CSCI 132: Basic Data Structures and Algorithms

Queues (Array Implementation)

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Spring 2024

https://www.cs.montana.edu/pearsall/classes/spring2024/132/main.html



Announcements

Lab 9 due **tomorrow** 4/2 @ 11:59 PM (Queues)

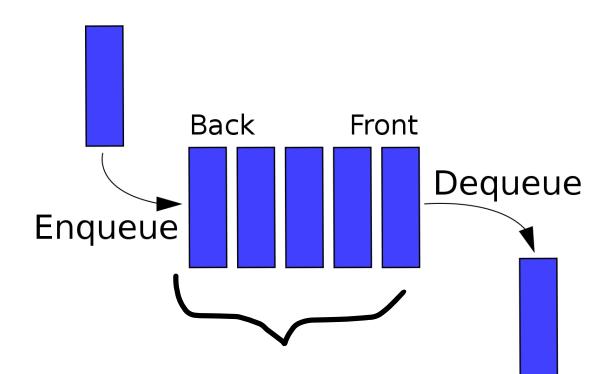
No office hours tomorrow ©

Program 3 due this Friday 4/5





A **Queue** is a data structure that holds data, but operates in a First-in First-out (**FIFO**) fashion



Once again, we need a data structure to hold the data of the queue

- Linked List
- Array

Elements get added to the **Back** of the Queue.

Elements get removed from the **Front** of the queue





A **Queue** is a data structure that holds data, but operates in a First-in First-out (**FIFO**) fashion

The Queue ADT has the following methods:

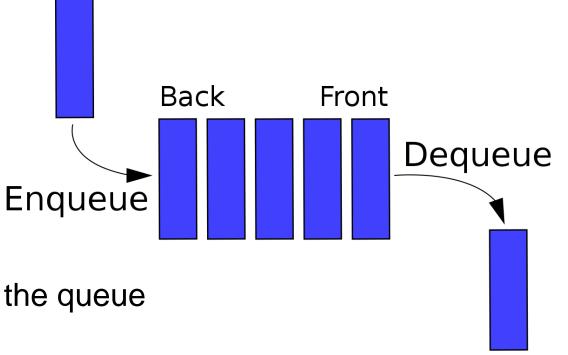
Enqueue- Add new element to the queue

Dequeue- Remove element from the queue

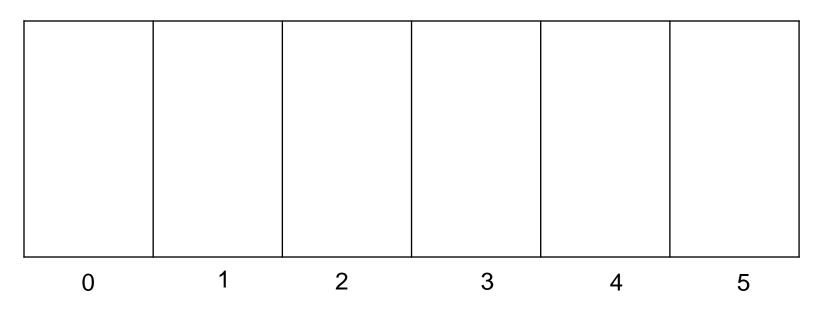
** Always remove the front-most element

Peek()- Return the element that is at the front of the queue

IsEmpty() – Returns true if queue is empty, returns false is queue is not empty

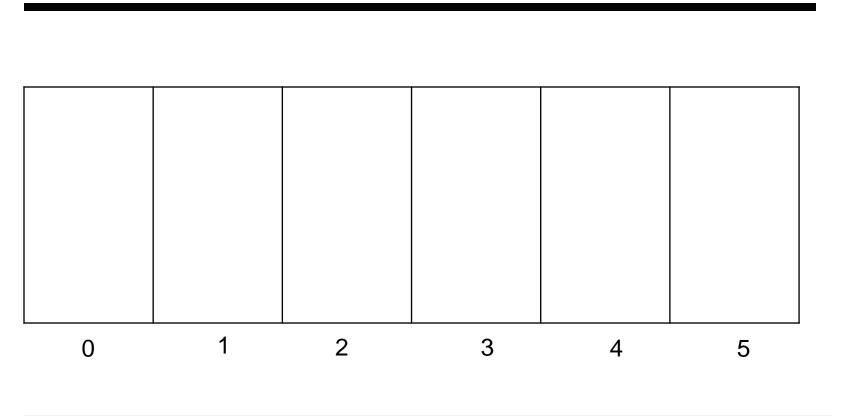


4





Suppose that we have a queue that can hold 6 elements



We need to keep track of a few things:

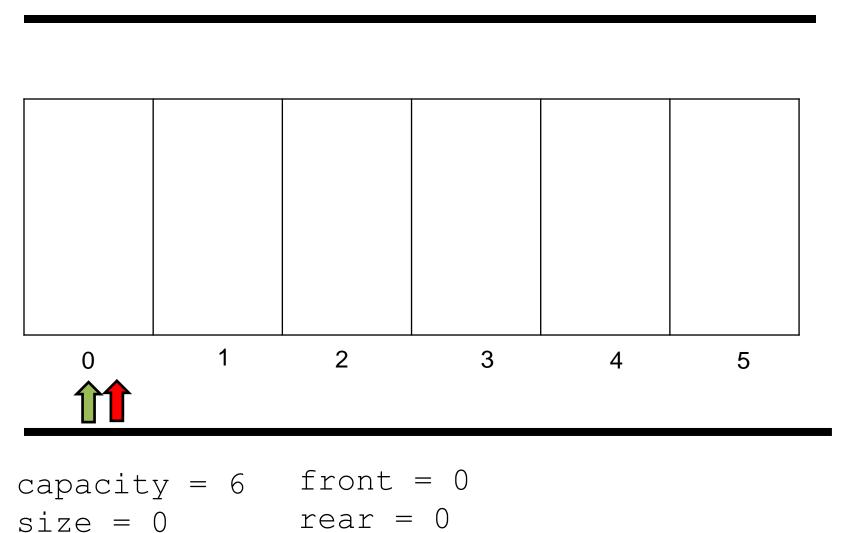
1. The index of the **front** of the queue

2. The index of the rear of the queue

- 3. The size of the queue
- 4. The capacity of the queue



Suppose that we have a queue that can hold 6 elements



We need to keep track of a few things:

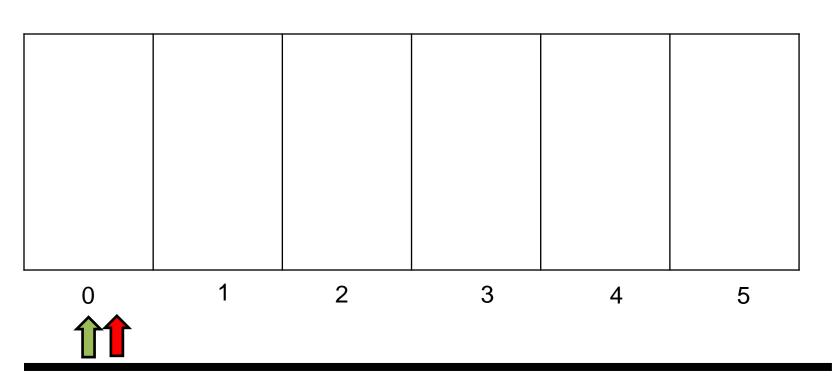
1. The index of the **front** of the queue

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- 3. The size of the queue
- 4. The capacity of the queue



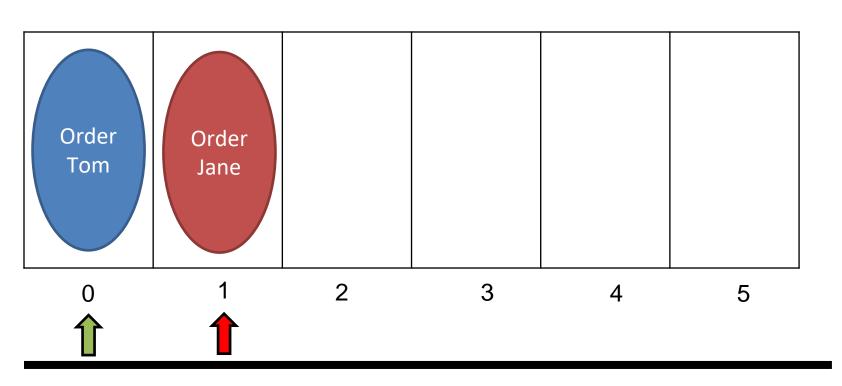
Suppose that we have a queue that can hold 6 elements



capacity = 6 front = 0 size = 0 rear = 0



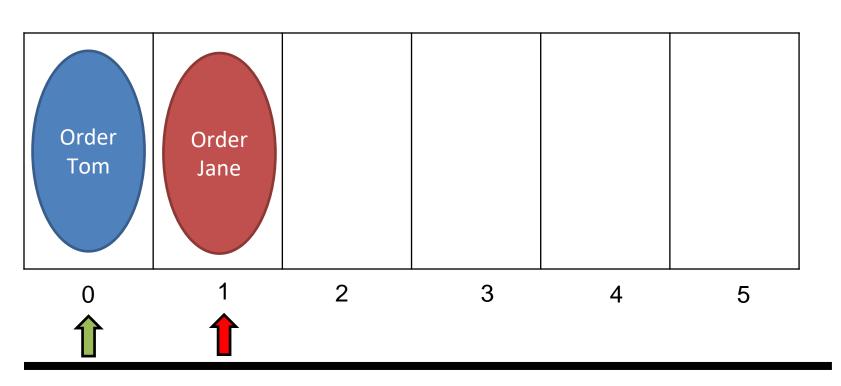
Suppose that we have a queue that can hold 6 elements



capacity = 6 front = 0size = 2 rear = 1



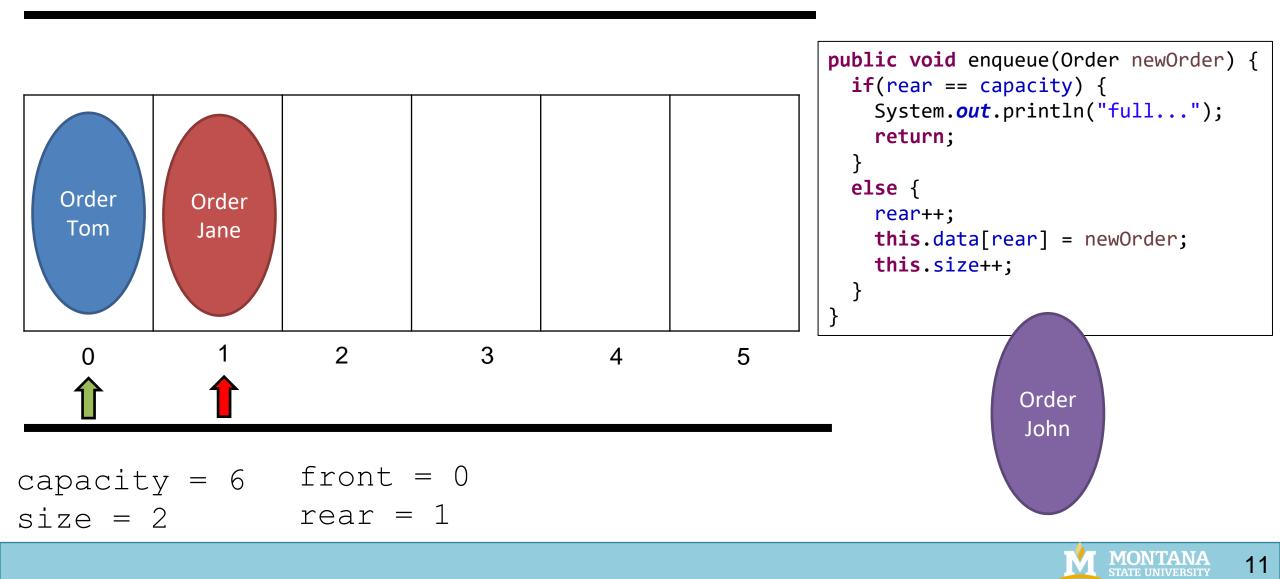
Suppose that we have a queue that can hold 6 elements



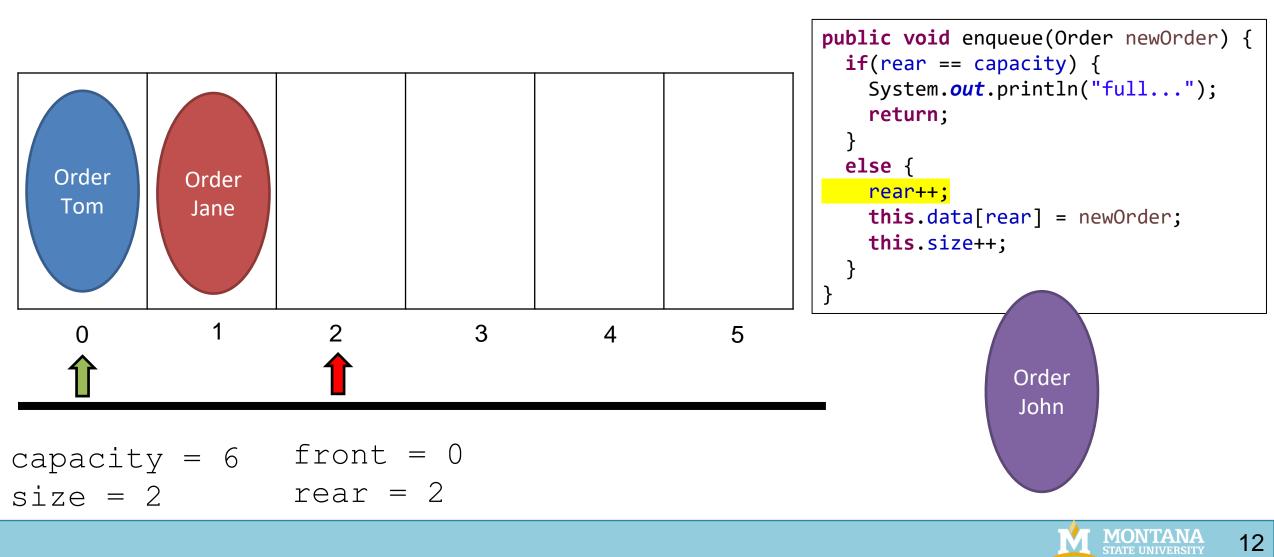
Enqueue?

