

CSCI 232:

Data Structures and Algorithms

Dynamic Programming (Part 3)

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Spring 2025

Rod Cutting

Given a rod of length n inches, and an array of prices that includes prices of all pieces of size smaller than n , determine the maximum value obtainable by cutting up the rod and selling the pieces.

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |



$n = 8$
(no cuts)

Total profit
\$20



$n = 2$

$n = 2$

$n = 2$

$n = 2$

Total profit
\$20



$n = 3$

Total profit
\$18

$n = 5$



$n = 2$

Total profit
\$22

Optimal profit!

$n = 6$

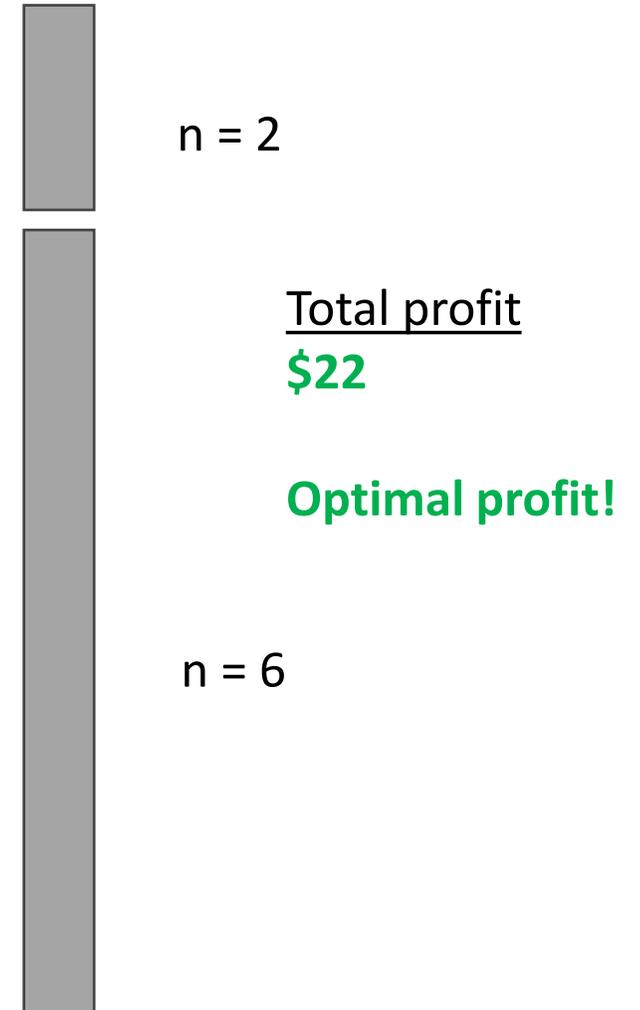
Rod Cutting

Given a rod of length n inches, and an array of prices that includes prices of all pieces of size smaller than n , determine the maximum value obtainable by cutting up the rod and selling the pieces.

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

Optimal Substructure

Our solution for a rod length of $n=8$, has the optimal solution for rod length of $n = 6$, and $n = 2$



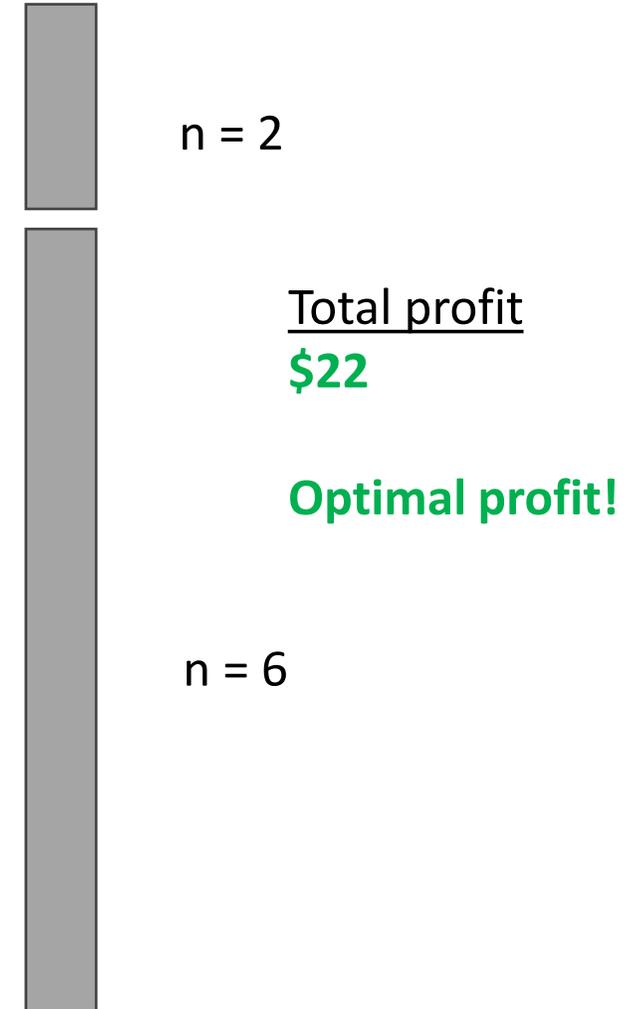
Rod Cutting

Given a rod of length n inches, and an array of prices that includes prices of all pieces of size smaller than n , determine the maximum value obtainable by cutting up the rod and selling the pieces.

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

General Approach:

Compute all possible ways to cut the rod using dynamic programming, and return which one had the highest profit



Rod Cutting

Given a rod of length n inches, and an array of prices that includes prices of all pieces of size smaller than n , determine the maximum value obtainable by cutting up the rod and selling the pieces.

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |



$n = 2$



$n = 2$

Total profit
\$20



$n = 2$



$n = 2$

Overlapping subproblems

We will compute the optimal way to cut a rod of length $n=2$ many times. We will use memoization to make sure we don't compute problems that we have already solved.

Rod Cutting

| | | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|---|
| | | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 | |

↑
index

n = 8



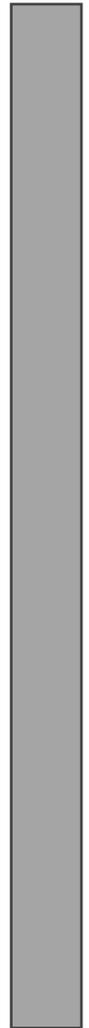
Technically, our algorithm will consider making a cut of length 8 first, but we will skip over this part to avoid confusion

Rod Cutting

| | | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|---|
| | | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 | |


index

n = 8



Rod Cutting

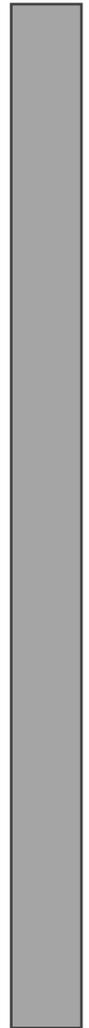
| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

0 1 2 3 4 5 6 7



index

n = 8



Two options

n = 8



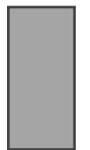
Don't Cut

n = 7



Make cut of length **index**

n = 1

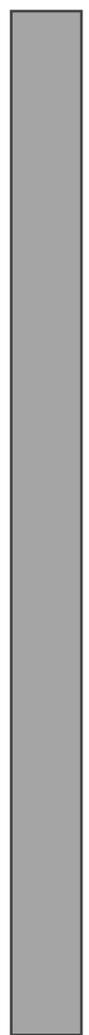


Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |


index

n = 8

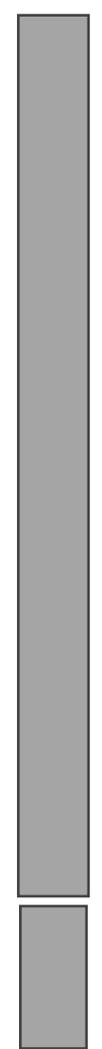


n = 8



Don't Cut

n = 7



n = 1

Make cut of length **index**

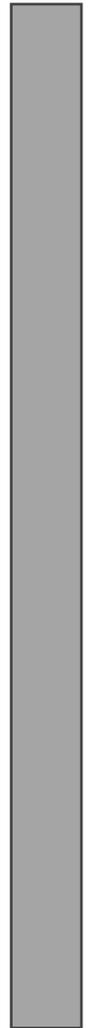
Two options

We want to select the option that yield the highest profit

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



Two options

n = 8



Don't Cut

Now we recurse, and check a new cut value



index



n = 7

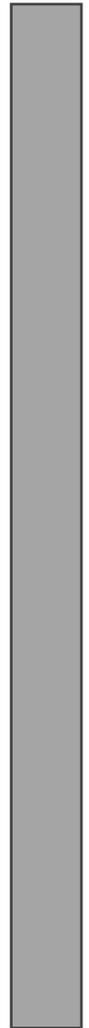
n = 1

Make cut of length **index**

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



Two options

n = 8



Don't Cut

Now we recurse, and check a new cut value

(index - 1)



index



n = 7

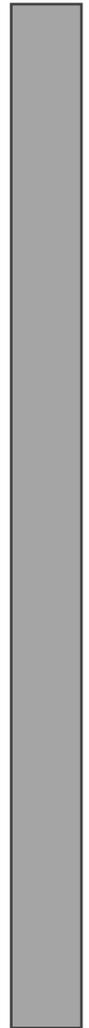
n = 1

Make cut of length **index**

Rod Cutting

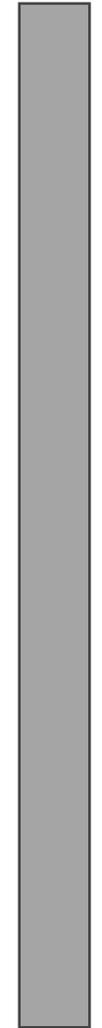
| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



Two options

n = 8



Don't Cut

n = 2



n = 6

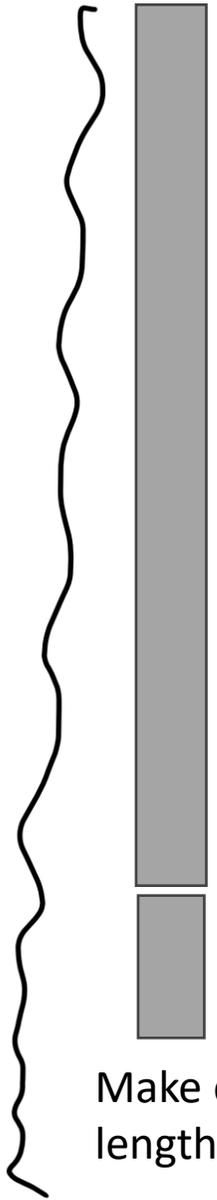


Make cut of length **index**



index

n = 7



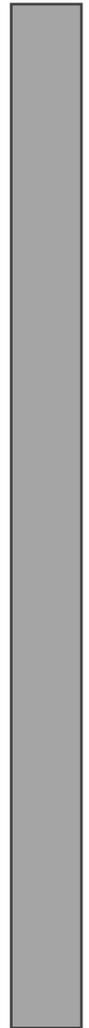
n = 1

Make cut of length **index**

Rod Cutting

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



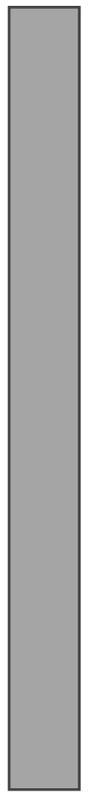
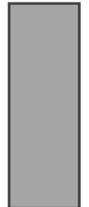
Two options

