CSCI 132: Basic Data Structures and Algorithms

Recursion (Part 3)

Reese Pearsall & Illiana Castillon Fall 2024

https://www.cs.montana.edu/pearsall/classes/fall2024/132/main.html



Announcements

Lab 10 due tomorrow

Lowest Lab Grade gets dropped at the end of the semester





Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?



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D = [1, 5, 10, 25] K = 37



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Answer = 4

(Quarter, dime, two pennies)



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(Quarter, dime, two pennies)

Use as many quarters as possible, then as many dimes as possible, ...



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

D = [1, 5, 10, 25] K = 37

Answer = 4

(Quarter, dime, two pennies)

Use as many quarters as possible, then as many dimes as possible, ...

This is known as the greedy approach



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 25]

K = 37

Use as many quarters as possible, then as many dimes as possible, ...



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 18, 25] Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ... K = 37

What if there were also an 18-cent coin?



Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 18, 25] Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ... K = 37

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?



K = 37

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

D = [1, 5, 10, 18, 25] Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ...

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?

Real Answer = 18, 18, 1 (3 coins)



K = 37

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

```
D = [1, 5, 10, 18, 25]
Use as many quarters as possible, then as many 18
cent pieces as possible, then dimes , ...
```

```
25, 10, 1, 1 (4 coins)
```

What if there were also an 18-cent coin?

```
Real Answer = 18, 18, 1 (3 coins)
```

Lesson Learned: The Greedy approach works for the United States denominations, but not for a general set of denominations



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

25 + 25 + 10 + 1 + 1 + 1 = 63



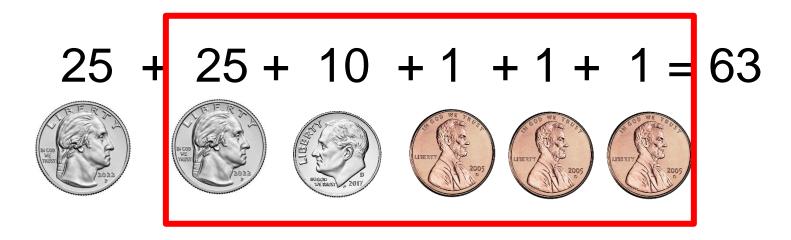
What can you conclude?

Does this provide an answer to any other change making problems?



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

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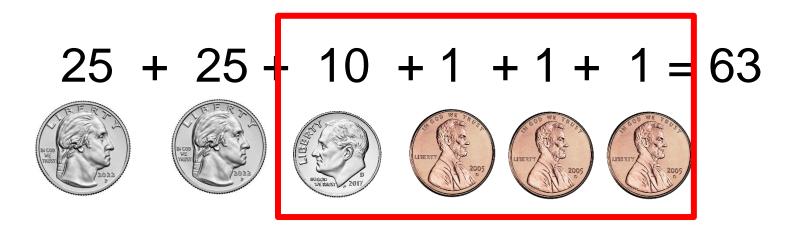


This is the minimum coins needed to make 38 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

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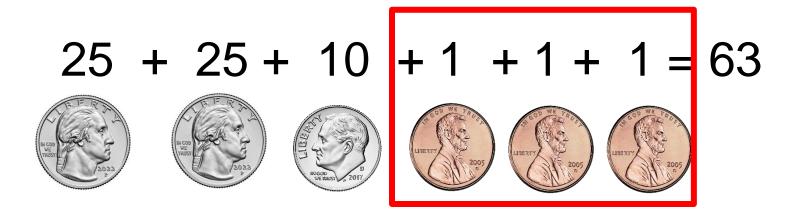


This is the minimum coins needed to make 13 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

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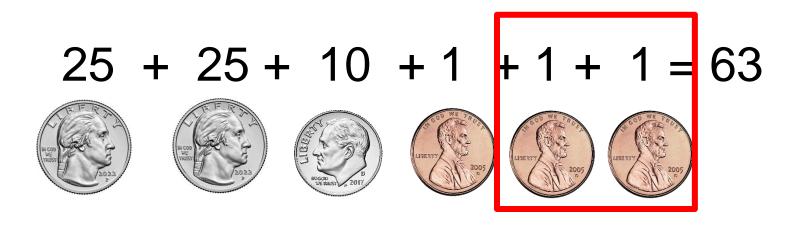


This is the minimum coins needed to make 3 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

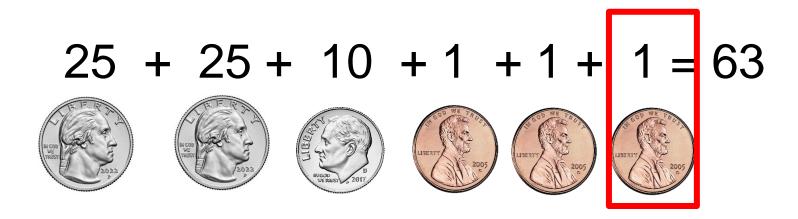


This is the minimum coins needed to make 2 cents



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)



This is the minimum coins needed to make 1 cent



Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE)

25 + 25 + 10 + 1 + 1 + 1 = 63



The solution to the change making problems consists of solutions to smaller change making problems

We can use **recursion** to solve this problem



In general, suppose a country has coins with denominations:

 $1 = d_1 < d_2 < \dots < d_k$ (US coins: $d_1 = 1, d_2 = 5, d_3 = 10, d_4 = 25$)

Algorithm: To make change for p cents, we are going to figure out change for every value x < p. We will build solution for p out of smaller solutions.



C(p) – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.



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$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$
Now find the minimum number of coins needed to make 12 cents



C(p) – minimum number of coins to make p cents.

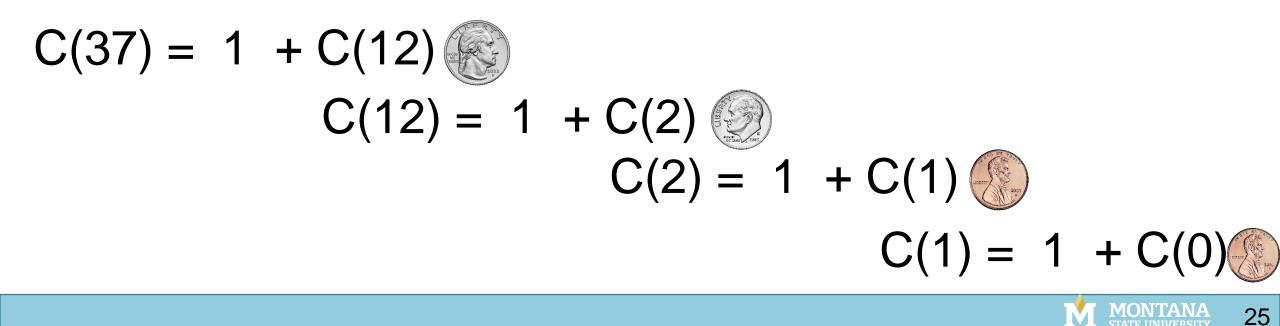
x – value (e.g. \$0.25) of a coin used in the optimal solution.

C(p) = 1 + C(p - x). C(37) = 1 + C(12)We used one dime C(12) = 1 + C(2)

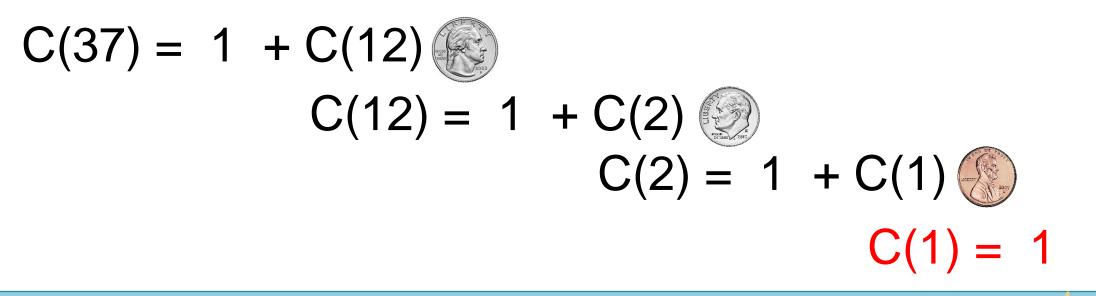
Now find the minimum number of coins needed to make 2 cents



C(p) – minimum number of coins to make p cents.

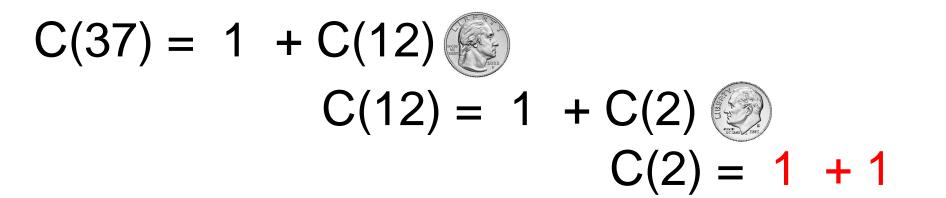


C(p) – minimum number of coins to make p cents.



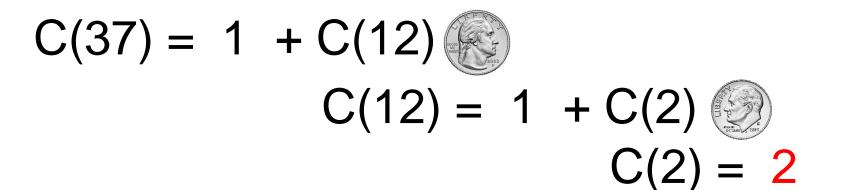


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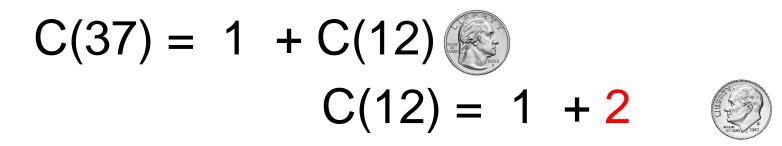


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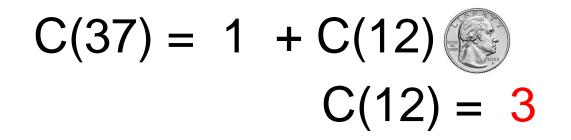


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C(37) = 1 + 3



C(p) – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution. C(p) = 1 + C(p - x).

C(37) = 4

The minimum number of coins needed to make 37 cents is 4



In general, suppose a country has coins with denominations:

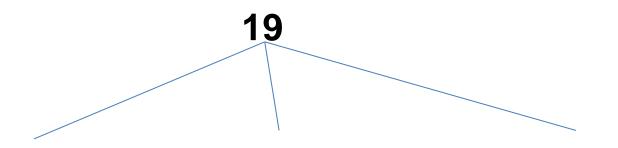
 $1 = d_1 < d_2 < \dots < d_k \qquad (\text{US coins: } d_1 = 1, d_2 = 5, d_3 = 10, d_4 = 25)$

(This algorithm must work for ALL denominations)

Algorithm: To make change for p cents, we are going to figure out change for every value x < p. We will build solution for p out of smaller solutions.

