

CSCI 132:

Basic Data Structures and Algorithms

Recursion (Part 3)

Reese Pearsall
Spring 2024

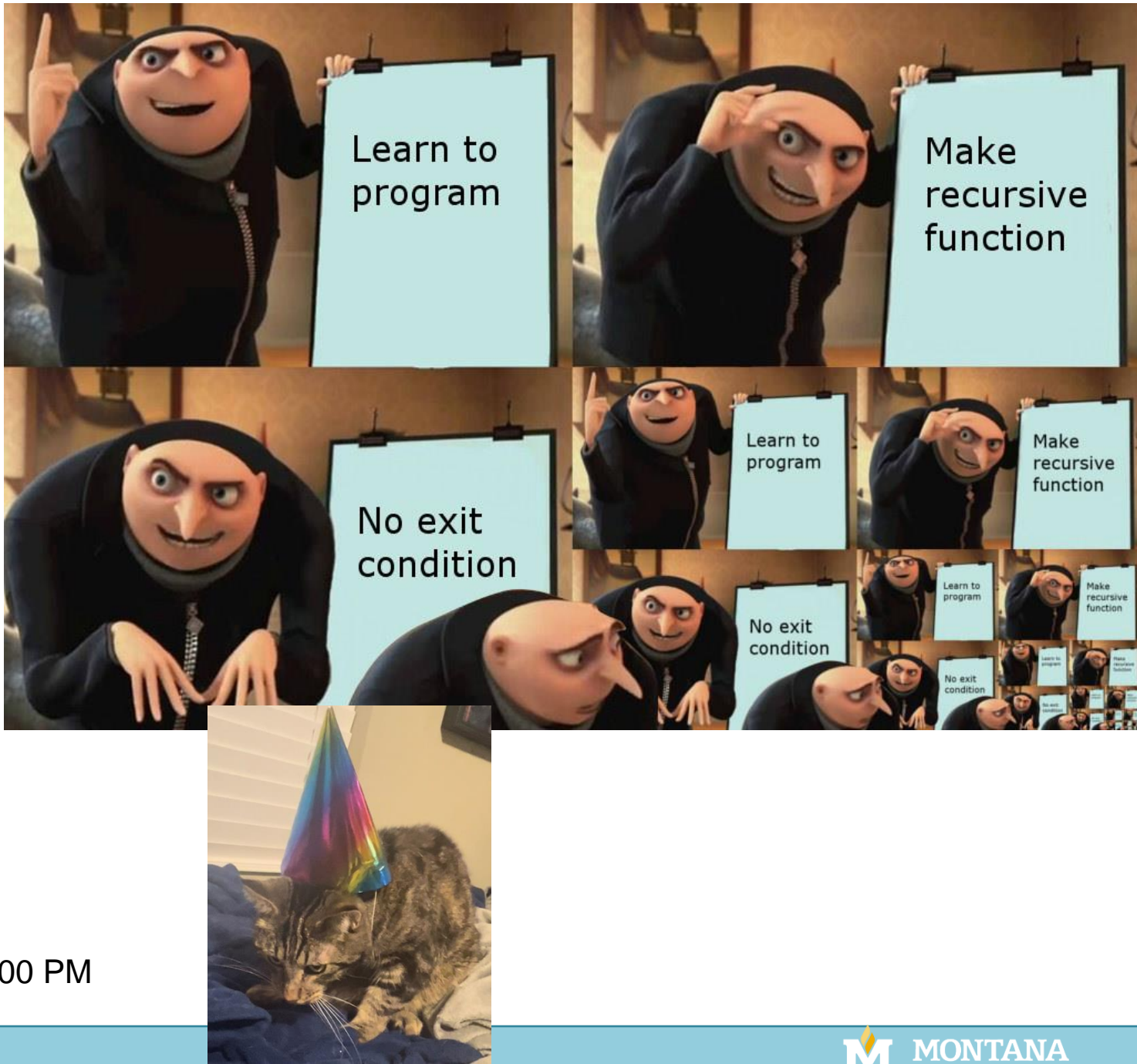
Announcements

No class on Wednesday
(4/17)

Program 4 due next
Friday



Gianforte Hall Groundbreaking Ceremony: 4/17 @ 2:00 PM



Change Making Problem

Given a set of coin denominations \mathbf{D} , how can you represent \mathbf{K} cents with the smallest number of coins?

Change Making Problem

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

$$D = [1, 5, 10, 25]$$

$$K = 37$$

Change Making Problem

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

$$D = [1, 5, 10, 25]$$

$$K = 37$$

Answer = 4

(Quarter, dime, two pennies)

Change Making Problem

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

$$D = [1, 5, 10, 25]$$

$$K = 37$$

Answer = 4

(Quarter, dime, two pennies)

Algorithm?

Change Making Problem

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

$$D = [1, 5, 10, 25]$$

$$K = 37$$

Answer = 4

(Quarter, dime, two pennies)

Use as many quarters as possible, then as many dimes as possible, ...

Change Making Problem

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

$$D = [1, 5, 10, 25]$$

$$K = 37$$

Answer = 4

(Quarter, dime, two pennies)

Use as many quarters as possible, then as many dimes as possible, ...

This is known as the **greedy** approach

Change Making Problem

Given a set a coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

$D = [1, 5, 10, 25]$

$K = 37$

Use as many quarters as possible, then as many dimes as possible, ...

Change Making Problem

Given a set of coin denominations **D**, how can you represent **K** cents with the smallest number of coins?

Greedy Algorithm

$D = [1, 5, 10, 18, 25]$

$K = 37$

Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ...

What if there were also an 18-cent coin?

Change Making Problem

Given a set of coin denominations \mathbf{D} , how can you represent \mathbf{K} cents with the smallest number of coins?

Greedy Algorithm

$D = [1, 5, 10, 18, 25]$

$K = 37$

Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes , ...

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?

Change Making Problem

Given a set of coin denominations D , how can you represent K cents with the smallest number of coins?

Greedy Algorithm

$D = [1, 5, 10, 18, 25]$

$K = 37$

Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes, ...

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?

Real Answer = 18, 18, 1 (3 coins)

Change Making Problem

Given a set of coin denominations \mathbf{D} , how can you represent \mathbf{K} cents with the smallest number of coins?

Greedy Algorithm

$D = [1, 5, 10, 18, 25]$

$K = 37$

Use as many quarters as possible, then as many 18 cent pieces as possible, then dimes, ...

25, 10, 1, 1 (4 coins)

What if there were also an 18-cent coin?

Real Answer = 18, 18, 1 (3 coins)

Lesson Learned: The Greedy approach works for the United States denominations, but not for a general set of denominations

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))

$$25 + 25 + 10 + 1 + 1 + 1 = 63$$



What can you conclude?

Does this provide an answer to any other change making problems?

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))



This is the minimum coins needed to make 38 cents

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))

$$25 + 25 + 10 + 1 + 1 + 1 = 63$$



This is the minimum coins needed to make 13 cents

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))

$$25 + 25 + 10 + 1 + 1 + 1 = 63$$


The image shows three US coins: two quarters (25 cents) and one dime (10 cents), followed by three pennies (1 cent each). The three pennies are enclosed in a red rectangular box. The coins are arranged horizontally below the equation.

This is the minimum coins needed to make 63 cents

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))

$$25 + 25 + 10 + 1 + 1 + 1 = 63$$



This is the minimum coins needed to make 63 cents

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))

$$25 + 25 + 10 + 1 + 1 + 1 = 63$$



This is the minimum coins needed to make 1 cent

Change Making Problem

Suppose I tell you that 2 quarters, 1 dime, and 3 pennies are the minimum number of coins needed to make **63 cents**

(We will assume we have the standard US denominations [1, 5, 10, 25] (NO 50 CENT PIECE))

$$25 + 25 + 10 + 1 + 1 + 1 = 63$$



The solution to the change making problems consists of solutions to smaller change making problems

We can use **recursion** to solve this problem

Change Making Problem

In general, suppose a country has coins with denominations:

$$1 = d_1 < d_2 < \cdots < d_k \quad (\text{US coins: } d_1 = 1, d_2 = 5, d_3 = 10, d_4 = 25)$$

Algorithm: To make change for p cents, we are going to figure out change for every value $x < p$. We will build solution for p out of smaller solutions.

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$

We used one quarter

Now find the minimum number
of coins needed to make 12
cents

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + \underbrace{C(12)}_{\text{We used one dime}}$$
$$C(12) = 1 + C(2)$$

Now find the minimum number of coins needed to make 2 cents

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12) \img alt="Quarter coin" data-bbox="345 565 398 655"/>$$

$$C(12) = 1 + C(2) \img alt="Dime coin" data-bbox="555 660 602 745"/>$$

$$C(2) = 1 + C(1) \img alt="Penny coin" data-bbox="755 745 802 830"/>$$

$$C(1) = 1 + C(0) \img alt="Penny coin" data-bbox="955 855 1000 930"/>$$

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12) \img alt="Dime coin" data-bbox="345 565 398 655"/>$$

$$C(12) = 1 + C(2) \img alt="Dime coin" data-bbox="555 660 602 745"/>$$

$$C(2) = 1 + C(1) \img alt="Penny coin" data-bbox="755 745 802 830"/>$$

$$C(1) = 1$$

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$



$$C(12) = 1 + C(2)$$



$$C(2) = 1 + 1$$

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$



$$C(12) = 1 + C(2)$$



$$C(2) = 2$$

Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$



$$C(12) = 1 + 2$$



Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + C(12)$$
$$C(12) = 3$$



Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 1 + 3$$



Change Making Problem

$C(p)$ – minimum number of coins to make p cents.

x – value (e.g. \$0.25) of a coin used in the optimal solution.

$$C(p) = 1 + C(p - x).$$

$$C(37) = 4$$

The minimum number of coins needed to make 37 cents is 4

Change Making Problem

In general, suppose a country has coins with denominations:

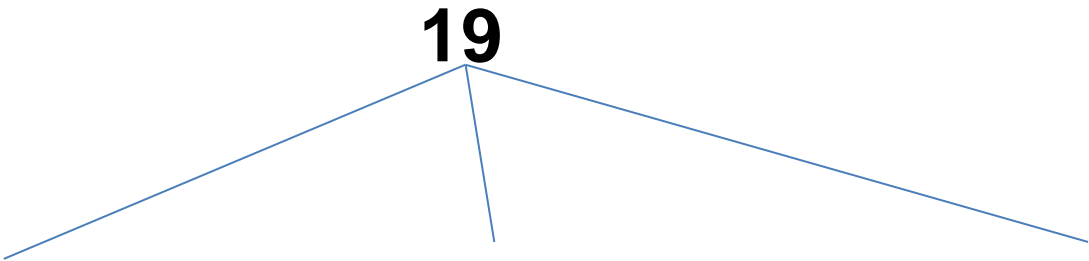
$$1 = d_1 < d_2 < \cdots < d_k \quad (\text{US coins: } d_1 = 1, d_2 = 5, d_3 = 10, d_4 = 25)$$

(This algorithm must work for ALL denominations)

Algorithm: To make change for p cents, we are going to figure out change for every value $x < p$. We will build solution for p out of smaller solutions.

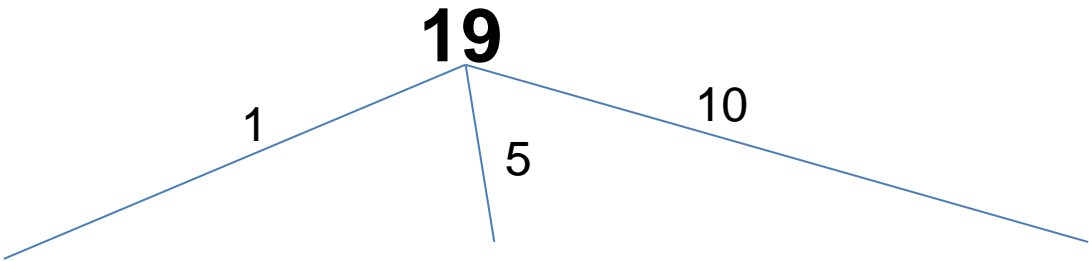
Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10