n = 8



Don't Cut

n = 2



Make cut of length **index**

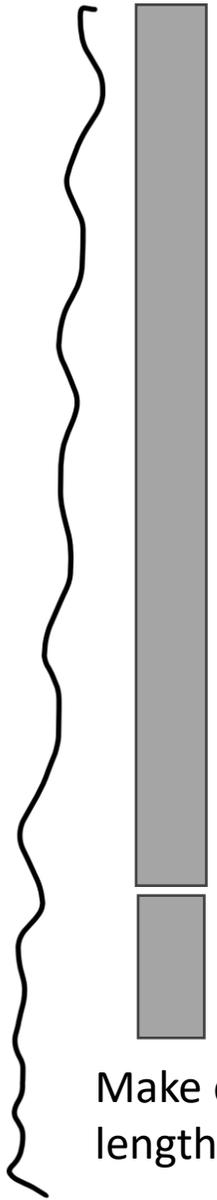
n = 6



index

We want to select the option that yield the highest profit

n = 7



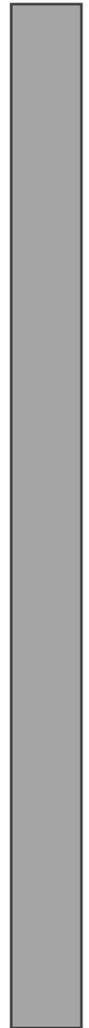
n = 1

Make cut of length **index**

Rod Cutting

| | | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|---|
| | | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 | |

n = 8



Two options

n = 8



Don't Cut



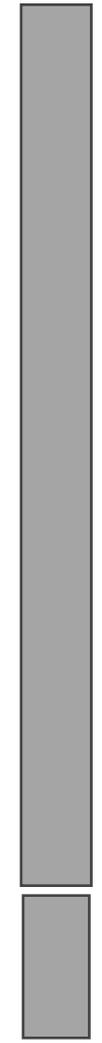
index



n = 2

n = 6

Make cut of length **index**



n = 7

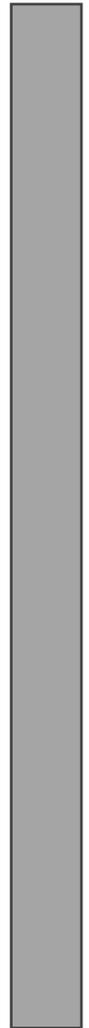
n = 1

Make cut of length **index**

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



Two options

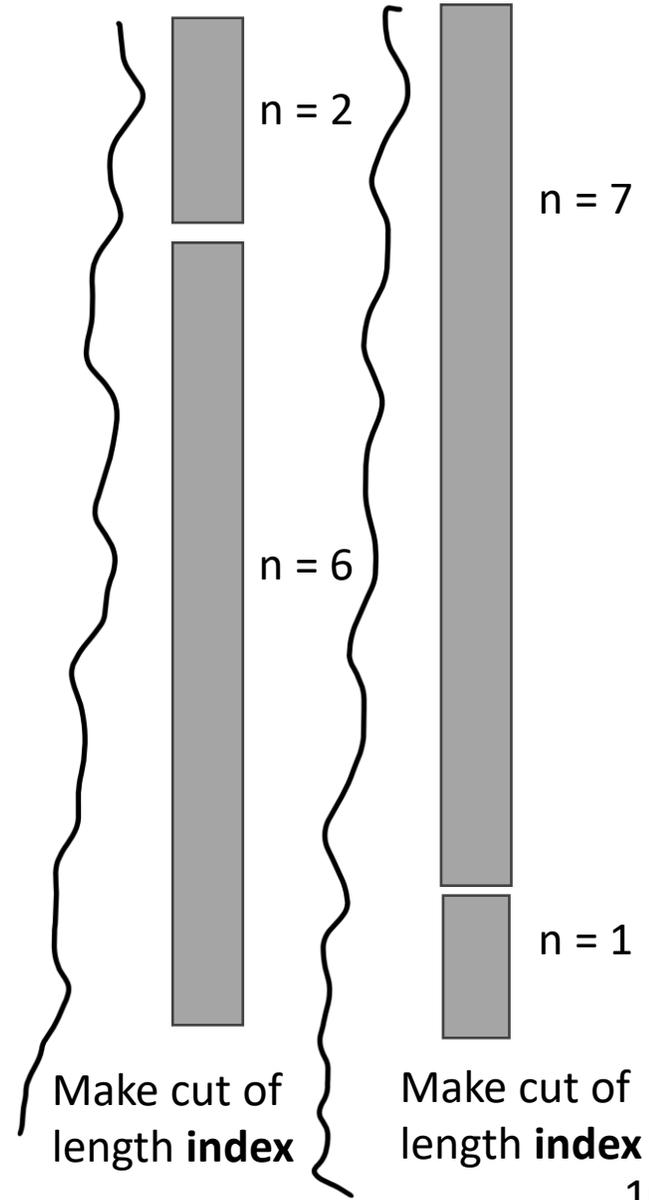
n = 8



Don't Cut



index



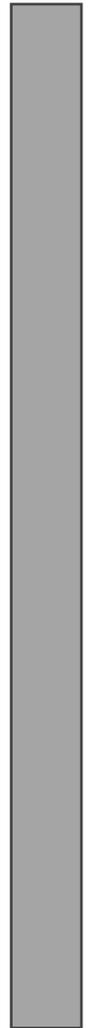
Make cut of length **index**

Make cut of length **index**

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8

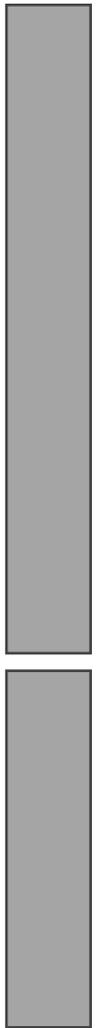


Two options

n = 8



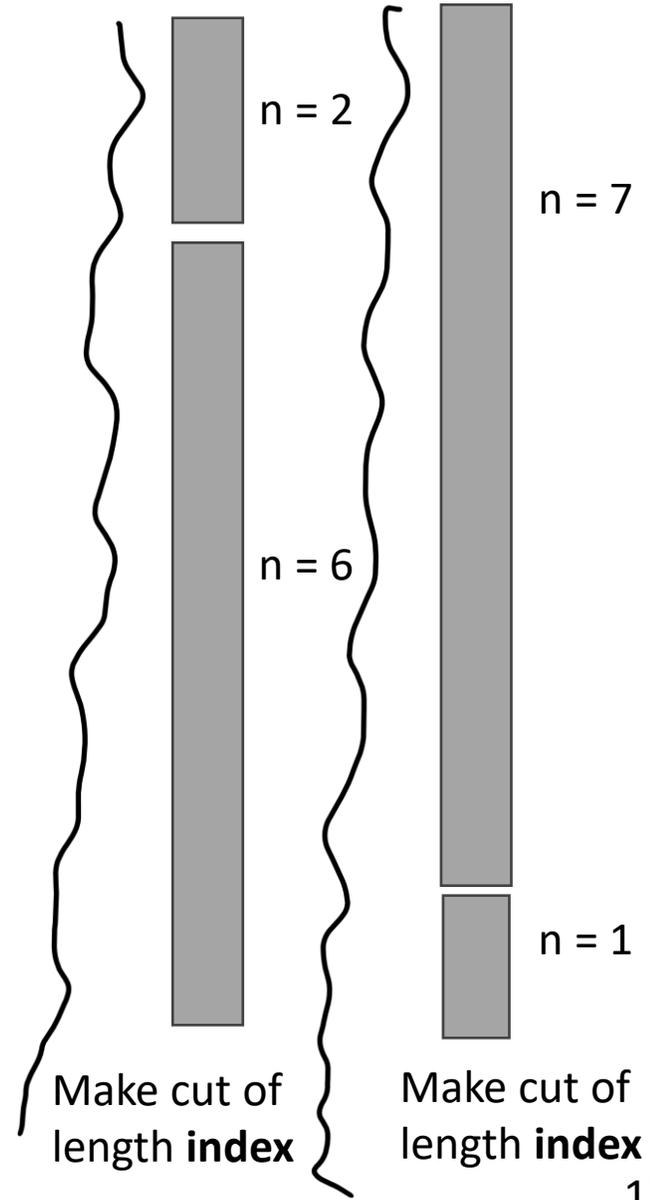
Don't Cut



N = 5

N = 3

Make cut of length **index**



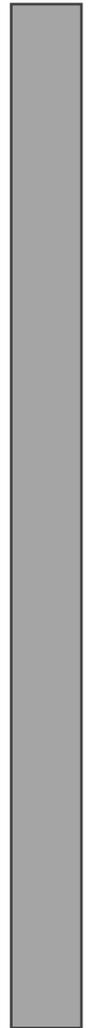
Make cut of length **index**

Make cut of length **index**

Rod Cutting

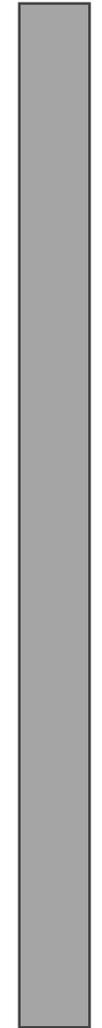
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|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



Two options

n = 8

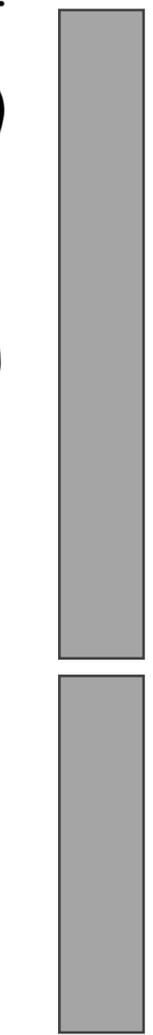


Don't Cut

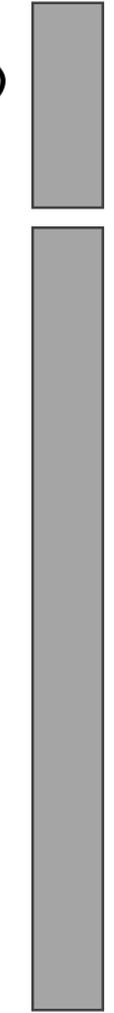
Whenever we don't make the cut, we don't adjust the size of the rod, but we check the next cut length



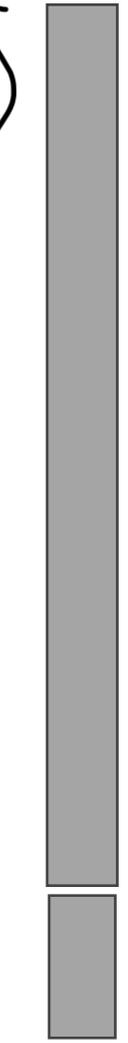
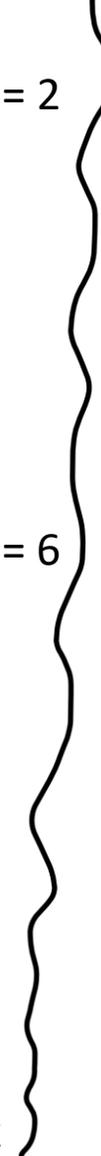
index



Make cut of length **index**



Make cut of length **index**



Make cut of length **index**

Rod Cutting

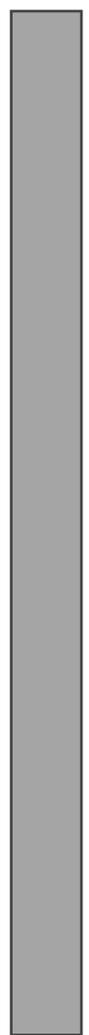
| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

