

CSCI 232:

Data Structures and Algorithms

Heaps

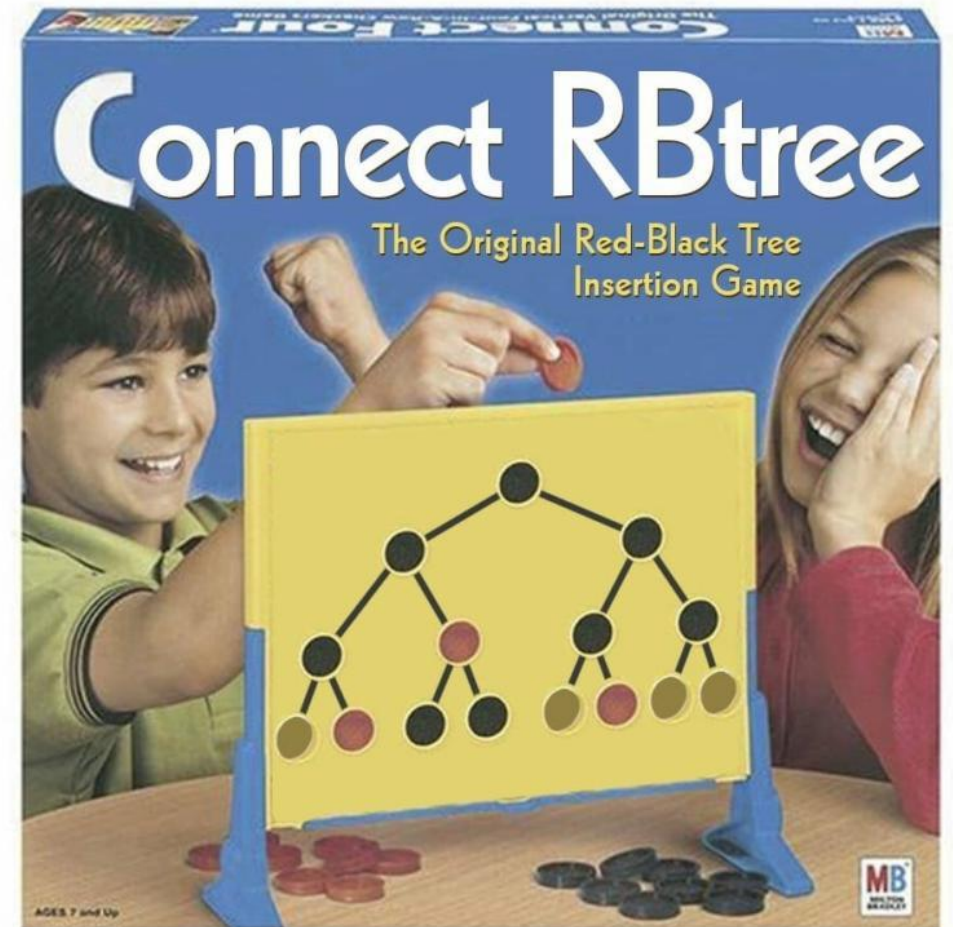
Reese Pearsall
Spring 2024

Announcements

Lab 7 due **tomorrow** at 11:59 PM

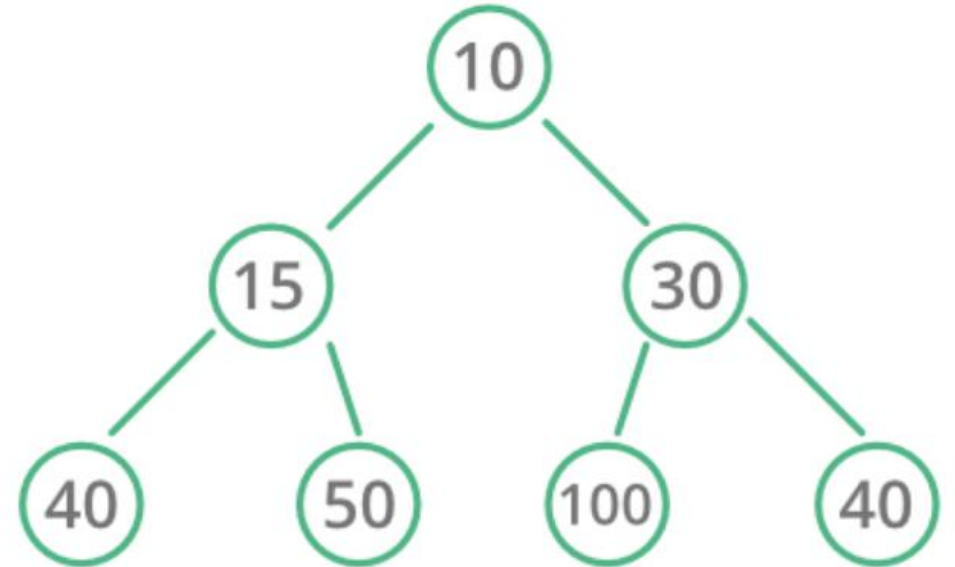
Program 2 due **Sunday** at 11:59 PM

No class next week



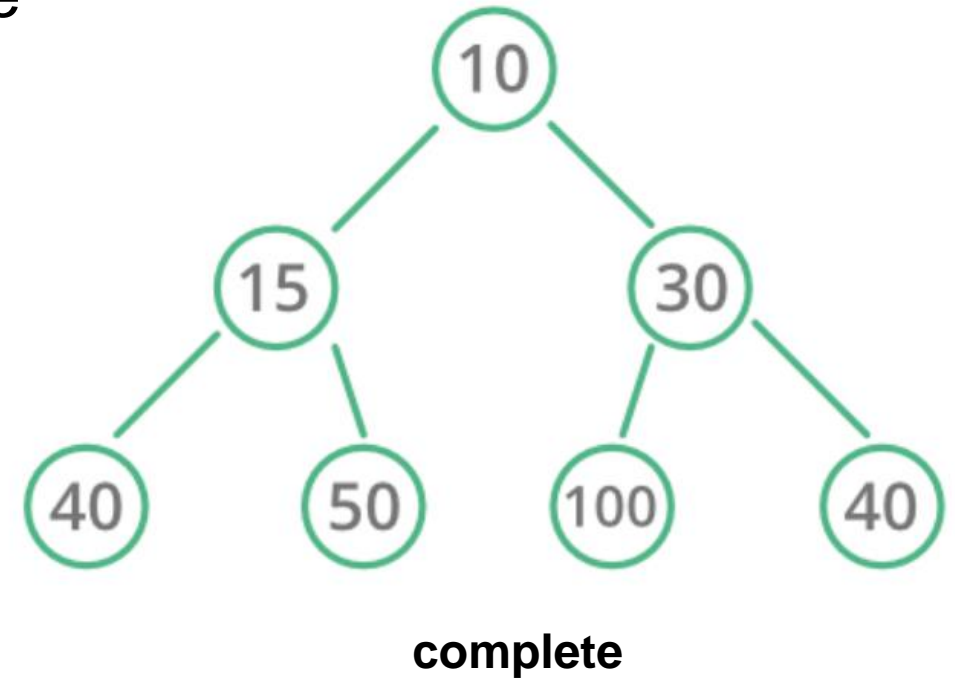
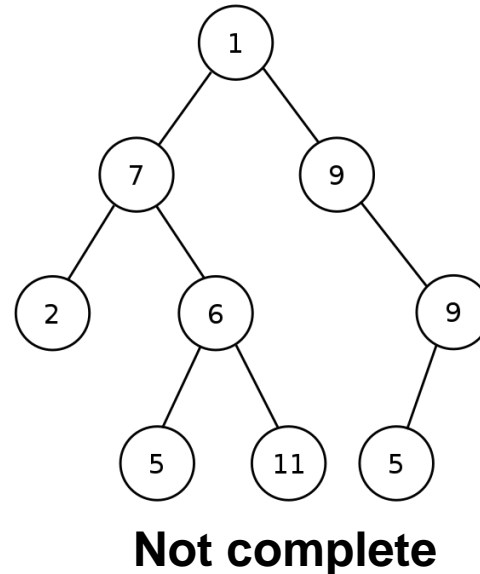
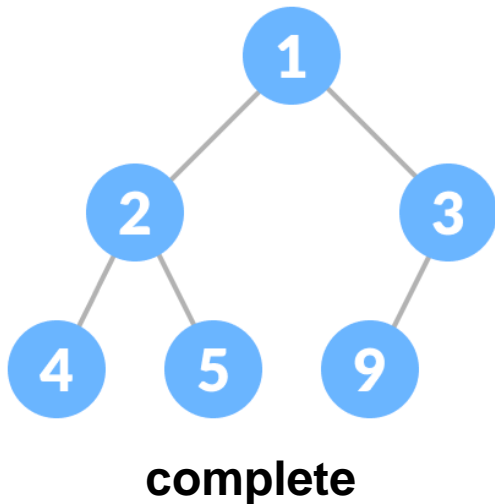
the game that you play in nightmares ^

The **Heap** data structure is complete binary tree that follows the heap property



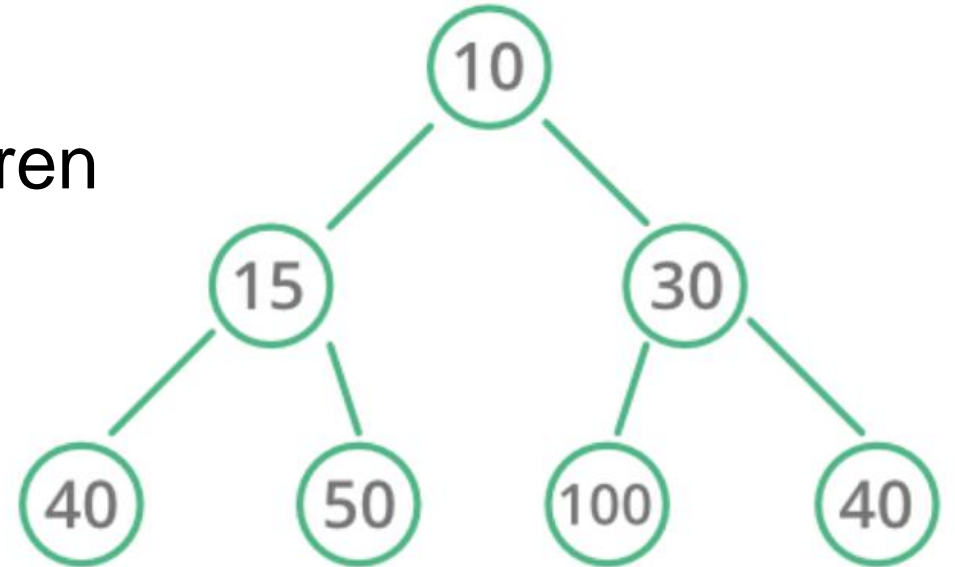
The **Heap** data structure is **complete** binary tree that follows the heap property

Complete tree - Every level, except possibly the last, is completely filled, and all nodes in the last level are as far left as possible



The **Heap** data structure is complete **binary** tree that follows the heap property

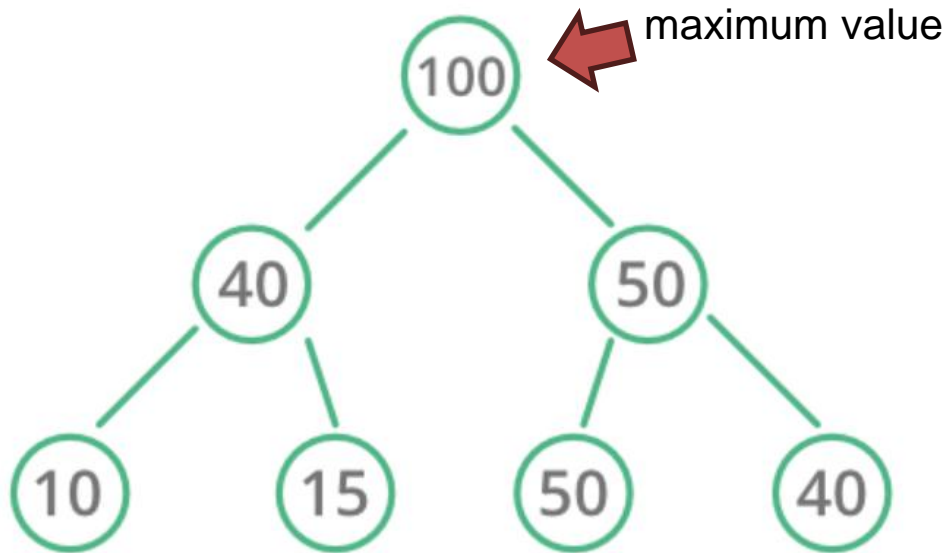
Binary – cannot have more than two children



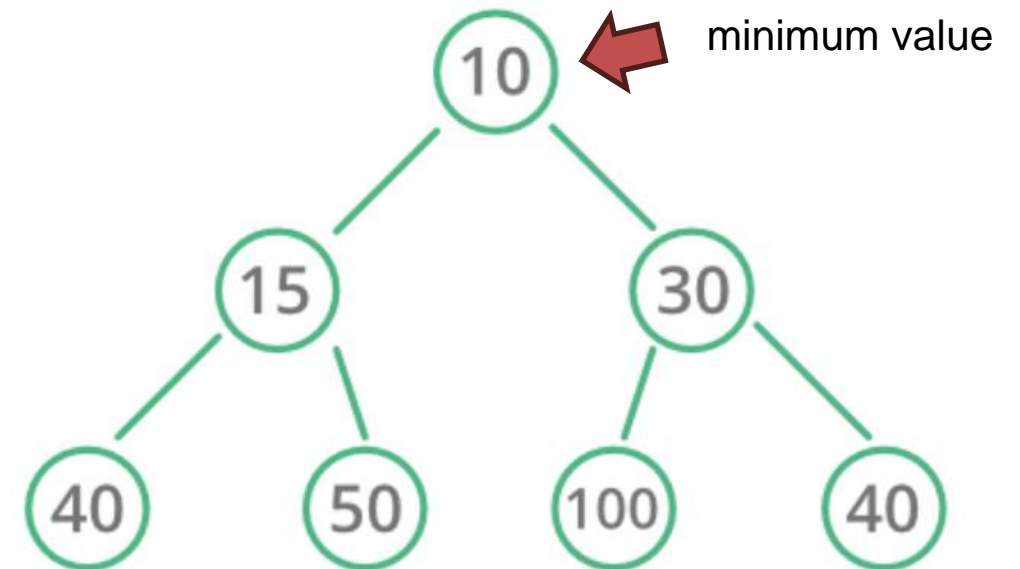
The **Heap** data structure is complete binary tree that follows the **heap property**

Two types of heaps

Max Heap – Parent nodes are greater than both of its children

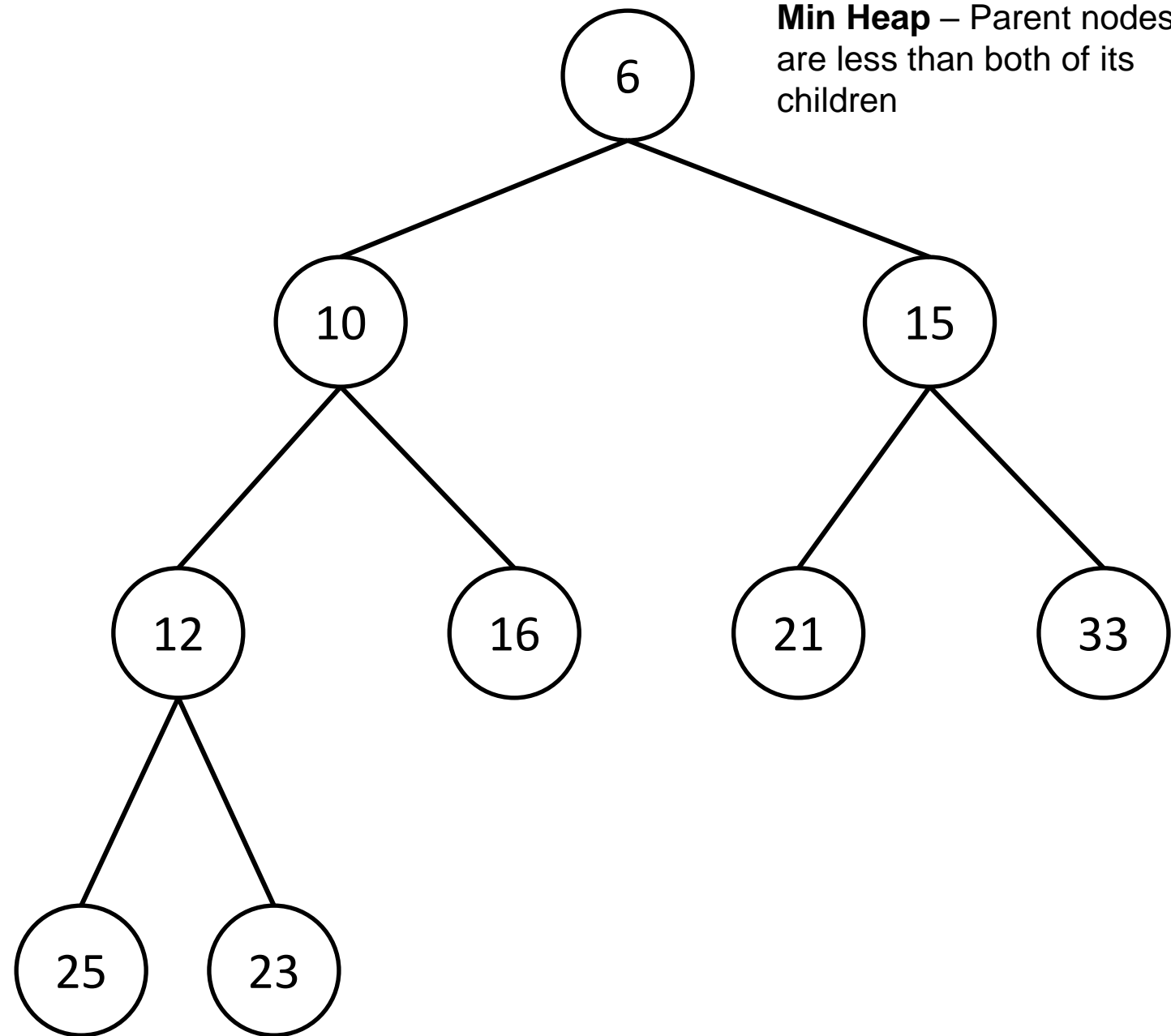


Min Heap – Parent nodes are less than both of its children



Heap Operations - Insert

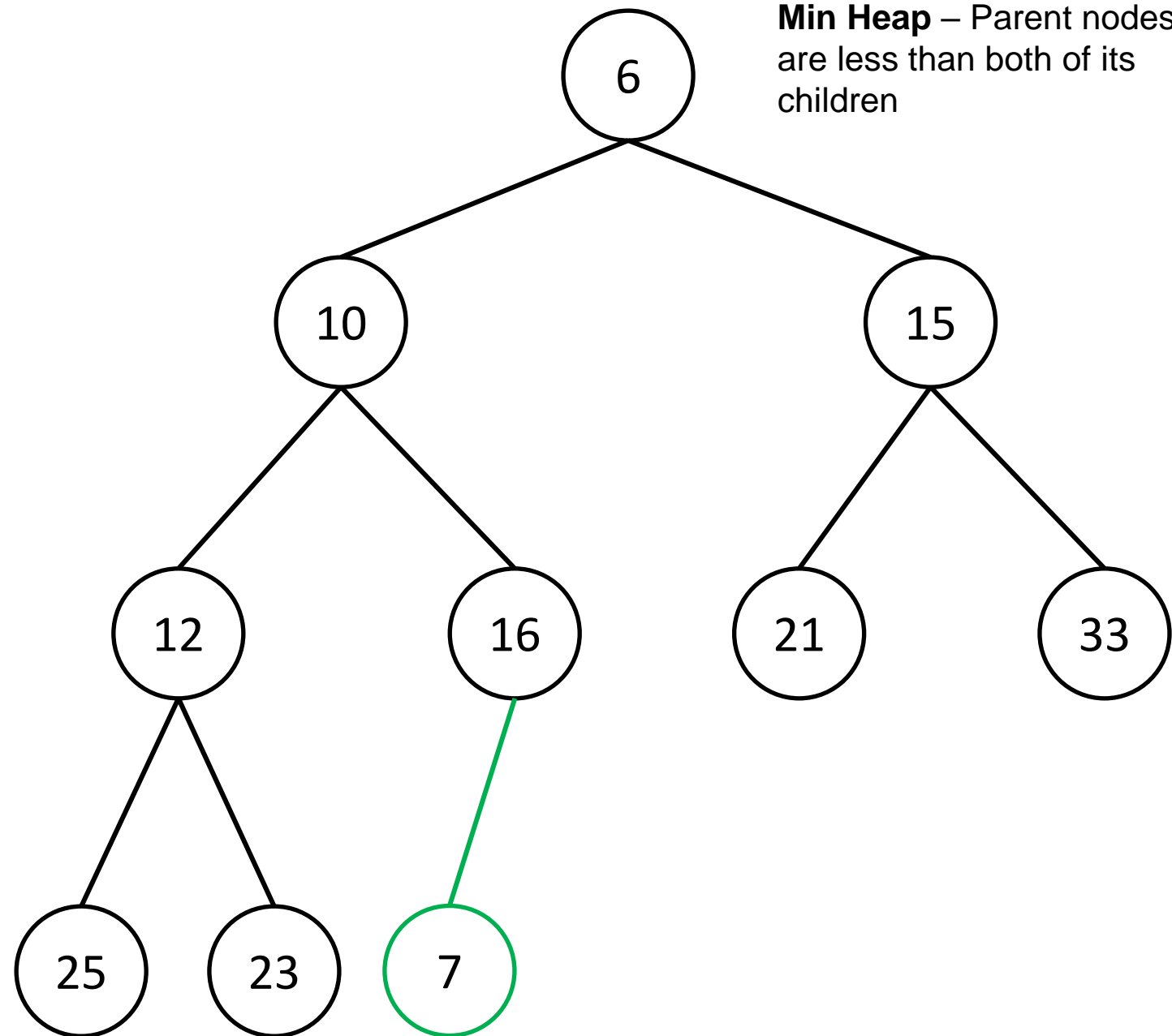
`add(7);`



Heap Operations - Insert

`add(7);`

Because this is a complete binary tree, this is the only place a new node can go

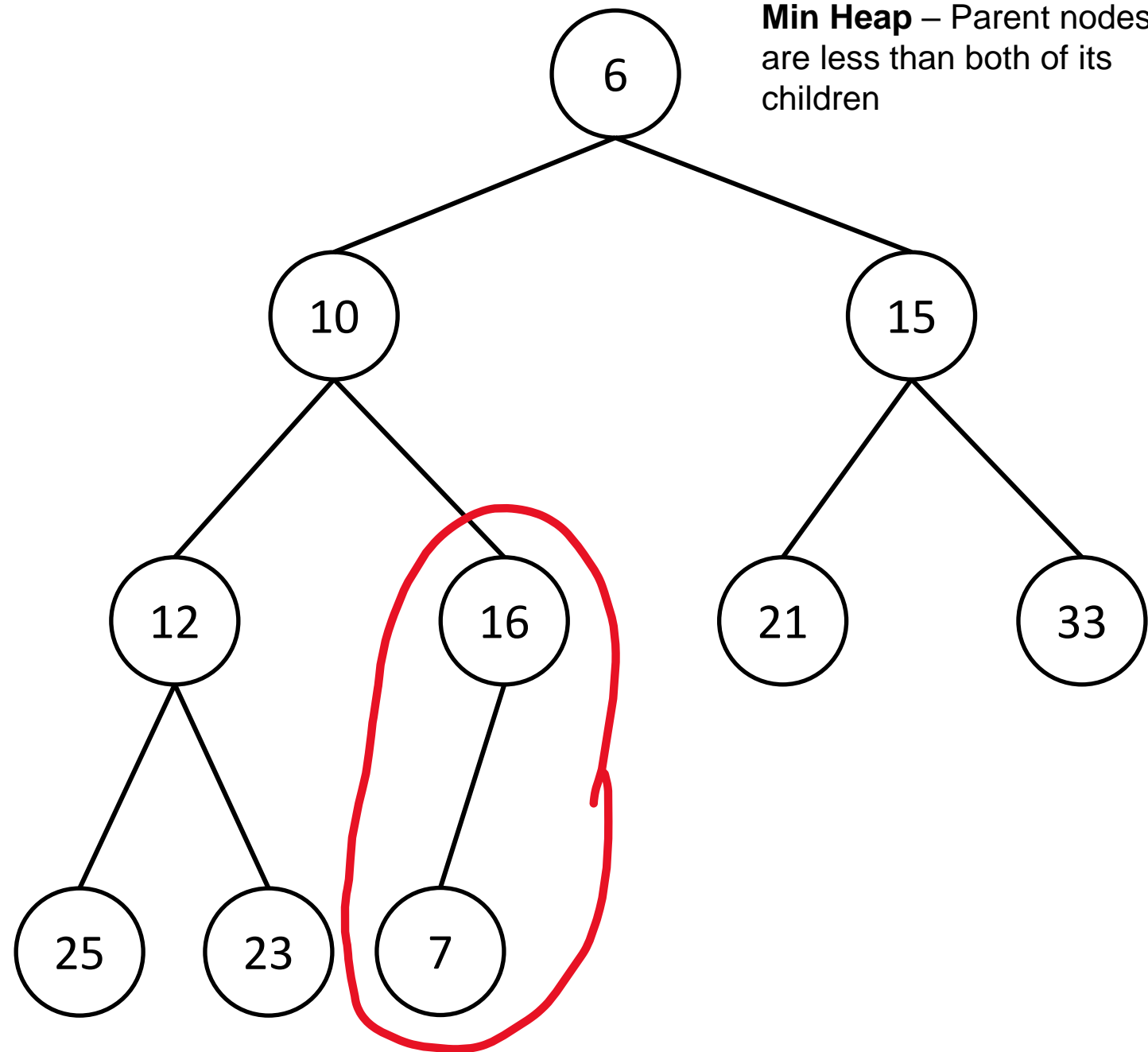


Heap Operations - Insert

`add(7);`

Because this is a complete binary tree, this is the only place a new node can go

However, we are now violating the heap property



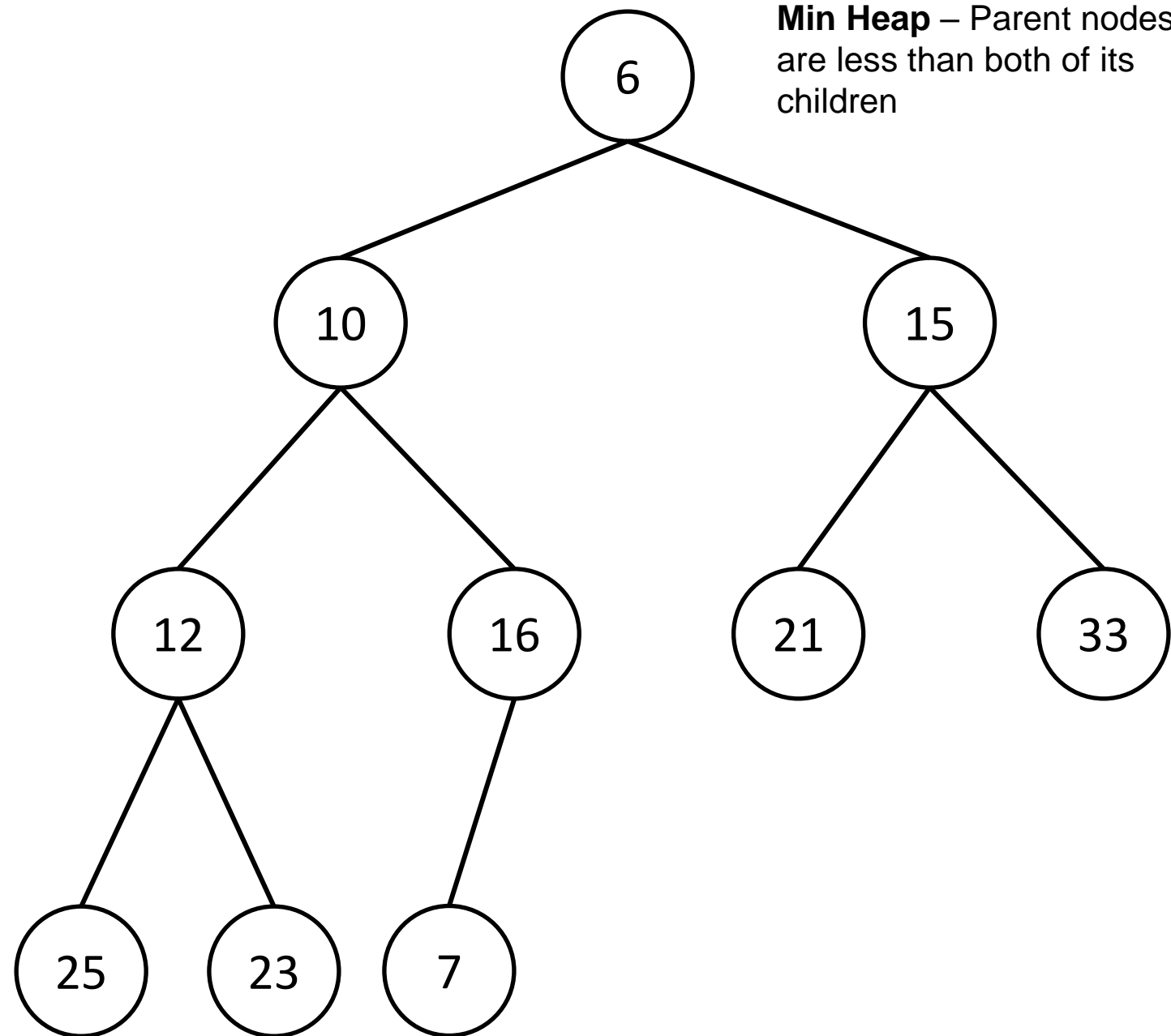
Heap Operations - Insert

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Because this is a complete binary tree, this is the only place a new node can go