capacity = 6 front = 0 size = 2 rear = 1



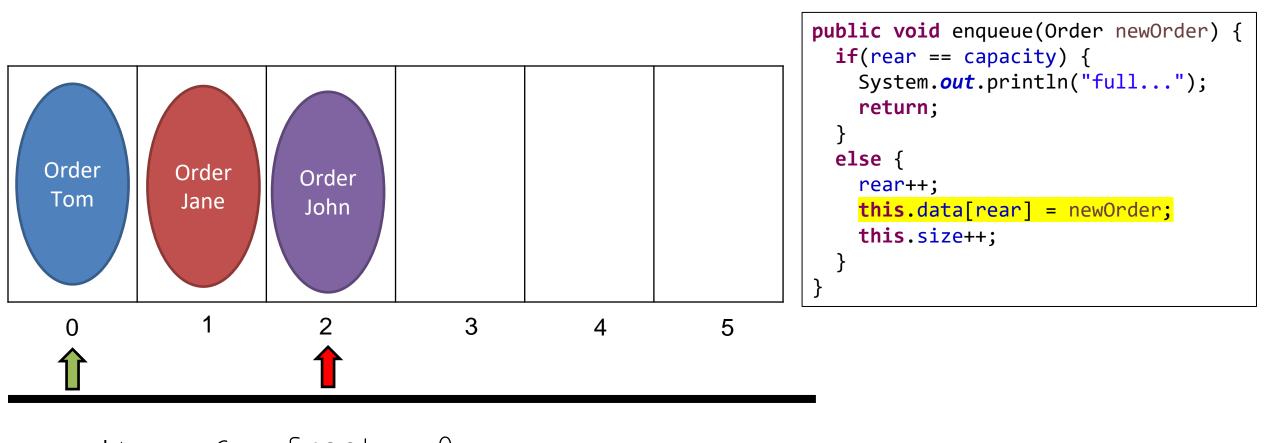


Suppose that we have a queue that can hold 6 elements



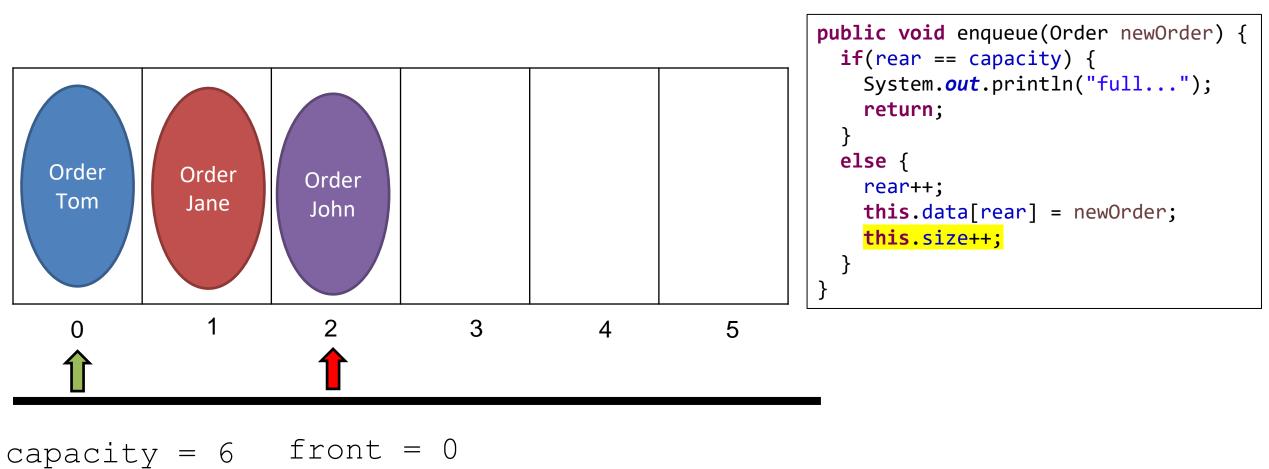
12

Suppose that we have a queue that can hold 6 elements



capacity = 6 front = 0 size = 2 rear = 2

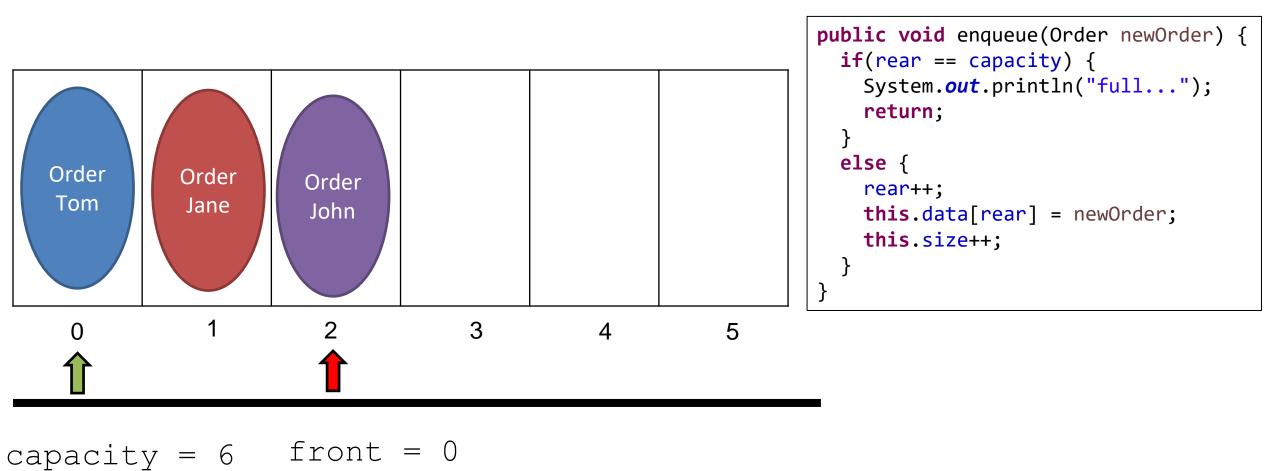
Suppose that we have a queue that can hold 6 elements



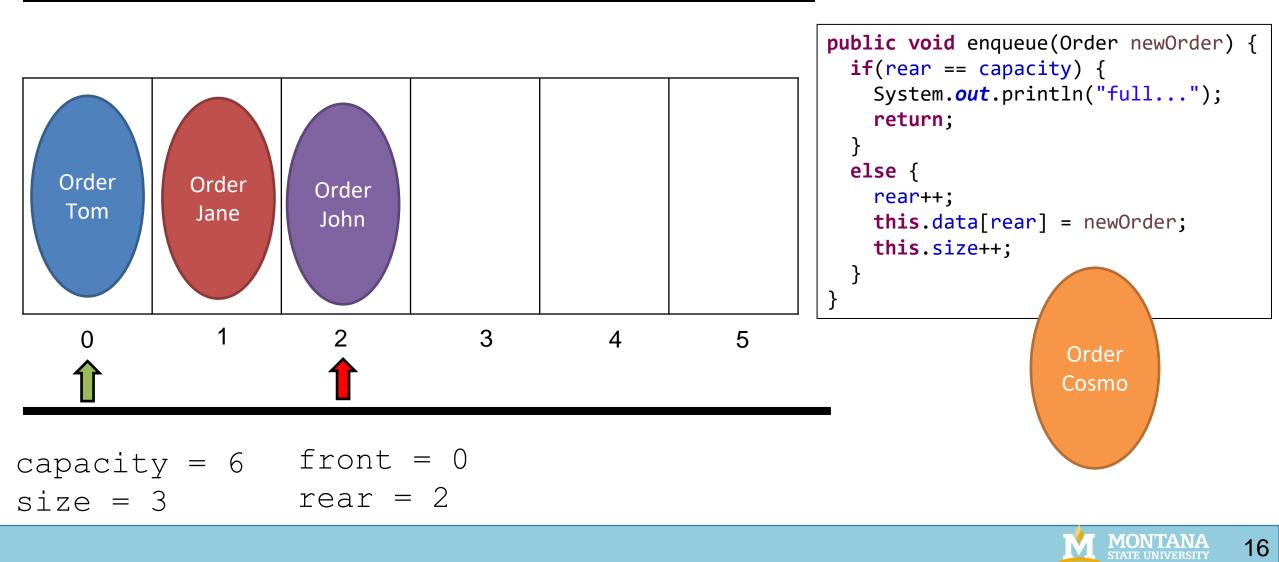
size = 3 rear = 2



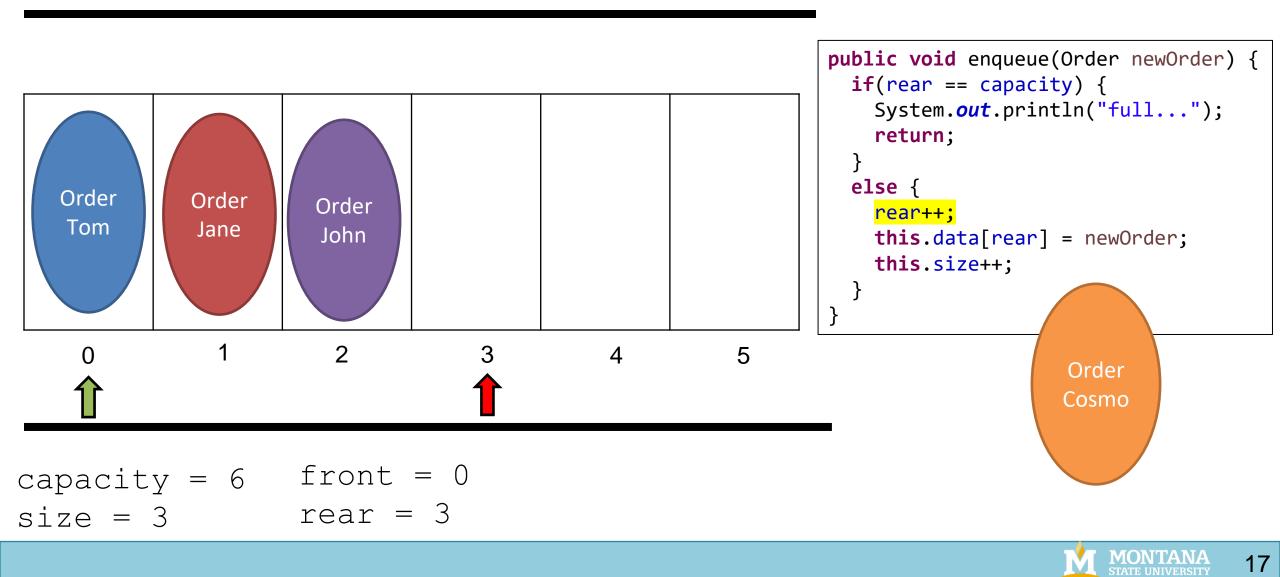
Suppose that we have a queue that can hold 6 elements

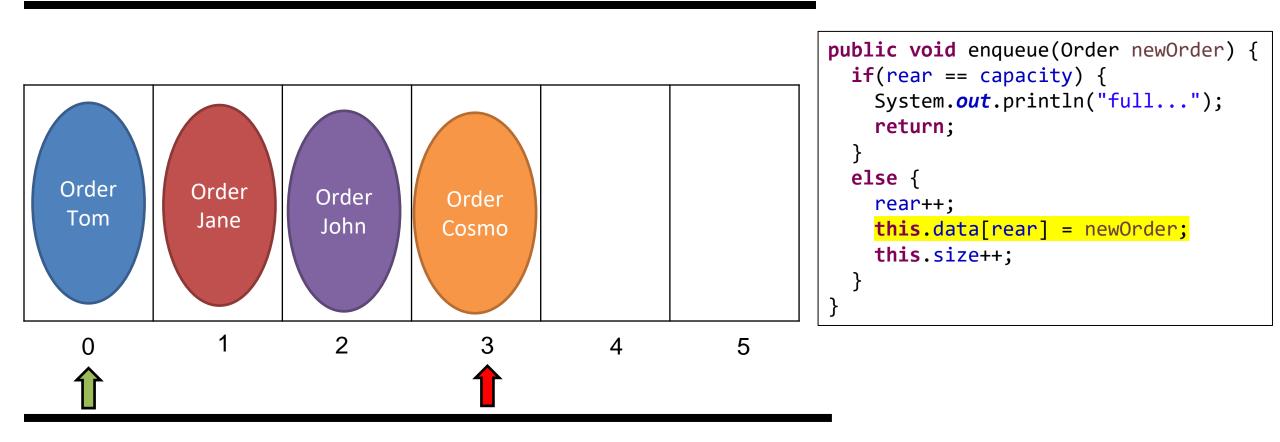


size = 3 rear = 2





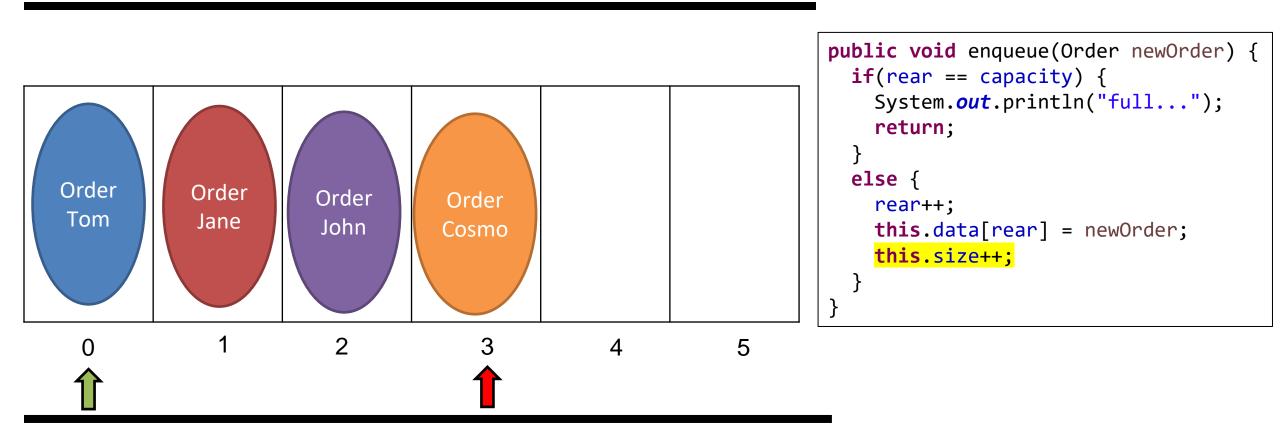




$$capacity = 6 \quad front = 0$$

size = 3
$$rear = 3$$

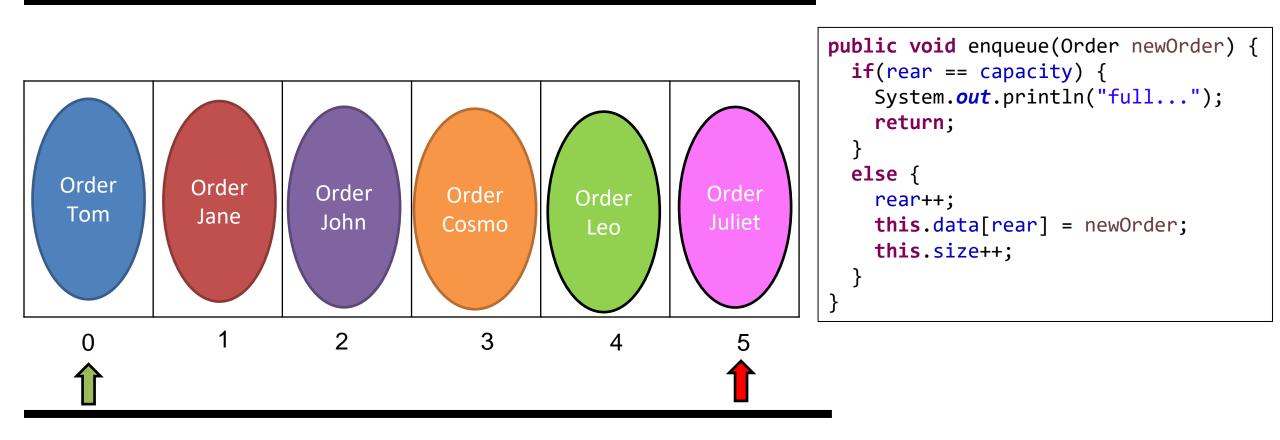
Suppose that we have a queue that can hold 6 elements