Change Making Problem

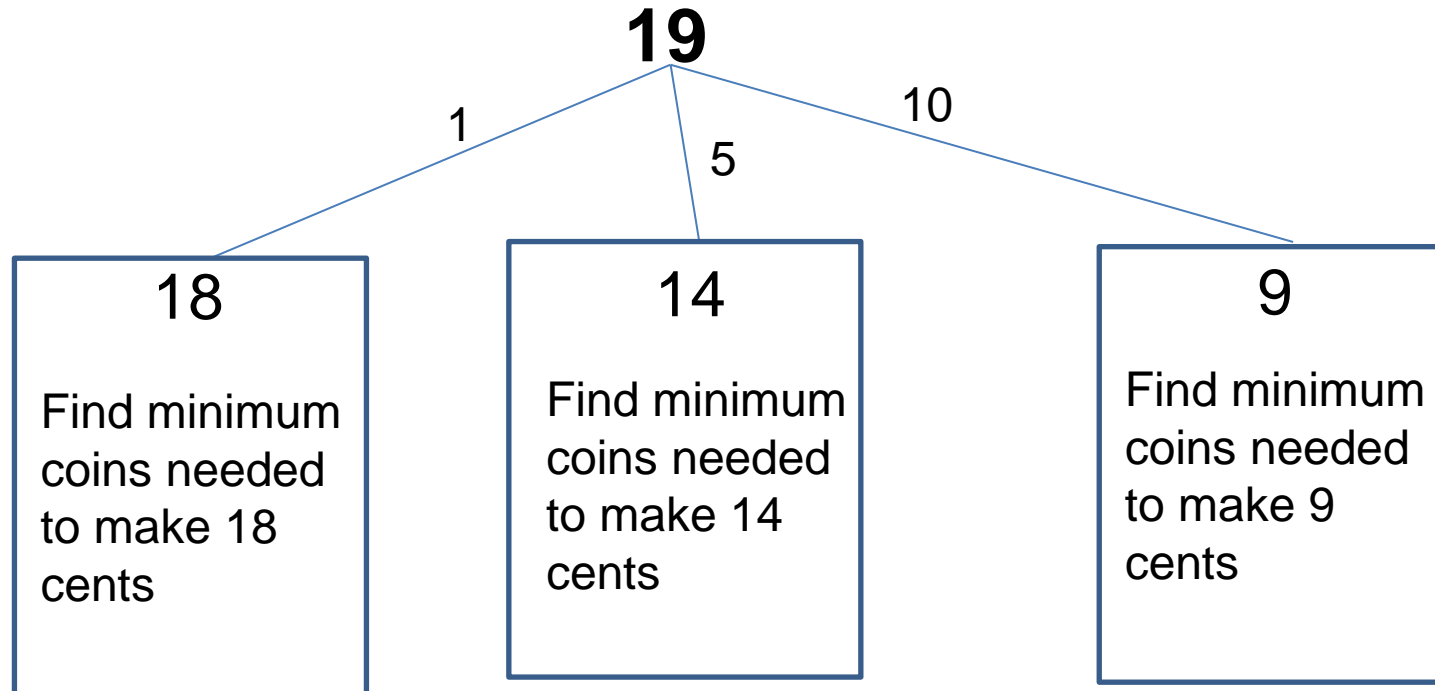
Make \$0.19 with \$0.01, \$0.05, \$0.10



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations

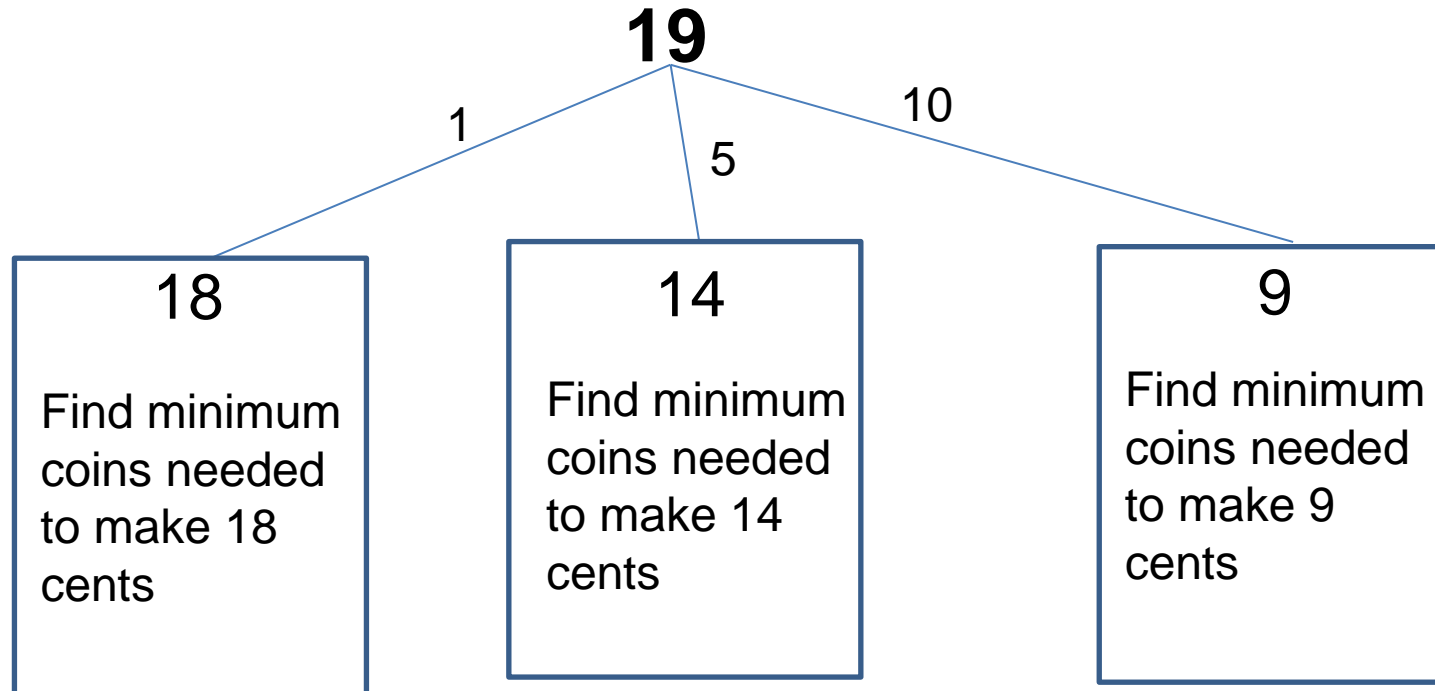


To find the minimum number of coins needed to create 19 cents, we generate k subproblems

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations

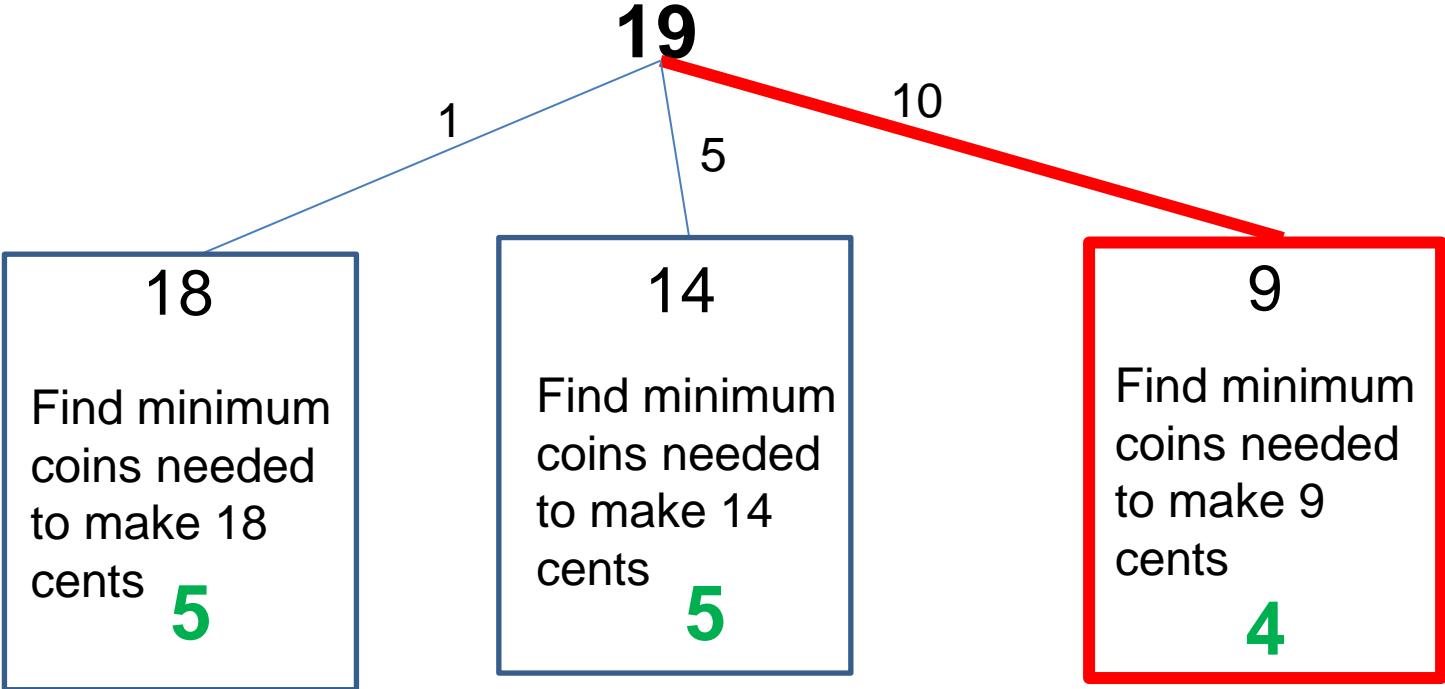


We want to select the **minimum** solution of these three subproblems

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations

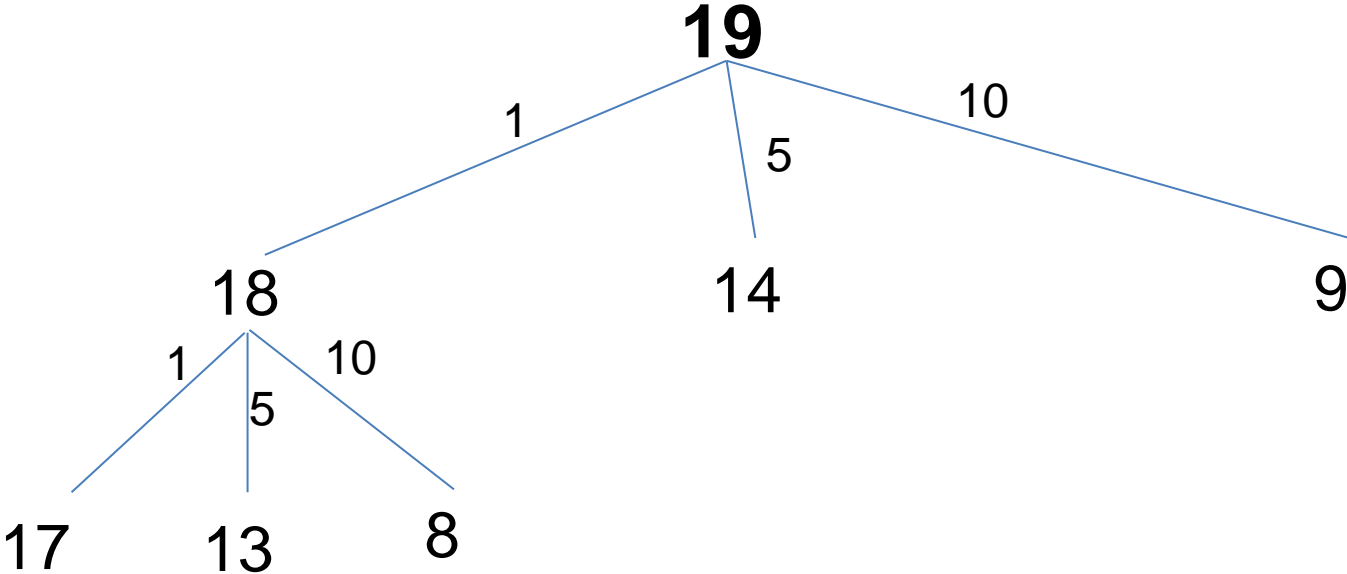


For the solution of our original problem (19), we want to select this branch (one dime used)

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations



Find minimum
coins needed
to make 17
cents

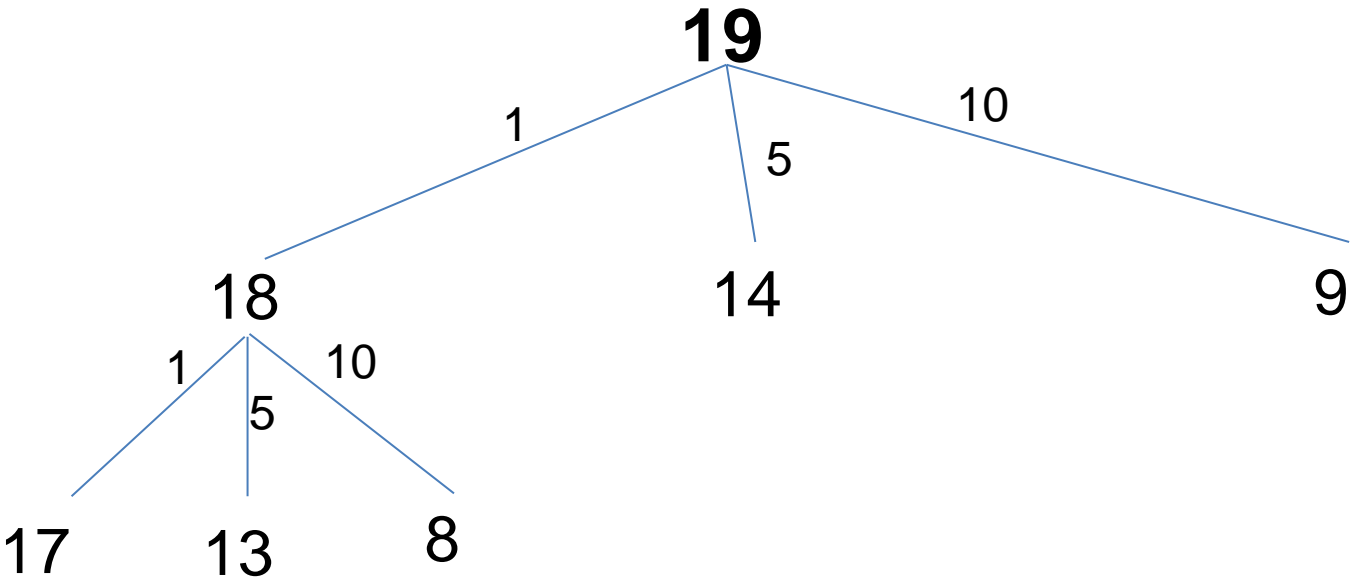
Find minimum
coins needed
to make 13
cents

