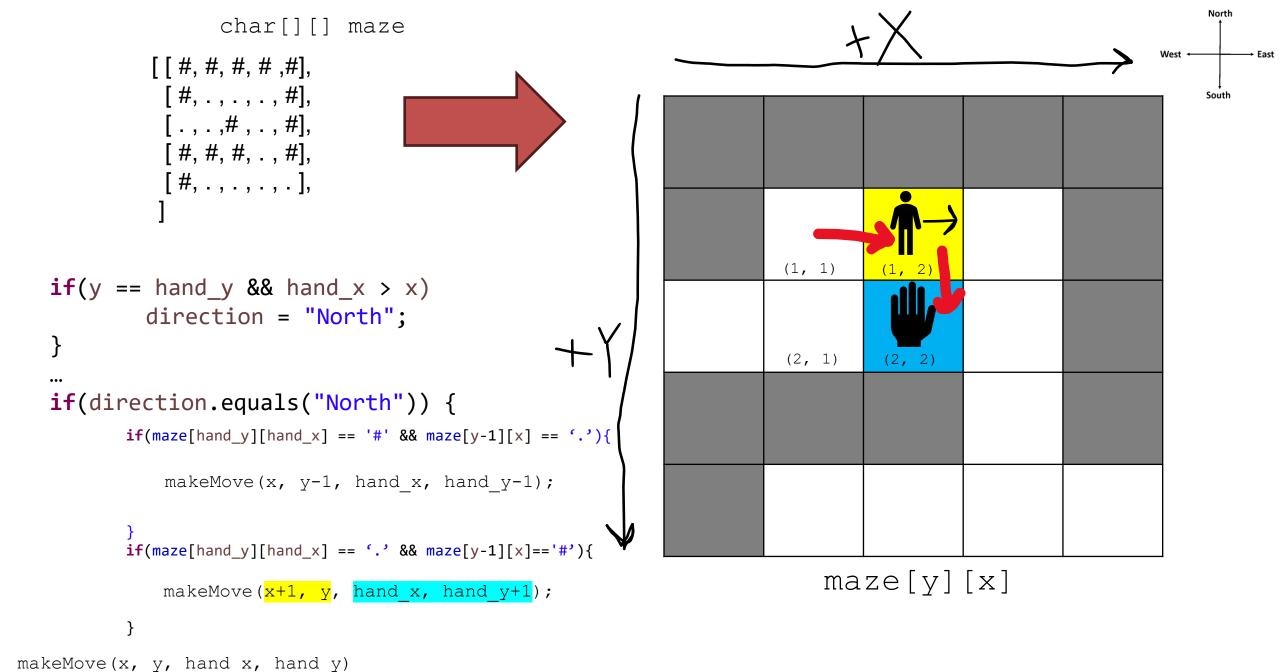
Turn right and move forward one spot?

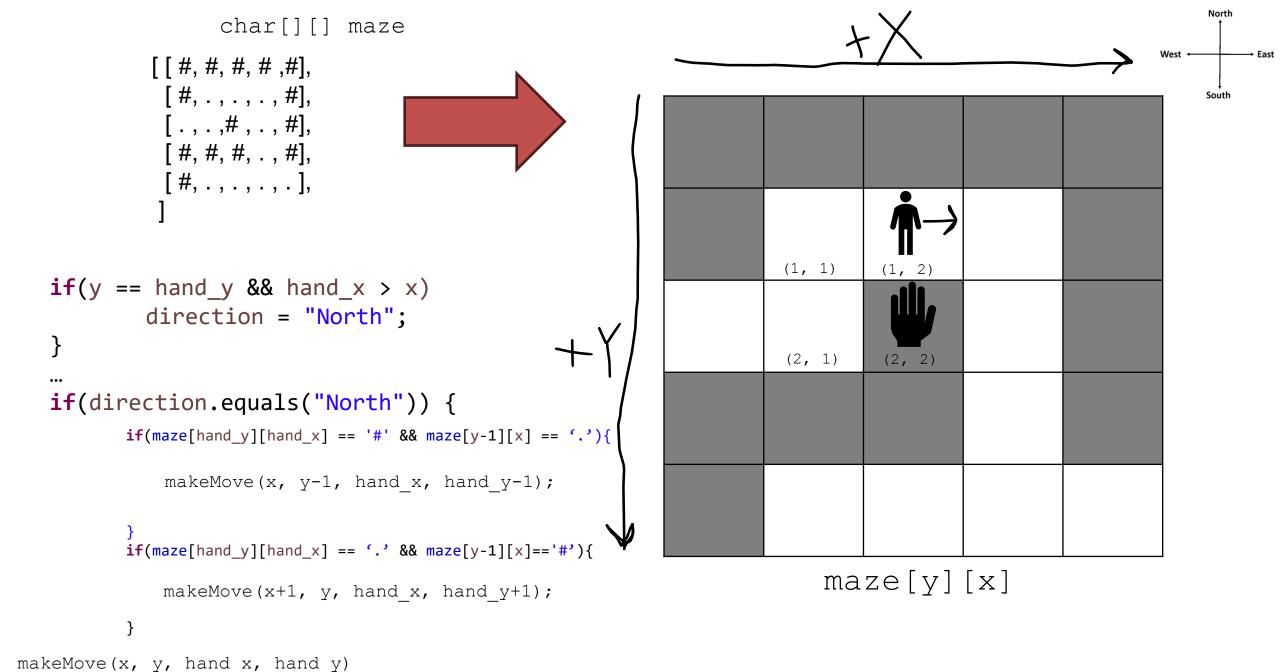
maze[y][x]

