CSCI 232: Data Structures and Algorithms

Shortest Path (Part 1)

Reese Pearsall Spring 2024

Announcements

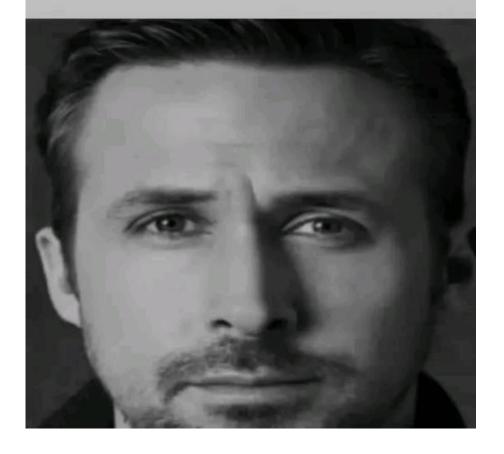
Lab 10 due **Friday** (Part 1 of Program 3)

No Office hours Tomorrow

All of program 3 has been posted

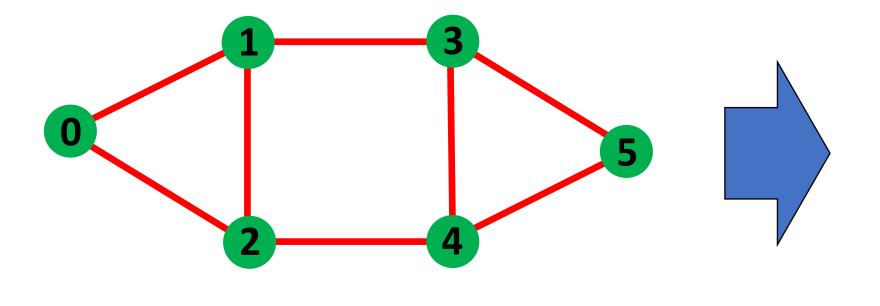
Survey results

Me at 7am choosing between my future or my Bed

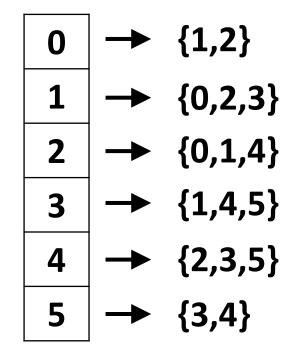


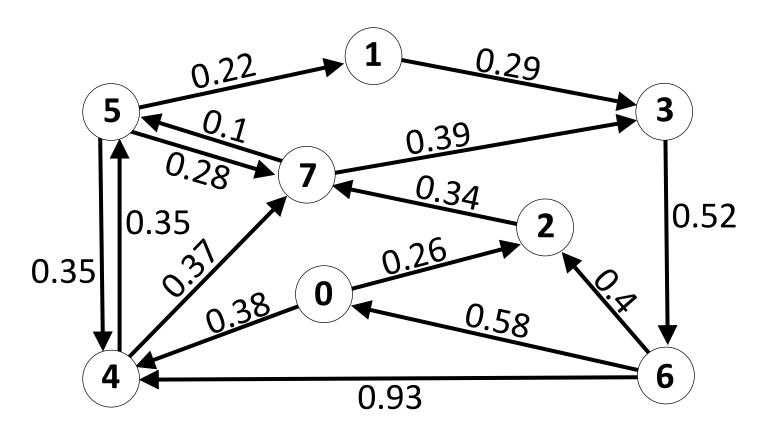
Graphs

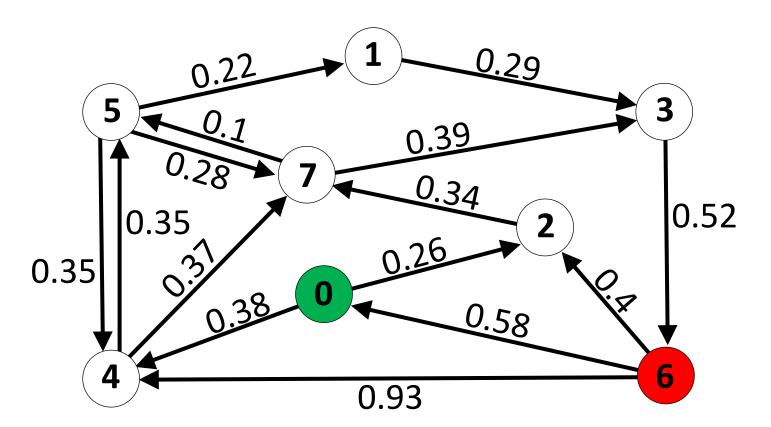
$$G = (V, E)$$



Adjacency List

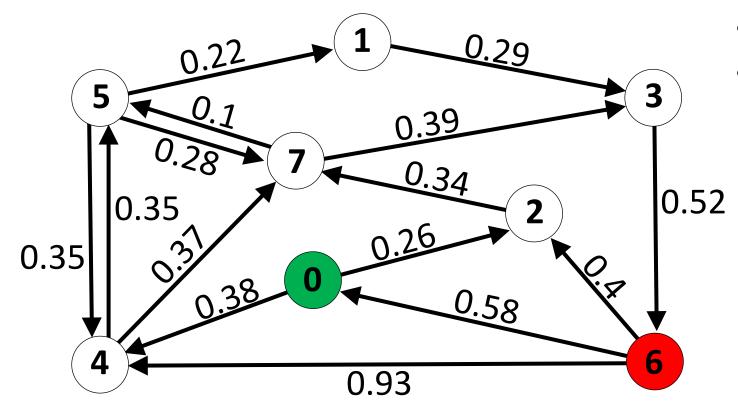






Path with the smallest sum of edge weights.

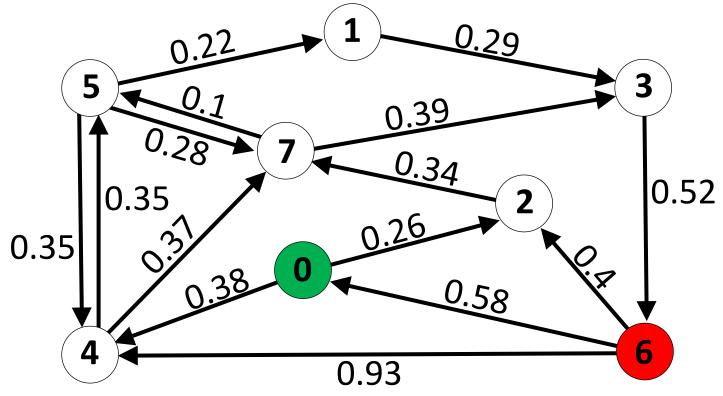
What is the shortest path between vertex 0 and vertex 6?



Assumptions:

- Graph is directed.
- Graph is edge-weighted.
- Edge weights are non-negative.
- Graph need not be simple (though our example will be).

What is the shortest path between vertex 0 and vertex 6?

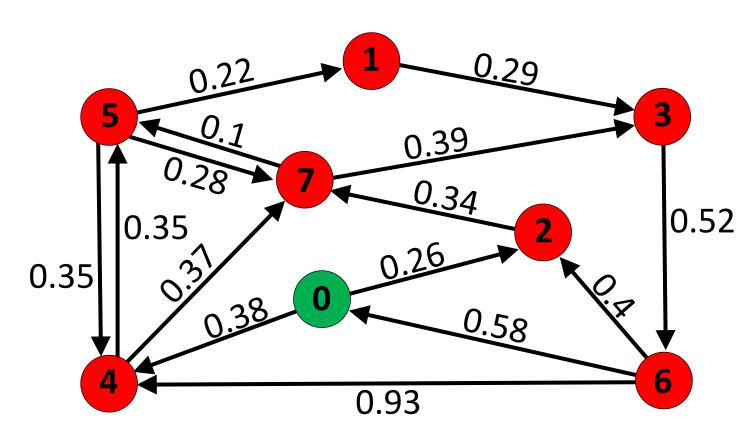


Assumptions:

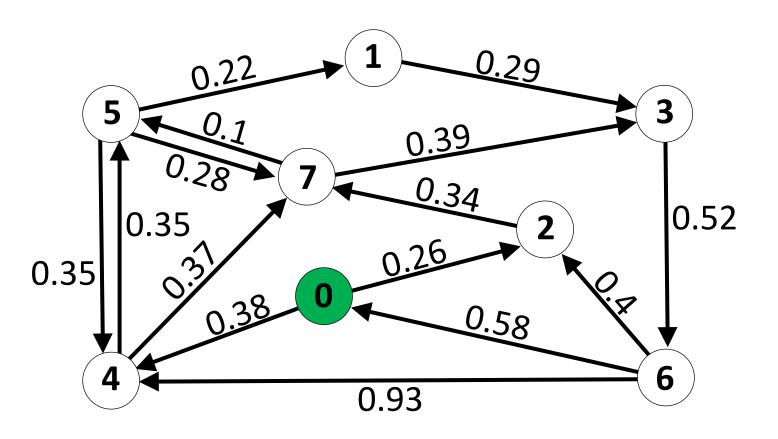
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What is the shortest path between vertex 0 and vertex 6?

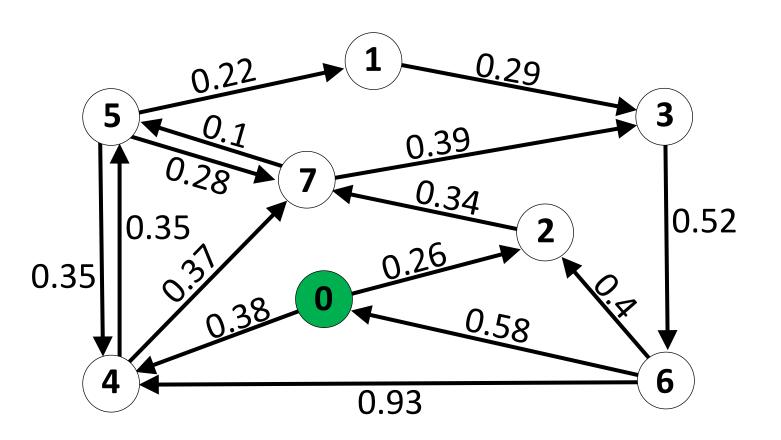


We are going to find the shortest path between vertex 0 and every other vertex, flooding out from 0.

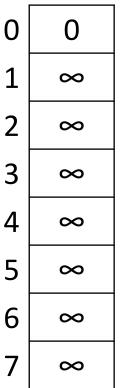


Distance from 0

?
?
?
?
?
?
?



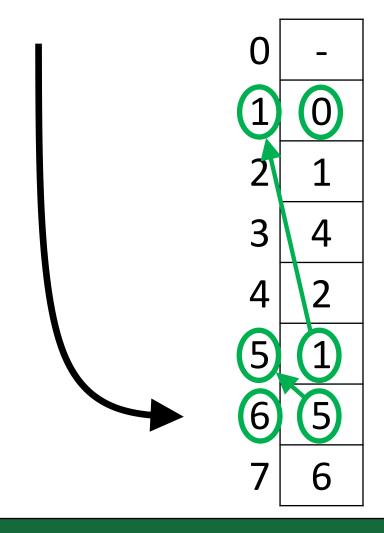
Distance from 0

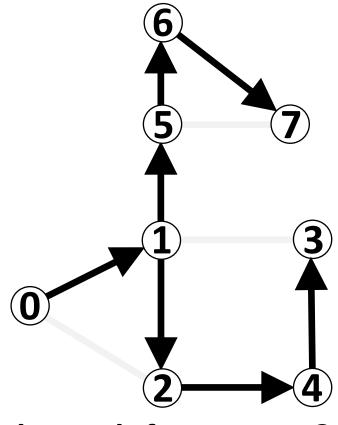


How can we keep track of routes?

Graphs - Paths

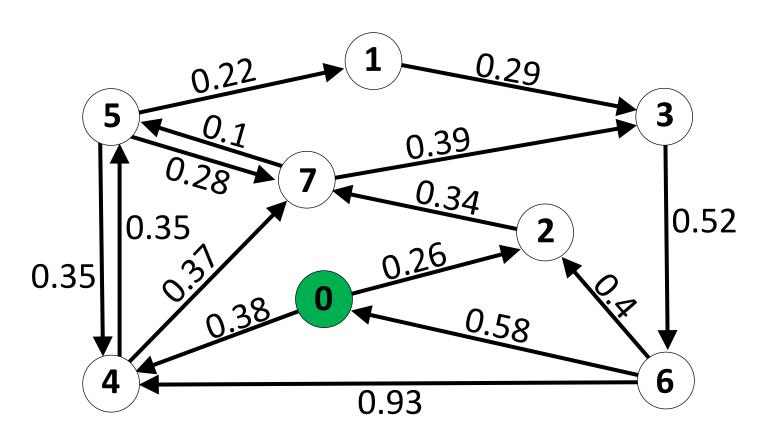
int[] previousVertex





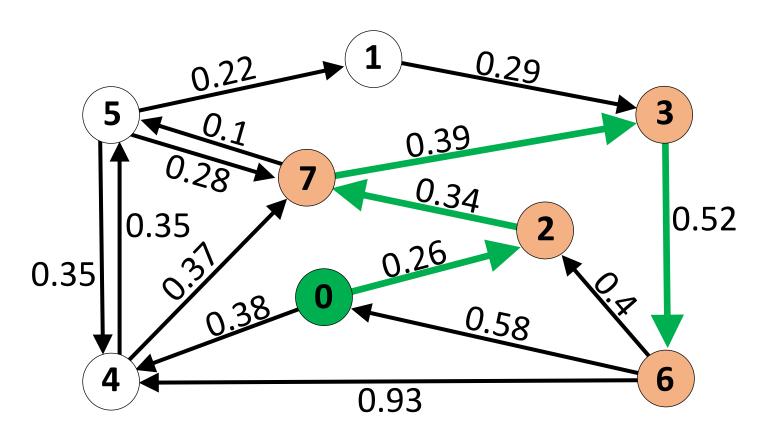
How do we determine the path from 0 to 6?

Start at vertex 6. Find its previous vertex. Find its previous vertex... until we get back to the start (0).



Distance **Previous** from 0 vertex 0 0 0 ∞ ∞ 3 3 ∞ 4 4 ∞ 5 5 ∞ 6 ∞ 6 7 ∞

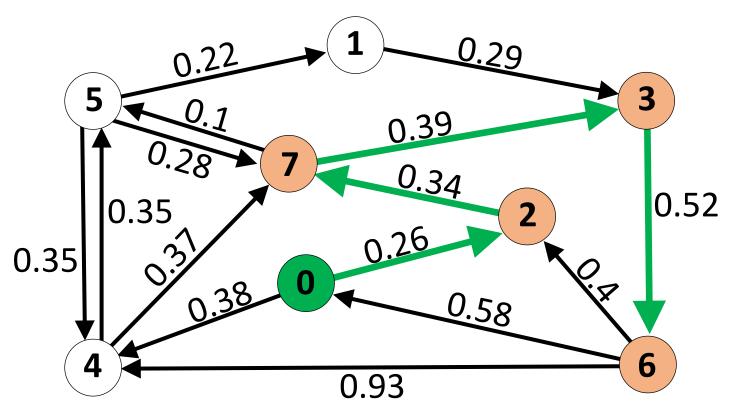
How can we keep track of routes?