 $capacity = 6 \quad front = 0$ size = 4 rear = 3



Suppose that we have a queue that can hold 6 elements



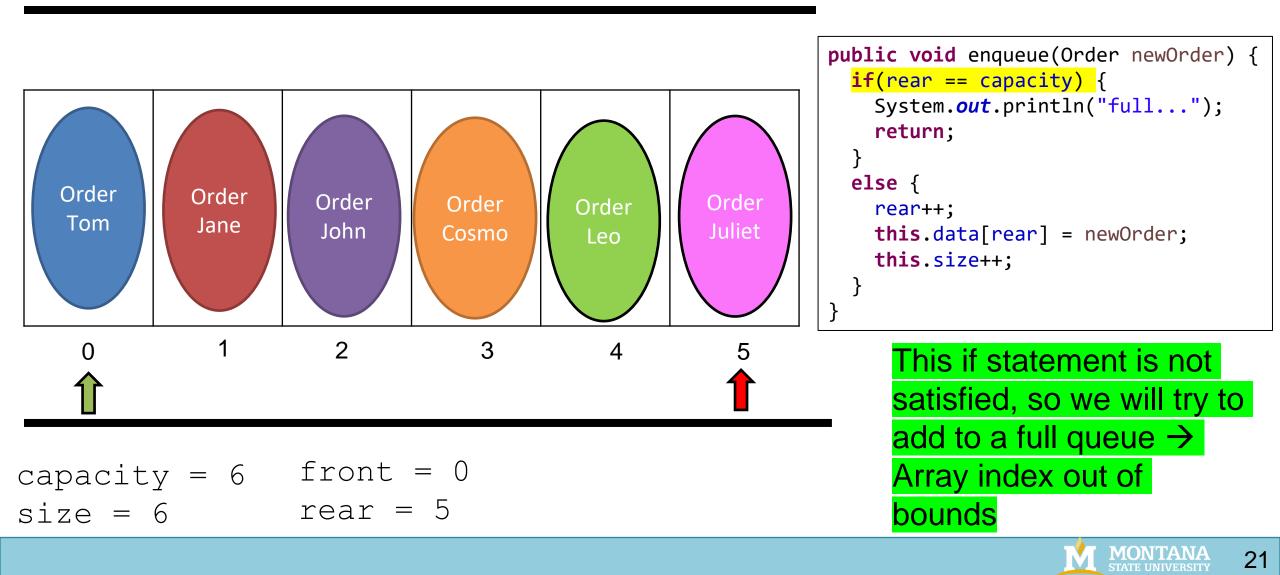
$$capacity = 6 \quad front = 0$$

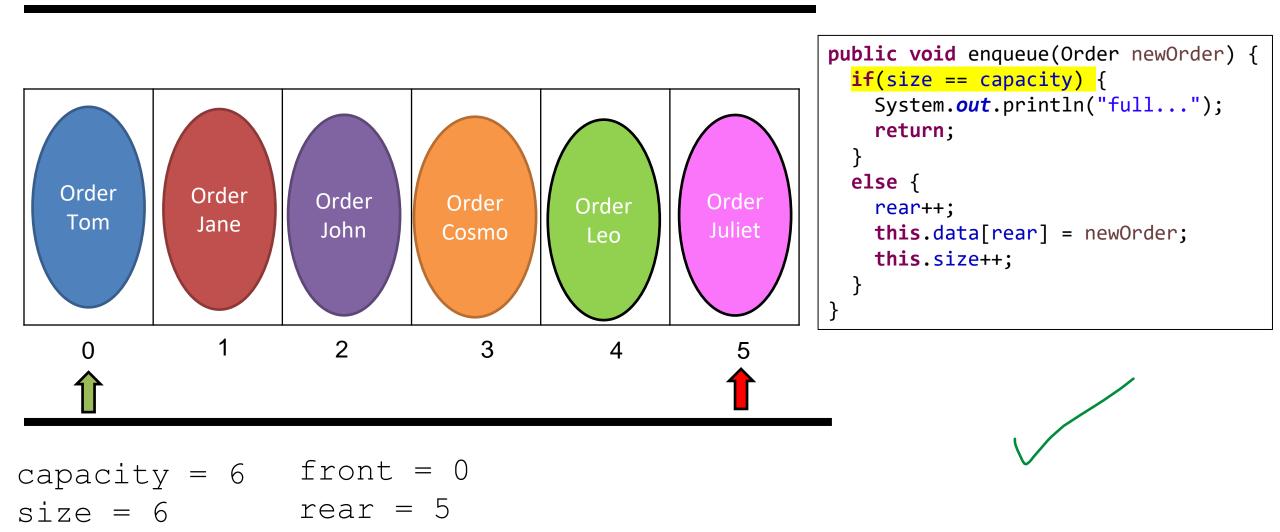
size = 6 \quad rear = 5

Issues with this?

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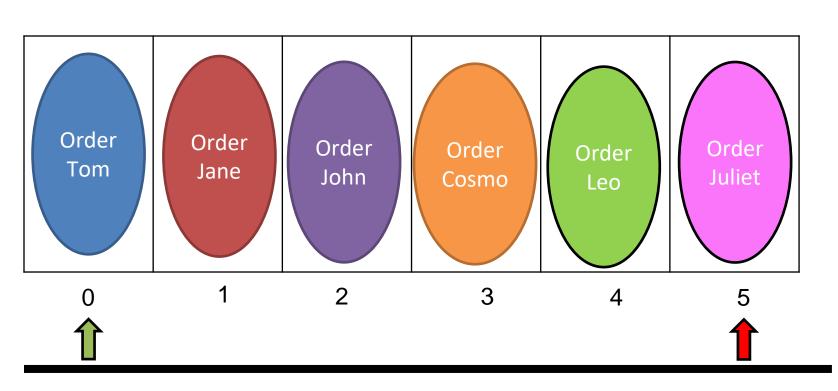
20







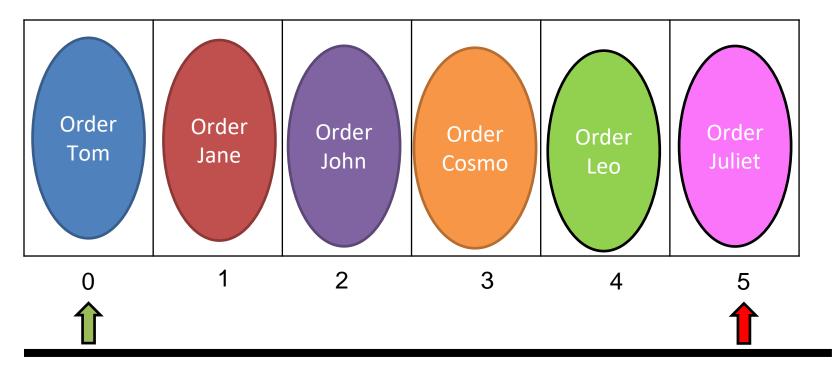
Suppose that we have a queue that can hold 6 elements





capacity = 6 front = 0 size = 6 rear = 5

Suppose that we have a queue that can hold 6 elements

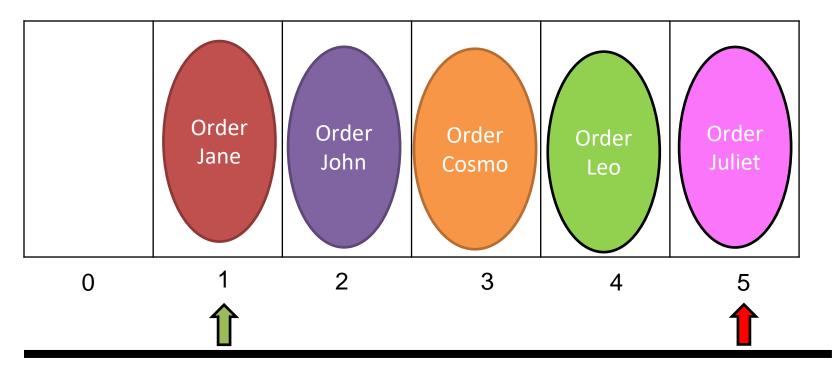


Remove the front element, move front pointer forward one spot





Suppose that we have a queue that can hold 6 elements

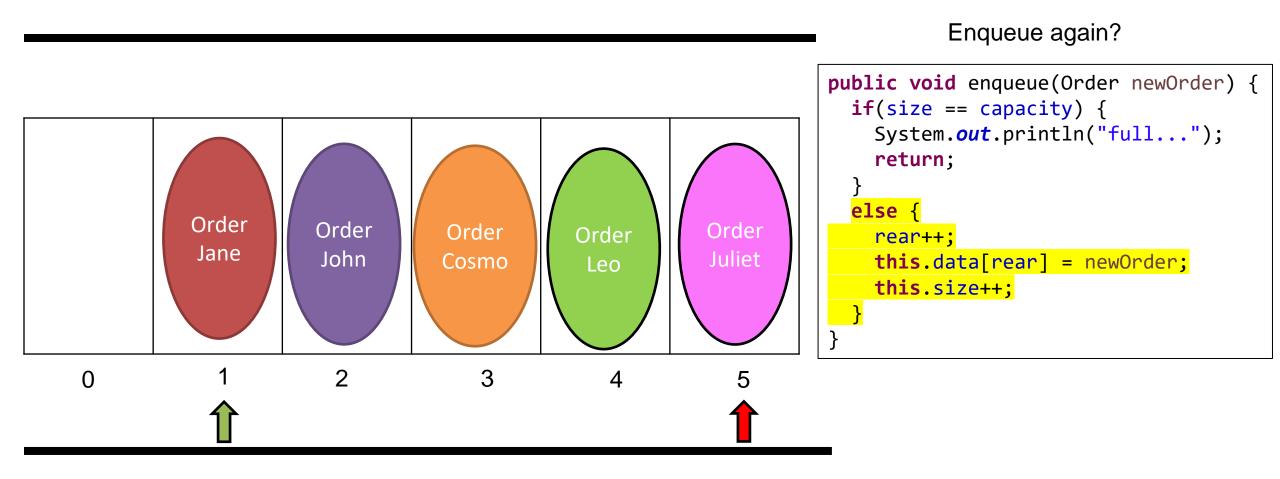


Remove the front element, move front pointer forward one spot

capacity = 6 front = 0 size = 6 rear = 5



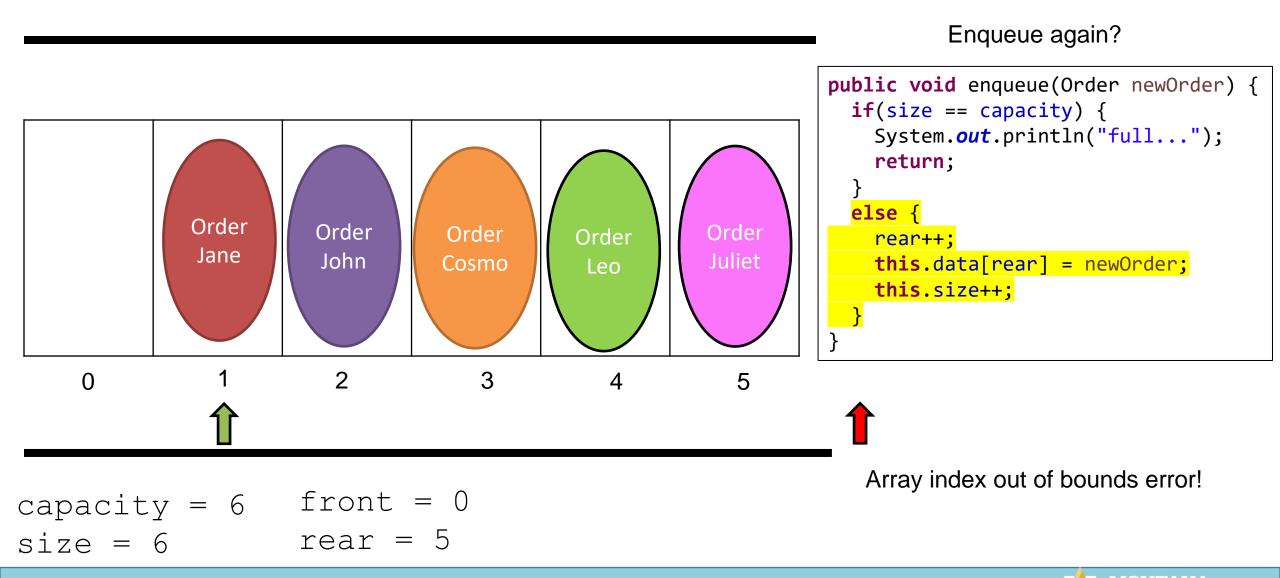
Suppose that we have a queue that can hold 6 elements



 $capacity = 6 \quad front = 0$ size = 6 \quad rear = 5

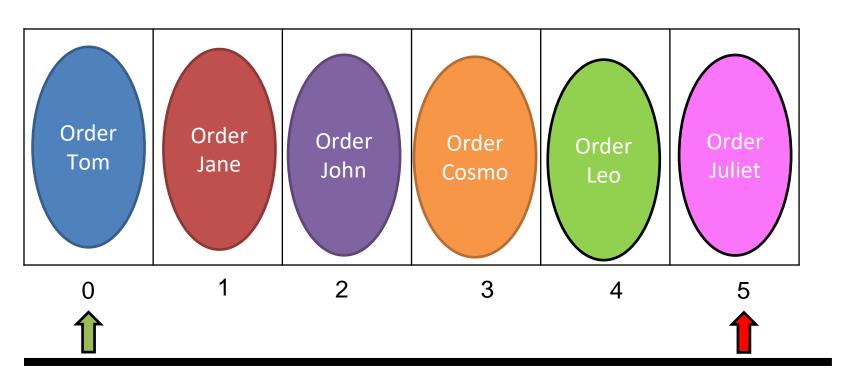


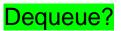
Suppose that we have a queue that can hold 6 elements



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Suppose that we have a queue that can hold 6 elements



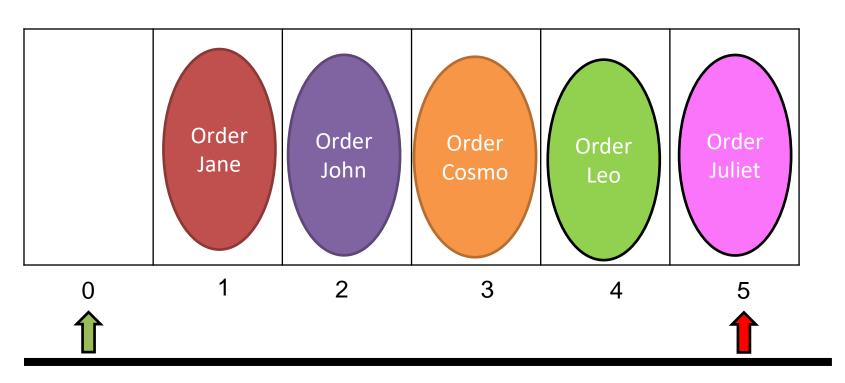


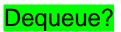
- 1. Remove the front element
- 2. Make some room in the back

capacity = 6 front = 0size = 6 rear = 5



Suppose that we have a queue that can hold 6 elements



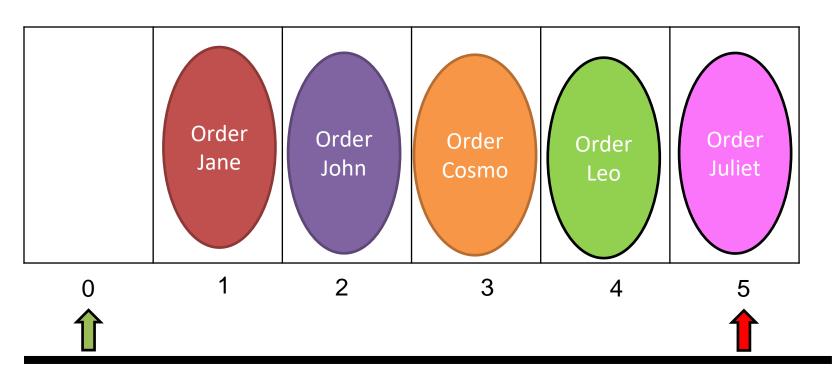


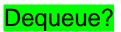
- 1. Remove the front element
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capacity = 6 front = 0size = 6 rear = 5



Suppose that we have a queue that can hold 6 elements





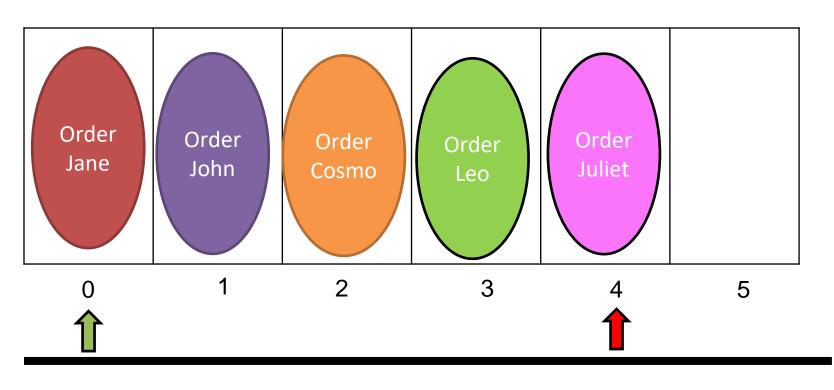
- 1. Remove the front element
- 2. Make some room in the back

Shift all of our data over one spot

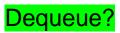
capacity = 6 front = 0size = 6 rear = 5



Suppose that we have a queue that can hold 6 elements



$$capacity = 6$$
 front = 0
size = 5 rear = 4



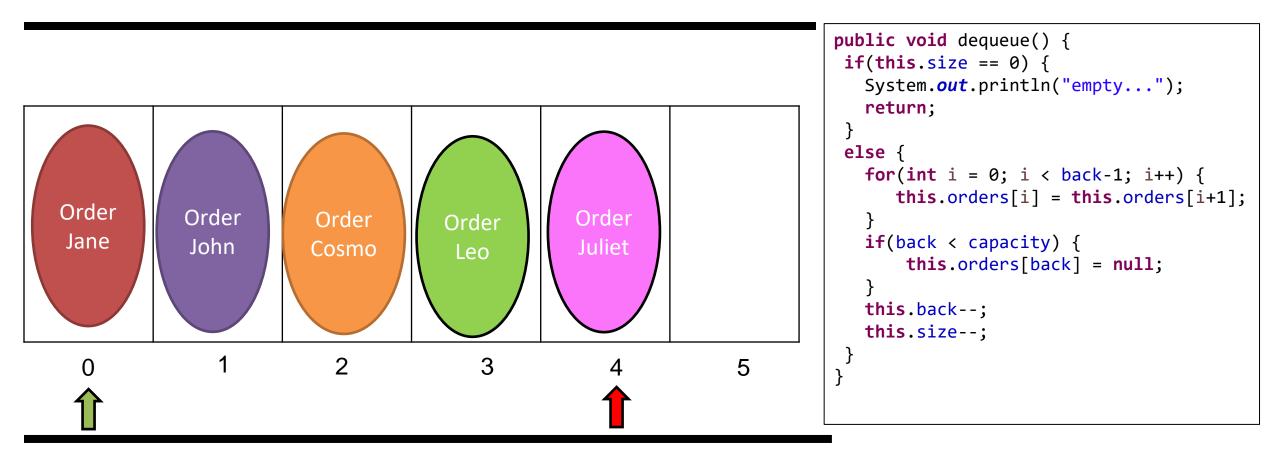
- 1. Remove the front element
- 2. Make some room in the back