0 1 2 3 4 5 6 7



index

n = 8



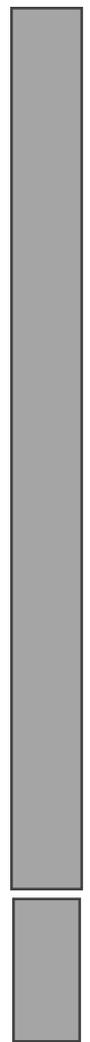
Two options

n = 8



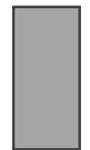
Don't Cut

n = 7



Make cut of length **index**

n = 1



Rod Cutting

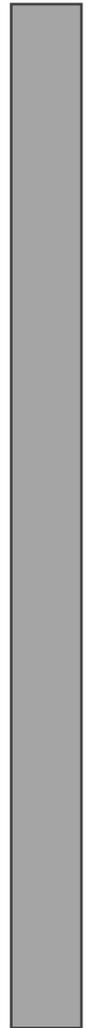
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0 1 2 3 4 5 6 7



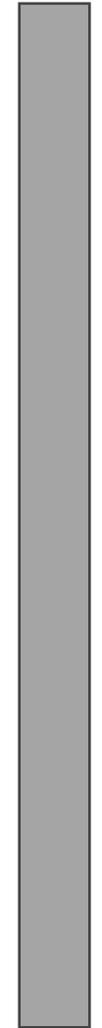
index

n = 8



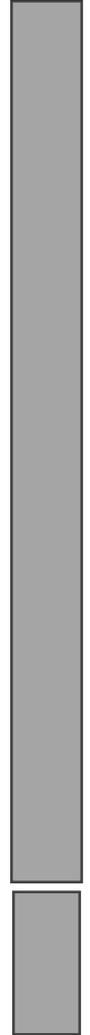
Two options

n = 8

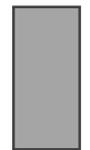


Don't Cut

n = 7



n = 1



Make cut of length **index**

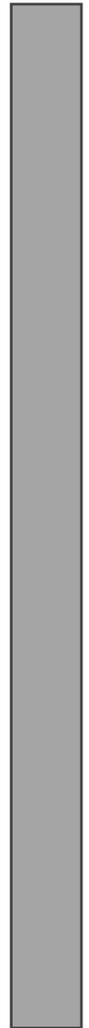
We made a cut of length index, so let's figure out how much that piece is worth!

`prices[index]`

Rod Cutting

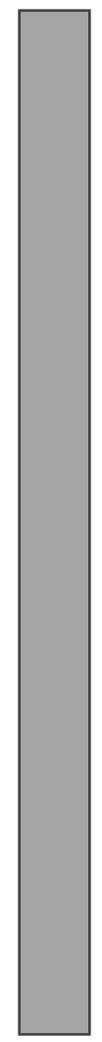
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|--------|---|---|---|---|----|----|----|----|
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| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

n = 8



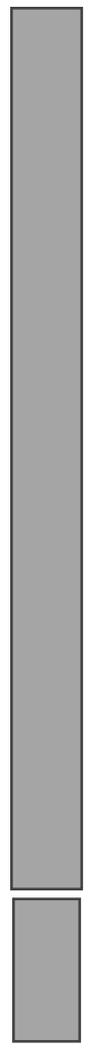
Two options

n = 8



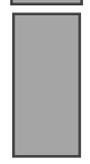
Don't Cut

n = 7



Make cut of length **index**

n = 1



We made a cut of length index, so let's figure out how much that piece is worth!

`prices[index]`

We have 1 inch of rod left, so we need to now figure out the optimal way to cut this



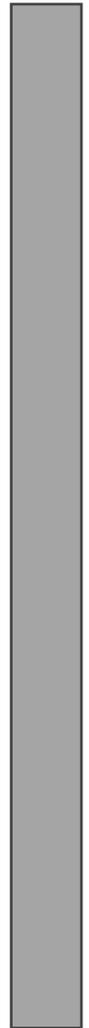
index

Rod Cutting

| | | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|---|
| | | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 | |

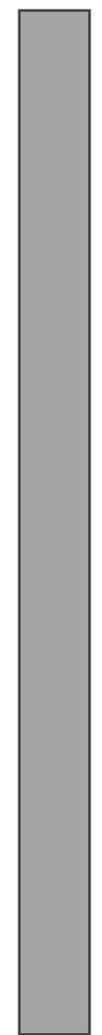


n = 8



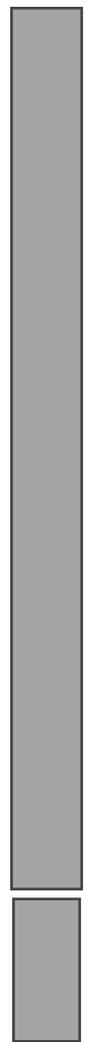
Two options

n = 8



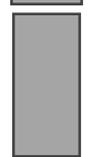
Don't Cut

n = 7



Make cut of length **index**

n = 1



We made a cut of length index, so let's figure out how much that piece is worth!

`prices[index]`

Length of cut made = (index + 1)

We have 1 inch of rod left, so we need to now figure out the optimal way to cut this --Recurse!

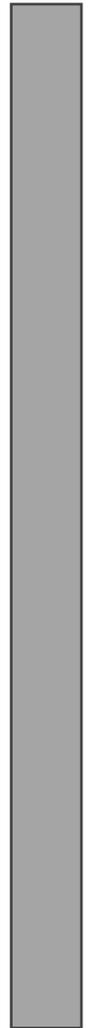
Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
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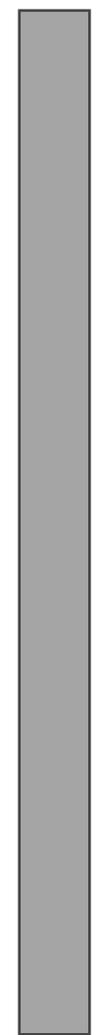
index

n = 8



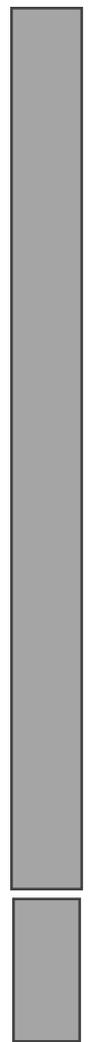
Two options

n = 8



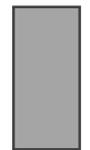
Don't Cut

n = 7



Make cut of length **index**

n = 1



We made a cut of length index, so lets figure out how much that piece is worth!

`prices[index]`

Length of cut made = $(index + 1)$

New subproblem = $n - \text{length_of_cut}$

We have 1 inch of rod left, so we need to now figure out the optimal way to cut this --Recurse!

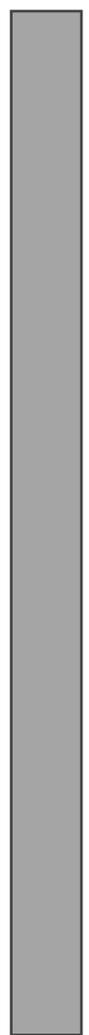
Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |



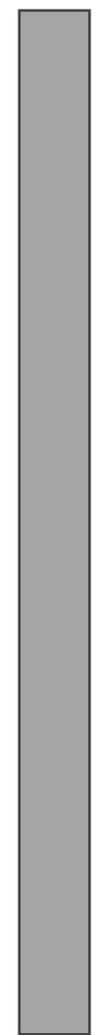
index

n = 8



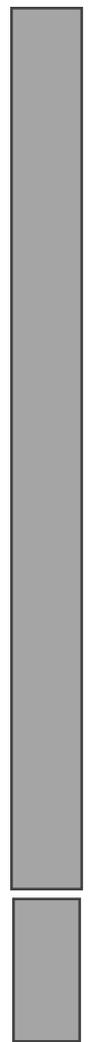
Two options

n = 8



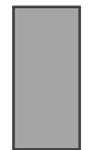
Don't Cut

n = 7



Make cut of length **index**

n = 1



We made a cut of length index, so lets figure out how much that piece is worth!

`prices[index]`

Length of cut made = $(index + 1)$

New subproblem = $n - \text{length_of_cut}$

We have 1 inch of rod left, so we need to now figure out the optimal way to cut this --Recurse!

Rod Cutting

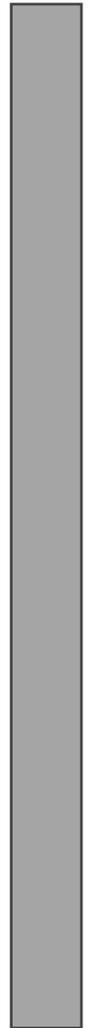
| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

0 1 2 3 4 5 6 7



index

n = 8



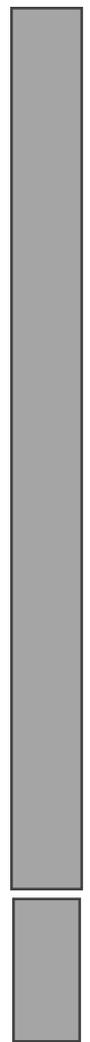
Two options

n = 8



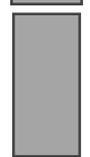
Don't Cut

n = 7



Make cut of length **index**

n = 1



Whenever we make the cut, we adjust the size of the rod, but keep the same index

Rod Cutting

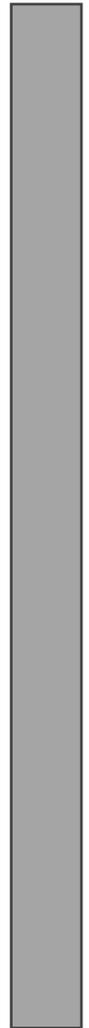
| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

0 1 2 3 4 5 6 7



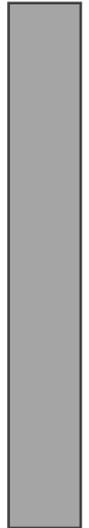
index

n = 8

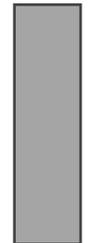


Two options

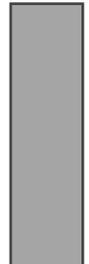
n = 4



Don't Cut



n = 2



n = 2

Make cut of length **index**

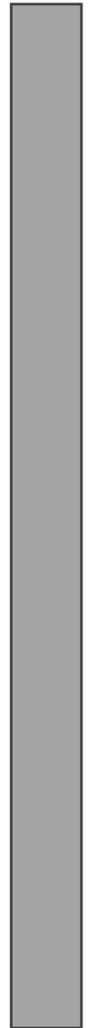
Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |



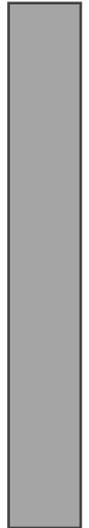
index

n = 8



Two options

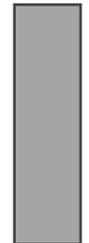
n = 4



Don't Cut

Profit: 9

n = 2



n = 2



Make cut of length **index**

Profit: 10

Given a rod of length 4 and a potential cut value of length 2, the optimal solution is to **make the cut**

Rod Cutting

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |



index

n = 8



Two options

n = 4



Don't Cut

Profit: 9

n = 2



n = 2



Make cut of length **index**

Profit: 10

Given a rod of length 4 and a potential cut value of length 2, the optimal solution is to **make the cut**

If we ever encounter this same subproblem again, we want to make sure we don't recompute it