Find minimum
coins needed
to make 8
cents

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

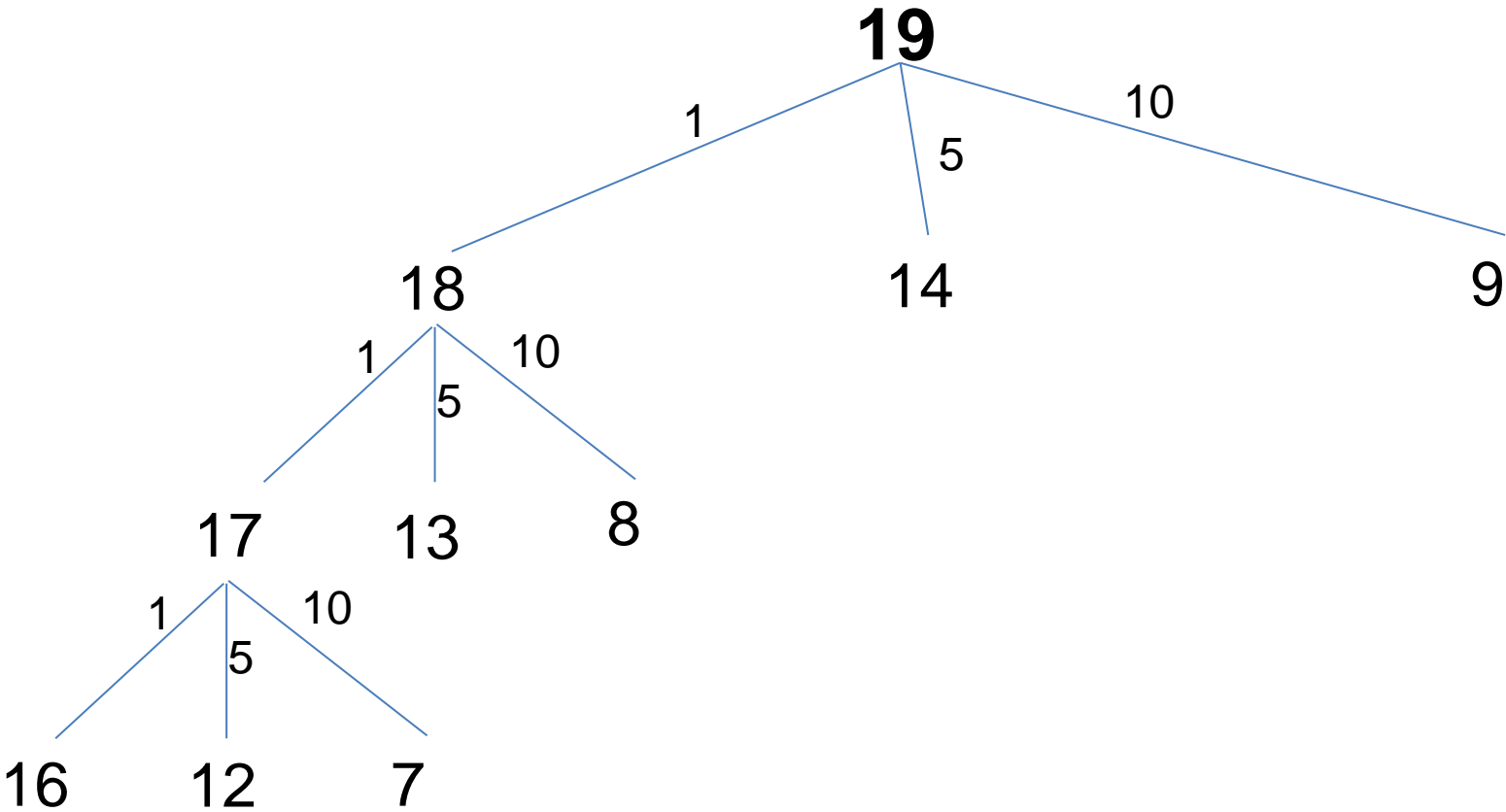
k = # denominations



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

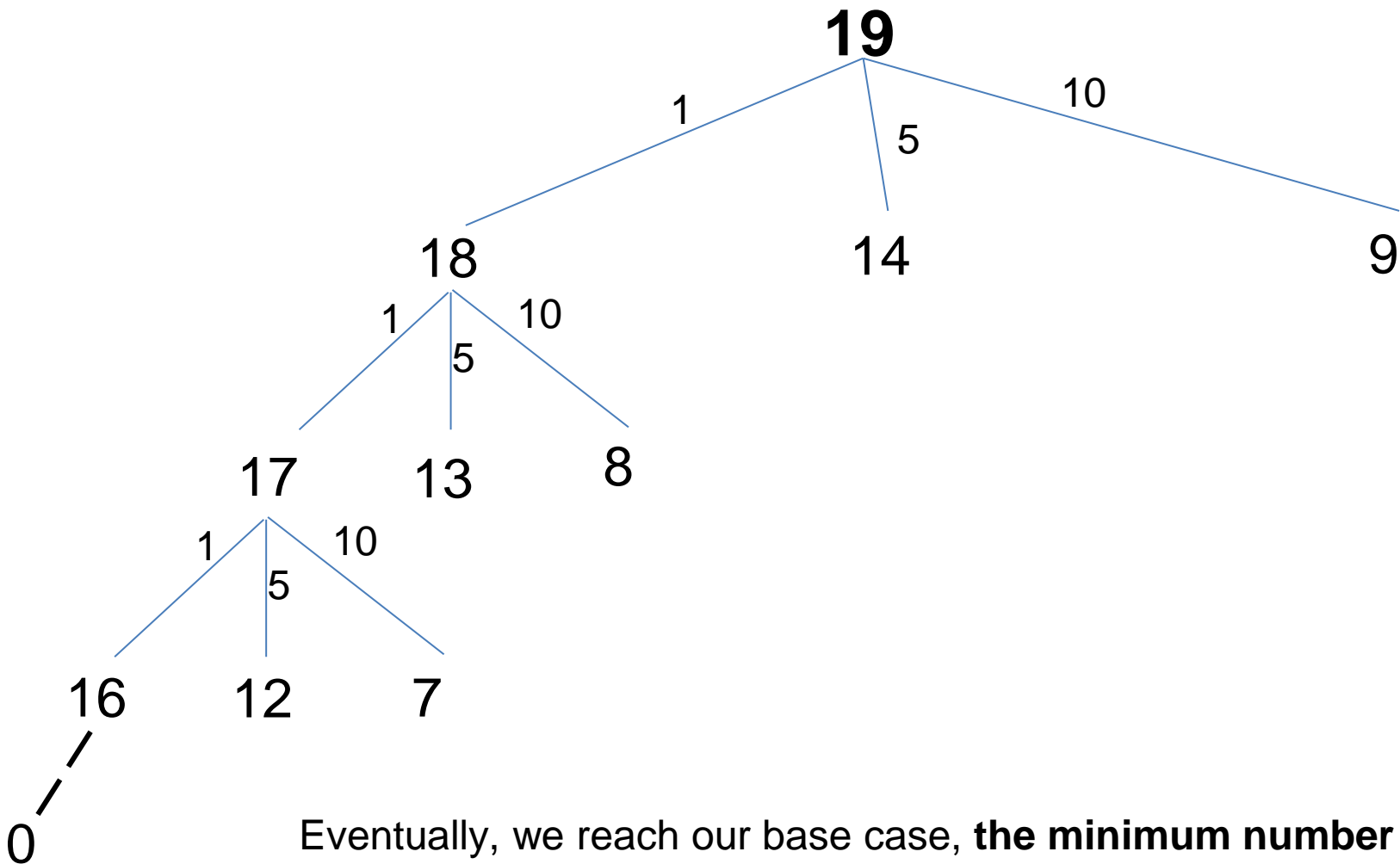
k = # denominations



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations

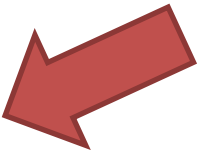
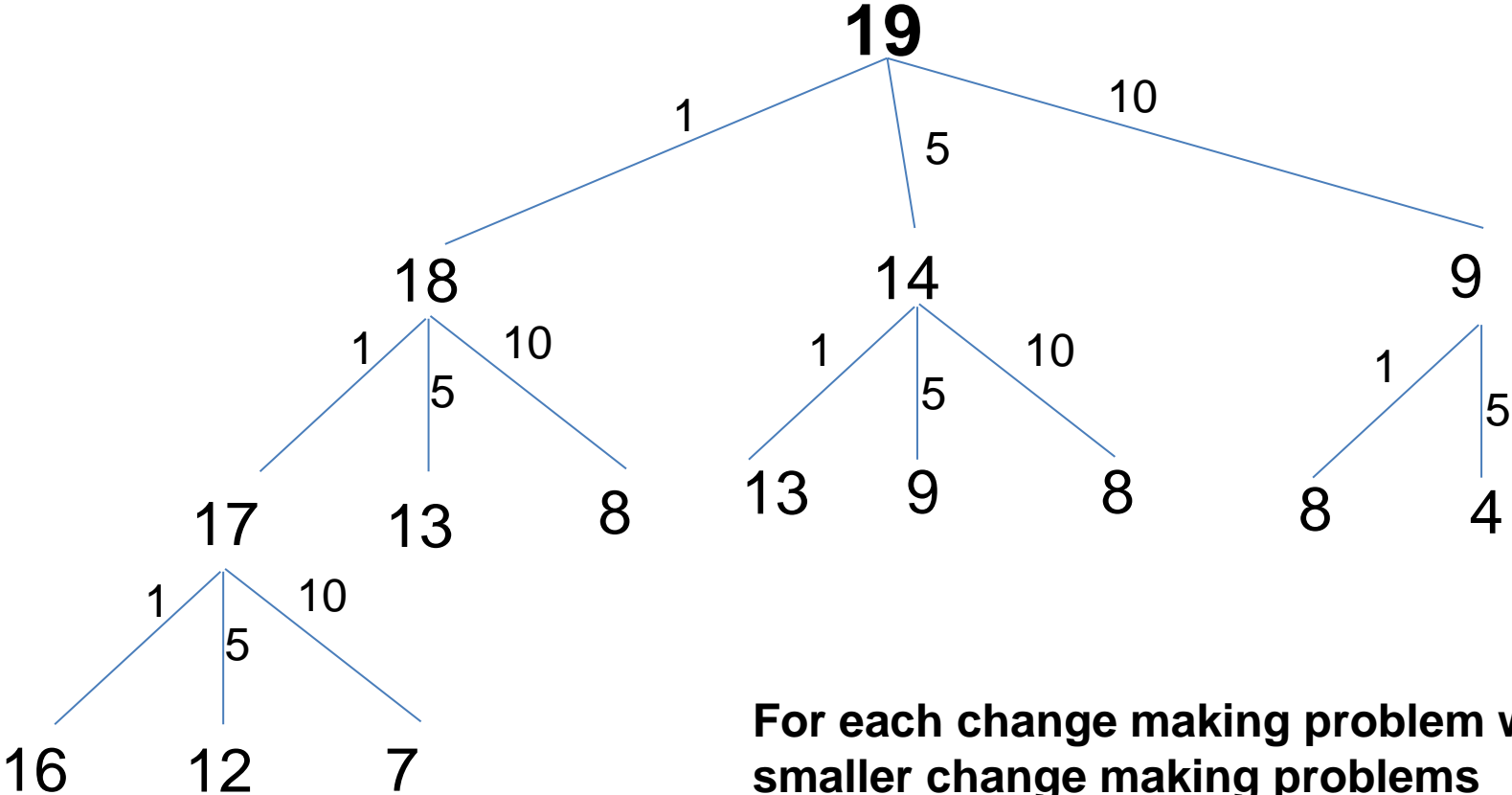


Eventually, we reach our base case, **the minimum number of coins needed to make 0 cents**

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations



When C(9), we cant use a 10 cent piece...

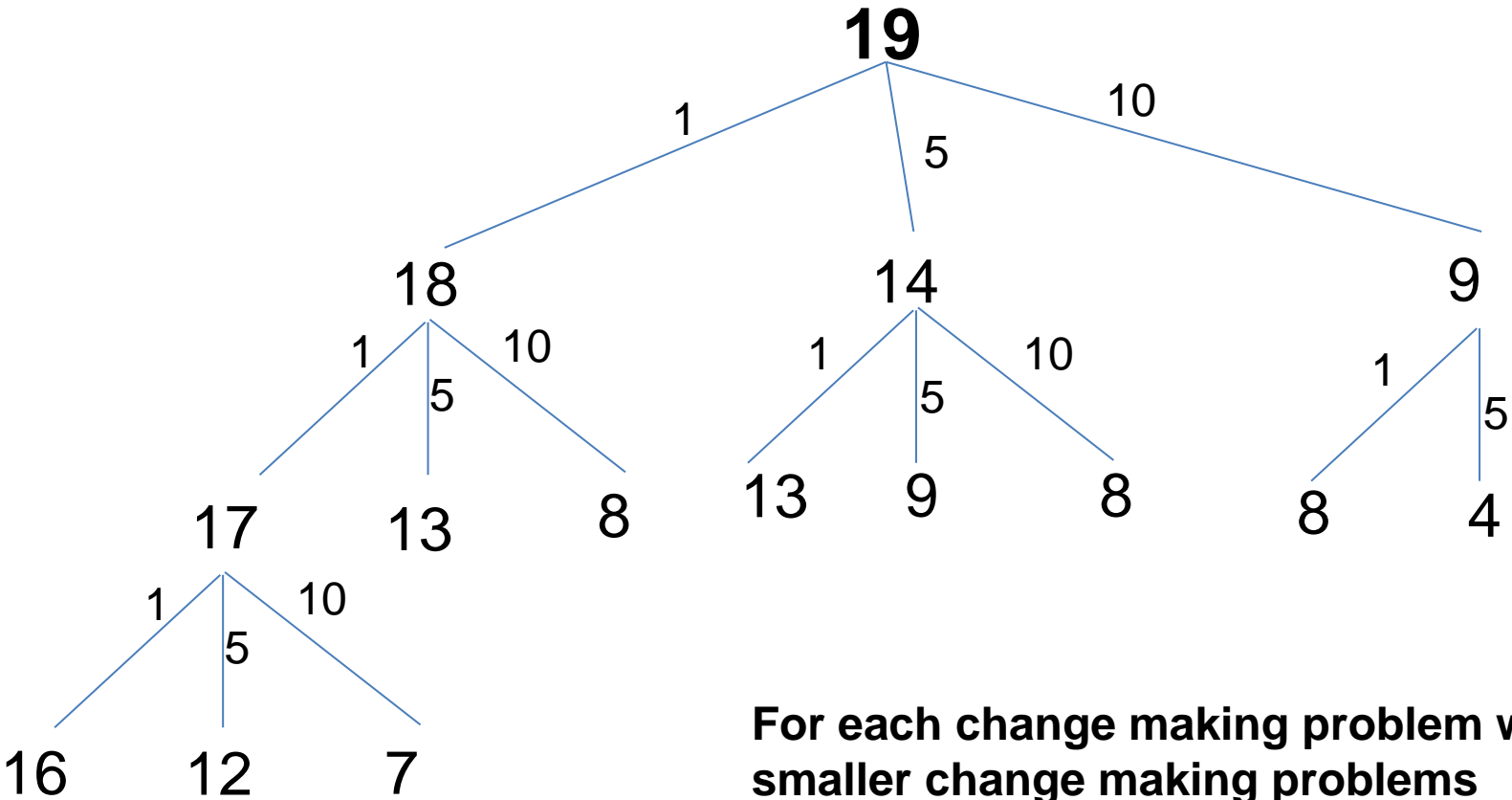
For each change making problem we solve, we must solve at most 3 smaller change making problems

Once we solve the smaller problems, we must select the branch that has the minimum value

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations



When C(9), we cant use a 10 cent piece...

For each change making problem we solve, we must solve at most 3 smaller change making problems

Once we solve the smaller problems, we must select the branch that has the minimum value



Change Making Problem

$$C(p) = \begin{cases} \min_{i:d_i \leq p} C(p - d_i) + 1, & p > 0 \\ 0, & p = 0 \end{cases}$$

Least change for 19 cents = minimum of:

- least change for 19-10 = 9 cents
- least change for 19-5 = 14 cents
- least change for 19-1 = 18 cents

For each problem P, we will solve the problem for (P – d), where d represents each possible denomination

Change Making Problem

$$C(p) = \begin{cases} \min_{i:d_i \leq p} C(p - d_i) + 1, & p > 0 \\ 0, & p = 0 \end{cases}$$

Least change for 19 cents = minimum of:

- least change for 19-10 = 9 cents
- least change for 19-5 = 14 cents
- least change for 19-1 = 18 cents

For each problem P, we will solve the problem for (P – d), where d represents each possible denomination

We want to select only the branch that yields the minimum value

Change Making Problem

$$C(p) = \begin{cases} \min_{i:d_i \leq p} C(p - d_i) + 1, & p > 0 \\ 0, & p = 0 \end{cases}$$

Least change for 19 cents = minimum of:

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- least change for 19-5 = 14 cents
- least change for 19-1 = 18 cents

For each problem P, we will solve the problem for (P – d), where d represents each possible denomination

If we ever need to make change for 0 cents, return 0

We want to select only the branch that yields the minimum value

Change Making Problem

$$C(p) = \begin{cases} \min_{i:d_i \leq p} C(p - d_i) + 1, & p > 0 \\ 0, & p = 0 \end{cases}$$

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For each problem P, we will solve the problem for (P – d), where d represents each possible denomination

If we ever need to make change for 0 cents, return 0

We want to select only the branch that yields the minimum value

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
```

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
```

```
    if p == 0  
        return 0;
```

Base Case

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
```

```
if p == 0  
    return 0;
```

```
else  
    min =  $\infty$   
    a =  $\infty$ 
```

Base Case

```
int min = Integer.MAX_VALUE;  
int a = Integer.MAX_VALUE;;
```

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
```

```
if p == 0  
    return 0;
```

Base Case

```
else  
    min =  $\infty$   
    a =  $\infty$ 
```

```
int min = Integer.MAX_VALUE;  
int a = Integer.MAX_VALUE;;
```

```
for each  $d_i$  in D  
    if (p -  $d_i$ ) >= 0  
        a = min_coins(D, p -  $d_i$ )
```

Recurse, and find the minimum number of coins needed using each valid denomination

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
```

```
if p == 0  
    return 0;
```

Base Case

```
else
```

```
    min =  $\infty$   
    a =  $\infty$ 
```

```
int min = Integer.MAX_VALUE;  
int a = Integer.MAX_VALUE;;
```

```
    for each  $d_i$  in D  
        if  $(p - d_i) \geq 0$   
            a = min_coins(D, p -  $d_i$ )  
            if a < min  
                min = a
```

Recurse, and find the minimum number of coins needed using each valid denomination

Select the branch that has the minimum value

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
    if p == 0
        return 0;
    else
        min = ∞
        a = ∞
```

Base Case

```
int min = Integer.MAX_VALUE;
int a = Integer.MAX_VALUE;;
```

```
    for each  $d_i$  in D
        if  $(p - d_i) \geq 0$ 
             $a = \text{min\_coins}(D, p - d_i)$ 
        if  $a < \text{min}$ 
             $\text{min} = a$ 
```

Recurse, and find the minimum number of coins needed using each valid denomination