However, we are now violating the heap property

When new nodes are added, we may need to move it up in the tree



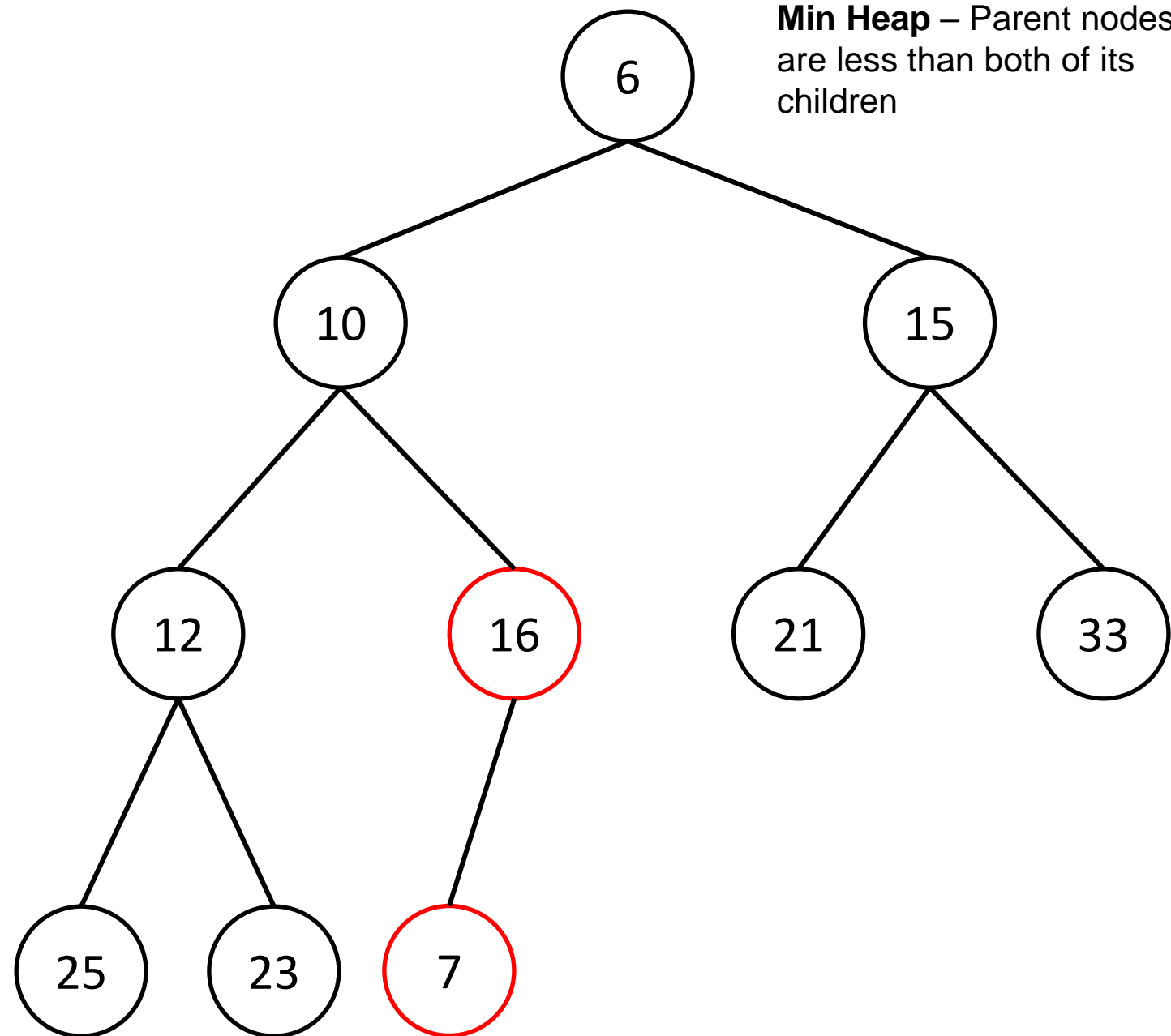
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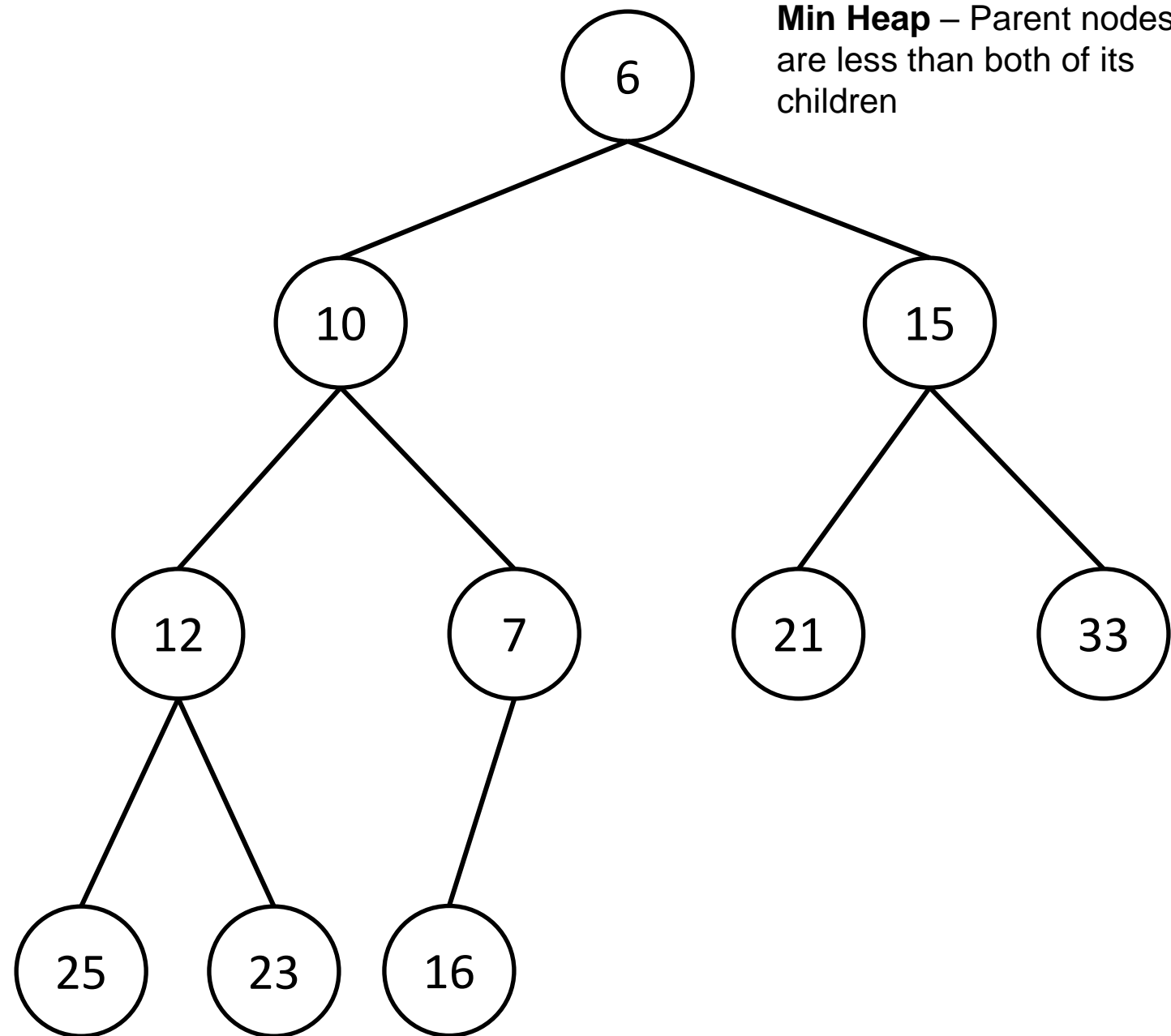
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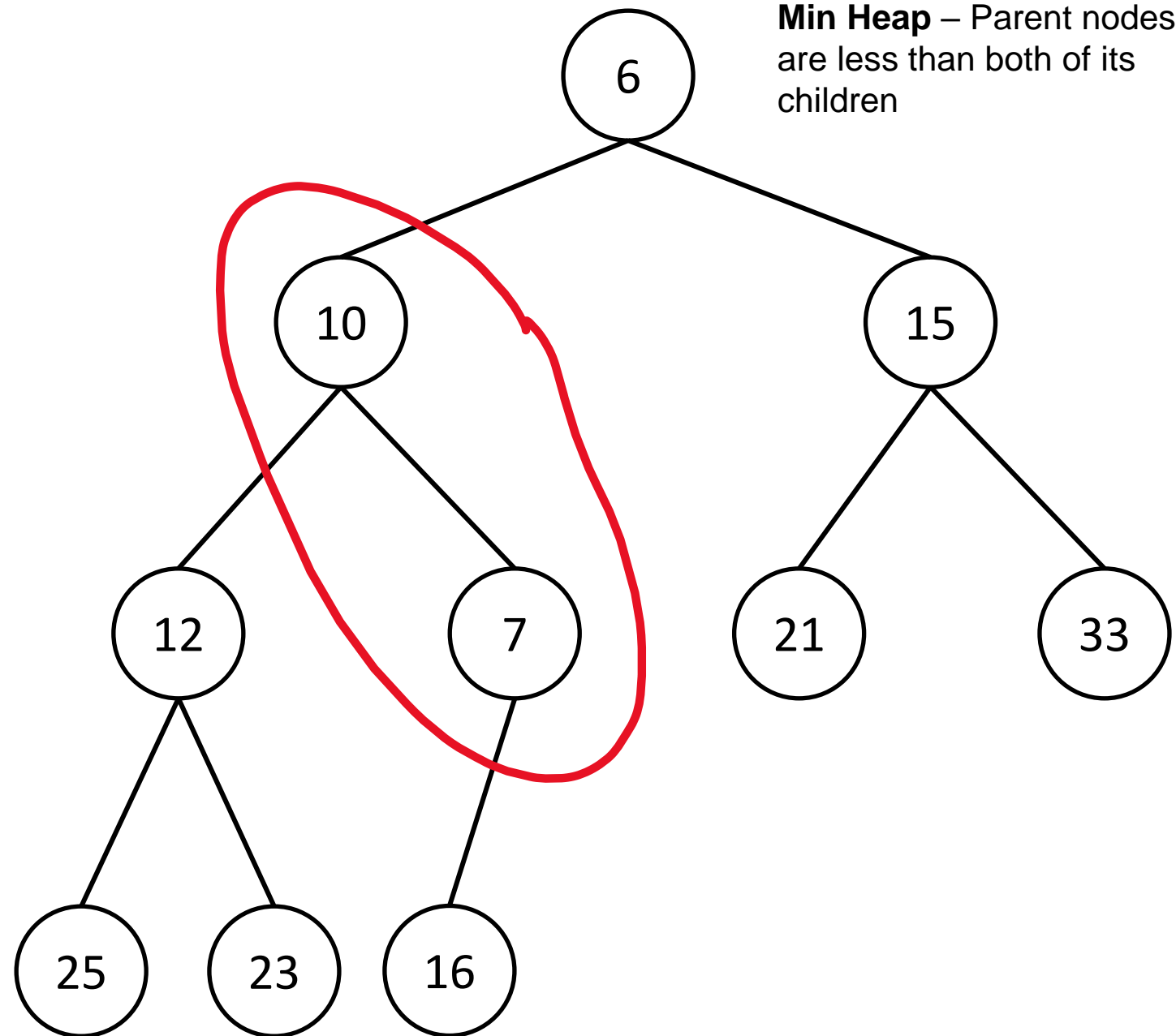
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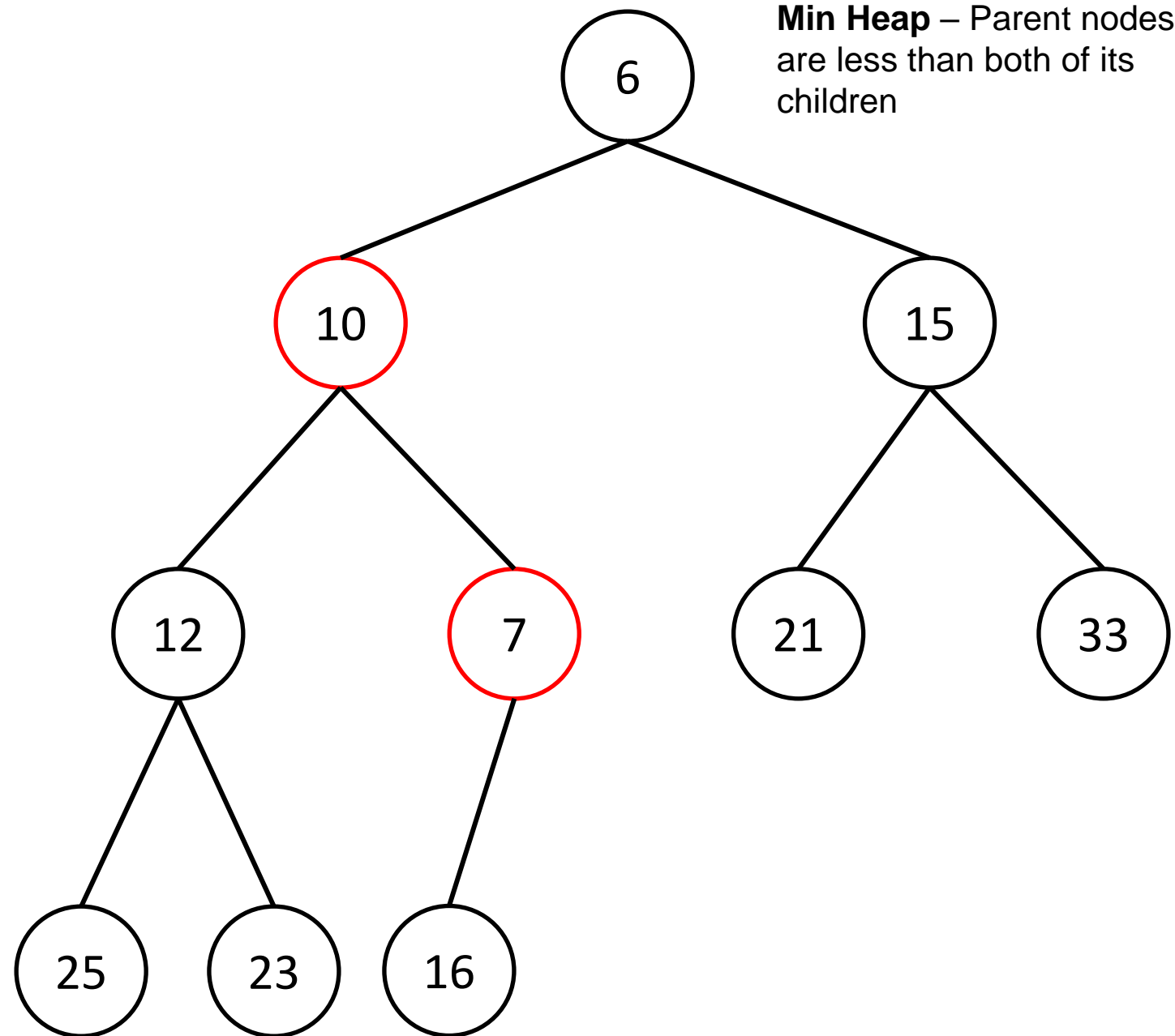
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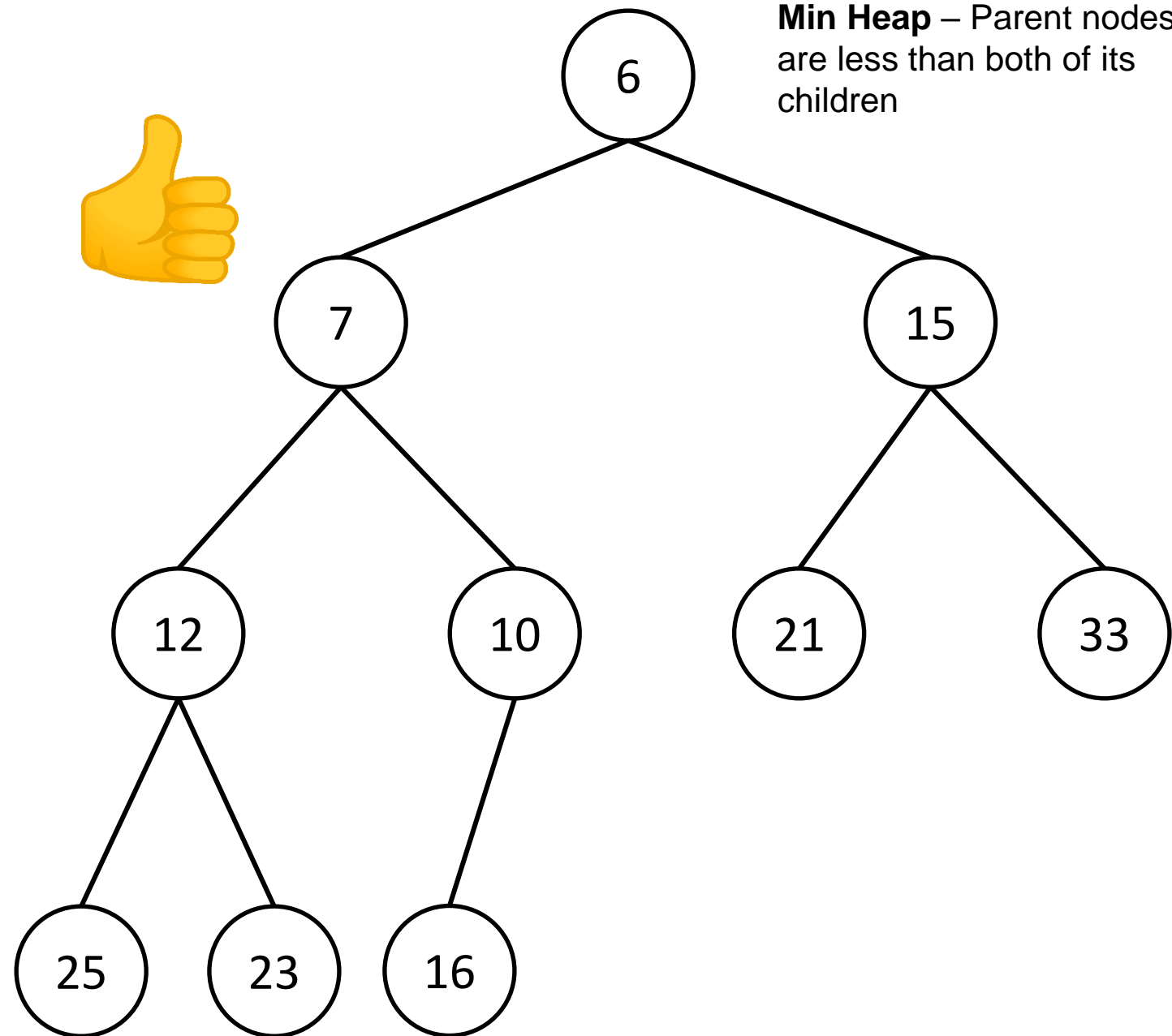
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Heap Operations - Insert

Min Heap – Parent nodes are less than both of its children

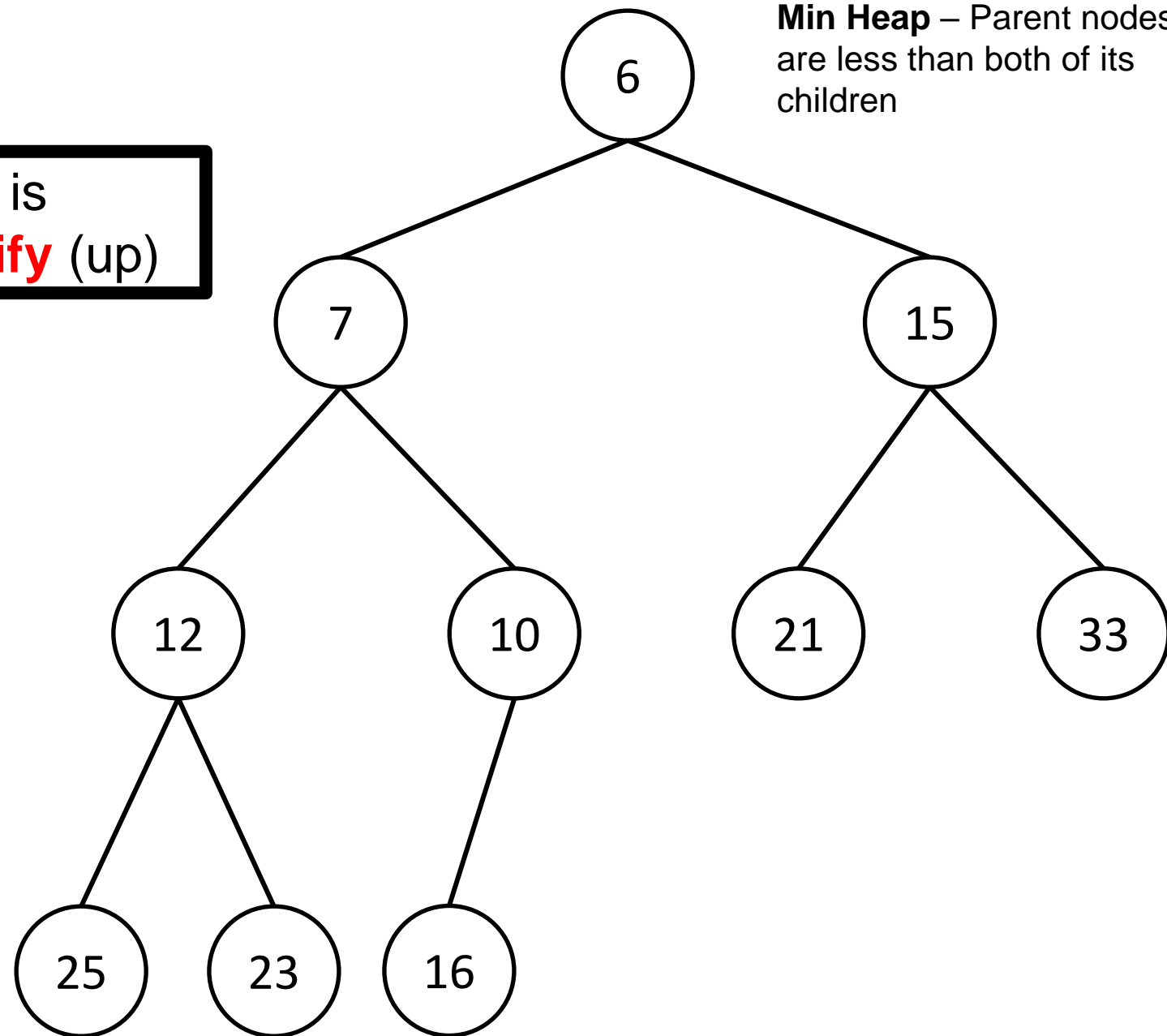
`add(7);`

This process is called **Heapify** (up)

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However, we are now violating the heap property

When new nodes are added, we may need to move it up in the tree



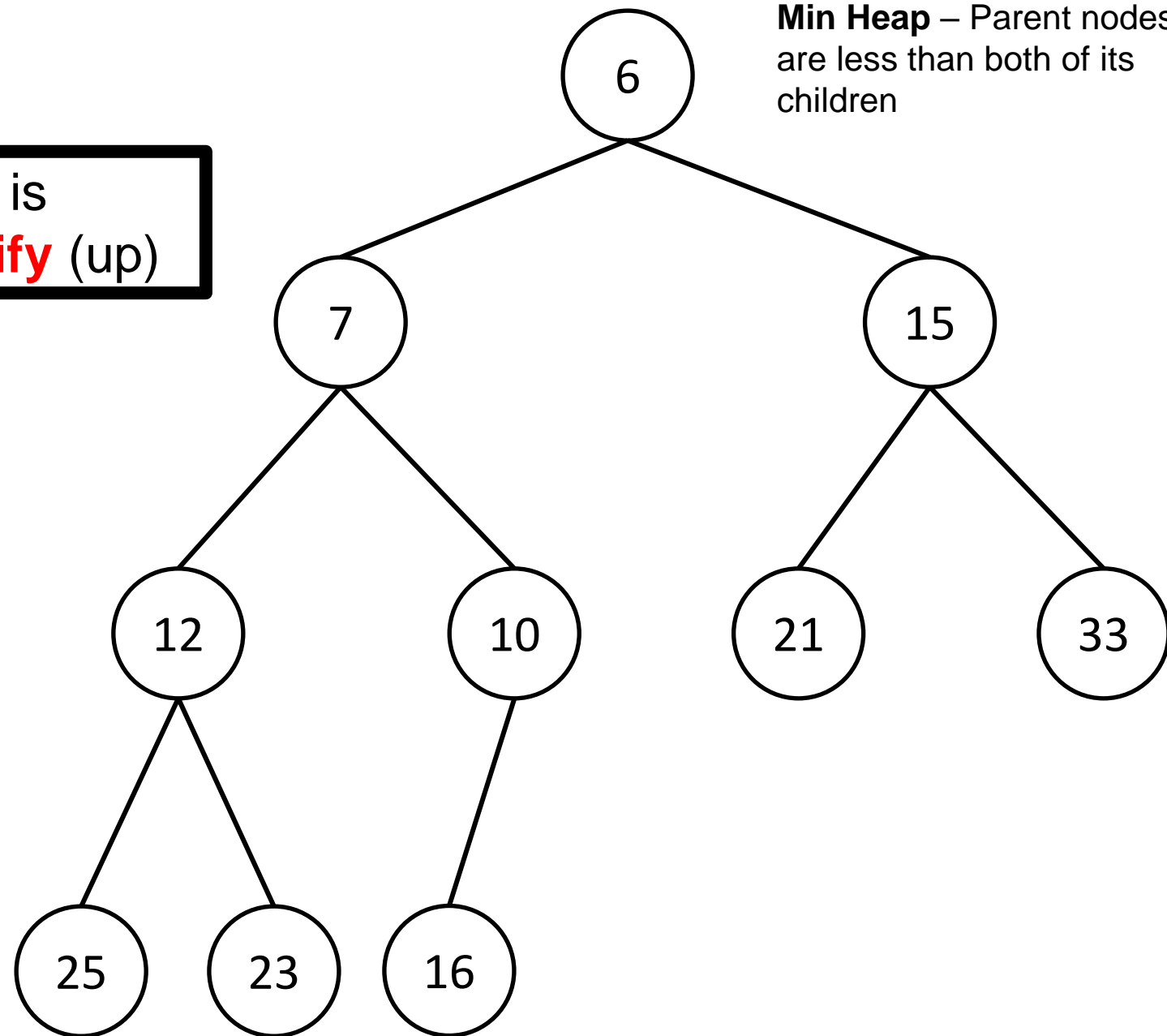
Heap Operations - Insert

add(7);

add(14);

This process is
called **Heapify** (up)

Min Heap – Parent nodes
are less than both of its
children



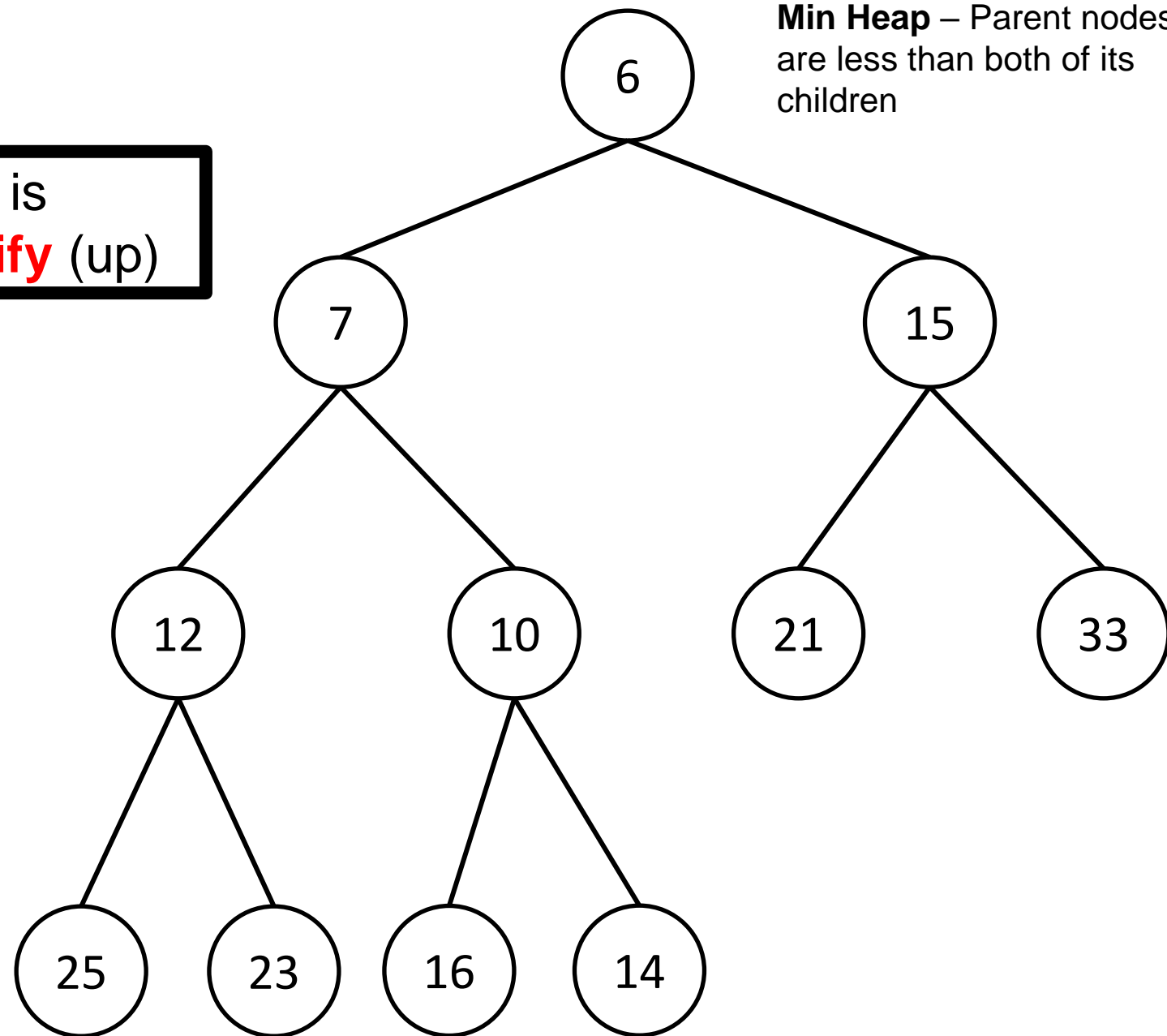
Heap Operations - Insert

`add(7);`

`add(14);`

This process is
called **Heapify** (up)

Min Heap – Parent nodes
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Heap Operations - Insert

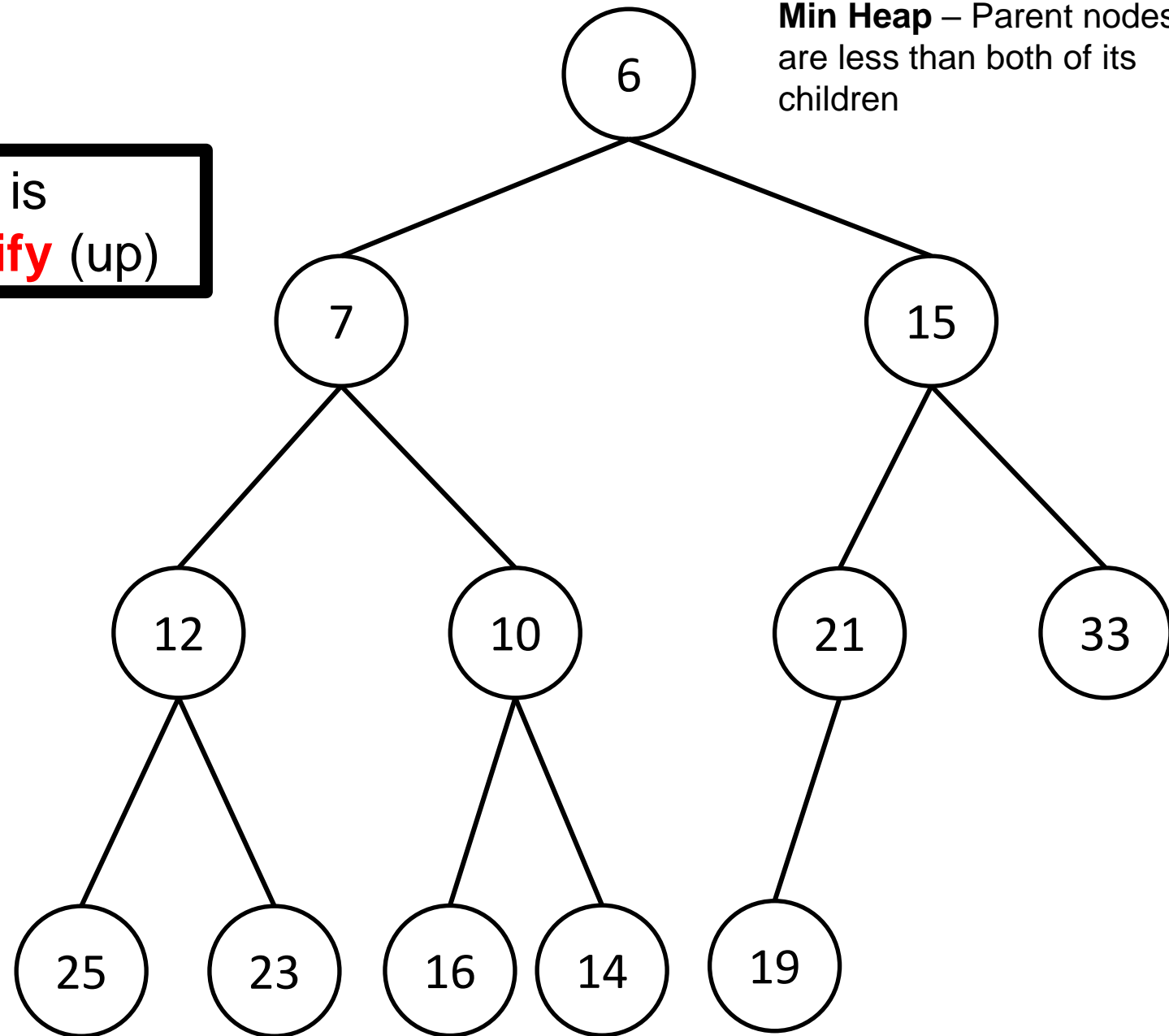
Min Heap – Parent nodes are less than both of its children

`add(7);`

`add(14);`

`add(19);`

This process is called **Heapify** (up)



Heap Operations - Insert

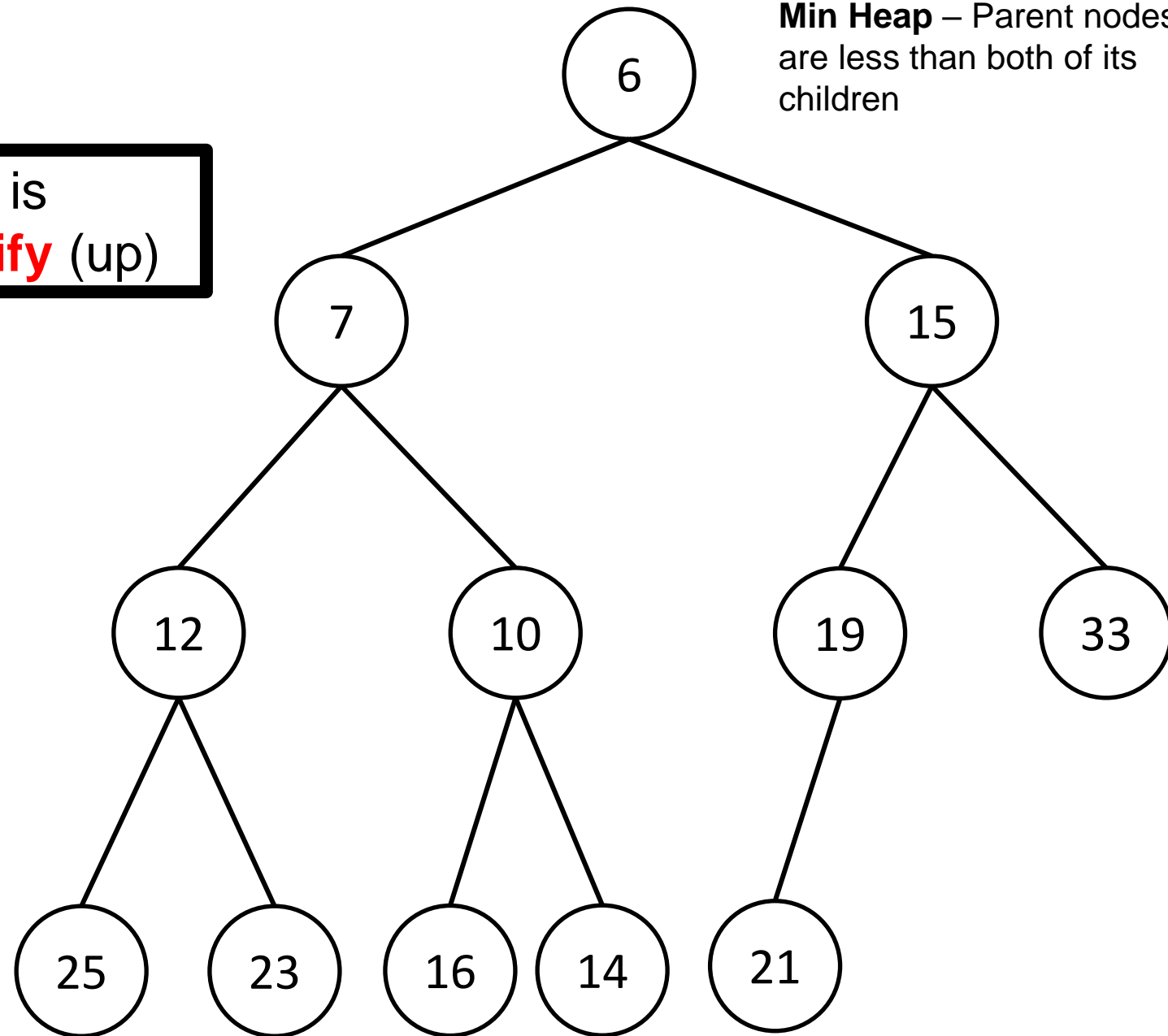
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Heap Operations - Insert

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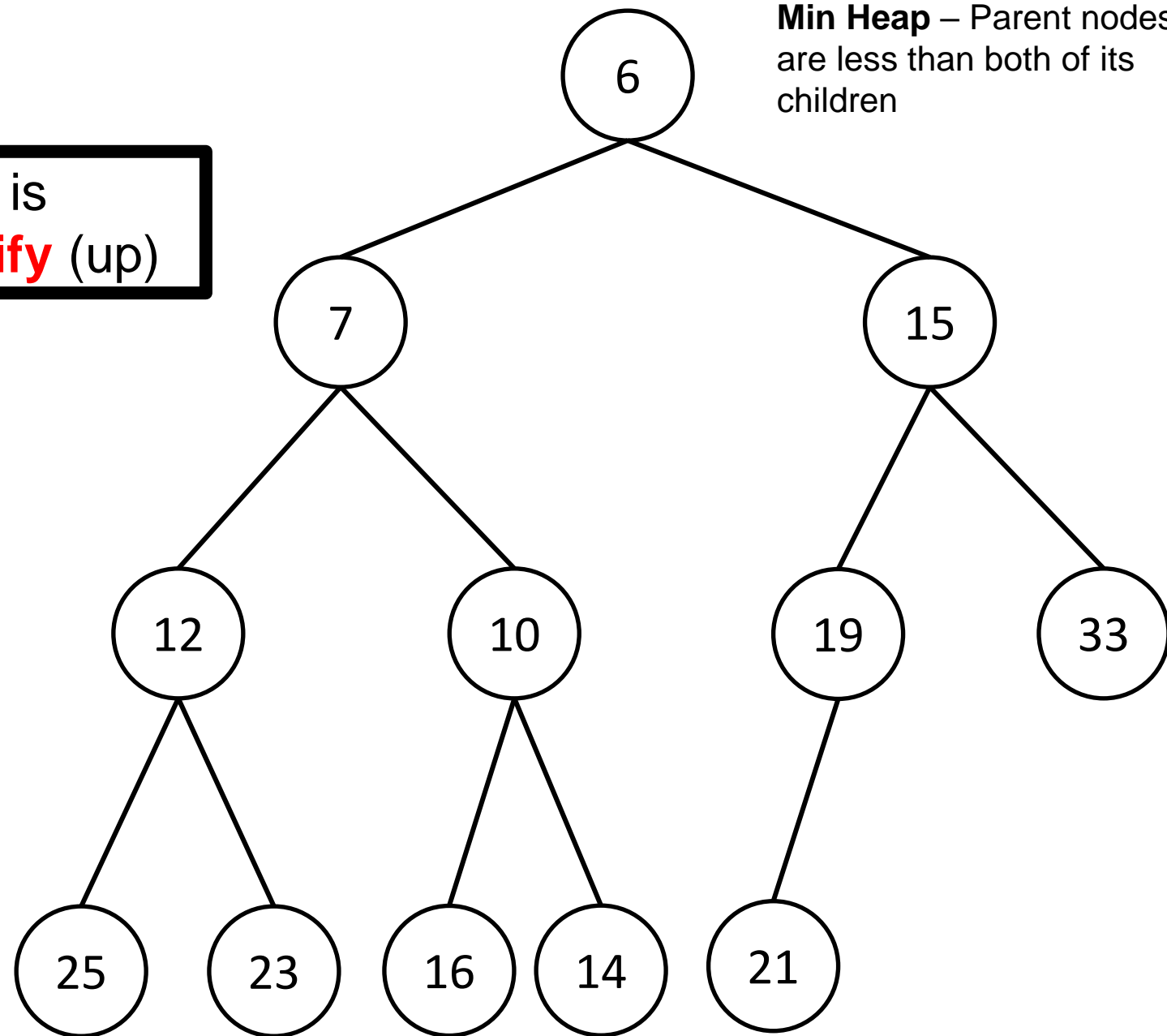
`add(14);`

`add(19);`

Running time?