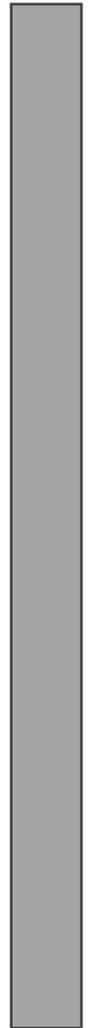
Rod Cutting

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

0 1 2 3 4 5 6 7

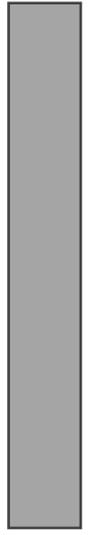
↑
index

n = 8



Two options

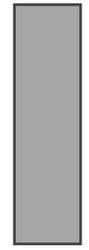
n = 4



Don't Cut

Profit: 9

n = 2



n = 2



Make cut of length **index**

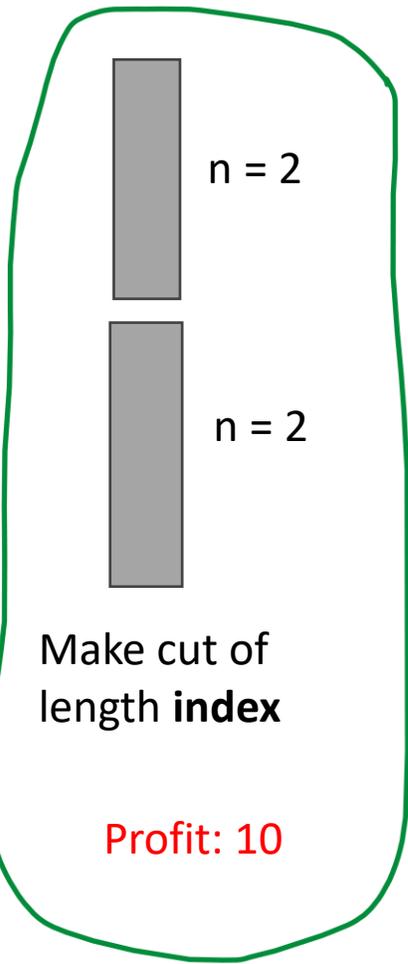
Profit: 10

We need to put this solution (10) into our memorization table

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

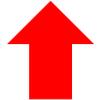

index

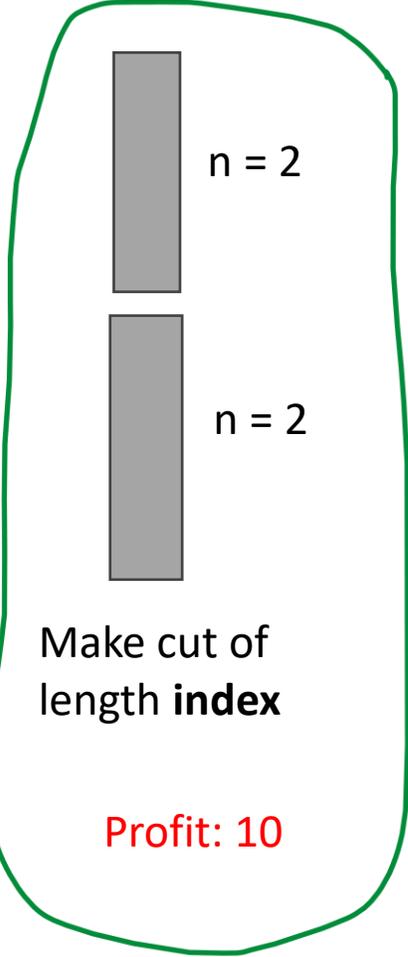


| | | Rod Length | | | | | | | |
|------------|---|------------|---|---|---|---|---|---|---|
| | | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Cut Length | 1 | | | | | | | | |
| | 2 | | | | | | | | |
| | 3 | | | | | | | | |
| | 4 | | | | | | | | |
| | 5 | | | | | | | | |
| | 6 | | | | | | | | |
| | 7 | | | | | | | | |
| | 8 | | | | | | | | |

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | <i>0</i> | <i>1</i> | <i>2</i> | <i>3</i> | <i>4</i> | <i>5</i> | <i>6</i> | <i>7</i> |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |


index



Rod Length

| | | | | | | | | |
|----------|----------|----------|----------|----------|----------|----------|----------|----------|
| | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| 1 | | | | | | | | |
| 2 | | | | | | | | |
| 3 | | | | | | | | |
| 4 | | | | | | | | |
| 5 | | | | | | | | |
| 6 | | | | | | | | |
| 7 | | | | | | | | |
| 8 | | | | | | | | |

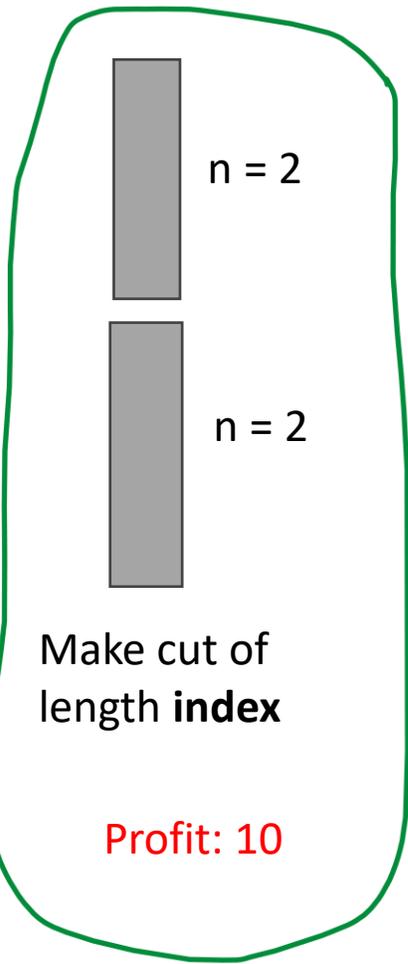
Cut Length

Rod Cutting

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

↑
index

dp[index][n] = 10



} n = 4

Cut Length

Rod Length

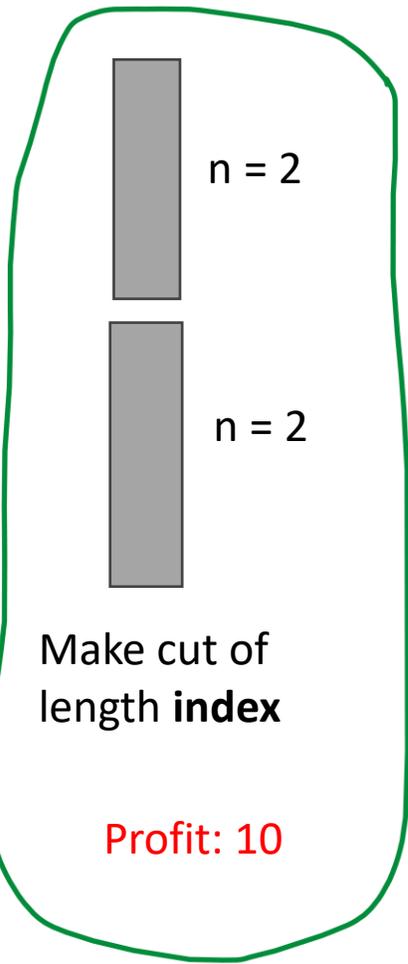
| | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---|---|---|---|----|---|---|---|---|
| 1 | | | | | | | | |
| 2 | | | | 10 | | | | |
| 3 | | | | | | | | |
| 4 | | | | | | | | |
| 5 | | | | | | | | |
| 6 | | | | | | | | |
| 7 | | | | | | | | |
| 8 | | | | | | | | |

Rod Cutting

| | | | | | | | | |
|--------|---|---|---|---|----|----|----|----|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |

↑
index

dp[index][n] = 10



} n = 4

Cut Length

Rod Length

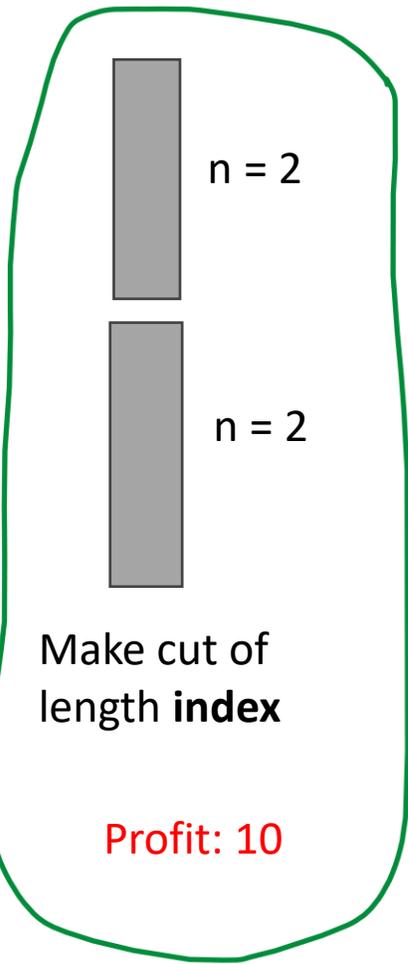
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---|---|---|---|---|----|---|---|---|---|
| 0 | | | | | | | | | |
| 1 | | | | | 10 | | | | |
| 2 | | | | | | | | | |
| 3 | | | | | | | | | |
| 4 | | | | | | | | | |
| 5 | | | | | | | | | |
| 6 | | | | | | | | | |
| 7 | | | | | | | | | |

Rod Cutting

| | | | | | | | | |
|--------|----------|----------|----------|----------|----------|----------|----------|----------|
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| Length | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| Price | 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |



dp[index][n] = 10



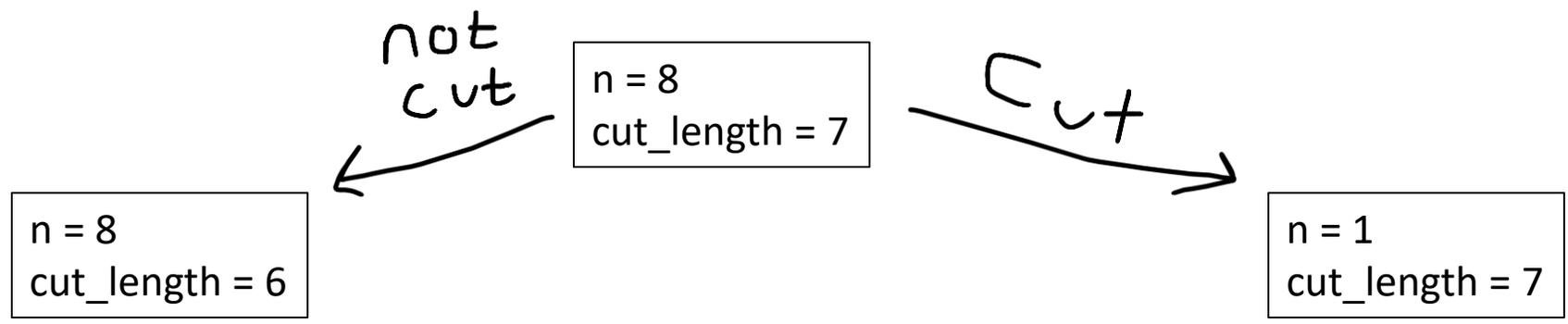
Rod Length

Cut Length

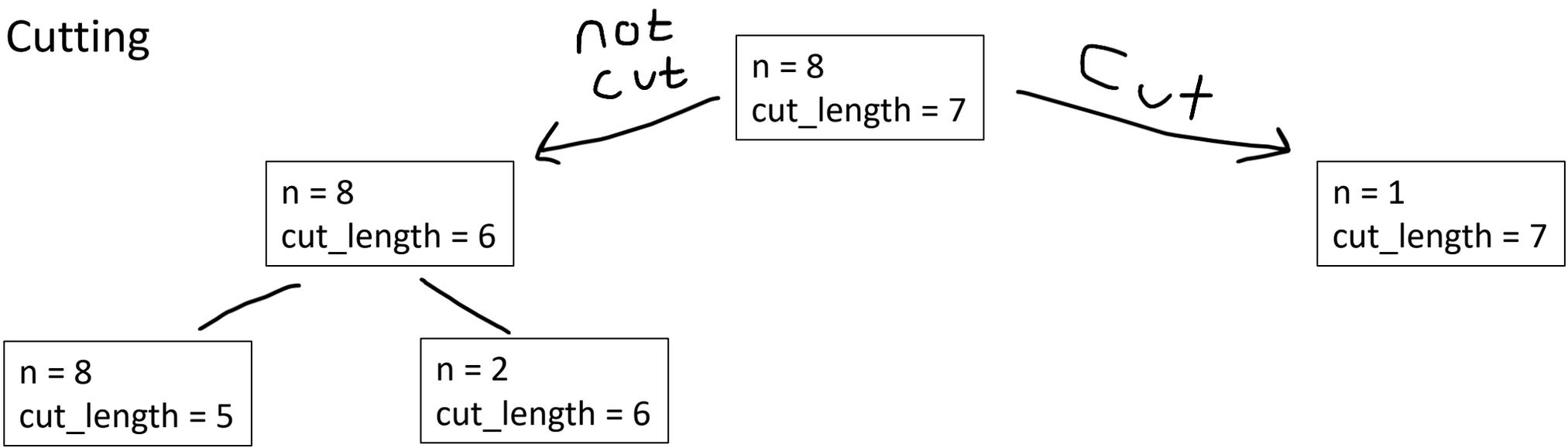
| | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---|---|---|---|----|---|---|---|---|
| 1 | | | | | | | | |
| 2 | | | | 10 | | | | |
| 3 | | | | | | | | |
| 4 | | | | | | | | |
| 5 | | | | | | | | |
| 6 | | | | | | | | |
| 7 | | | | | | | | |
| 8 | | | | | | | | |