Select the branch that has the minimum value

Change Making Problem

D = array of denominations [1, 5, 10, 18, 25]
p = desired change (37)

```
min_coins(D, p)
    if p == 0
        return 0;
    else
        min = ∞
        a = ∞
```

Base Case

```
int min = Integer.MAX_VALUE;
int a = Integer.MAX_VALUE;;
```

```
    for each  $d_i$  in D
        if  $(p - d_i) \geq 0$ 
             $a = \text{min\_coins}(D, p - d_i)$ 
            if  $a < \text{min}$ 
                 $\text{min} = a$ 
```

Recurse, and find the minimum number of coins needed using each valid denomination

Select the branch that has the minimum value

```
    return 1 + min
```

Once, our for loop finishes, we should know the branch that had the minimum, so return $(1 + \text{min})$, 1 because one coin was used in the current method call

Change Making Problem

```
min_coins(D, p)
    if p == 0
        return 0;
    else
        min =  $\infty$ 
        a =  $\infty$ 

        for each  $d_i$  in D
            if (p -  $d_i$ )  $\geq$  0
                a = min_coins(D, p -  $d_i$ )
                if a < min
                    min = a

    return 1 + min
```

Change Making Problem

```
min_coins(D, p)
    if p == 0
        return 0;
    else
        min =  $\infty$ 
        a =  $\infty$ 

        for each  $d_i$  in D
            if  $(p - d_i) \geq 0$ 
                a = min_coins(D, p -  $d_i$ )
                if a < min
                    min = a

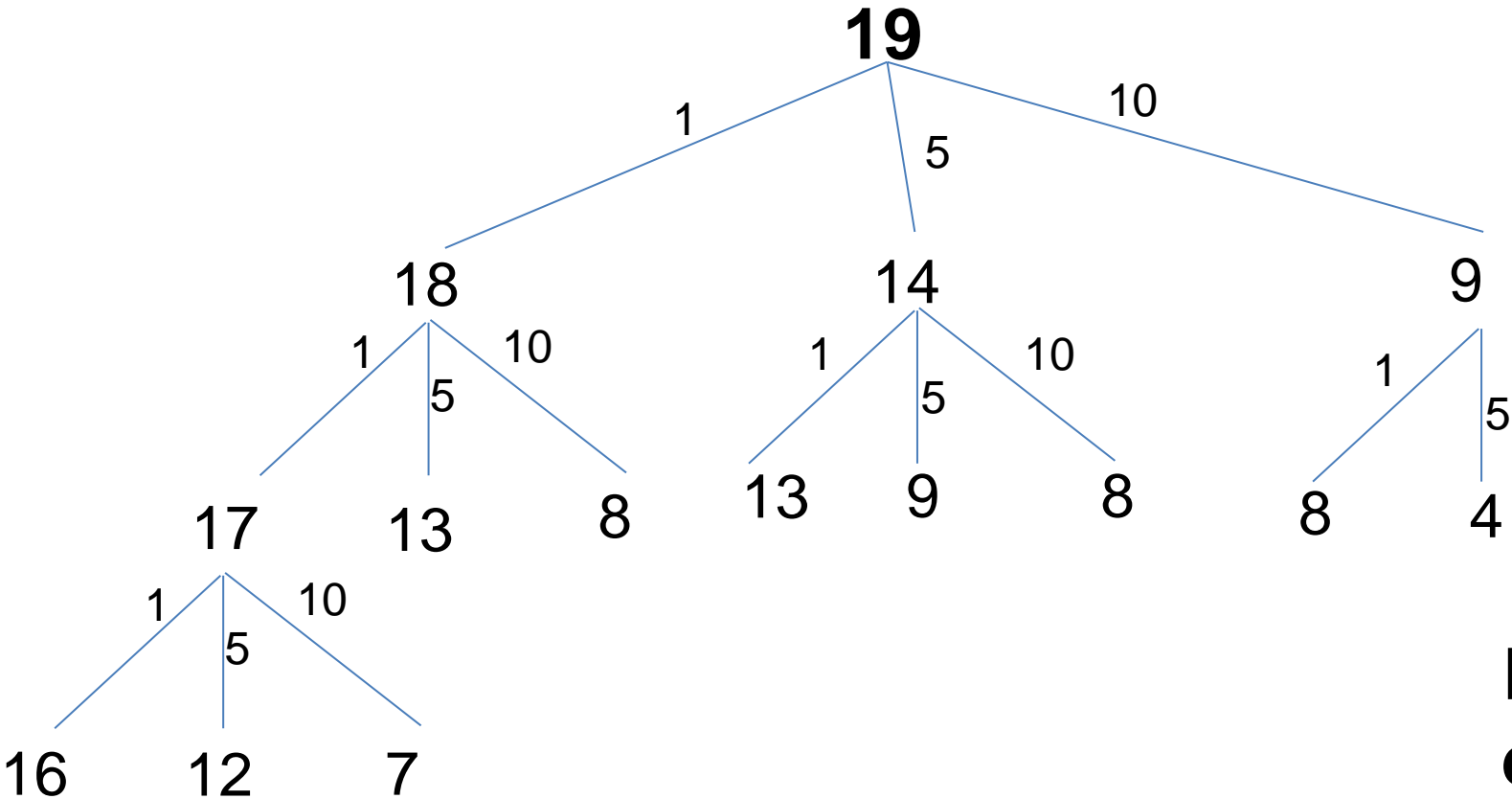
    return 1 + min
```

Running time?

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations
p = value to make change for

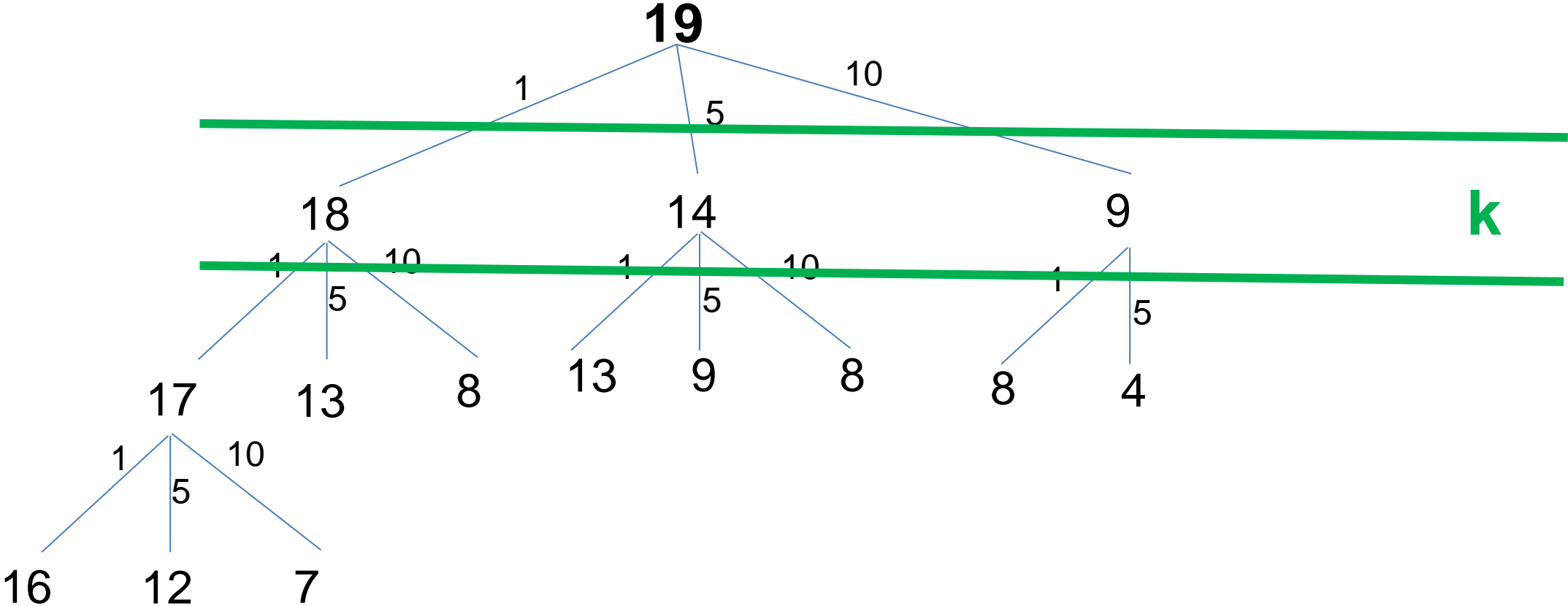


For sufficiently large p,
every permutation of
denominations is
included.

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

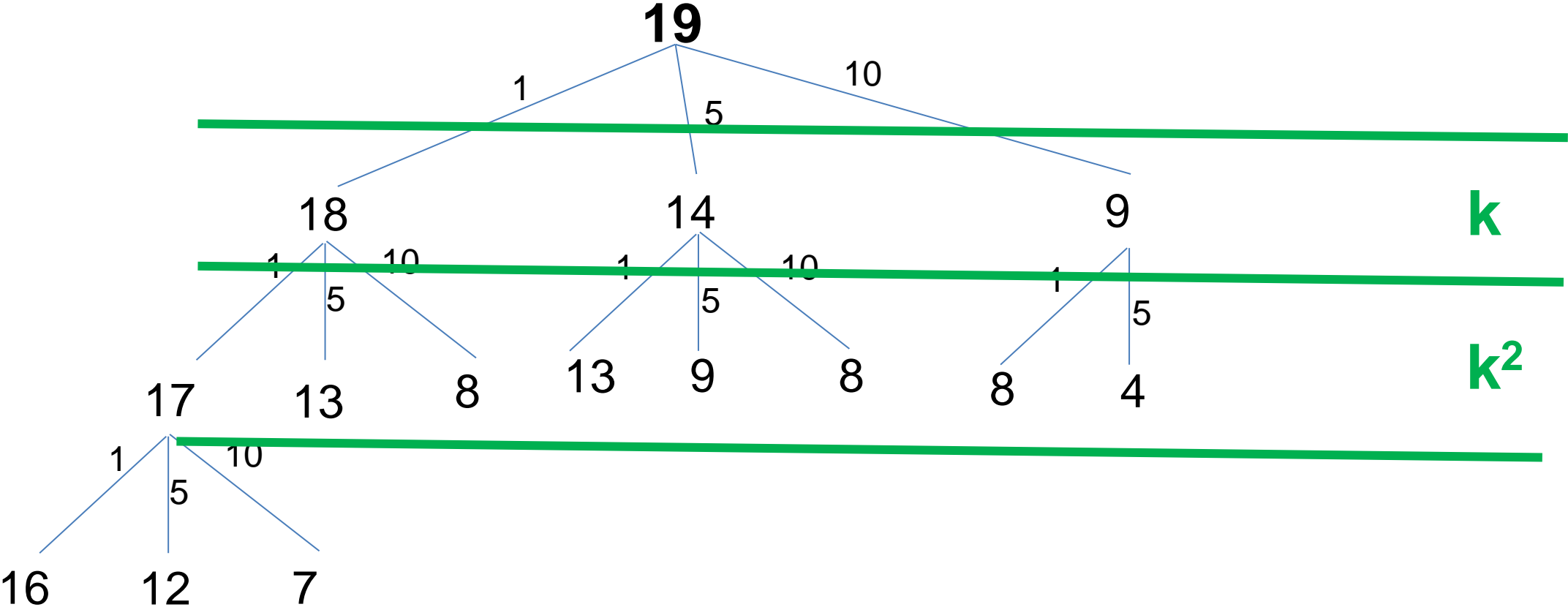
k = # denominations
p = value to make change for



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

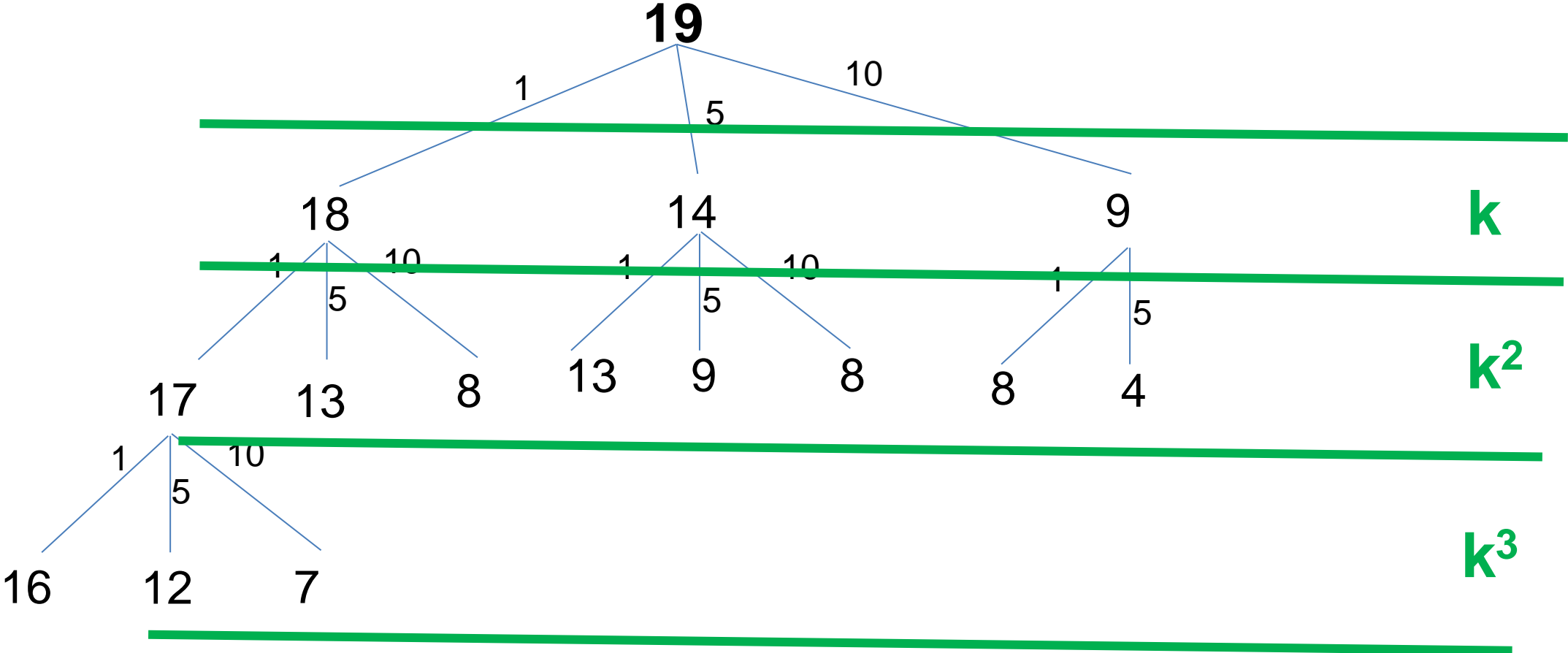
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 p = value to make change for