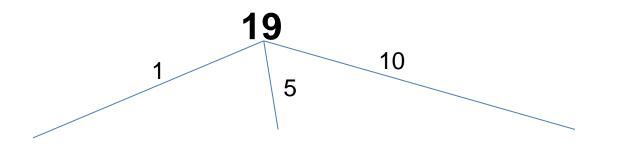


Make \$0.19 with \$0.01, \$0.05, \$0.10





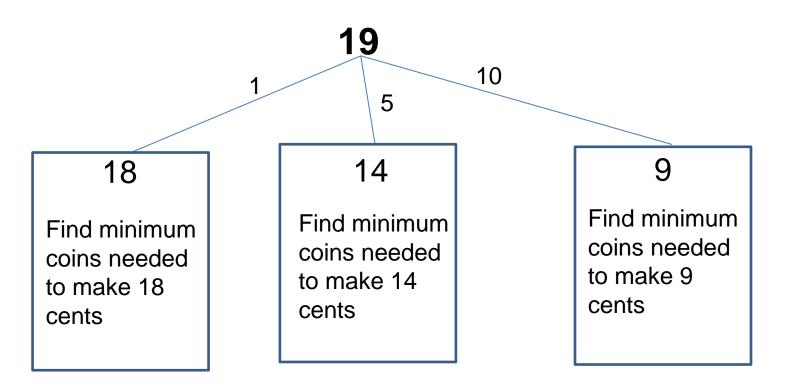
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Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations

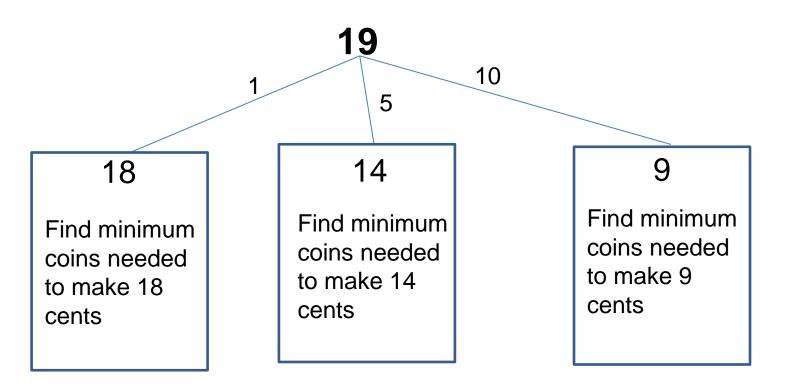


To find the minimum number of coins needed to create 19 cents, we generate **k** subproblems



Make \$0.19 with \$0.01, \$0.05, \$0.10

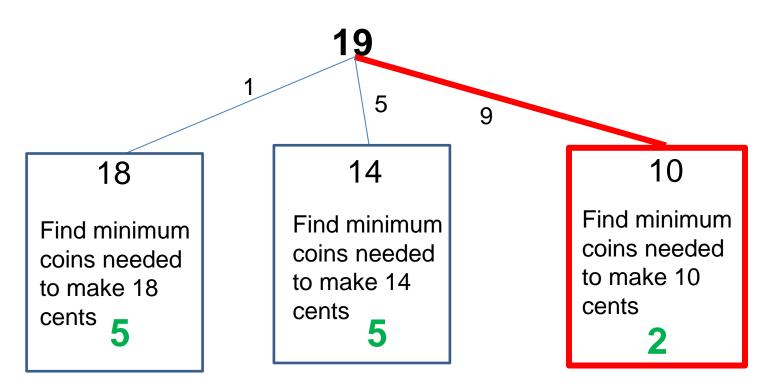
k = # denominations



We want to select the **minimum** solution of these three subproblems

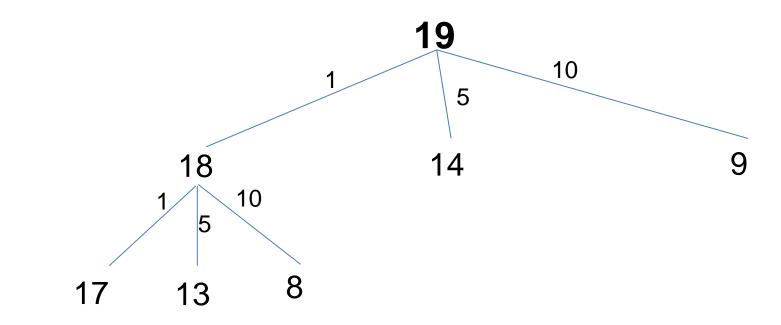






For the solution of our original problem (19), we want to select this branch (one nine cent used)

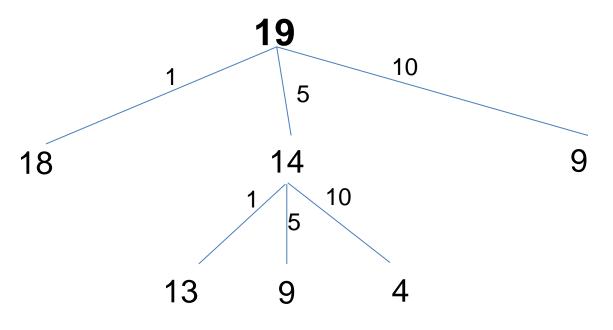
Make \$0.19 with \$0.01, \$0.05, \$0.10



Find minimum	Find minimum	Find minimum
coins needed	coins needed	coins needed
to make 17	to make 13	to make 8
cents	cents	cents



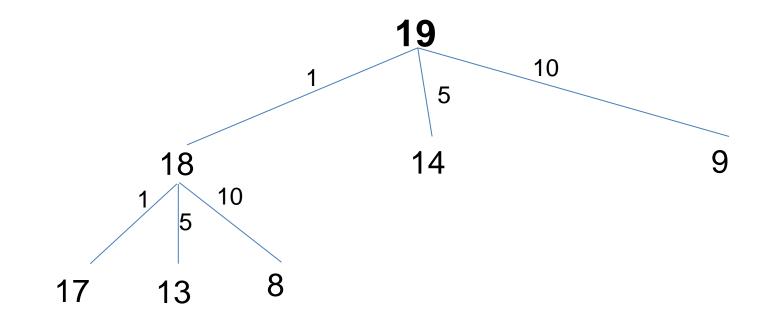
Make \$0.19 with \$0.01, \$0.05, \$0.10



Find minimum	Find minimum	Find minimum
coins needed	coins needed	coins needed
to make 13	to make 9	to make 4
cents	cents	cents

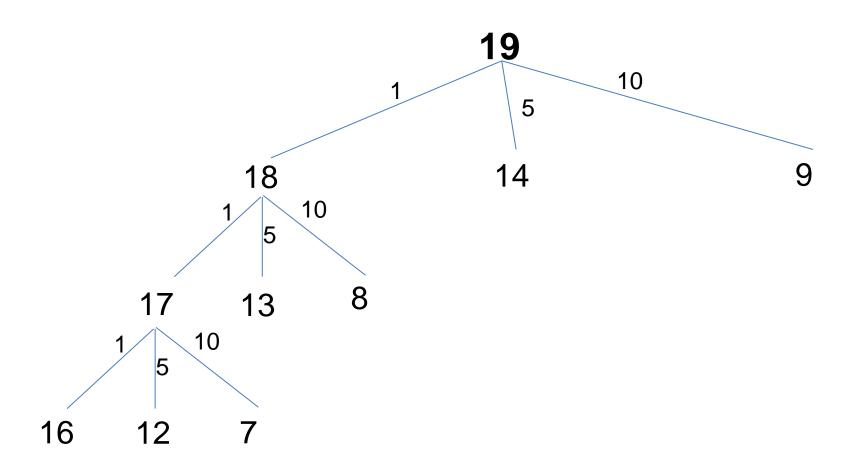


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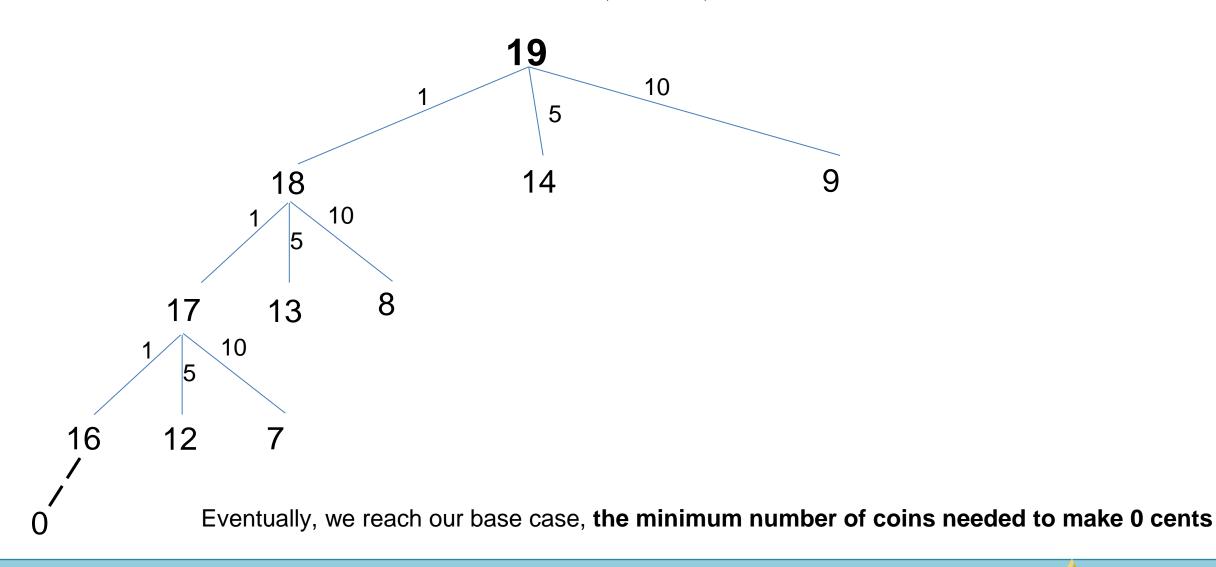