Whenever we solve a subproblem, remember to place it inside of our memoization table

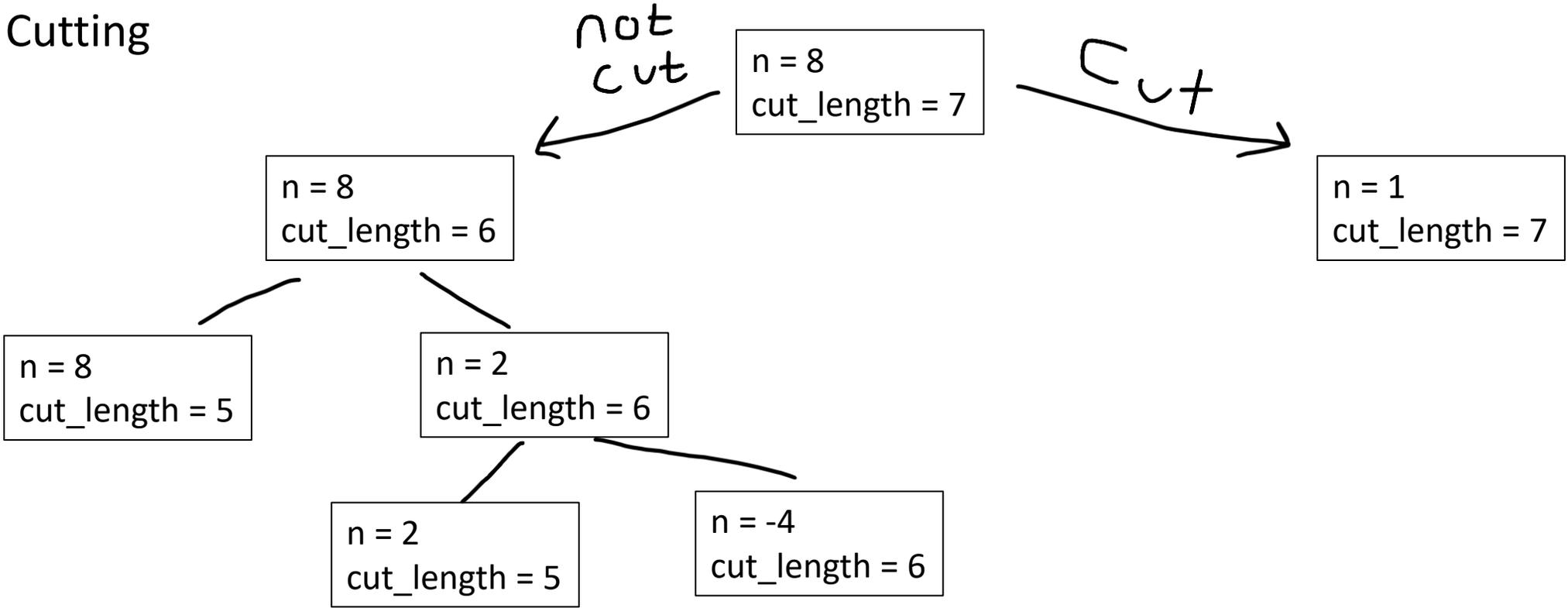
Rod Cutting



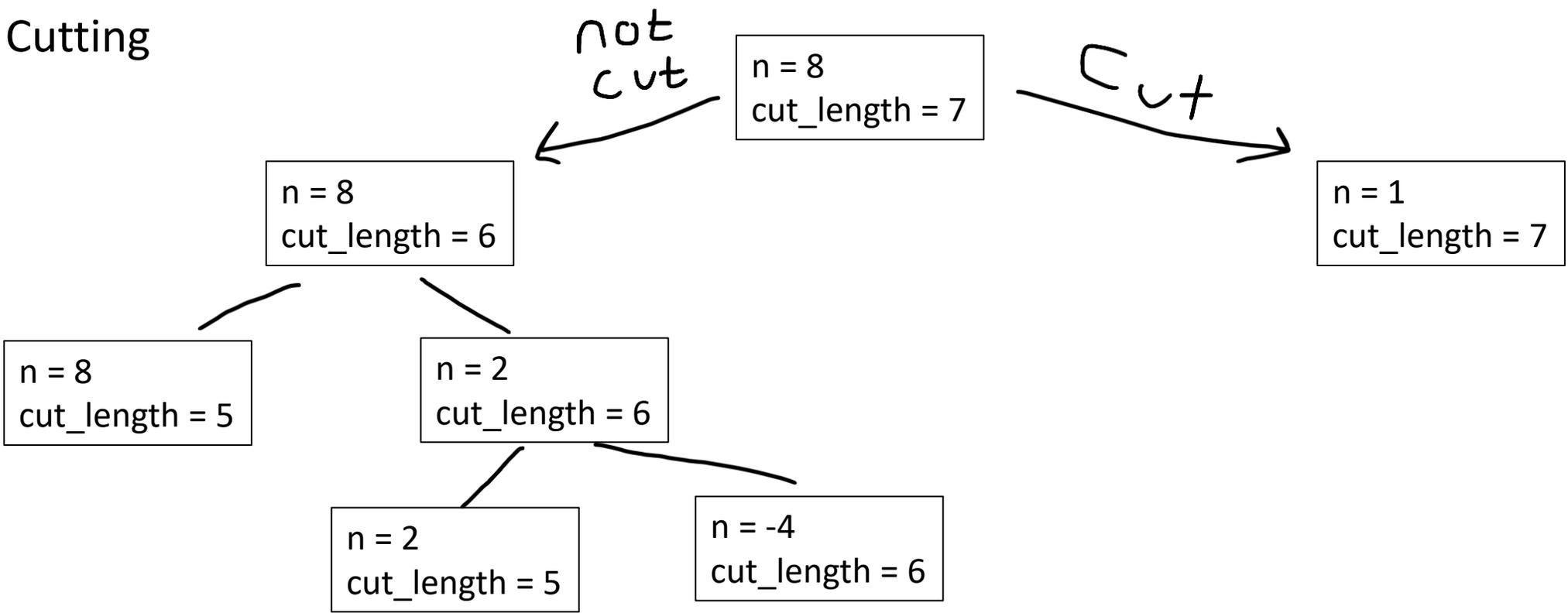
Rod Cutting



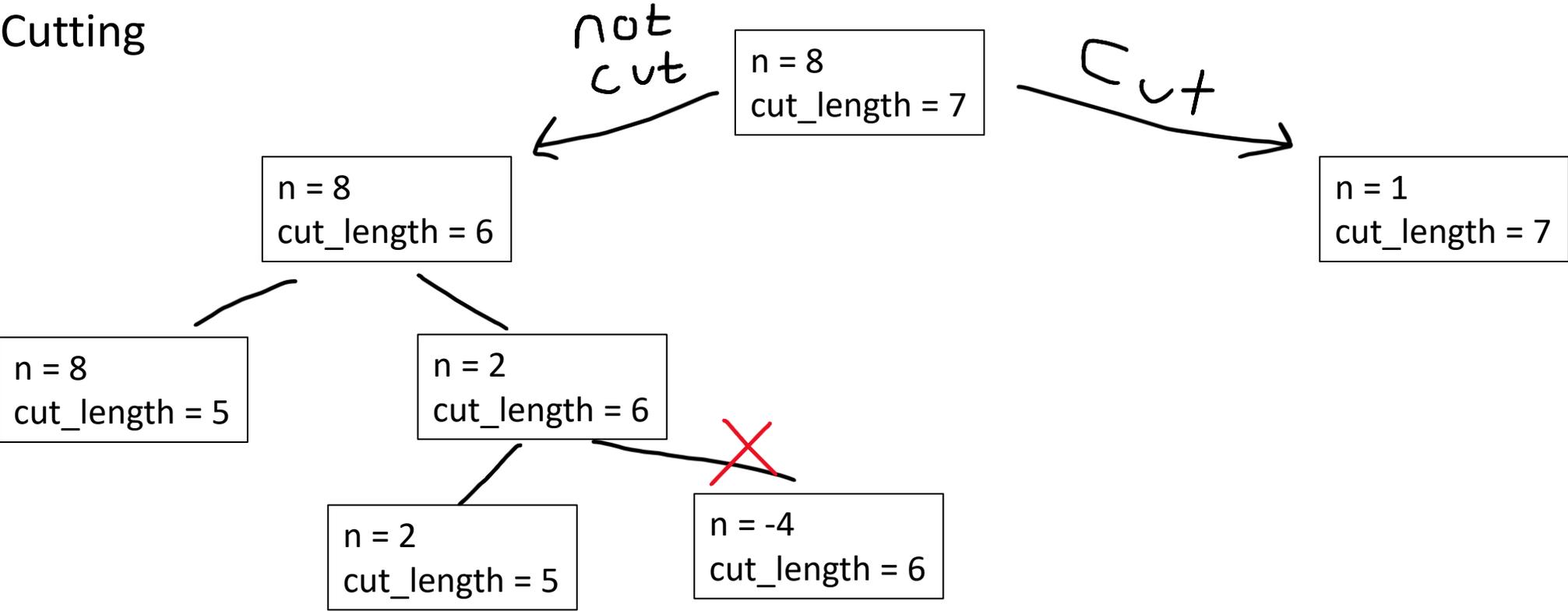
Rod Cutting



Rod Cutting

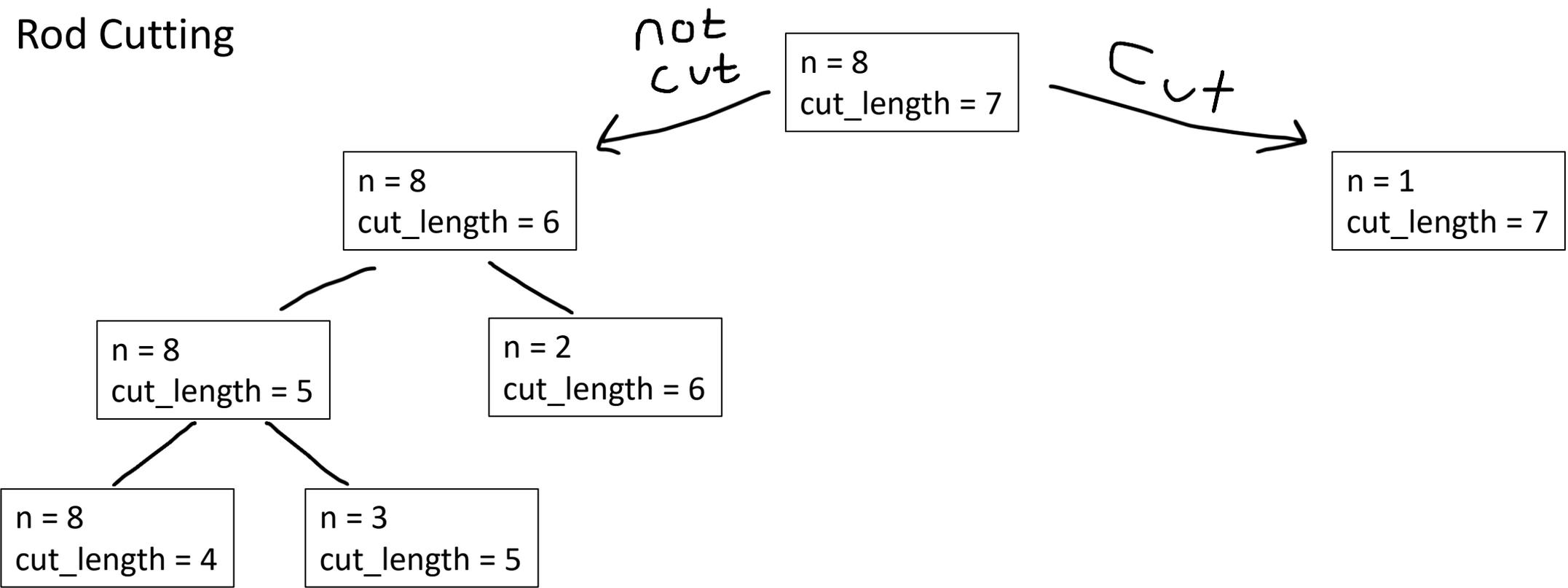


Rod Cutting

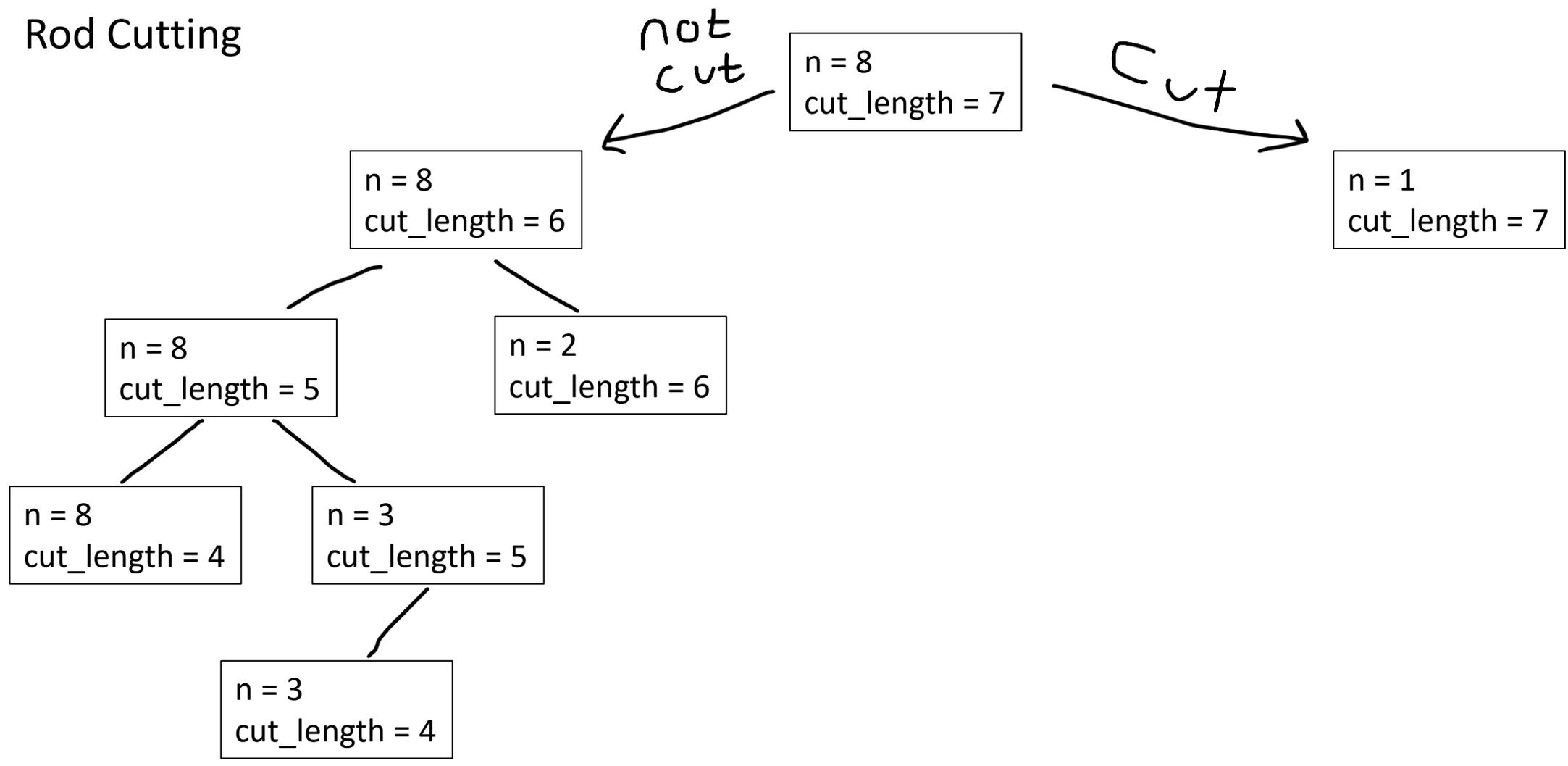


Only make the cut if its possible

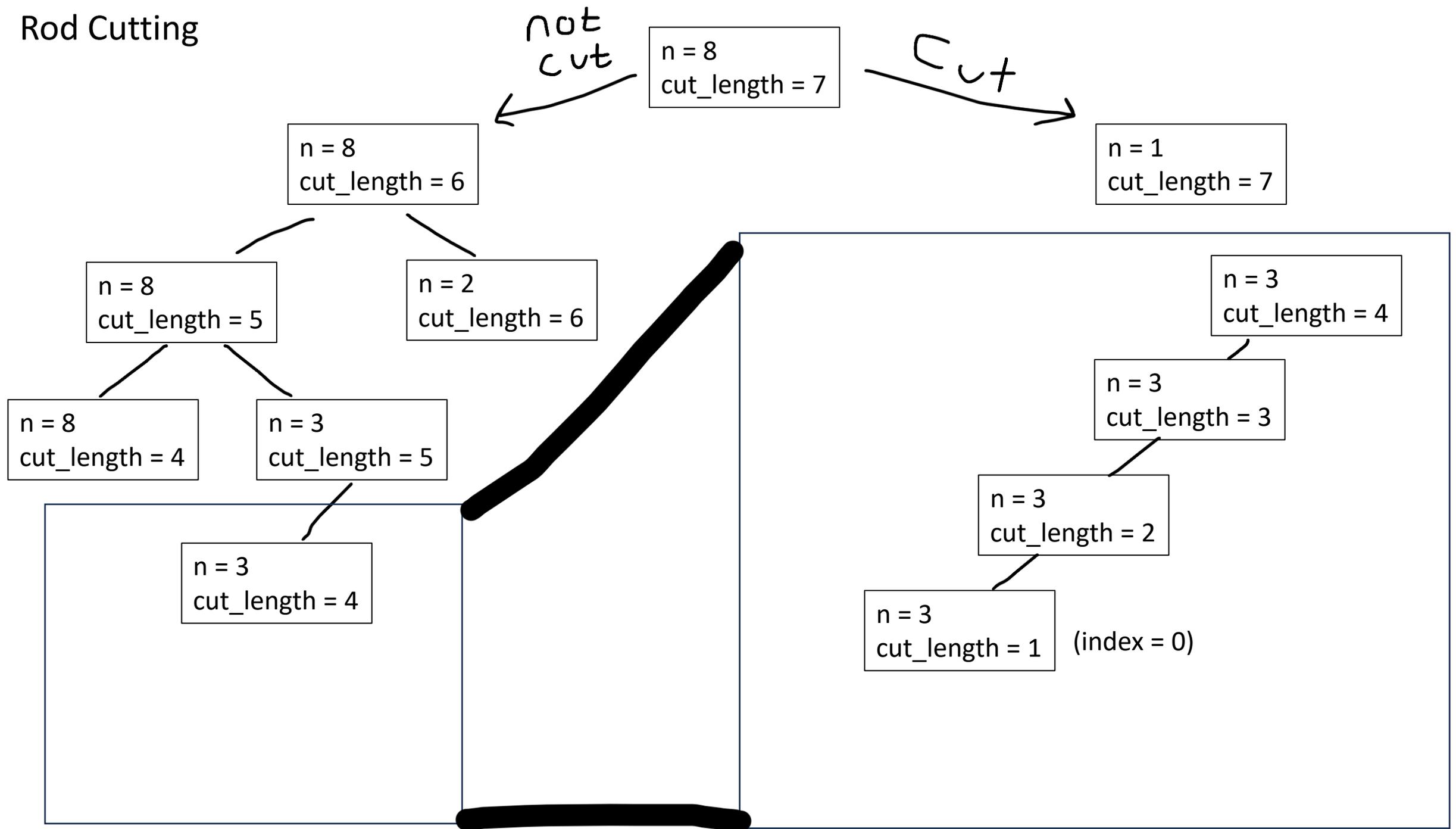
Rod Cutting



Rod Cutting



Rod Cutting



Rod Cutting

not cut

cut

n = 8
cut_length = 7

n = 8
cut_length = 6

n = 1
cut_length = 7

n = 8
cut_length = 5

n = 2
cut_length = 6

n = 8
cut_length = 4

n = 3
cut_length = 5

n = 3
cut_length = 4

n = 3
cut_length = 3

n = 3
cut_length = 2

n = 3
cut_length = 4

n = 3
cut_length = 1 (index = 0)

We can't chop into lengths less than 1, so we can compute a solution right here

Rod Cutting

not cut

Cut

n = 8
cut_length = 7

n = 8
cut_length = 6

n = 1
cut_length = 7

n = 8
cut_length = 5

n = 2
cut_length = 6

n = 8
cut_length = 4

n = 3
cut_length = 5

| | |
|--------|---|
| Length | 1 |
| Price | 1 |

↑ index

n = 3
cut_length = 4

n = 3
cut_length = 3

n = 3
cut_length = 2

n = 3
cut_length = 1 (index = 0)

Profit made into chopping into rods of length 1:
- n * prices[0] = 3

n = 3
cut_length = 4

Rod Cutting

not cut

Cut

n = 8
cut_length = 7

n = 8
cut_length = 6

n = 1
cut_length = 7

n = 8
cut_length = 5

n = 2
cut_length = 6

n = 8
cut_length = 4

n = 3
cut_length = 5

| | |
|--------|---|
| Length | 1 |
| Price | 1 |

↑ index

n = 3
cut_length = 4

n = 3
cut_length = 3

n = 3
cut_length = 2

n = 3
cut_length = 1 (index = 0)