Shift all of our data over one spot

The front of our queue will always stay at zero



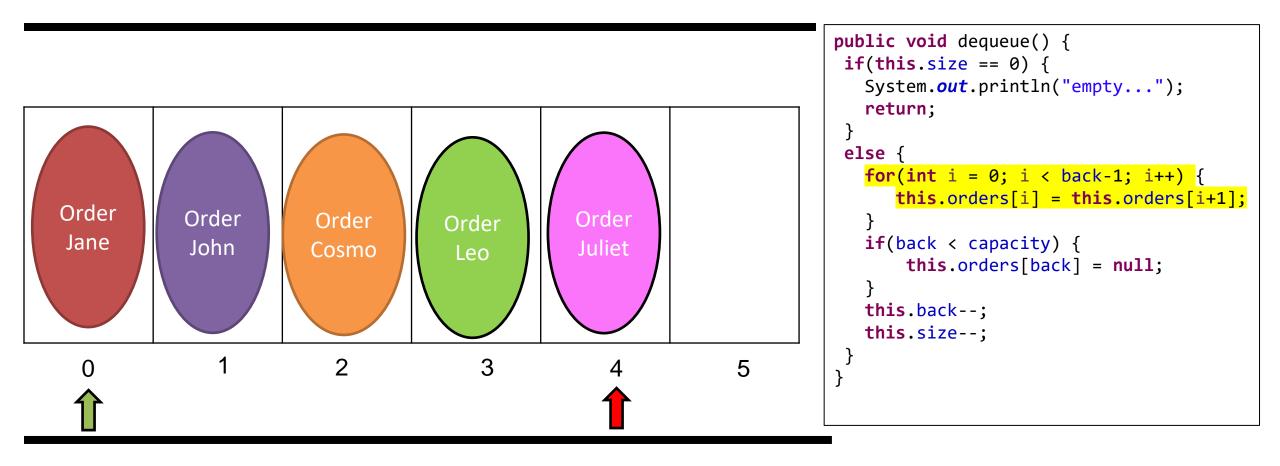
Suppose that we have a queue that can hold 6 elements



 $capacity = 6 \quad front = 0$ size = 5 rear = 4



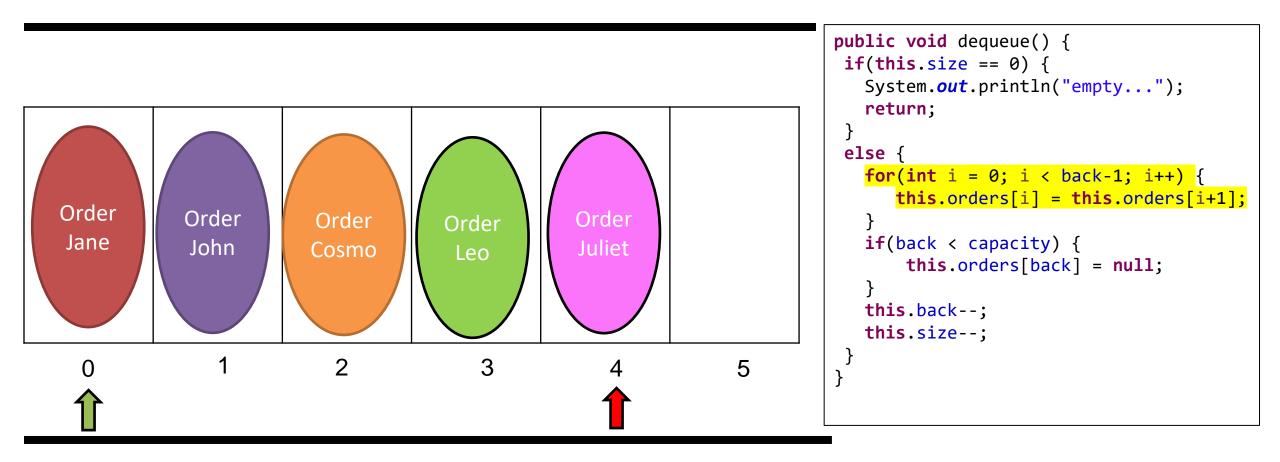
Suppose that we have a queue that can hold 6 elements



 $capacity = 6 \quad front = 0$ size = 5 rear = 4



Suppose that we have a queue that can hold 6 elements

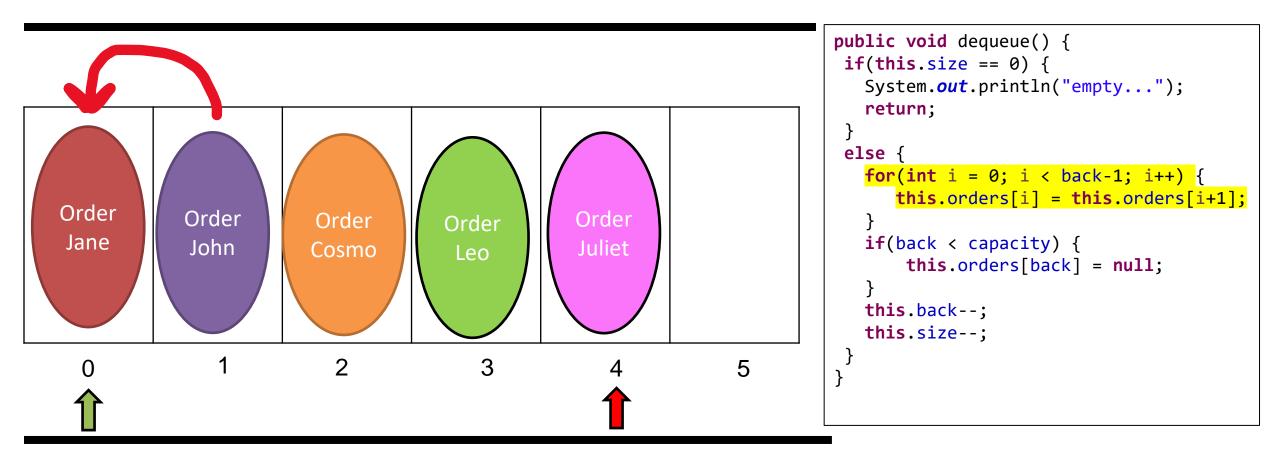


 $capacity = 6 \quad front = 0$ size = 5 rear = 4

Shift everything over one spot



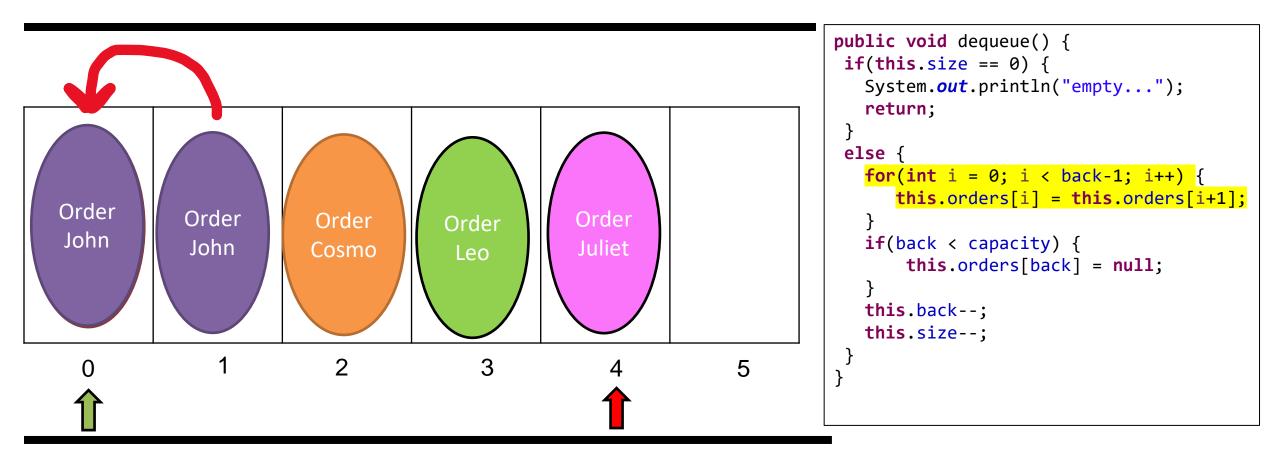
Suppose that we have a queue that can hold 6 elements



 $capacity = 6 \quad front = 0$ size = 5 rear = 4



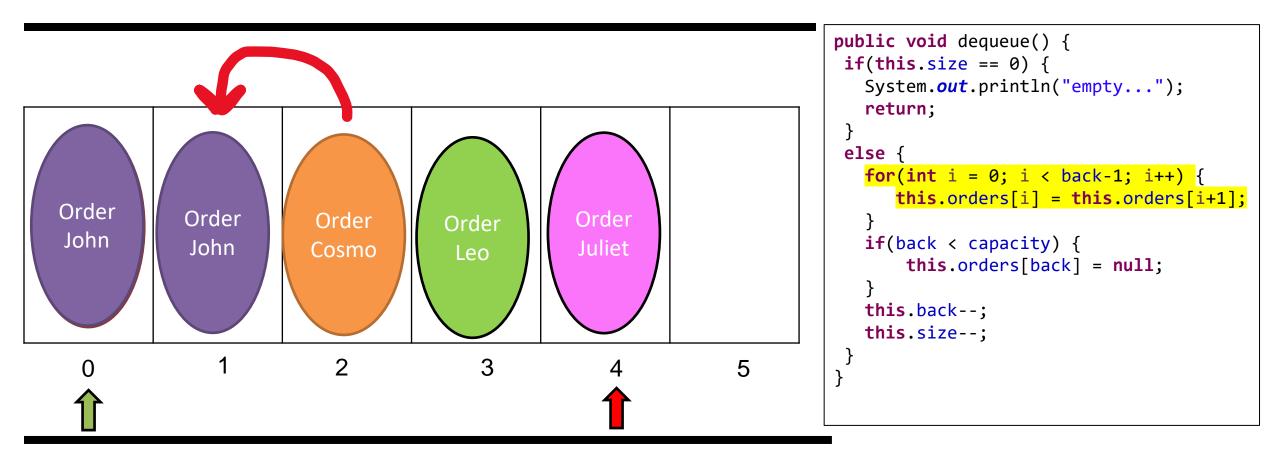
Suppose that we have a queue that can hold 6 elements