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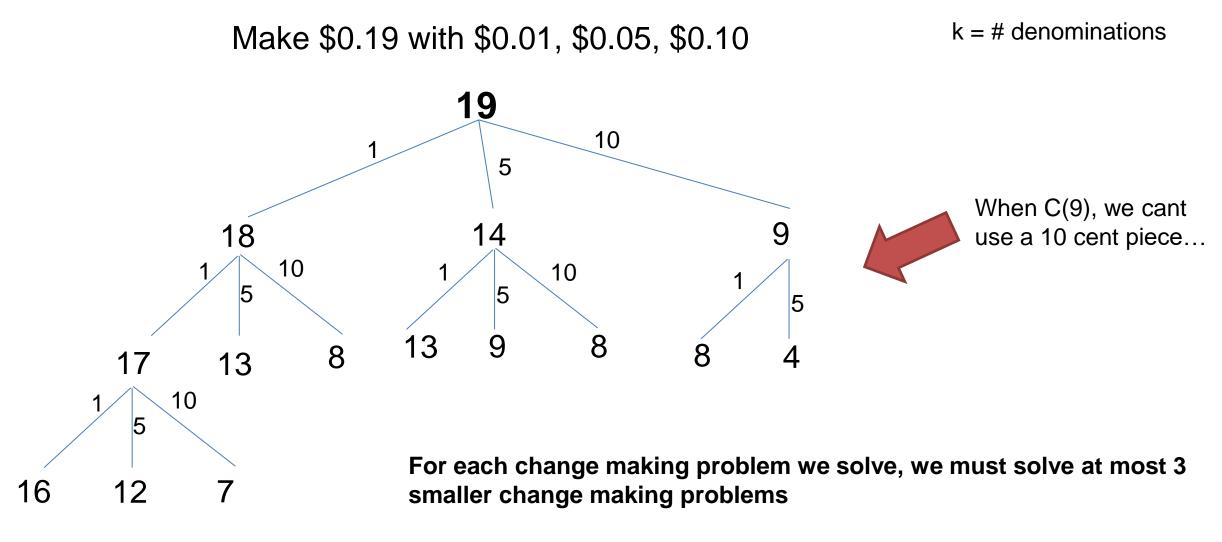




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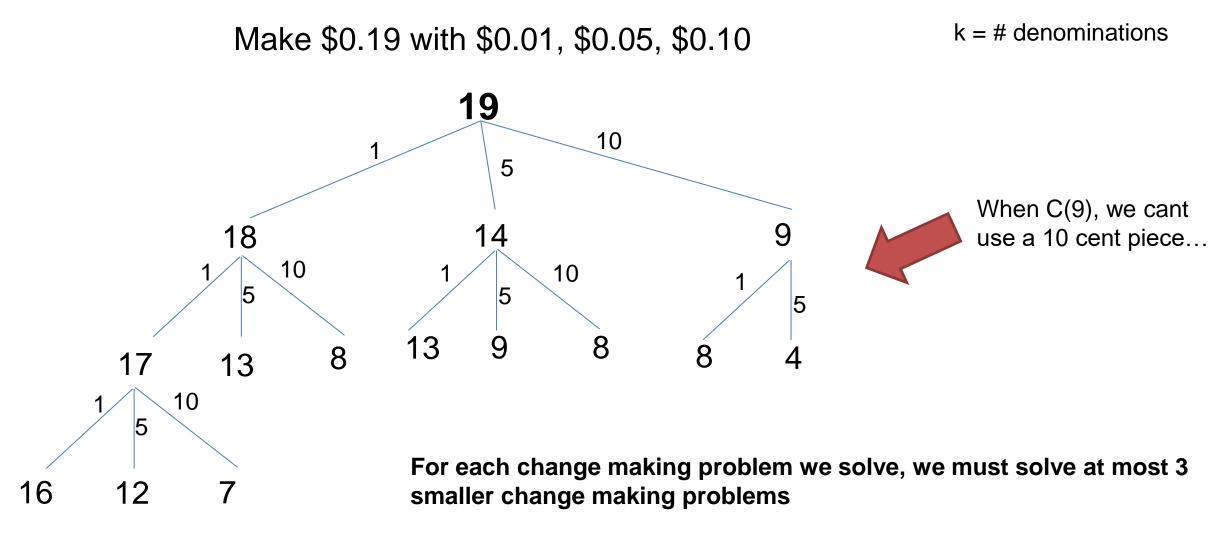






Once we solve the smaller problems, we must select the branch that has the minimum value





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$$C(p) = \begin{cases} \min_{i:d_i \le p} C(p - d_i) + 1, p > 0\\ 0, p = 0 \end{cases}$$

Least change for 19 cents = minimum of:

- least change for 19-10 = 9 cents
- least change for 19-5 = 14 cents
- least change for 19-1 = 18 cents

For each problem P, we will solve the problem for (P - d), where d represents each possible denomination



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We want to select only the branch the yields the minimum value



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If we ever need to make change for 0 cents, return 0



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D = array of denominations [1, 5, 10, 18, 25]p = desired change (37)

min_coins(D, p)



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if p == 0
 return 0;

Base Case



D = array of denominations [1, 5, 10, 18, 25] p = desired change (37)

min_coins(D, p)
if p == 0
 return 0;
else
 min = ∞
 a = ∞

Base Case
int min = Integer.MAX_VALUE;
int a = Integer.MAX_VALUE;;



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min coins(D, p) if p == 0Base Case return 0; else $min = \mathbf{0}$ int min = Integer.MAX_VALUE; int a = Integer.MAX_VALUE;; a = **00** for each d_i in D $if (p - d_i) >= 0$

 $a = min_coins(D, p - d_i)$

Recurse, and find the minimum number of coins needed using each valid denomination

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MONTANA 55

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min_coins(D, p)		
if p == 0	Base Case	
return 0;		
else	· · · · · · · · · · · · · · · · · · ·	
min = ၸ	<pre>int min = Integer.MAX_VALUE; int a = Integer.MAX_VALUE;;</pre>	
a = co	Int a - Integer MAN_MALOL,	

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Recurse, and find the minimum number of coins needed using each valid denomination

Select the branch that has the minimum value

return 1 + min

Once, our for loop finishes, we should know the branch that had the minimum, so return (1 + min), 1 because one coin was used in the current method call