Profit made into chopping into rods of length 1:
- $n * \text{prices}[0] = 3$ This will be our base case

n = 3
cut_length = 4

LETS TRY TO CODE THIS

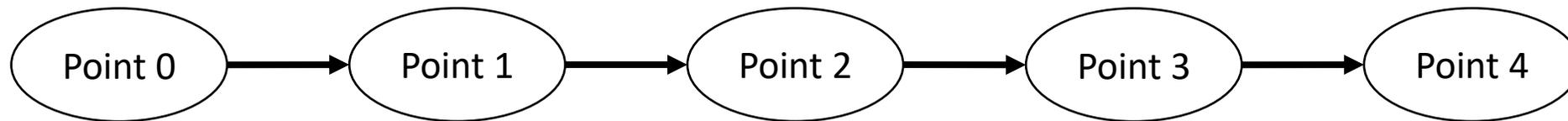
If you are confused are the recursion is set up, don't stress out about it. Its not a big deal.

n = 8
cut_l

The goal here is to show how we are using dynamic programming to solve this problem

Taxi Profit

Given a street that goes from point 1 to point N, we are able to pick up customers on the street and take them to the other end of the street. Our taxi is only able to go one direction

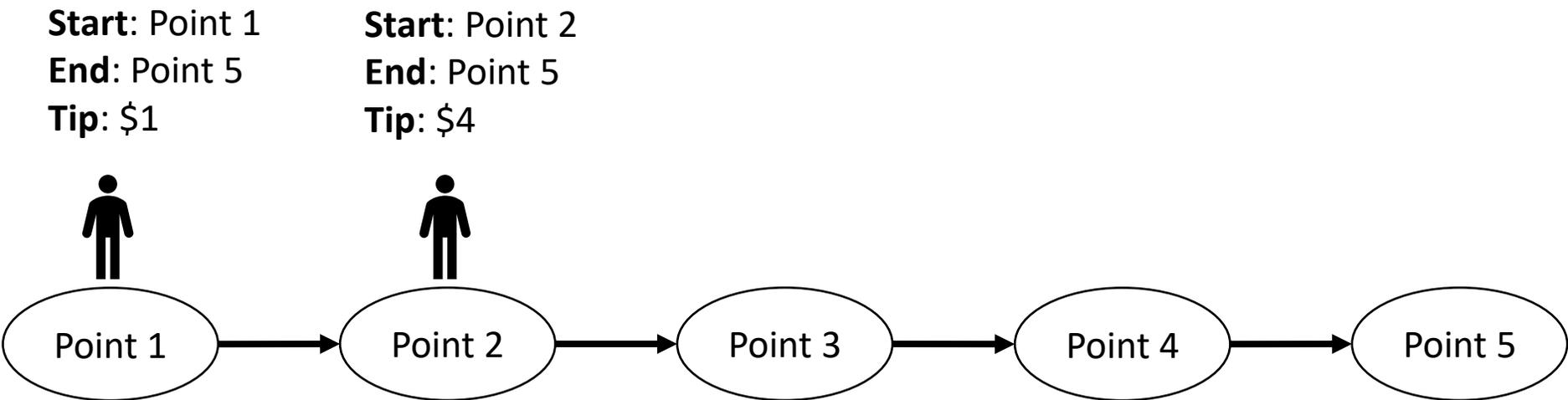


Taxi Profit

Given a street that goes from point 1 to point N, we are able to pick up customers on the street and take them to the other end of the street. Our taxi is only able to go one direction

Each customer has their starting point (pick-up spot), their ending point (destination), and the tip you will receive if you pick them up

The profit for selecting a customer is $(END - START) + TIP$



Taxi Profit

Given a street that goes from point 1 to point N, we are able to pick up customers on the street and take them to the other end of the street. Our taxi is only able to go one direction

Each customer has their starting point (pick-up spot), their ending point (destination), and the tip you will receive if you pick them up

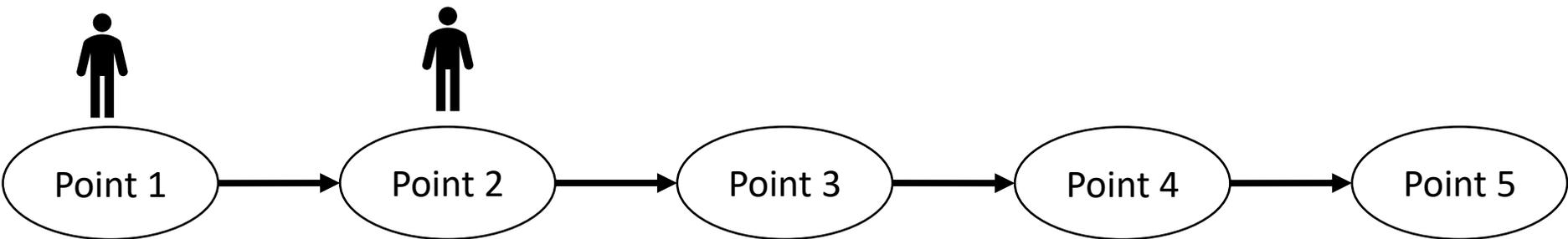
The profit for selecting a customer is $(END - START) + TIP$

$$(5 - 1) + 1 = 5$$

$$(5 - 2) + 4 = 7$$

Start: Point 1
End: Point 5
Tip: \$1

Start: Point 2
End: Point 5
Tip: \$4



Taxi Profit

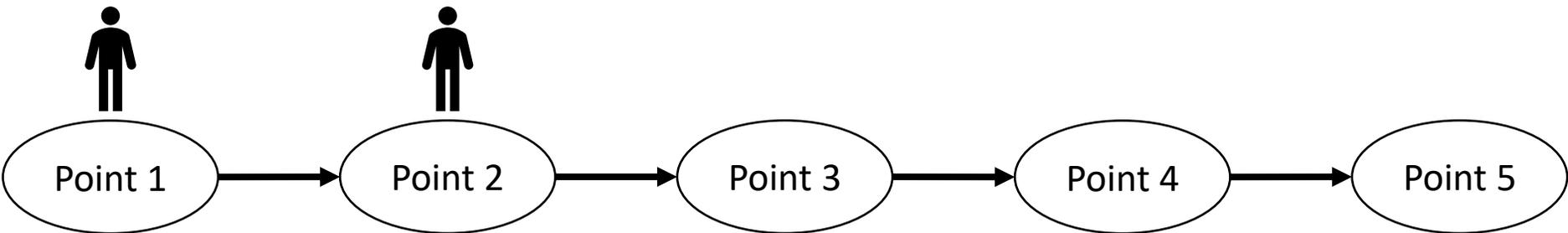
Given a street that goes from point 1 to point N, we are able to pick up customers on the street and take them to the other end of the street. Our taxi is only able to go one direction

Each customer has their starting point (pick-up spot), their ending point (destination), and the tip you will receive if you pick them up

The profit for selecting a customer is $(END - START) + TIP$

```
rides = [[1, 5, 1]      [2, 5, 4]      ]
```

Start: Point 1 **Start:** Point 2
End: Point 5 **End:** Point 5
Tip: \$1 **Tip:** \$4



What is the **maximum profit** that the taxi can make when going from point 1 to point N?

Taxi Profit

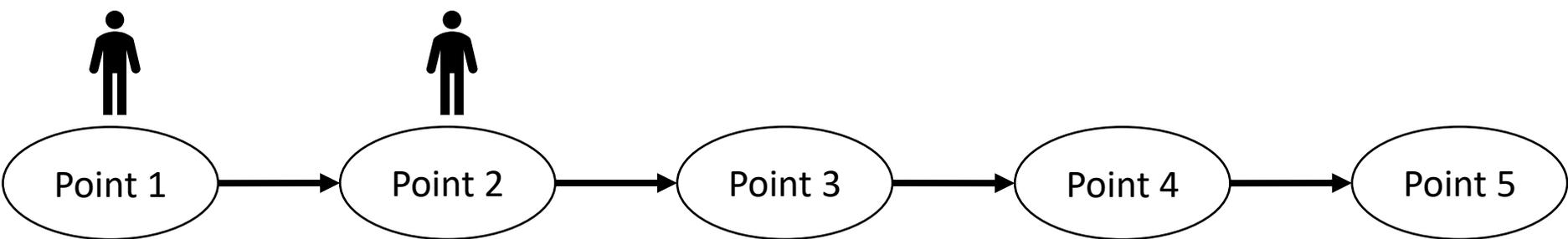
Given a street that goes from point 1 to point N, we are able to pick up customers on the street and take them to the other end of the street. Our taxi is only able to go one direction

Each customer has their starting point (pick-up spot), their ending point (destination), and the tip you will receive if you pick them up

The profit for selecting a customer is $(END - START) + TIP$

```
rides = [[1, 5, 1]      [2, 5, 4] ]  
        Start: Point 1  Start: Point 2  
        End: Point 5    End: Point 5  
        Tip: $1         Tip: $4
```

Is there another way we could represent this problem?

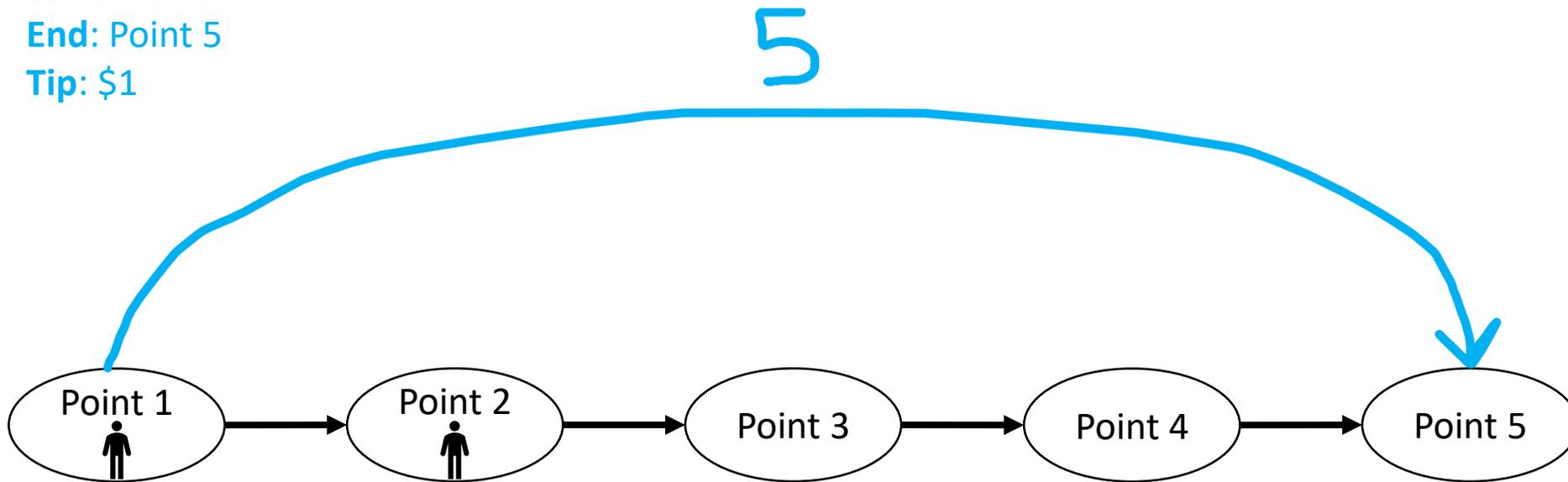


What is the **maximum profit** that the taxi can make when going from point 1 to point N?