capacity = 6 front = 0 size = 5 rear = 4

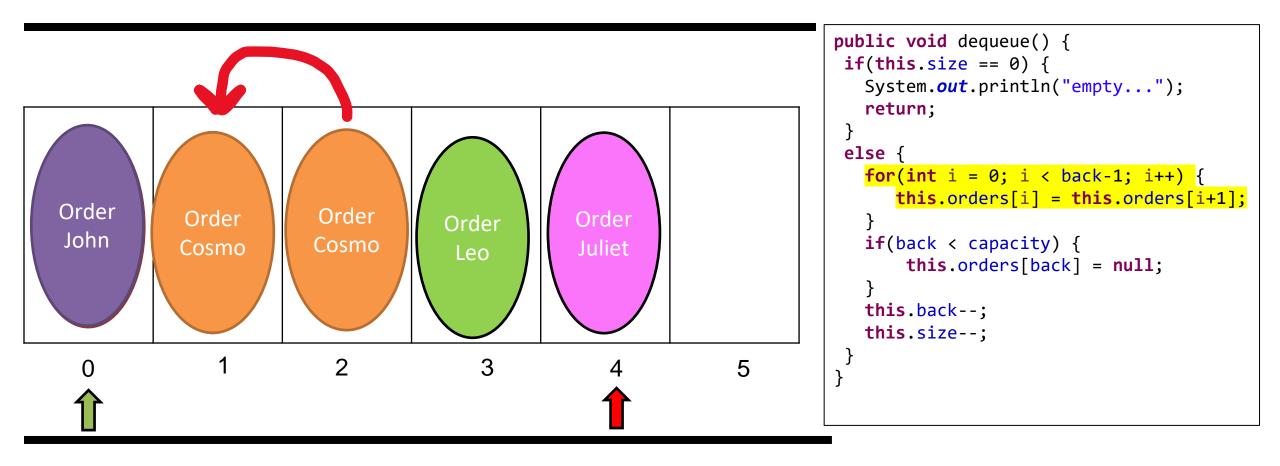


Suppose that we have a queue that can hold 6 elements



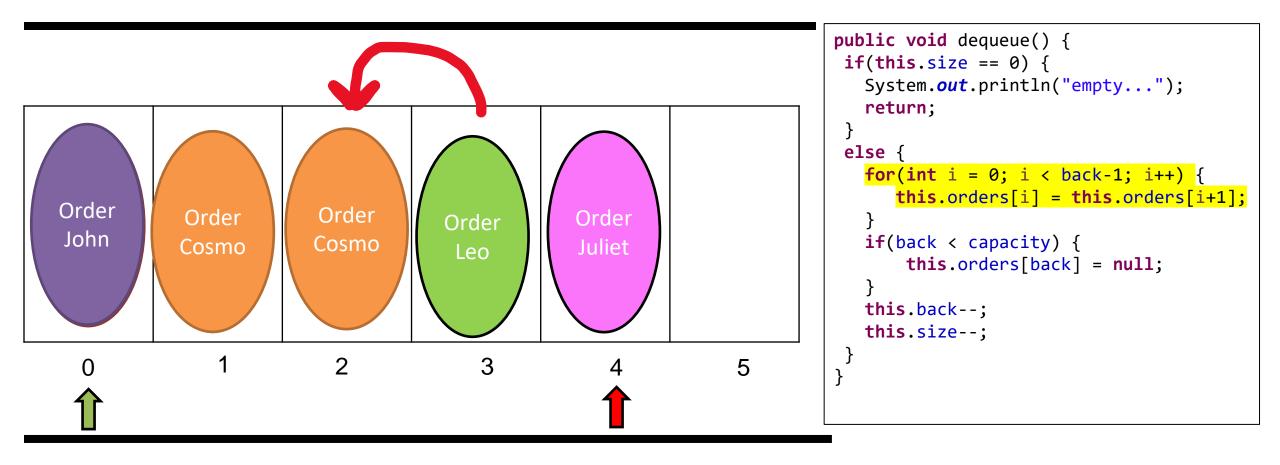


Suppose that we have a queue that can hold 6 elements



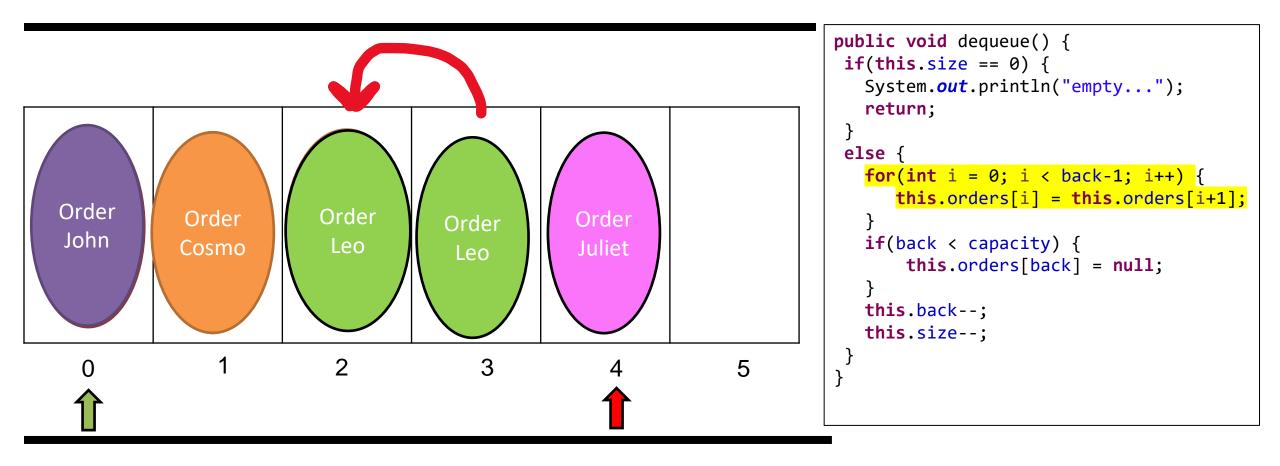


Suppose that we have a queue that can hold 6 elements



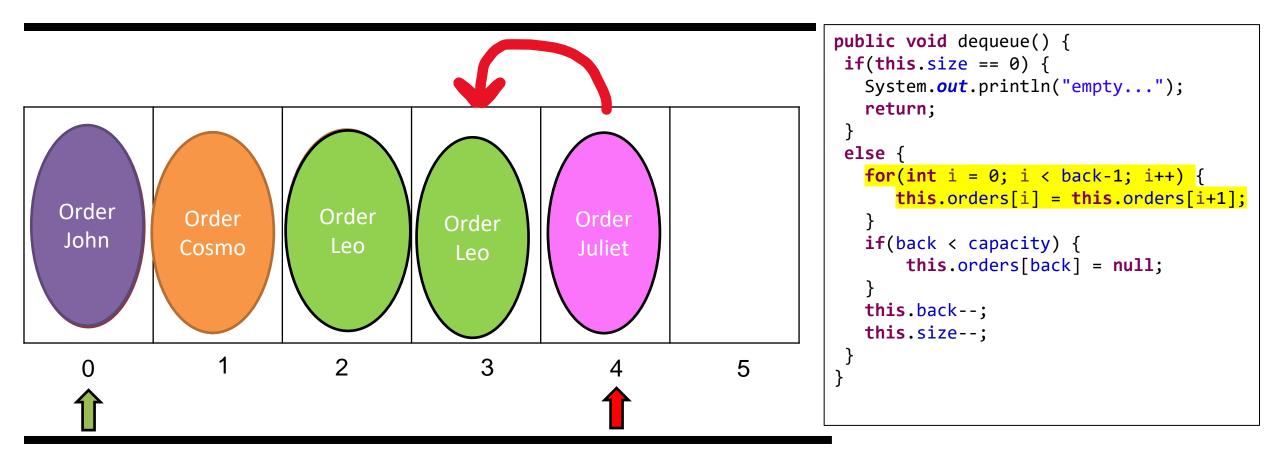


Suppose that we have a queue that can hold 6 elements



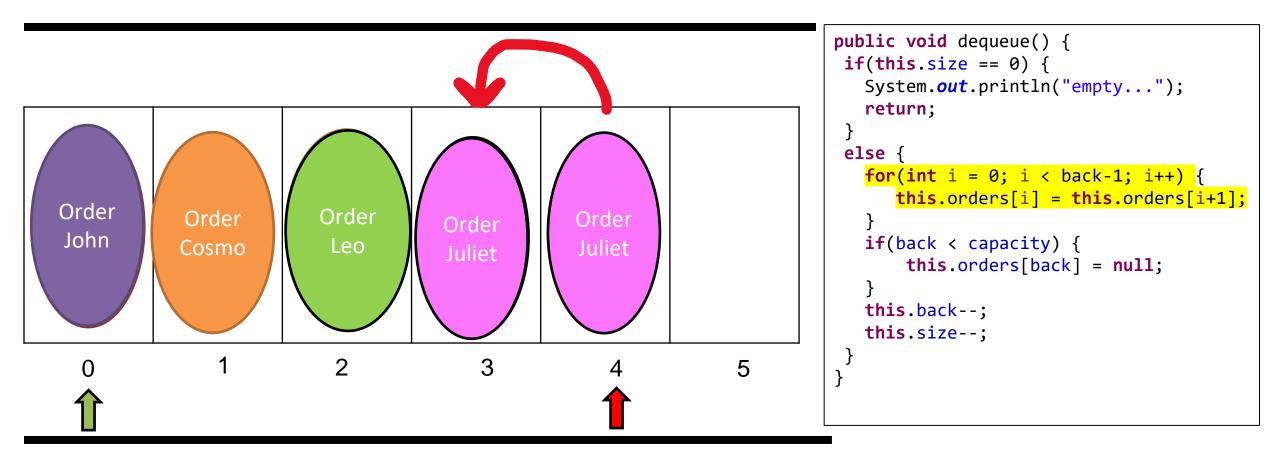


Suppose that we have a queue that can hold 6 elements



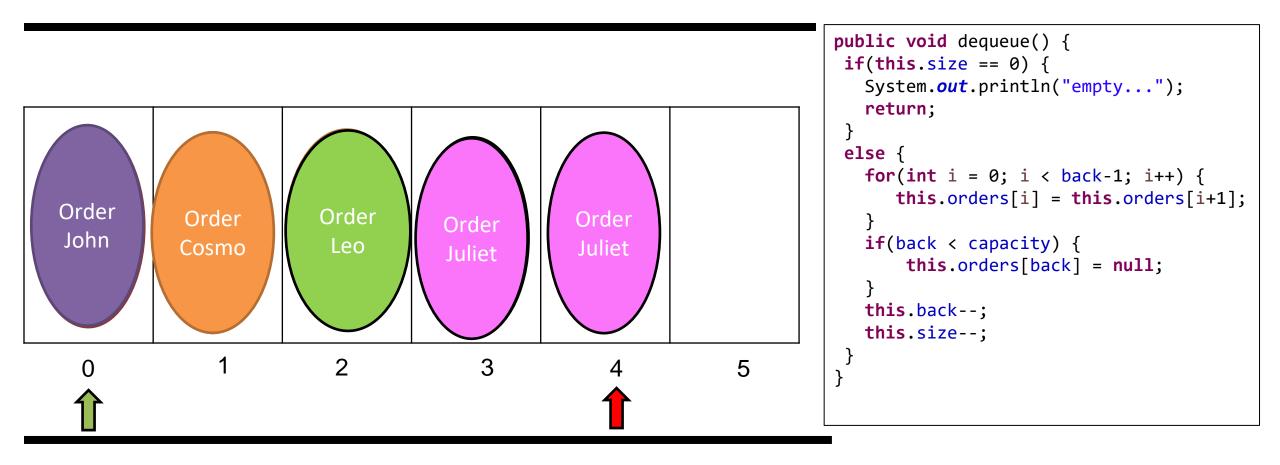


Suppose that we have a queue that can hold 6 elements



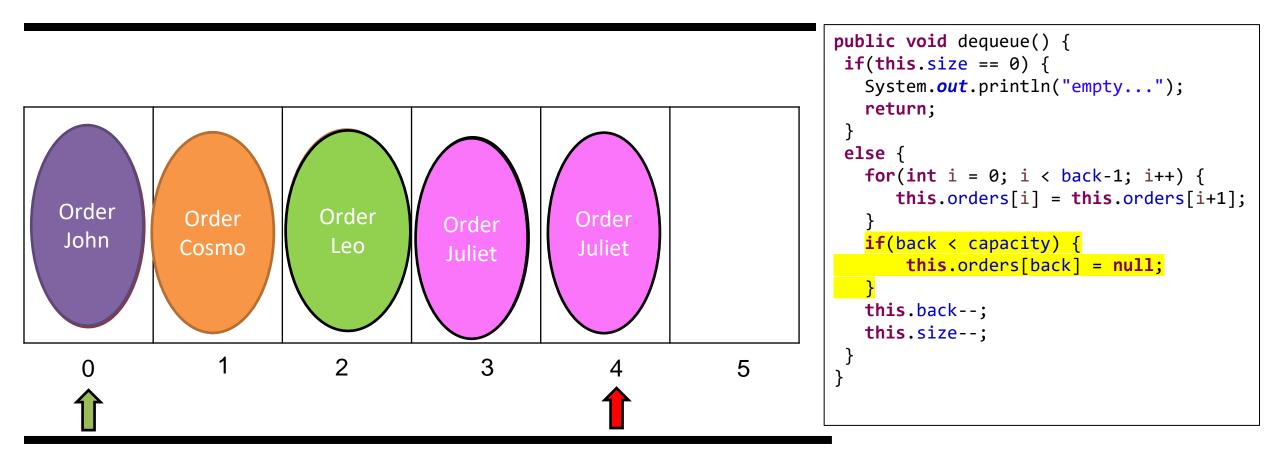


Suppose that we have a queue that can hold 6 elements



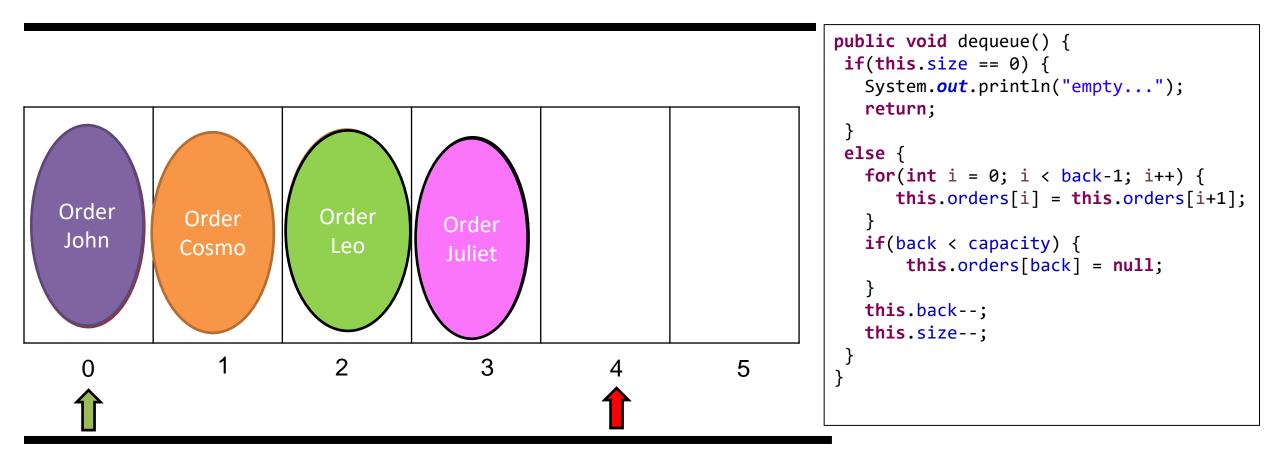


Suppose that we have a queue that can hold 6 elements



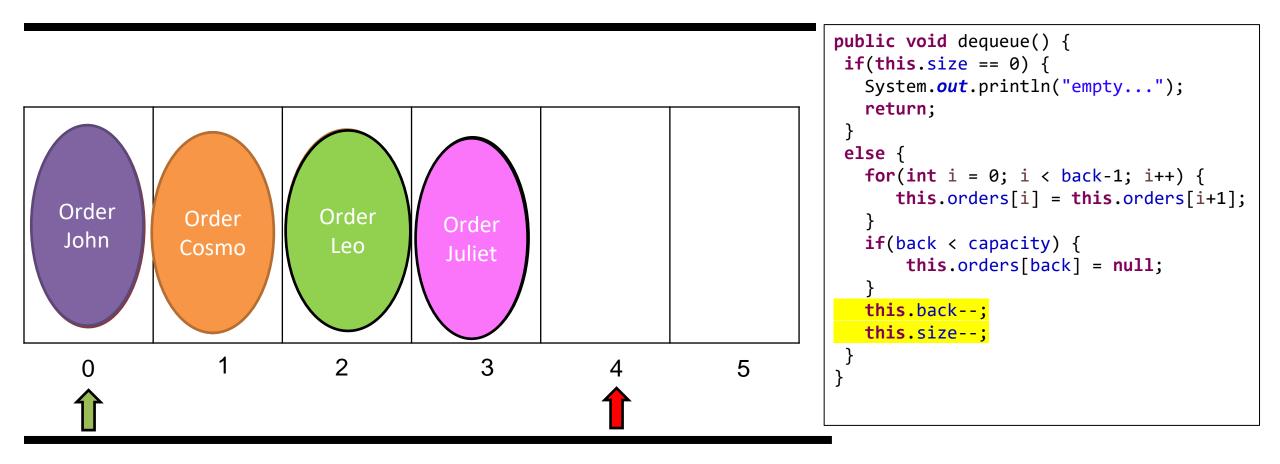


Suppose that we have a queue that can hold 6 elements



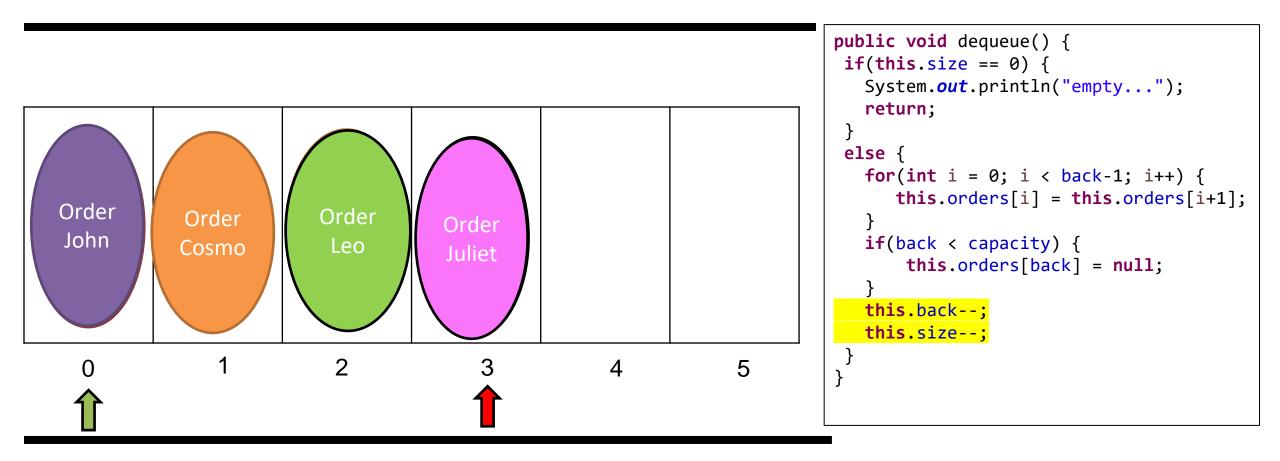


Suppose that we have a queue that can hold 6 elements



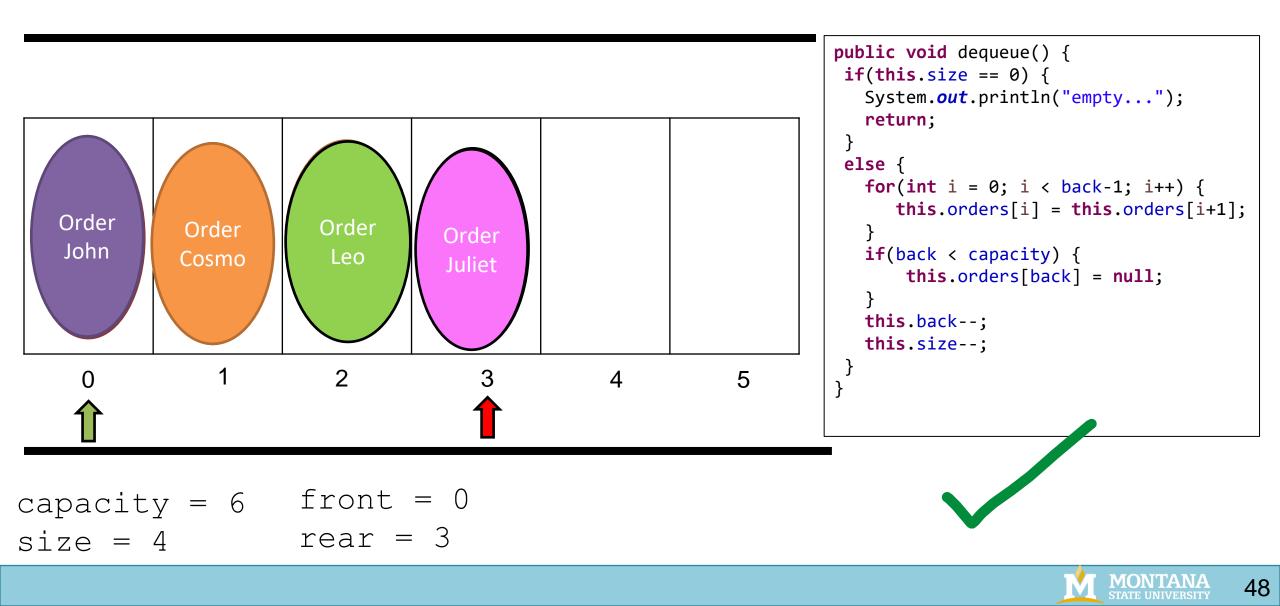


Suppose that we have a queue that can hold 6 elements

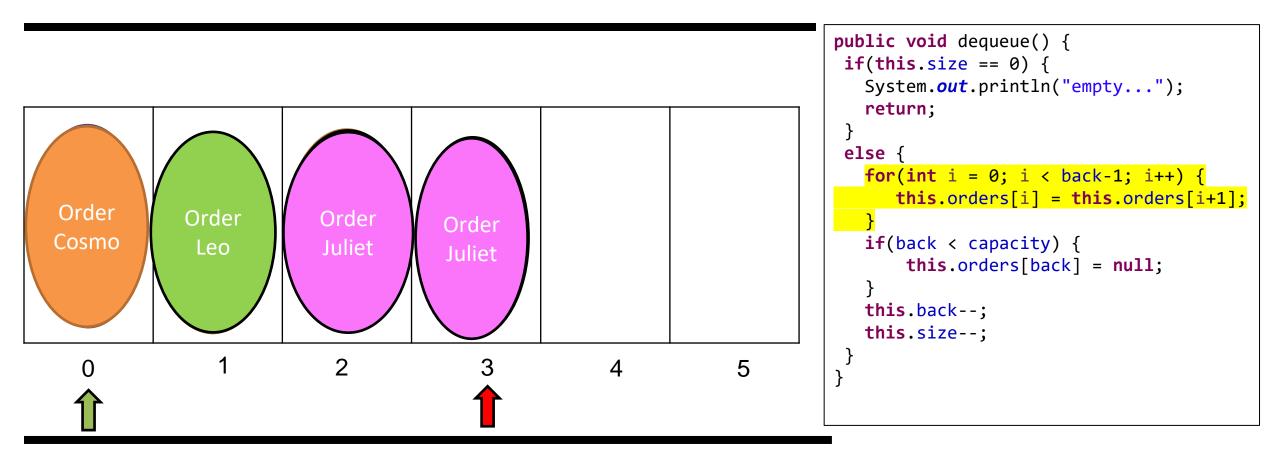




Suppose that we have a queue that can hold 6 elements



Suppose that we have a queue that can hold 6 elements

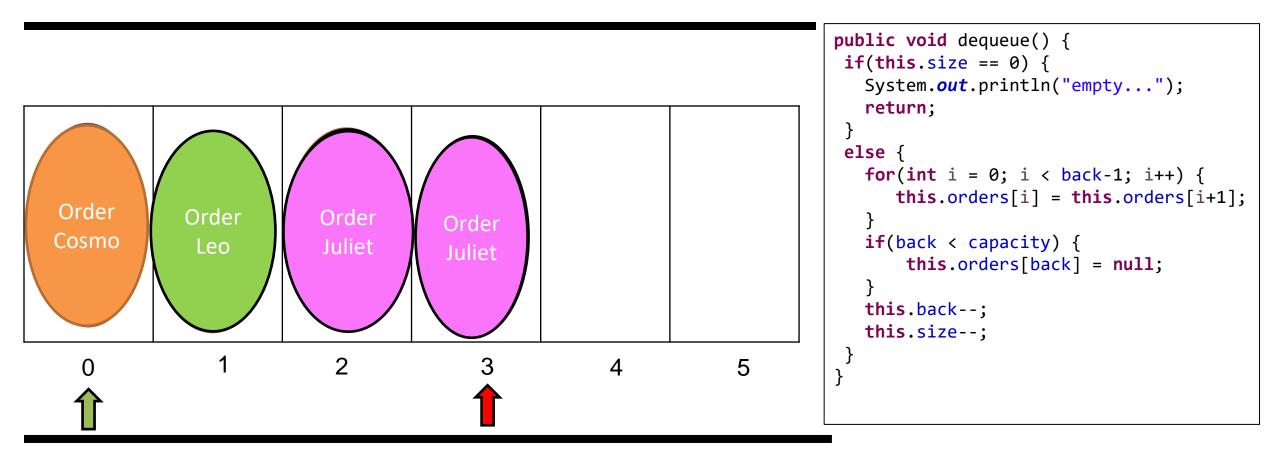


$$capacity = 6 \quad front = 0$$

size = 4
$$rear = 3$$



Suppose that we have a queue that can hold 6 elements

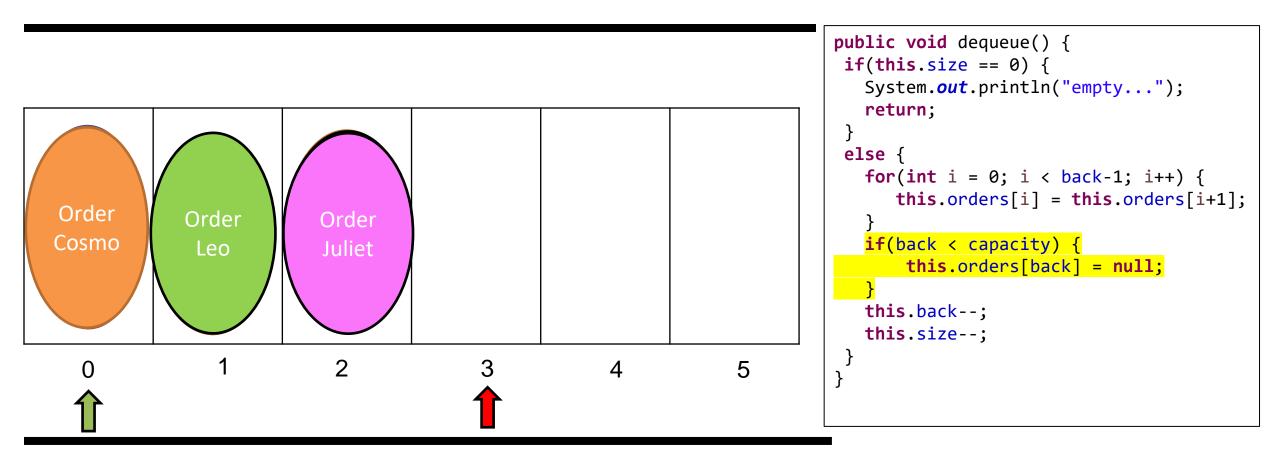


$$capacity = 6 \quad front = 0$$

size = 4
$$rear = 3$$

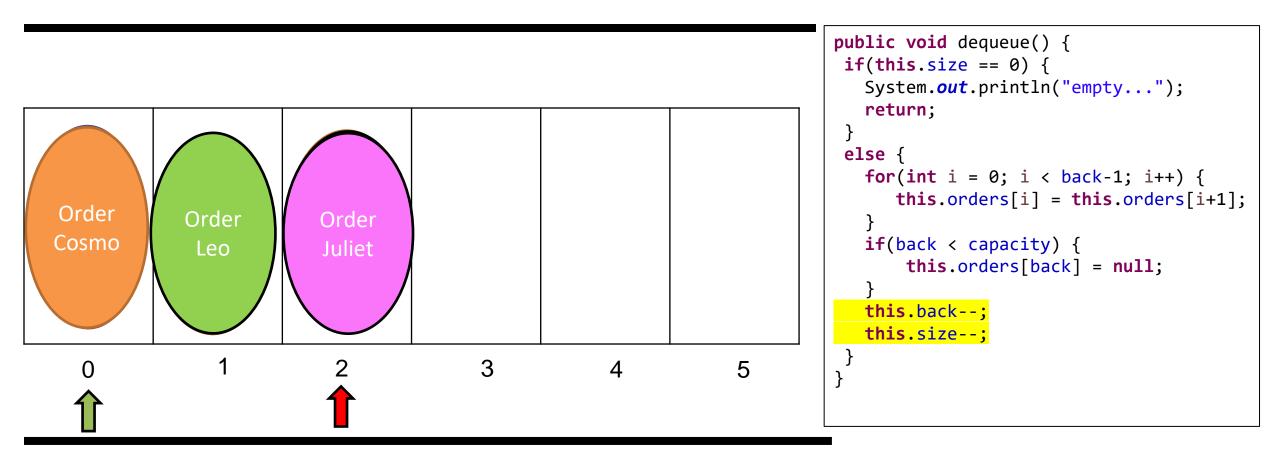


Suppose that we have a queue that can hold 6 elements

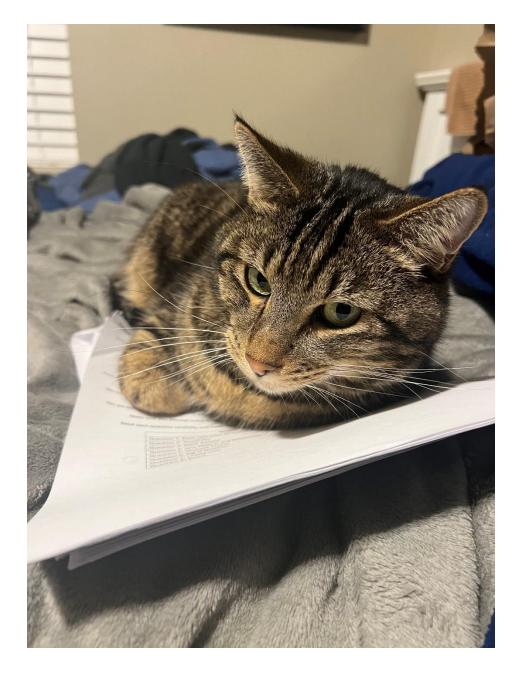




Suppose that we have a queue that can hold 6 elements



capacity = 6 front = 0 size = 3 rear = 2





```
public void enqueue(Order newOrder) {
  if(rear == capacity) {
    System.out.println("full...");
    return;
 else {
    rear++;
    this.data[rear] = newOrder;
    this.size++;
```



```
public void enqueue(Order newOrder) {
  if(rear == capacity) { 0(1)
    System.out.println("full..."); o(1)
    return; O(1)
  }
  else {
    rear++; O(1)
    this.data[rear] = newOrder; o(1)
    this.size++; o(1)
```



```
public void enqueue(Order newOrder) {
  if(rear == capacity) { 0(1)