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

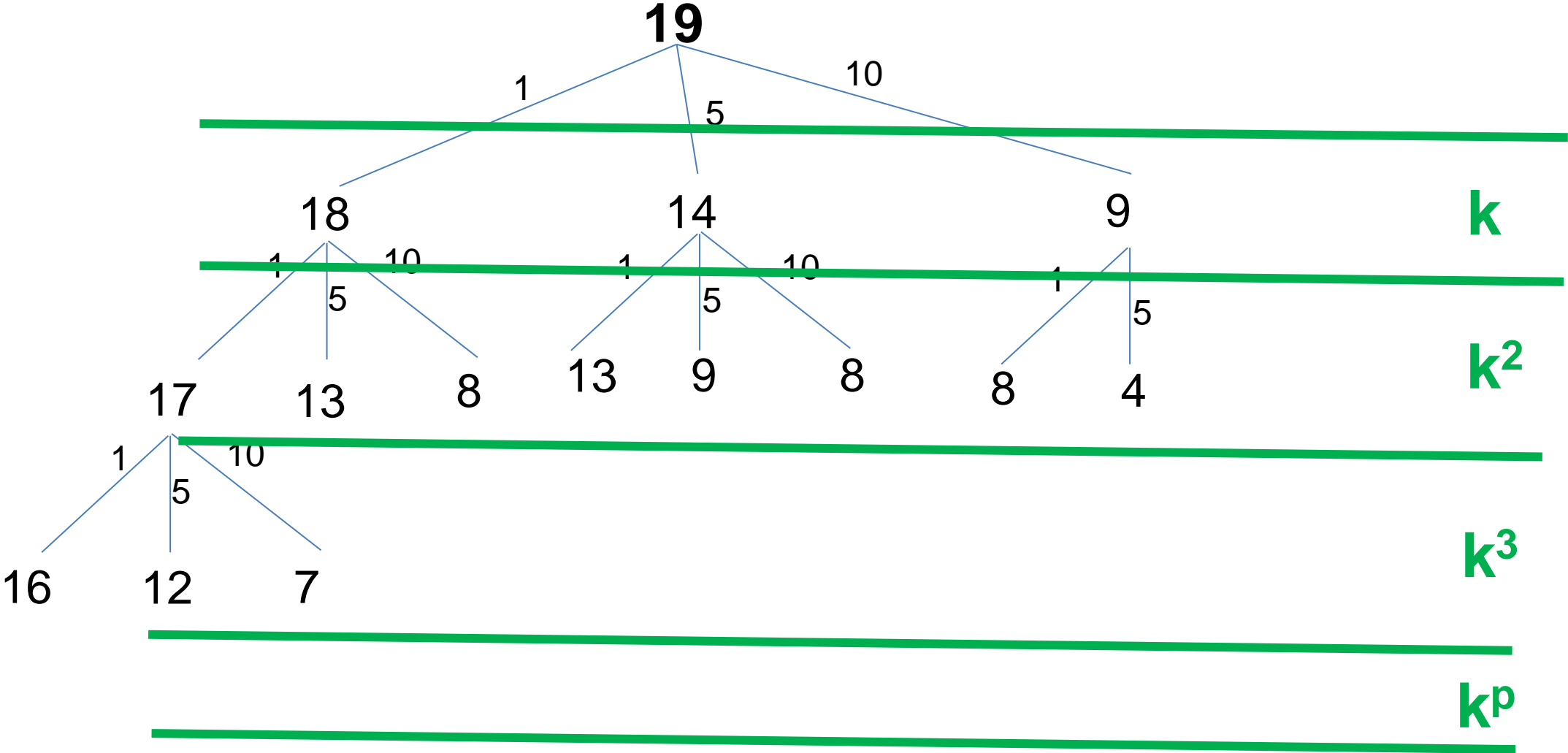
k = # denominations
 p = value to make change for



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations
 p = value to make change for



Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations
 p = value to make change for

10

As k and p both grow, the number of recursive calls being made will grow **exponentially**

If we have a lot of coin denominations, we will have **a lot** of branching

k

k^2

k^3

k^p

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations
 p = value to make change for

10

As k and p both grow, the number of recursive calls being made will grow **exponentially**

Running time: $O(\text{skull})$

k

k^2

k^3

k^p

Change Making Problem

Make \$0.19 with \$0.01, \$0.05, \$0.10

k = # denominations
 p = value to make change for

10

As k and p both grow, the number of recursive calls being made will grow **exponentially**

Running time: probably $O(k^p)$ or $O(k!)$

For a large set of denominations, or a large p , this algorithm will take a long time to run

k

k^2

k^3

k^p

Change Making Problem

Combinations for 7 cents

Using $D = [1, 5, 10]$

[1, 1, 5]

[1, 1, 1, 1, 1, 1, 1]

[1, 1, 5] and [5, 1, 1] is the
same combination...

Permutations for 7 cents

Using $D = [1, 5, 10]$

[1, 1, 5]

[5, 1, 1]

[1, 5, 1]

...

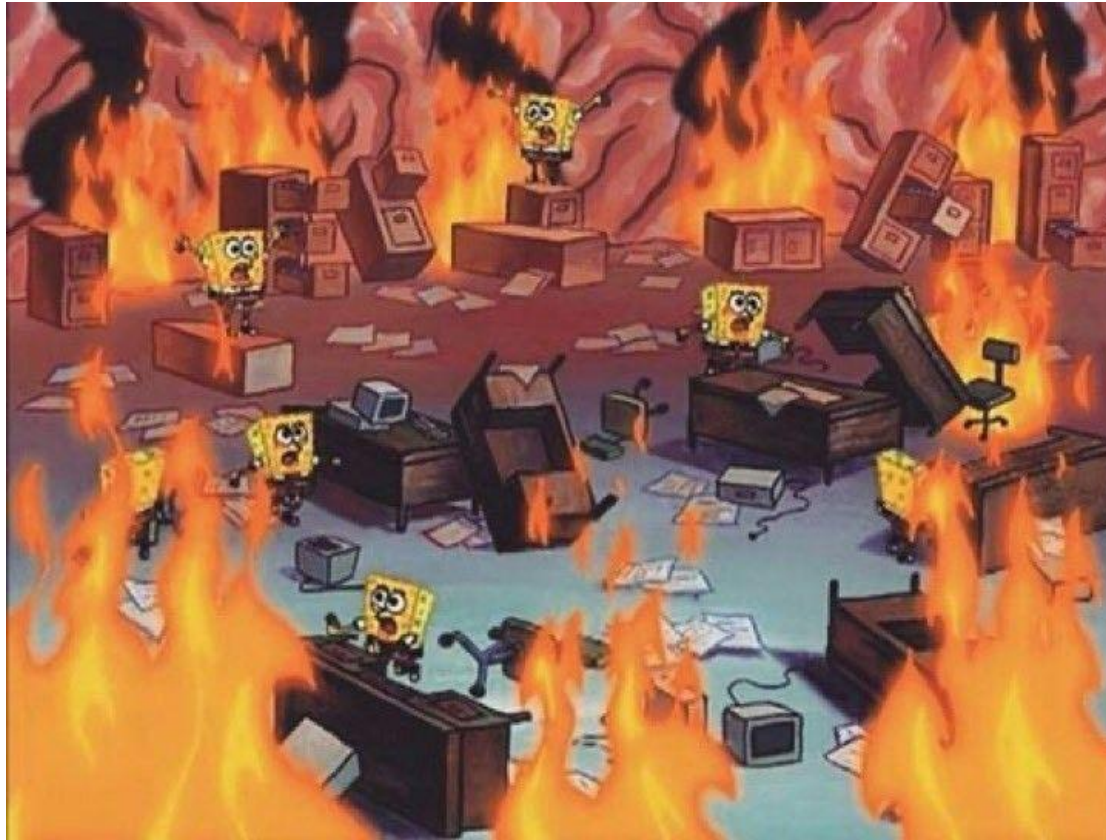
[1, 1, 1, 1, 1, 1, 1]

Order does not matter.

[1, 1, 5] and [5, 1, 1] are
considered different solutions

In our change making algorithm, we are calculating every possible permutation (bad)

Let's try 81 cents!



Change Making Problem

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D = Array of coin denominations [1, 5, 10, 25]

p = value to make change for

n = minimum number of coins used to make p cents

Goal: Find an **n -length** combination of coins from **D** that were used to make **p**

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Goal: Find an **n -length** combination of coins from **D** that were used to make **p**

To do this, we will compute **all n -length combinations**, but only return the combinations that add up to be **p** (not very efficient)

For $n=3$, these are the combinations to be generated:

- [1, 1, 1]
- [1, 1, 5]
- [1, 1, 10]
- [1, 5, 5]
- [1, 5, 10]
- [1, 10, 10]
- [5, 5, 5]
- [5, 5, 10]
- [5, 10, 10]
- [10, 10, 10]

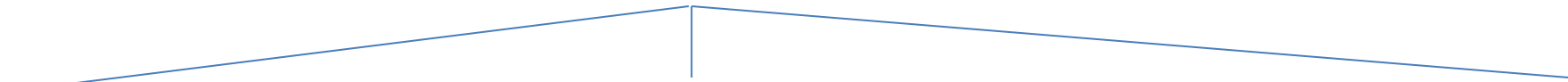
Note:

[5, 1, 1] is not a “valid” combination, because it is the same thing as [1, 1, 5]

Generating Combinations for finding coins

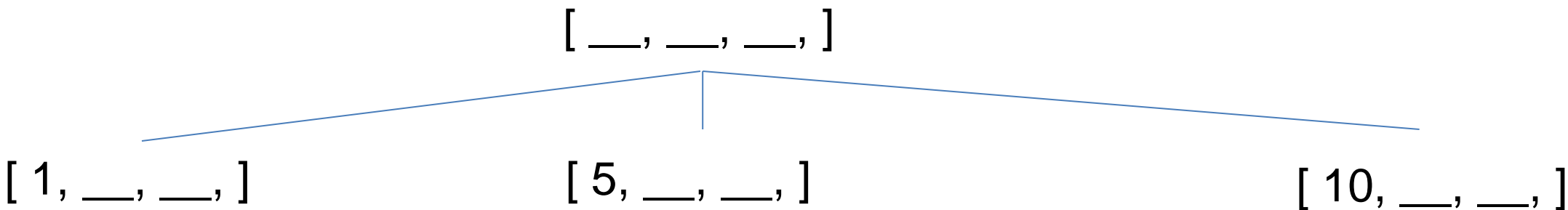
Denominations (D) = [1, 5, 10]
n = 3

[__, __, __,]



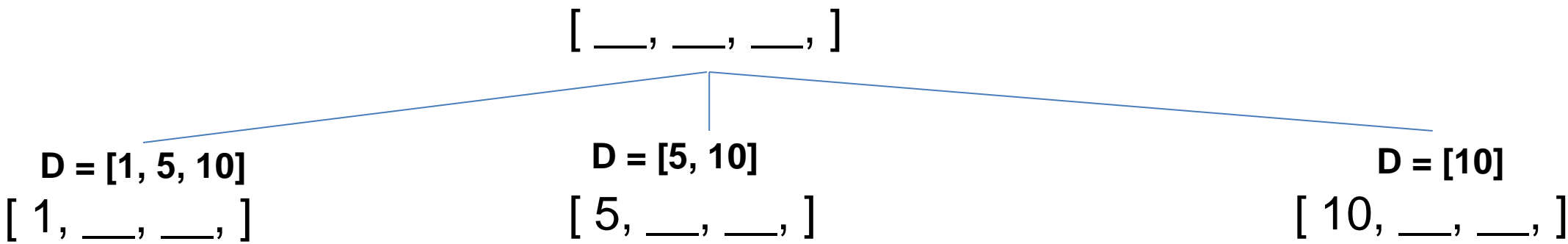
Generating Combinations for finding coins

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Generating Combinations for finding coins

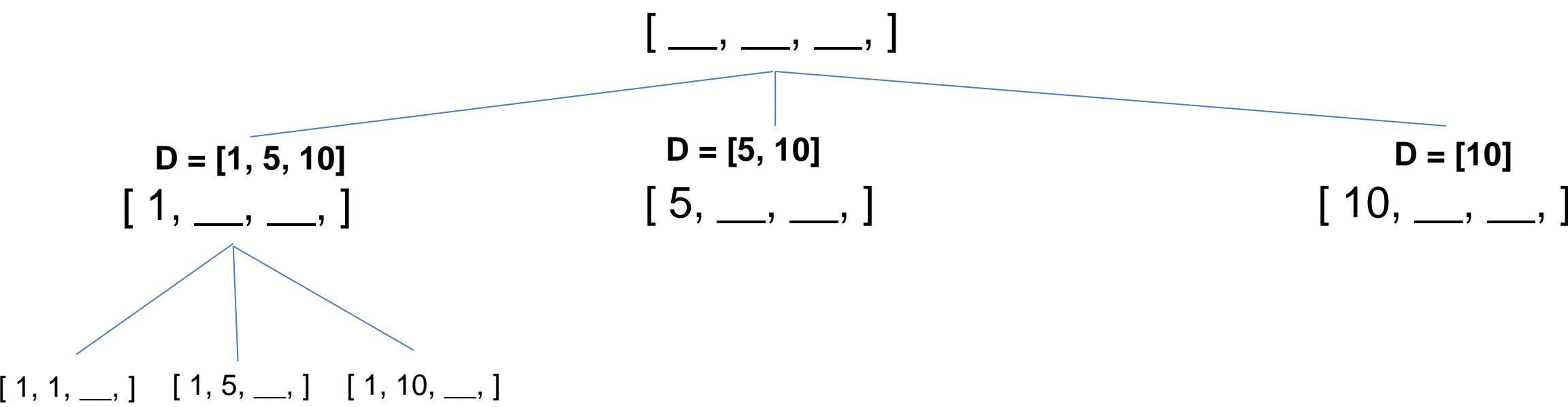
Denominations (D) = [1, 5, 10]
n = 3



↑
This branch does not get 1 included in the denomination set, because all the combinations involving 1 will be handled by the left tree

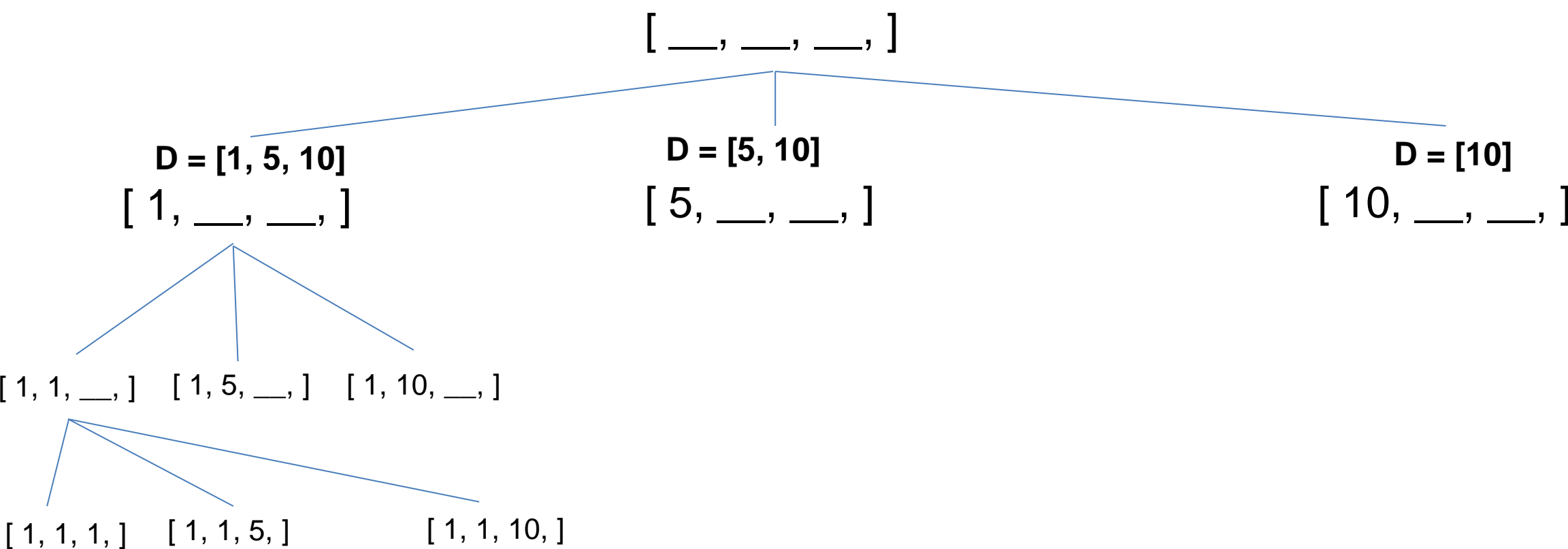
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Generating Combinations for finding coins