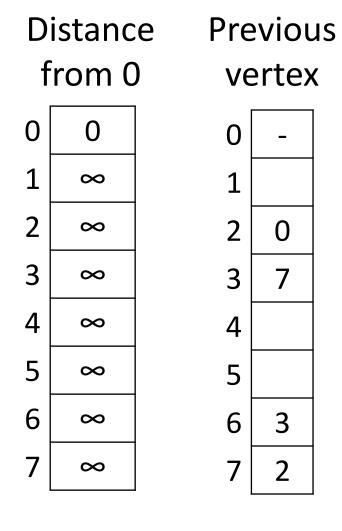


Distance from 0 0 0 ∞ 0.26 0.99 4 ∞ 5 ∞ 1.51 6 0.60

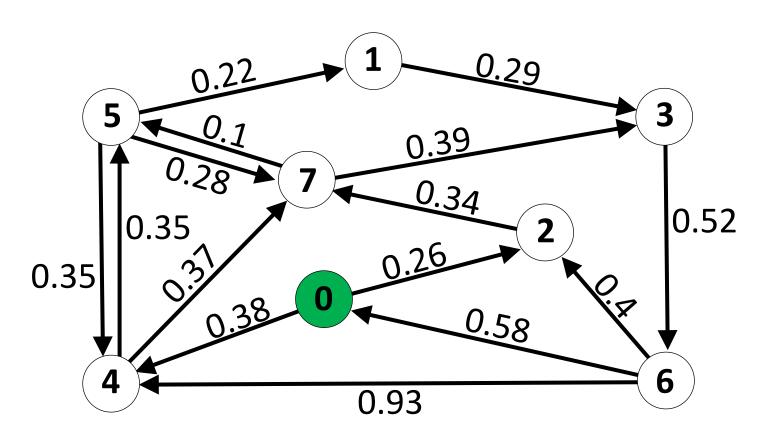
Previous vertex

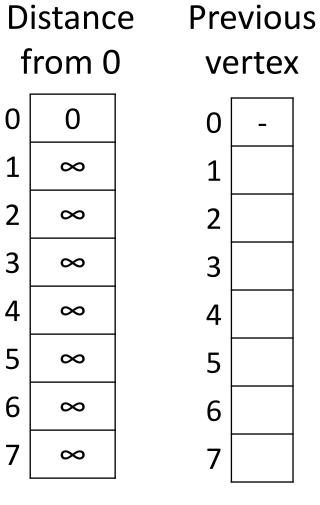
How can we keep track of routes?

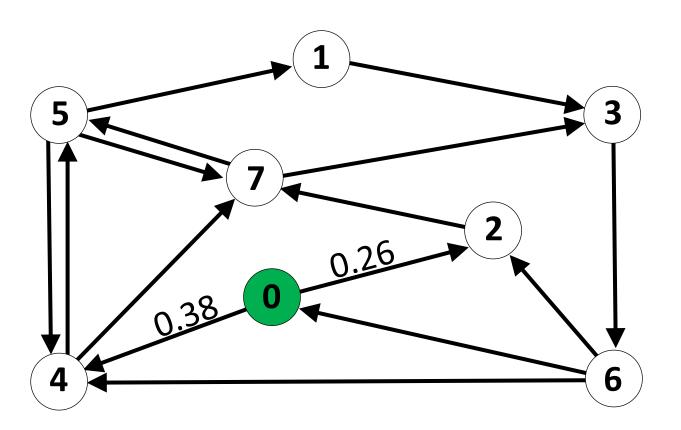


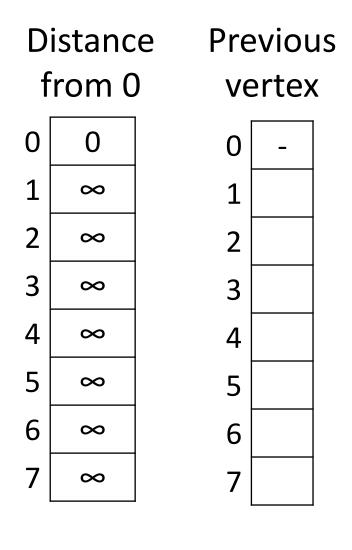


If this is the shortest path from 0 to 6, what can we say about the shortest path from 0 to 3?

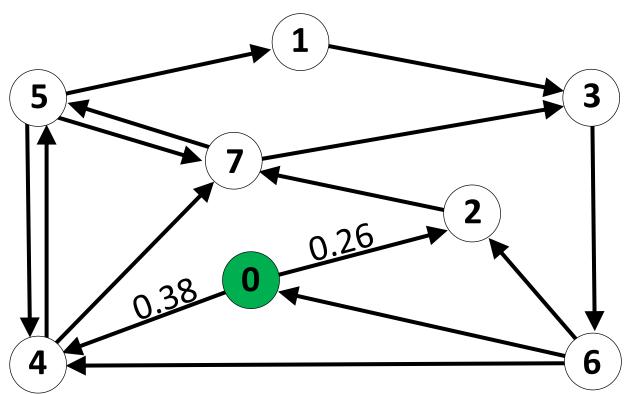


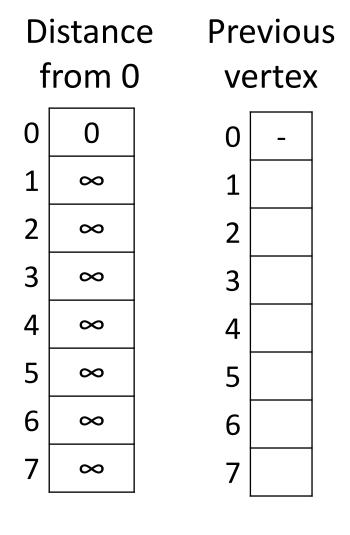




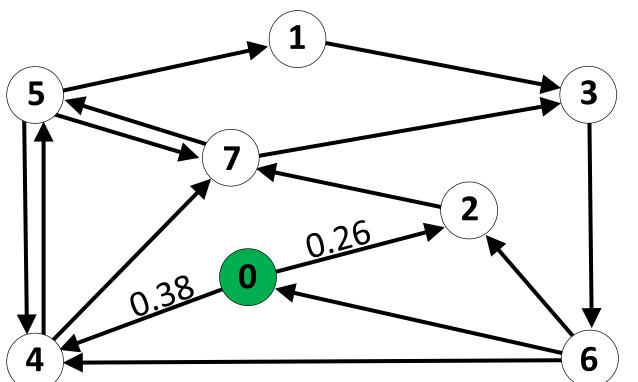


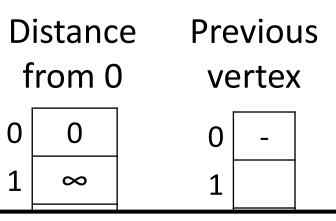
Claim: There cannot possibly be a shorter path from 0 to 2 than the edge from 0 to 2 because...?





Claim: There cannot possibly be a shorter path from 0 to 2 than the edge from 0 to 2 because non-negative edge weights mean every other path is at least 0.38 or 0.26.

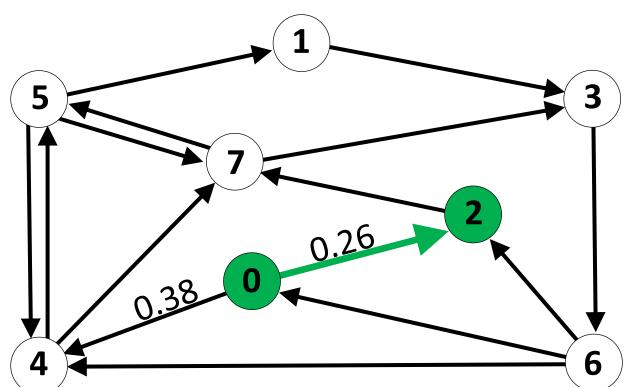


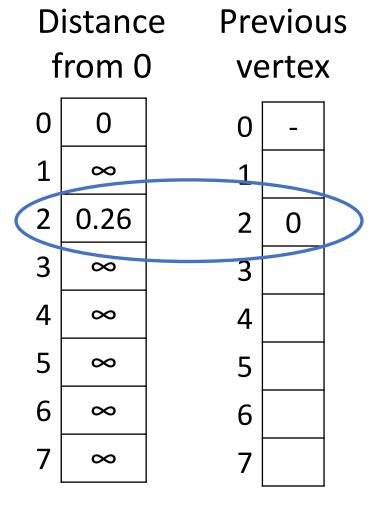


Can we say the same thing about the edge from 0 to 4?

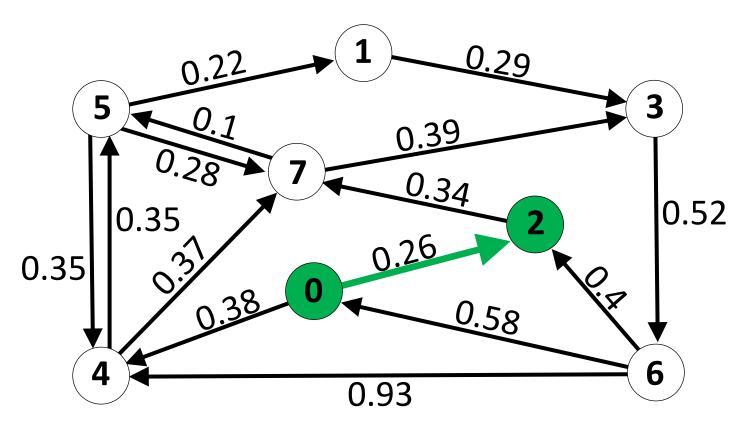
I.e., Could there be a shorter path from 0 to 4 other than the edge from 0 to 4?

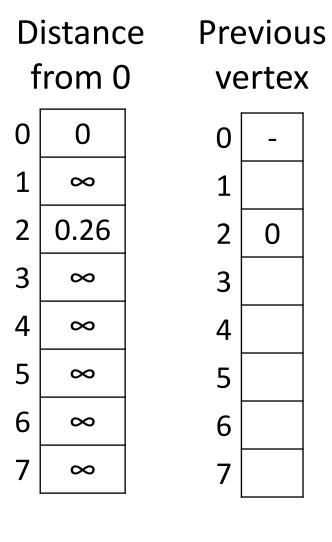
Claim: There cannot possibly be a shorter path from 0 to 2 than the edge from 0 to 2 because non-negative edge weights mean every other path is at least 0.38 or 0.26.



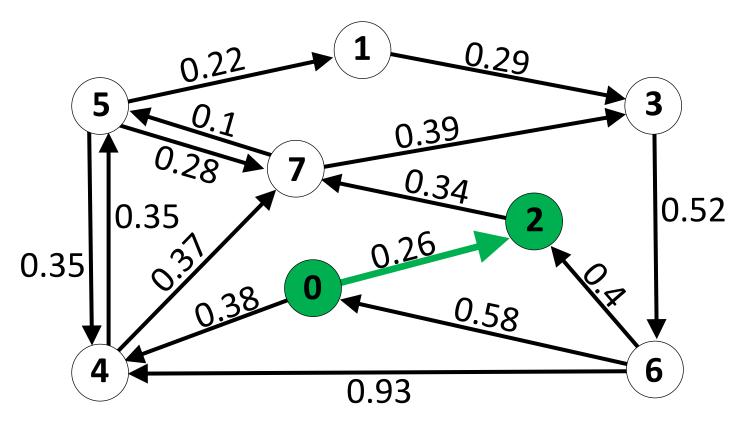


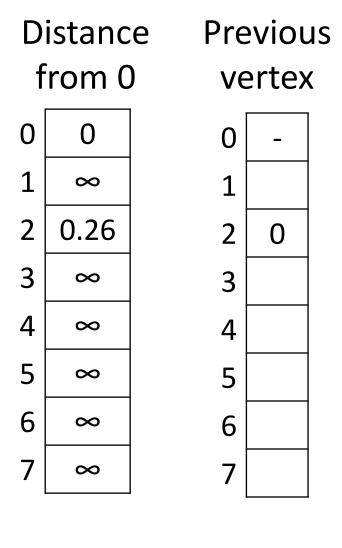
Claim: There cannot possibly be a shorter path from 0 to 2 than the edge from 0 to 2 because non-negative edge weights mean every other path is at least 0.38 or 0.26.



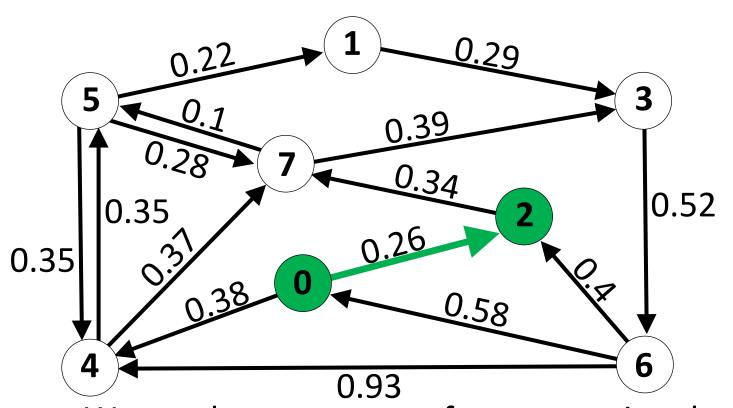


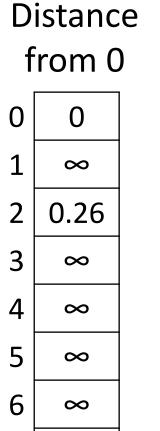
We need some process for progressing through the graph.





We need some process for progressing through the graph. What if we prioritized neighbors based on path (not edge) distance?





 ∞

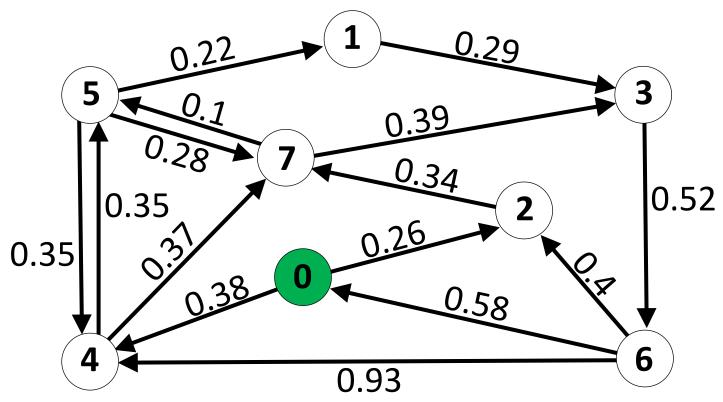
Previous **Priority** vertex queue

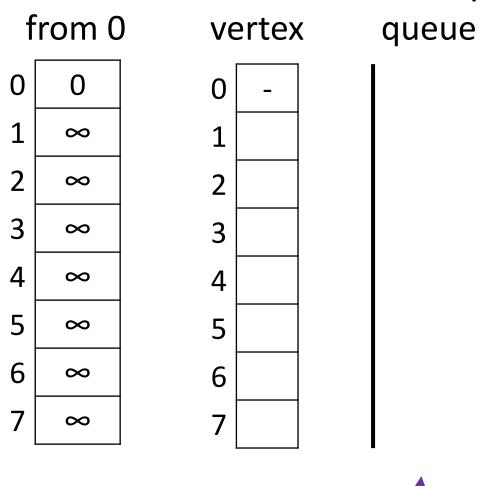
6

We need some process for progressing through the graph.