min coins(D, p) if p == 0return 0; else min = 🚥 a = **co** for each d_i in D $if (p - d_i) >= 0$ $a = min coins(D, p - d_i)$ if a < min min = a

```
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```

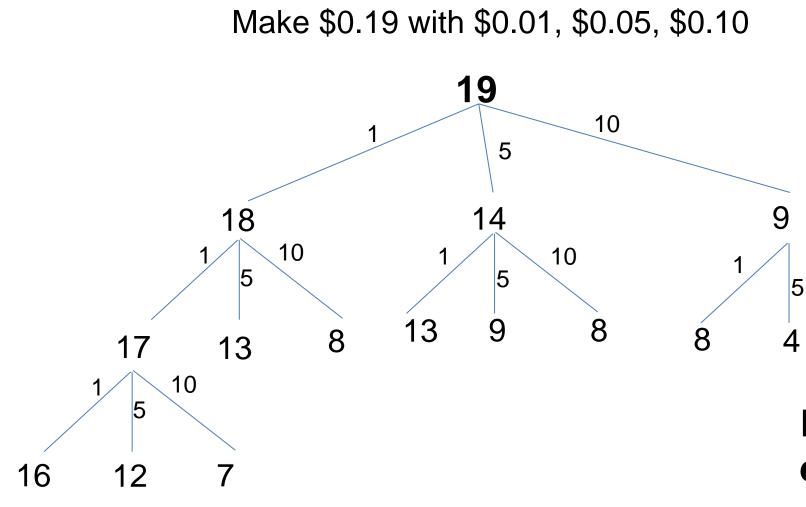


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return 1 + min



Running time?

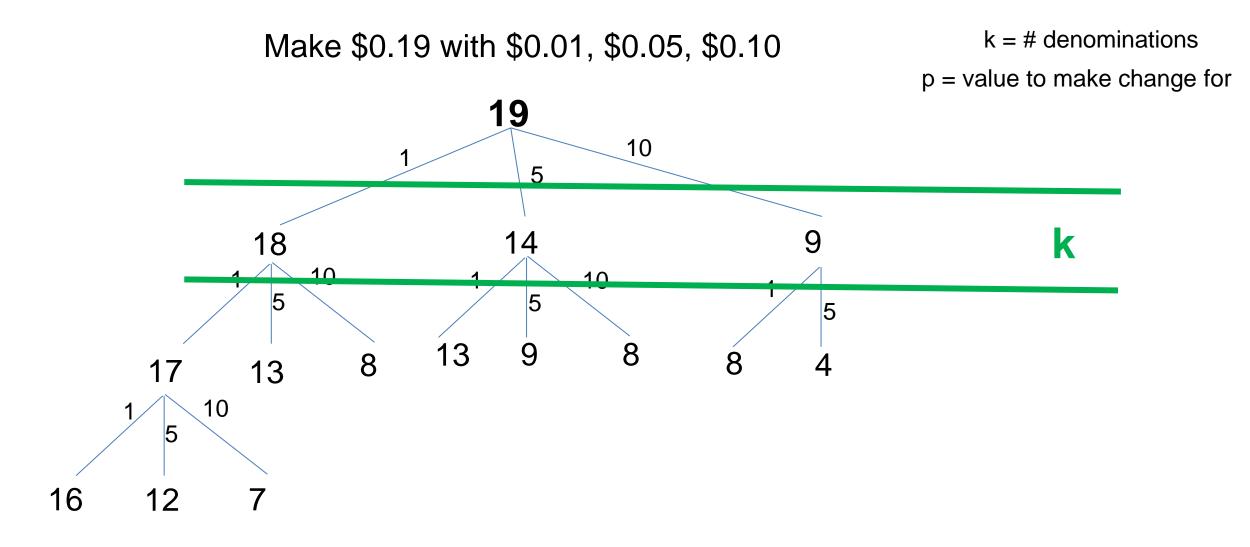


k = # denominations

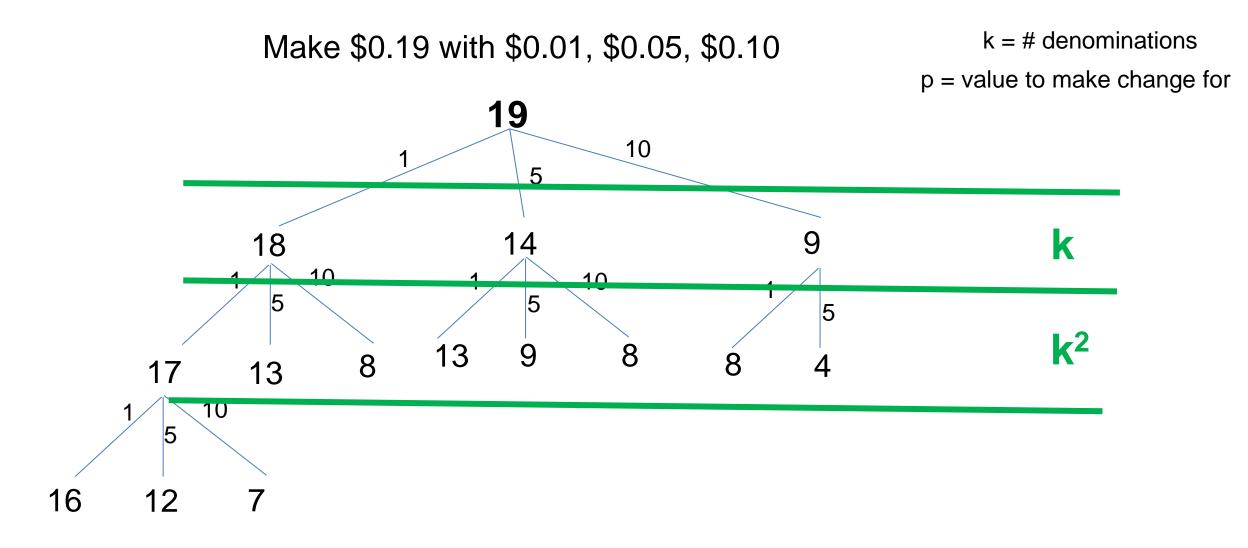
p = value to make change for

For sufficiently large p, every permutation of denominations is included.

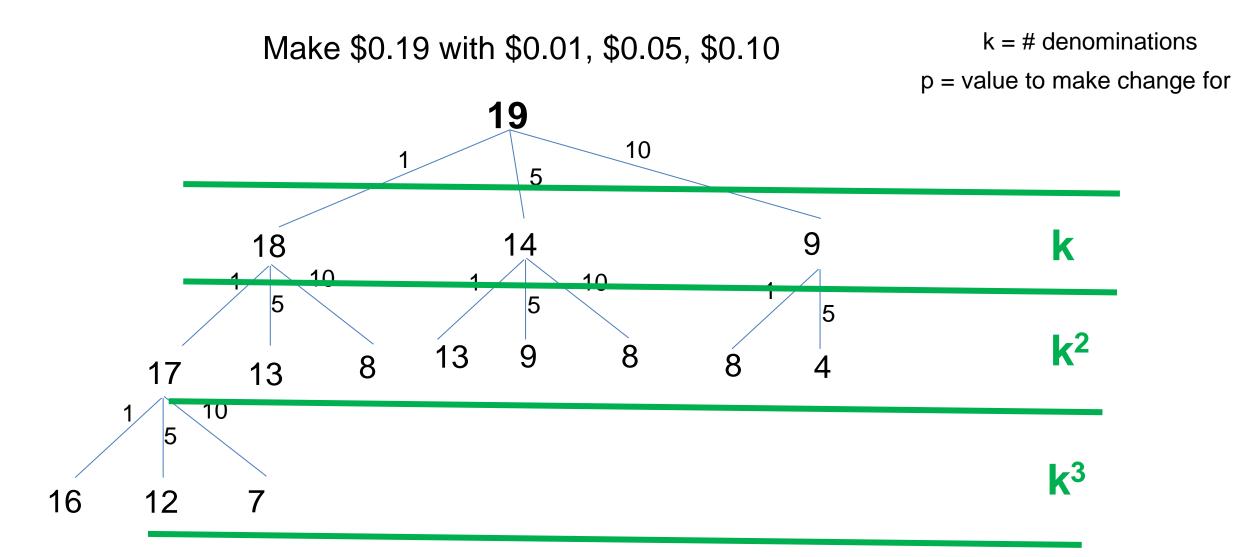




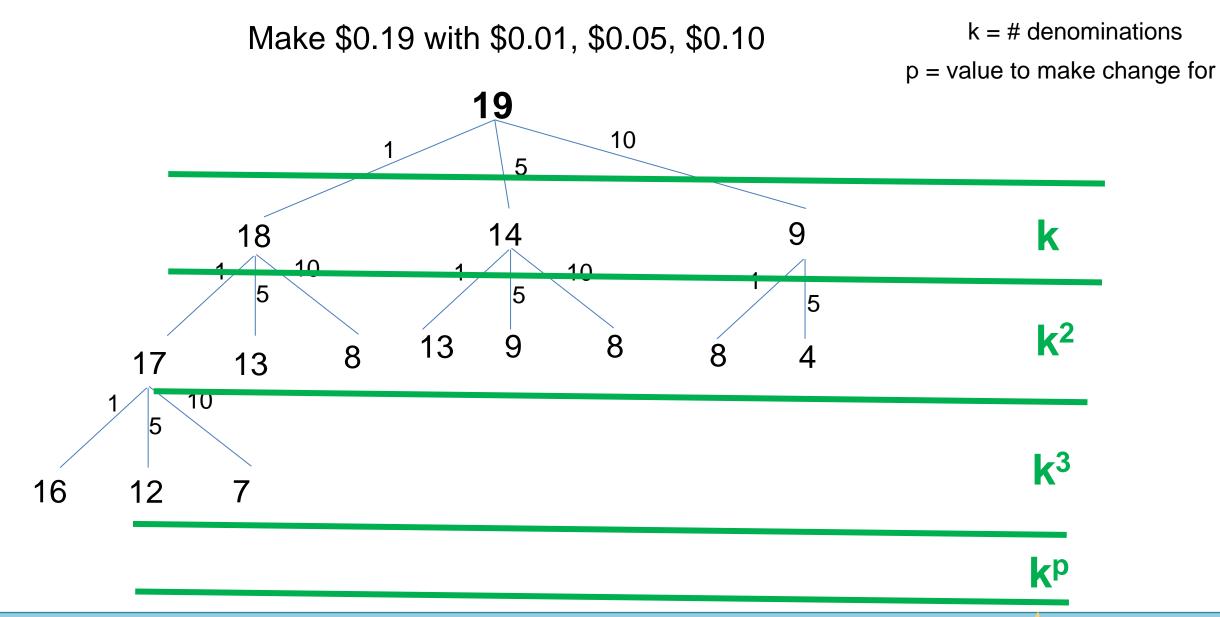














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Make \$0.19 with \$0.01, \$0.05, \$0.10



p = value to make change for

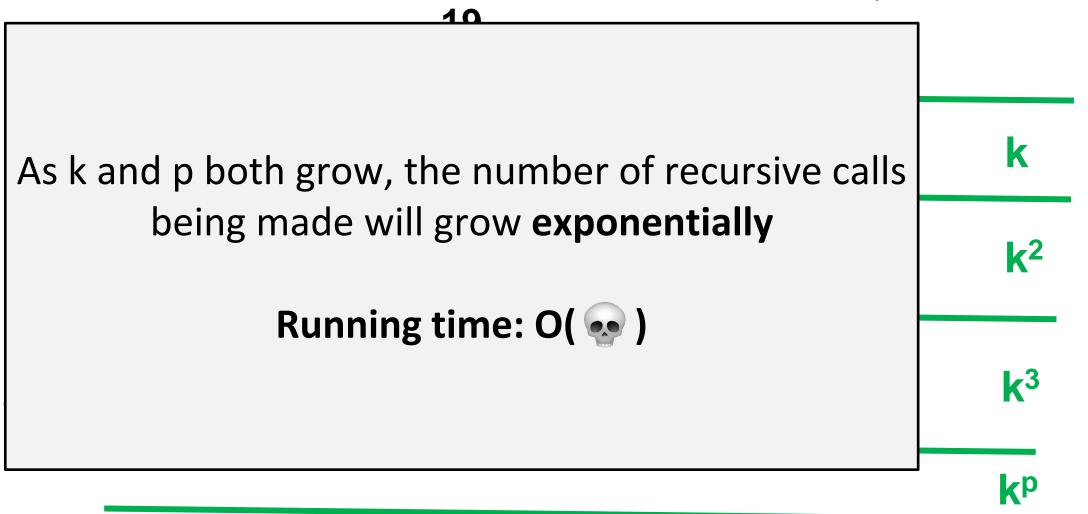
10 As k and p both grow, the number of recursive calls k being made will grow **exponentially k**² If we have a lot of coin denominations, we will have a lot of branching **k**³ kp







p = value to make change for

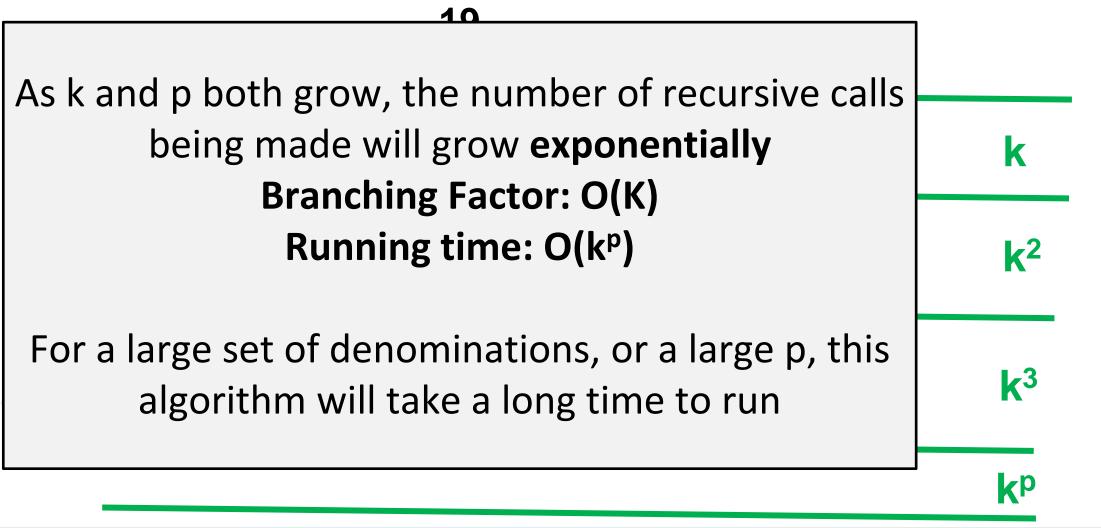




Make \$0.19 with \$0.01, \$0.05, \$0.10



p = value to make change for



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Combinations for 7 cents Using D = [1, 5, 10]

[1, 1, 5] [1, 1, 1, 1, 1, 1, 1]

[1, 1, 5] and [5, 1, 1] is the same combination...

Permutations for 7 cents Using D = [1, 5, 10]