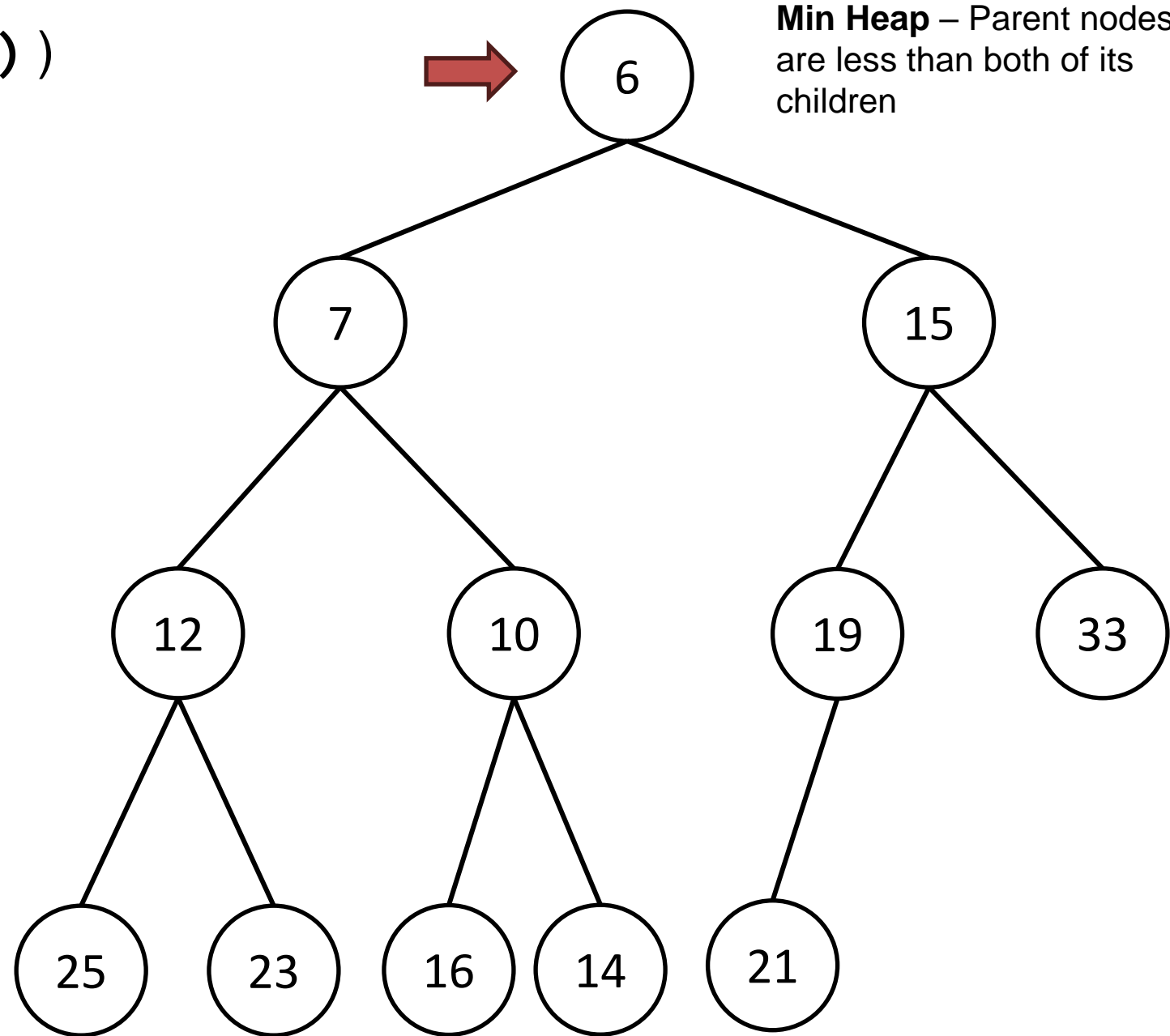
- Finding where to place new node: **$O(1)$** (this will make sense later)
- Insertion – **$O(1)$**
- Heapify Up – **$O(\log n)$**

Total Running Time: **$O(\log n)$**



Heap Operations – Removal (`poll()`)

When using a Heap, we only remove the root node, which will be either the maximum value or minimum value

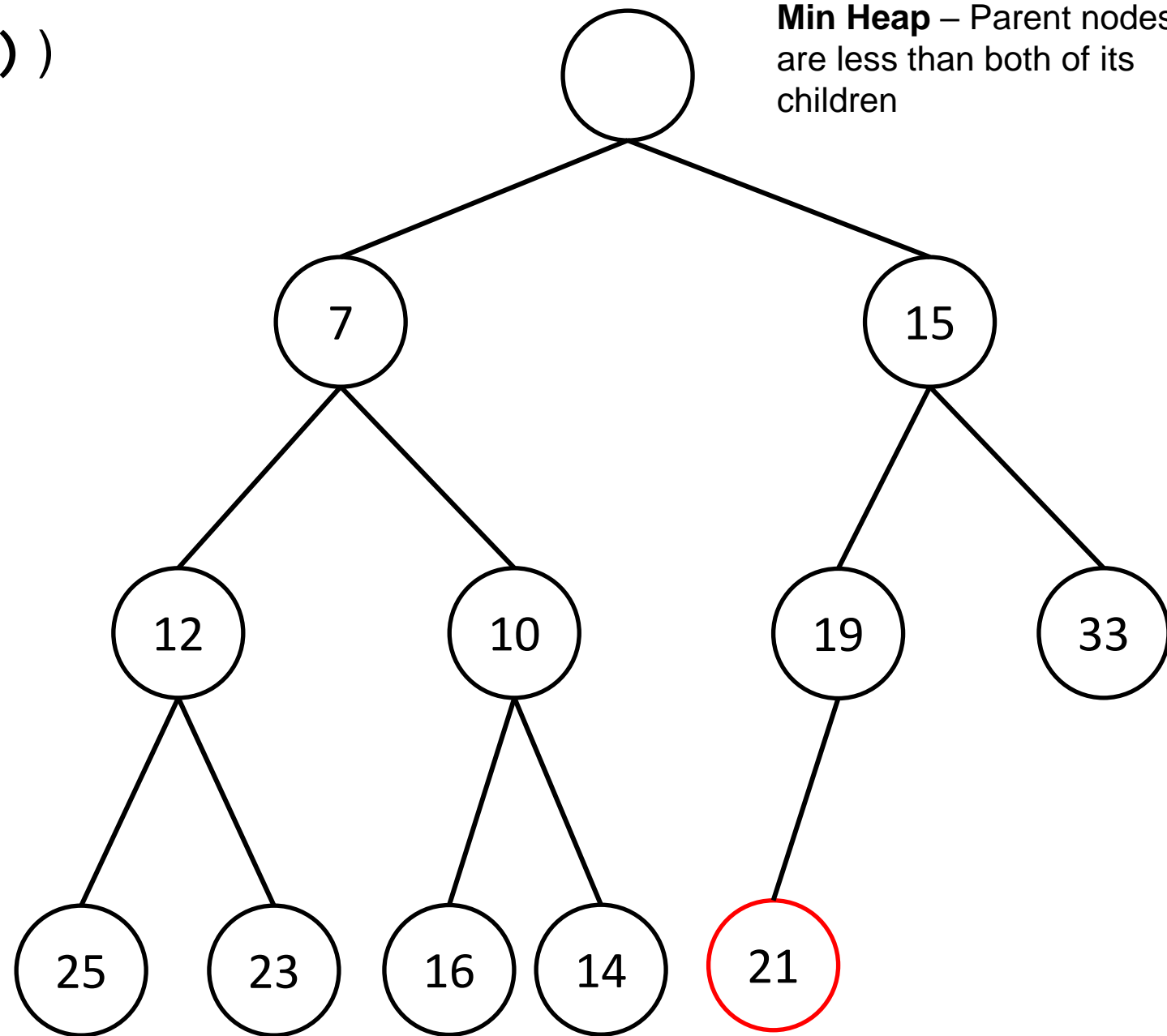


Heap Operations – Removal (`poll()`)

Min Heap – Parent nodes are less than both of its children

When using a Heap, we only remove the root node, which will be either the maximum value or minimum value

When the root is removed, we replace it with **the last node that was added to the heap**

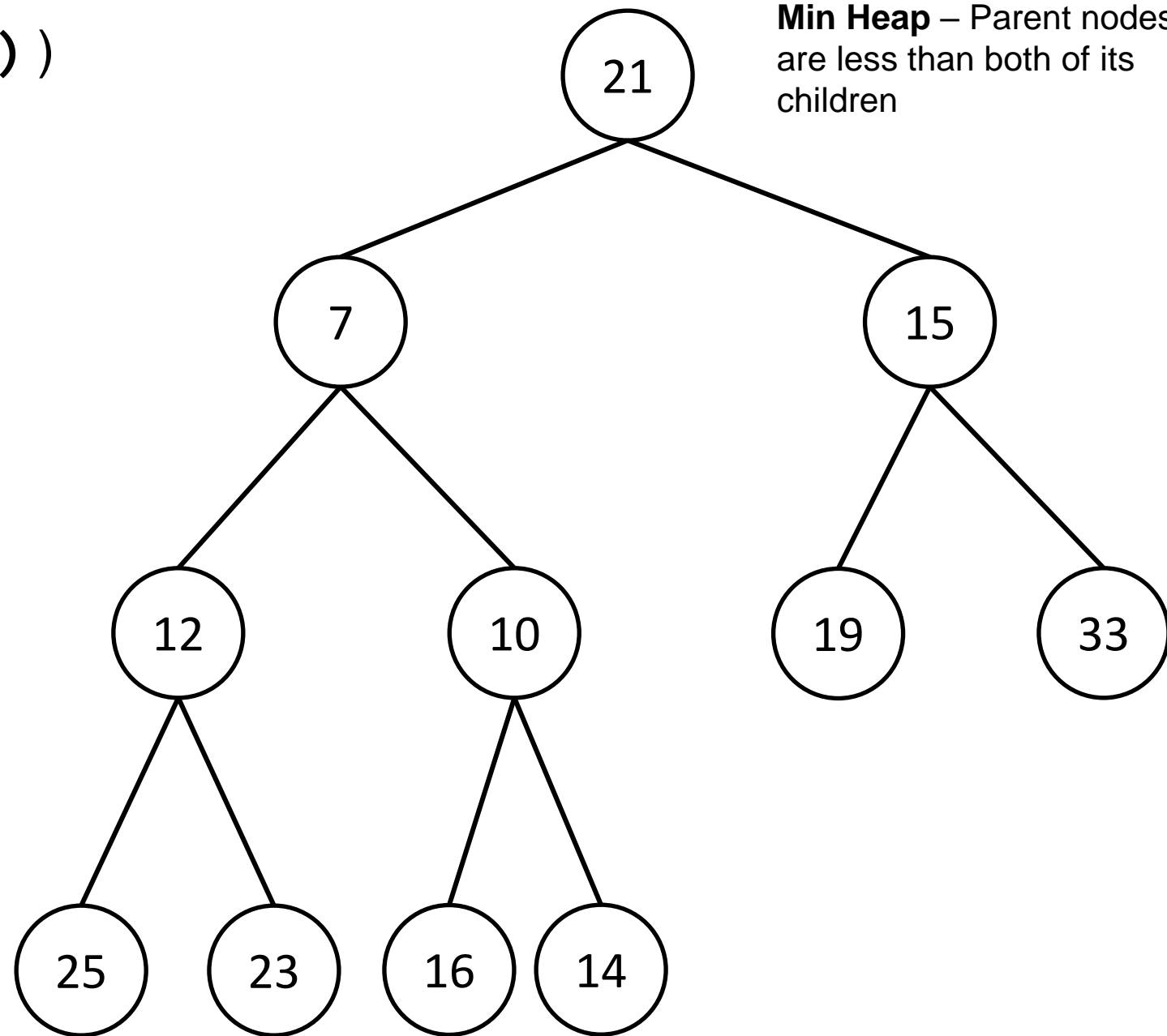


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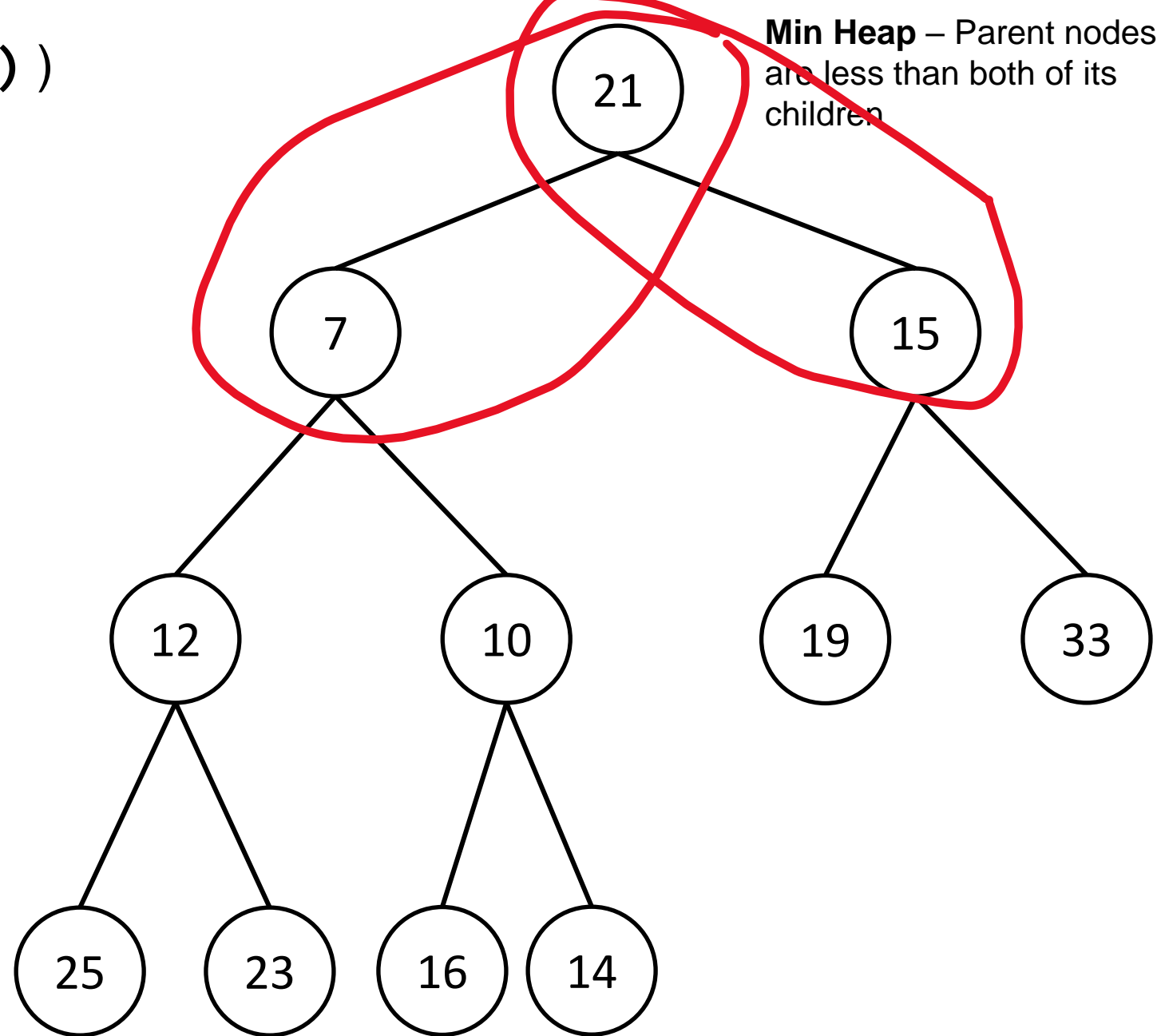
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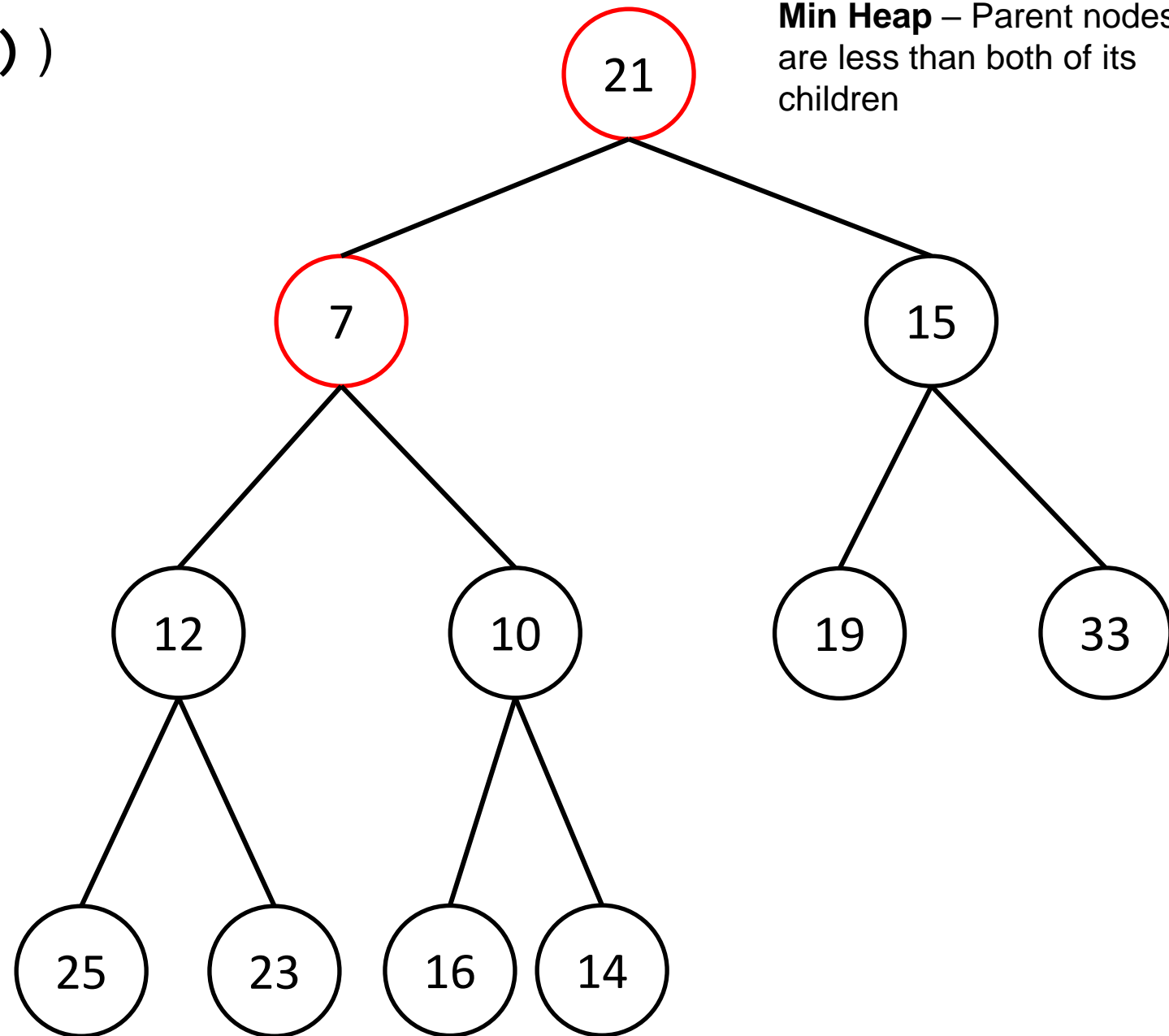
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When the root is replaced, it may need to be moved down in the tree



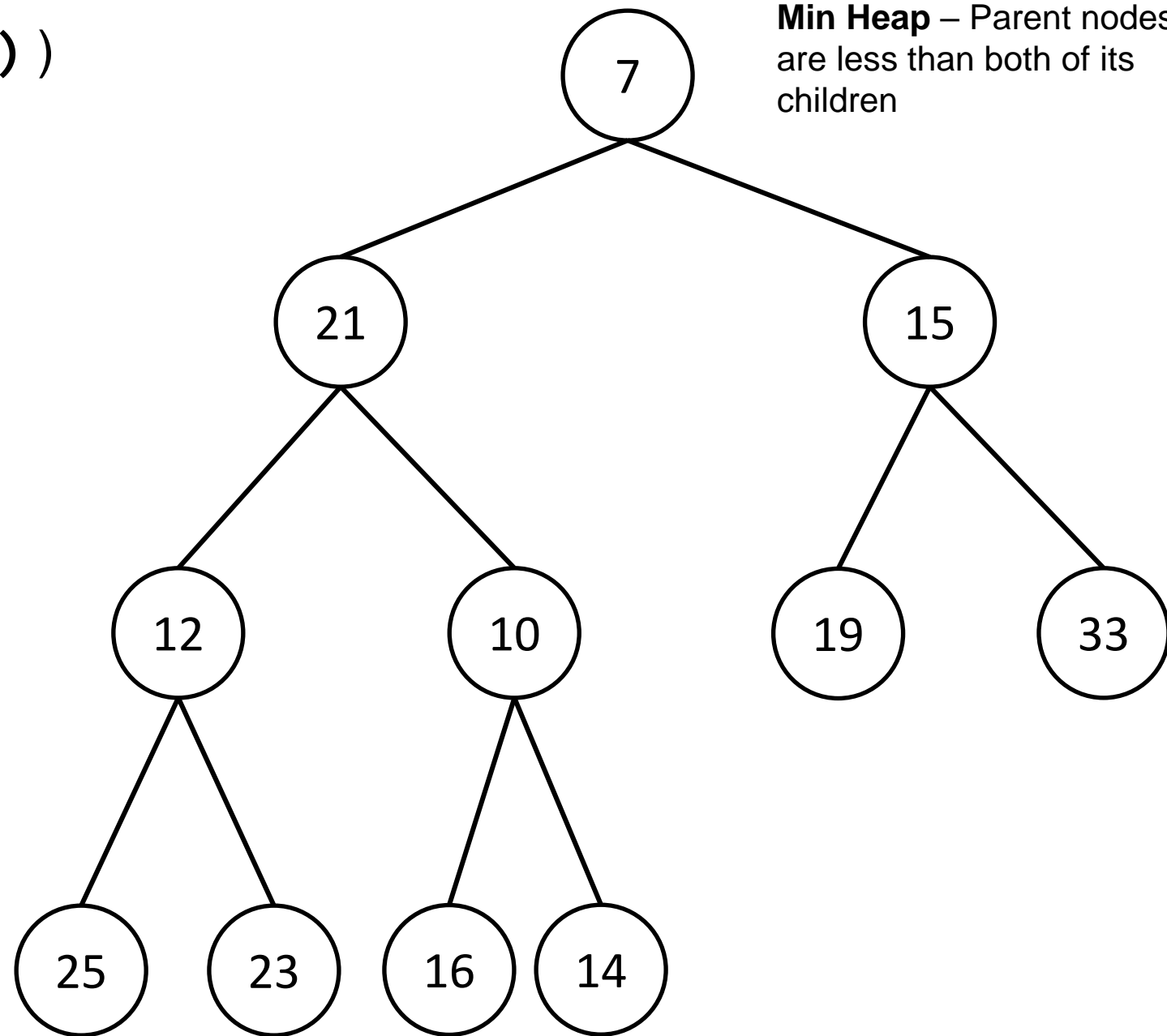
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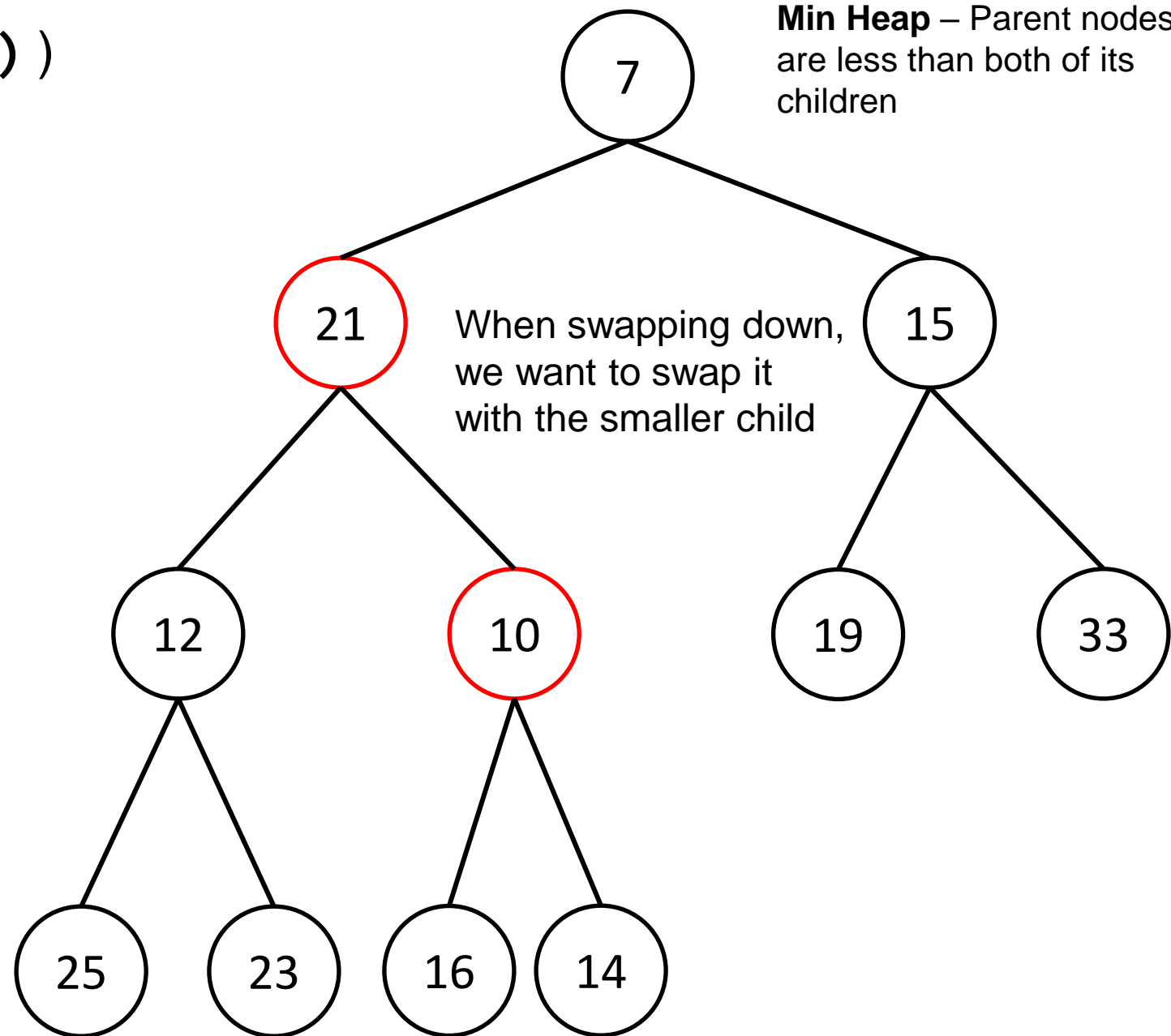