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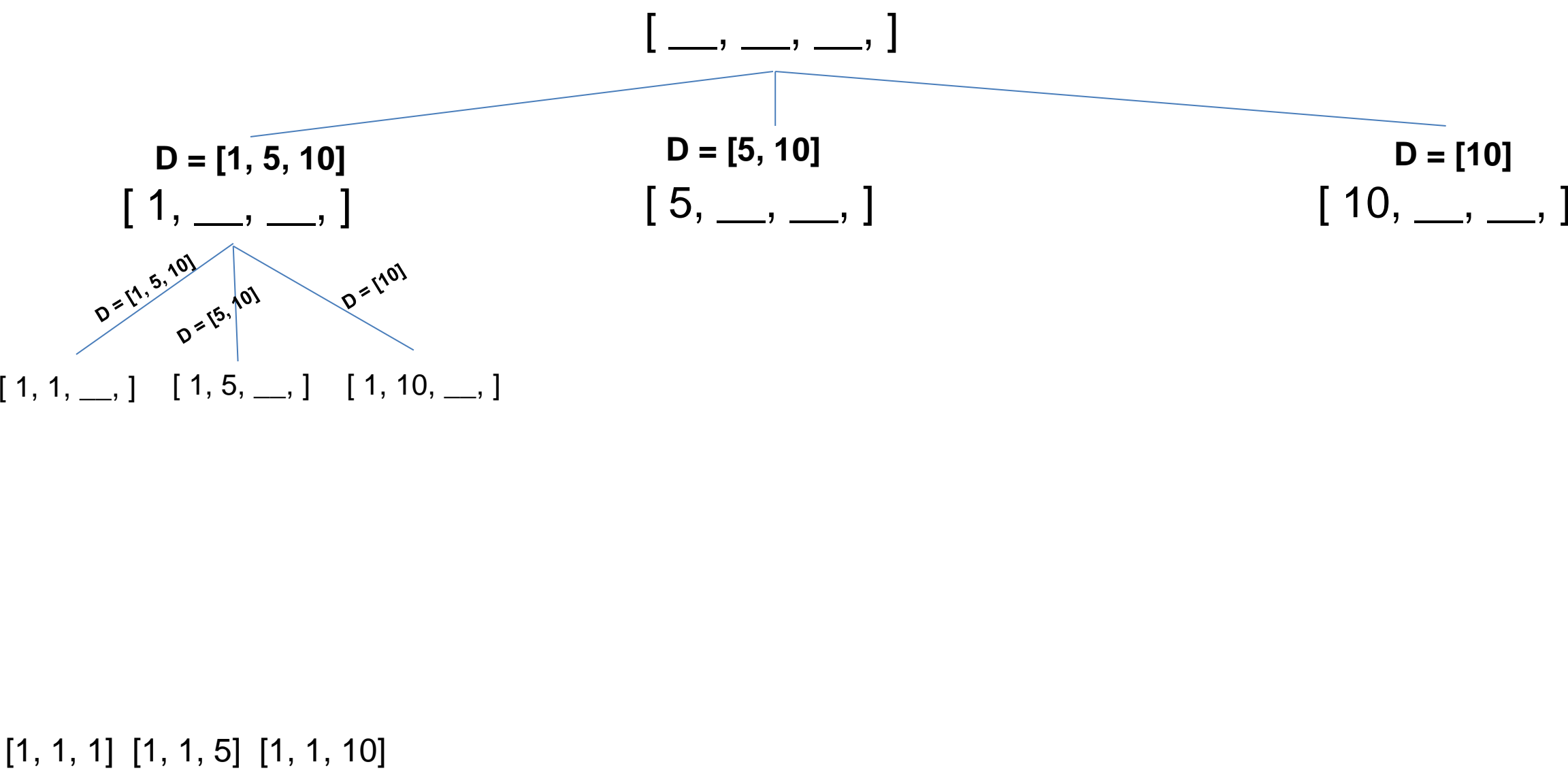


(Base case) When are combinations reach a length of 3, we will stop recursing

[1, 1, 1] [1, 1, 5] [1, 1, 10]

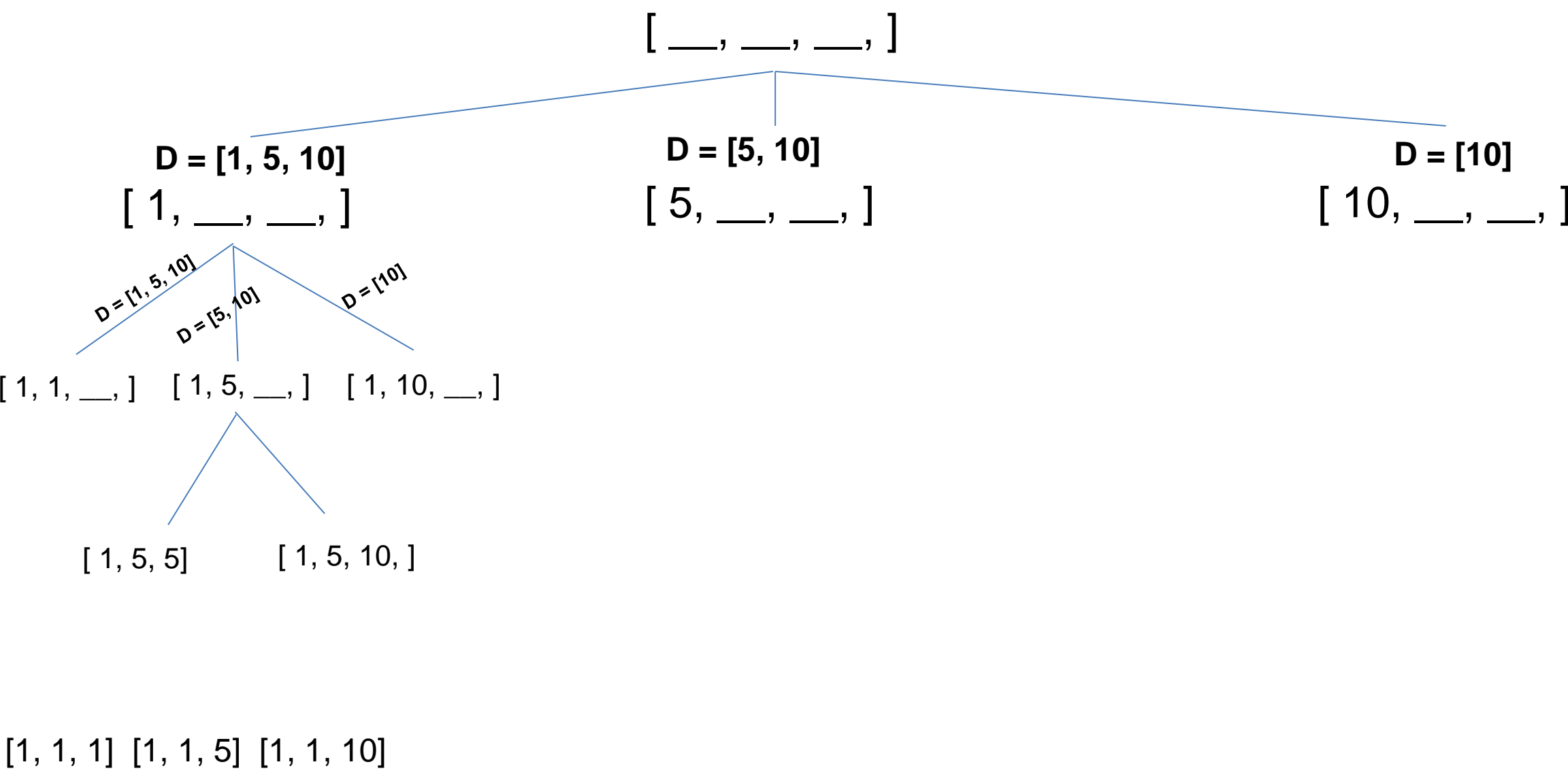
Generating Combinations for finding coins

Denominations (D) = [1, 5, 10]
n = 3



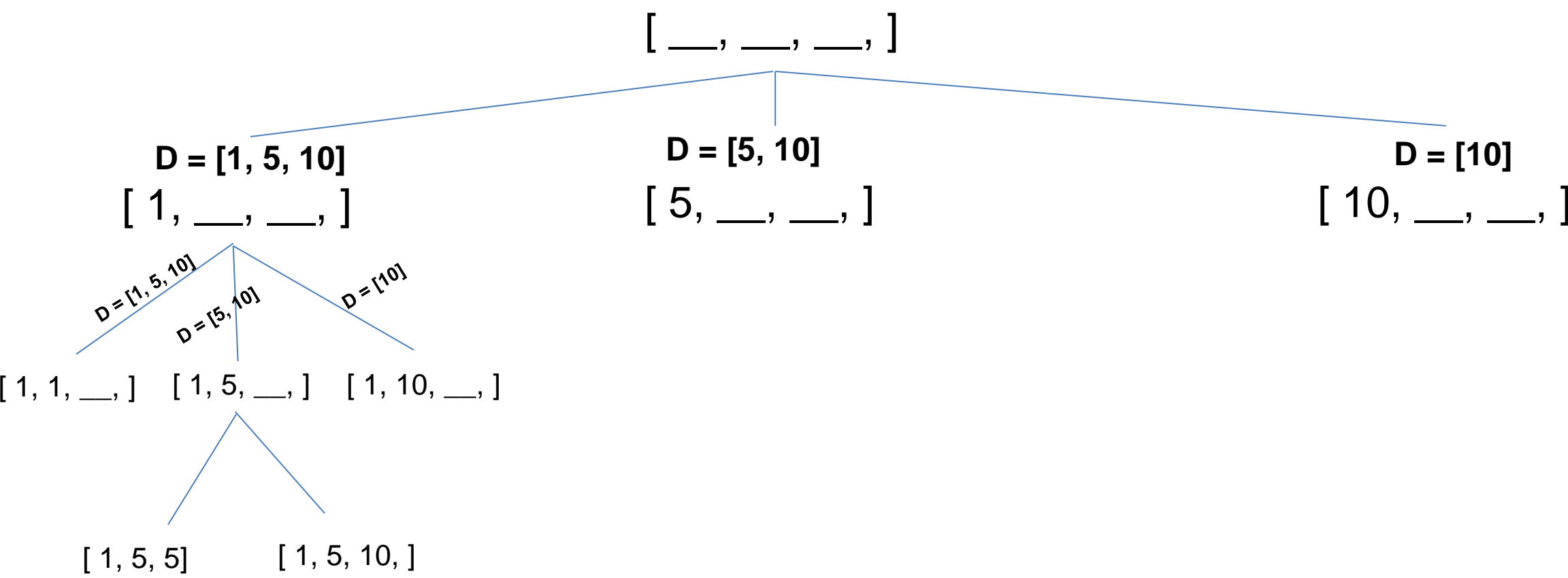
Generating Combinations for finding coins

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Generating Combinations for finding coins

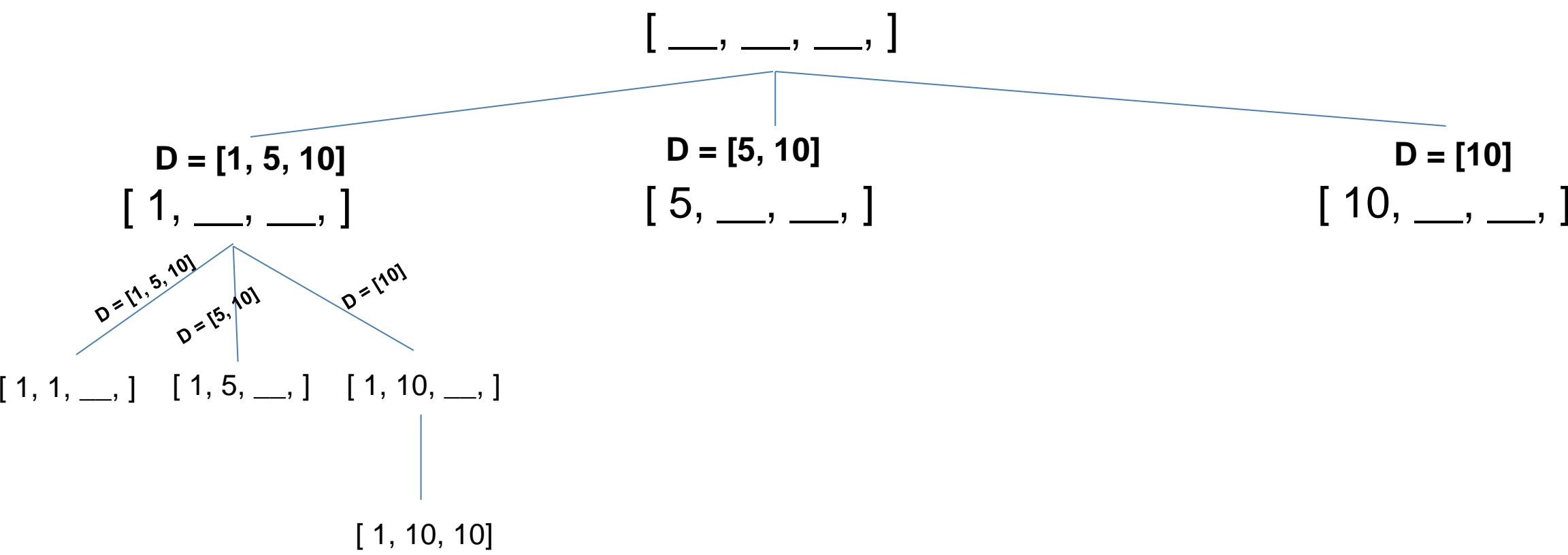
Denominations (D) = [1, 5, 10]
n = 3



[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10]

