CSCI 132: Basic Data Structures and Algorithms

Lessons Learned so far + Importing Linked Lists

Reese Pearsall

Spring 2024

https://www.cs.montana.edu/pearsall/classes/spring2024/132/main.html



Announcements

Come get your midterm exam

- Exam average was 83
- Don't stress if you didn't do well
- Make sure I calculated your score correctly

Lab 7 due tomorrow

Class Registration





Lab 7



Class Registration



Next Class to Register for:

CSCI 232- Data Structures and Algorithms (Offered in Summer 2024 and Fall 2024)

Other classes that may be of interest:

CS 145 – Web Design CSCI 112- Programming in C

CSCI 215 – Social and Ethical issues in Computer Science CSCI 204 – Multimedia Dev Methods (Game Design) CSCI 291- Introduction to Data Science (NEW)

Term:	2024	Fall Sen	nestei	r	``	•	
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Days:	Mon	Tues	Wed	Thur	Fri	Sat	Sun
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Resources Available to you

Student Success Center - Spring 2024

Tutoring Schedule - Barnard Hall 259: Monday, January 22nd - Friday, May 3rd

Schedule	Monday	Tuesday	Wednesday	Thursday	Friday
8:00 a.m.					
9:00 a.m.			Alex Ellingsen	Muzhou Chen	Kaden Bach
10:00 a.m.	Sultan Yarylgassimov	Ruby Martin Katie Harmon		Muzhou Chen	Gerard Shu Fuhnwi
11:00 a.m.	Sultan Yarylgassimov	Riley Slater Sage James	Jack Ruder	Nicholas Addotey Ryan Johnson	Gerard Shu Fuhnwi
Noon	Asibul Islam Shahnaj Mou	Riley Slater	Jack Ruder	Nicholas Addotey	Jared Matury Matthew Phillips
1:10 p.m.	Jake Rivers	Joshua Bowen	Karishma Rahman	Muhammad Bhatti	
2:10 p.m.	Angelo Porcello Gideon Popoola	Racquel Bowen Muhammad Arju	Gideon Popoola Karishma Rahman	Nishu Nath Muhammad Bhatti	
3:10 p.m.	Angelo Porcello Brayden Miller	Muhammad Arju Justin Mau	Shama Maganur Fatima Ododo	Nishu Nath	
4:10 p.m.		Justin Mau	Shama Maganur Fatima Ododo		

CS Tutoring Center: Barnard Hall 259



Teaching Assistants/Graders

- Section 005- Fatima Ododo
- Email: fatima.ododo@student.montana.edu
- Office Hours: Wednesdays 3:10 5:10 PM in Barnard Hall 259
- Section 004- Shama Maganur
- · Email: shama.maganur@gmail.com
- Office Hours: Wednesdays 3:10 5:10 PM in Barnard Hall 259
- Section 003- Shama Maganur
- Email: shama.maganur@gmail.com
- Office Hours: Wednesdays 3:10 5:10 PM in Barnard Hall 259
- Section 006- Fatima Ododo
- Email: fatima.ododo@student.montana.edu
- Office Hours: Wednesdays 3:10 5:10 PM in Barnard Hall 259

TA Office Hours and Email

Smarty Cats Tutoring

Reese Pearsall Email: reese.pearsall@montana.edu Office Hours: Monday and Wednesday 1:00 - 2:00 PM, Tuesday and Friday 12:10 - 1:00 PM Office: Barnard Hall 361 Discord: @reese_p	My email, Discord, office hours
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Big-O

Big-O notation is a way to describe the running-time/time complexity of an algorithm regarding the number of operations that are executed in the algorithm (in relation to some input **n**)

• Focus on worst-case scenario, and how the algorithm scales as n gets really big



Big-O

Big-O notation is a way to describe the running-time/time complexity of an algorithm regarding the number of operations that are executed in the algorithm (in relation to some input n)

• Focus on worst-case scenario, and how the algorithm scales as n gets really big

A very powerful computer and a very weak computer running the same algorithm will both execute the same number of operations (the speed at which they execute these operations will be different)

Takeaway: the asymptotic running time (the big-o running time) will be the same for each computer



8

Big-O notation is a way to describe the running-time/time complexity of an algorithm regarding the number of operations that are executed in the algorithm (in relation to some input n)

• Focus on worst-case scenario, and how the algorithm scales as n gets really big

To find the total running time of an algorithm, we calculate the runningtime of each operation in the algorithm and then add everything together

In Big-O, we can drop non-dominant factors and multiplicative constants (coefficients)

O(n) + O(n) + O(n): Total running time = $O(3n) \in O(n)$ Where n is _____



Big-O

























Can hold one data type

• Can also store objects, which allow for multiple data types

Entire array is stored at a **contiguous** spot in memory

Nodes in the linked list can hold multiple data types

Linked list nodes are stored at **non-contiguous** spots in memory









Traversing a linked list requires more work than traversing an array

Both data structures can grow dynamically, and new elements can be added, but they way they add new elements is **drastically** different

Create a brand-new array, copy everything over from old array

Update pointers

Create a brand-new array, copy everything over from old array O(n)

Update pointers

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Takeaway: Adding a new element to an ArrayList requires much more work than adding a new element to a Linked List

Arrays are generally much easier to sort than Nodes in a Linked List

If you are constantly needing to add new elements to the data structure, using a Linked List requires much less work in the long run

Arrays are more memory efficient (adding is not very memory efficient though)

When to use each data structure?