Order does not matter. [1, 1, 5] and [5, 1, 1] are considered different solutions

In our change making algorithm, we are calculating every possible permutation (bad)



Let's try 81 cents!





This algorithm returns the minimum number of coins needed (ie 4), but it does not tell us what coins were used in that solution



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D = Array of coin denominations [1, 5, 10, 25]p = value to make change forn = minimum number of coins used to make p cents

Goal: Find an **n-length** combination of coins from **D** that were used to make **p**



This algorithm returns the minimum number of coins needed (ie 4), but it does not tell us what coins were used in that solution

D = Array of coin denominations [1, 5, 10, 25]p = value to make change forn = minimum number of coins used to make p cents

Goal: Find an **n-length** combination of coins from **D** that were used to make **p**

To do this, we will compute **all n-length combinations**, but only return the combinations that add up to be **p** (not very efficient)



Denominations (D) = [1, 5, 10]

For **n=3**, these are the combinations to be generated:

- [1, 1, 1]
- [1, 1, 5]
 [1, 1, 10]

• [1, 5, 10]

- [1, 1, 10]
 [1, 5, 5]
- [5, 10, 10]
 [10, 10, 10]

• [1, 10, 10]

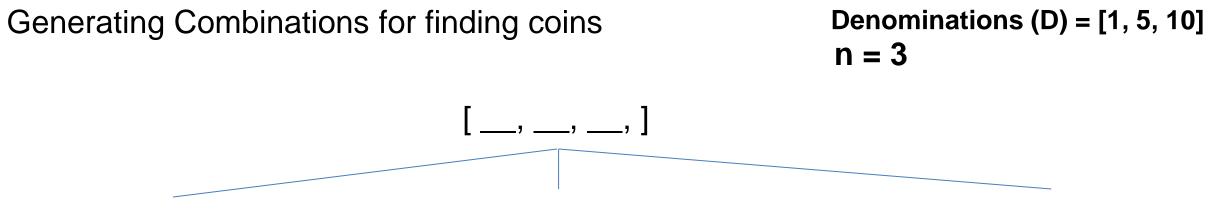
• [5, 5, 5]

• [5, 5, 10]