What if we prioritized neighbors based on path (not edge) distance?

vertex (distance)





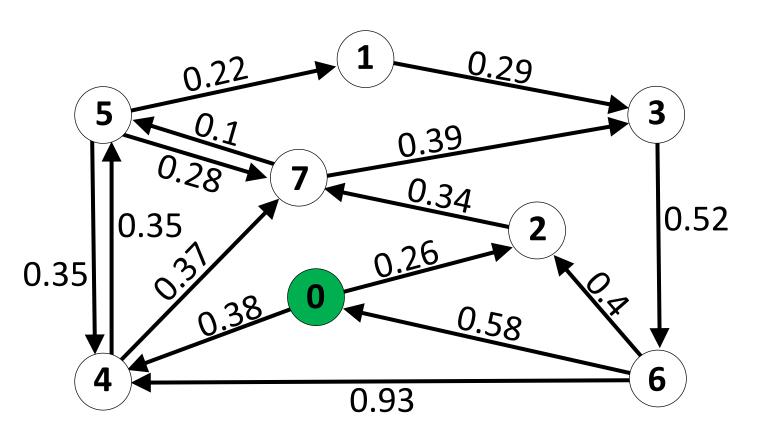
Previous

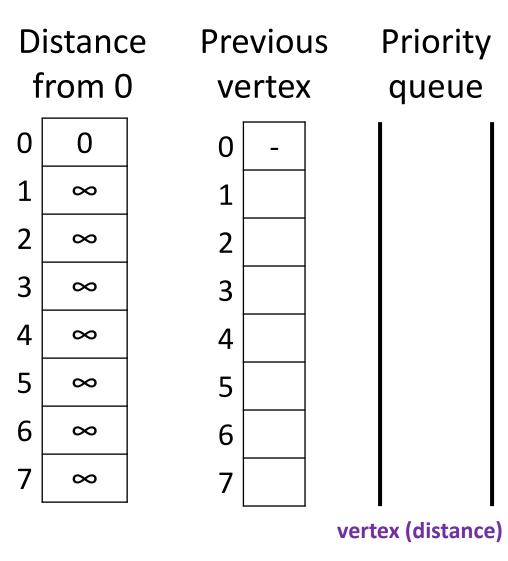
Priority

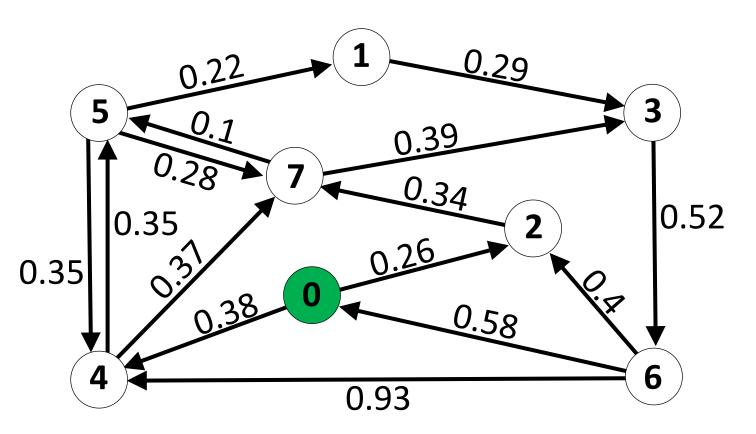
We need some process for progressing through the graph.
What if we prioritized neighbors based on path (not edge) distance?

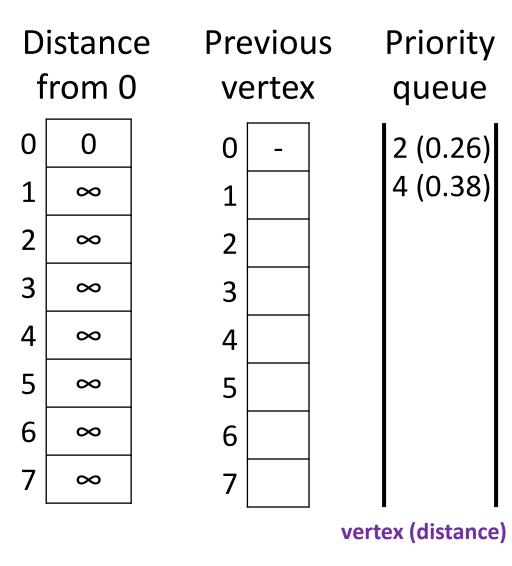
vertex (distance)

Distance









Distance Previous **Priority** Shortest Path from 0 vertex queue 2 (0.26) 0 4 (0.38) ∞ ∞ 3 3 ∞ 0.34 4 ∞ 0.52 0.35

5

6

 ∞

 ∞

 ∞

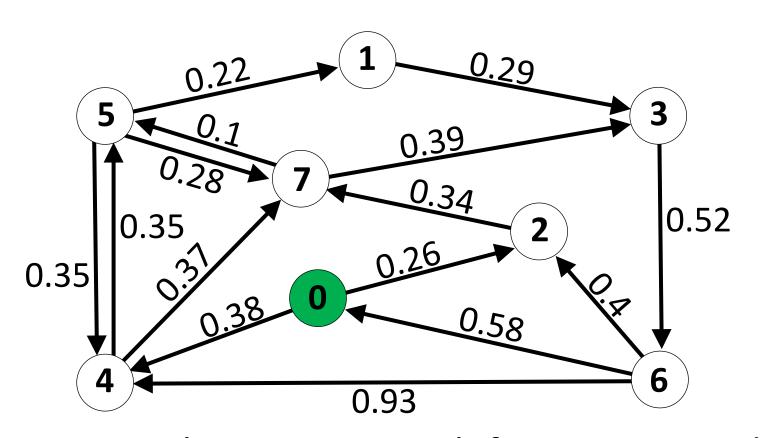
6

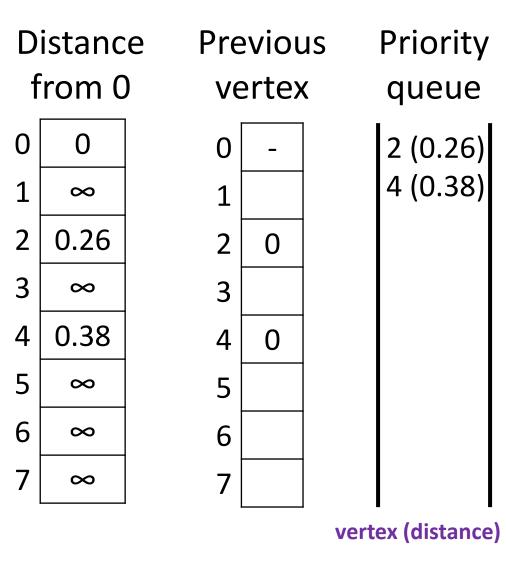
What can we reach from connected vertices and at what distance (from 0)?

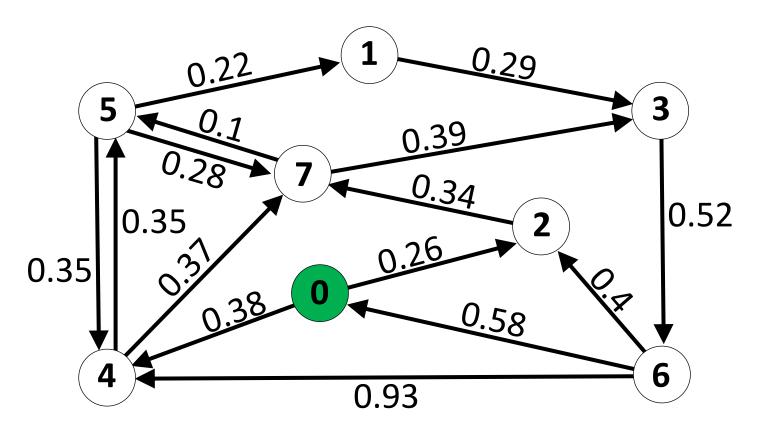
0.93

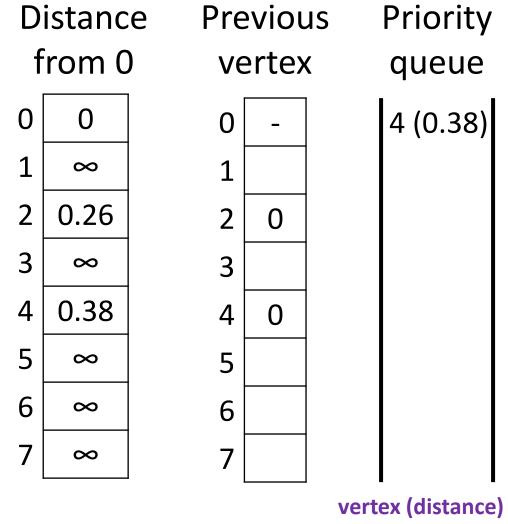
0.35

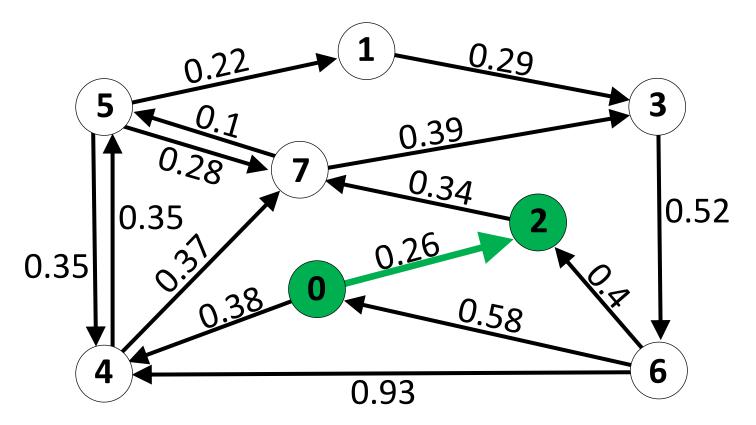
vertex (distance)