    System.out.println("full..."); o(1)
    return; O(1)
  }
  else {
                                            Total running time:
    rear++; 0(1)
    this.data[rear] = newOrder; o(1)
                                             O(1)
    this.size++; O(1)
```



```
public void dequeue() {
 if(this.size == 0) {
   System.out.println("empty...");
   return;
 else {
   for(int i = 0; i < back-1; i++) {</pre>
      this.orders[i] = this.orders[i+1];
   if(back < capacity) {</pre>
       this.orders[back] = null;
   this.back--;
   this.size--;
```



```
public void dequeue() {
 if(this.size == 0) { O(1)
   System.out.println("empty..."); O(1)
   return; O(1)
 }
 else {
   for(int i = 0; i < back-1; i++) { O(N-1)</pre>
      this.orders[i] = this.orders[i+1];0(1)
   }
   if(back < capacity) { O(1)</pre>
       this.orders[back] = null; o(1)
   }
   this.back--; O(1)
   this.size--; O(1)
```

```
N = # elements
in our queue
```

```
public void dequeue() {
 if(this.size == 0) { O(1)
   System.out.println("empty..."); O(1)
   return; O(1)
 }
 else {
   for(int i = 0; i < back-1; i++) { O(N-1)</pre>
                                                           N = # elements
      this.orders[i] = this.orders[i+1];0(1)
                                                           in our queue
   }
   if(back < capacity) { O(1)</pre>
       this.orders[back] = null; o(1)
                                                   Total running time:
   }
   this.back--; O(1)
                                                   O(N)
   this.size--; O(1)
```



```
public void dequeue() {
 if(this.size == 0) { O(1)
   System.out.println("empty..."); O(1)
   return; O(1)
 }
 else {
   for(int i = 0; i < back-1; i++) { O(N-1)</pre>
      this.orders[i] = this.orders[i+1];0(1)
   }
   if(back < capacity) { O(1)</pre>
       this.orders[back] = null; o(1)
   }
   this.back--; O(1)
   this.size--; O(1)
```

Total running time:

O(N)

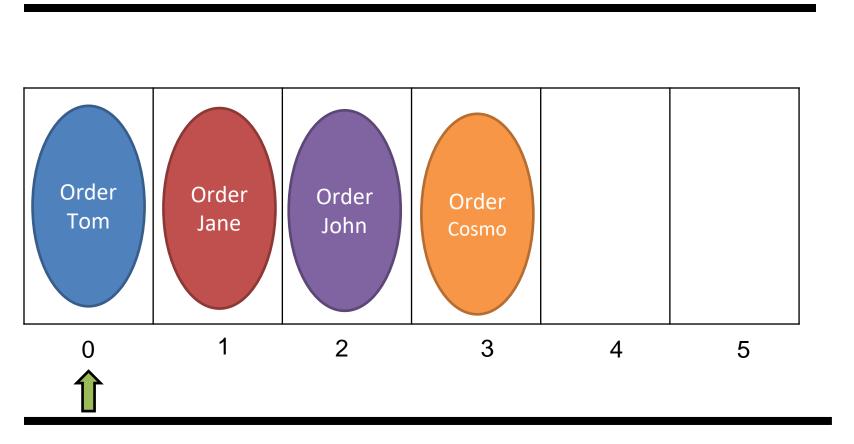
This algorithm works fine, but the issue is that shifting data can be costly

(think about if this queue has 1000000 things in it—) we must shift 999999 elements!)



How to improve our queue?





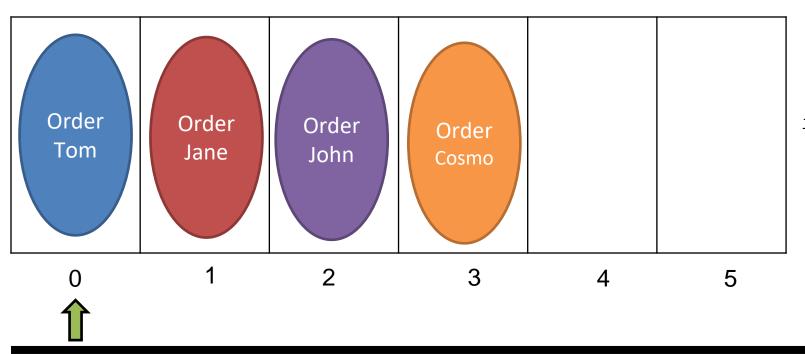
We are going to make use of the **modulus** (%) operator !