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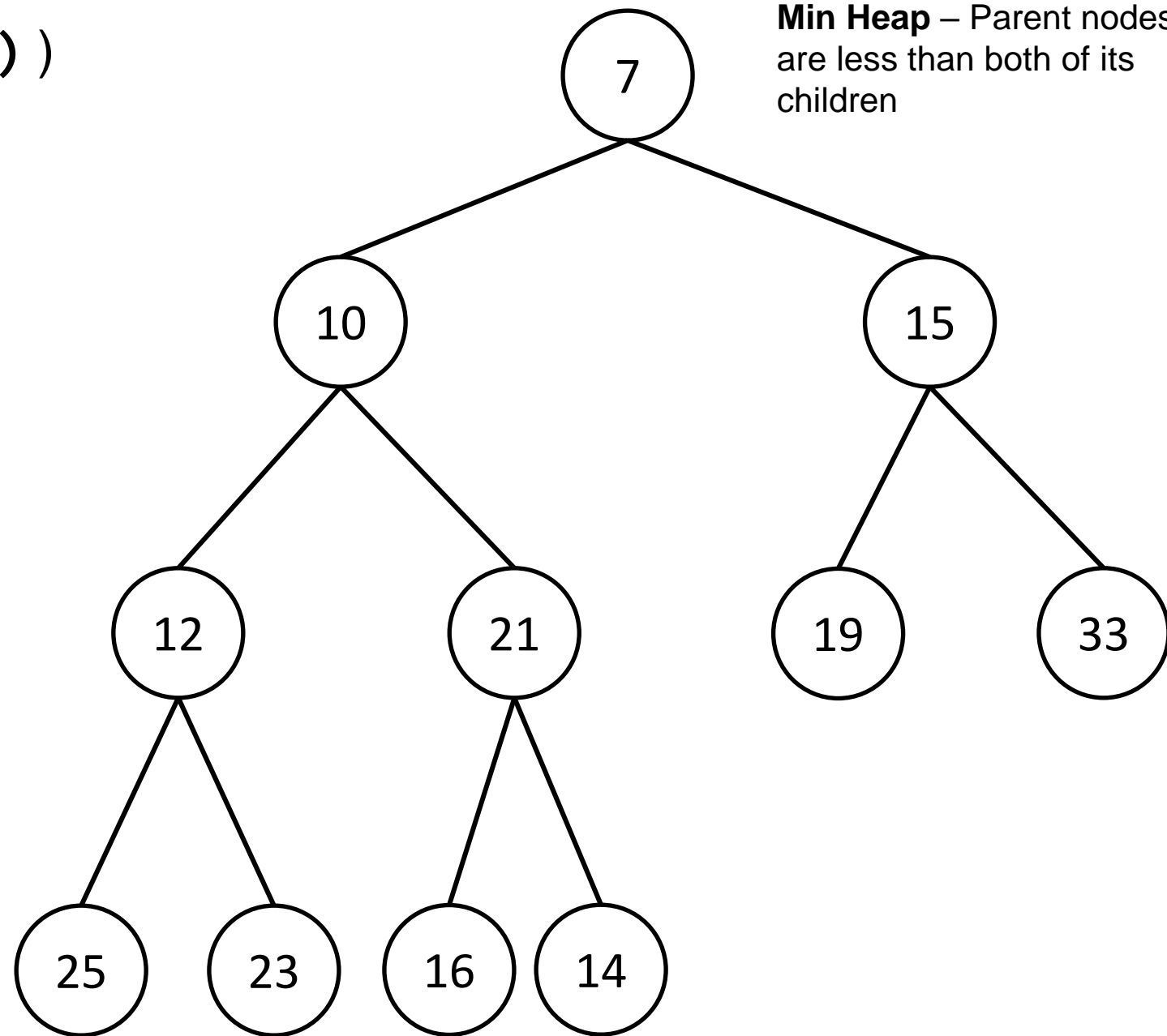
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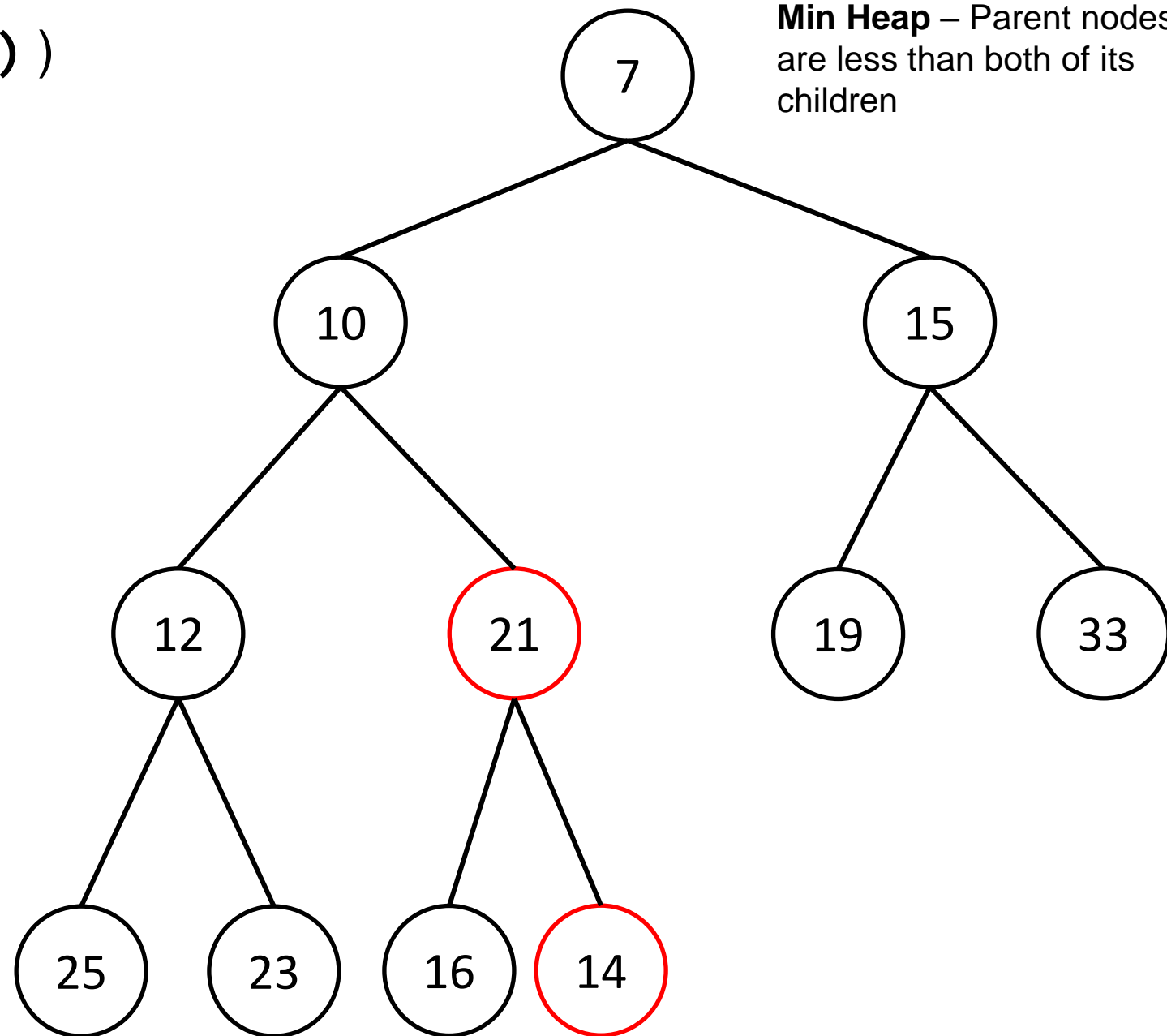
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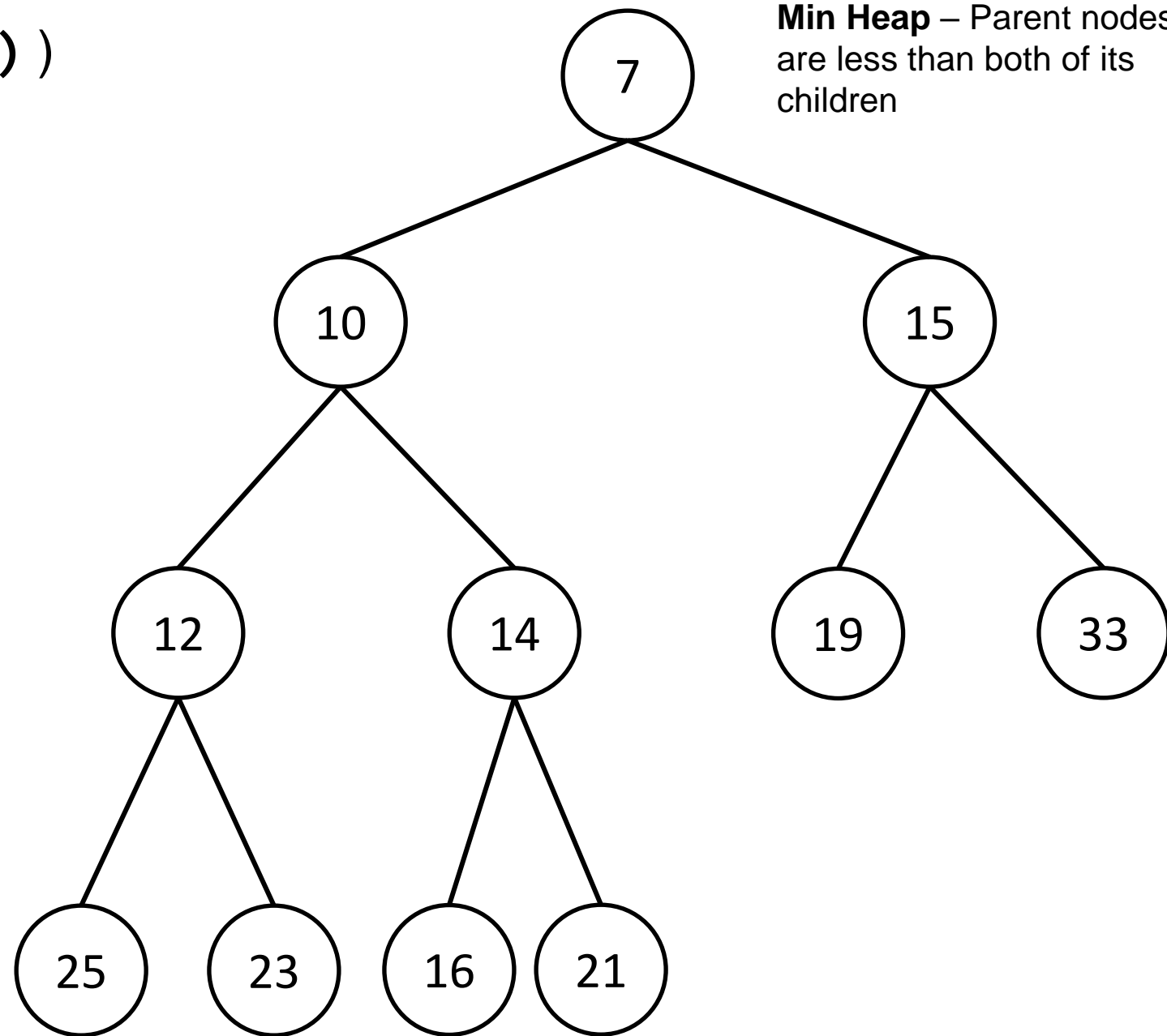
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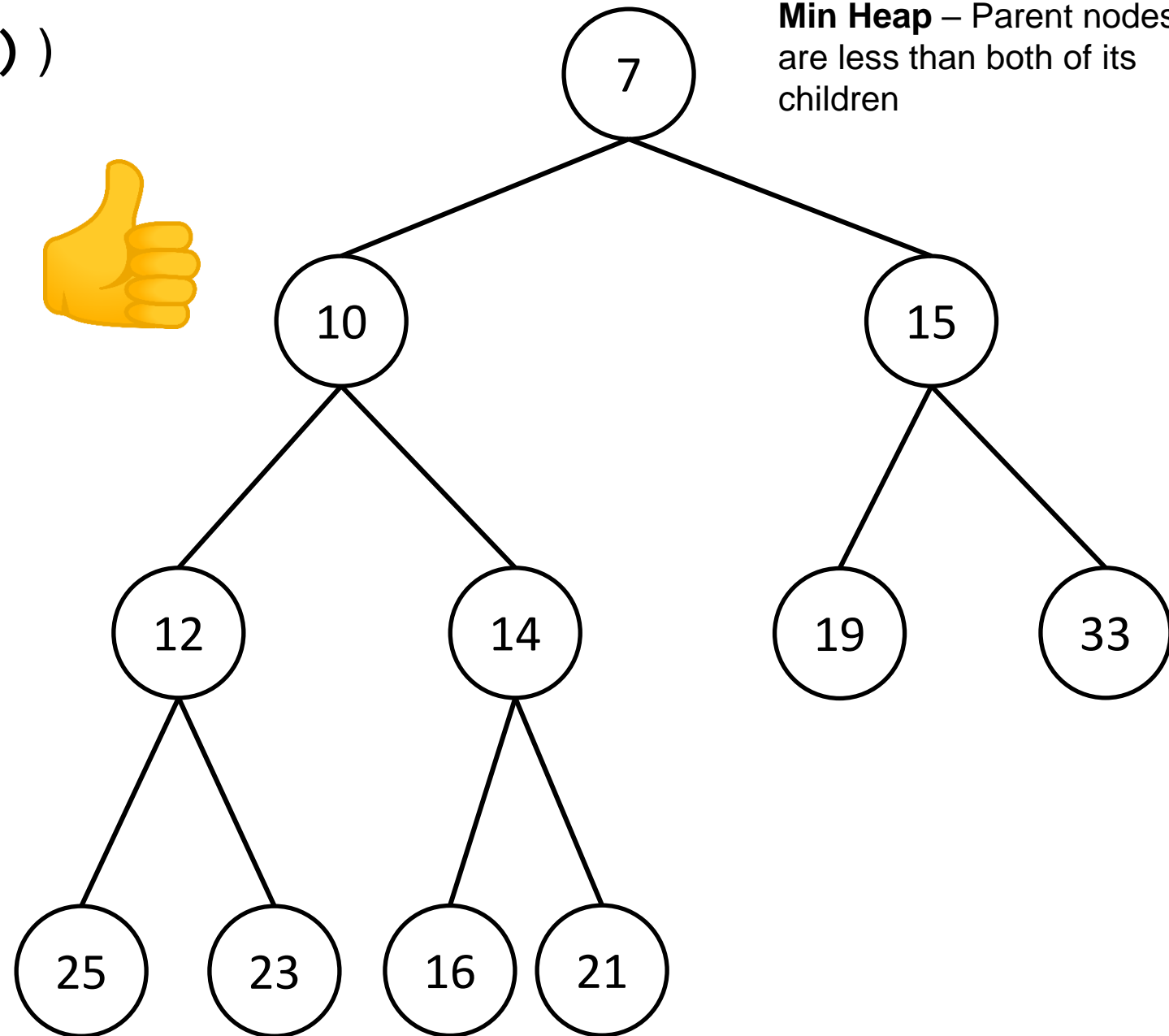
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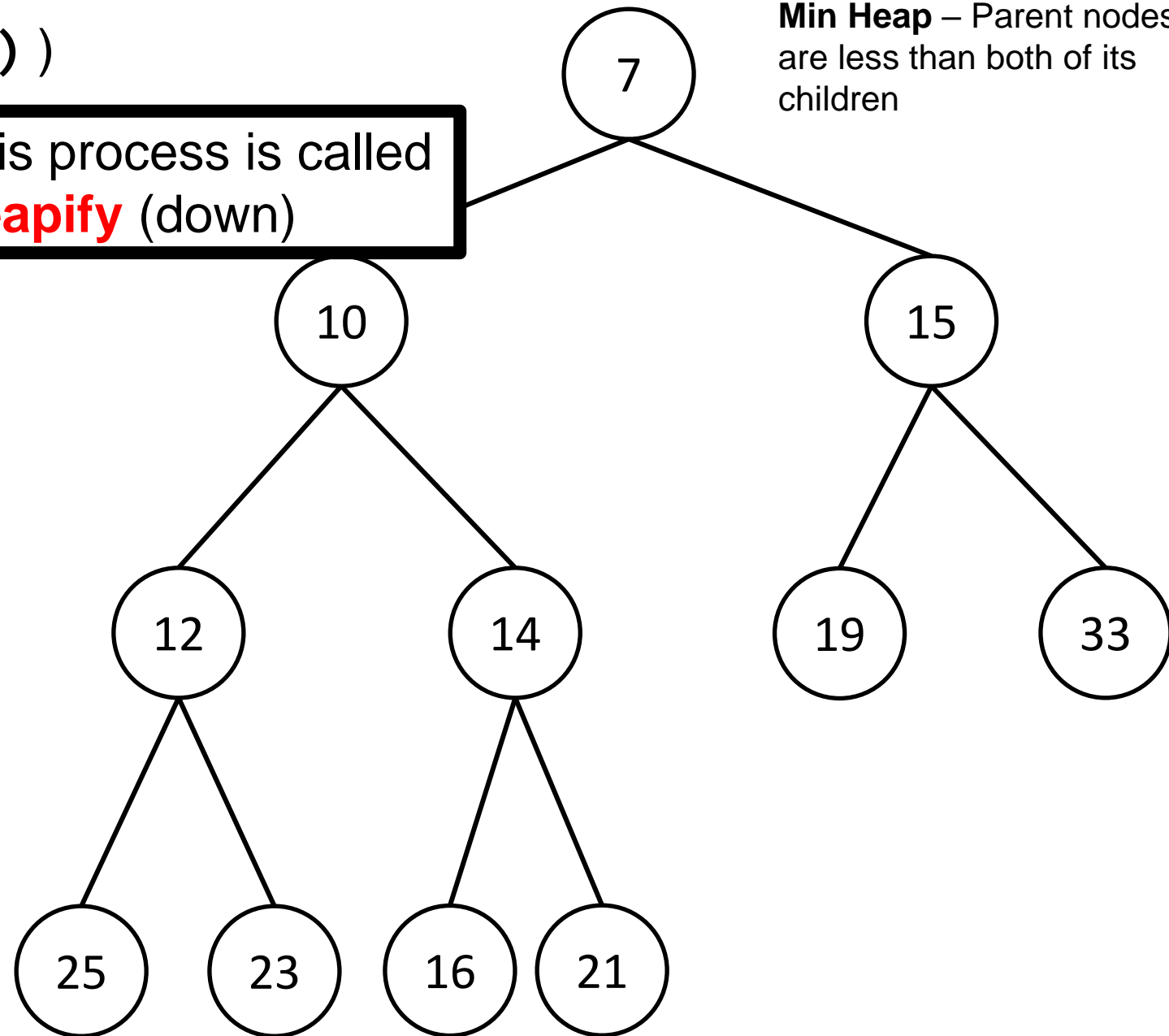
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This process is called **Heapify** (down)

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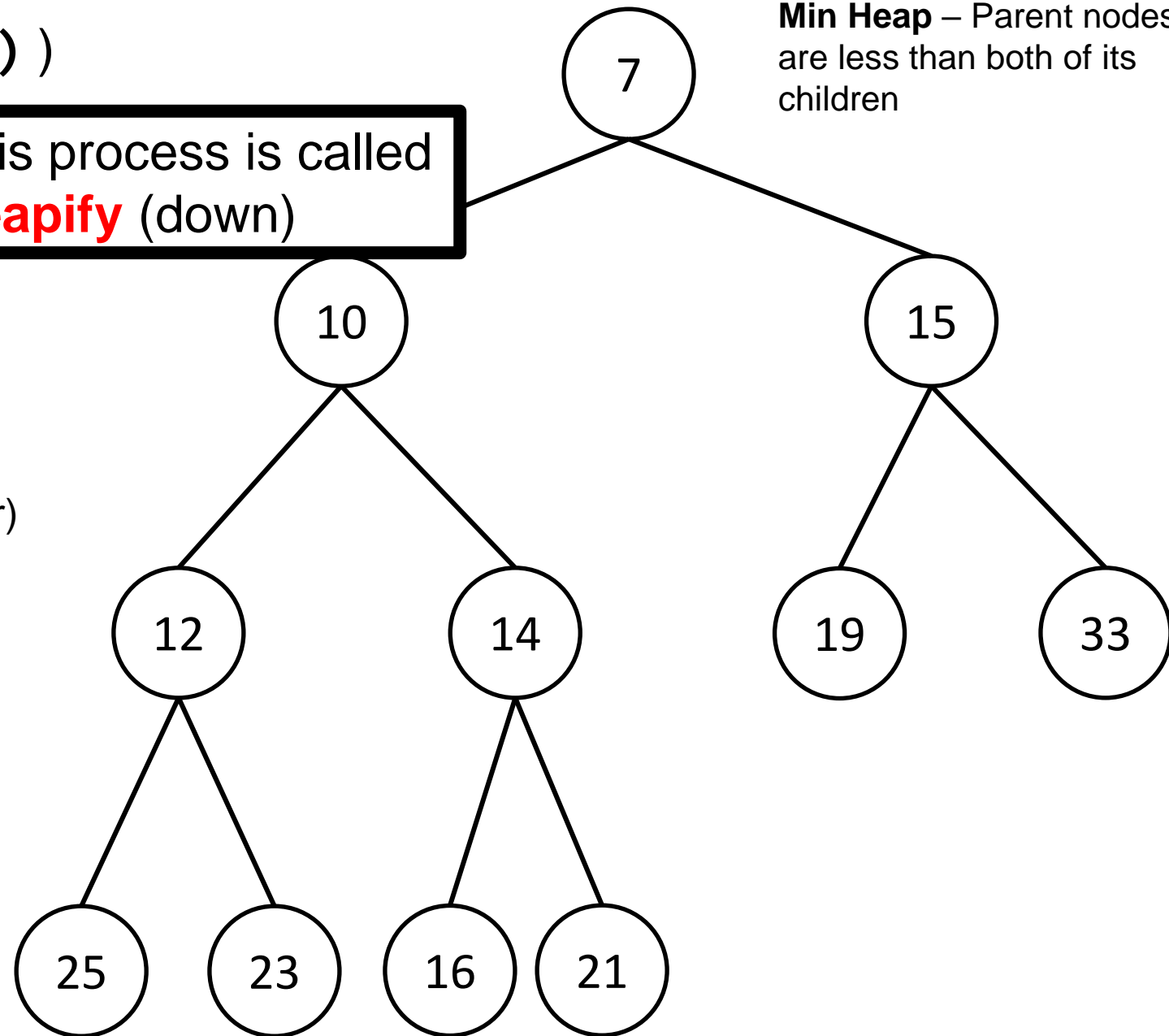
When the root is replaced, it may need to be moved down in the tree



Heap Operations – Removal (`poll()`)

Min Heap – Parent nodes are less than both of its children

This process is called **Heapify** (down)



Running time?

- Removing root: **$O(1)$**
- Replacing root: **$O(1)$** (this will make sense later)
- Heapify down: **$O(\log n)$**

Total running time: **$O(\log n)$**

Heap Operations – Removal (`poll()`)

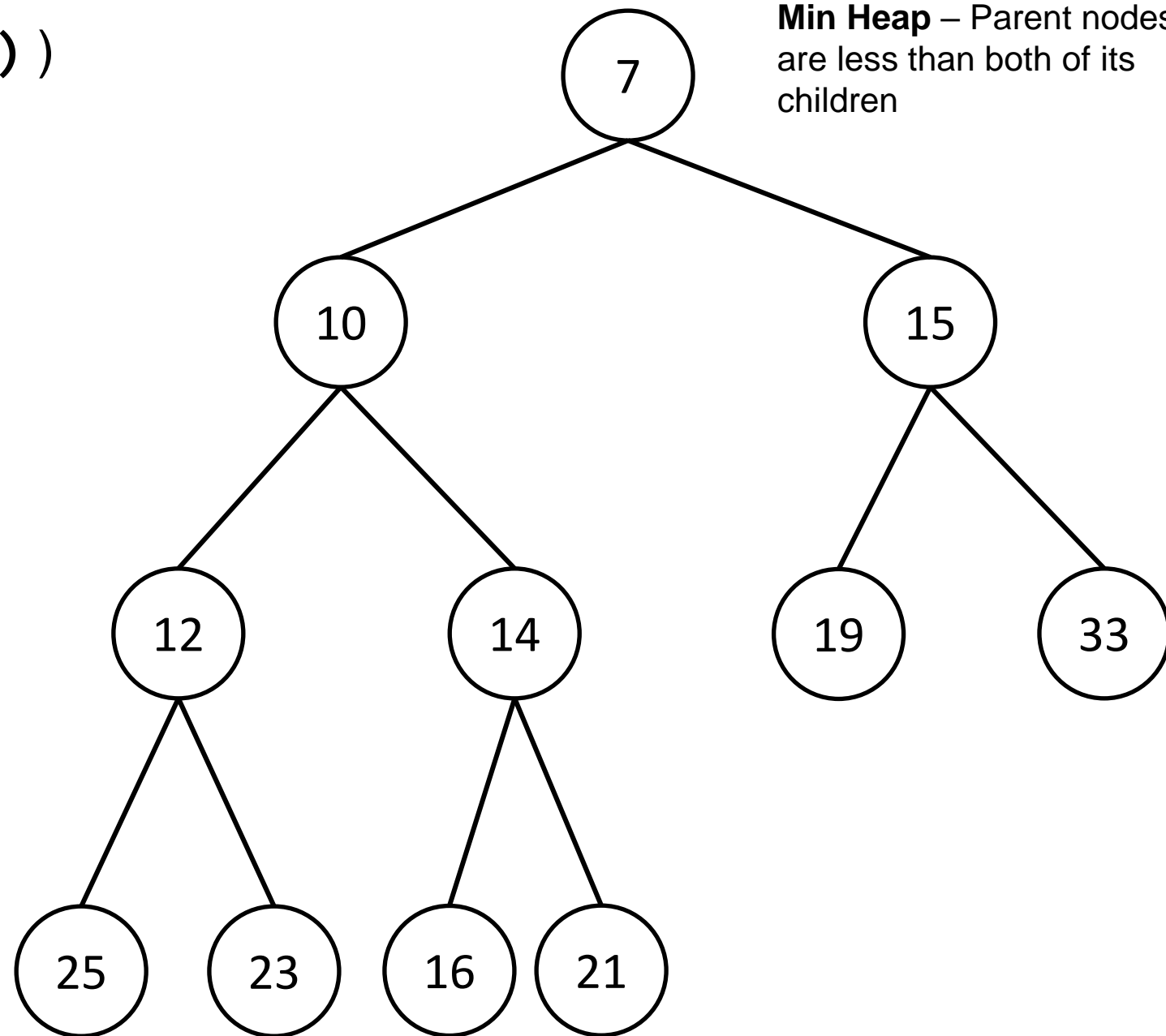
Min Heap – Parent nodes are less than both of its children

Heapify (up)

Moving the new leaf node **up** in the tree

Heapify (down)

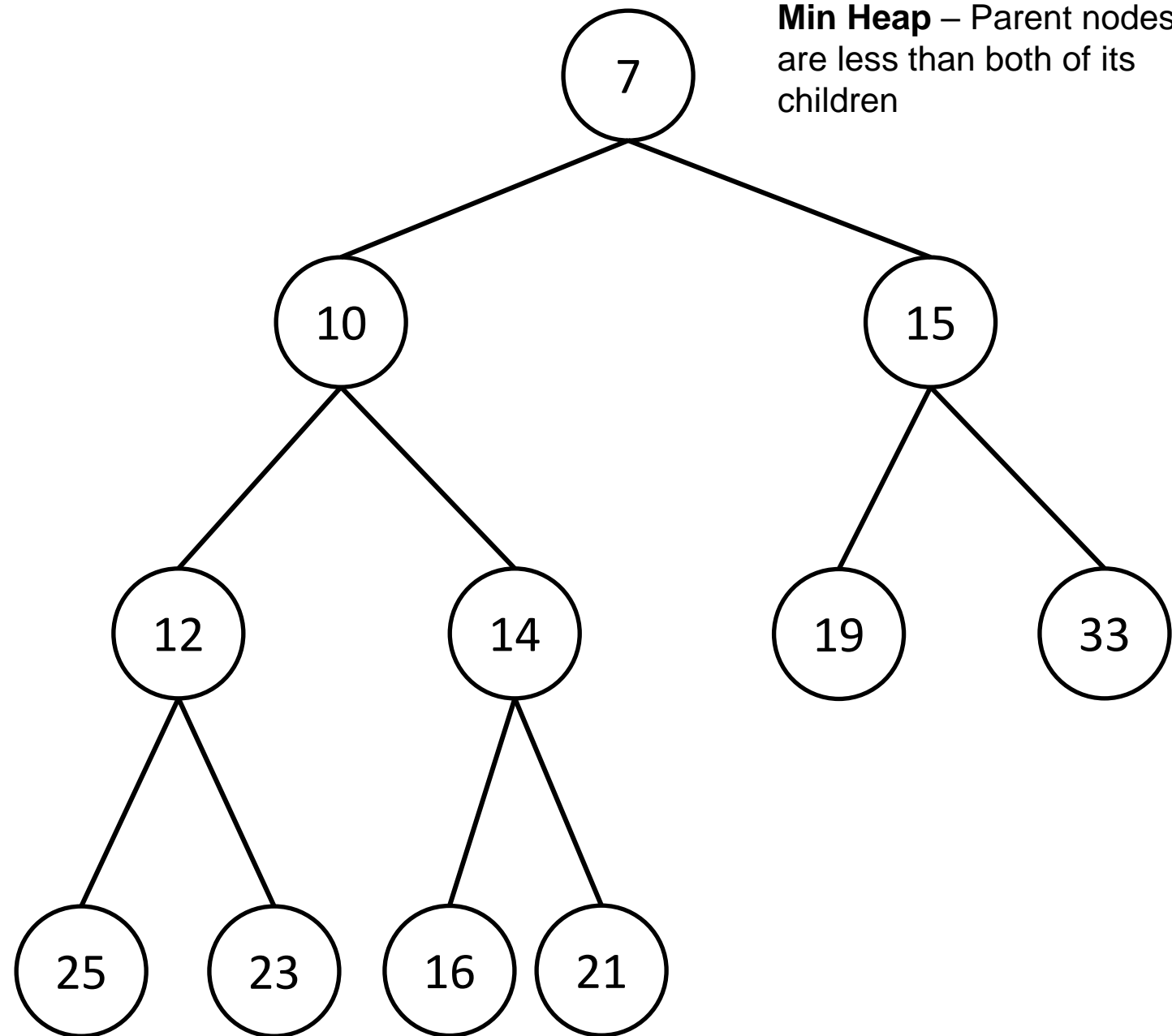
Moving the new root node **down** in the tree



Heap Representation

How to represent a heap?

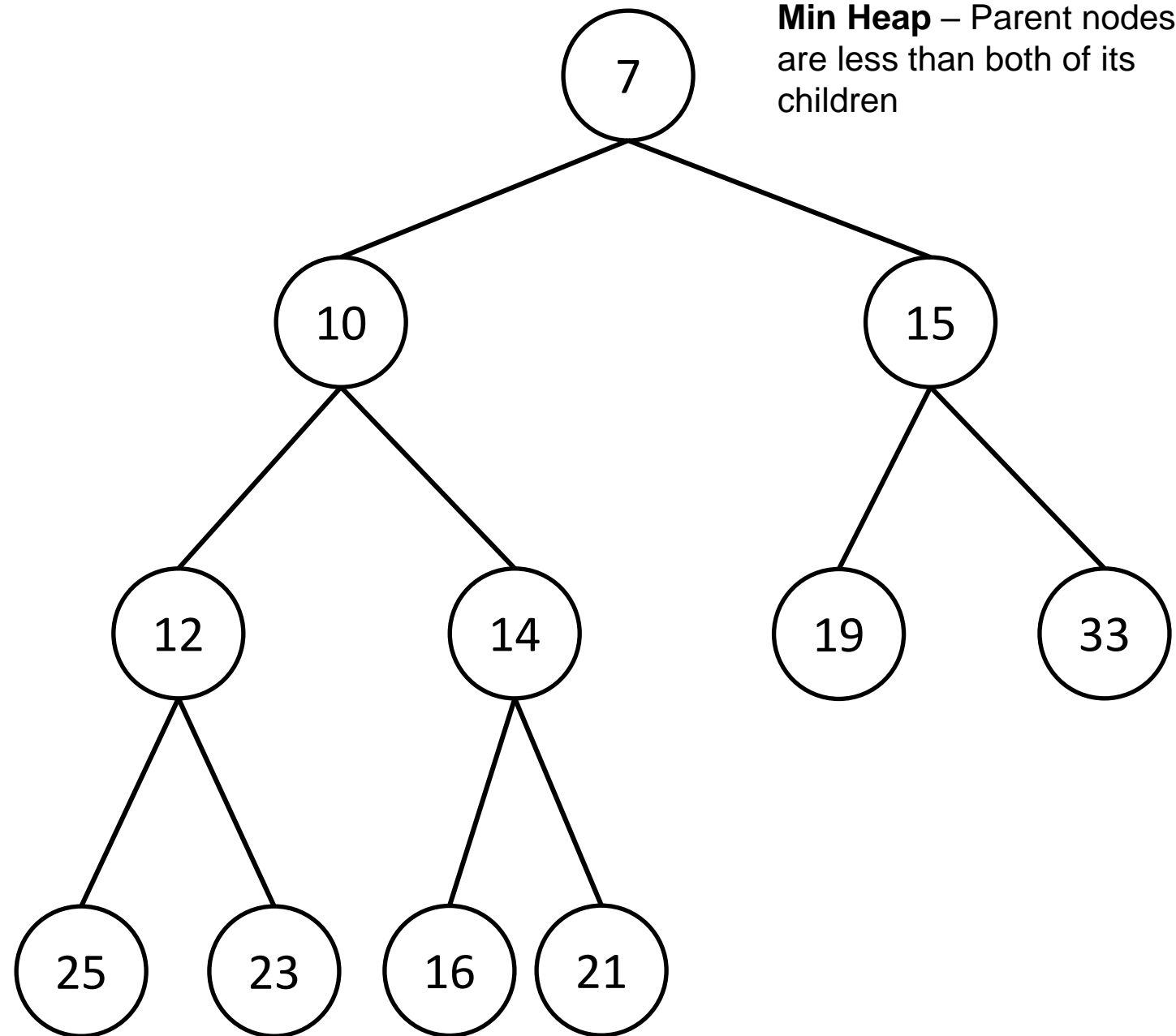
```
public class HeapNode{  
    Node leftChild;  
    Node rightChild;  
    Node parent;  
    (...)  
}
```



Heap Representation

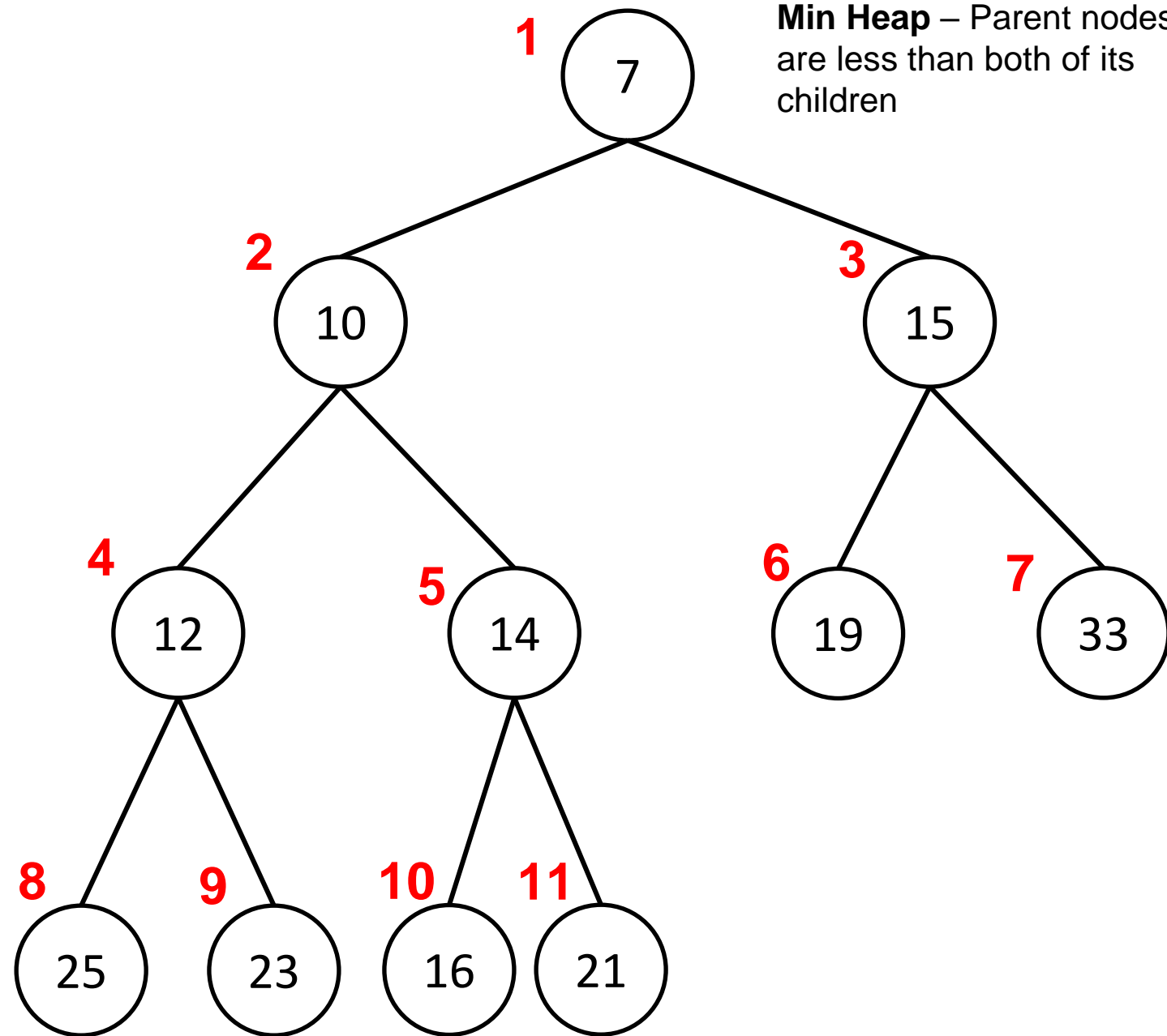
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Heap Representation