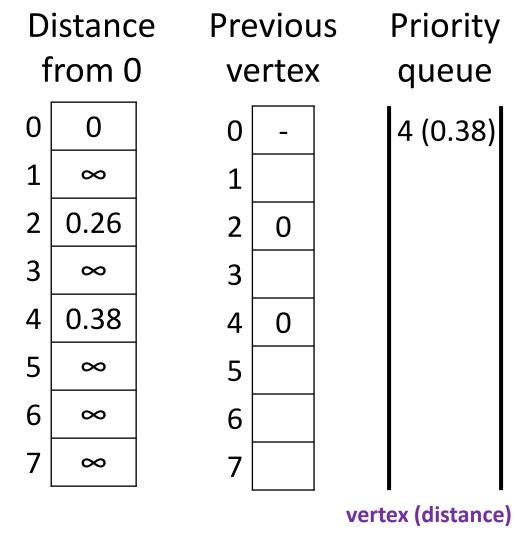


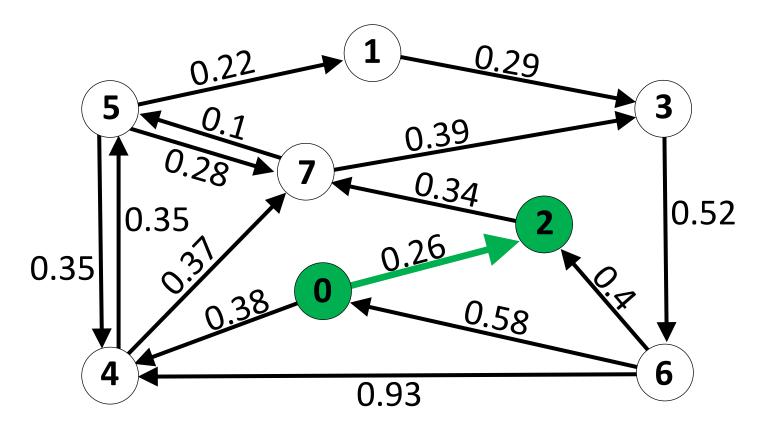


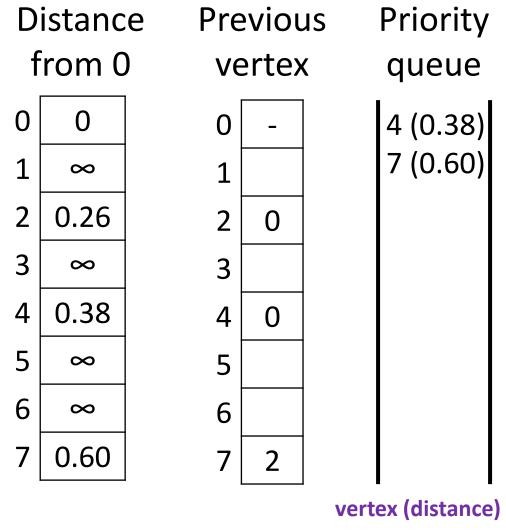


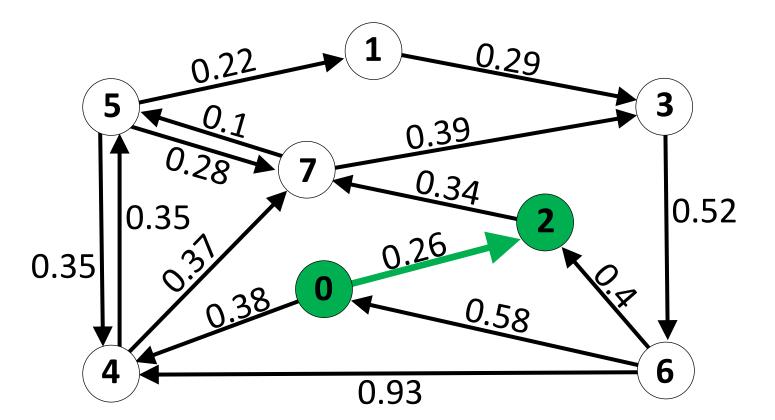




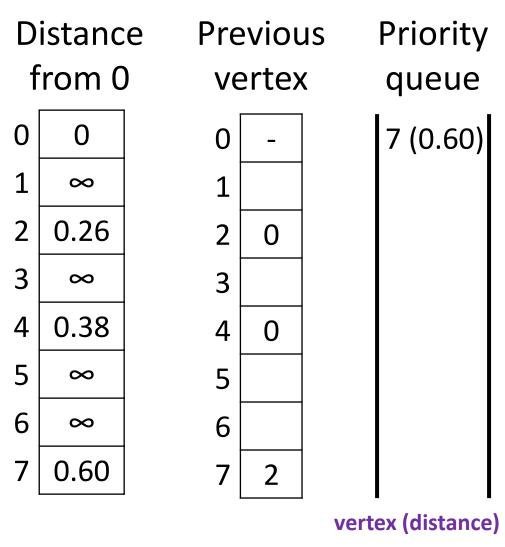


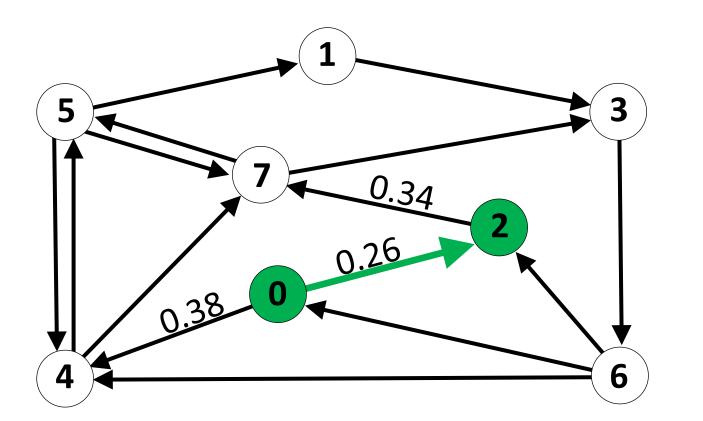


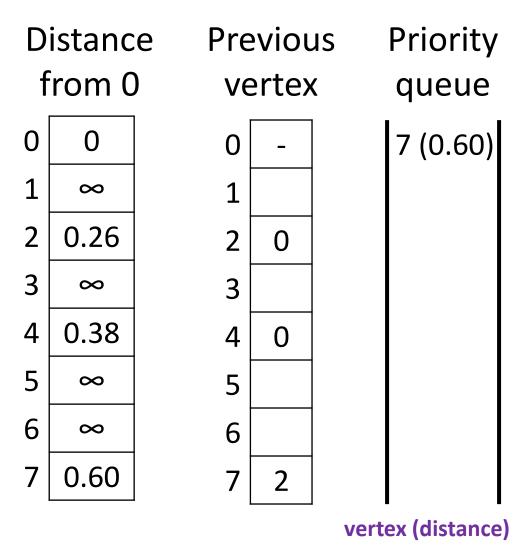




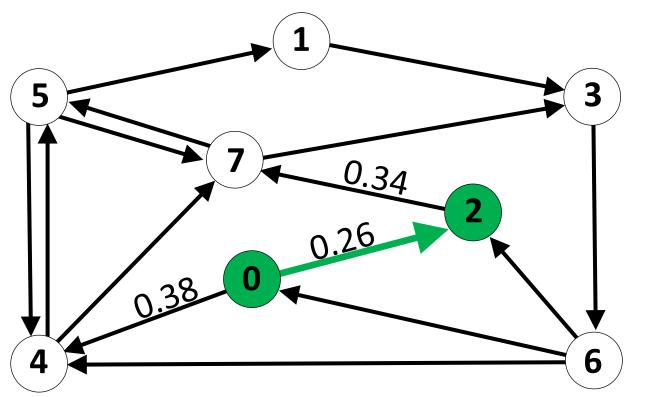
Repeat.

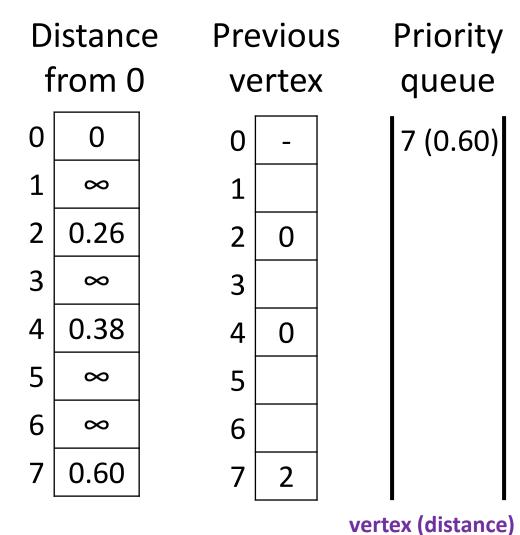




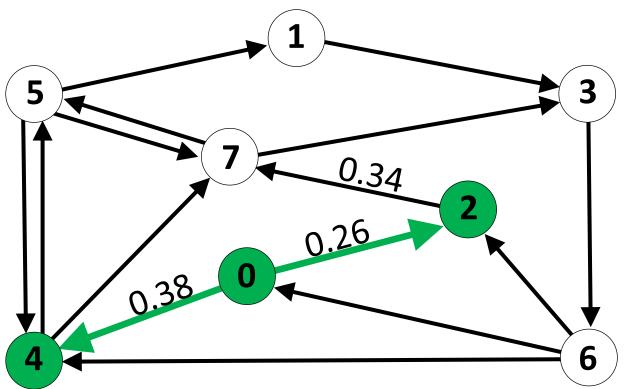


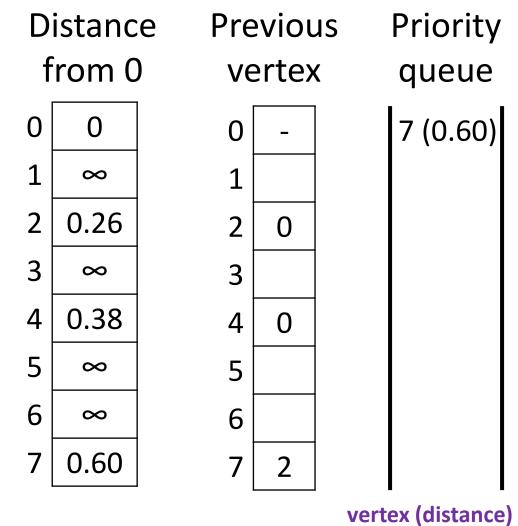
What can we say about the shortest path from 0 to 4?



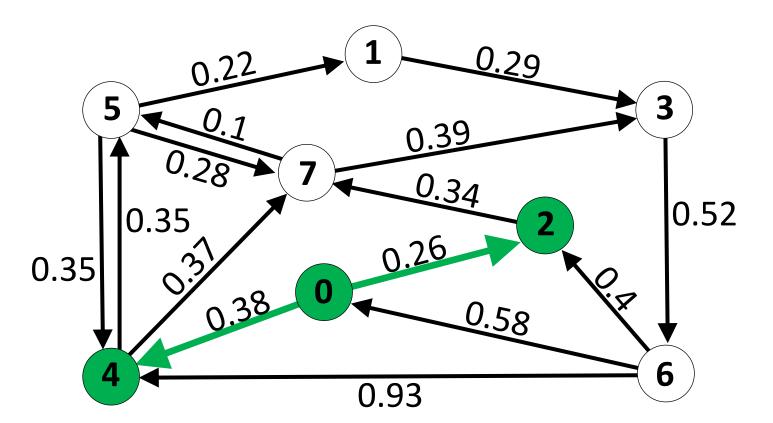


The 0 to 4 edge has to be the shortest path between 0 and 4, since any other path would go from $0 \rightarrow 2 \rightarrow 7 \rightarrow ?$ at cost at least 0.26 + 0.34 = 0.6 > 0.38

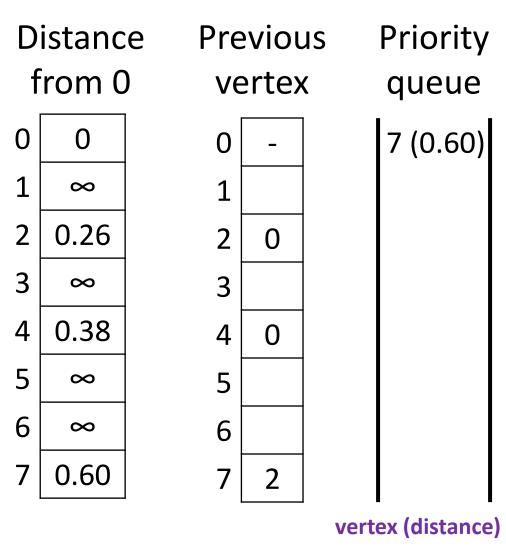


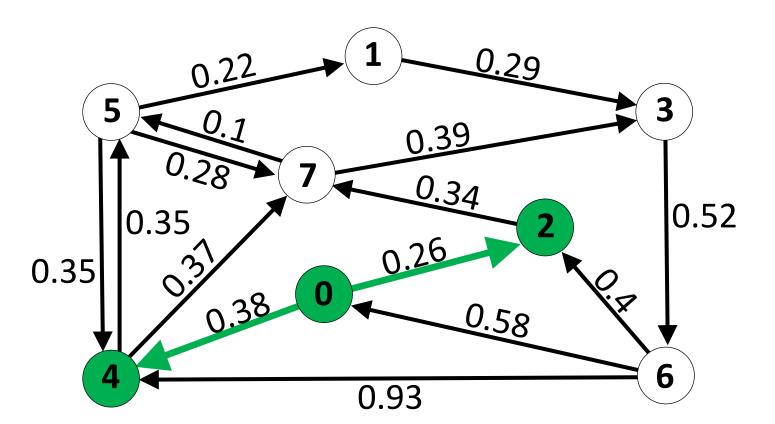


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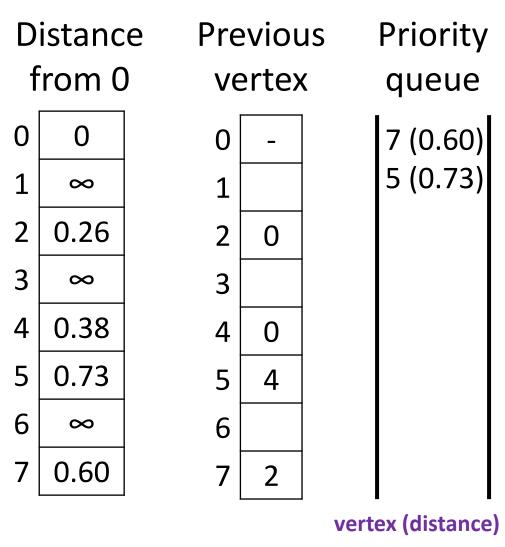


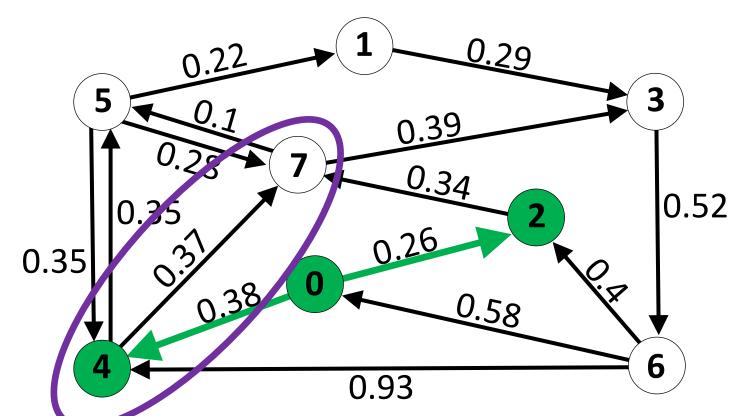
Add neighbors to queue/previous.





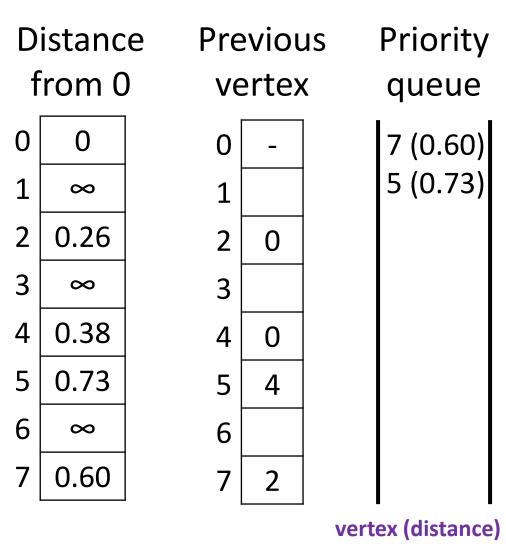
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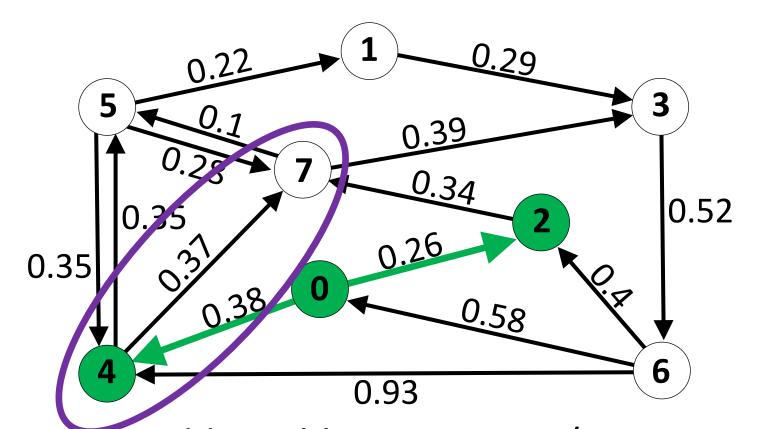


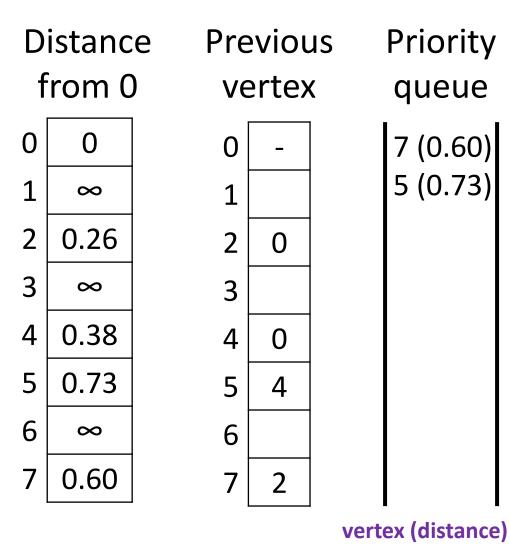


Add neighbors to queue/previous.

We have another route to 7!