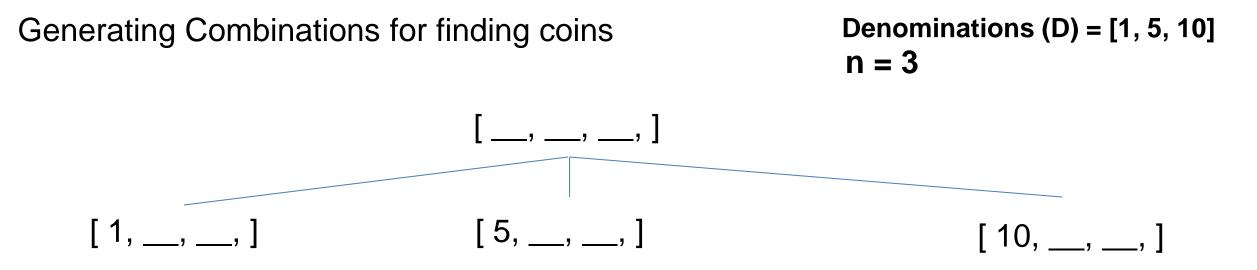
Note:

[5, 1, 1] is not a "valid" combination, because it is the same thing as [1, 1, 5]

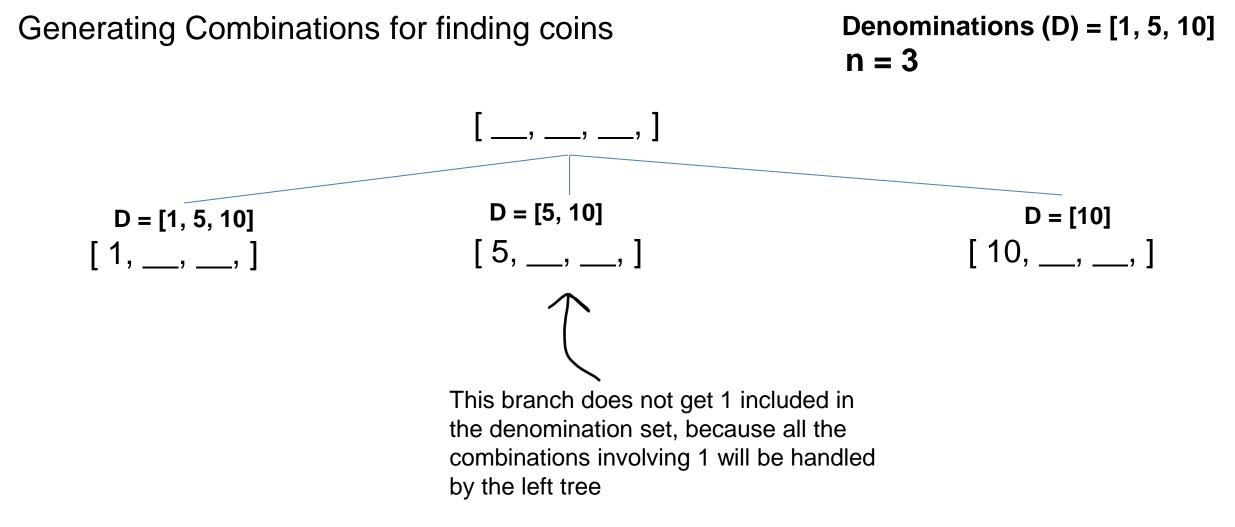




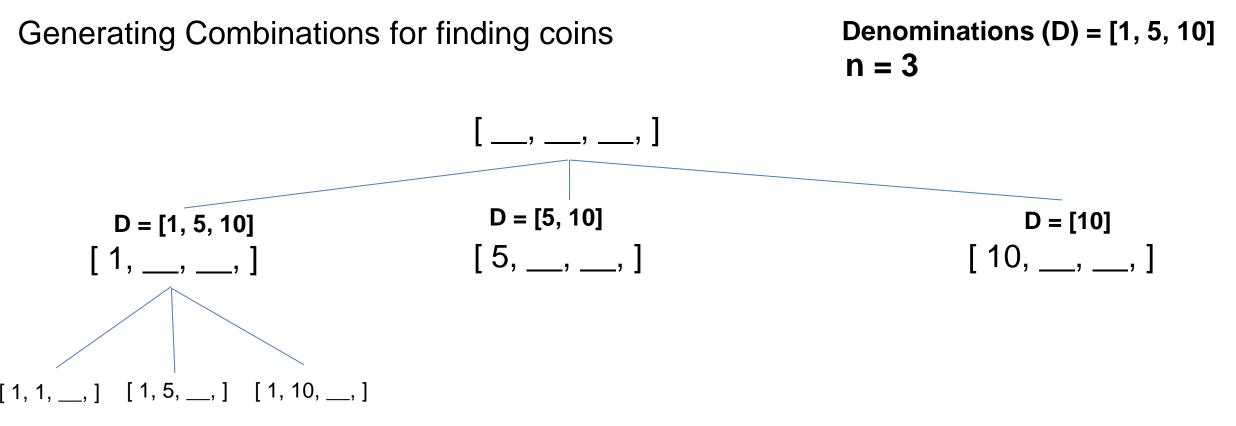




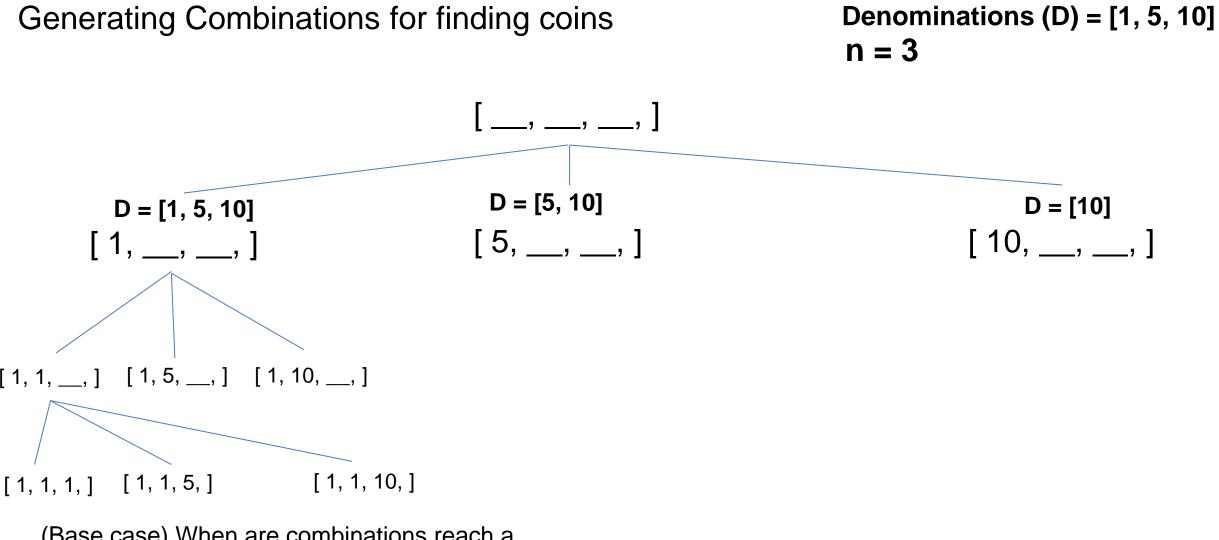








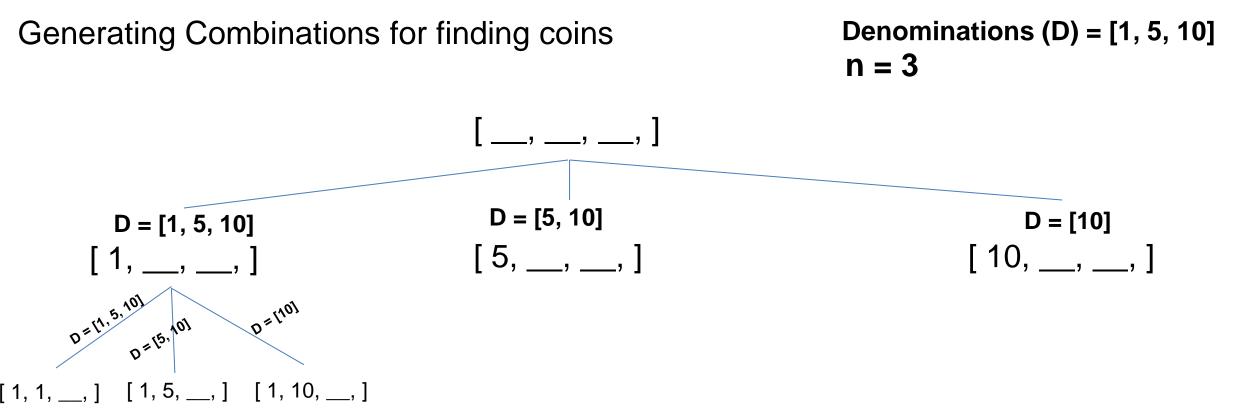




(Base case) When are combinations reach a length of 3, we will stop recursing

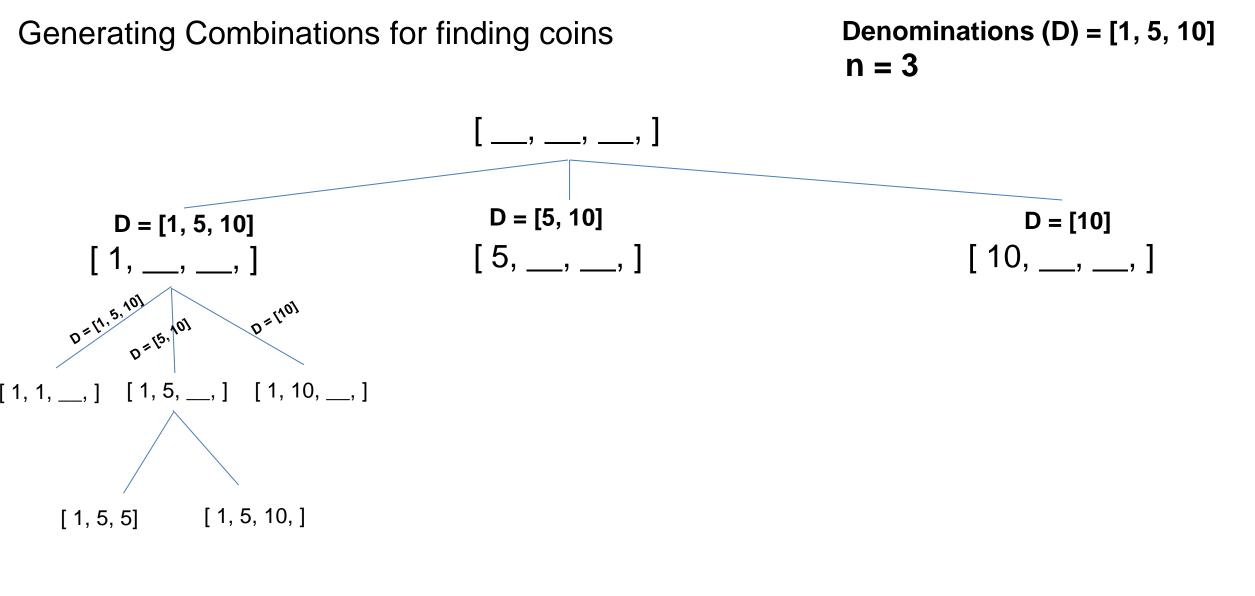
[1, 1, 1] [1, 1, 5] [1, 1, 10]





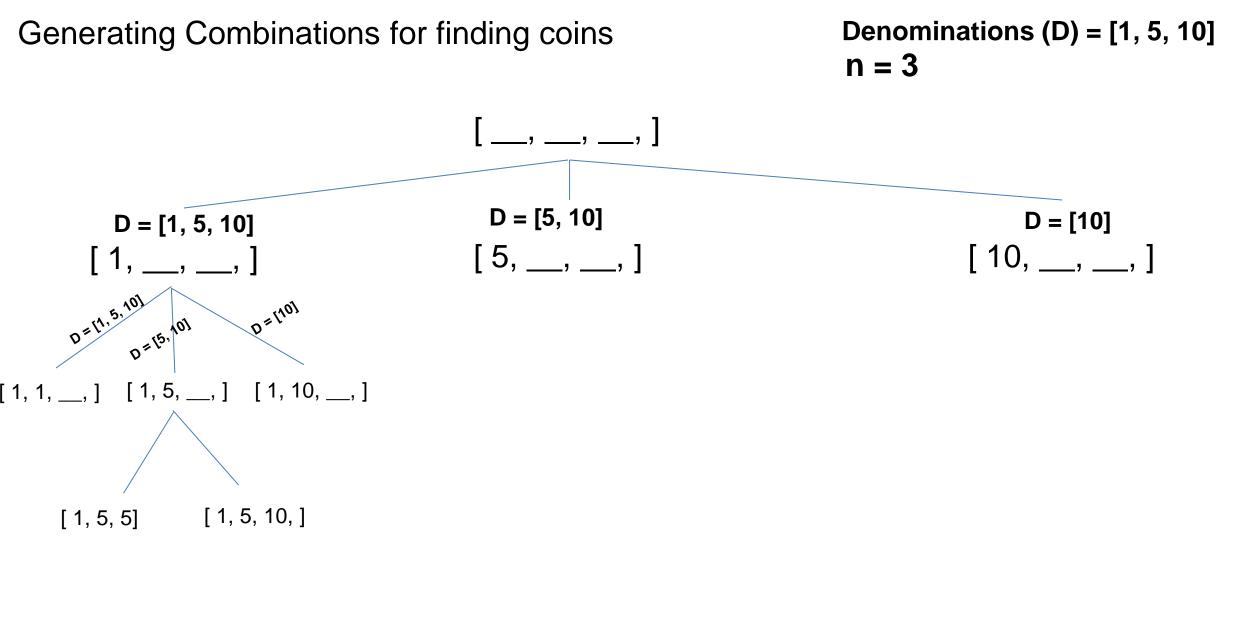
[1, 1, 1] [1, 1, 5] [1, 1, 10]





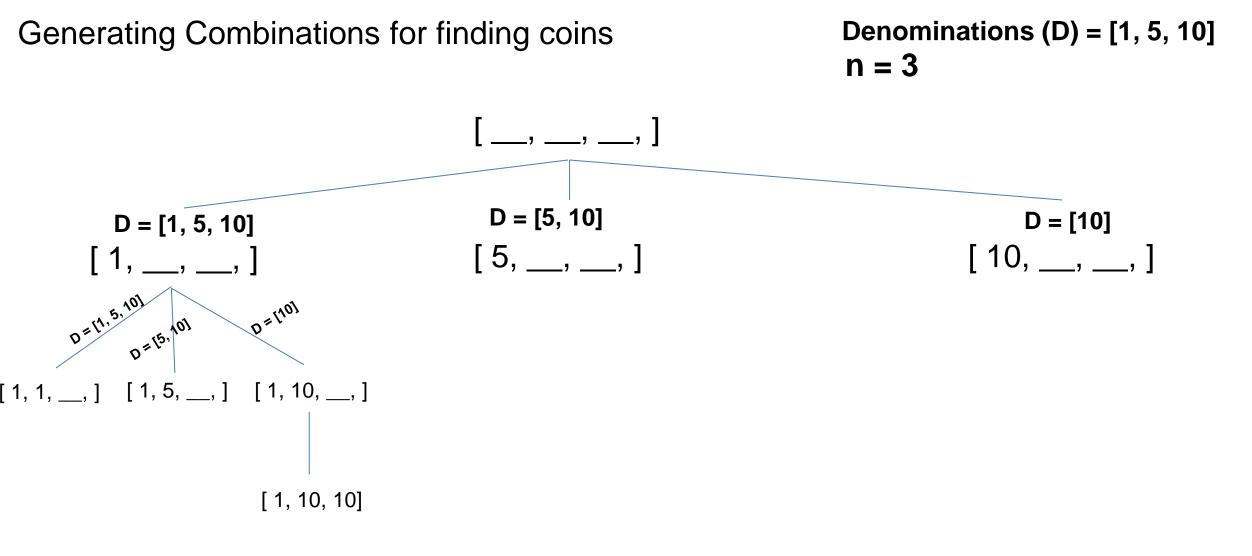
[1, 1, 1] [1, 1, 5] [1, 1, 10]





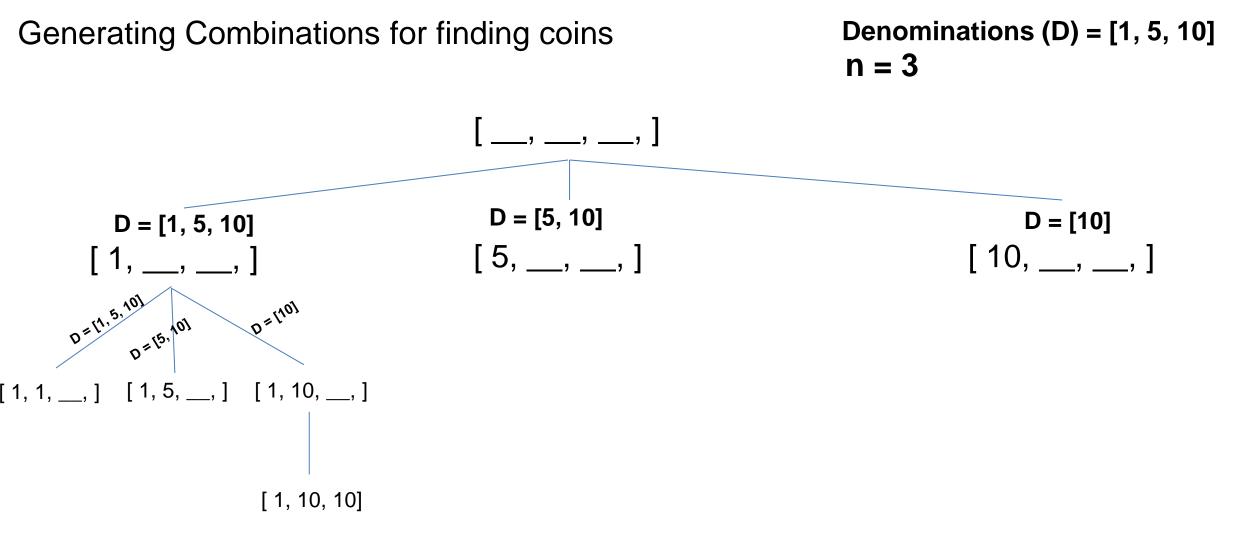
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10]





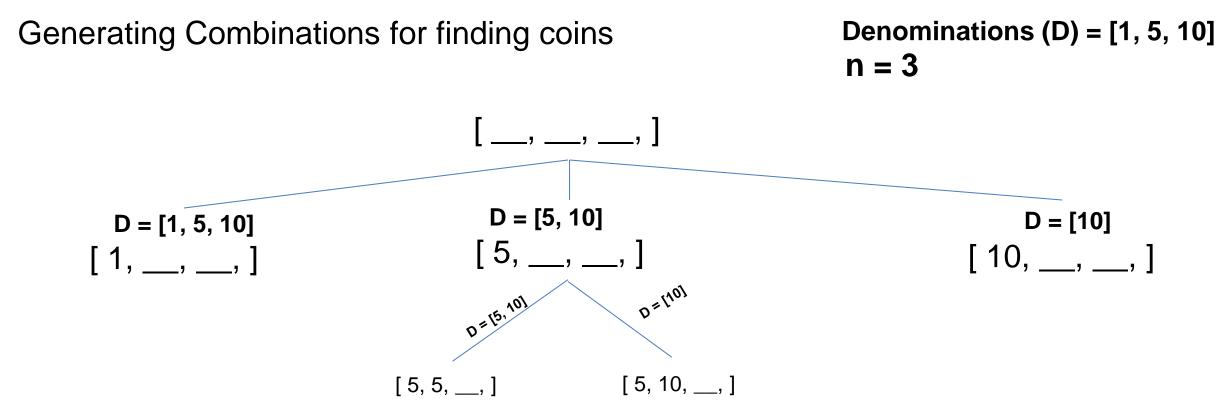
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10]





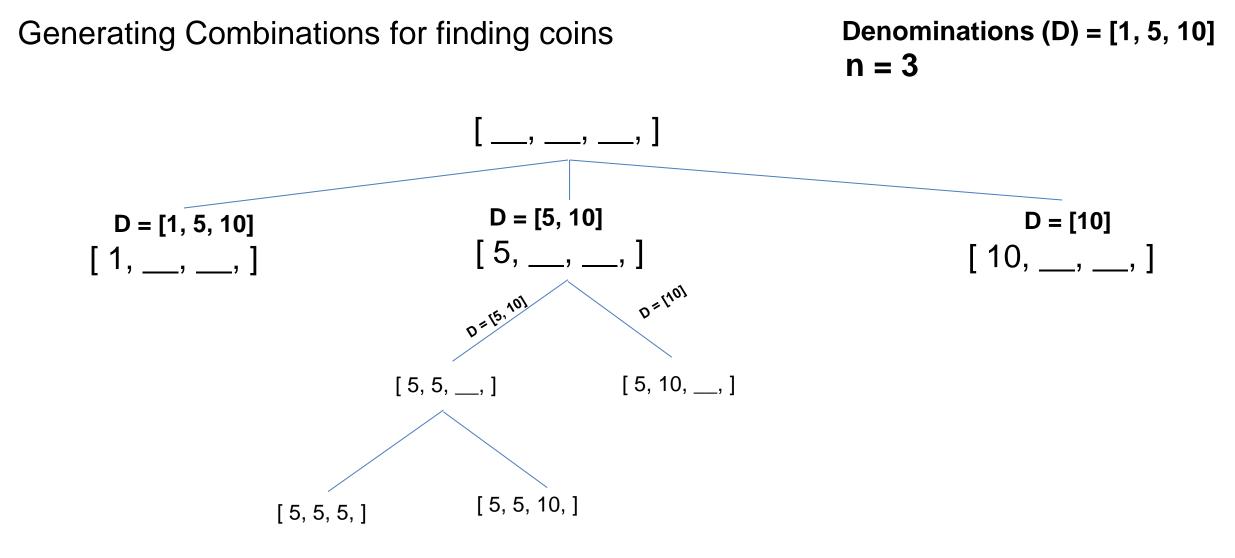
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]





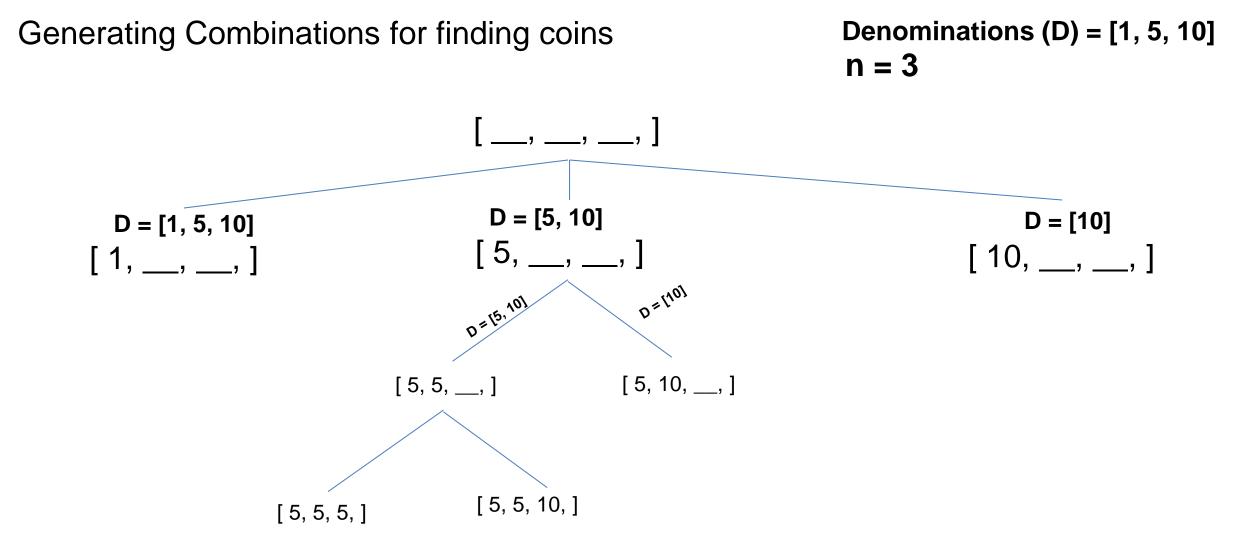
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]





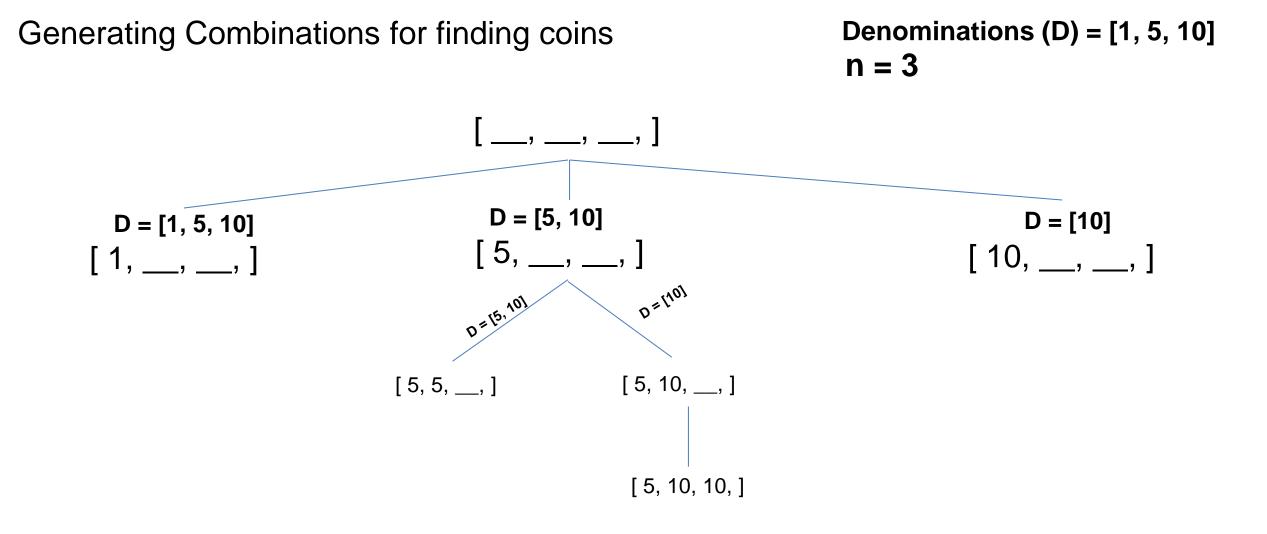
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]





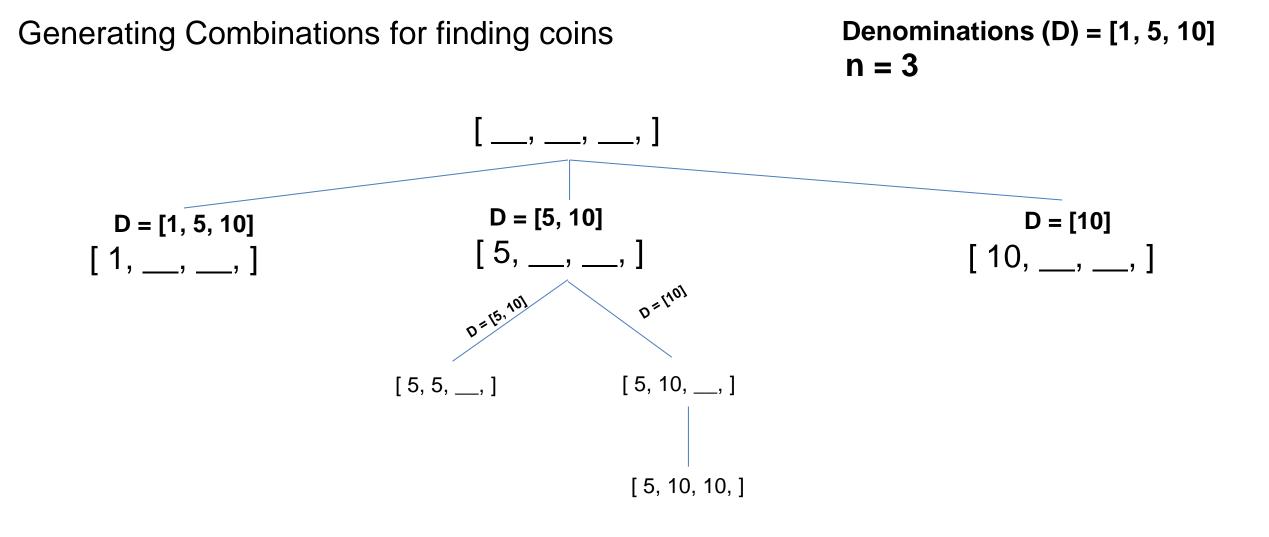
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10]





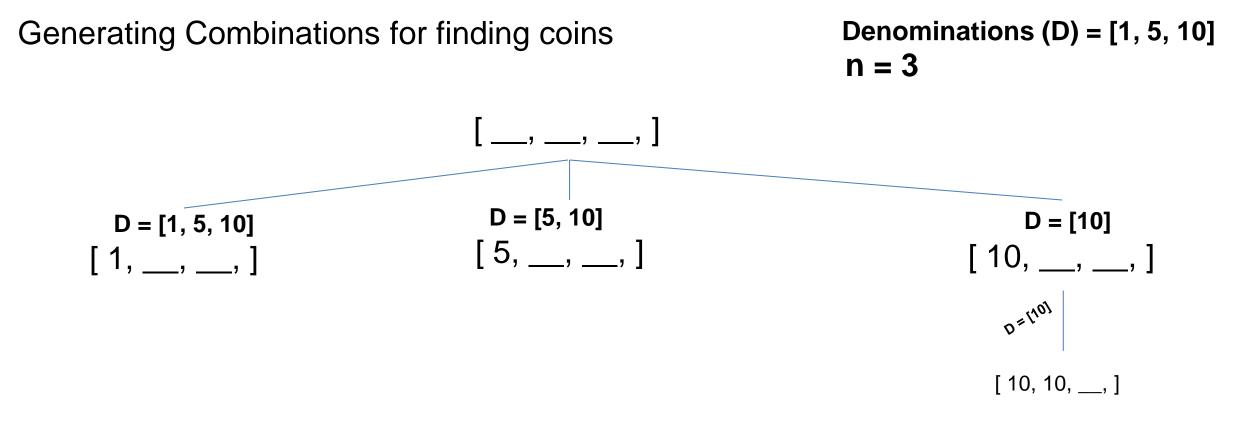
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10]





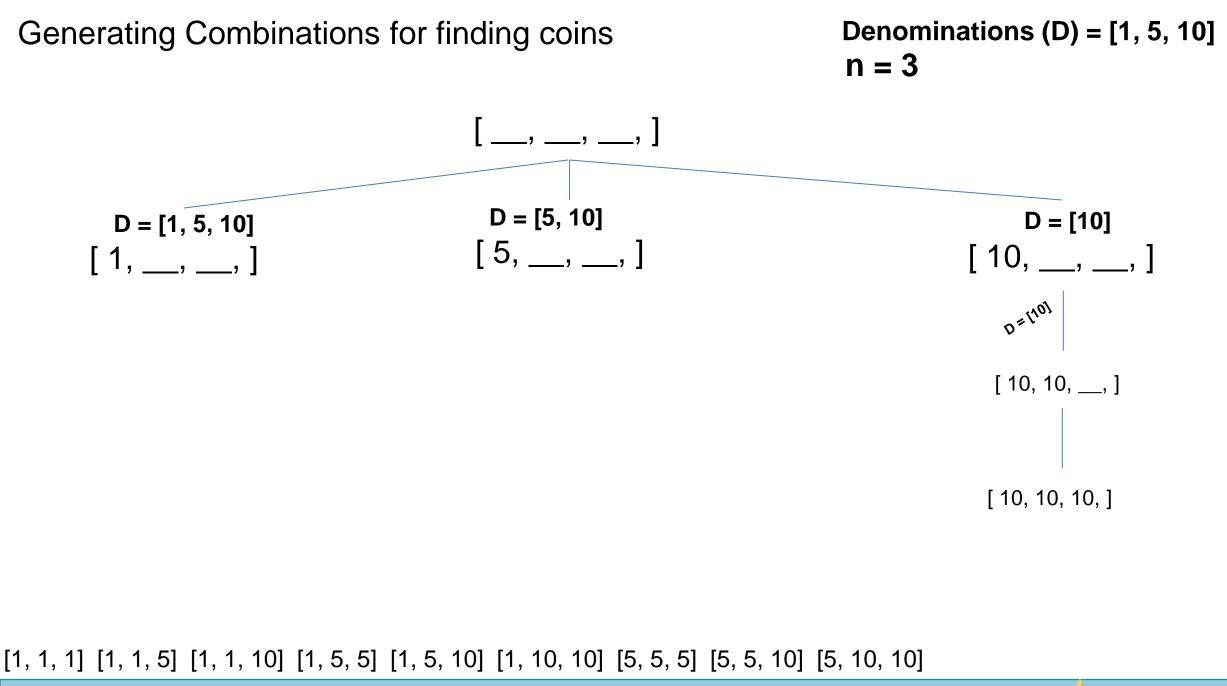