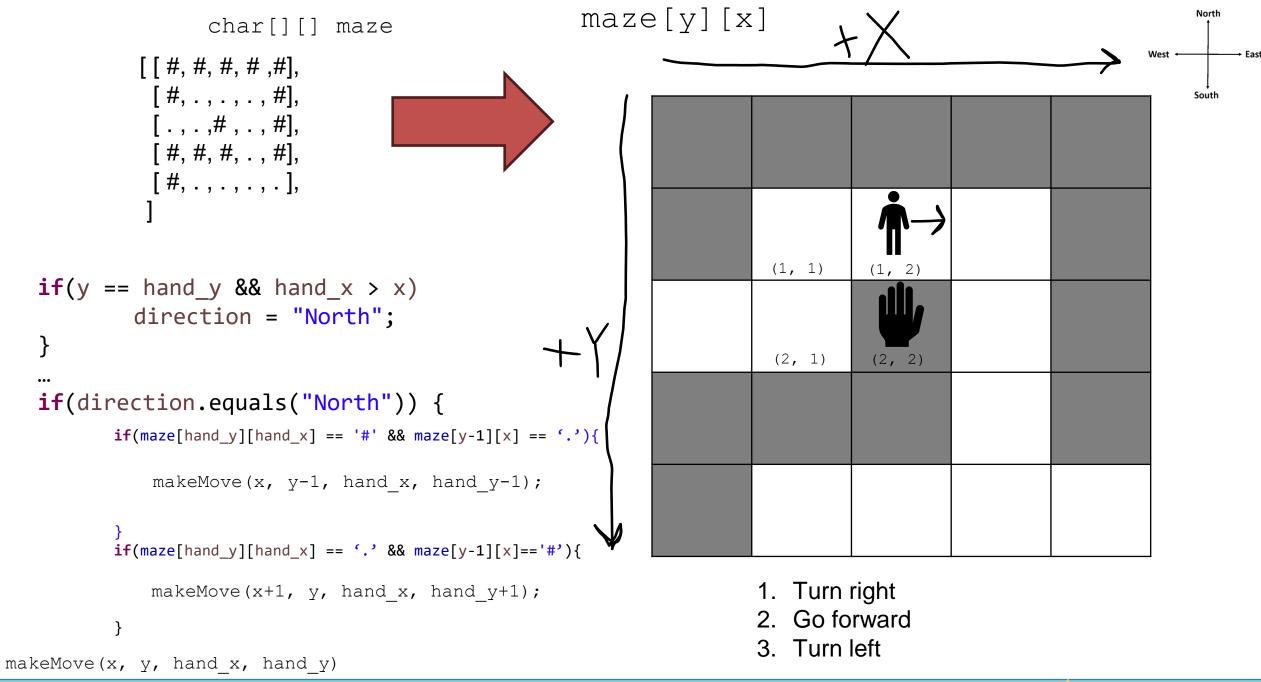
```
char[][]
                          maze
         [[#, #, #, #, #],
          [#,.,.,#],
          [.,.,#,.,#],
          [#, #, #, ., #],
          [#,.,.,.],
                                                                           (1, 2)
if(y == hand_y \&\& hand_x > x)
        direction = "North";
                                                                   (2, 1)
                                                                            (2, 2)
if(direction.equals("North")) {
      if(maze[hand_y][hand_x] == '#' && maze[y-1][x] == '.'){
          makeMove(x, y-1, hand_x, hand_y-1);
      if(maze[hand_y][hand_x] == '.' && maze[y-1][x]=='#'){
                                                                      maze[y][x]
```