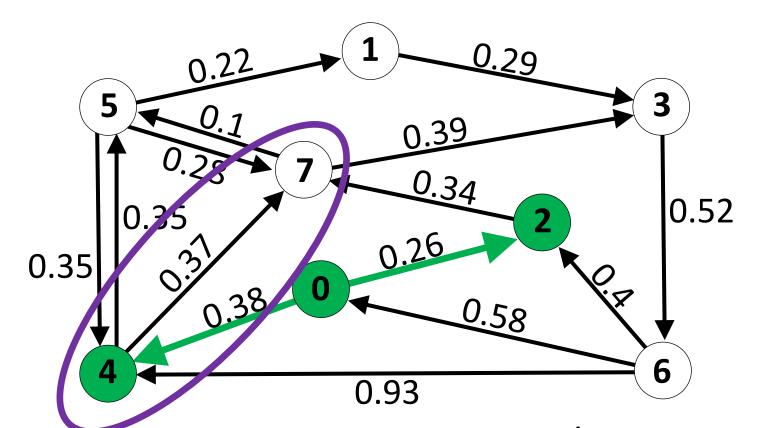


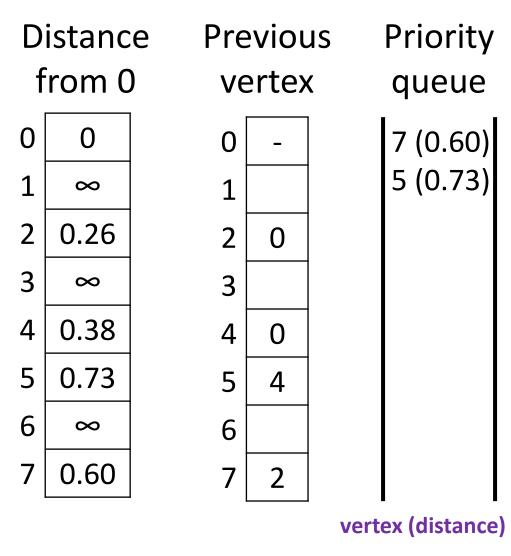




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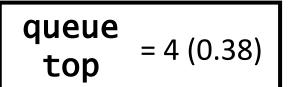
We have another route to 7! Check to see if it is shorter!

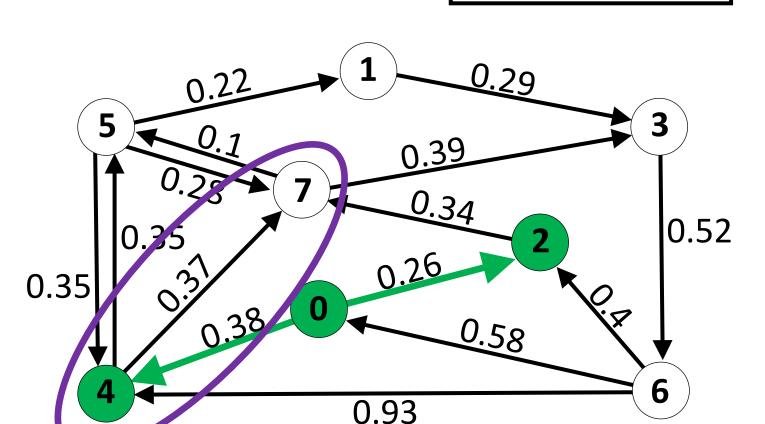


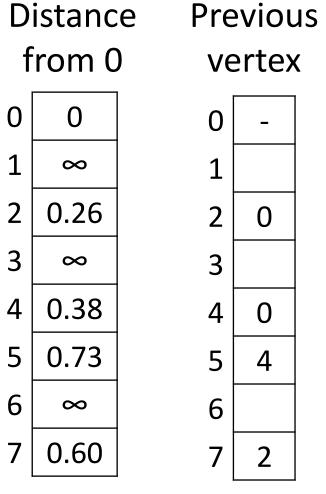


Add neighbors to queue/previous.

We have another route to 7! Check to see if it is shorter! It's not (0.38 + 0.37 = 0.75 > 0.60).







vertex (distance)

Priority

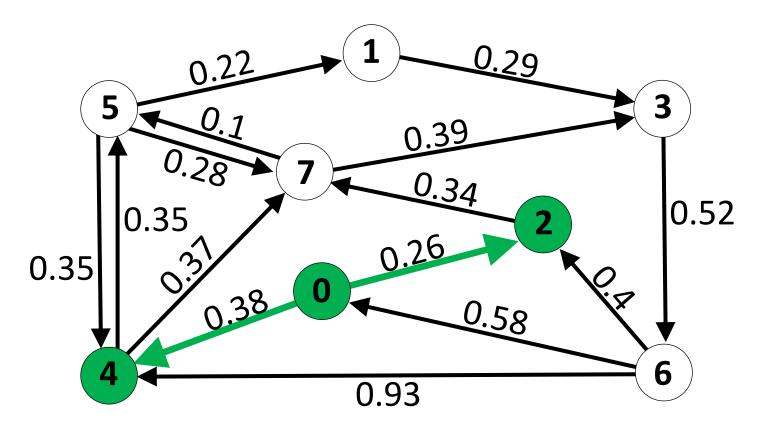
queue

7 (0.60)

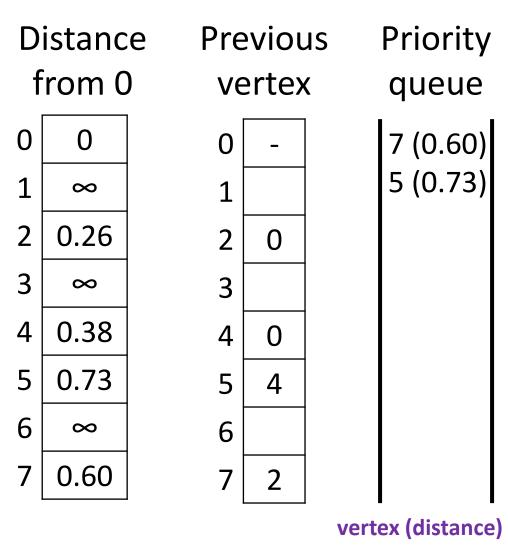
5 (0.73)

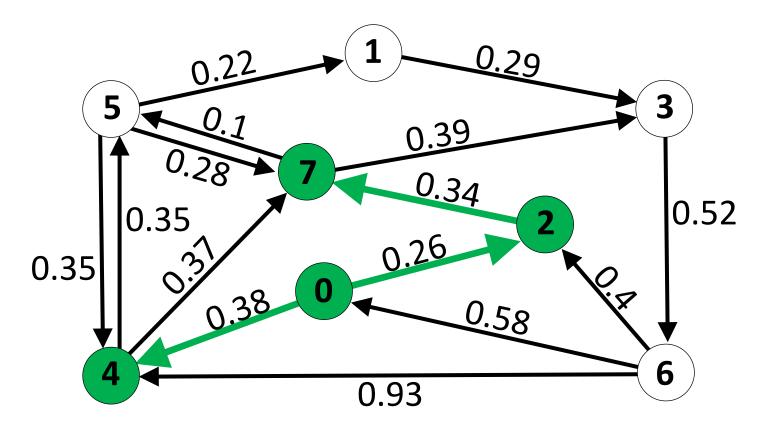
Rule: When processing vertex V, only add/modify queue for neighbor U if and only if:

distance[v] + weight(v, u) < distance[u]</pre>

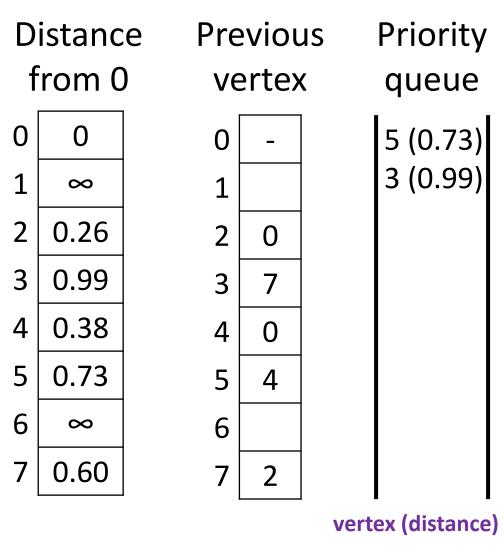


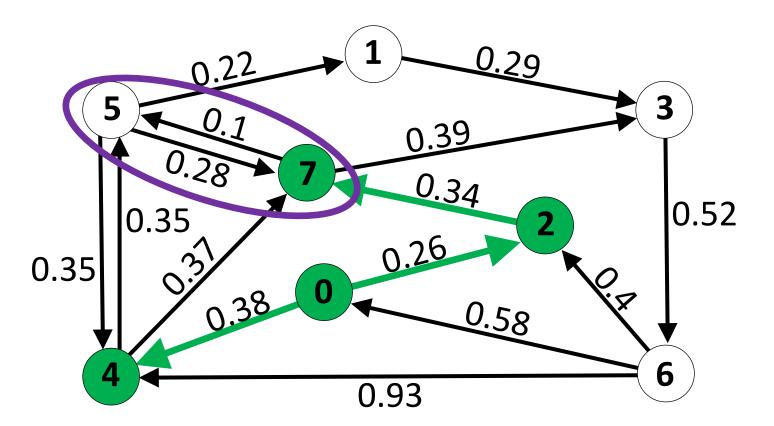
Repeat.

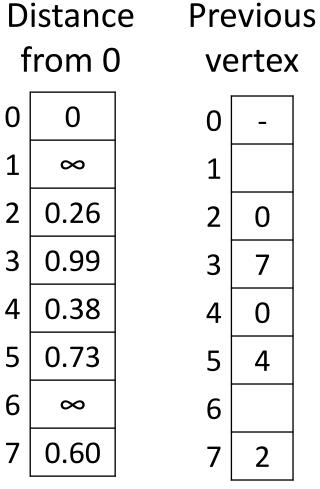




Repeat.





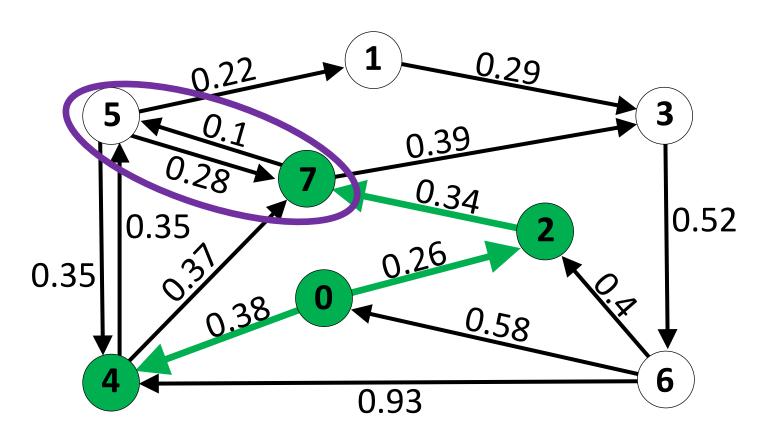


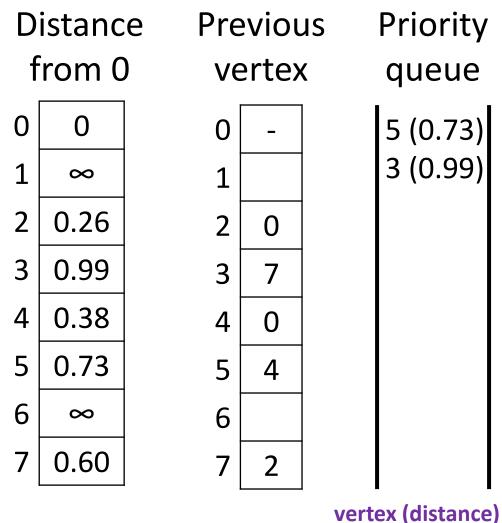
Previous		Priority
ertex	(queue
-		5 (0.73)
		3 (0.99)
0		
7		
0		
4		
2		
	- 0 7 0 4	- 0 7 0 4

Repeat.

We have another route to 5, and at cost 0.7 < 0.73.

vertex (distance)





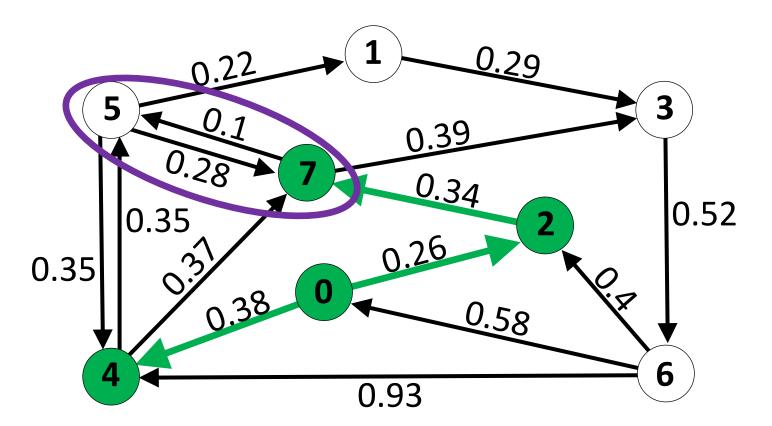
Repeat. We have another route to 5, and at cost 0.7 < 0.73. i.e., distance[v] + weight(v, u) < distance[u]

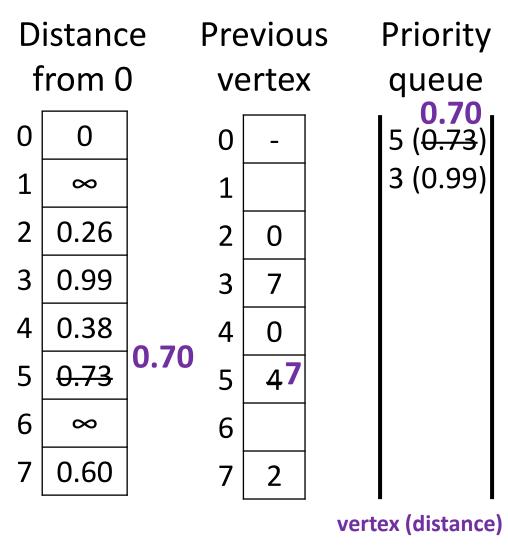
Priority

queue

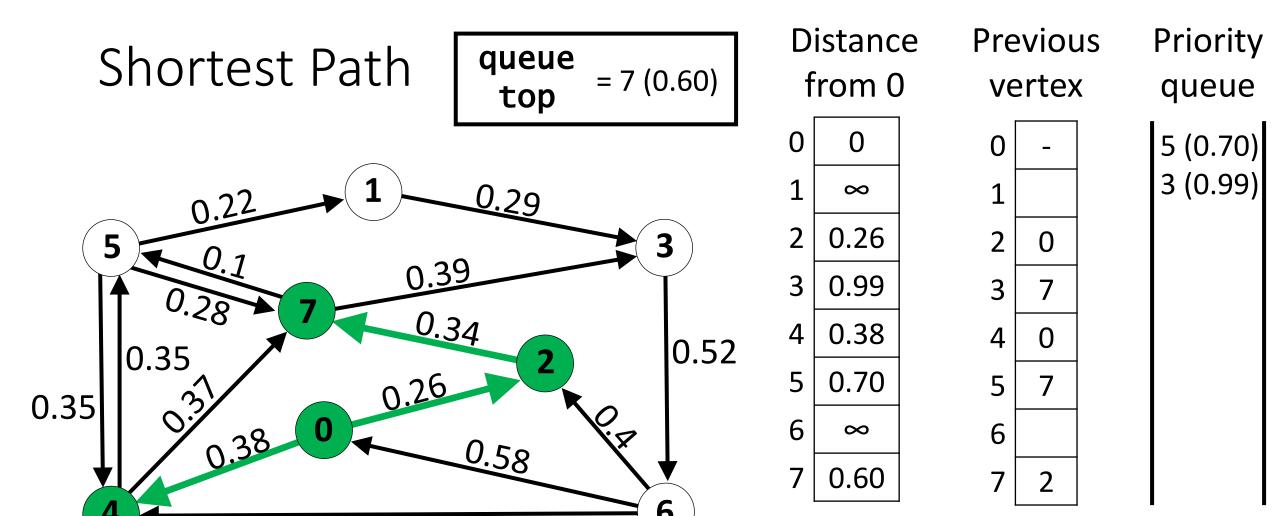
5 (0.73)

3 (0.99)





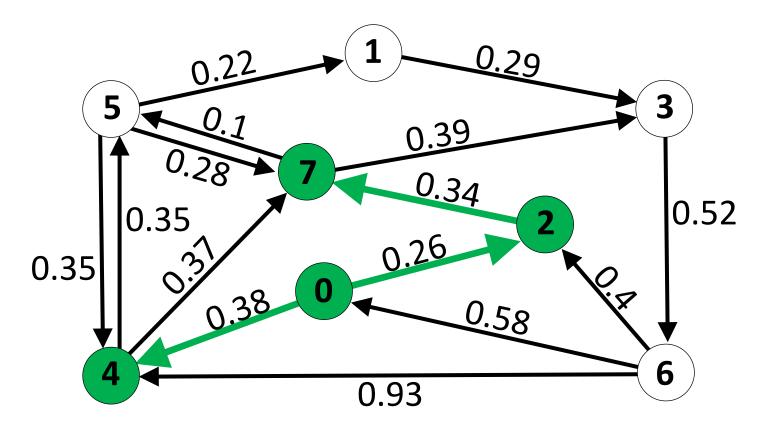
Repeat. We have another route to 5, and at cost 0.7 < 0.73. So updated queue/previous/distance.