[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10] [5, 10, 10]



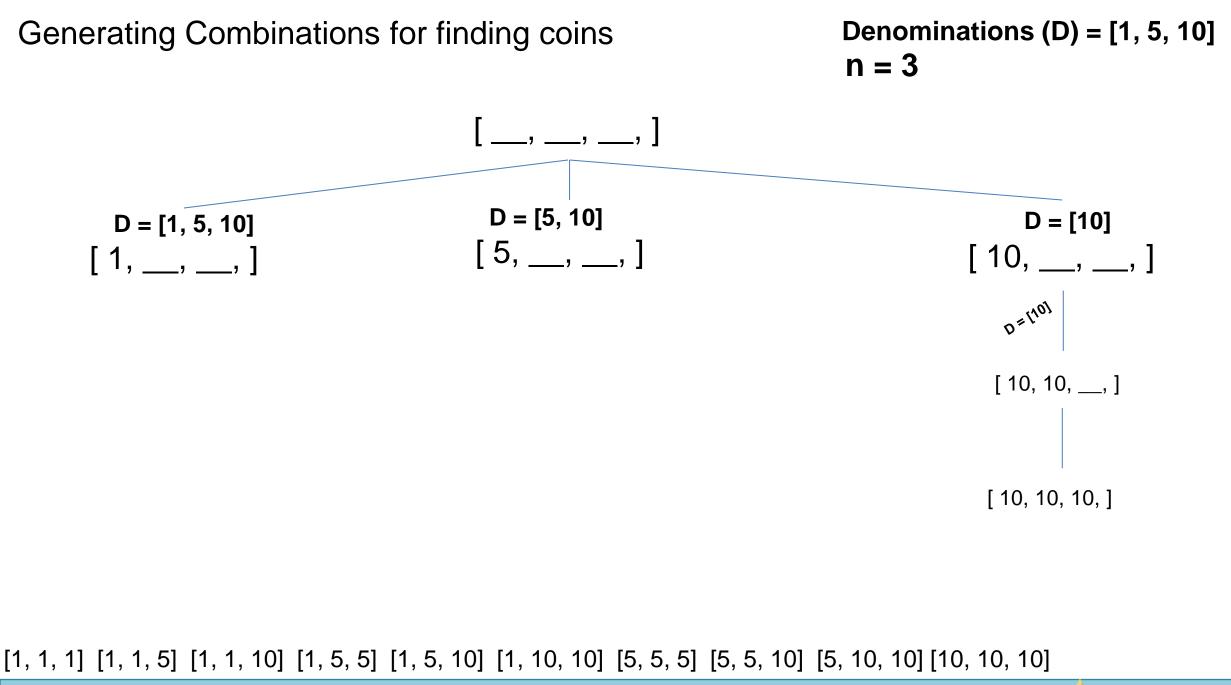


[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10] [5, 10, 10]





MONTANA 90



MONTANA 91

Denominations (D) = [1, 5, 10] n = 3

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]



We've generated all combinations of length 3



Denominations (D) = [1, 5, 10] n = 3

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]

We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to ${\bf K}$



Denominations (D) = [1, 5, 10] n = 3 K = 16

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]



We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to K

Suppose K = 16 (a minimum of 3 coins is needed to make 16 cents)



Denominations (D) = [1, 5, 10] n = 3 K = 16

1. [1, 1, 1] 2. [1, 1, 5] 3. [1, 1, 10] 4. [1, 5, 5] 5. [1, 5, 10] 6. [1, 10, 10] 7. [5, 5, 5] 8. [5, 5, 10] 9. [5, 10, 10] 10.[10, 10, 10]



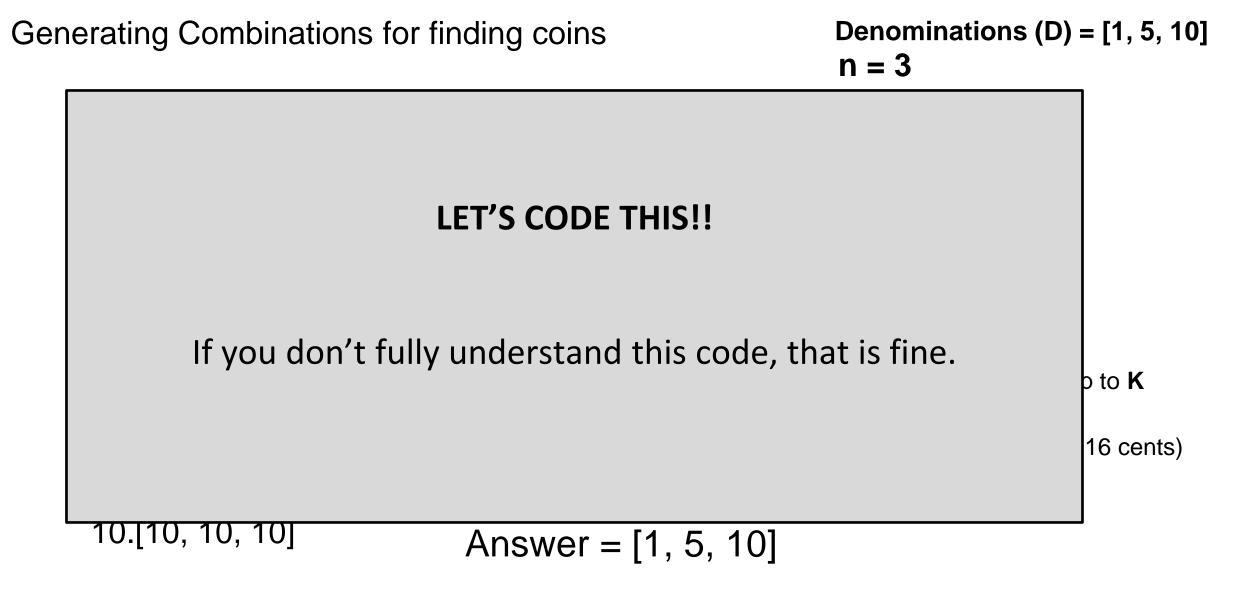
We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to K

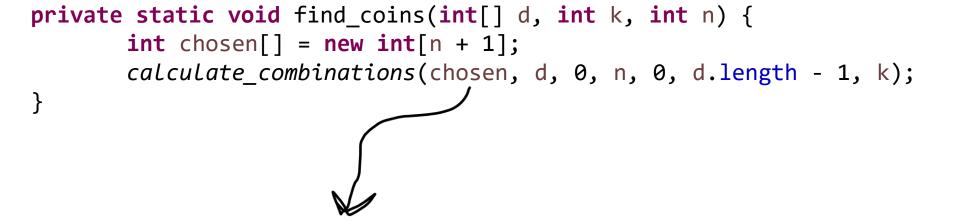
Suppose K = 16 (a minimum of 3 coins is needed to make 16 cents)

Answer = [1, 5, 10]









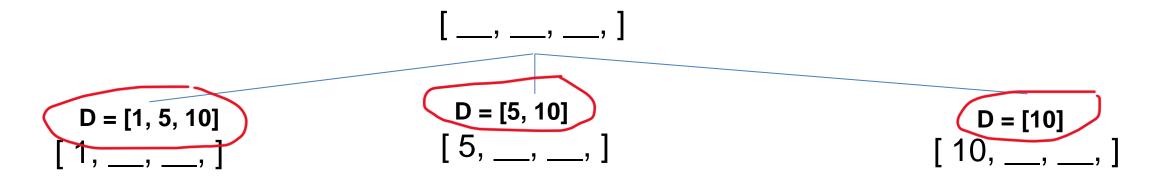
Array that we build up over time. Holds indices of currently selected denominations for some combination

[1,___, ___] [1,1, ___] [1,1, 5]