Min Heap – Parent nodes are less than both of its children

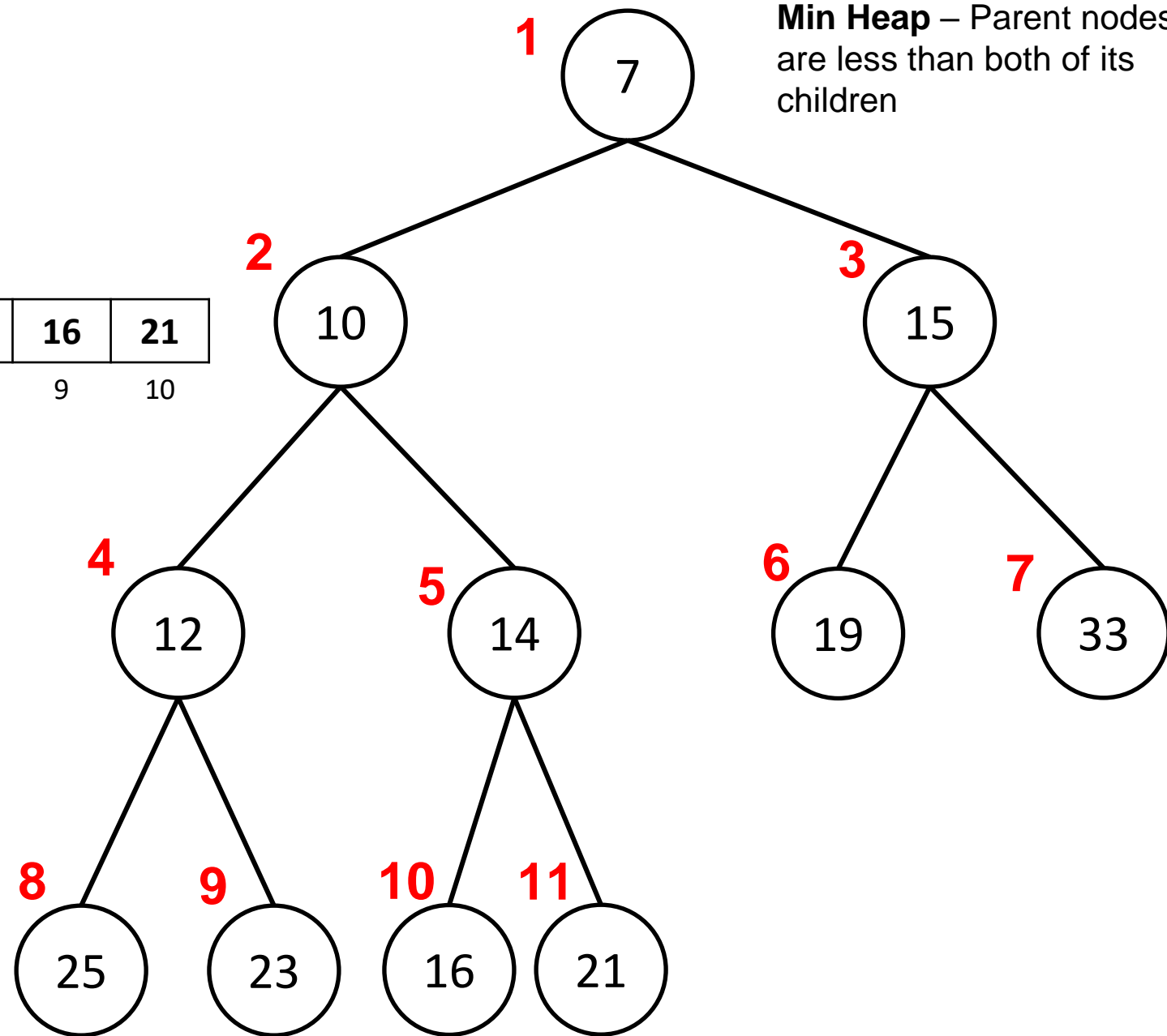


Heap Representation

Min Heap – Parent nodes are less than both of its children

Array

7	10	15	12	14	19	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10



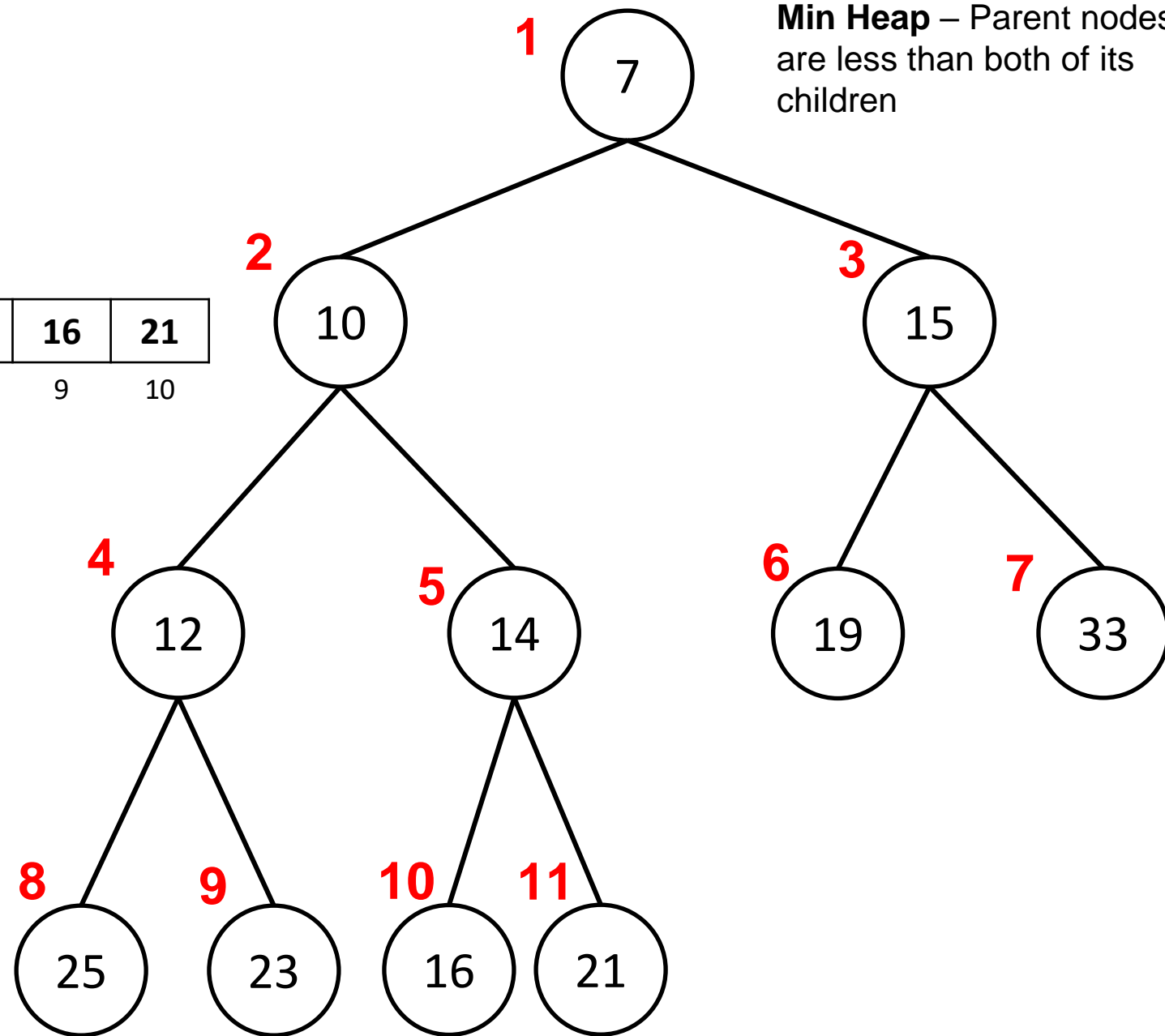
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Given a spot in the array, how can we find its children?



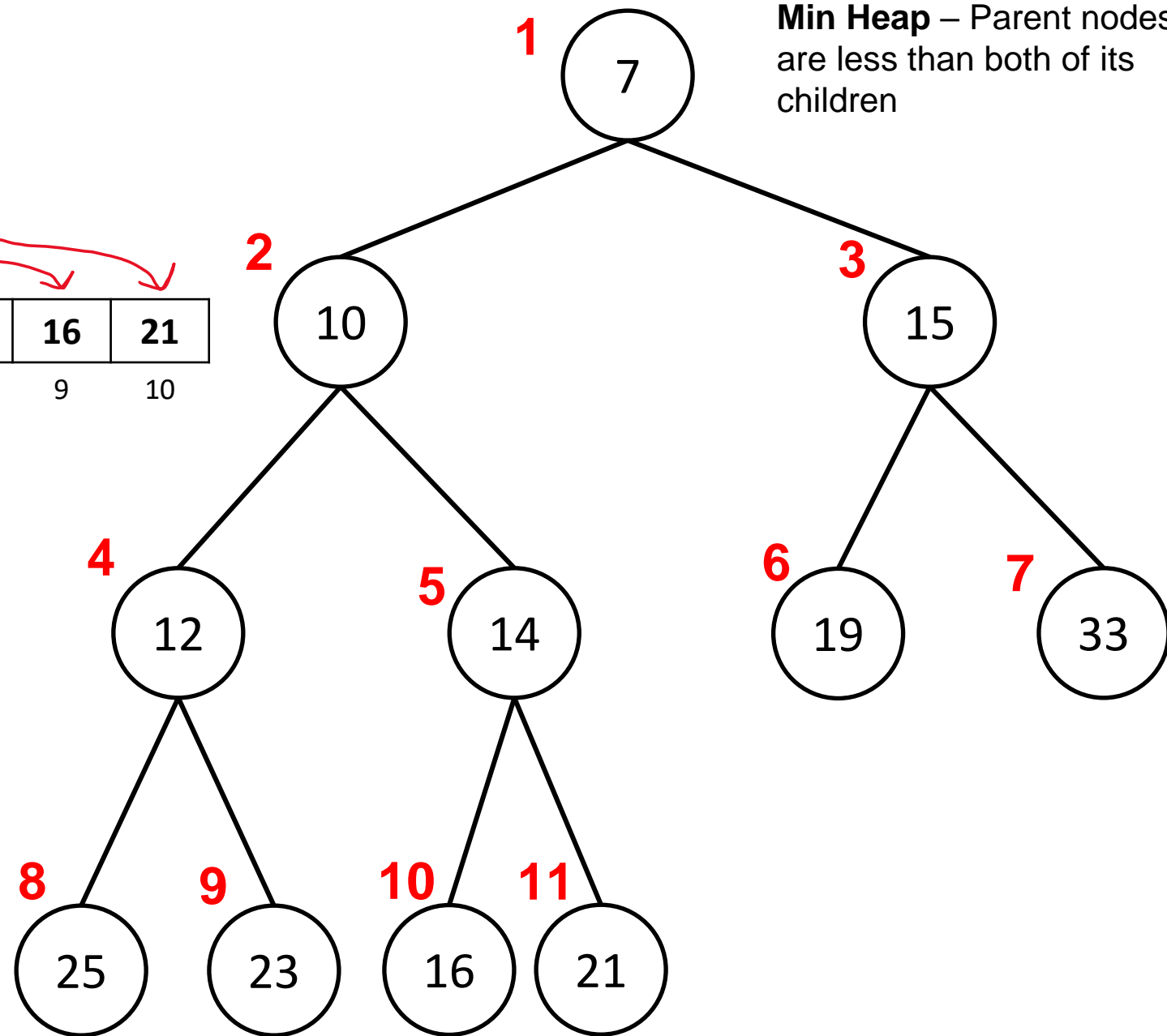
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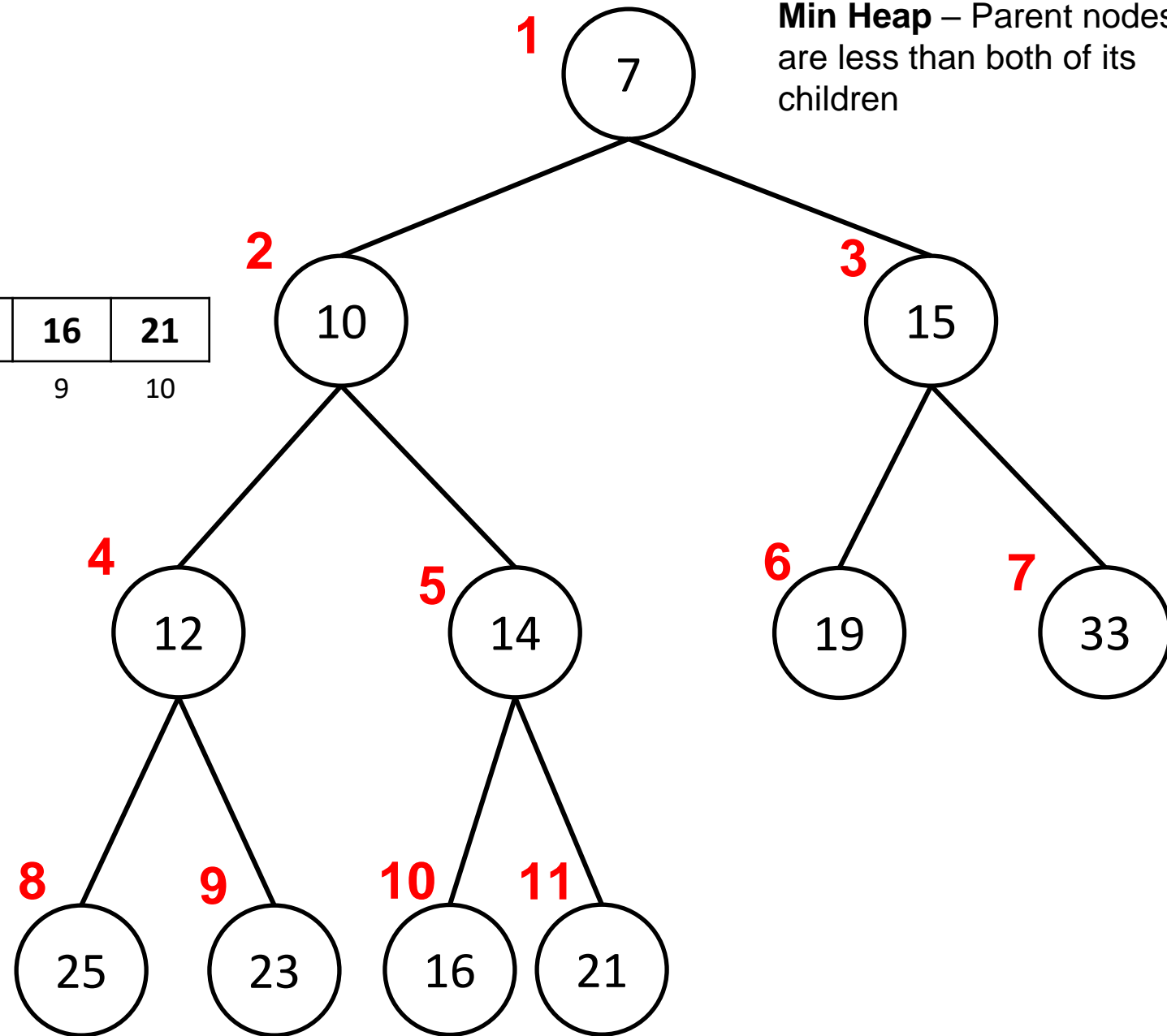
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Because this is a complete binary tree, there is a pretty nifty formula for this



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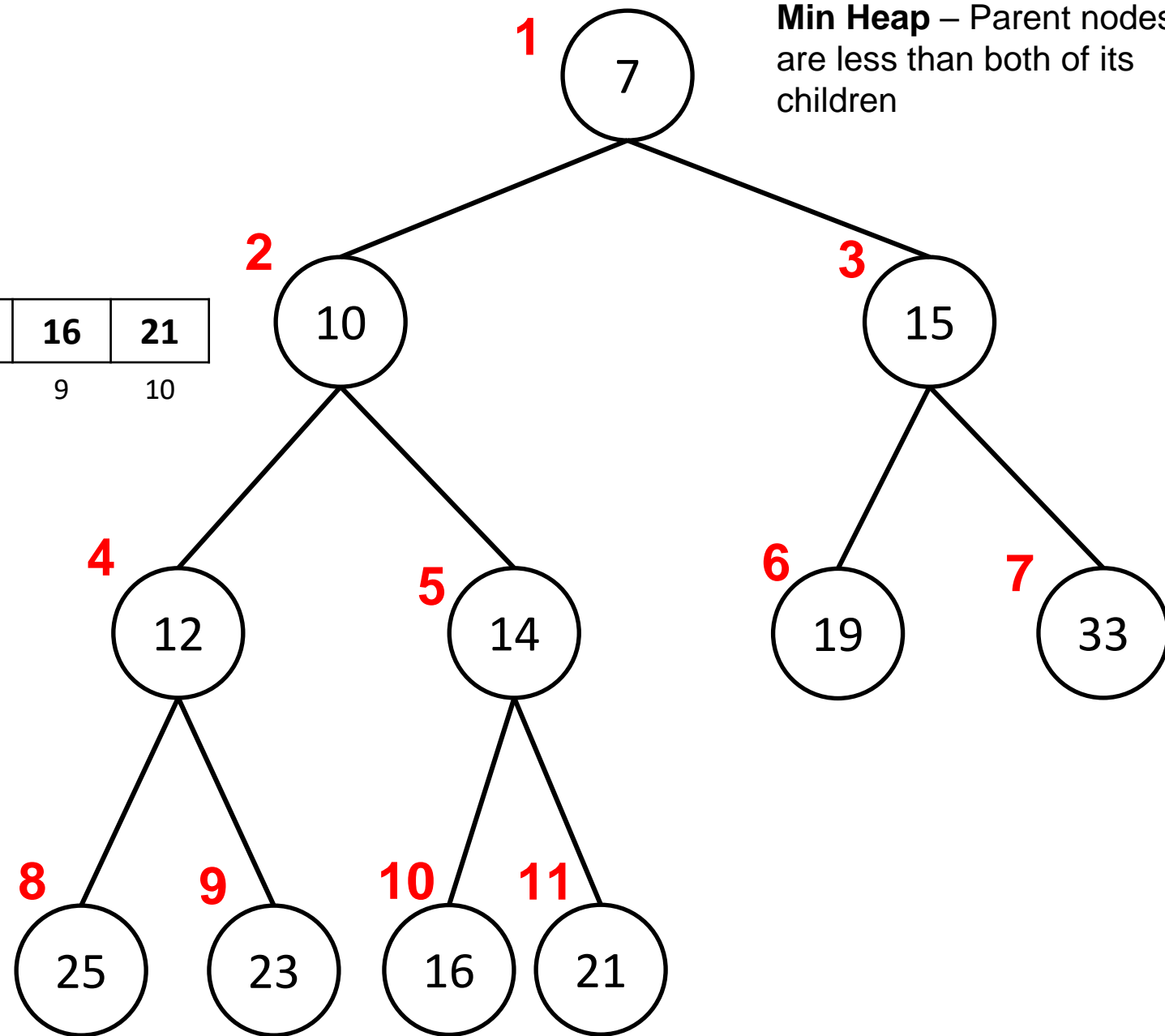
For a given element at index i

Its left child will be located at index:

$$2 * i + 1$$

Its right child will be located at index:

$$2 * i + 2$$



Heap Representation

Left Child = $2 * 4 + 1 = \text{index } 9 !$

Array

7	10	15	12	14	19	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

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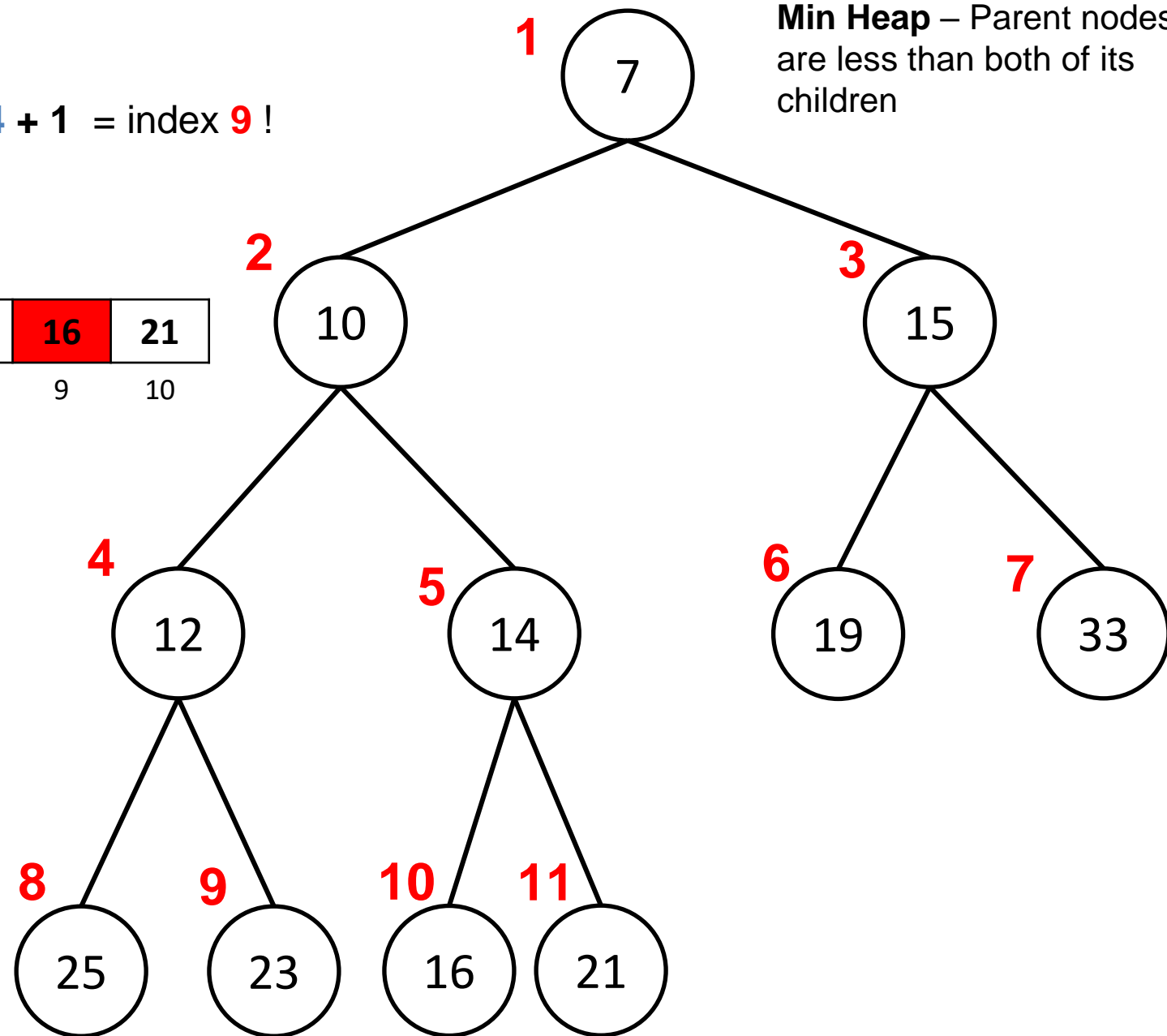
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Min Heap – Parent nodes are less than both of its children



Heap Representation

Left Child = $2 * 4 + 1 = \text{index } 9 !$
Right Child = $2 * 4 + 2 = \text{index } 10 !$

Array

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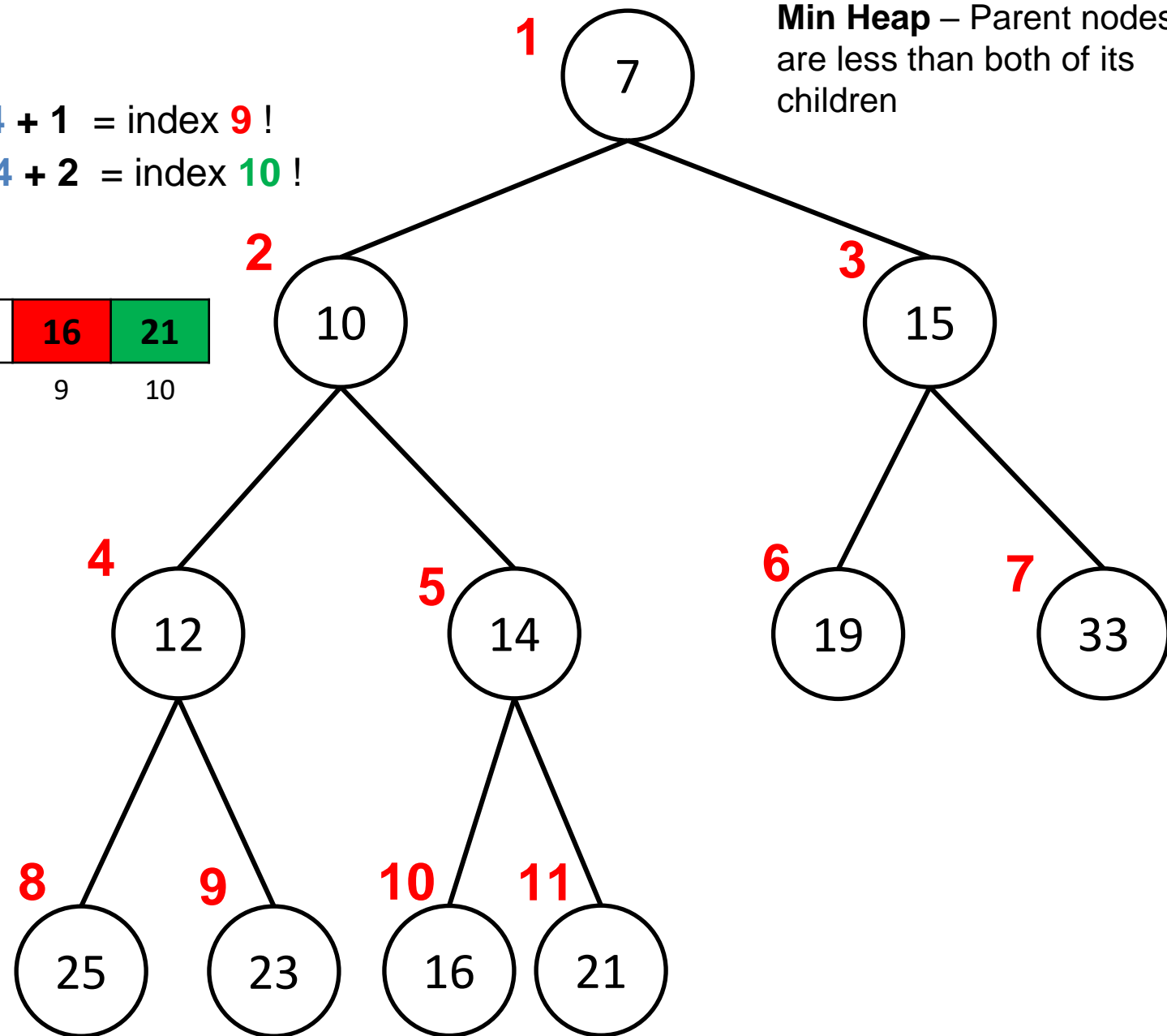
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Min Heap – Parent nodes are less than both of its children



Heap Representation

Left Child = $2 * 0 + 1 = \text{index } 1!$
Right Child = $2 * 0 + 2 = \text{index } 2!$

Array

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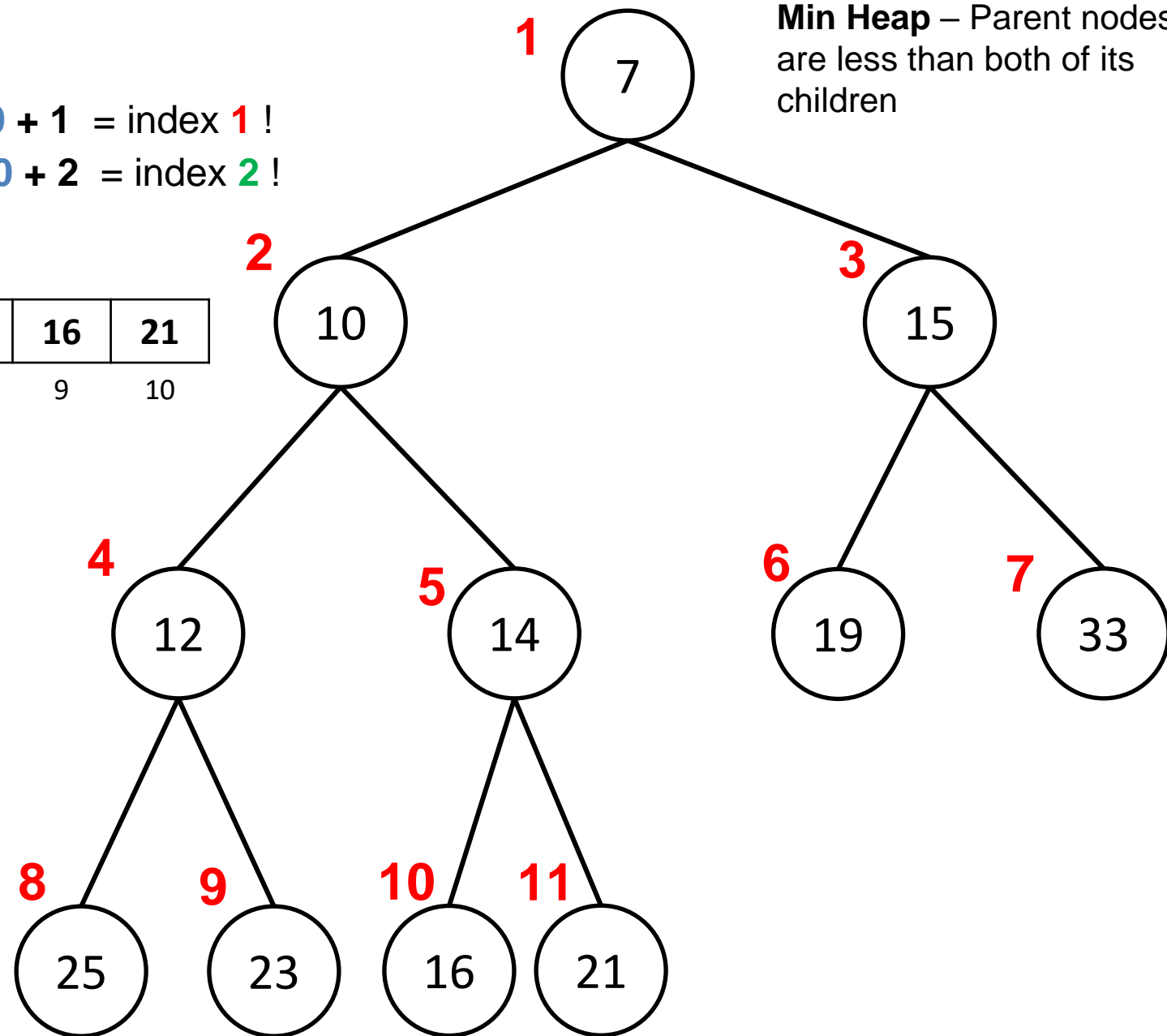
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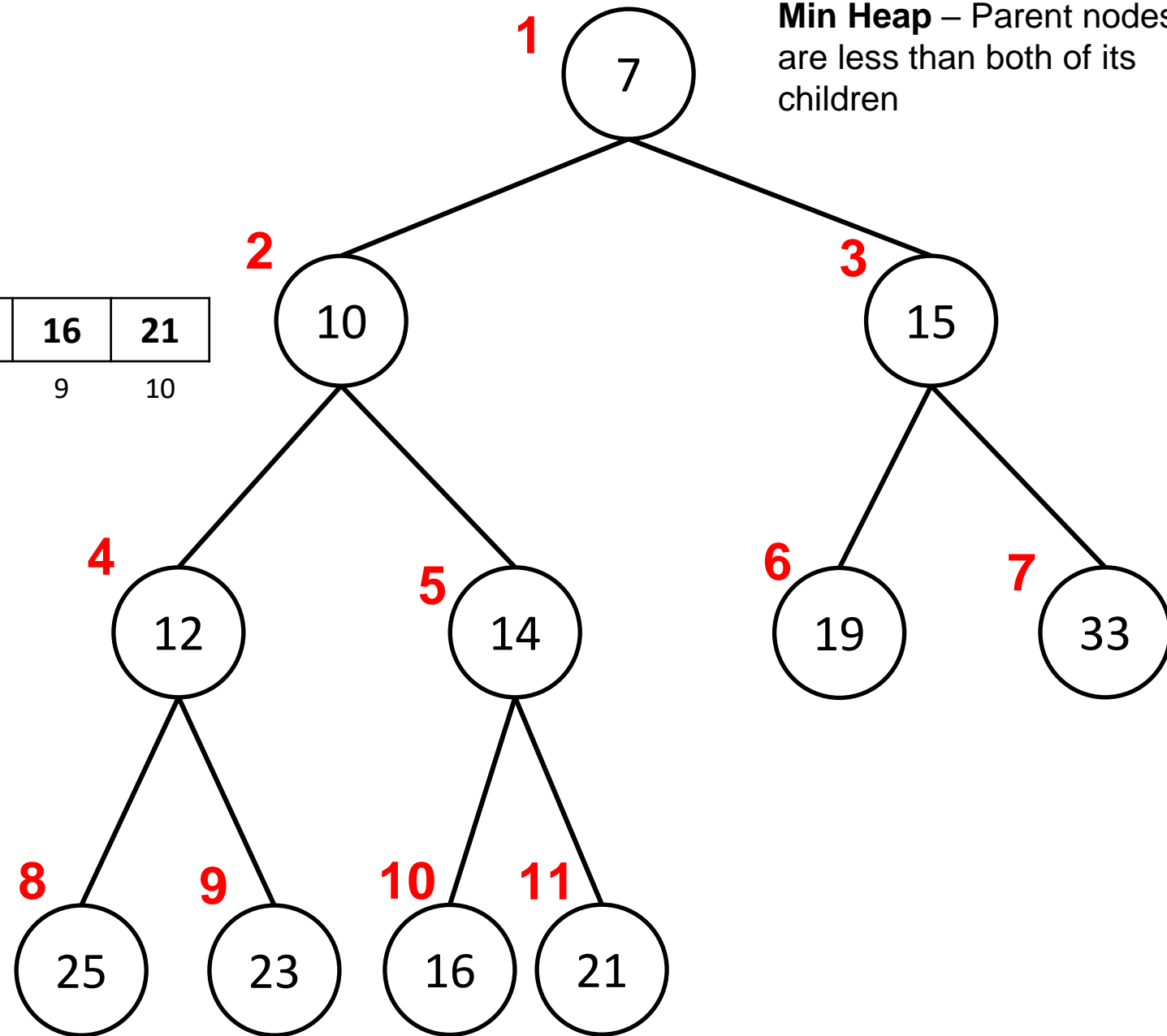
Heap Representation

Min Heap – Parent nodes are less than both of its children

Array

7	10	15	12	14	19	33	25	23	16	21
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Given a spot in the array, how can we find its parent?