Generating Combinations for finding coins

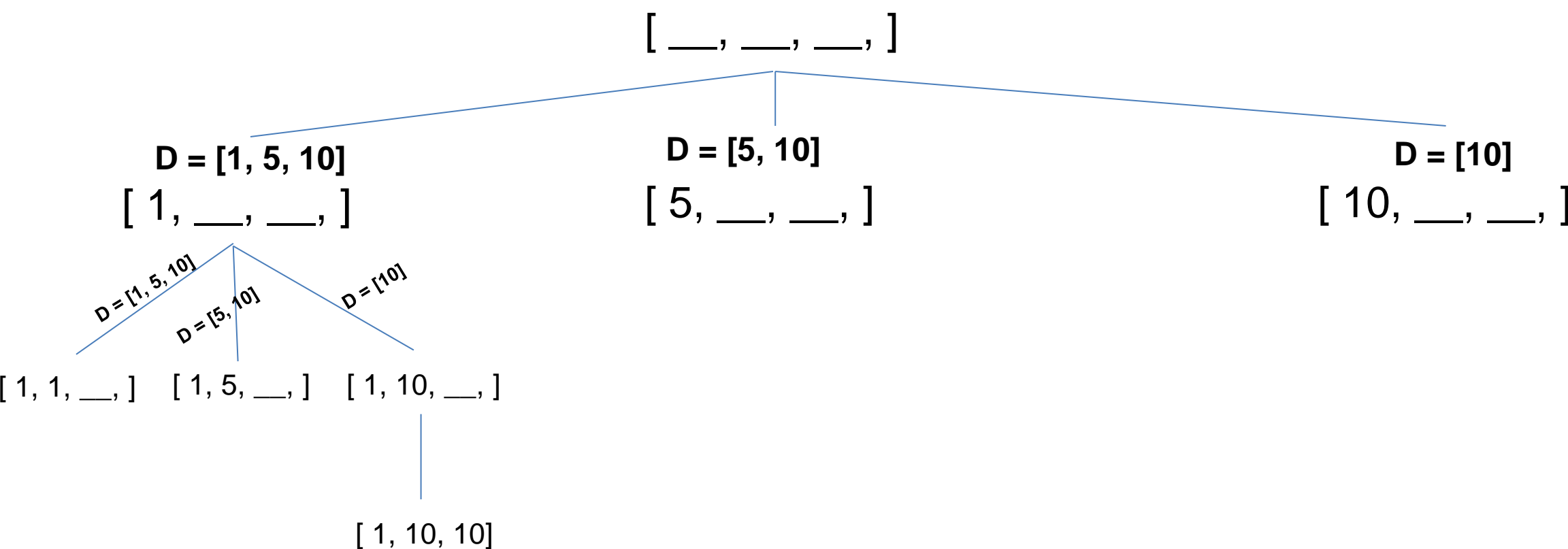
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[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10]

Generating Combinations for finding coins

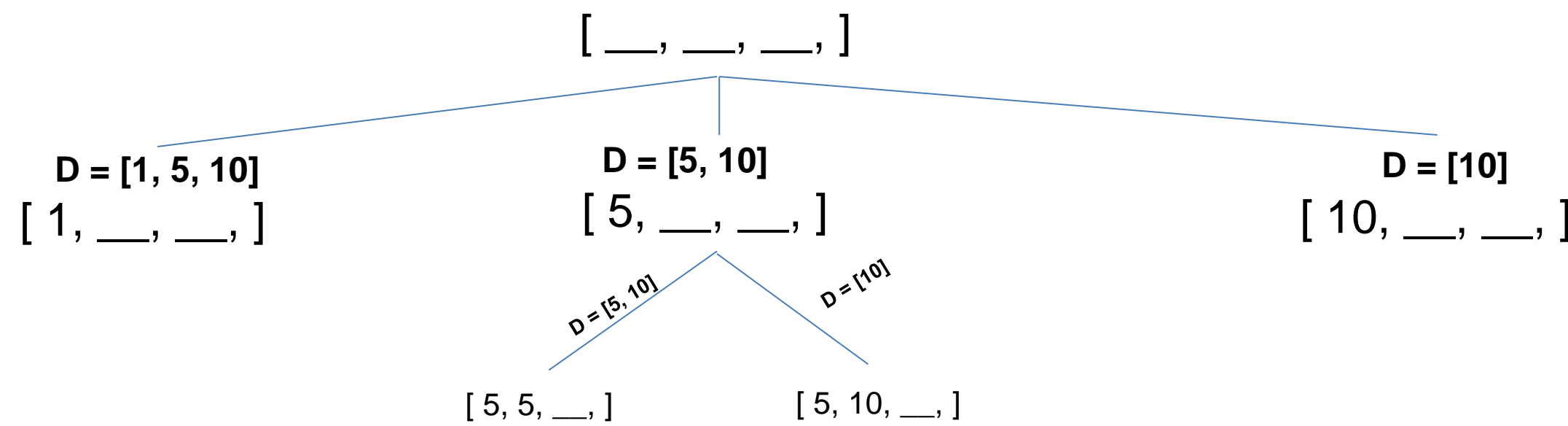
Denominations (D) = [1, 5, 10]
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[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]

Generating Combinations for finding coins

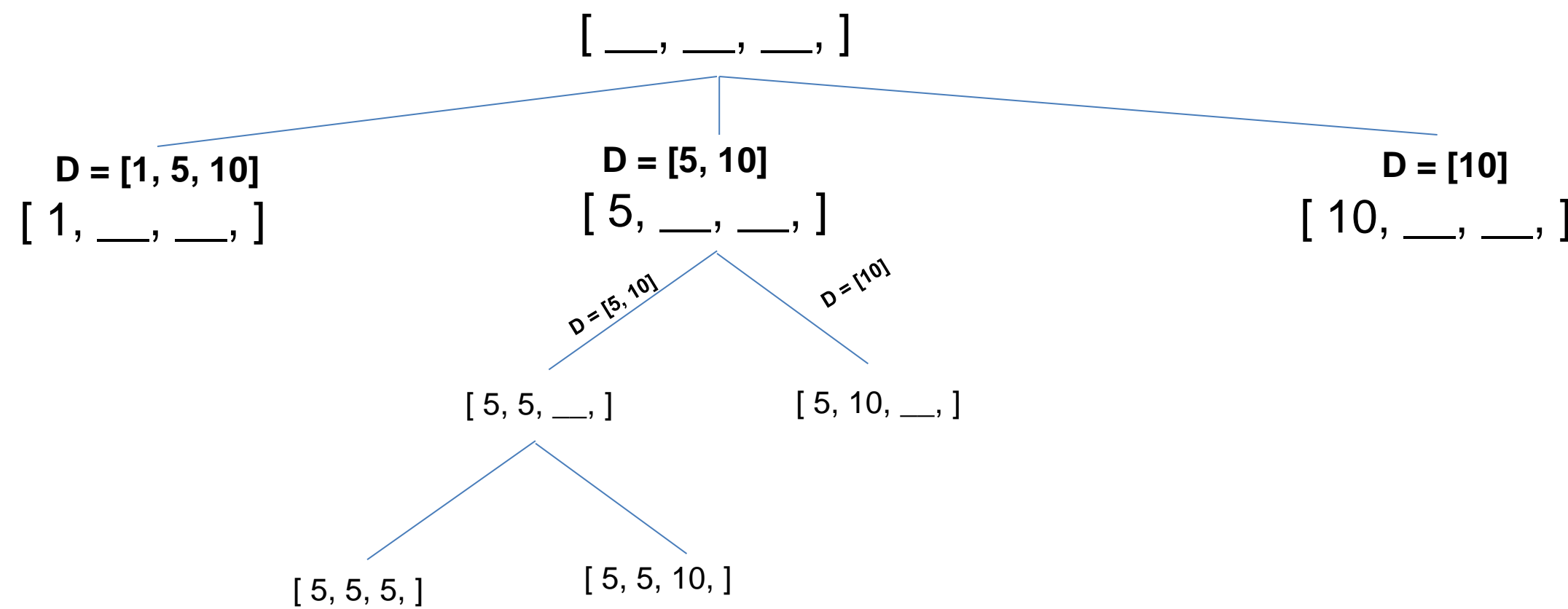
Denominations (D) = [1, 5, 10]
n = 3



[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]

Generating Combinations for finding coins

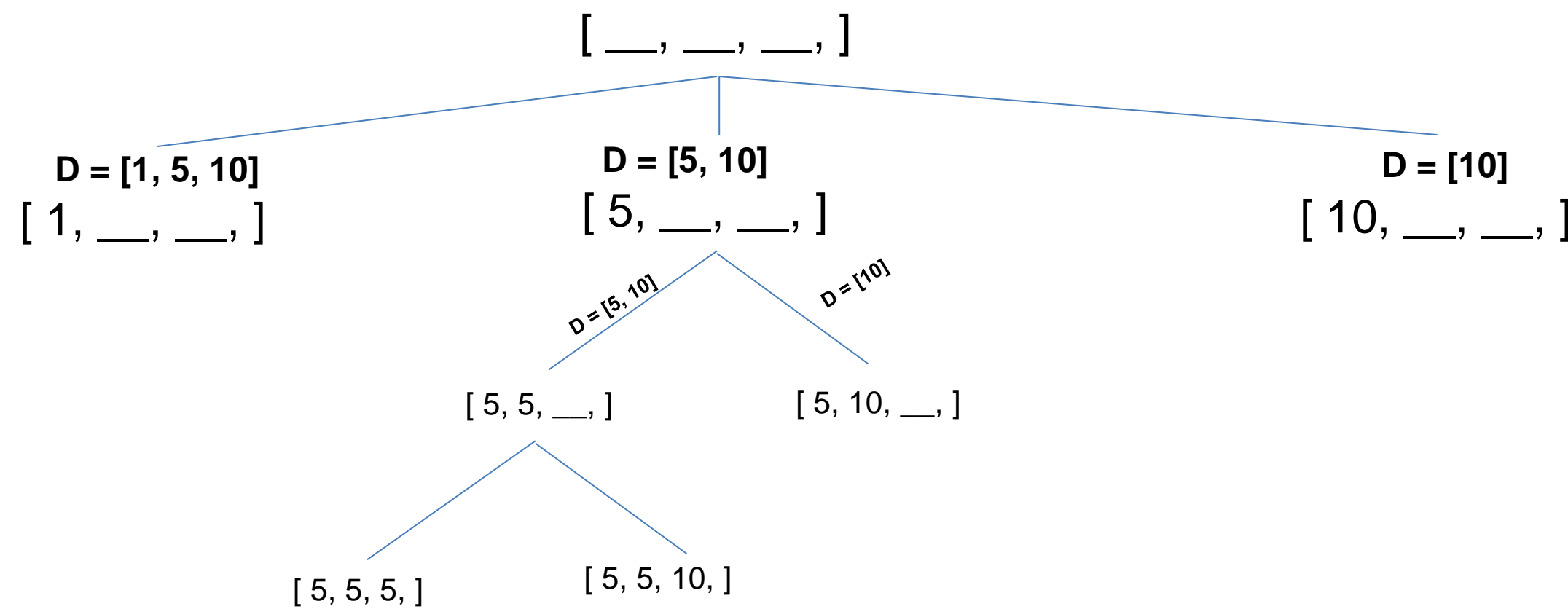
Denominations (D) = [1, 5, 10]
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[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10]

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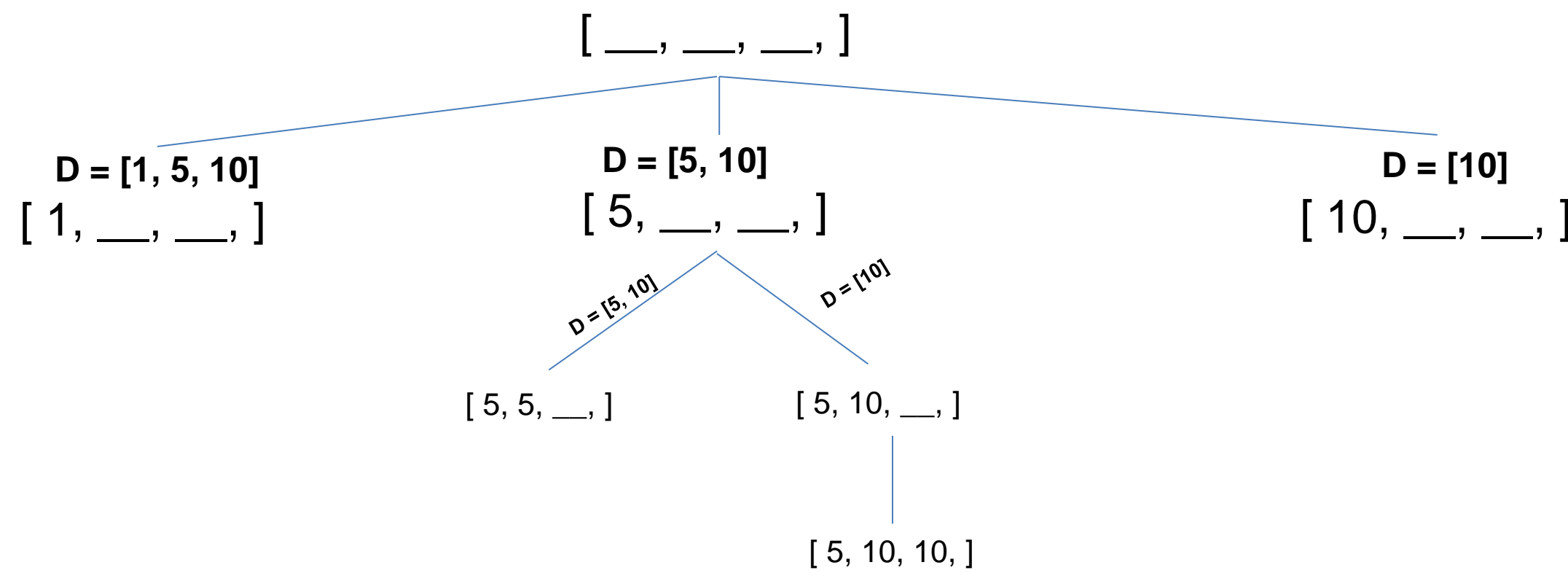
Denominations (D) = [1, 5, 10]
n = 3



[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10]

Generating Combinations for finding coins

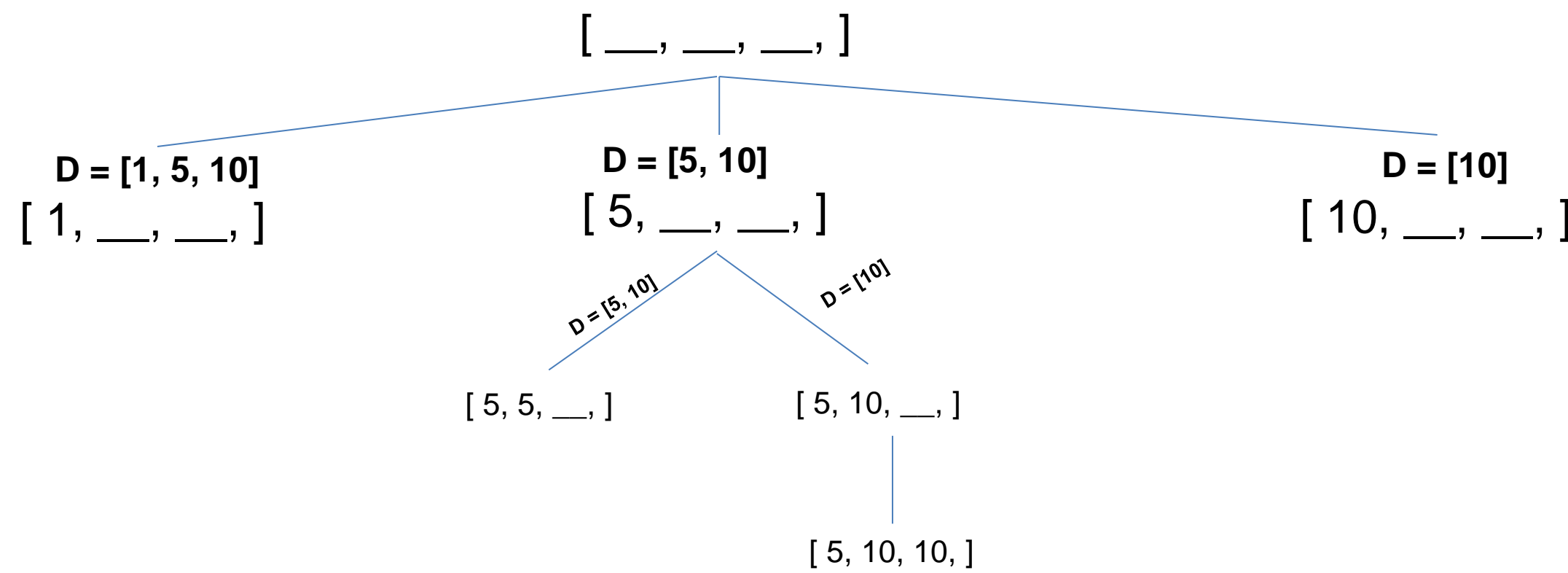
Denominations (D) = [1, 5, 10]
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[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10]