It depends on how you are using your data and if you know how much data you have

If you don't know how much data you need to store, or if you are constantly needing to add new elements to the data structure \rightarrow Linked Lists

If you know how much data you need to store, and if you can add all your data at once \rightarrow Arrays/ArrayLists

These two data structures are <u>implementations</u> of a **List** Abstract Data Type (ADT)

ADT is a class whose behavior is defined by a set of operations and how a user interacts with it.

A list data type must be able to **get** an element, **add** an element, **remove** an element, etc \rightarrow How they do these operations is up to the subclass (LL and AL)

As programmers, we use handy methods that were written by other people that allows us to use these data structures

Car car1 = new Car("Ferrari","Black");

Car car1 = <u>new</u> Car("Ferrari","Black");

Car car1 = <u>new</u> Car("Ferrari","Black"); Car car2 = <u>new</u> Car("Toyota","Blue");


```
Car car1 = <u>new</u> Car("Ferrari","Black");
Car car2 = <u>new</u> Car("Toyota","Blue");
Car car3 = car1;
```



```
Car car1 = <u>new</u> Car("Ferrari","Black");
Car car2 = <u>new</u> Car("Toyota","Blue");
Car car3 = car1;
```

The new keyword is not used, so a new object is **not** created. Instead, it will point to the same object that car1 is pointing to


```
Car car1 = <u>new</u> Car("Ferrari", "Black");
Car car2 = <u>new</u> Car("Toyota","Blue");
Car car3 = car1;
```

car3.set_color("Red");


```
Car car1 = <u>new</u> Car("Ferrari","Black");
Car car2 = <u>new</u> Car("Toyota","Blue");
Car car3 = car1;
```

```
car3.set_color("Red");
System.out.println(car1.getColor());
```



```
Car car1 = <u>new</u> Car("Ferrari","Black");
Car car2 = <u>new</u> Car("Toyota","Blue");
Car car3 = car1;
```

car3.set_color("Red");
System.out.println(car1.getColor());

Red

```
Abstract Classes are restricted classes that cannot be used to create objects. To access it, it must be inherited from another class.
```

```
public abstract class Employee {
```

```
Employee e = new Employee("Sally", 4444, 123456);
```

You cannot create instances of an abstract class.

```
Accountant kevin = new Accountant("Kevin Malone", 4444, 42000, 'C');
```

Instead, we use objects from another class that inherits from the abstract class

We will no longer be writing our own Linked List class, instead we will now import the Java-provided Linked List Class

import java.util.LinkedList;

We will no longer be writing our own Linked List class, instead we will now import the Java-provided Linked List Class

```
import java.util.LinkedList;
```

```
LinkedList<String> names = new LinkedList<String>();
The data type the
linked list will be
holding
```


The Linked List Class

https://docs.oracle.com/javase/7/docs/api/java/util/LinkedList.html

public class LinkedList<E>
extends AbstractSequentialList<E>
implements List<E>, Deque<E>, Cloneable, Serializable

Doubly-linked list implementation of the List and Deque interfaces. Implements all optional list operations, and permits all elements (including null).

All of the operations perform as could be expected for a doubly-linked list. Operations that index into the list will traverse the list from the beginning or the end, whichever is closer to the specified index.

The **documentation** describe how the LinkedList class was implemented, and all the methods/operations we can do with the Linked List class

Methods	
Modifier and Type	Method and Description
boolean	add(E_e) Appends the specified element to the end of this list.
void	add(int index, E element) Inserts the specified element at the specified position in this list.
boolean	addAll(Collection extends E c) Appends all of the elements in the specified collection to the end of this list, in the order that they are returned by the sp
boolean	addAll(int index, Collection extends E c) Inserts all of the elements in the specified collection into this list, starting at the specified position.
void	addFirst(E_e) Inserts the specified element at the beginning of this list.
void	addLast(E_e) Appends the specified element to the end of this list.
void	clear() Removes all of the elements from this list.
Object	clone() Returns a shallow copy of this LinkedList.
boolean	contains(Object o) Returns true if this list contains the specified element.
Iterator <e></e>	descendingIterator() Returns an iterator over the elements in this deque in reverse sequential order.
E	element() Retrieves, but does not remove, the head (first element) of this list.
E	get(int index) Returns the element at the specified position in this list.
E	getFirst() Raturne the first alament in this list

when you start coding in a new language without reading the documentation:

The Linked List Class

```
import java.util.LinkedList;
public class march20demo {
    public static void main(String[] args) {
        LinkedList<String> names = new LinkedList<String>();
        names.add("Reese");
        names.add("Spencer");
        names.add("Susan");
        System.out.println(names);
    }
```