Heap Representation

Array

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0	1	2	3	4	5	6	7	8	9	10

Given a spot in the array, how can we find its parent?

Because this is a complete binary tree, there is a pretty nifty formula for this

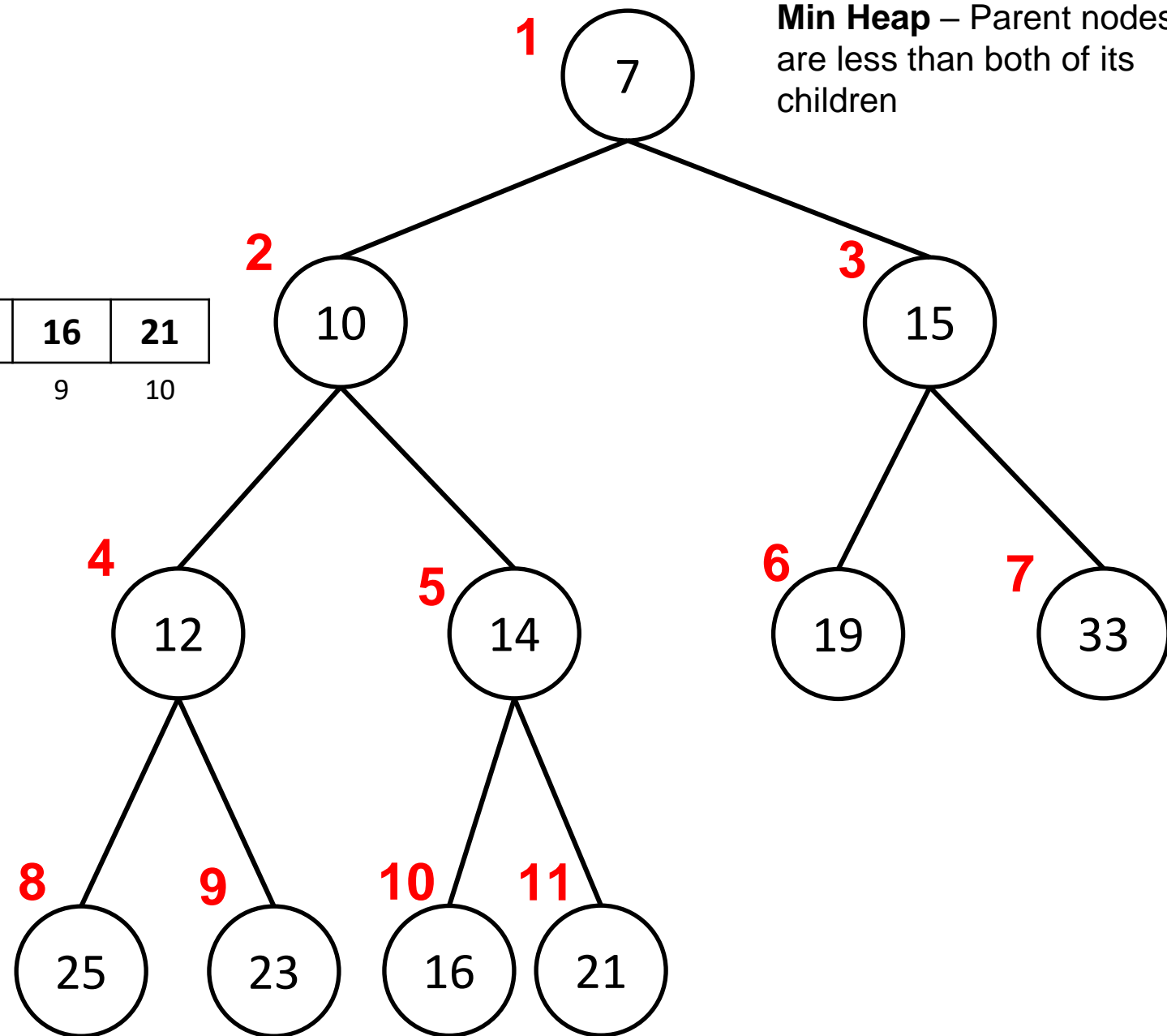
Given an index i

Its parent will be located at index:

$$(i - 1) / 2$$

(remember that the / operator will **floor** the answer)

Min Heap – Parent nodes are less than both of its children



Heap Representation

Min Heap – Parent nodes are less than both of its children

$$\text{Parent} = (6 - 1) / 2 = \text{Index } 2$$

Array

7	10	15	12	14	19	33	25	23	16	21
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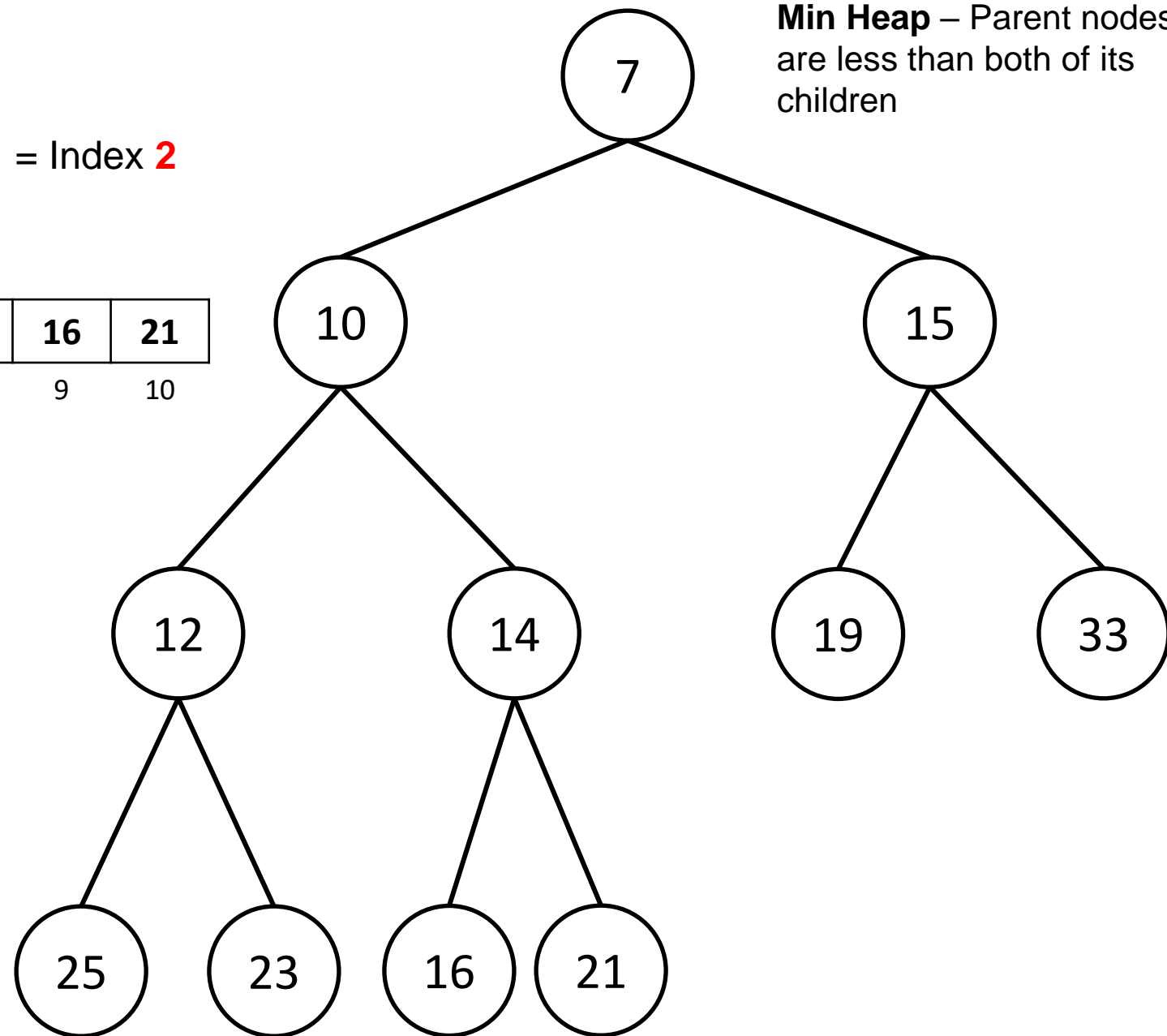
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Heap Representation

Min Heap – Parent nodes are less than both of its children

$$\text{Parent} = (3 - 1) / 2 = \text{Index } 1$$

Array

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0	1	2	3	4	5	6	7	8	9	10

Given a spot in the array, how can we find its parent?

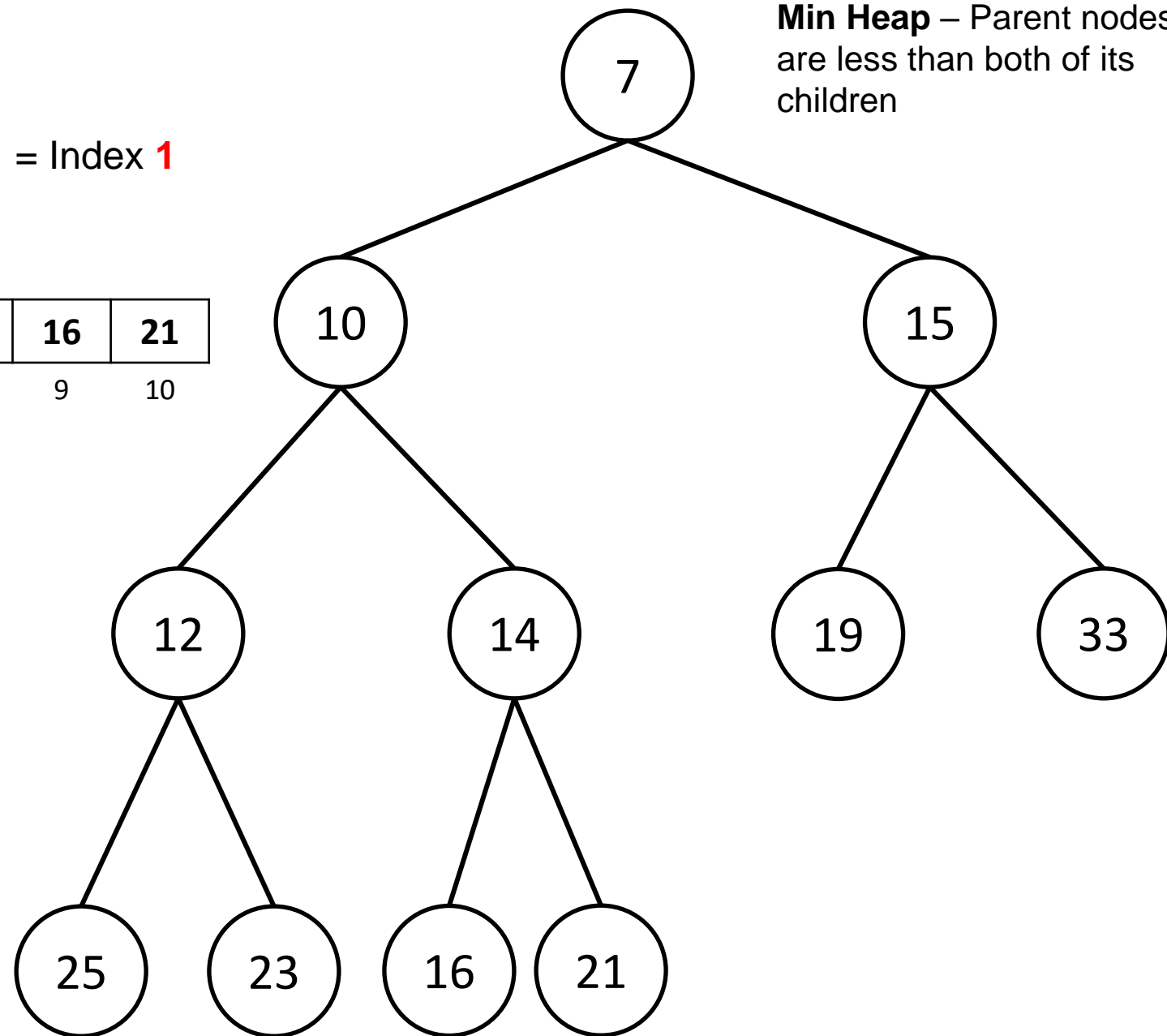
Because this is a complete binary tree, there is a pretty nifty formula for this

Given an index i

Its parent will be located at index:

$$(i - 1) / 2$$

(remember that the / operator will **floor** the answer)



Heap Representation

Min Heap – Parent nodes are less than both of its children

Array

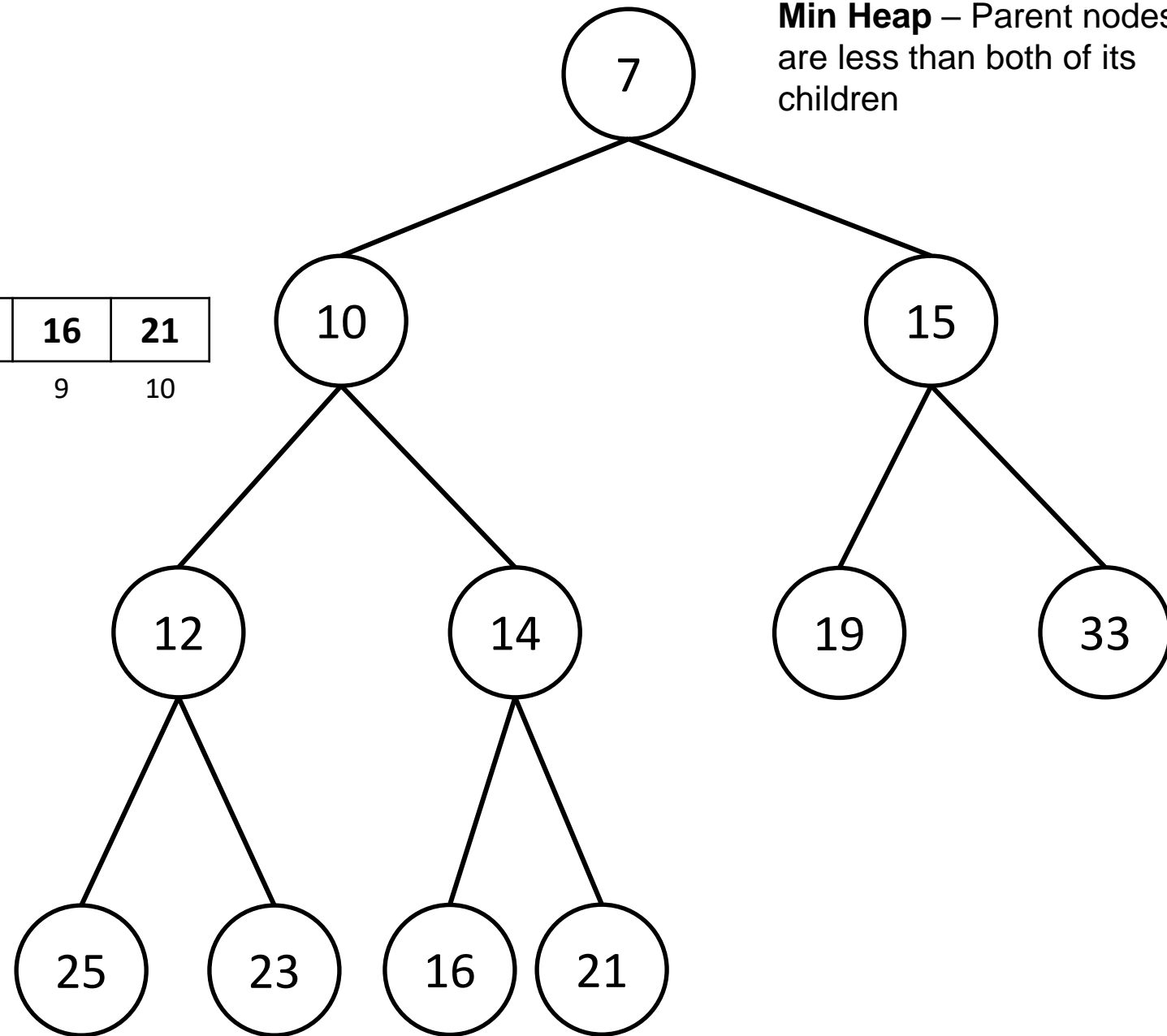
7	10	15	12	14	19	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

We can represent our tree with an array!
We have formulas to find the left child, right child, and parent for a given node

Left Child $2 * i + 1$

Right Child $2 * i + 2$

Parent $(i - 1) / 2$



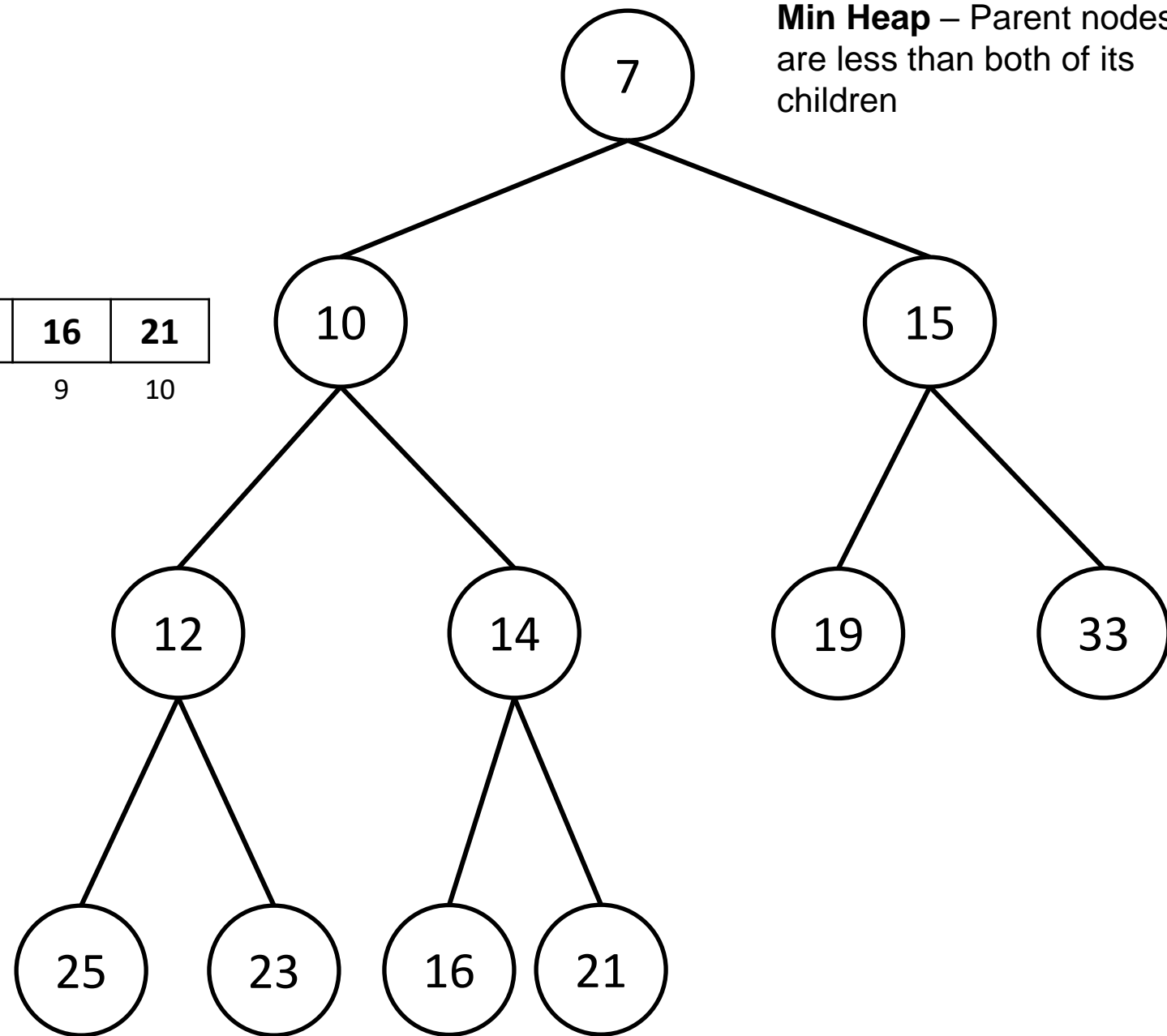
Heap Representation

Min Heap – Parent nodes are less than both of its children

Array

7	10	15	12	14	19	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

`insert(11);`



Heap Representation

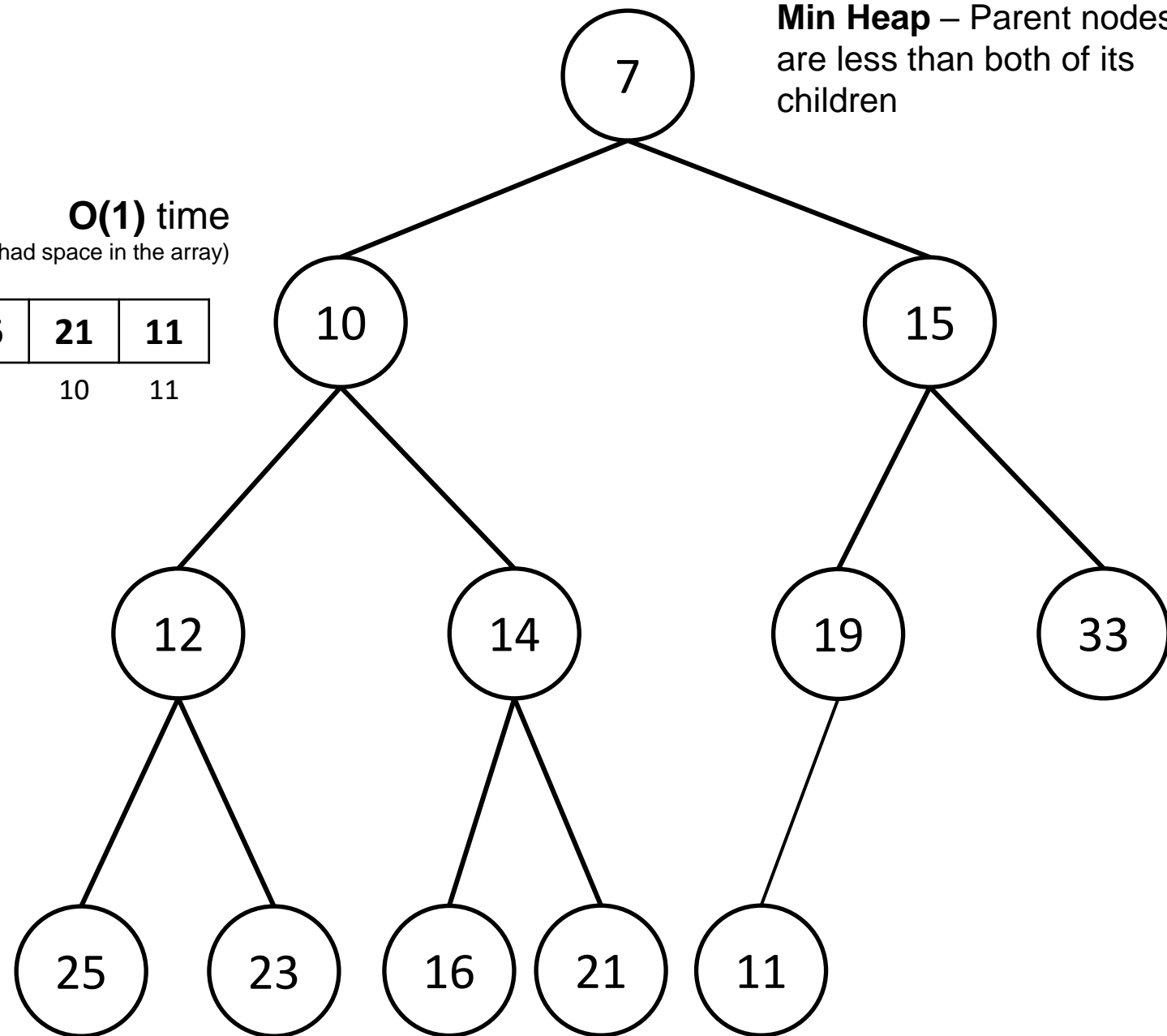
Min Heap – Parent nodes are less than both of its children

Array

$O(1)$ time
(assuming we had space in the array)

7	10	15	12	14	19	33	25	23	16	21	11
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`



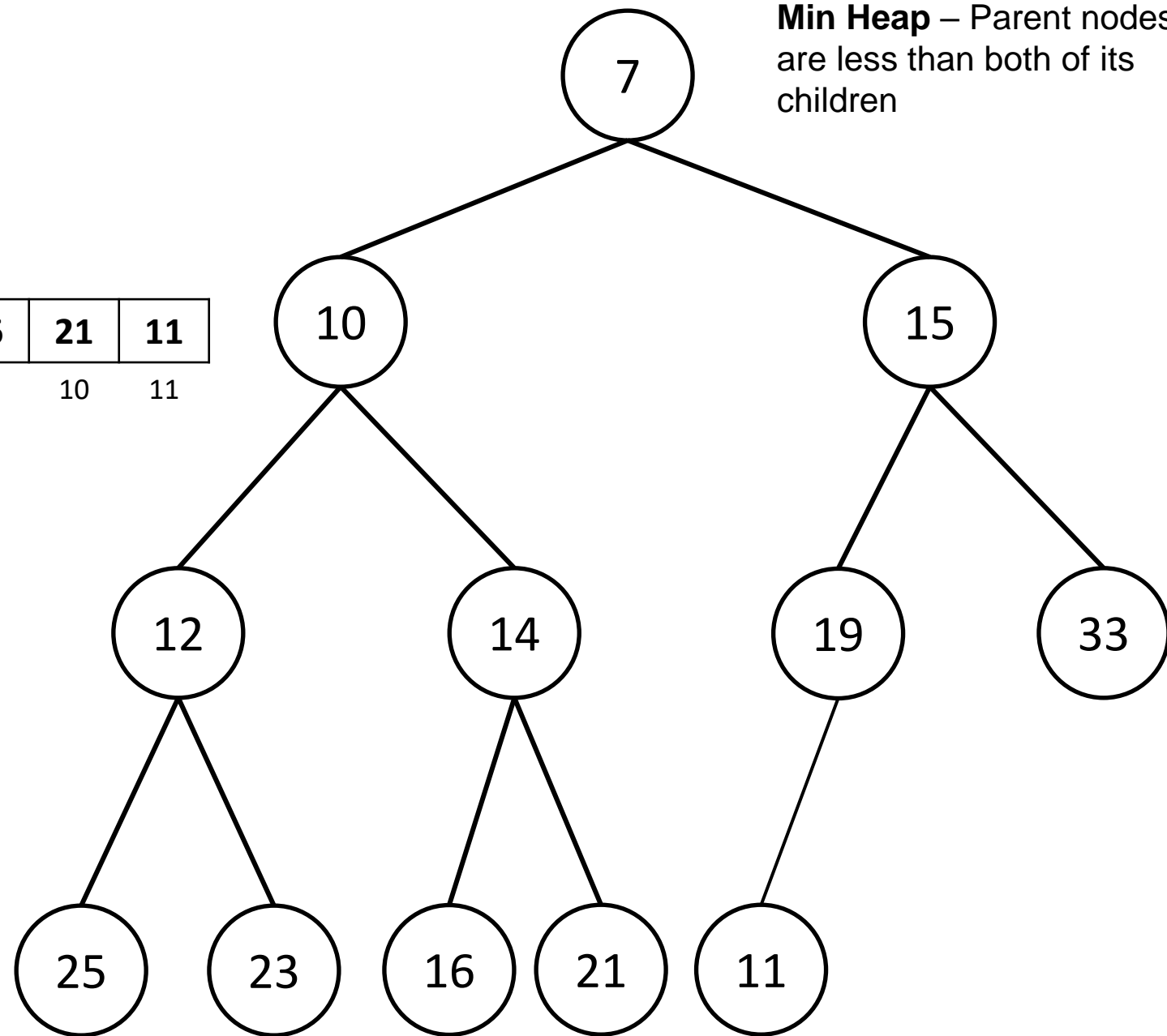
Heap Representation

Min Heap – Parent nodes are less than both of its children

Array

7	10	15	12	14	19	33	25	23	16	21	11
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`
Time to Heapify Up!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

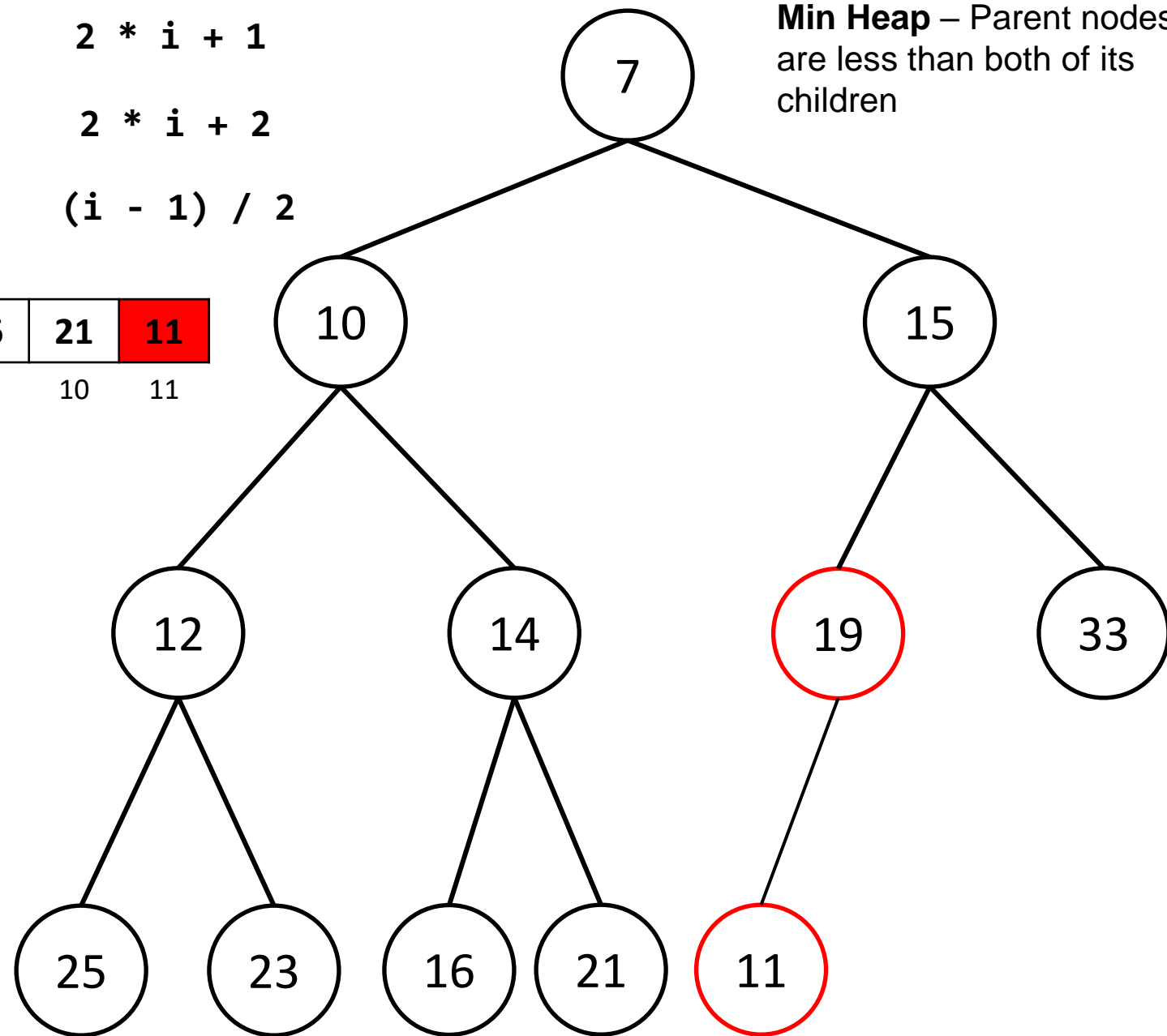
Array

7	10	15	12	14	19	33	25	23	16	21	11
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`

Time to Heapify Up!