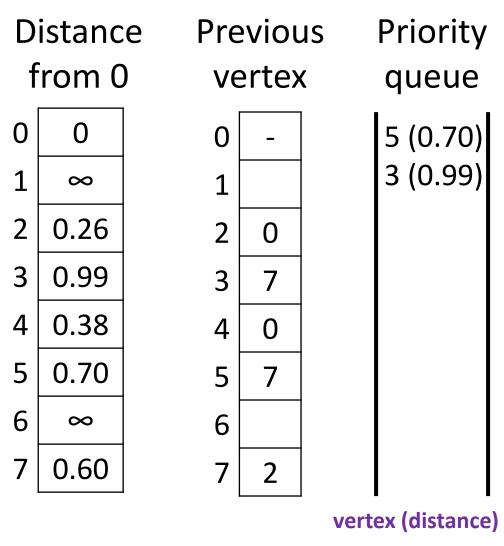
Repeat. We have another route to 5, and at cost 0.7 < 0.73. So updated queue/previous/distance.

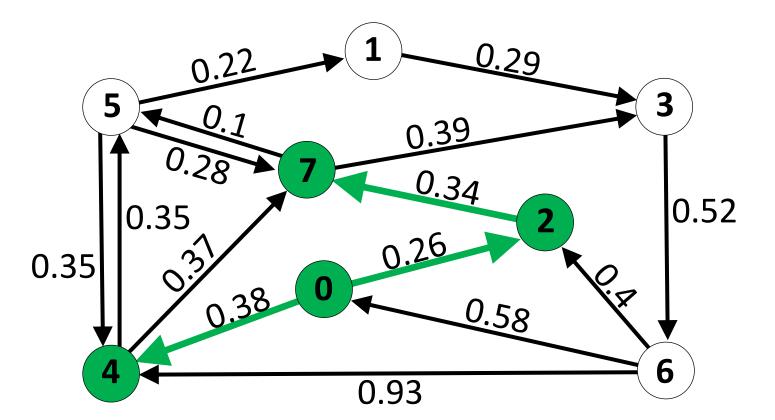
0.93

vertex (distance)

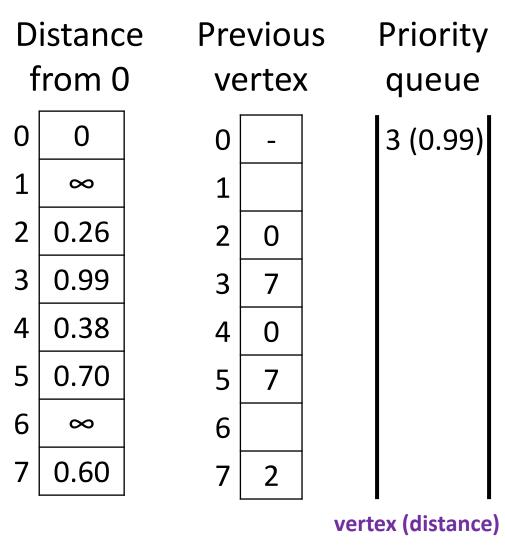


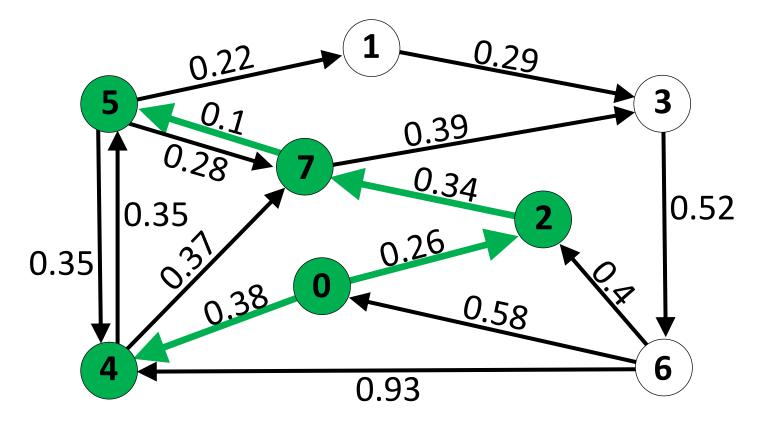
Repeat.



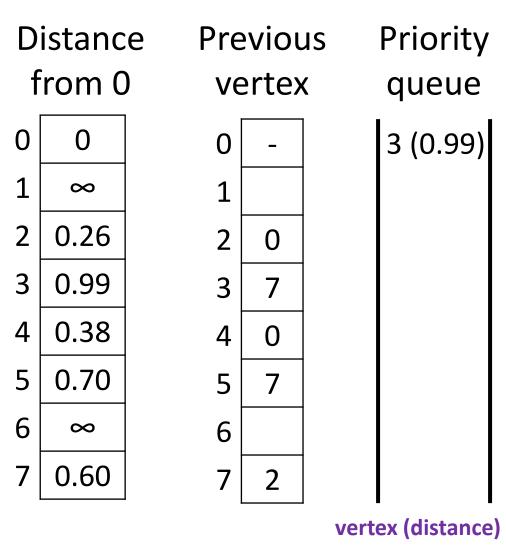


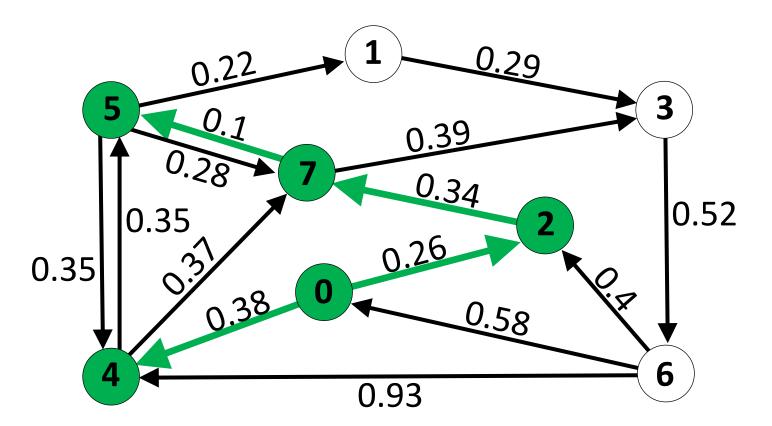
Repeat.



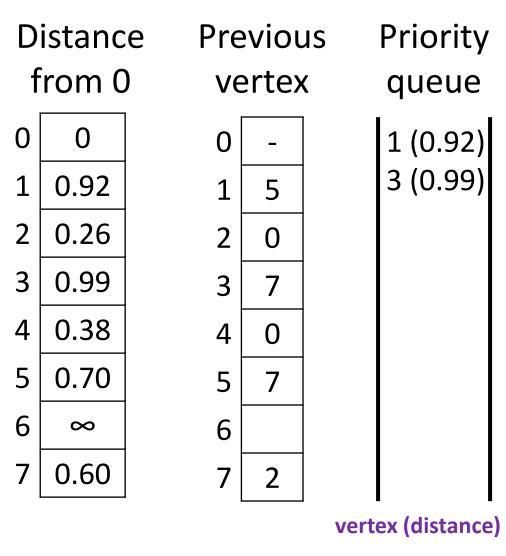


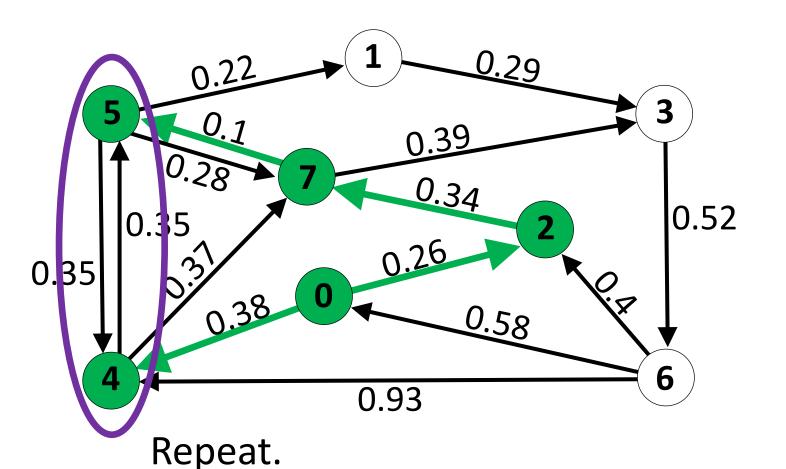
Repeat.

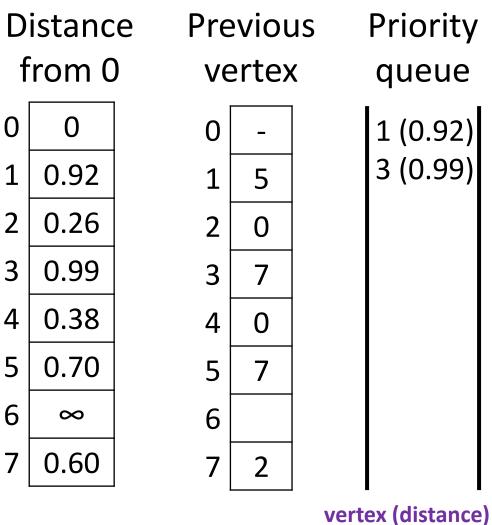




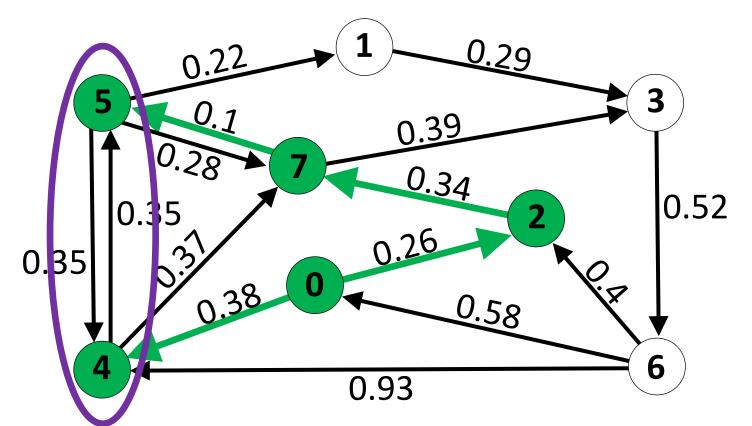
Repeat.

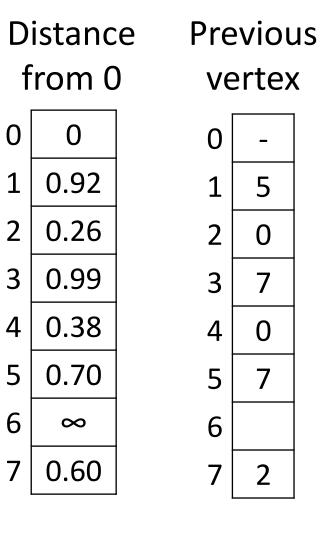






What about neighbor 4?





vertex (distance)

Priority

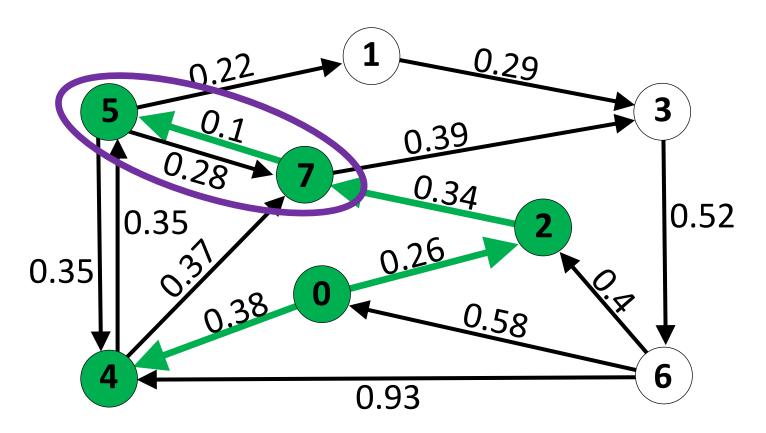
queue

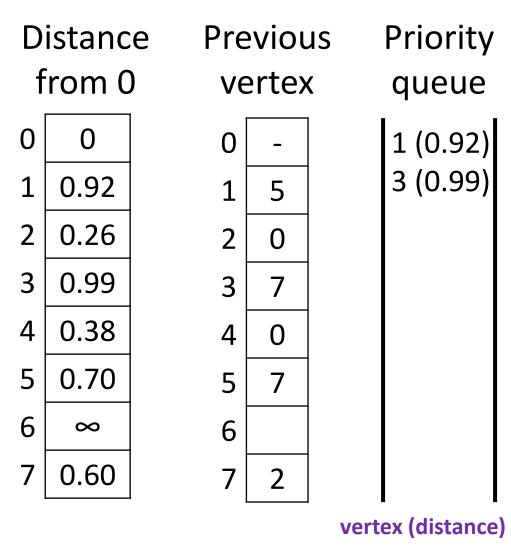
1 (0.92)

3 (0.99)

Repeat.