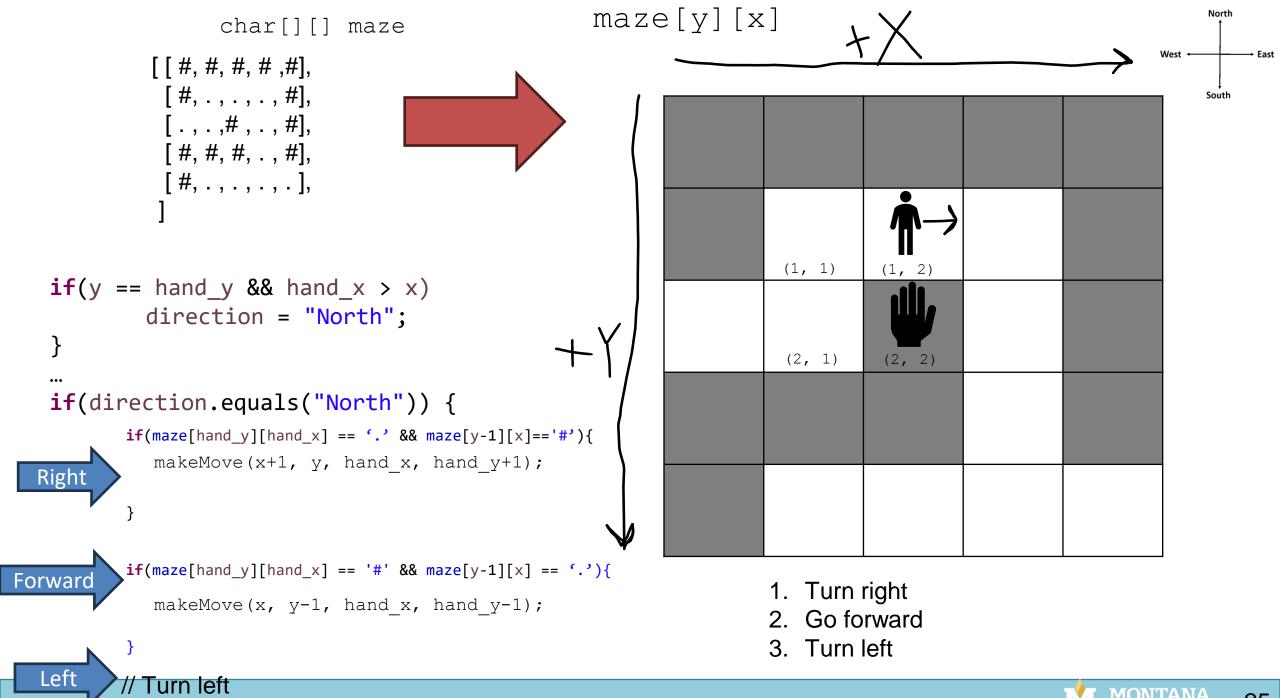


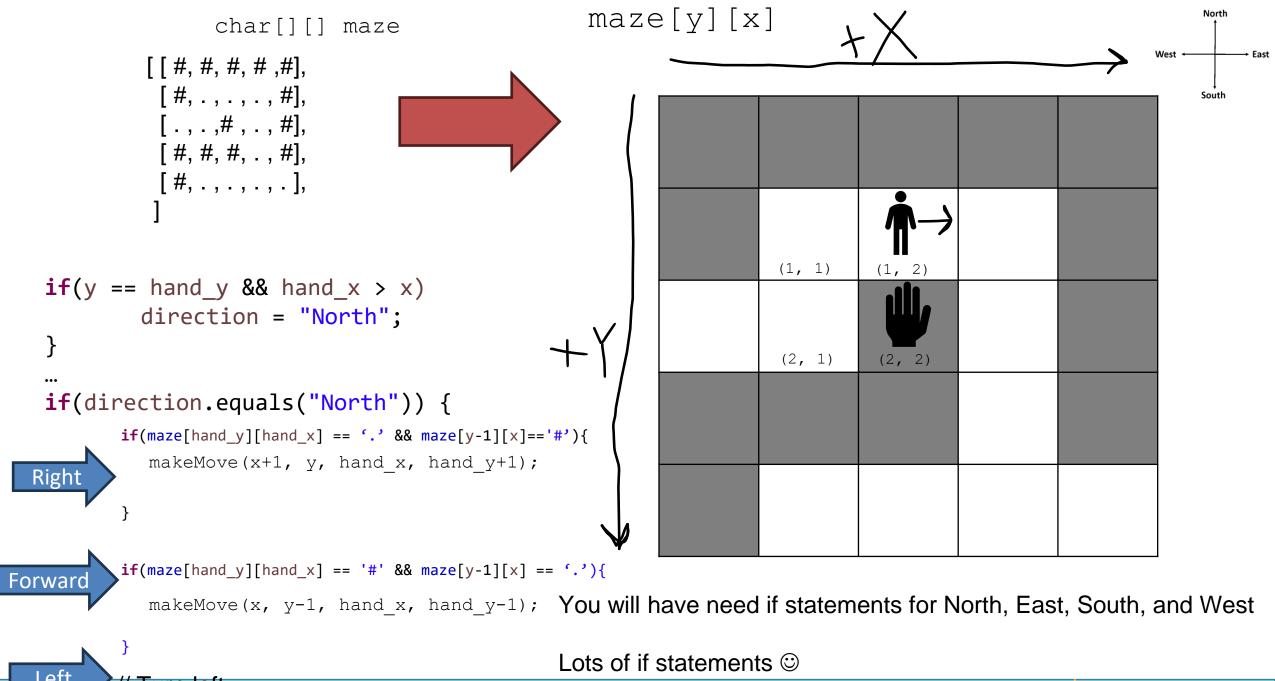


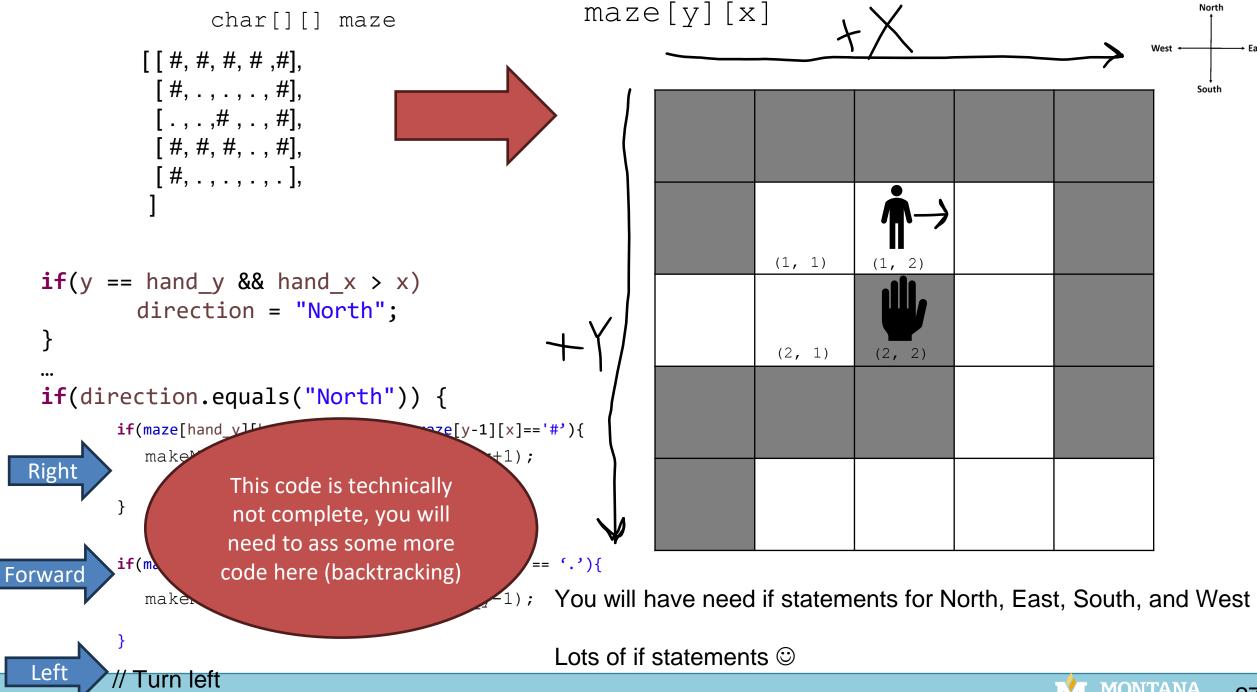


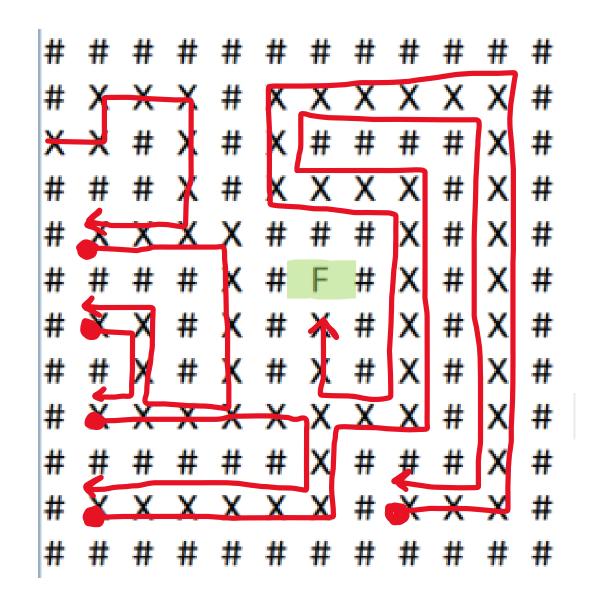












= Backtracking path



Searching

We store values in data structures, but we also need to retrieve/search for values!

Today, we will discuss techniques for how to search for a value in a data structure

(We will be using arrays, but these techniques could also be used on Linked Lists, queues, stacks, etc)



Searching

Option 1: Linear Search

Check every spot until one by one until we find what we are looking for

```
public int linear_search(int[] array, int s) {
     for(int i = 0; i < array.length; i++) {</pre>
          if(array[i] == s) {
                return i;
     return -1;
```

Searching

Option 1: Linear Search

Check every spot until one by one until we find what we are looking for

Not efficient for large data structures. O(n) running time

```
public int linear_search(int[] array, int s) {
     for(int i = 0; i < array.length; i++) {</pre>
          if(array[i] == s) {
                return i;
     return -1;
```



 0

 1
 2
 9
 10
 11
 15
 18
 21
 27
 31
 41
 43
 50

What if our array is sorted?