11's parent is located at $(11 - 1) / 2 = 5$



Heap Representation

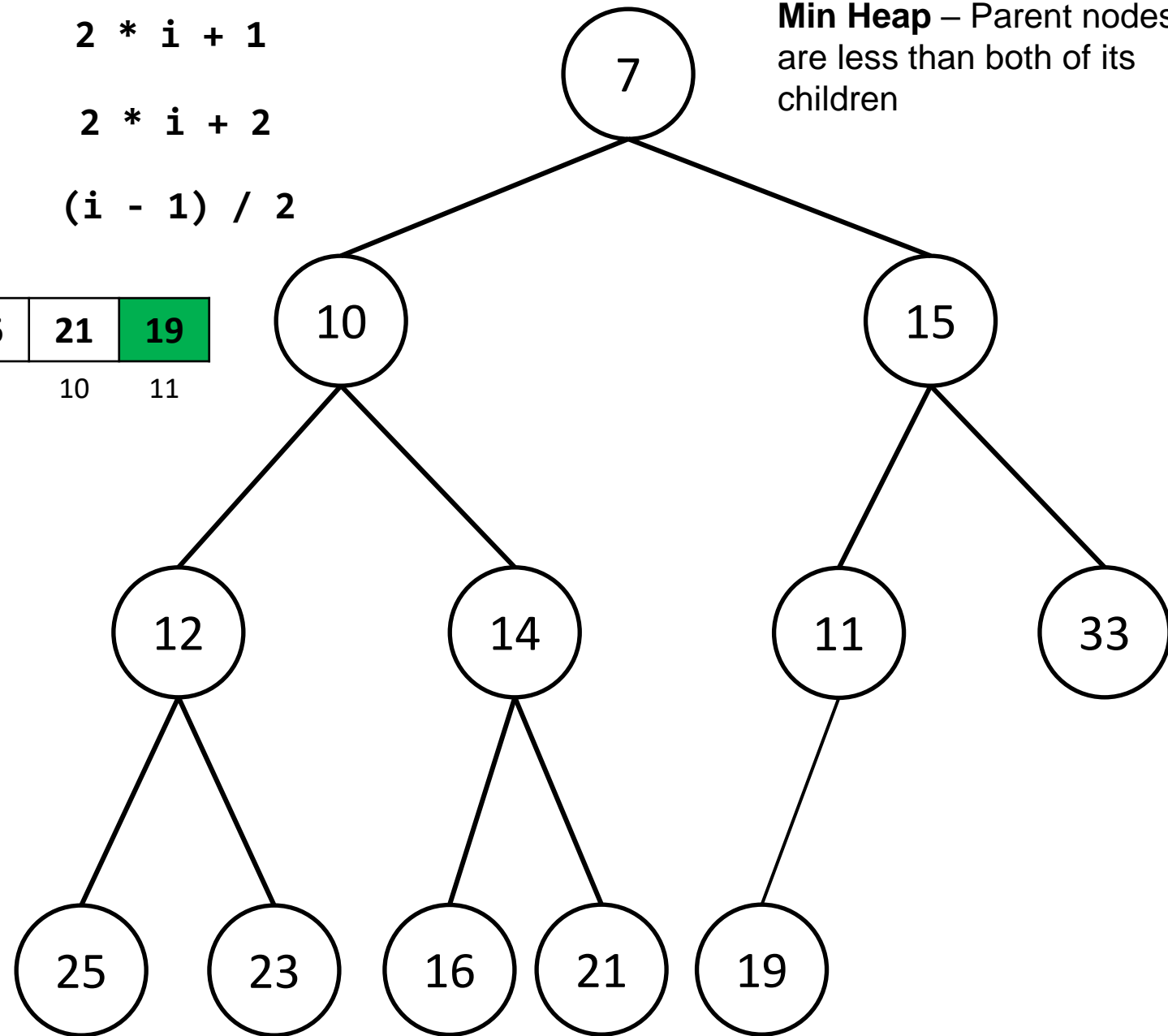
Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

Array

7	10	15	12	14	11	33	25	23	16	21	19
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`
Time to Heapify Up!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

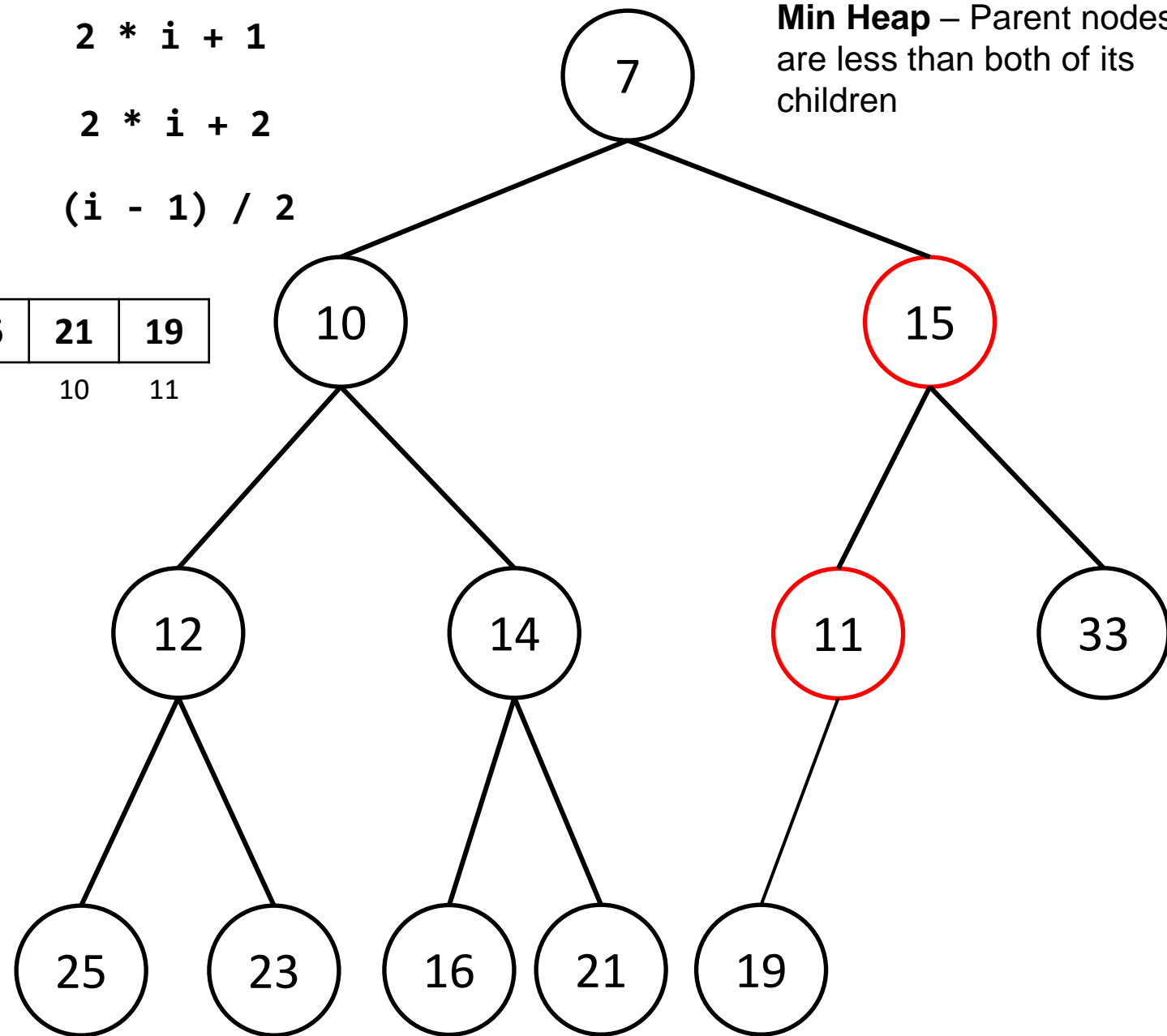
Array

7	10	15	12	14	11	33	25	23	16	21	19
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`

Time to Heapify Up!

11's parent is located at $(5 - 1) / 2 = 2$



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

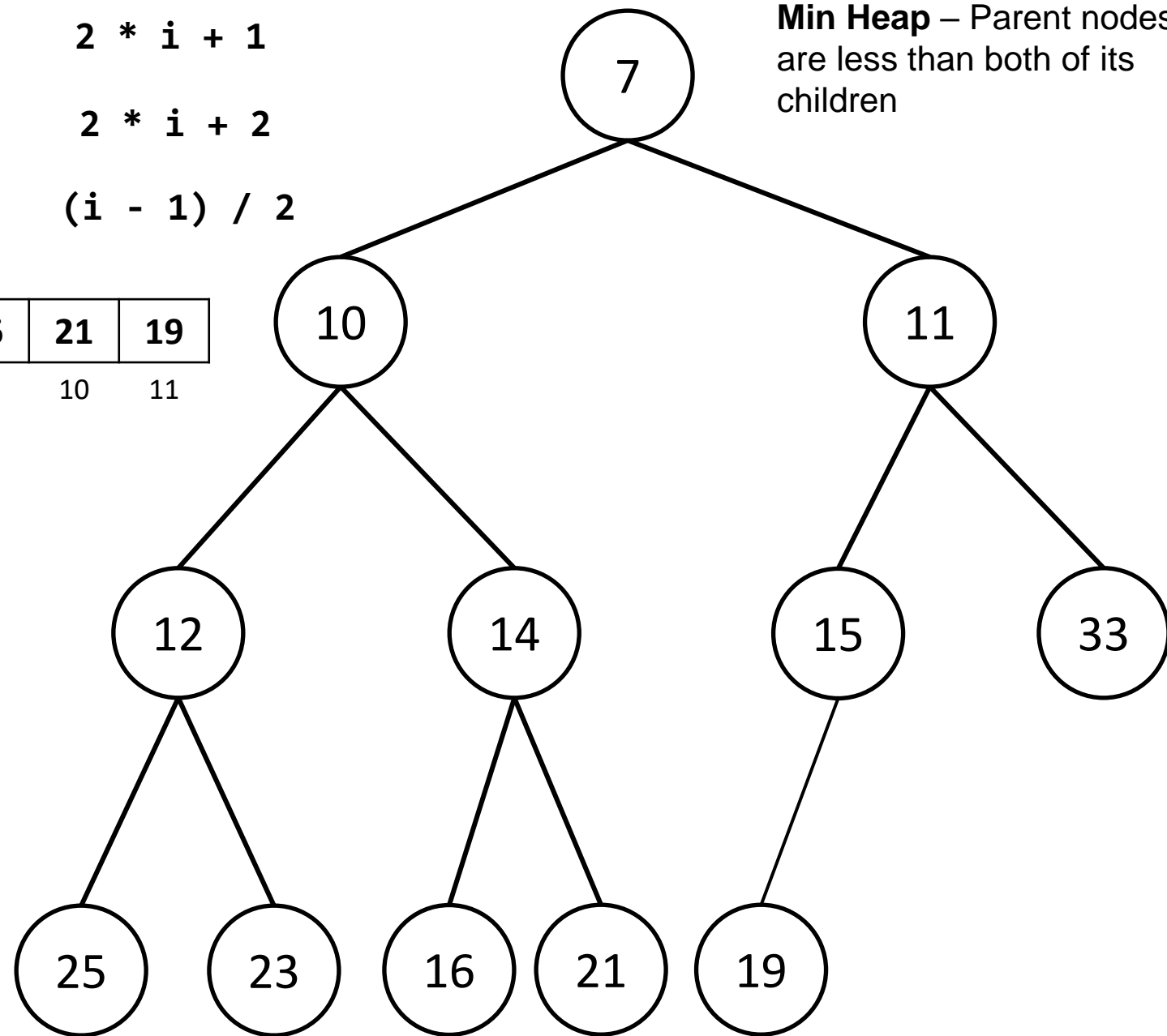
Array

7	10	11	12	14	15	33	25	23	16	21	19
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`

Time to Heapify Up!

11's parent is located at $(5 - 1) / 2 = 2$



Heap Representation

Left Child

$$2 * i + 1$$

Right Child

$$2 * i + 2$$

Parent

$$(i - 1) / 2$$

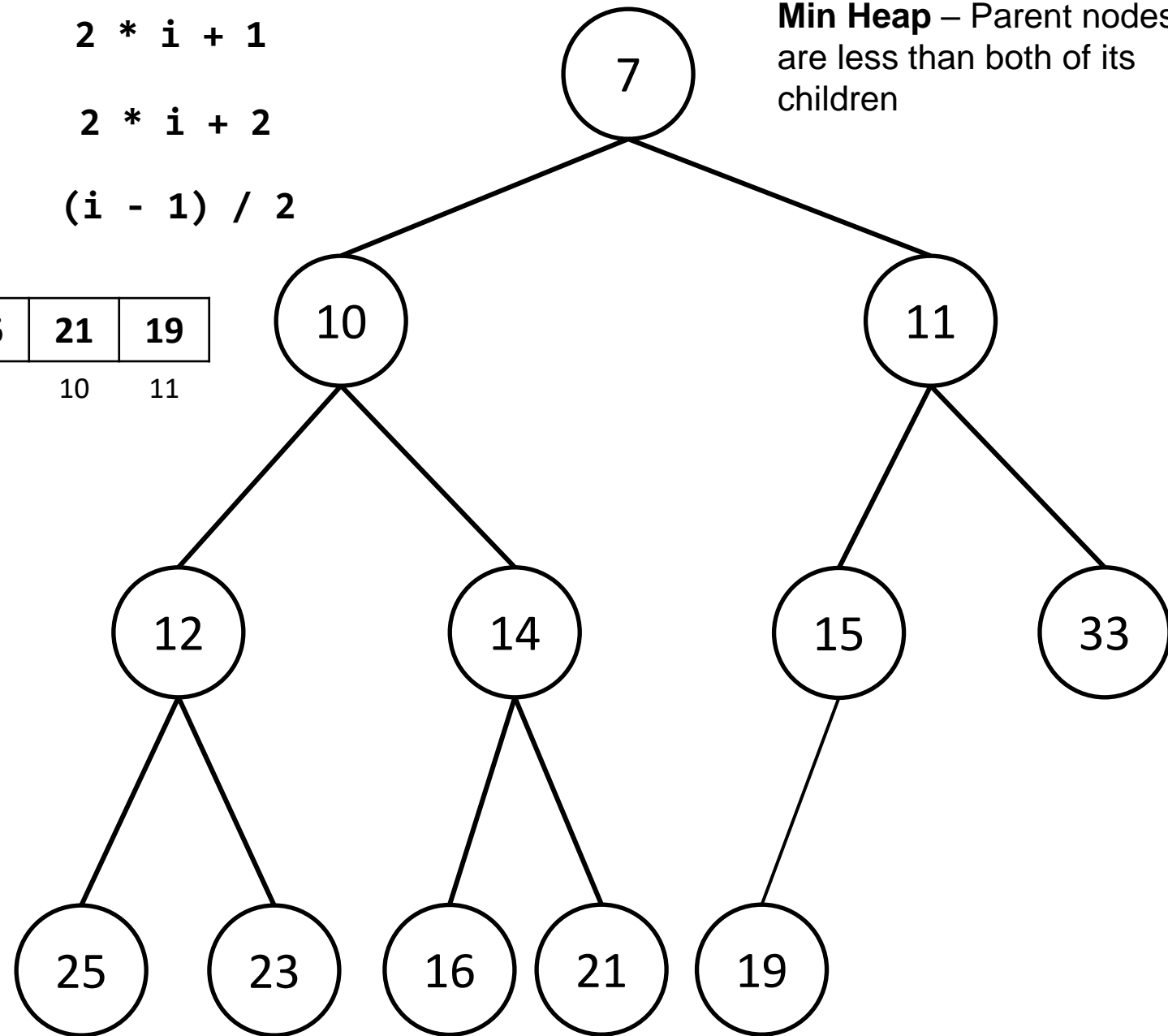
Min Heap – Parent nodes are less than both of its children

Array

7	10	11	12	14	15	33	25	23	16	21	19
0	1	2	3	4	5	6	7	8	9	10	11

`insert(11);`

Time to Heapify Up!



Heap Representation

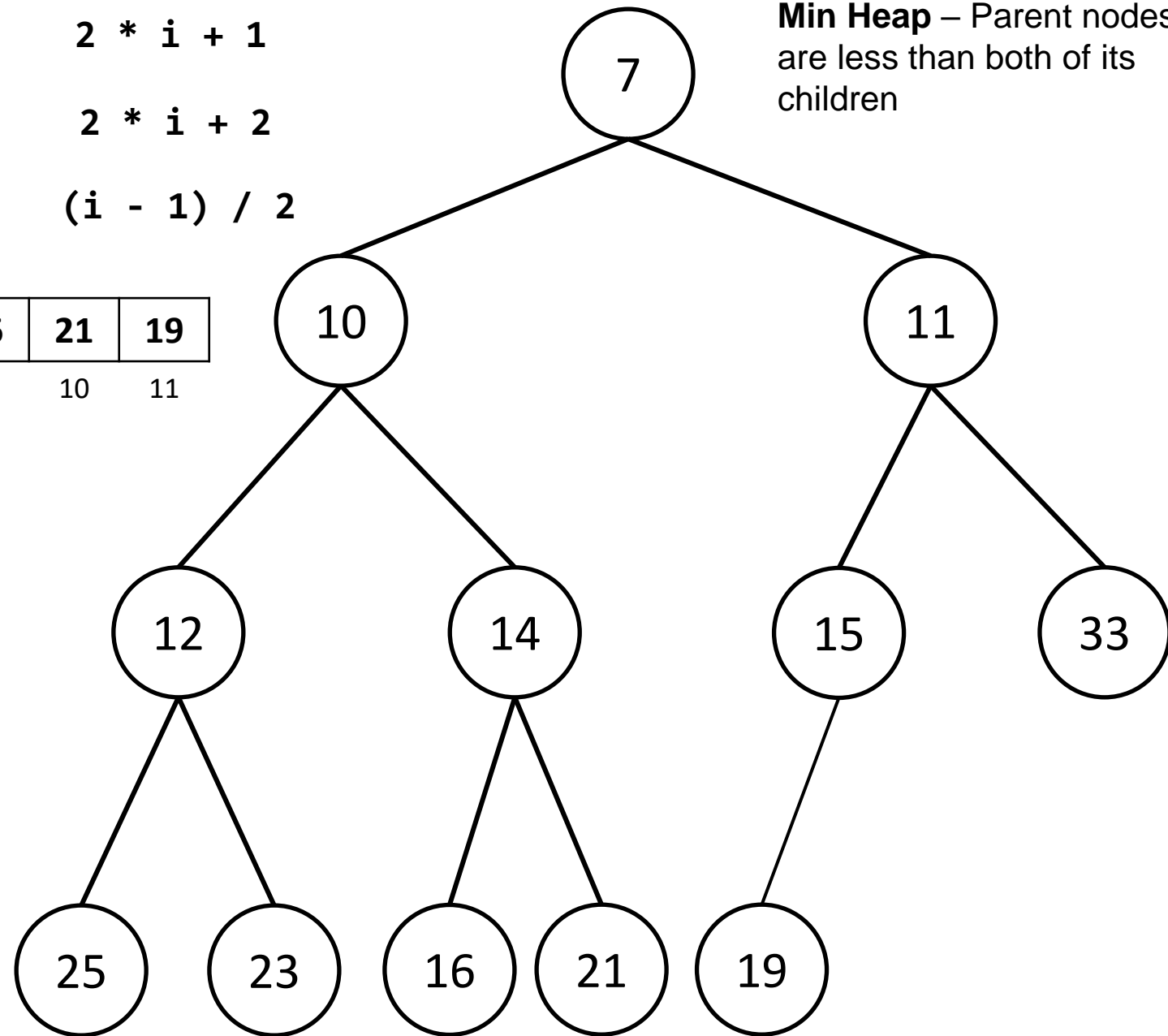
Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

Array

7	10	11	12	14	15	33	25	23	16	21	19
0	1	2	3	4	5	6	7	8	9	10	11

`poll();`



Heap Representation

Left Child $2 * i + 1$

Right Child $2 * i + 2$

Parent $(i - 1) / 2$

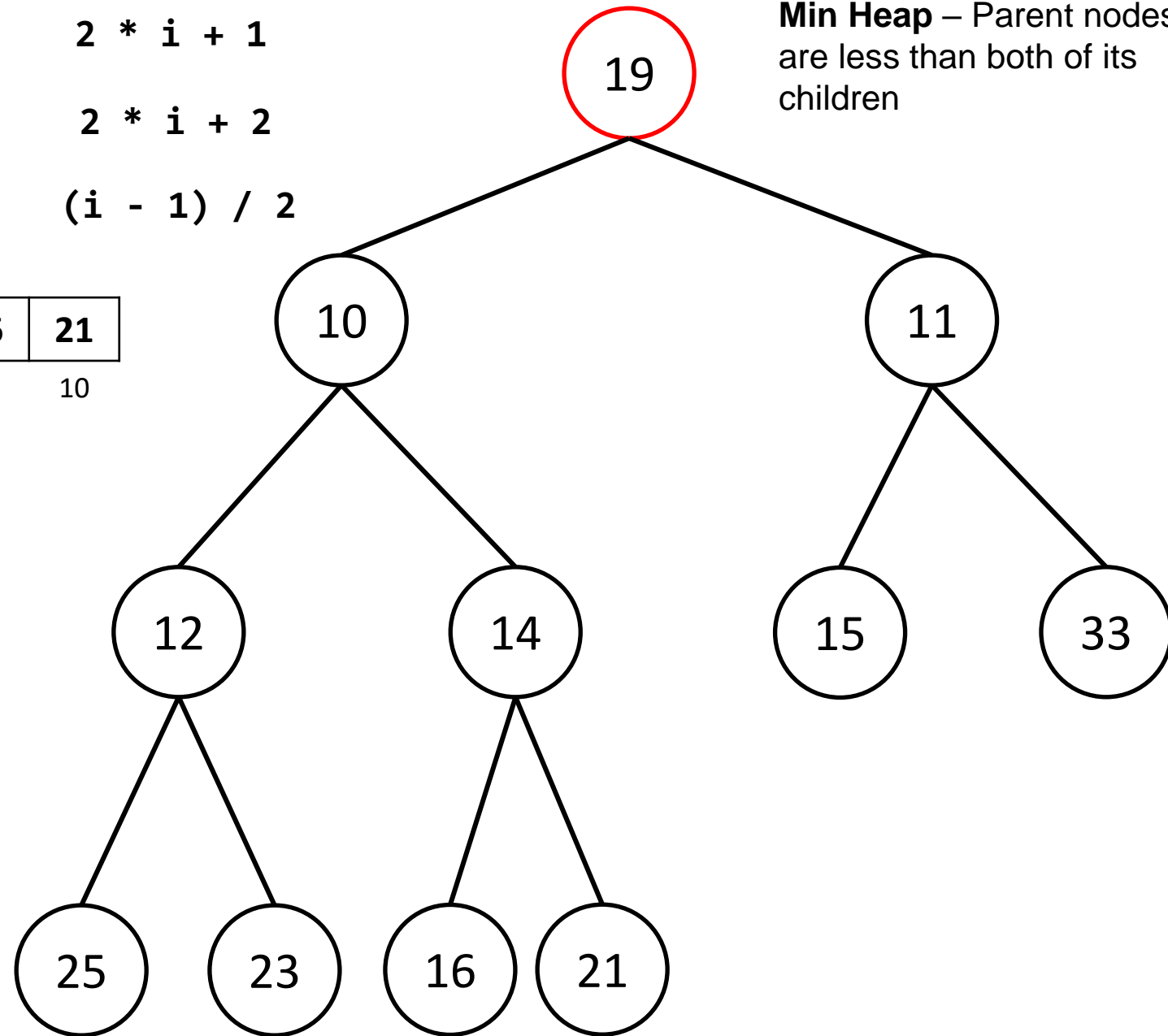
Min Heap – Parent nodes are less than both of its children

Array

$O(1)$ time

19	10	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

poll();



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

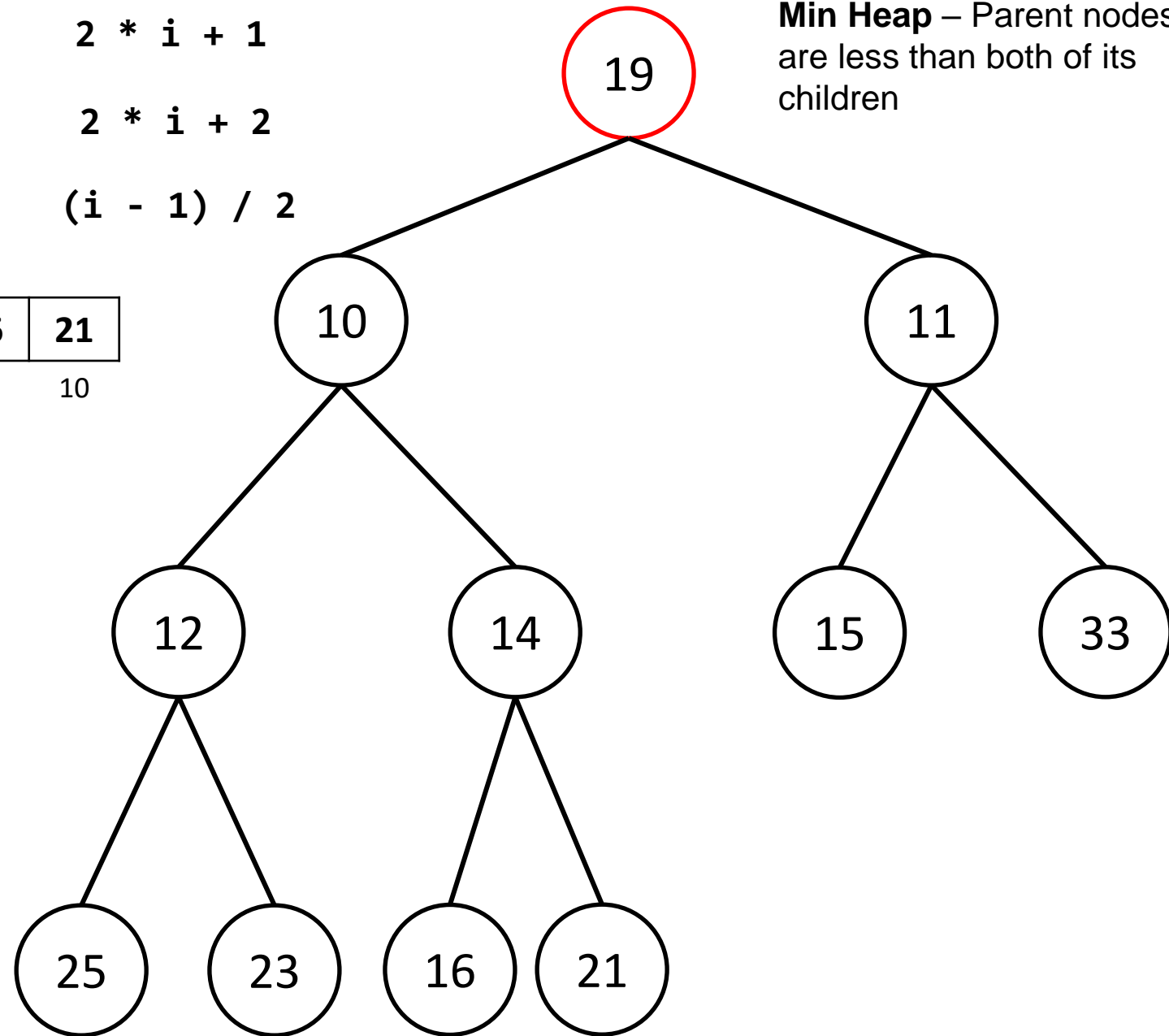
Min Heap – Parent nodes are less than both of its children

Array

19	10	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

`poll();`

Time to Heapify down!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

Array

19	10	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

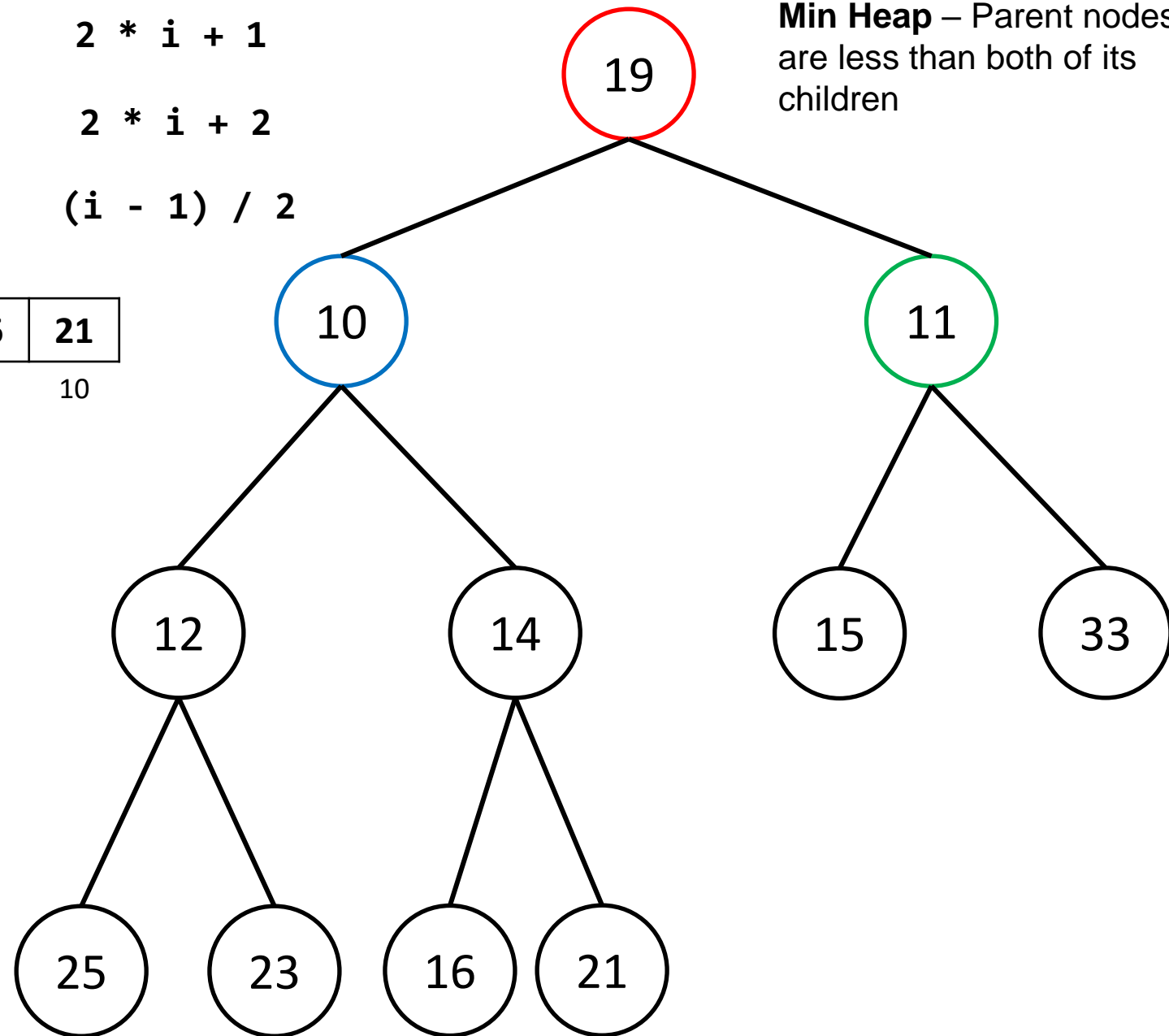
`poll();`

Time to Heapify down!