What about neighbor 4?distance[5] + weight(5, 4) = 0.70 + 0.35 = 1.05 ⊄ 0.38 = distance[4]

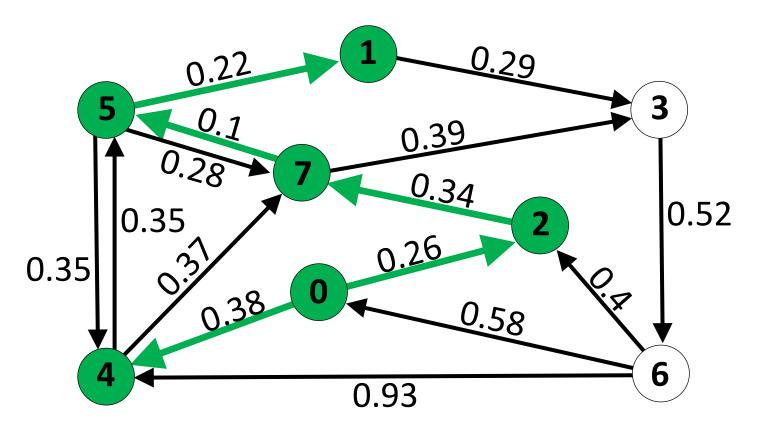




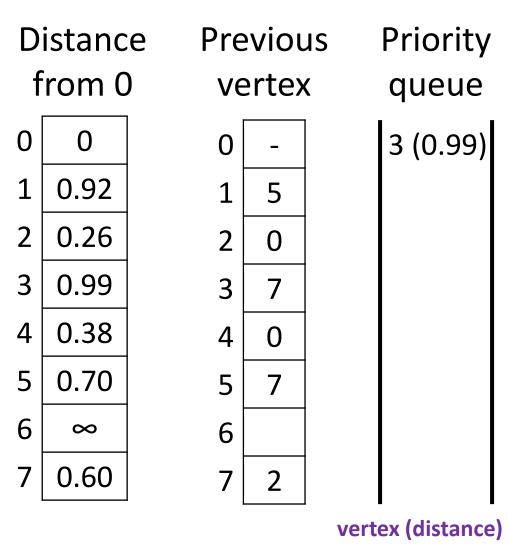
Repeat. What about neighbor 7?

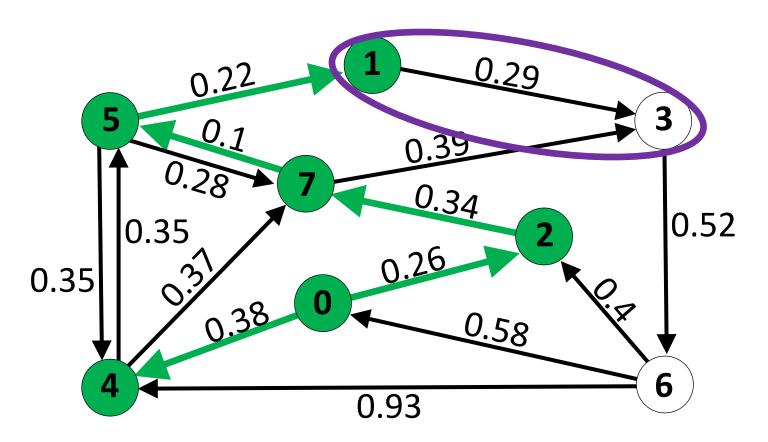
distance[5] + weight(5, 7) = 0.70 + 0.28 = 0.98 < 0.60 = distance[7]

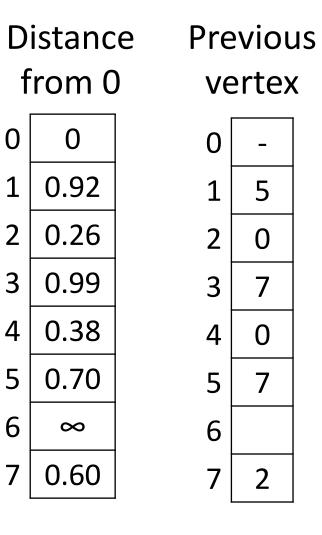




Repeat.







Repeat.

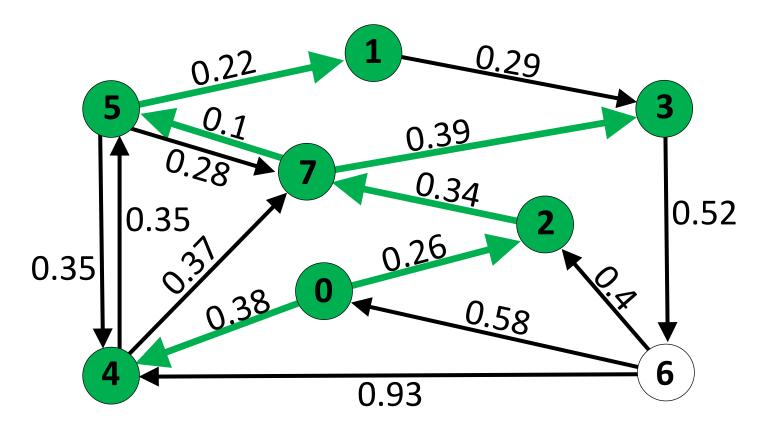
What about neighbor 3? 0.92 + 0.29 = 1.21 > 0.99

vertex (distance)

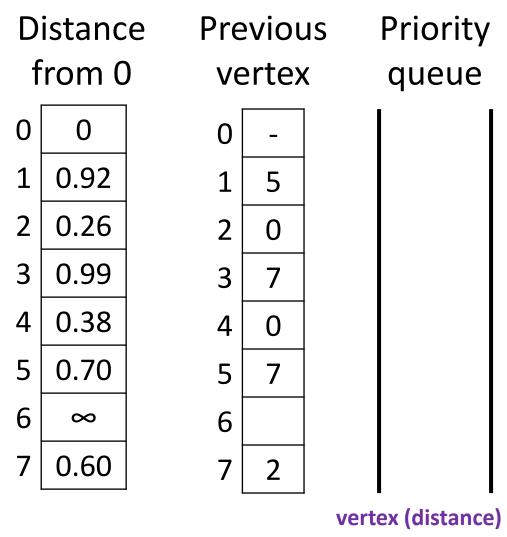
Priority

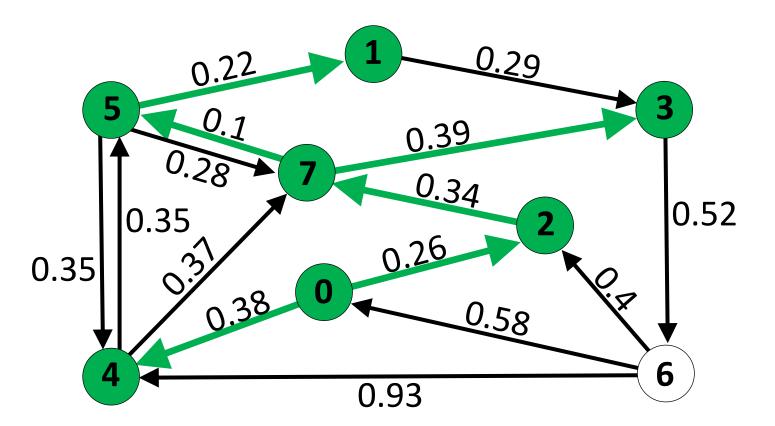
queue

3 (0.99)

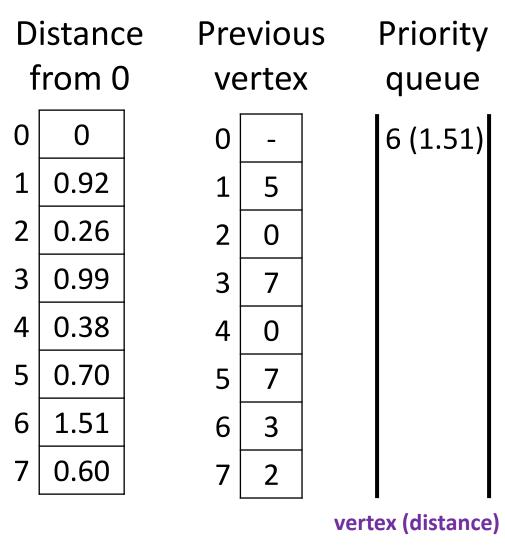


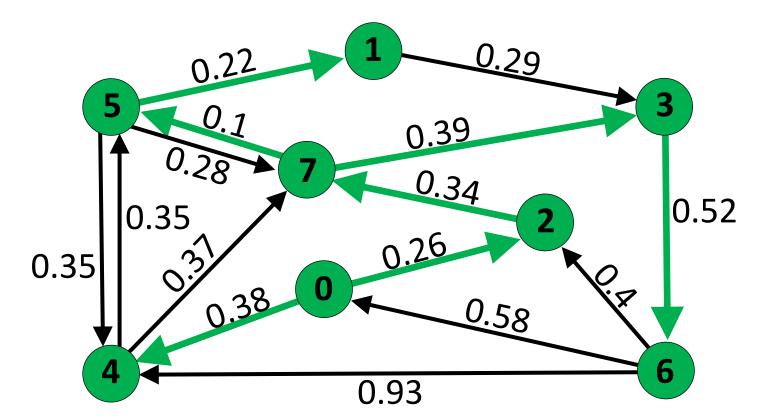
Repeat.



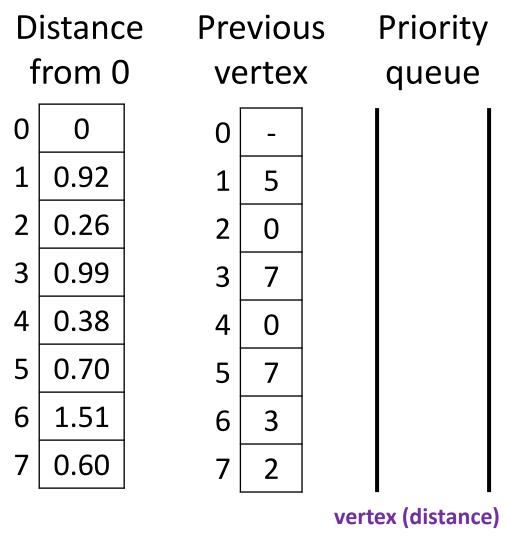


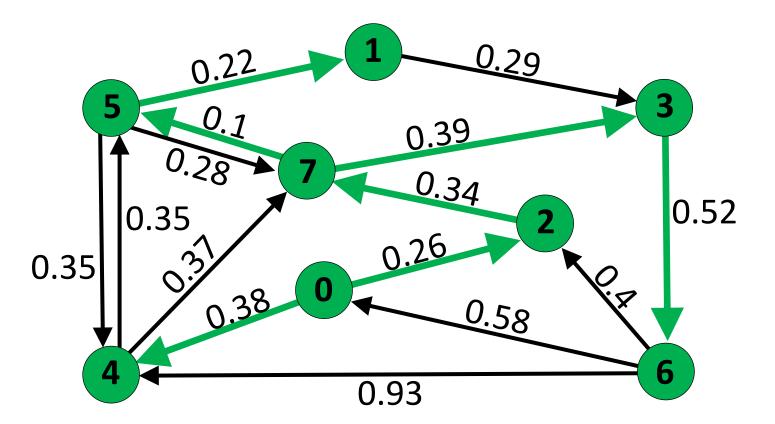
Repeat.



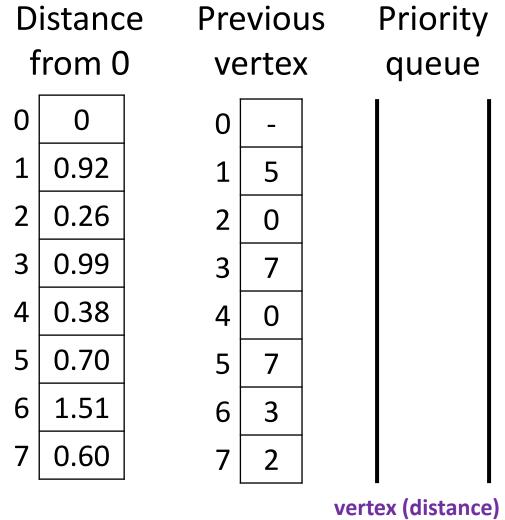


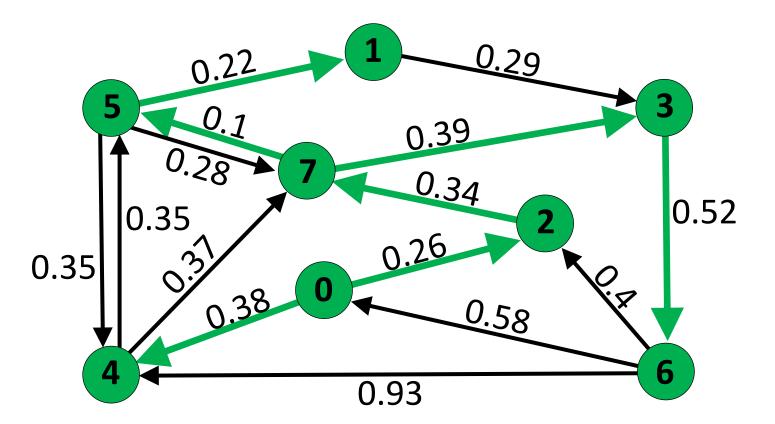
Repeat.





Repeat?



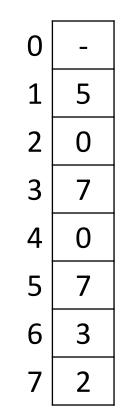


Distance from 0

0.60

Previous vertex

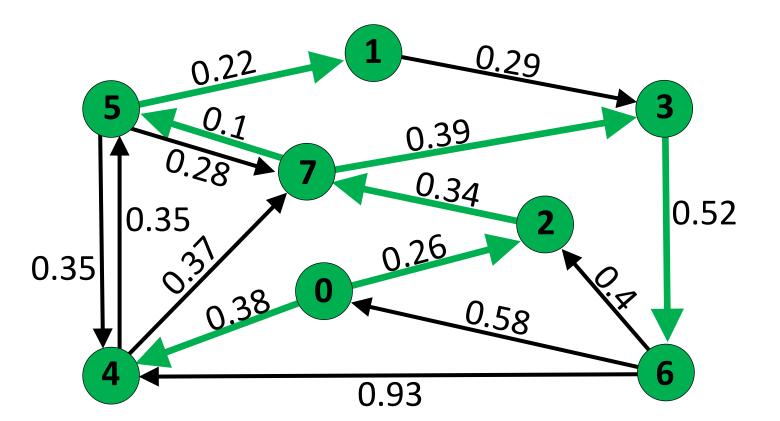
Priority queue





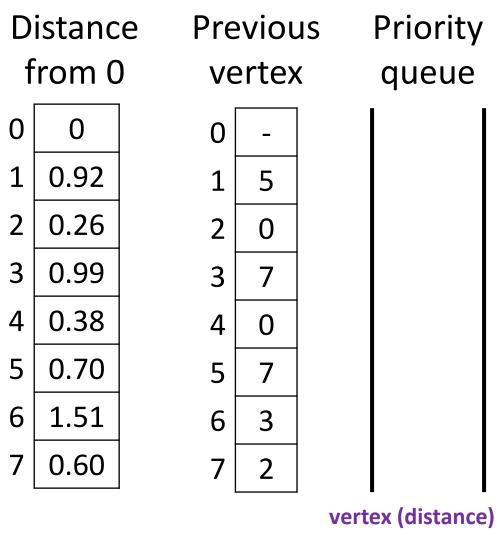
Repeat?