Target Value: 27

0												12	
1	2	9	10	11	15	18	21	27	31	41	43	50	

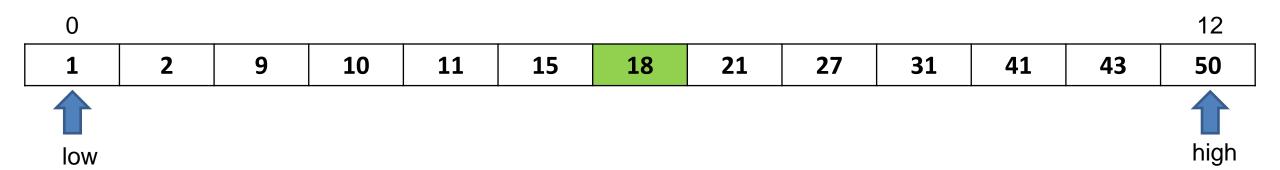
We can leverage the fact that this array is sorted to make searching more efficient Target Value: 27

 1
 2
 9
 10
 11
 15
 18
 21
 27
 31
 41
 43
 50

1. Start at the middle of the array

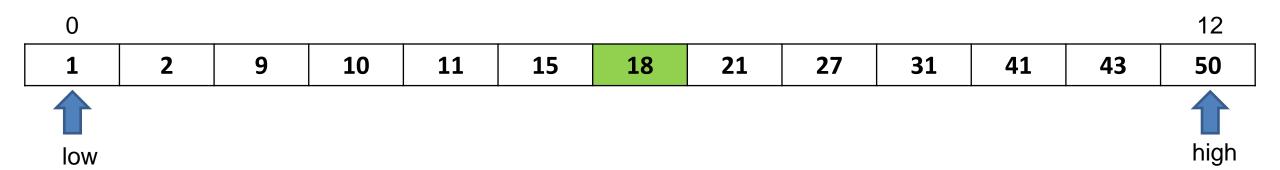
0												12	
1	2	9	10	11	15	18	21	27	31	41	43	50	

- 1. Start at the middle of the array
- 2. Compare to target value:
- → If the value is the target value, return
- → If the target value is greater than the middle, discard the "left section" of the array
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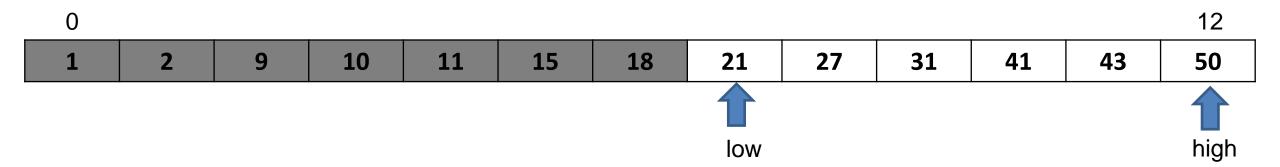
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We will define two pointers, low and high that point to the possible bounds of the target value



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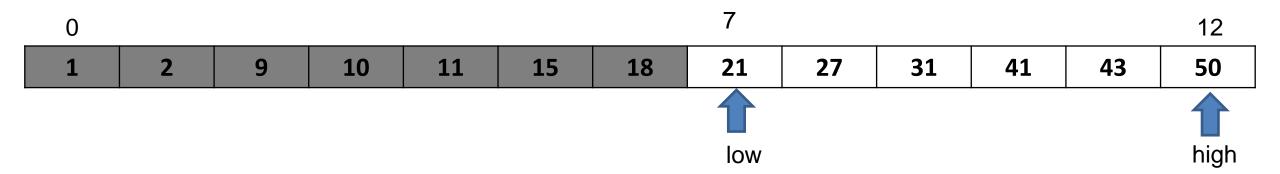
We will define two pointers, low and high that point to the possible bounds of the target value



- 1. Start at the middle of the array
- 2. Compare to target value:
- → If the value is the target value, return
- → If the target value is greater than the middle, discard the "left section" of the array (move the low pointer)
- → If the target value is less than the middle, discard the "right section" of the array (move the high pointer)

Because we know the array is sorted, and the target value is greater than our mid point, then we know the target value must be located somewhere to the right.

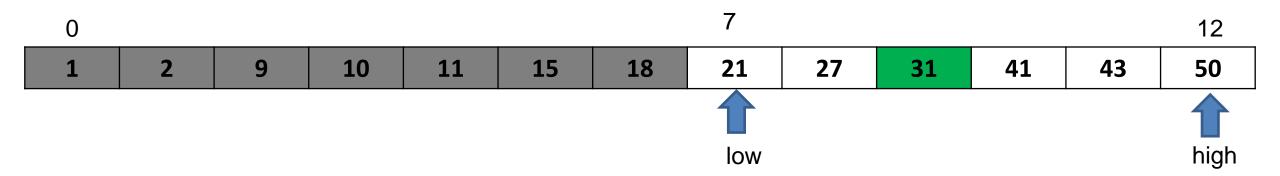
We can eliminate half of the array!!!