19's left child is located at $2 * 0 + 1 = 1$

19's right child is located at $2 * 0 + 2 = 2$

(We want to swap it with the lower value)



Heap Representation

Left Child

$$2 * i + 1$$

Right Child

$$2 * i + 2$$

Parent

$$(i - 1) / 2$$

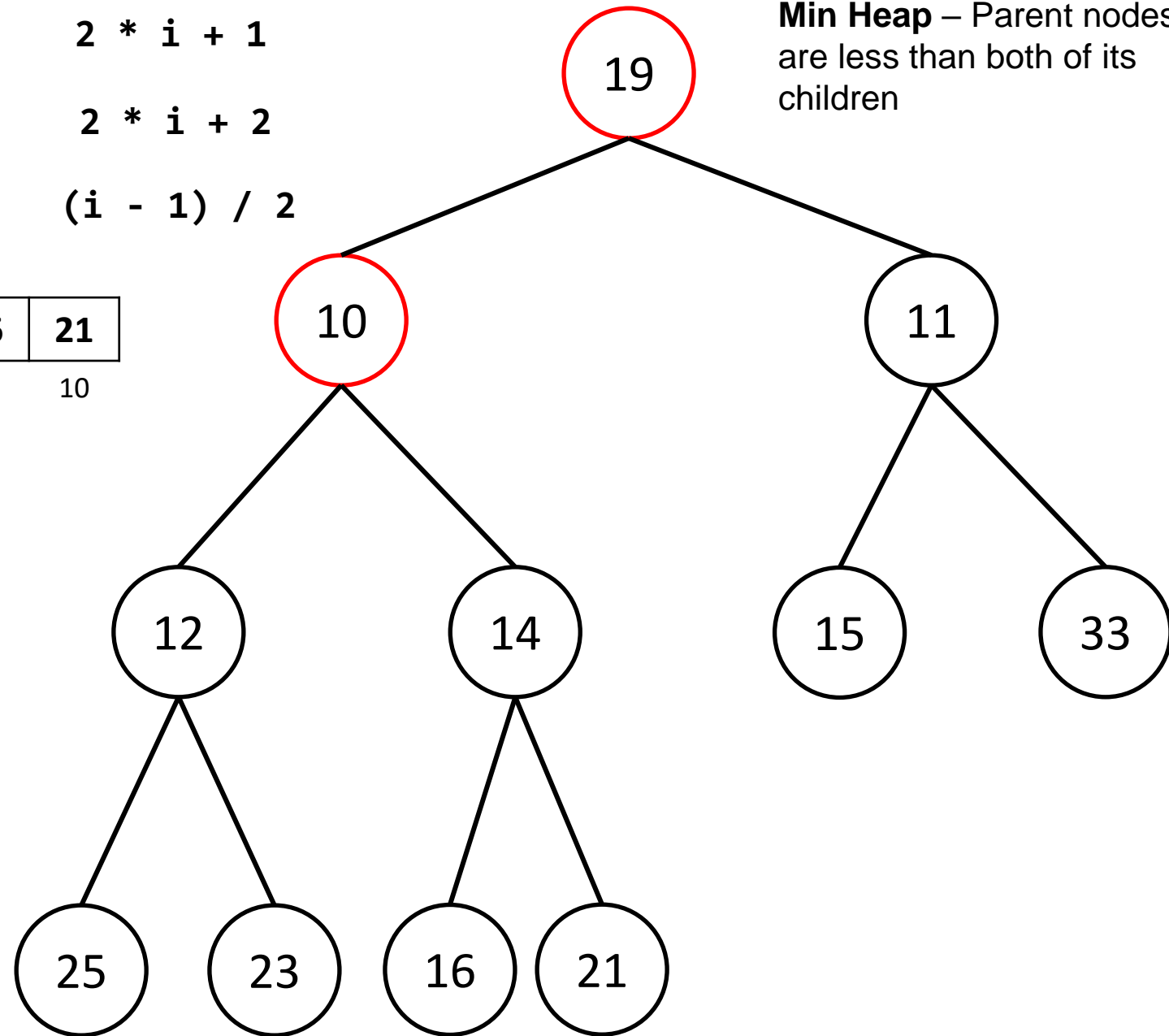
Min Heap – Parent nodes are less than both of its children

Array

19	10	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

poll();

Time to Heapify down!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

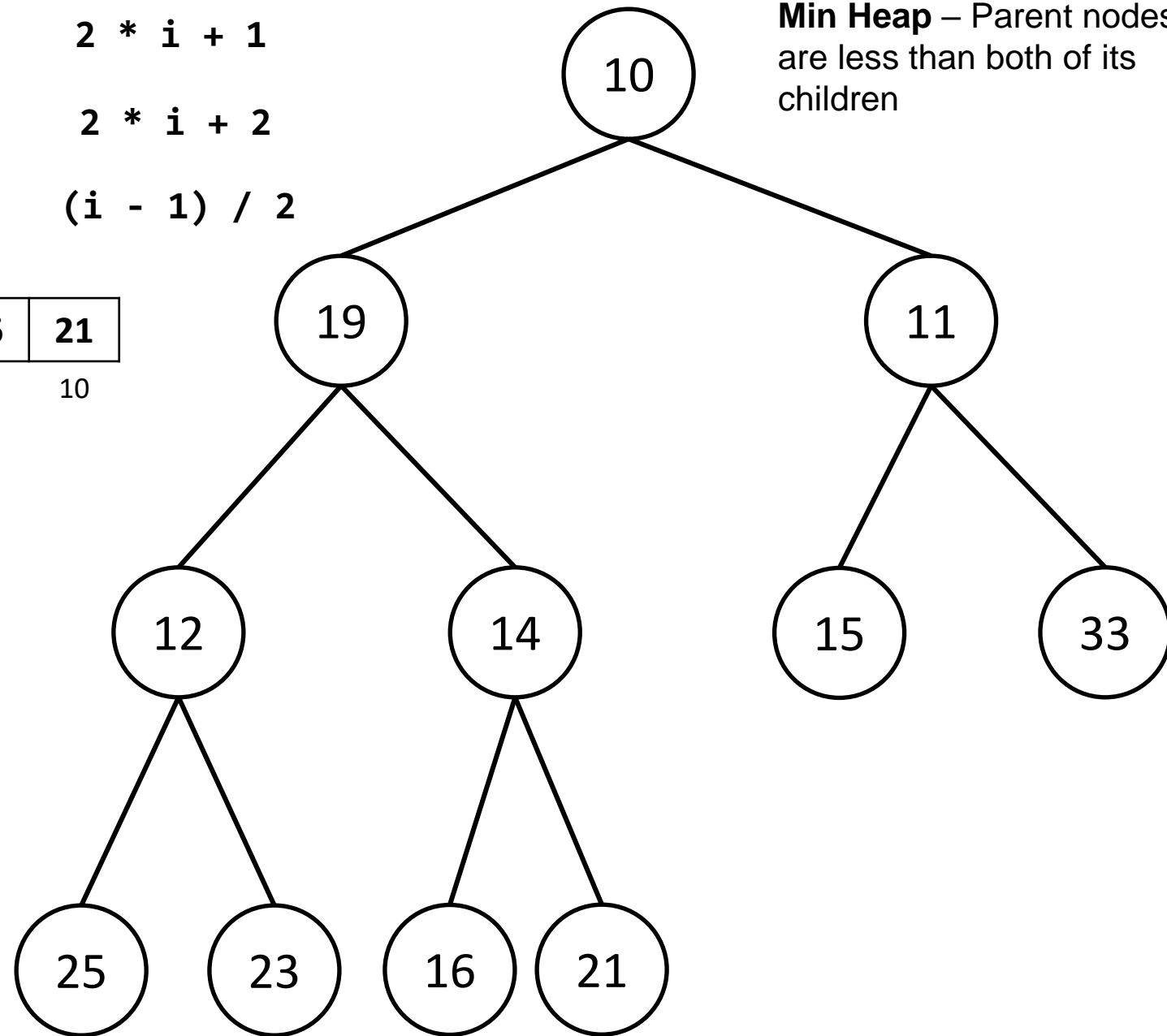
Min Heap – Parent nodes are less than both of its children

Array

10	19	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

`poll();`

Time to Heapify down!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

Array

10	19	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

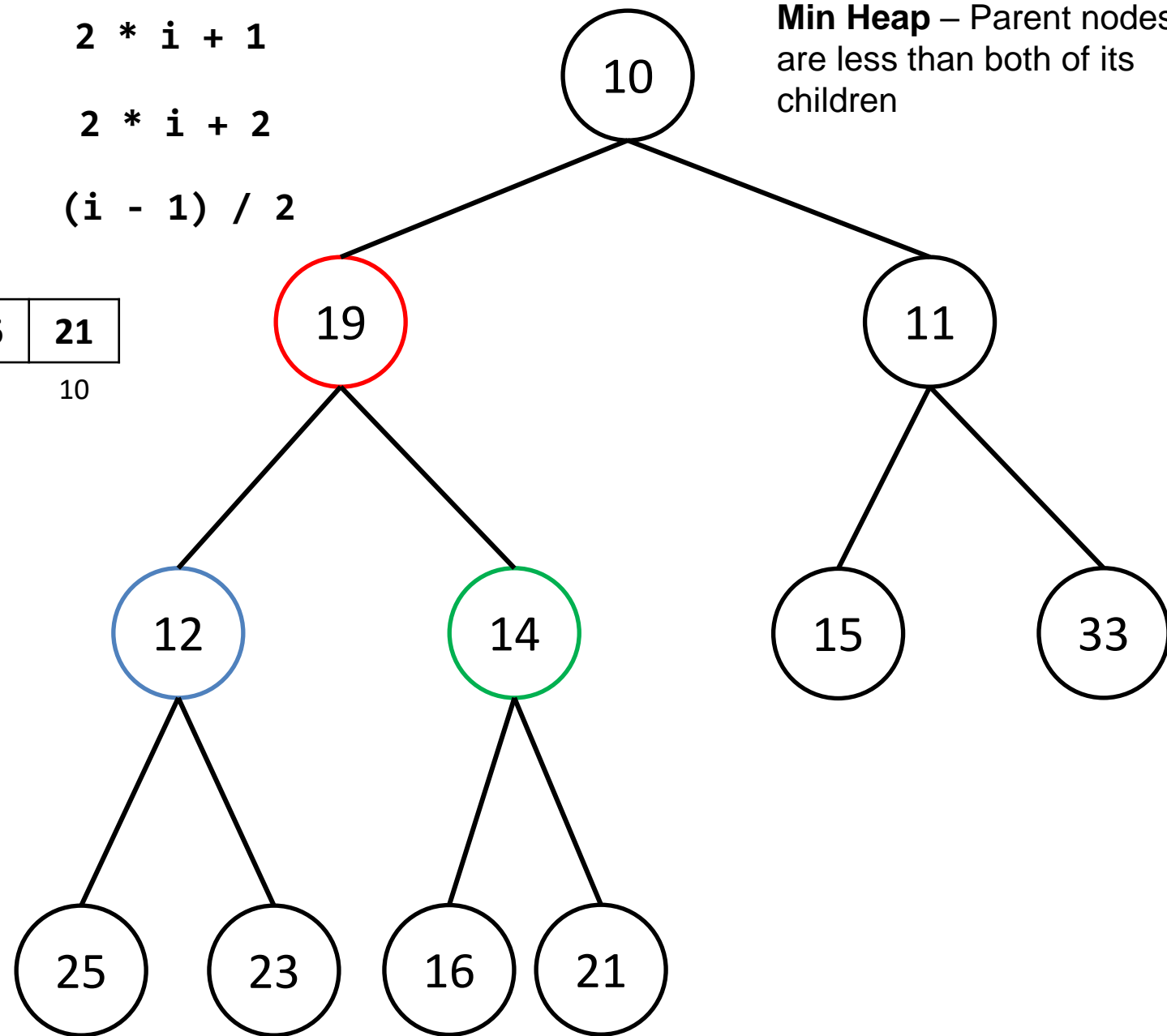
`poll();`

Time to Heapify down!

19's left child is located at $2 * 1 + 1 = 3$

19's right child is located at $2 * 1 + 2 = 4$

(We want to swap it with the lower value)



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

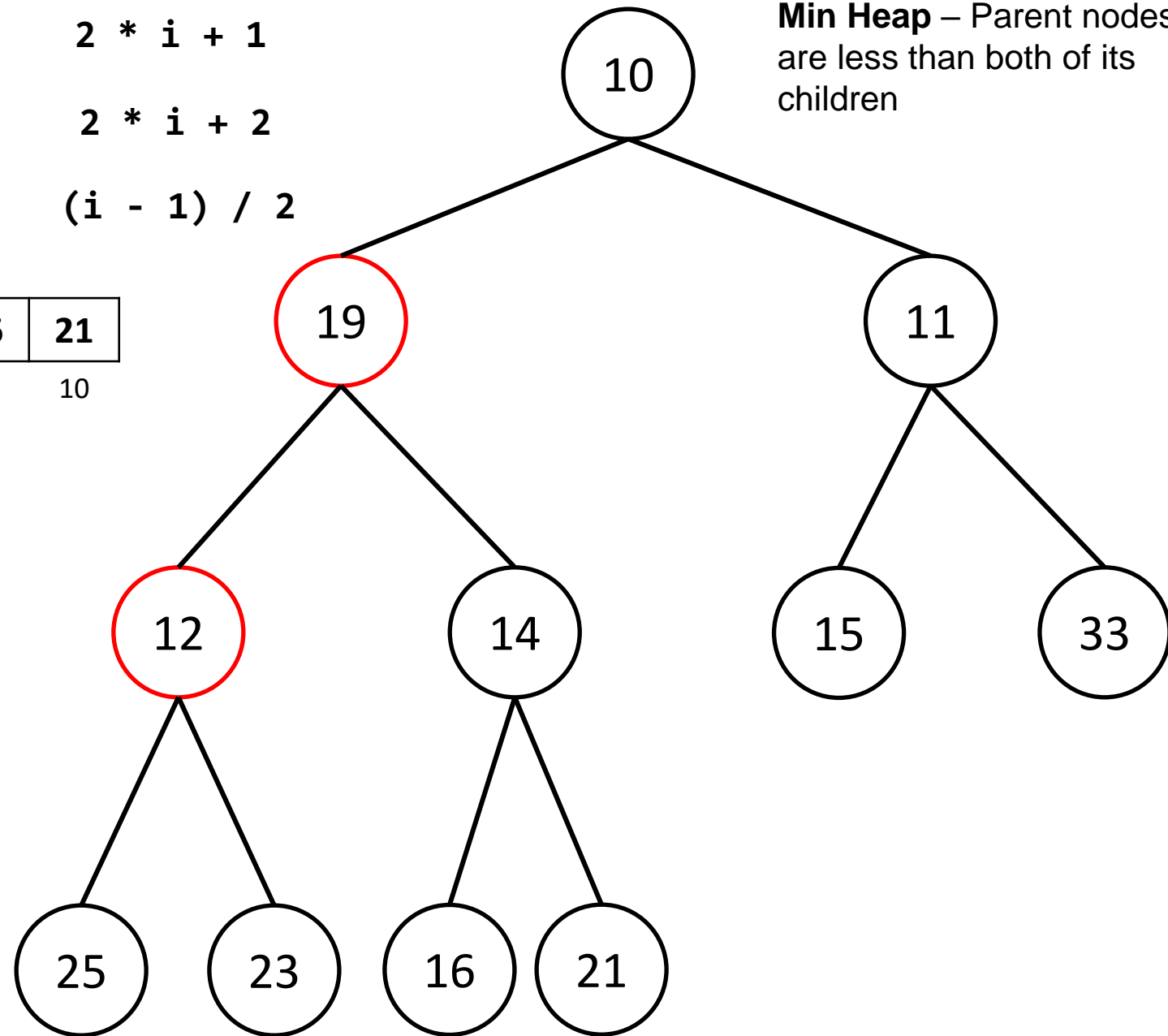
Min Heap – Parent nodes are less than both of its children

Array

10	19	11	12	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

`poll();`

Time to Heapify down!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

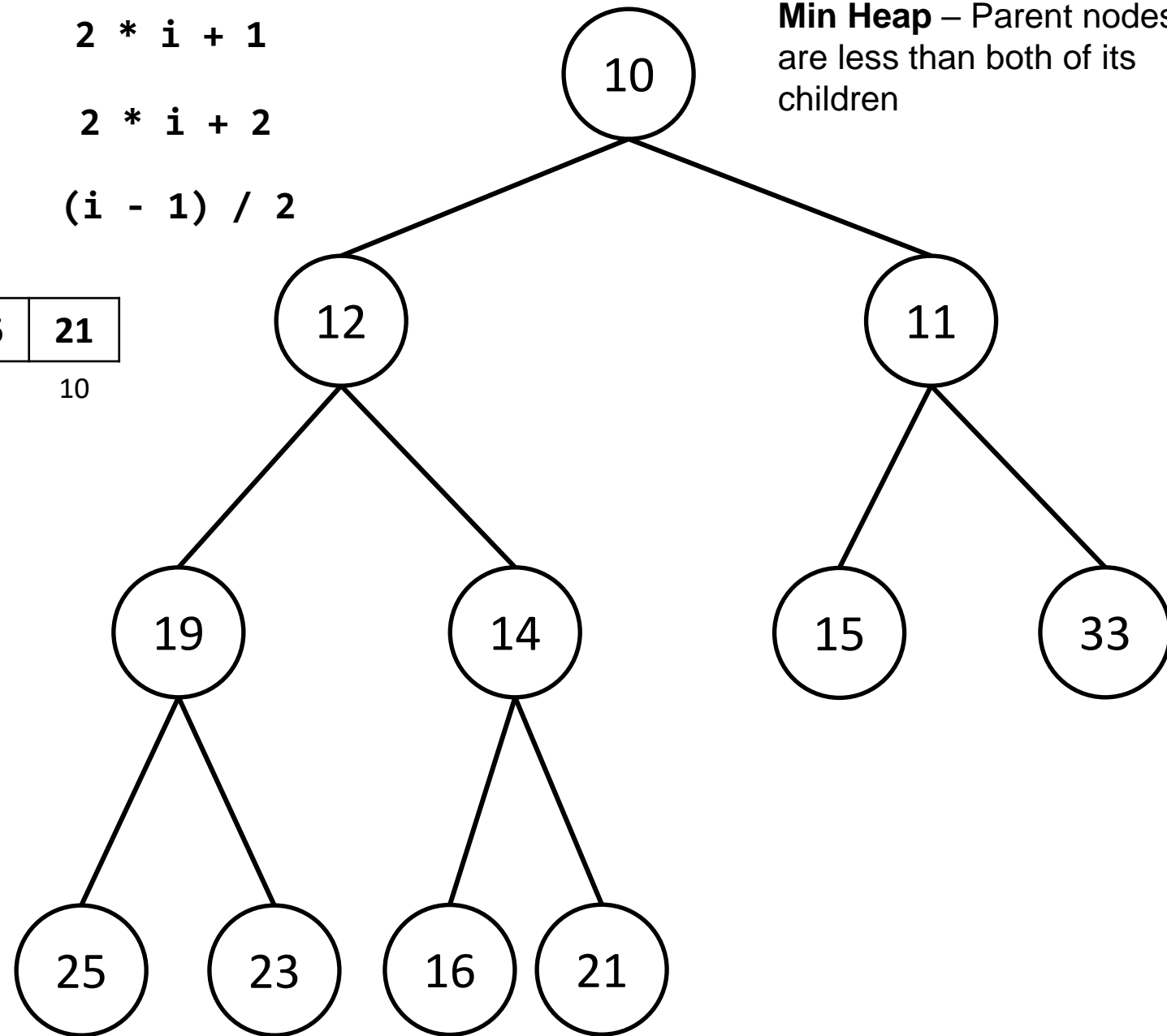
Min Heap – Parent nodes are less than both of its children

Array

10	12	11	19	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

`poll();`

Time to Heapify down!



Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

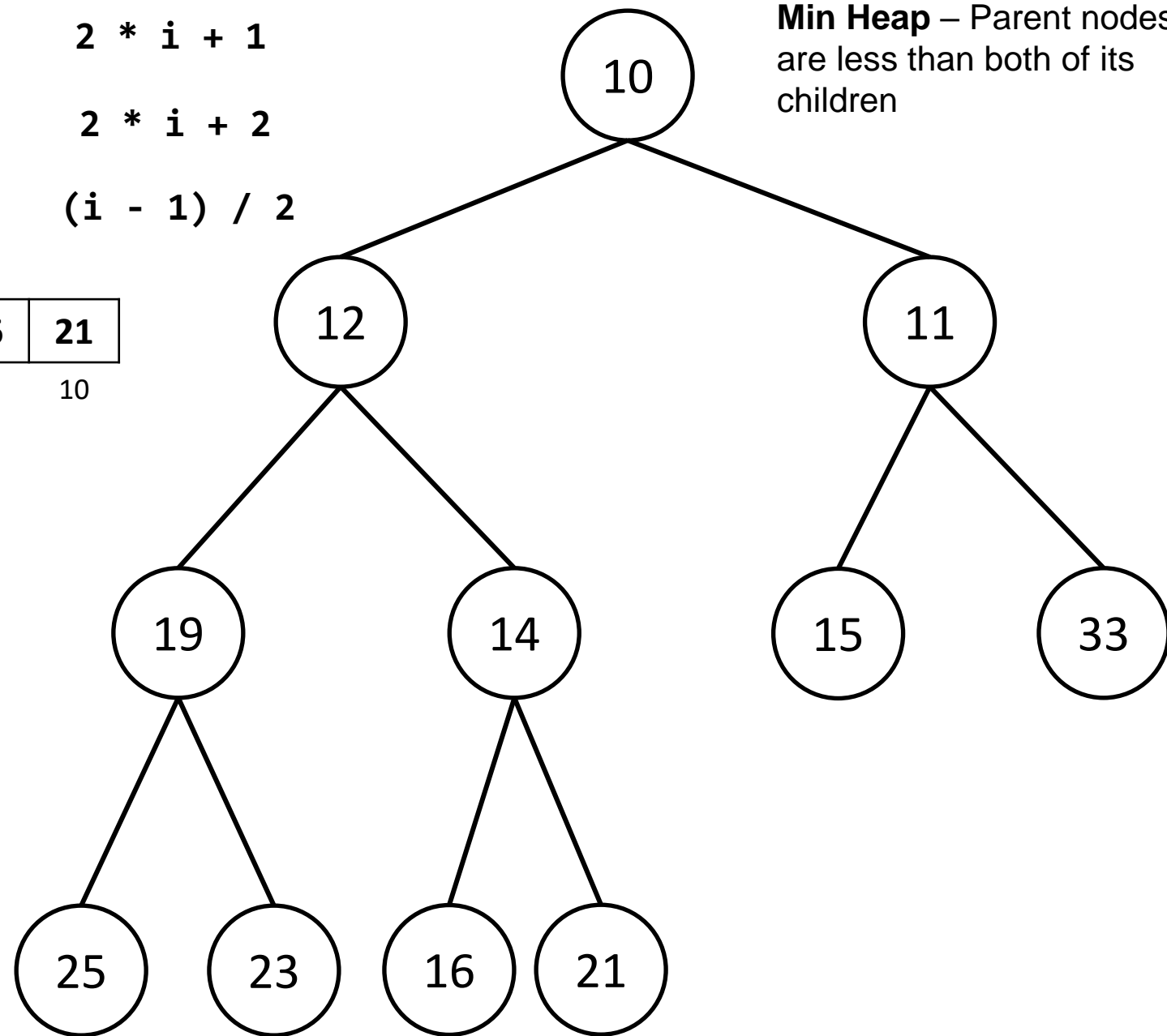
Min Heap – Parent nodes are less than both of its children

Array

10	12	11	19	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10

`poll();`

Time to Heapify down!



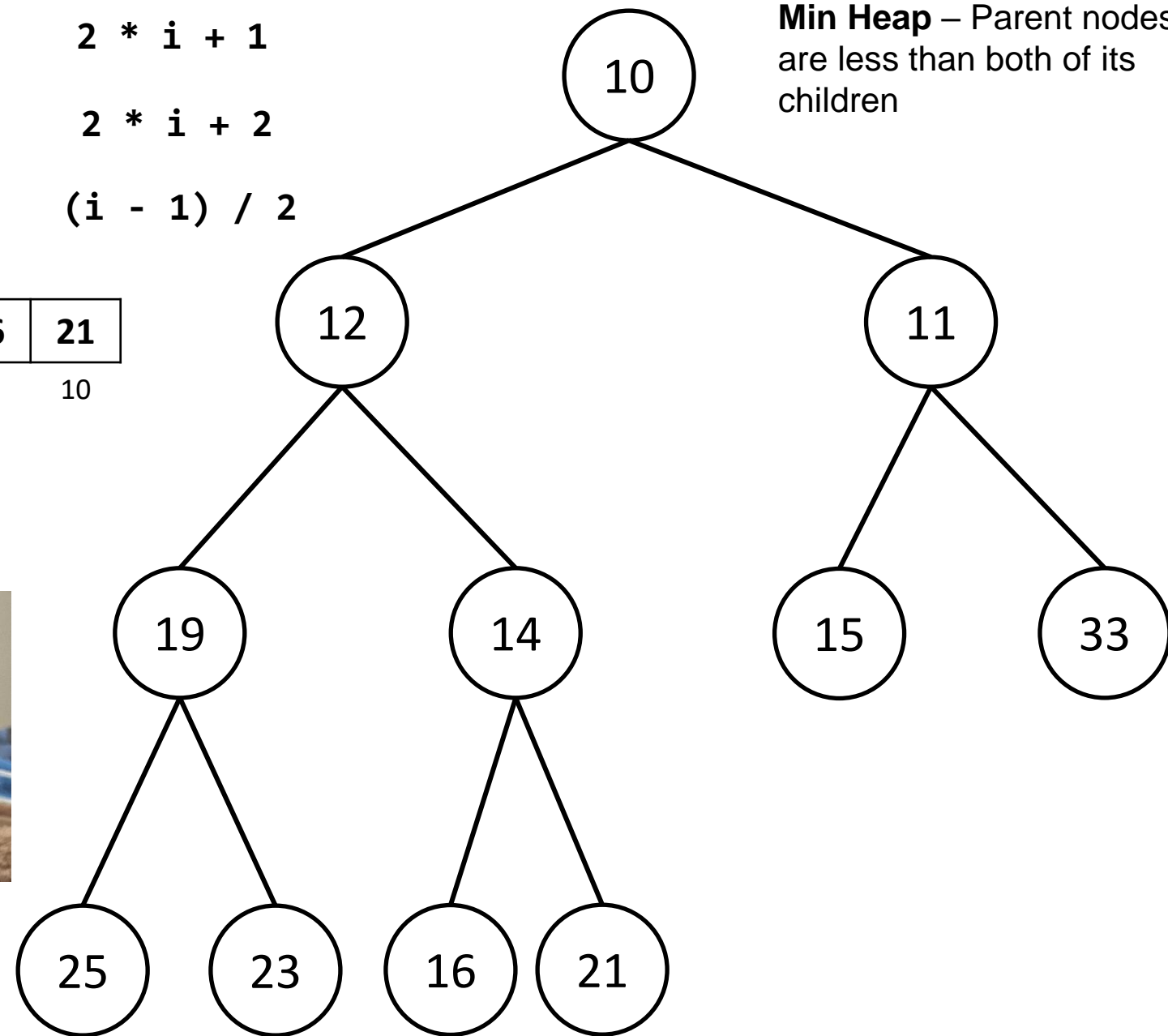
Heap Representation

Left Child $2 * i + 1$
Right Child $2 * i + 2$
Parent $(i - 1) / 2$

Min Heap – Parent nodes are less than both of its children

Array

10	12	11	19	14	15	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10



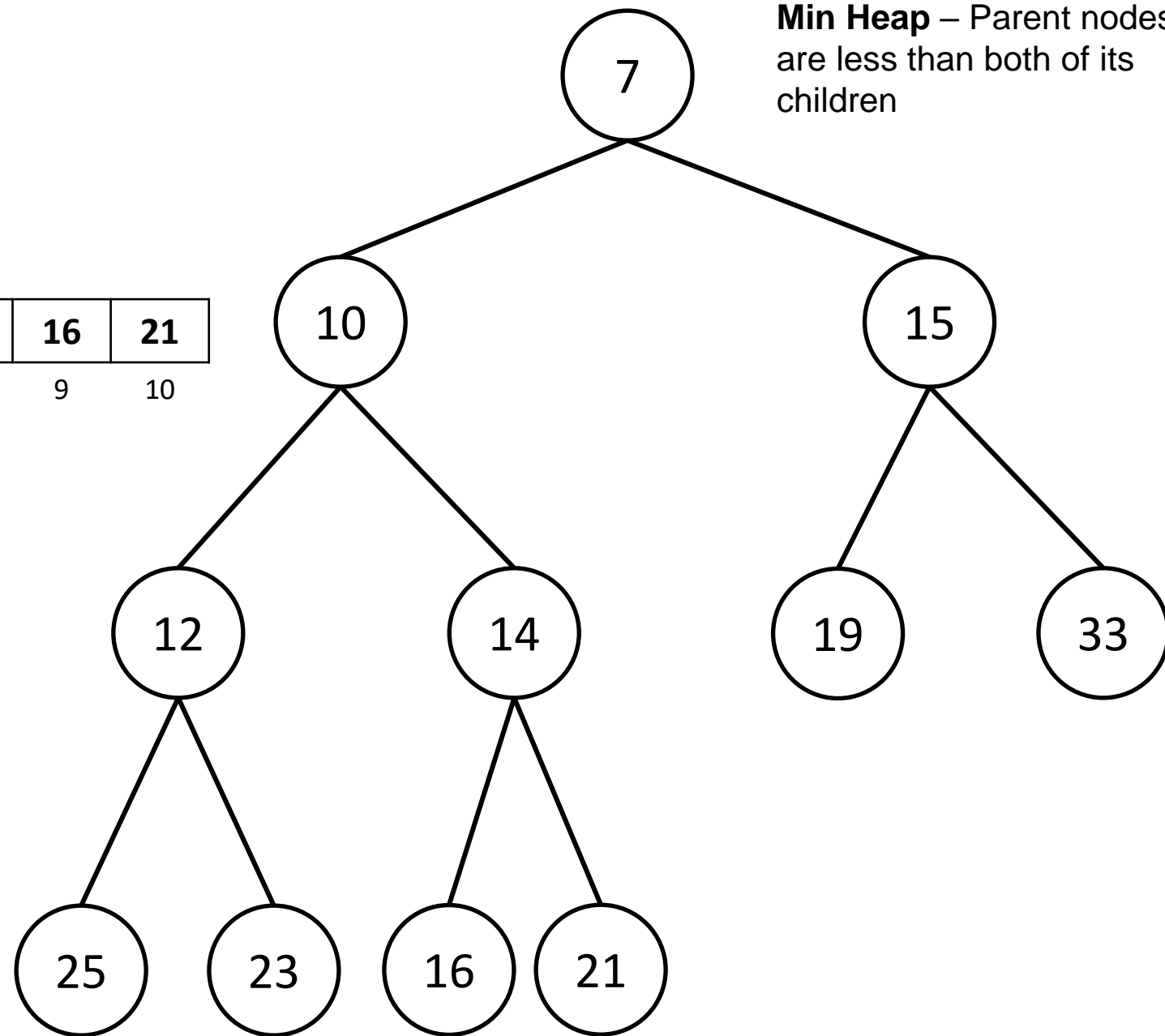
Let's code
this!!!



Min Heap – Parent nodes are less than both of its children

Array

7	10	15	12	14	19	33	25	23	16	21
0	1	2	3	4	5	6	7	8	9	10



What can a Heap do well that other data structures cannot as well?

What can a Heap do well that other data structures cannot as well?

Finding the largest/smallest element happens in **$O(1)$** time

Because we use an array, it might be more memory efficient than a standard tree

Does a Heap remind you of any other data structures?

Does a Heap remind you of any other data structures?

Priority Queue

Does a Heap remind you of any other data structures?

Priority Queue

Whenever we remove an element, we always remove the smallest/largest value (**poll()**)

Whenever we add an element, it initially gets added to the back of the array, and then swaps itself within the array

Takeaways

A Heap is a priority queue

Whenever we remove an element, we always remove the smallest/largest value (**poll()**)

Whenever we add an element, it initially gets added to the back of the array, and then swaps itself within the array

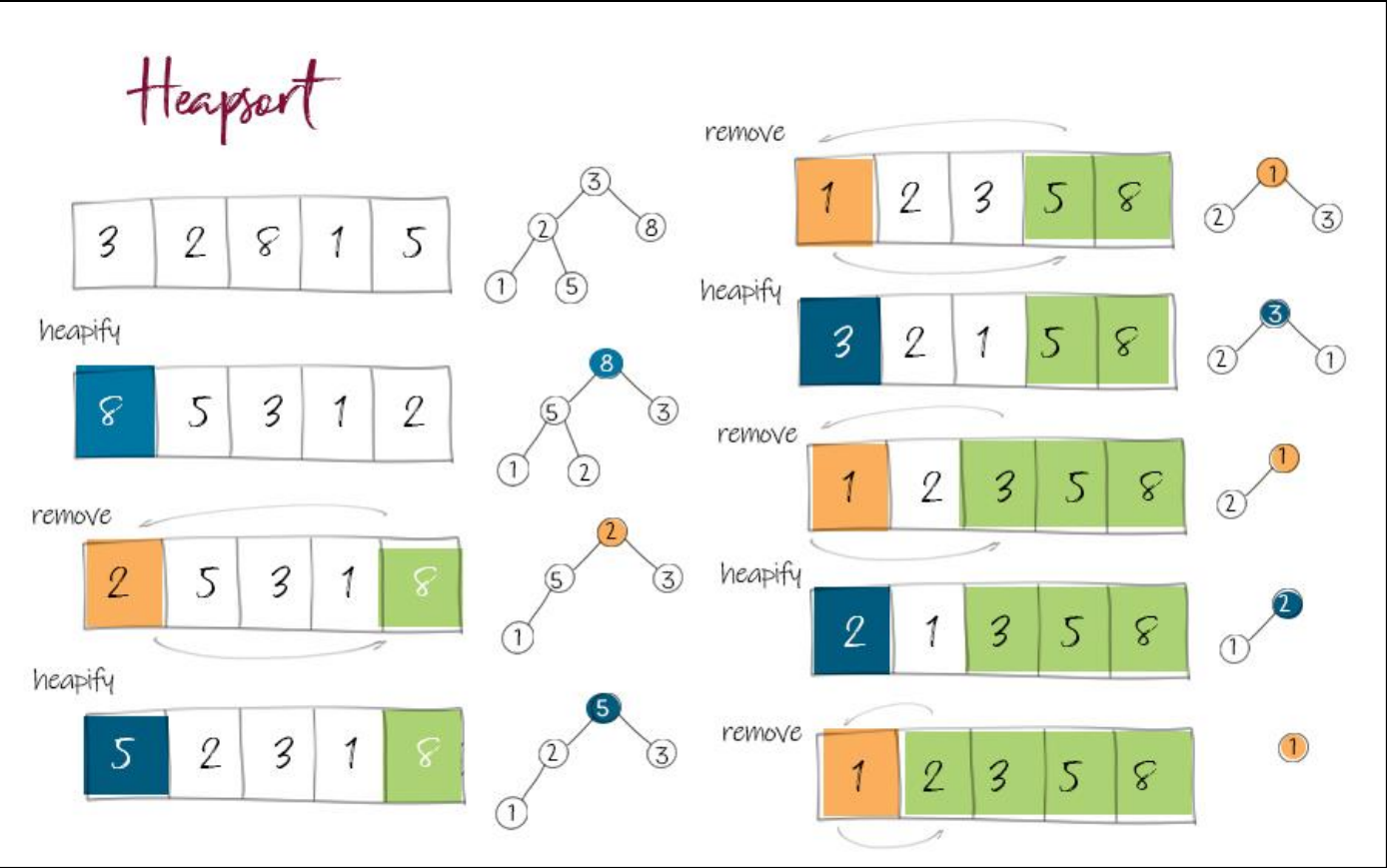
Getting the maximum/minimum value happens in $O(1)$ time

Class PriorityQueue<E>

There is a section of memory in your computer called “The Heap”, which is something totally unrelated to this data structure

Applications

Heapsort- Sorting algorithm that converts an unsorted array to a Heap, and then repeatedly remove the root node



Convert unsorted Array to Heap: $O(n \log n)$

Make array: sortedArray $O(n)$

while(!heap.isEmpty()): $O(n)$

$x = \text{heap.poll}()$ $O(\log n)$

insert x into sortedArray $O(1)$

return sortedArray $O(1)$

Total Running Time: **$O(n \log n)$**

(Same as merge sort, quick sort)