Generating Combinations for finding coins

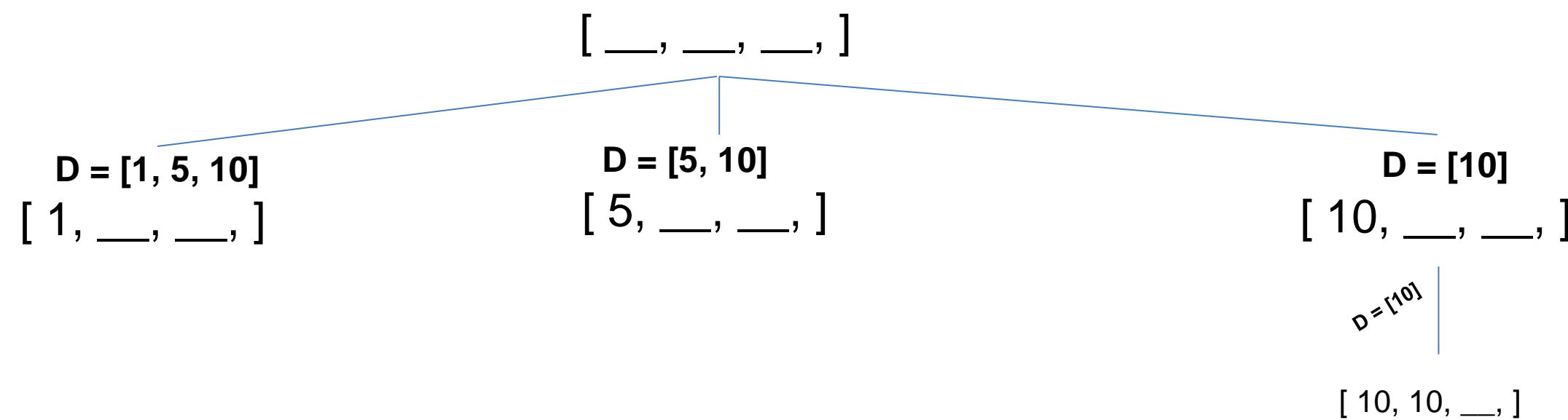
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[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10] [5, 10, 10]

Generating Combinations for finding coins

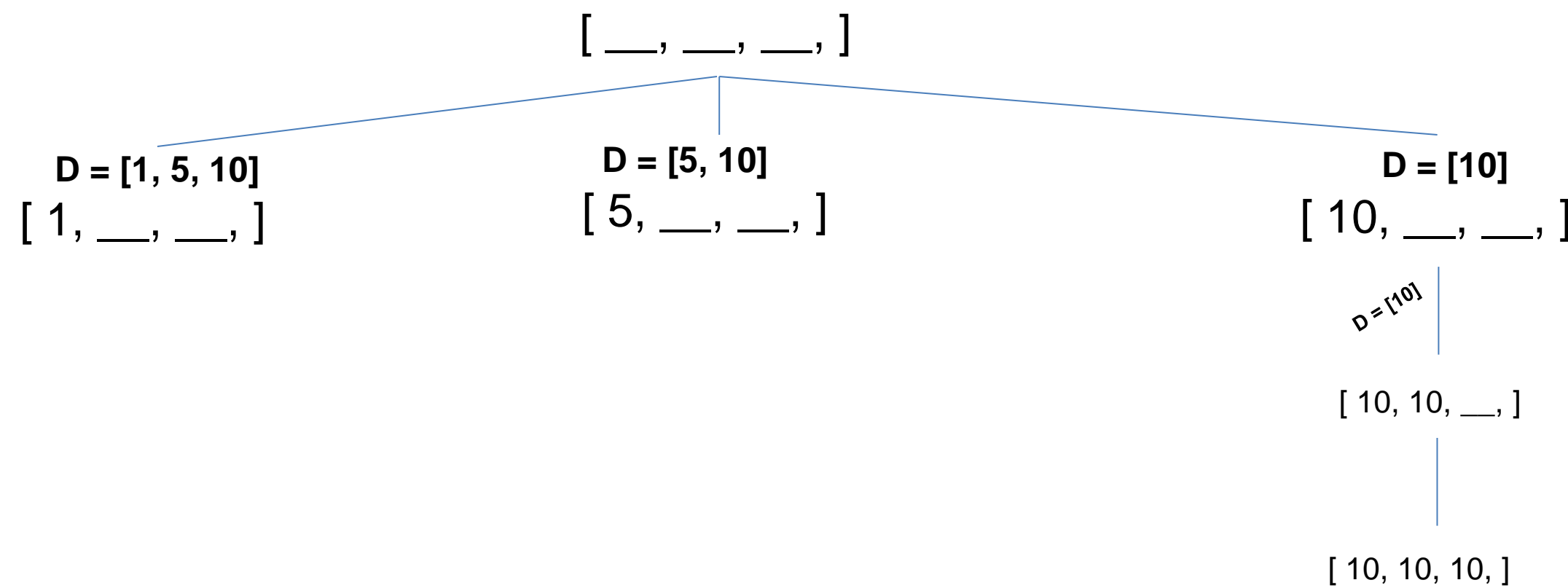
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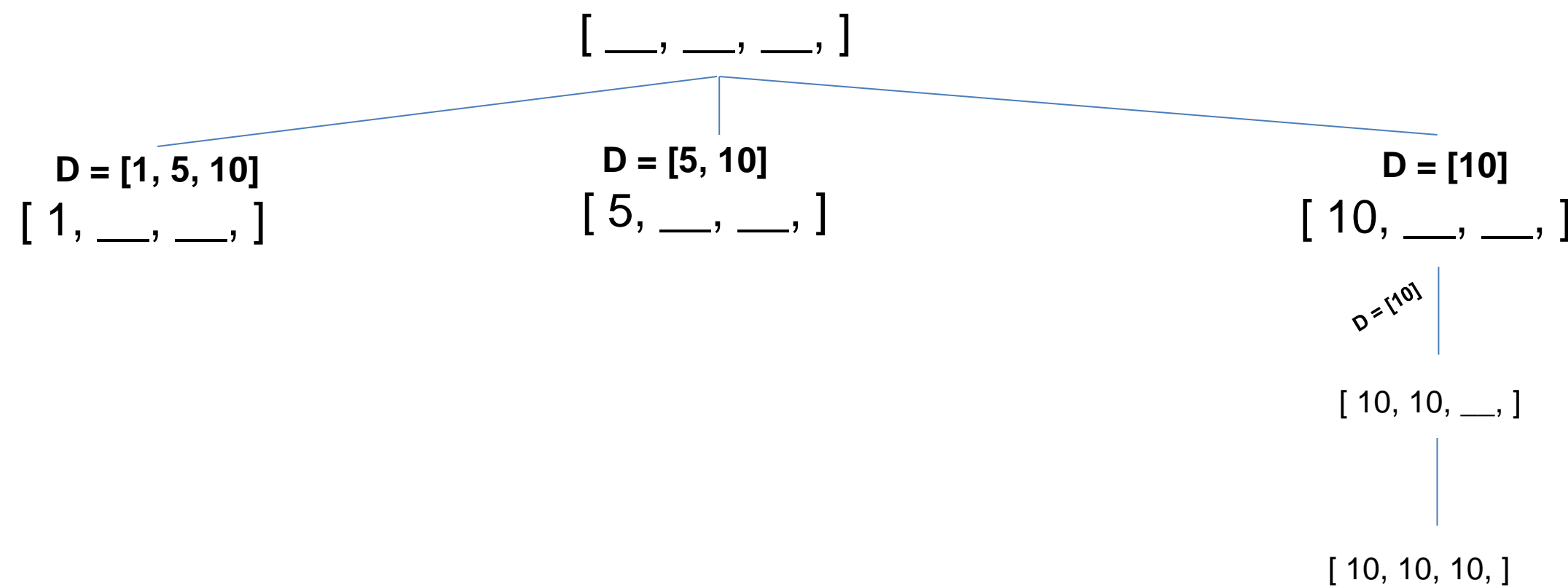
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Generating Combinations for finding coins

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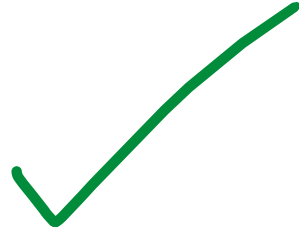


[1, 1, 1] [1, 1, 5] [1, 1, 10] [1, 5, 5] [1, 5, 10] [1, 10, 10] [5, 5, 5] [5, 5, 10] [5, 10, 10] [10, 10, 10]

Generating Combinations for finding coins

Denominations (D) = [1, 5, 10]
n = 3

1. [1, 1, 1]
2. [1, 1, 5]
3. [1, 1, 10]
4. [1, 5, 5]
5. [1, 5, 10]
6. [1, 10, 10]
7. [5, 5, 5]
8. [5, 5, 10]
9. [5, 10, 10]
10. [10, 10, 10]

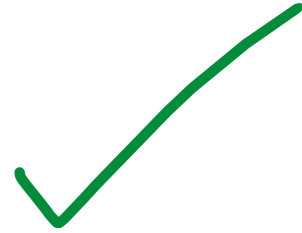


We've generated all combinations of length 3

Generating Combinations for finding coins

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We've generated all combinations of length 3

Now, we only want to print out the combinations that add up to **K**

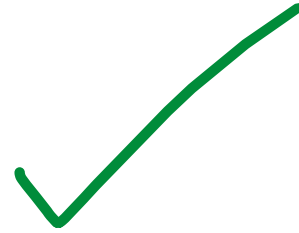
Generating Combinations for finding coins

Denominations (D) = [1, 5, 10]

n = 3

K = 16

1. [1, 1, 1]
2. [1, 1, 5]
3. [1, 1, 10]
4. [1, 5, 5]
5. [1, 5, 10]
6. [1, 10, 10]
7. [5, 5, 5]
8. [5, 5, 10]
9. [5, 10, 10]
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Suppose $K = 16$ (a minimum of 3 coins is needed to make 16 cents)

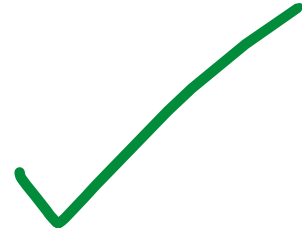
Generating Combinations for finding coins

Denominations (D) = [1, 5, 10]

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K = 16

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Answer = [1, 5, 10]

LET'S CODE THIS!!

If you don't fully understand this code, that is fine.

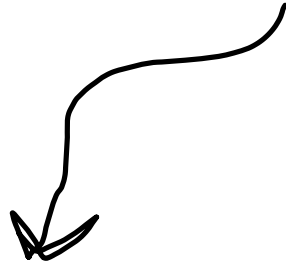
o to **K**

16 cents)

10.[10, 10, 10]

Answer = [1, 5, 10]

```
private static void find_coins(int[] d, int k, int n) {  
    int chosen[] = new int[n + 1];  
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);  
}
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
```

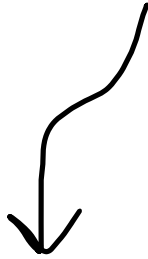
Array that we build up over time. Holds **indices** of currently selected denominations for some combination

[1, __, __]

[1, 1, __]

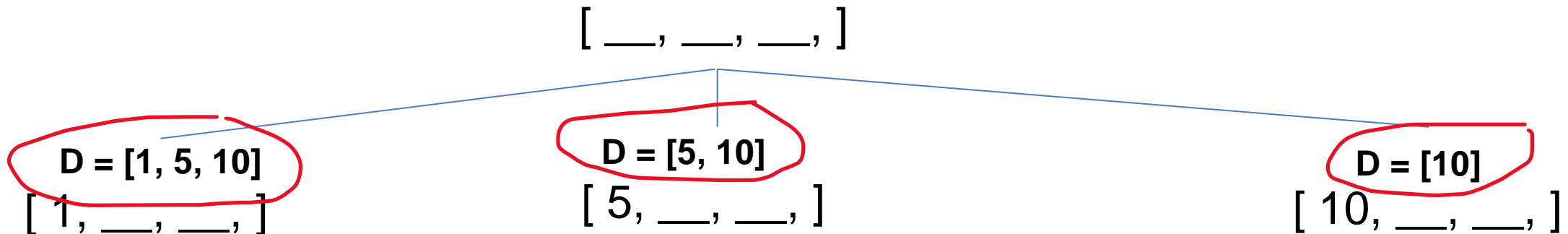
[1, 1, 5]


```
private static void find_coins(int[] d, int k, int n) {
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}
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
```

Array of denominations we pass for each recursive call



```
private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
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}
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
```

The next index that we need to insert at for `chosen` array

[__, __, __]



index

[1, __, __]



index

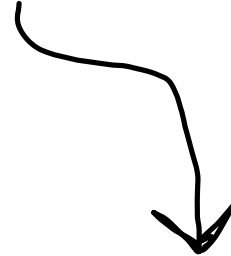
[1, 1, __]



index

[1, 1, 5]

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private static void find_coins(int[] d, int k, int n) {
    int chosen[] = new int[n + 1];
    calculate_combinations(chosen, d, 0, n, 0, d.length - 1, k);
}
```



```
void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
```

The desired size of the combination. When `index == r`, we have reached the desired combination size

[__, __, __]



index = 0

[1, __, __]



index = 1

[1, 1, __]



index = 2

[1, 1, 5]



index = 3

**We've hit our
desired size!**

```

void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
    if (index == r) {
        int counter = 0;
        ArrayList<Integer> coins = new ArrayList<Integer>();
        for (int i = 0; i < r; i++) {
            counter += arr[chosen[i]];
            coins.add(arr[chosen[i]]);
        }
        if(counter == target) {
            System.out.println(coins);
        }
        return;
    }
    for (int i = start; i <= end; i++) {
        chosen[index] = i;
        calculate_combinations(chosen, arr, index + 1, r, i, end, target);
    }
    return;
}

```

```

void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
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    for (int i = start; i <= end; i++) {
        chosen[index] = i;
        calculate_combinations(chosen, arr, index + 1, r, i, end, target);
    }
    return;
}

```

Base case

If we hit our base case, we know we have N things, so put them in an ArrayList and add them up

Only print out the combination if it adds up to target

```

void calculate_combinations(int[] chosen, int[] arr, int index, int r, int start, int end, int target) {
    if (index == r) {
        int counter = 0;
        ArrayList<Integer> coins = new ArrayList<Integer>();
        for (int i = 0; i < r; i++) {
            counter += arr[chosen[i]];
            coins.add(arr[chosen[i]]);
        }
        if(counter == target) {
            System.out.println(coins);
        }
        return;
    }
    for (int i = start; i <= end; i++) {
        chosen[index] = i;
        calculate_combinations(chosen, arr, index + 1, r, i, end, target);
    }
    return;
}

```

Recursive Case

Otherwise, insert selected coin into the `chosen` array

create $(end - start)$ branches, and give it a smaller section of D