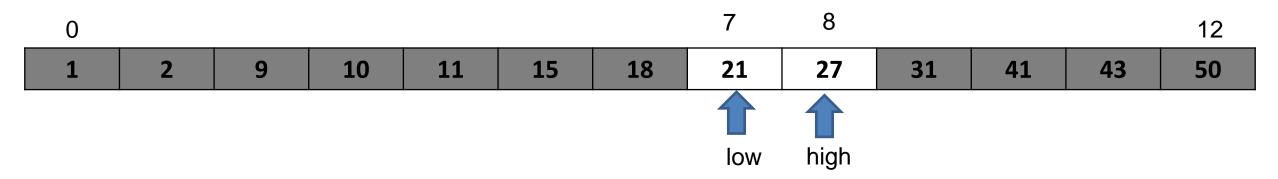
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- 3. Recalculate the mid point, and repeat loop back to step 2 until target value is found



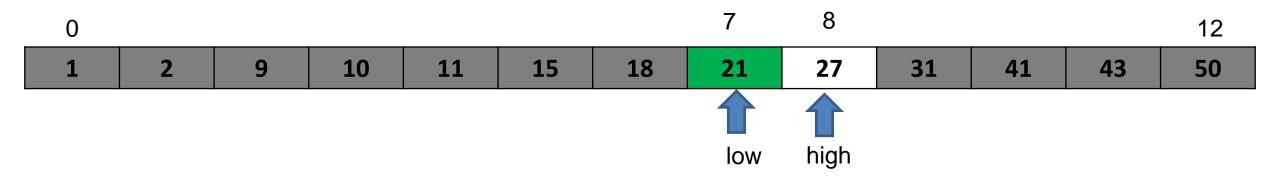
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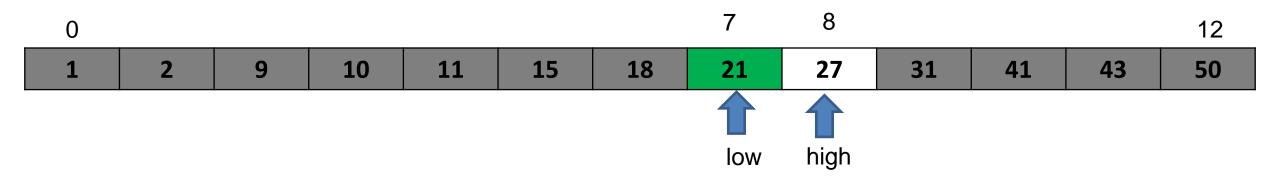
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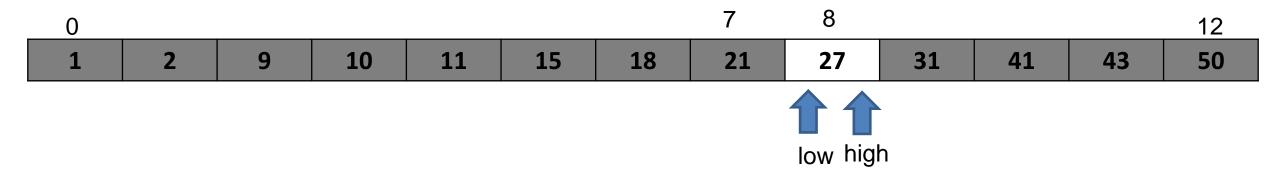
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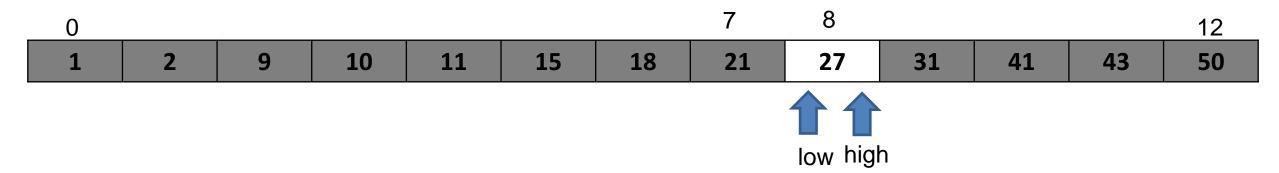
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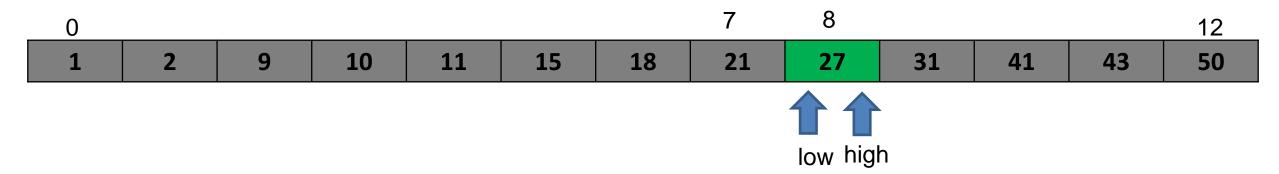
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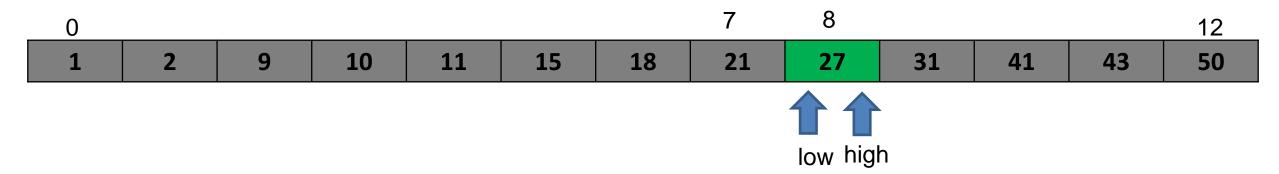
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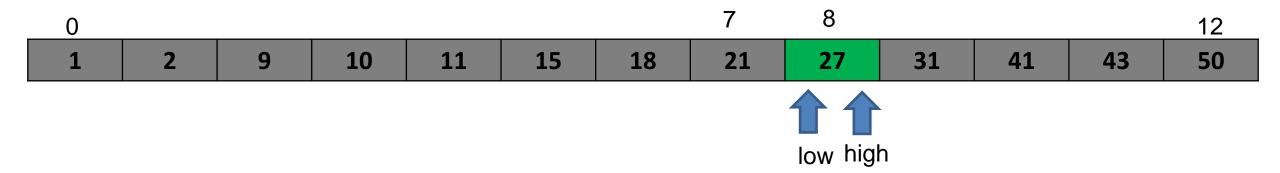
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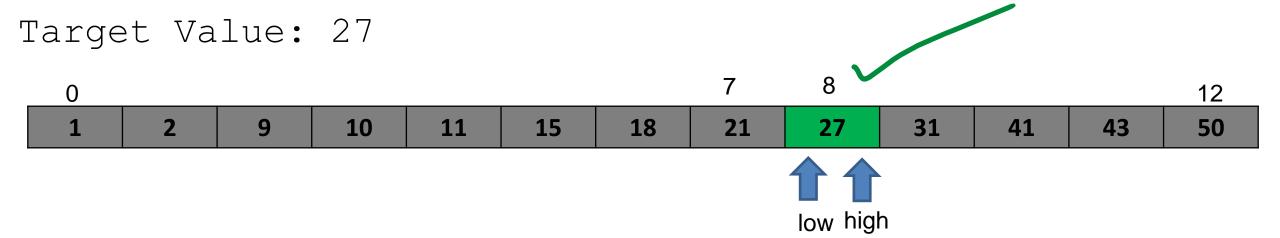