```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```

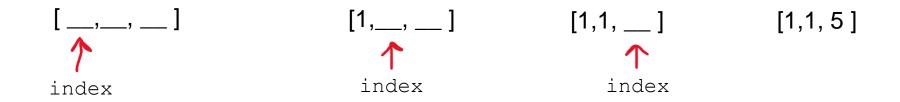
Array of denominations we pass for each recursive call





```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```

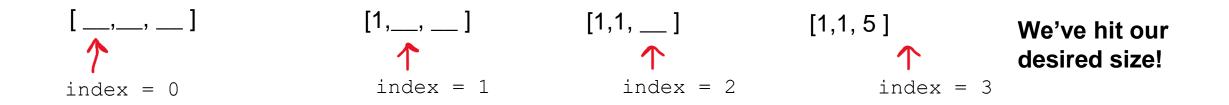
The next index that we need to insert at for chosen array





```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```

The desired size of the combination. When index == r, we have reached the desired combination size





```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
       if (index == r) {
               int counter = 0;
               ArrayList<Integer> coins = new ArrayList<Integer>();
               for (int i = 0; i < r; i++) {</pre>
                       counter += arr[chosen[i]];
                       coins.add(arr[chosen[i]]);
               if(counter == target) {
                       System.out.println(coins);
               return;
       for (int i = start; i <= end; i++) {</pre>
               chosen[index] = i;
               calculate combinations(chosen, arr, index + 1, r, i, end, target);
        }
       return;
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
       if (index == r) {
                                                                              If we hit our base
                int counter = 0;
                                                                              case, we know we
                ArrayList<Integer> coins = new ArrayList<Integer>();
                                                                              have N things, so
  Base case
                for (int i = 0; i < r; i++) {</pre>
                                                                              put them in an
                        counter += arr[chosen[i]];
                                                                              ArrayList and add
                        coins.add(arr[chosen[i]]);
                                                                              them up
                                                       Only print out the combination if it
                if(counter == target) { 
                                                       adds up to target
                        System.out.println(coins);
                return;
        for (int i = start; i <= end; i++) {</pre>
                chosen[index] = i;
                calculate combinations(chosen, arr, index + 1, r, i, end, target);
        return;
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
        if (index == r) {
                int counter = 0;
                ArrayList<Integer> coins = new ArrayList<Integer>();
                for (int i = 0; i < r; i++) {</pre>
                         counter += arr[chosen[i]];
                         coins.add(arr[chosen[i]]);
                if(counter == target) {
                         System.out.println(coins);
                return;
Recursive Case
        for (int i = start; i <= end; i++) {
    chosen[index] = i;</pre>
                calculate_combinations(chosen, arr, index + 1, r, i, end, target);
        return:
```

Otherwise, insert selected coin into the chosen array

create (end-start) branches, and give it a smaller section of D