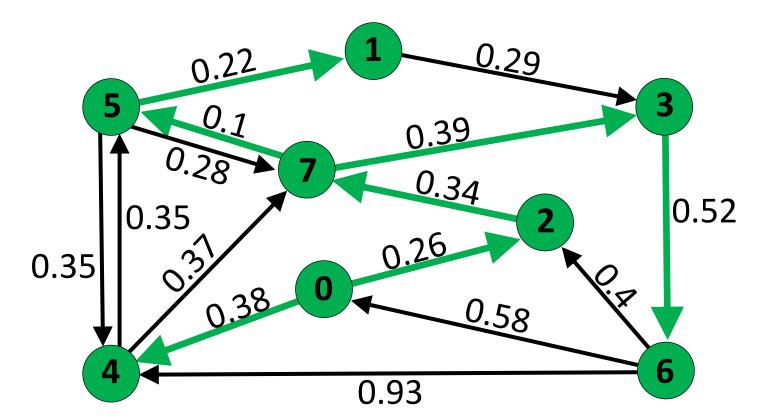
Neighbor 4? 1.51 + 0.93 > 0.38



Repeat?

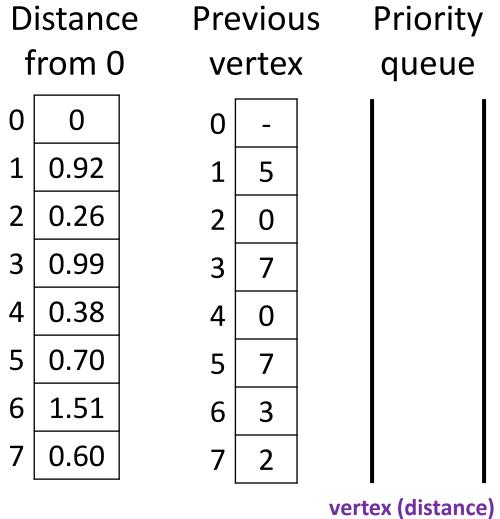
Neighbor 0? 1.51 + 0.58 > 0

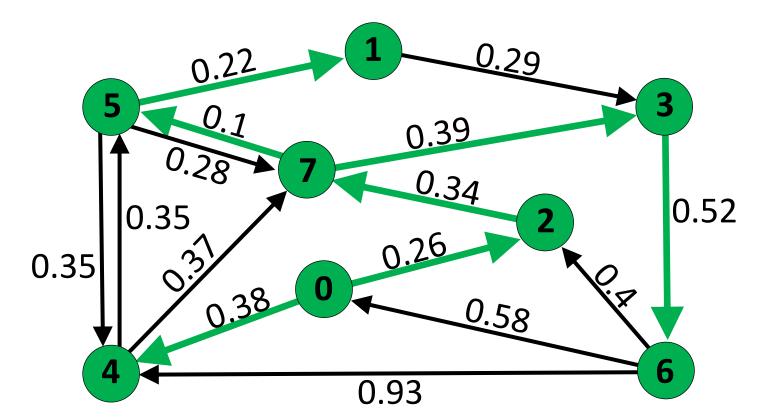




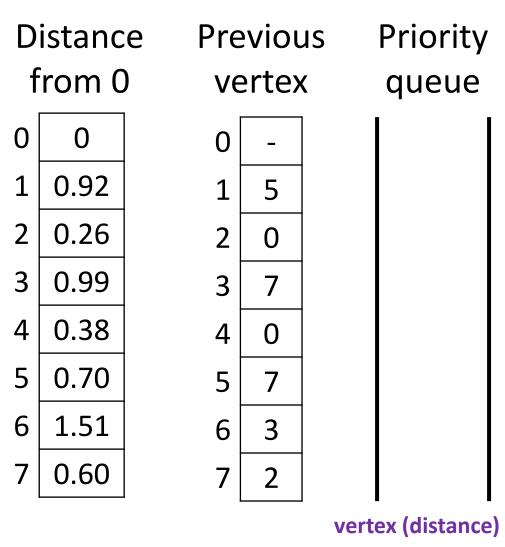
Repeat?

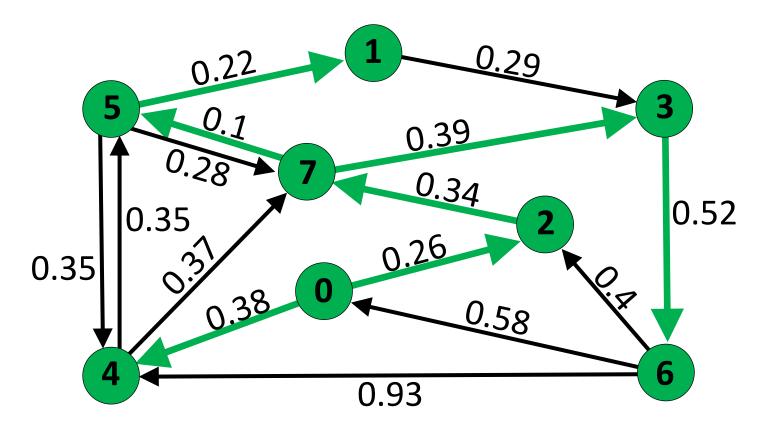
Neighbor 2? 1.51 + 0.4 > 0.26





When are we done?

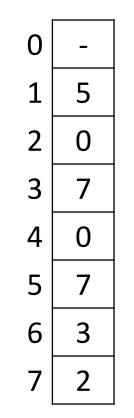




Distance from 0

Previous vertex

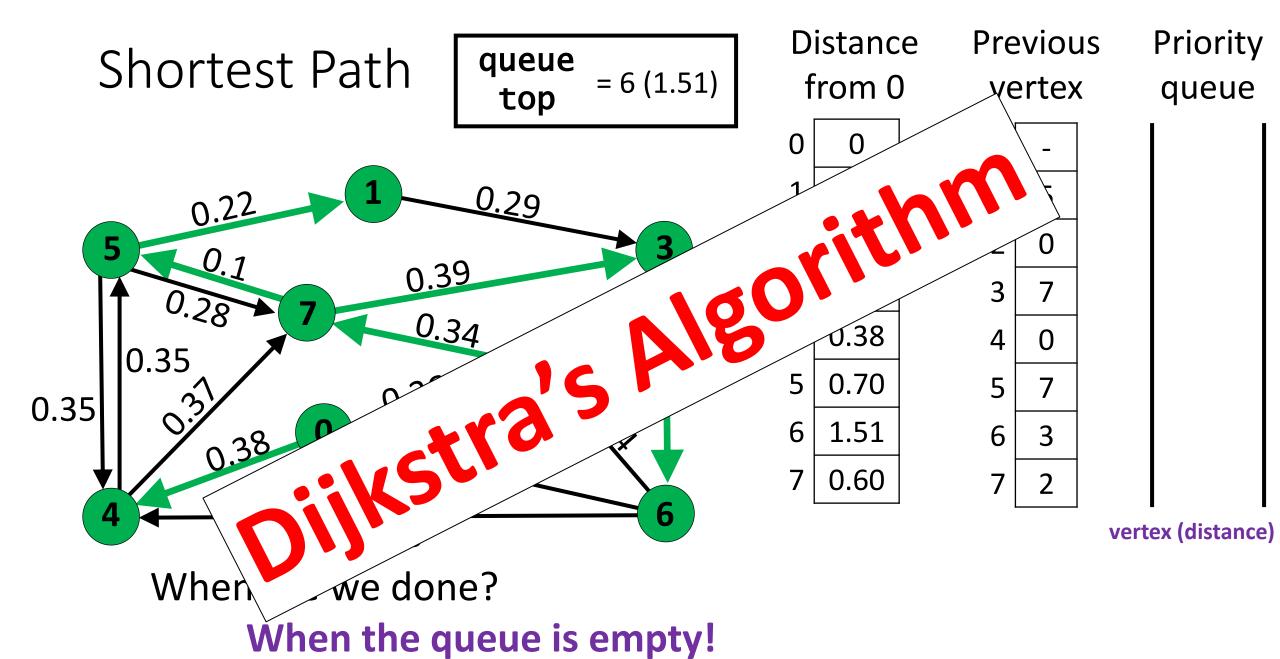
Priority queue

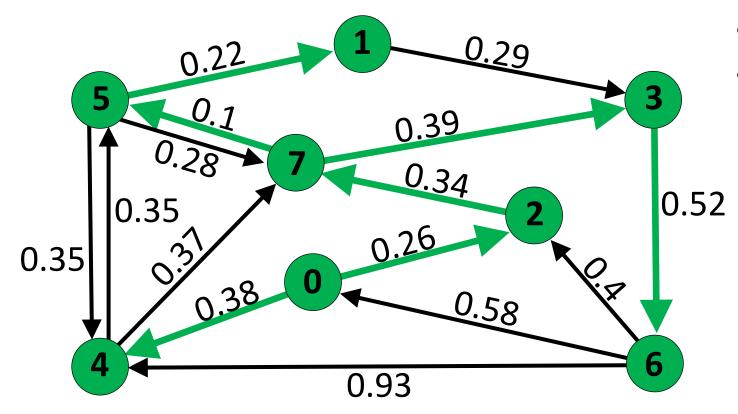


vertex (distance)

When are we done?

When the queue is empty!

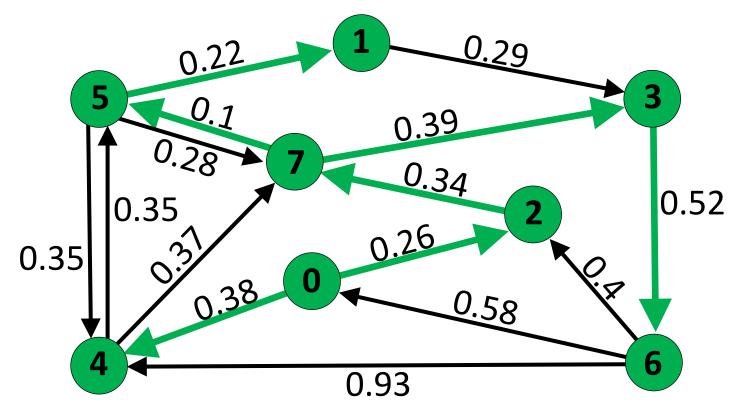




Assumptions:

- Graph is directed.
- Graph is edge-weighted.
- Edge weights are non-negative.
- Graph need not be simple (though our example will be).

What happens if there are self-loops?

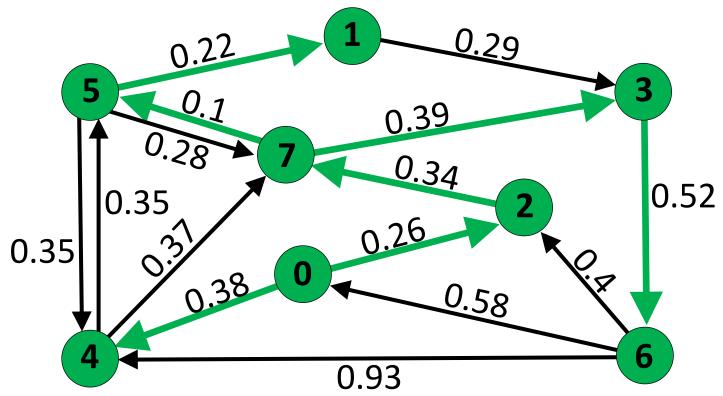


Assumptions:

- Graph is directed.
- Graph is edge-weighted.
- Edge weights are non-negative.
- Graph need not be simple (though our example will be).

What happens if there are self-loops?

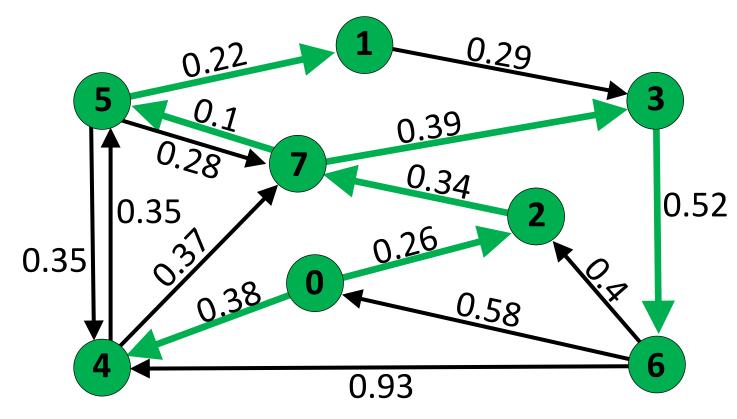
They are never taken, since they will never lower the cost of a path.



Assumptions:

- Graph is directed.
- Graph is edge-weighted.
- Edge weights are non-negative.
- Graph need not be simple (though our example will be).

What happens if there are parallel edges?

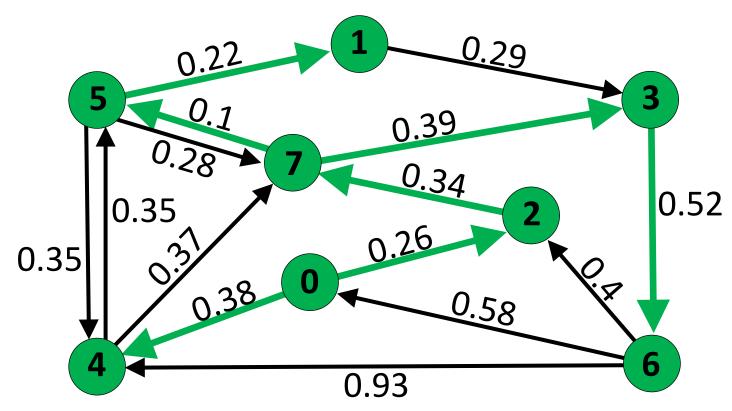


Assumptions:

- Graph is directed.
- Graph is edge-weighted.
- Edge weights are non-negative.
- Graph need not be simple (though our example will be).

What happens if there are parallel edges?

The cheapest one is taken and all others are ignored.



Assumptions:

- Graph is directed.
- Graph is edge-weighted.
- Edge weights are non-negative.
- Graph need not be simple (though our example will be).

What happens if there are negative weights?

```
public class Edge implements Comparable<Edge>{
                                         private int sourceVertex;
                                         private int destVertex;
                                         private double weight;
                                         public Edge(int vertex1, int vertex2, double weight) {
                                                this.sourceVertex = vertex1;
                                                this.destVertex = vertex2;
                                                this.weight = weight;
                                                 0.52
       0.35
0.35
                         0.93
```