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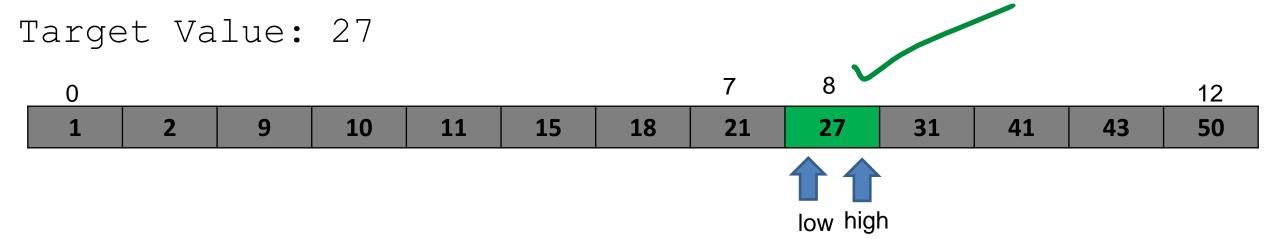
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This algorithm is known as Binary Search



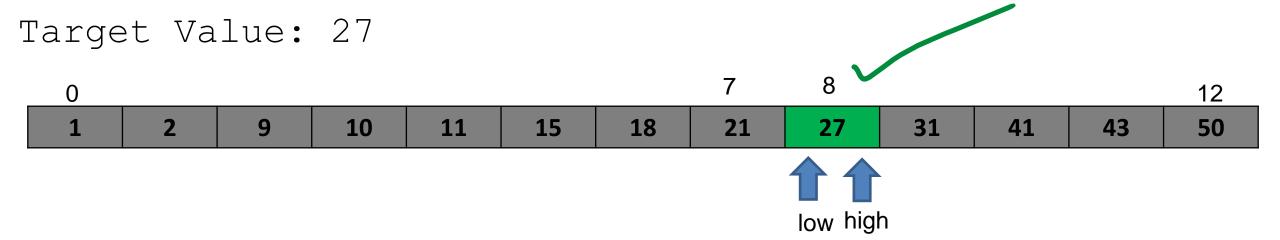
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How to calculate the mid point?