10 % 6 = 4

3 % 6 = 3

6 % 6 = 0

capacity = 6 front = 0 size = 4



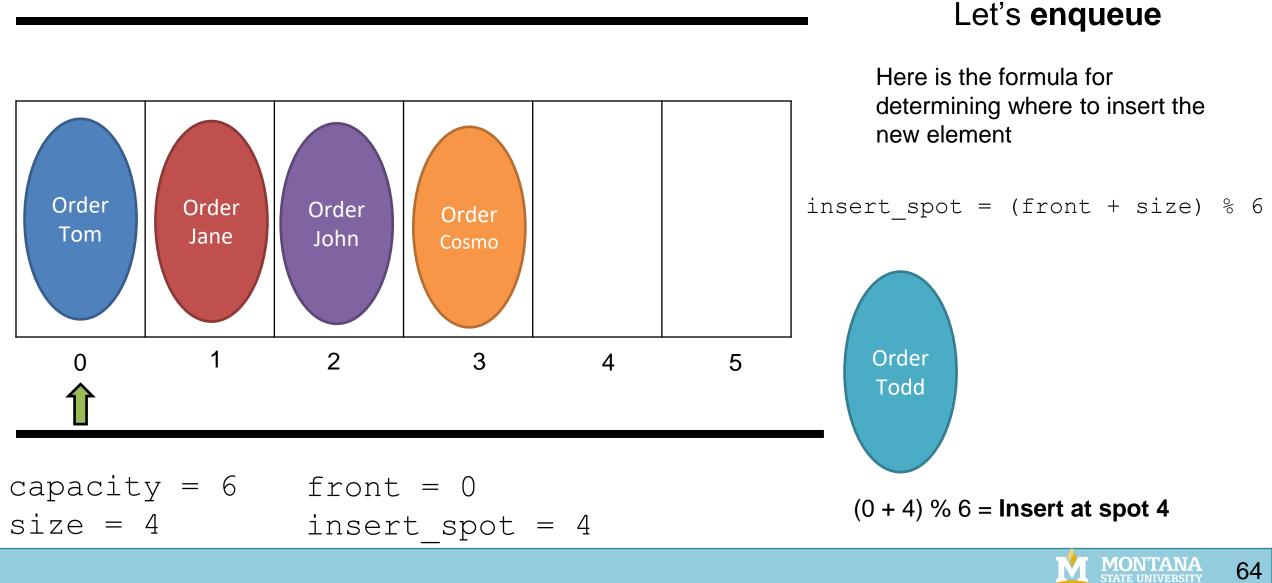
Let's enqueue

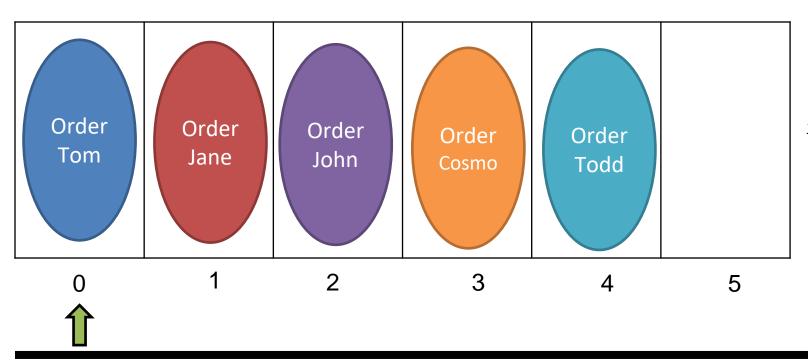
Here is the formula for determining where to insert the new element

```
insert spot = (front + size) % 6
```

capacity = 6 front = 0 size = 4







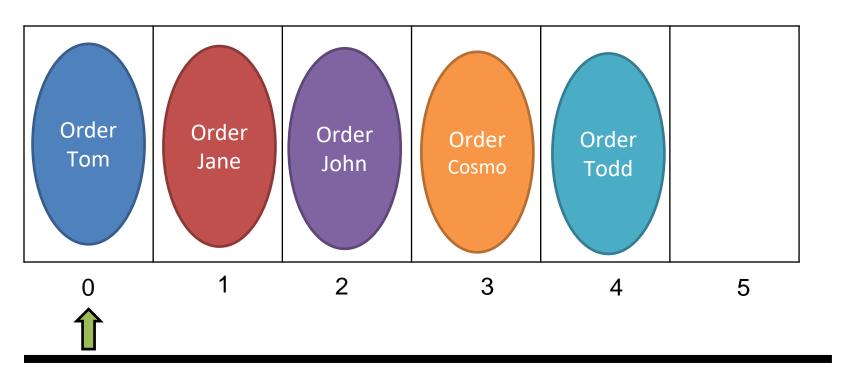
Let's enqueue

Here is the formula for determining where to insert the new element

```
insert_spot = (front + size) % 6
```

capacity = 6 front = 0
size = 4 insert_spot = 4





Let's **dequeue**

data[front] = null

capacity = 6 front = 0
size = 4 insert_spot = 4

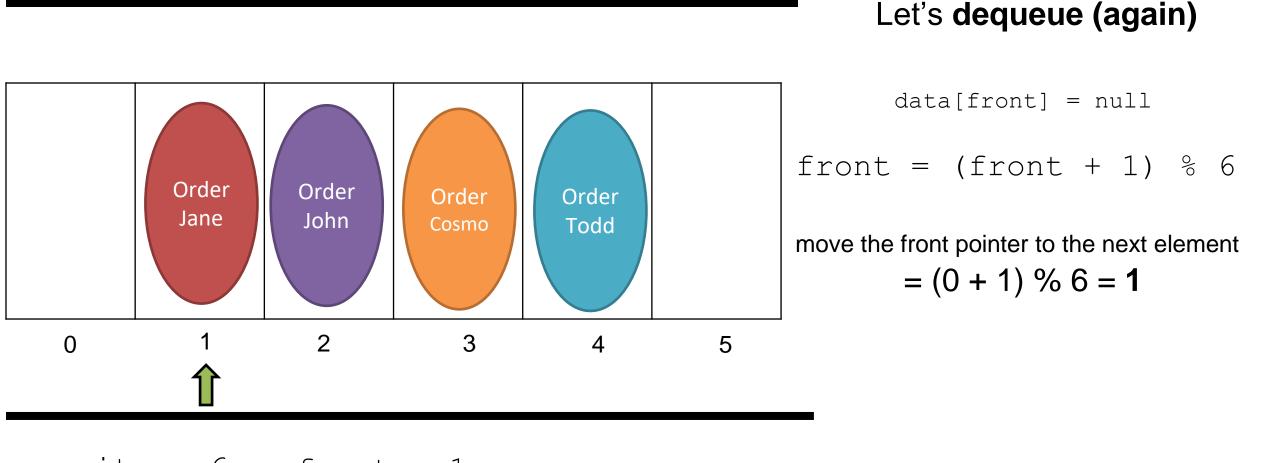


data[front] = null front = (front + 1) % 6 Order Order Order Order Jane John Cosmo Todd move the front pointer to the next element = (0 + 1) % 6 = **1** 2 3 0 4 5

capacity = 6 front = 1 size = 4 insert_spot = 4

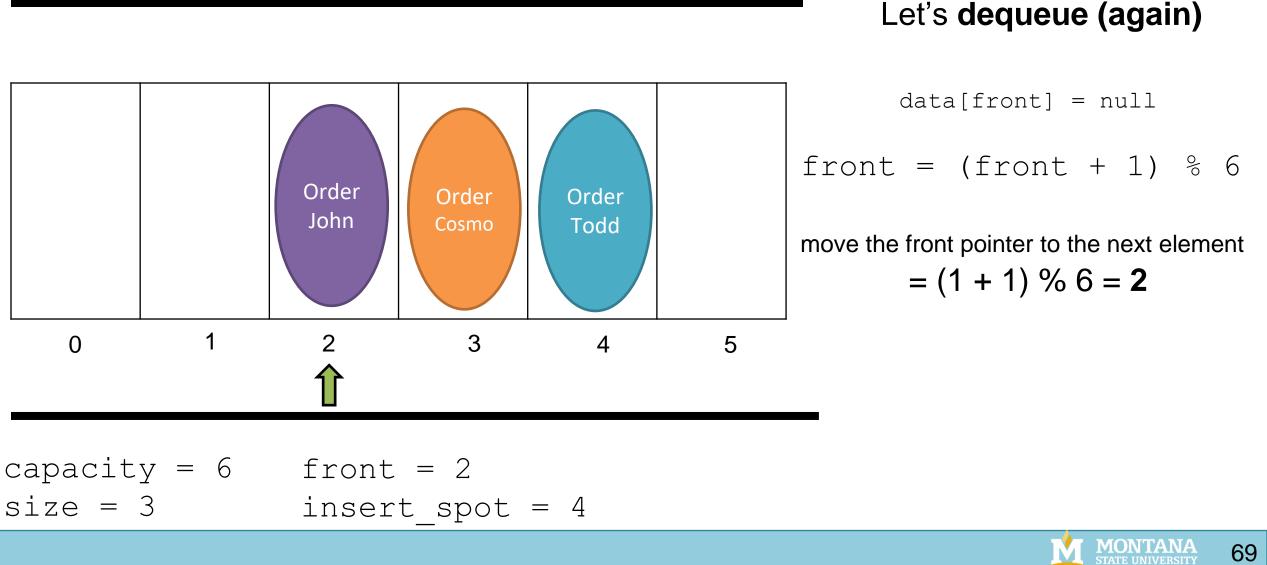


Let's **dequeue**

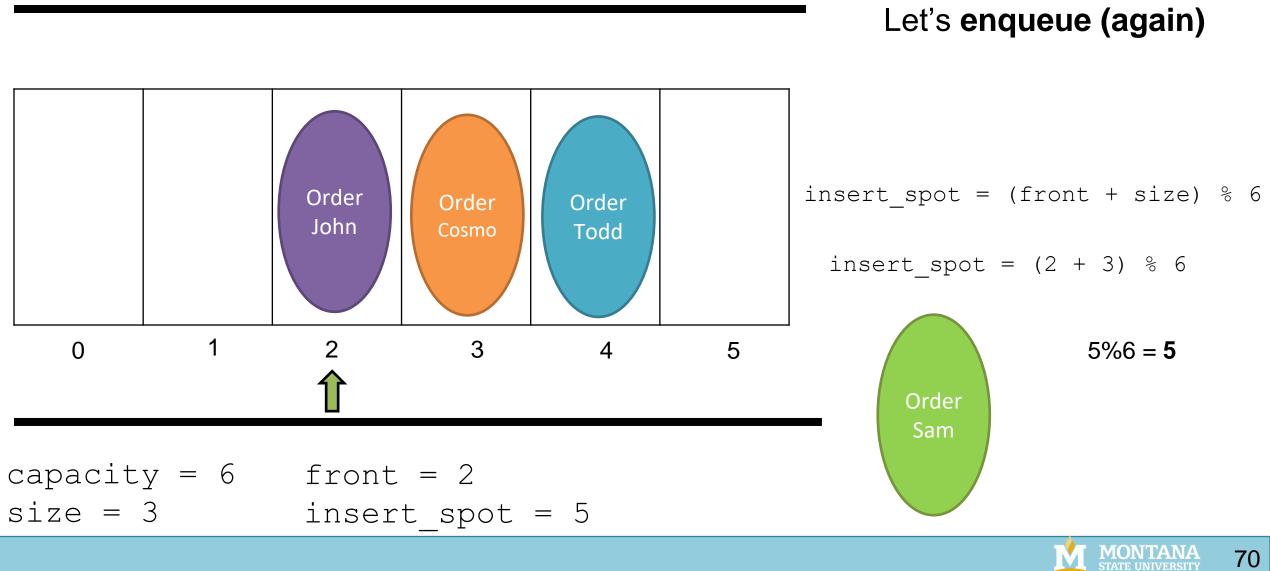


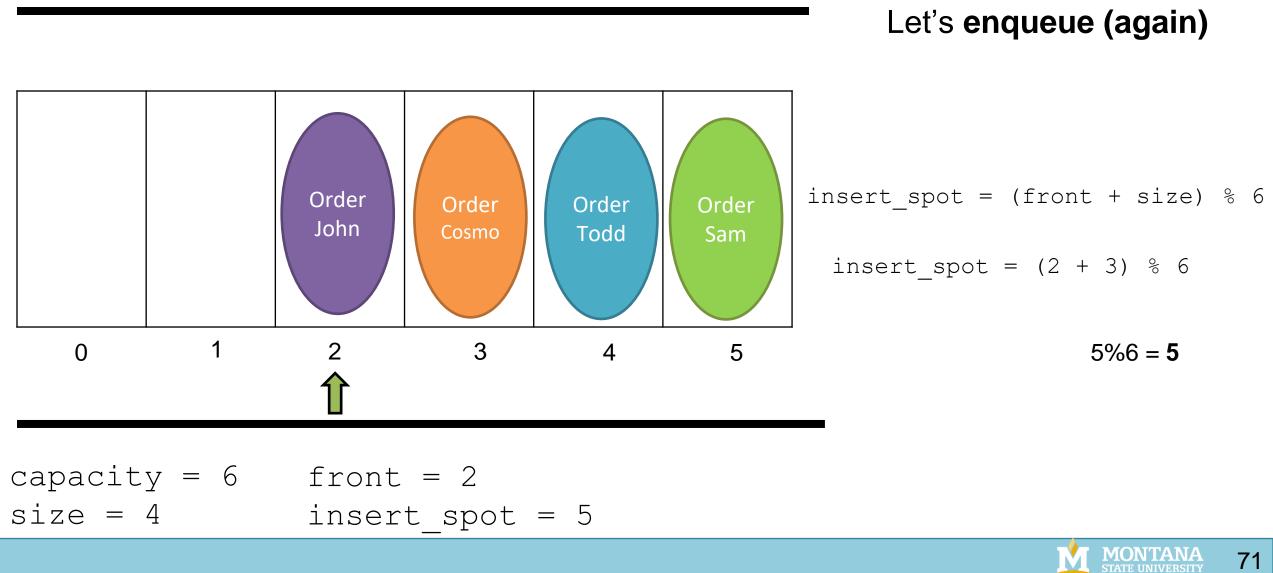
capacity = 6 front = 1 size = 4 insert_spot = 4



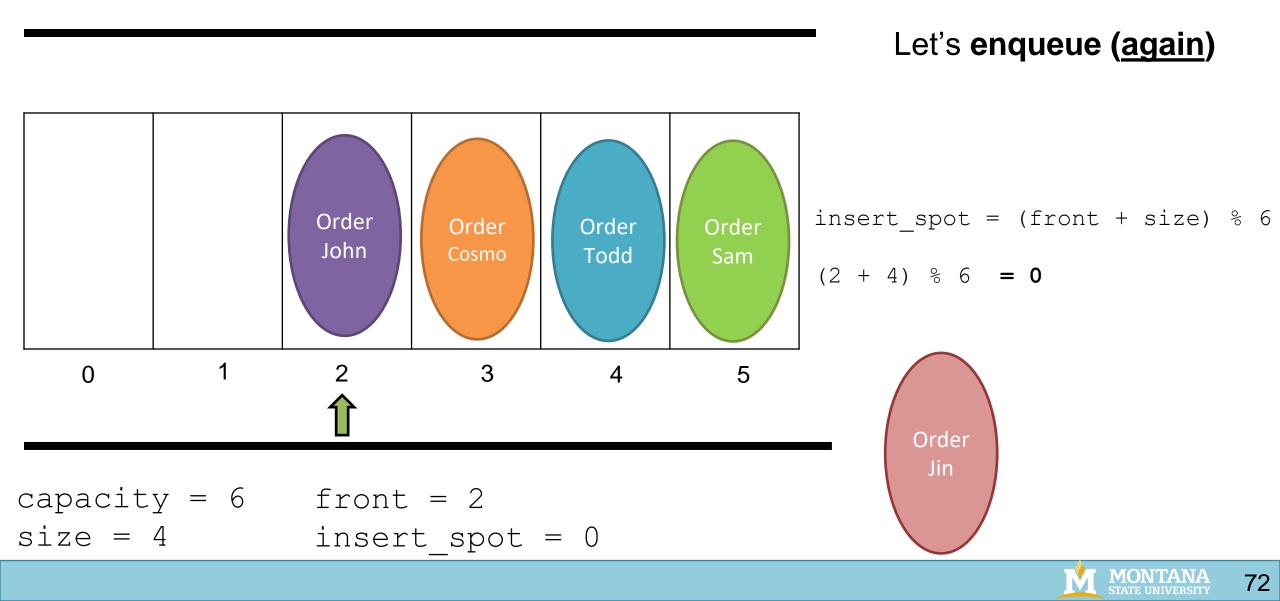


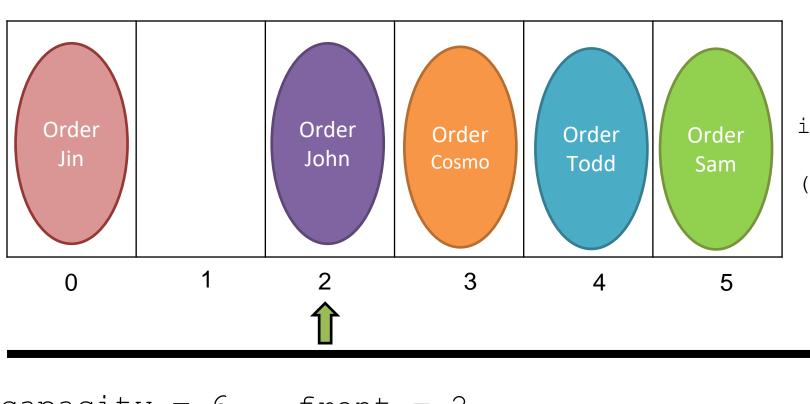
69





71

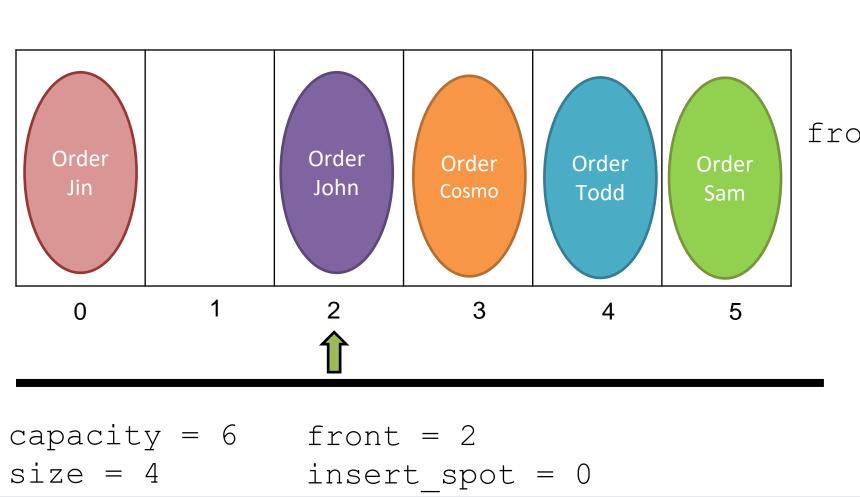




The modulus operator allows us to "**wrap around**" in our array!

capacity = 6 front = 2 size = 4 insert_spot = 0





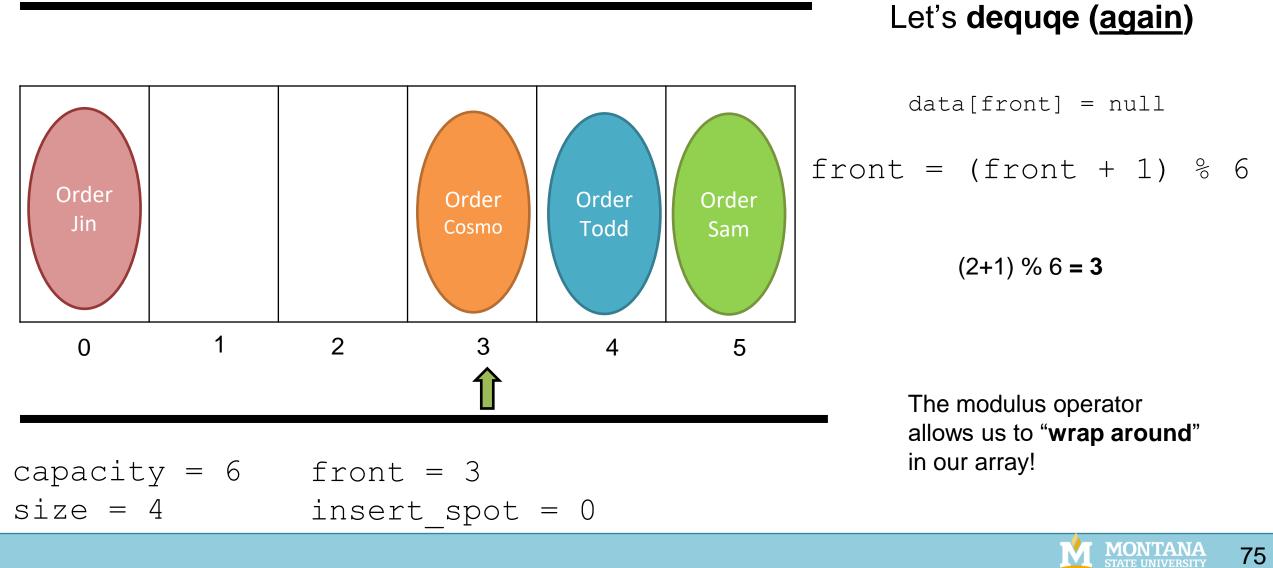
Let's dequqe (again)

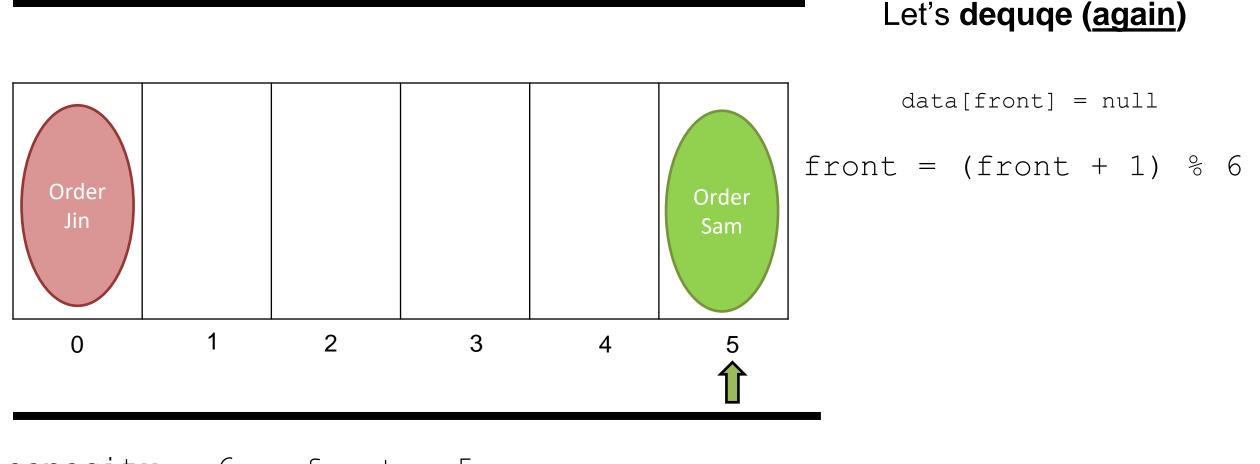
data[front] = null

front = (front + 1) % 6

The modulus operator allows us to "**wrap around**" in our array!

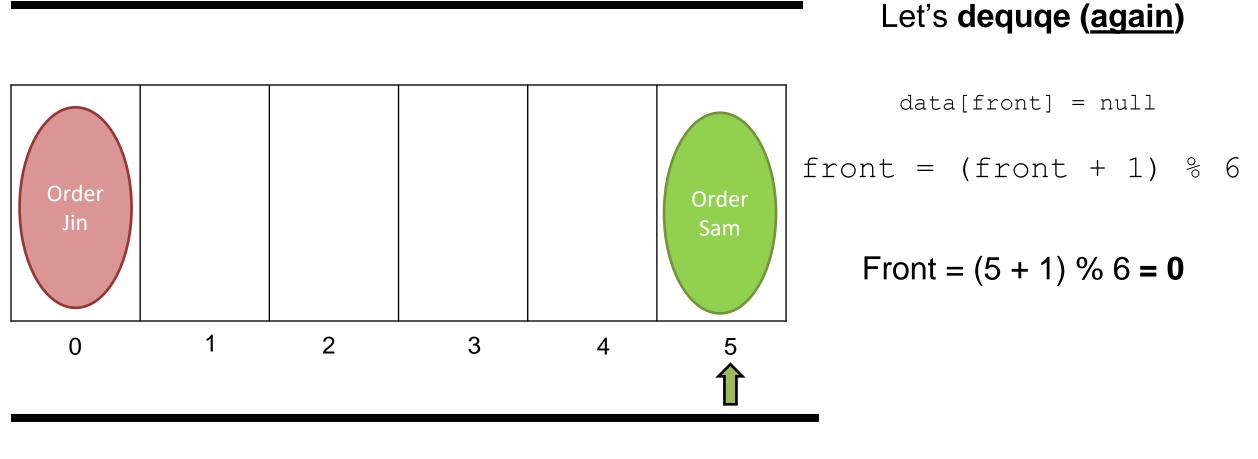




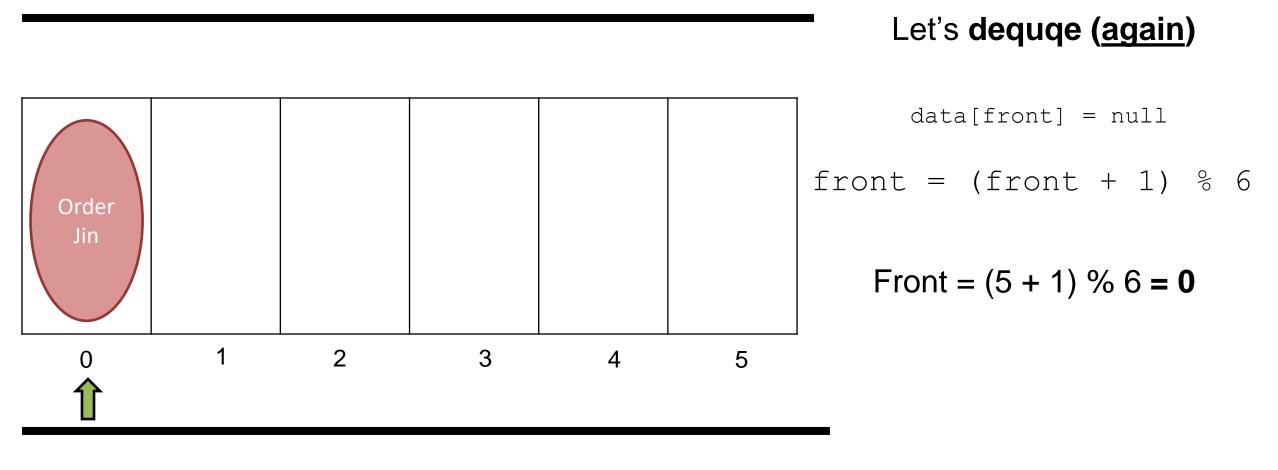


capacity = 6 front = 5
size = 2 insert_spot = 0









```
public void enqueue(Order newOrder) {
    if(this.size == this.data.length) {
        System.out.println("Queue is full");
    }
    int insert_spot = (front + size) % (this.data.length);
    data[insert_spot] = newOrder;
    this.size++;
    System.out.println("Added " +newOrder.getName() + " at index #" + insert_spot);
```

```
public void dequeue() {
    if(this.size == 0) {
        System.out.println("Queue is empty...");
        return;
    }
    else {
        Order o = this.data[front];
        this.data[front] = null;
        front = (front + 1) % this.data.length;
        this.size--;
        System.out.println(o.getName() + " order was removed ");
    }
}
```



```
public QueueLinkedList() {
    this.orders = new LinkedList<Order>();
    this.size = 0;
}
```

```
public QueueArray2() {
    this.orders = new Order[6];
    this.size = 0;
    this.front = 0;
    this.capacity = this.orders.length; //6
}
```

	Linked List	Array
Creation		
Enqueue		
Dequeue		
Peek		
Print Queue		

```
public QueueLinkedList() {
    this.orders = new LinkedList<Order>();
    this.size = 0;
} O(1)
```

```
public QueueArray2() {
    this.orders = new Order[6];
    this.size = 0;
    this.front = 0;
    this.capacity = this.orders.length; //6
} O(n), n = | array |
```

	Linked List	Array
Creation	O(1)	O(n)
Enqueue		
Dequeue		
Peek		
Print Queue		



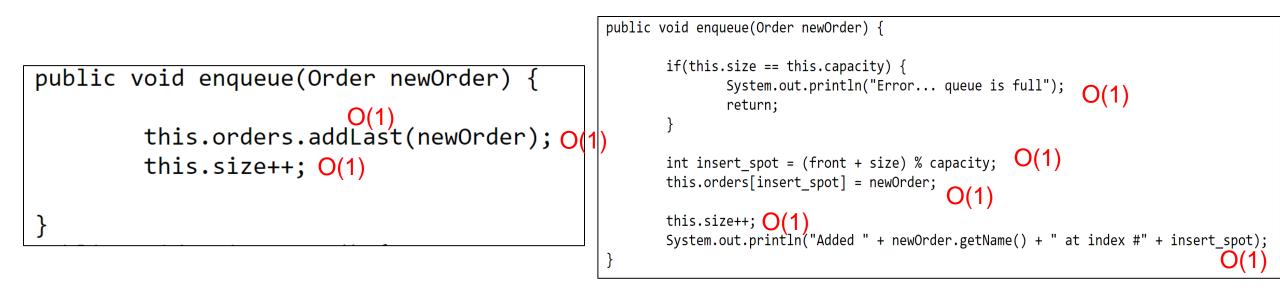
public void enqueue(Order newOrder) {

this.orders.addLast(newOrder);
this.size++;

```
public void enqueue(Order newOrder) {
    if(this.size == this.capacity) {
        System.out.println("Error... queue is full");
        return;
    }
    int insert_spot = (front + size) % capacity;
    this.orders[insert_spot] = newOrder;
    this.size++;
    System.out.println("Added " + newOrder.getName() + " at index #" + insert_spot);
}
```

	Linked List	Array
Creation	O(1)	O(n)
Enqueue		
Dequeue		
Peek		
Print Queue		





	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue		
Peek		
Print Queue		



```
public Order dequeue() {
    if(this.size != 0) {
        Order removed = this.orders.removeFirst();
        System.out.println(removed.getName() + "'s order
        size--;
        return removed;
    }
    else {
        return null;
    }
}
```

```
public void dequeue() {
    if(this.size == 0) {
        System.out.println("Error... queue is empty");
        return;
    }
    else {
        Order o = this.orders[front];
        this.orders[front] = null;
        front = (front + 1) % capacity;
        this.size--;
        System.out.println(o.getName() + "'s order was removed");
    }
}
```

	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue		
Peek		
Print Queue		



```
public Order dequeue() {
    if(this.size != 0) {
        Order removed = this.orders.removeFirst();
        O(1) System.out.println(removed.getName() + "'s order
        size--;
        return removed;
    }
    else {
        return null; O(1)
    }
}
```

```
public void dequeue() {
    if(this.size == 0) {
        System.out.println("Error... queue is empty"); O(1)
        return;
    }
    else {
        Order o = this.orders[front];
        this.orders[front] = null;
        front = (front + 1) % capacity; O(1)
        this.size--;
        System.out.println(o.getName() + "'s order was removed");
    }
}
```

	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue	O(1)	O(1)
Peek		
Print Queue		



return this.orders.getFirst()

return this.orders[front]

	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue	O(1)	O(1)
Peek		
Print Queue		



return this.orders.getFirst() O(1)

return this.orders[front] O(1)

	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue	O(1)	O(1)
Peek	O(1)	O(1)
Print Queue		

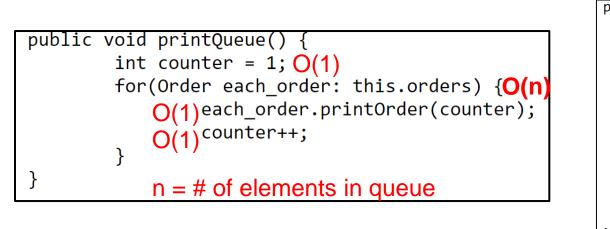


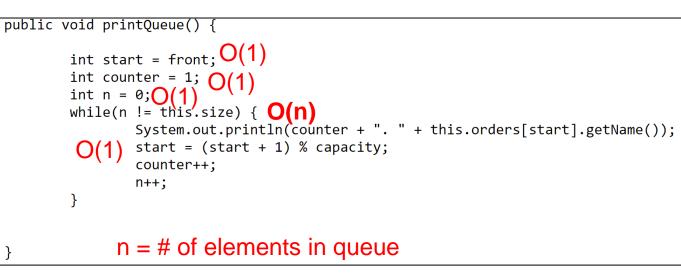
```
public void printQueue() {
    int counter = 1;
    for(Order each_order: this.orders) {
        each_order.printOrder(counter);
        counter++;
    }
}
```

```
public void printQueue() {
    int start = front;
    int counter = 1;
    int n = 0;
    while(n != this.size) {
        System.out.println(counter + ". " + this.orders[start].getName());
        start = (start + 1) % capacity;
        counter++;
        n++;
    }
}
```

	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue	O(1)	O(1)
Peek	O(1)	O(1)
Print Queue		







	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue	O(1)	O(1)
Peek	O(1)	O(1)
Print Queue	O(n)	O(n)



Takeaway: Adding and removing elements from a queue runs in constant time (O(1))

(FIFO)

(LIFO)

	Linked List	Array
Creation	O(1)	O(n)
Enqueue	O(1)	O(1)
Dequeue	O(1)	O(1)
Peek	O(1)	O(1)
Print Queue	O(n)	O(n)

Stack Runtime Analysis

Takeaway: Adding and removing elements from a **stack** runs in constant time $(\circ (1))$

	w/ Array	w/ Linked List
Creation	O(n)	O(1)
Push()	O(1)	O(1)
Pop()	O(1)	O(1)
peek()	O(1)	O(1)
Print()	O(n)	O(n)