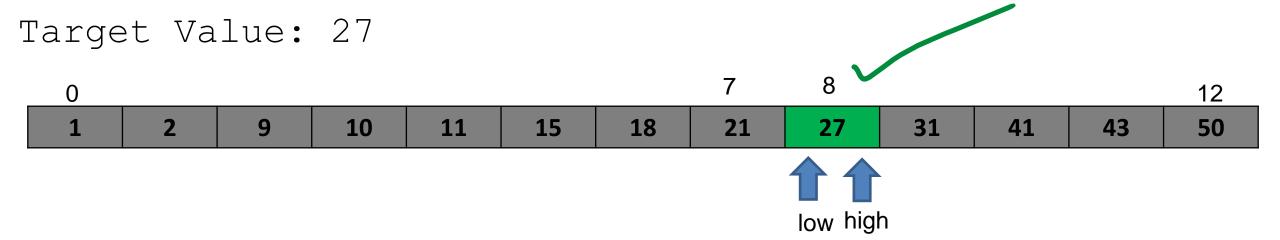
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How to calculate the mid point? (low + high) / 2



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How do we know when to stop looping?



- 1. Start at the middle of the array
- 2. Compare to target value:
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How do we know when to stop looping?

If we find the target value, or if low and high cross each other (low > high)



How do we know when to stop looping?

If we find the target value, or if low and high cross each other (low > high)

```
private static int binary_search(int[] array, int n) {
      int low = 0;
      int high = array.length - 1;
      while(low <= high) {</pre>
             int mid = (low + high) / 2;
             if(n == array[mid]) {
                    return mid;
             else if(n > array[mid]) {
                    low = mid + 1;
             else {
                    high = mid - 1;
      return -1;
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```

Running time? Each time we loop, we eliminate half the array

Initial length of array = n

Iteration 1 - Length of array =n/2

Initial length of array = n

Iteration 1 - Length of array =n/2

Iteration 2 - Length of array $=(n/2)/2=n/2^2$

Initial length of array = n

Iteration 1 - Length of array =n/2

Iteration 2 - Length of array $=(n/2)/2=n/2^2$

Iteration k - Length of array $= n/2^k$

Initial length of array = n

Iteration 1 - Length of array =n/2

Iteration 2 - Length of array $=(n/2)/2=n/2^2$

Iteration k - Length of array $=n/2^k$

After k iterations, eventually our array has been reduced to one element

Length of array
$$=n/2^k=1$$

$$n=2^k$$

"Two to what power makes n??"

After k iterations, eventually our array has been reduced to one element

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"Two to what power makes n??"

$$log_2(n) = log_2(2^k)$$

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"Two to what power makes n??"

$$log_2(n) = log_2(2^k)$$

$$log_2(n) = k * log_2 2$$

After k iterations, eventually our array has been reduced to one element

Length of array
$$=n/2^k=1$$

$$n=2^k$$

"Two to what power makes n??"

$$log_2(n) = log_2(2^k)$$

$$log_2(n) = k * log_2 2$$

$$log_2(n) = k$$

After K iterations, we will have done log(n) divisions

```
private static int binary_search(int[] array, int n) {
      int low = 0;
      int high = array.length - 1;
      while(low <= high) {</pre>
             int mid = (low + high) / 2;
             if(n == array[mid]) {
                    return mid;
             else if(n > array[mid]) {
                    low = mid + 1;
             else {
                    high = mid - 1;
      return -1;
```

Generally speaking, whenever we eliminate half of the problem each iteration, that will give us **O(logn)** running time

```
private static int binary_search(int[] array, int n) {
       int low = 0; O(1)
       int high = array.length - 1; O(1)
       while(low <= high) { O(log n)</pre>
              int mid = (low + high) / 2; O(1)
              if(n == array[mid]) { O(1)
                     return mid; O(1)
              else if(n > array[mid]) { O(1)
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              if(n == array[mid]) { O(1)
                     return mid; O(1)
              else if(n > array[mid]) { O(1)
                     low = mid + 1; O(1)
              else {
                    high = mid - 1; O(1)
       return -1; O(1)
```

Running time? O(log n)

```
private static int binary_search(int[] array, int n) {
      int low = 0;
      int high = array.length - 1;
      while(low <= high) {</pre>
             int mid = (low + high) / 2;
             int result = x.compareTo(array[mid])
             if(result = 0) {
                   return mid;
                                                We can do binary search
             else if(result > 0){
                                                on an array of Strings
                   low = mid + 1;
                                                using the compareTo()
             else {
                                                method
                   high = mid - 1;
      return -1;
```

```
private static int binary_search(??????????) {
      if(low <= high) {</pre>
             int mid = (low + high) / 2;
             if(n == array[mid]) {
                    return mid;
             else if(n > array[mid]) {
                    return binary_search(????????);
             else {
                    return binary_search(????????);
       else {
             return -1;
```

Binary Search can also be implemented using recursion