Taxi Profit (Graph Representation)

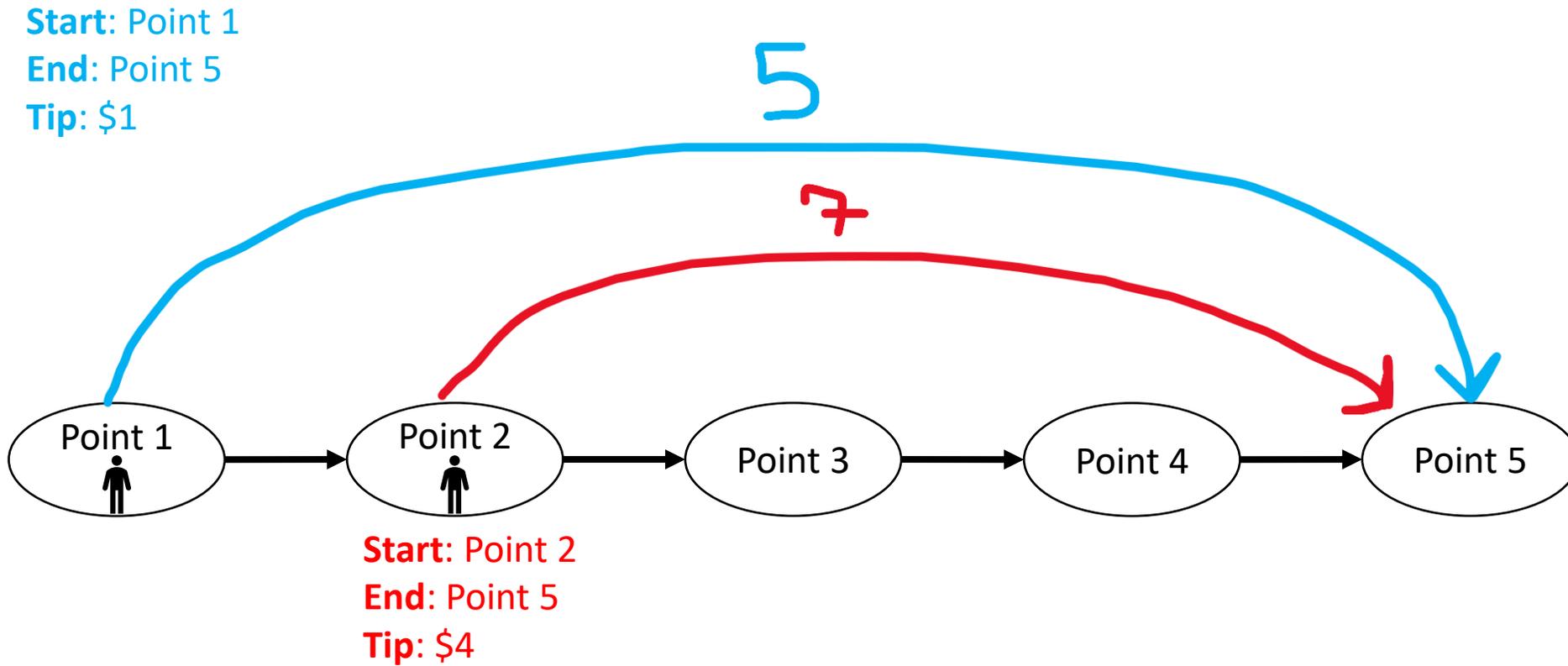
What is the **maximum profit** that the taxi can make when going from point 1 to point N?

Start: Point 1
End: Point 5
Tip: \$1



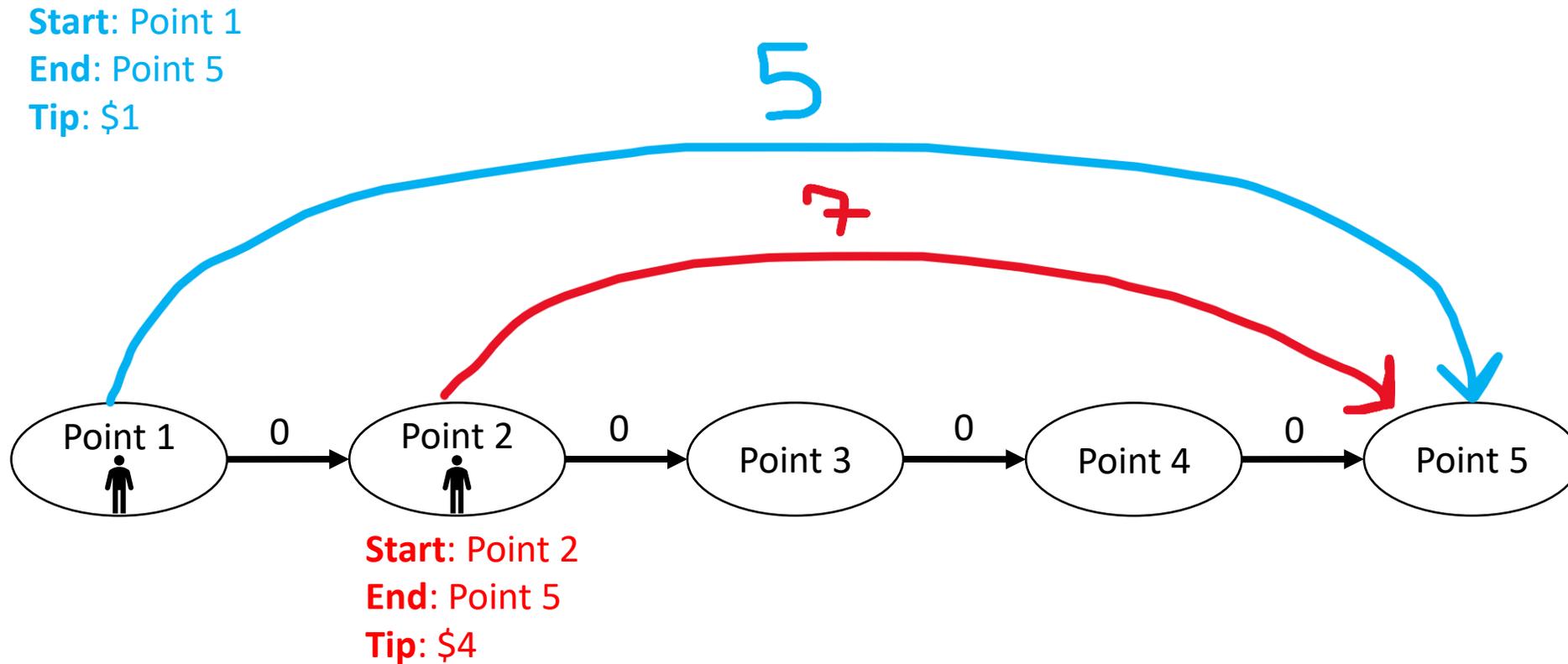
Taxi Profit (Graph Representation)

What is the **maximum profit** that the taxi can make when going from point 1 to point N?



Taxi Profit (Graph Representation)

What is the **maximum profit** that the taxi can make when going from point 1 to point N?



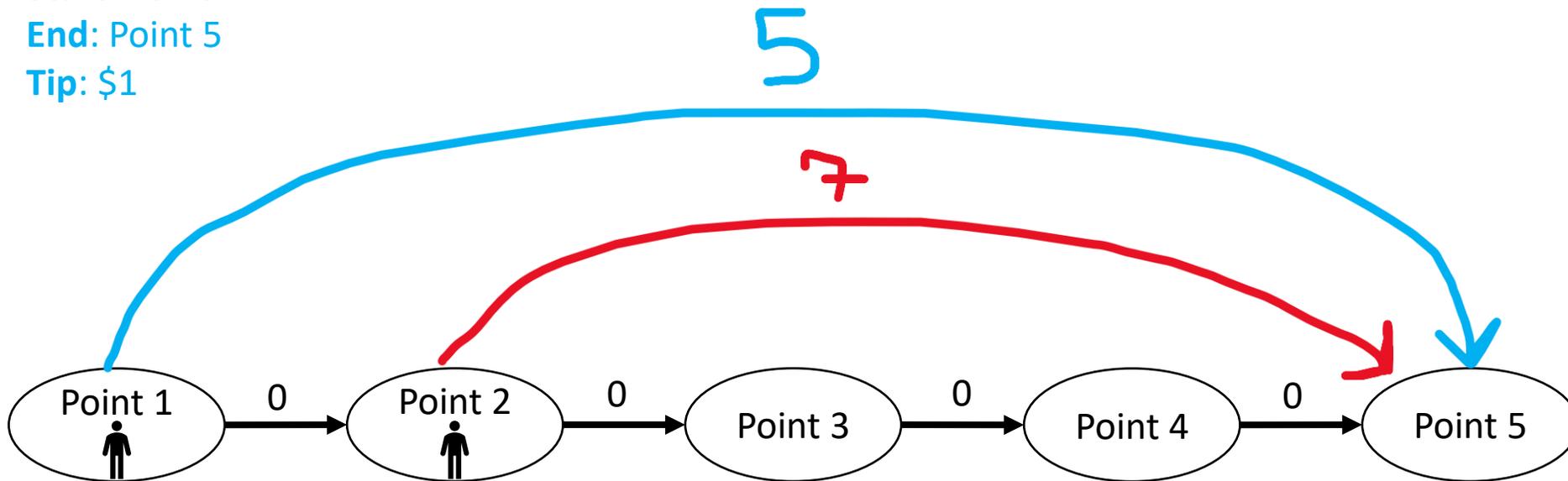
Taxi Profit (Graph Representation)

What is the **maximum profit** that the taxi can make when going from point 1 to point N?

What is the *longest* path from point 1 to point 5?

Directed Acyclic Graph (DAG)- a directed graph that contains no cycles

Start: Point 1
End: Point 5
Tip: \$1

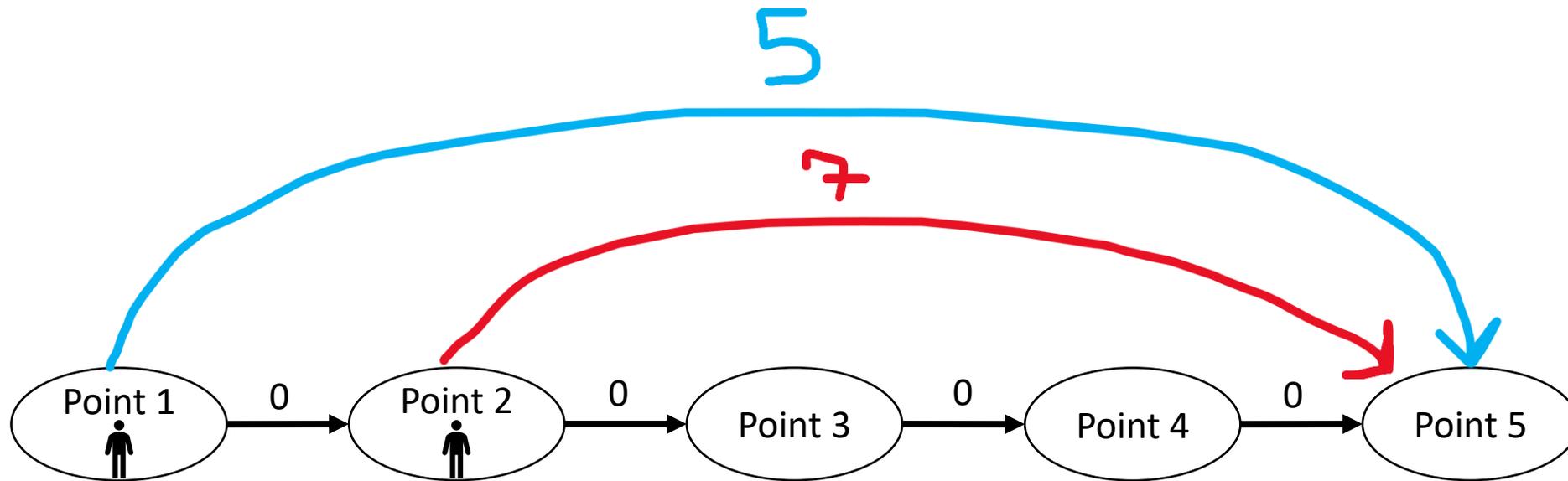


Start: Point 2
End: Point 5
Tip: \$4



Taxi Profit (Graph Representation)

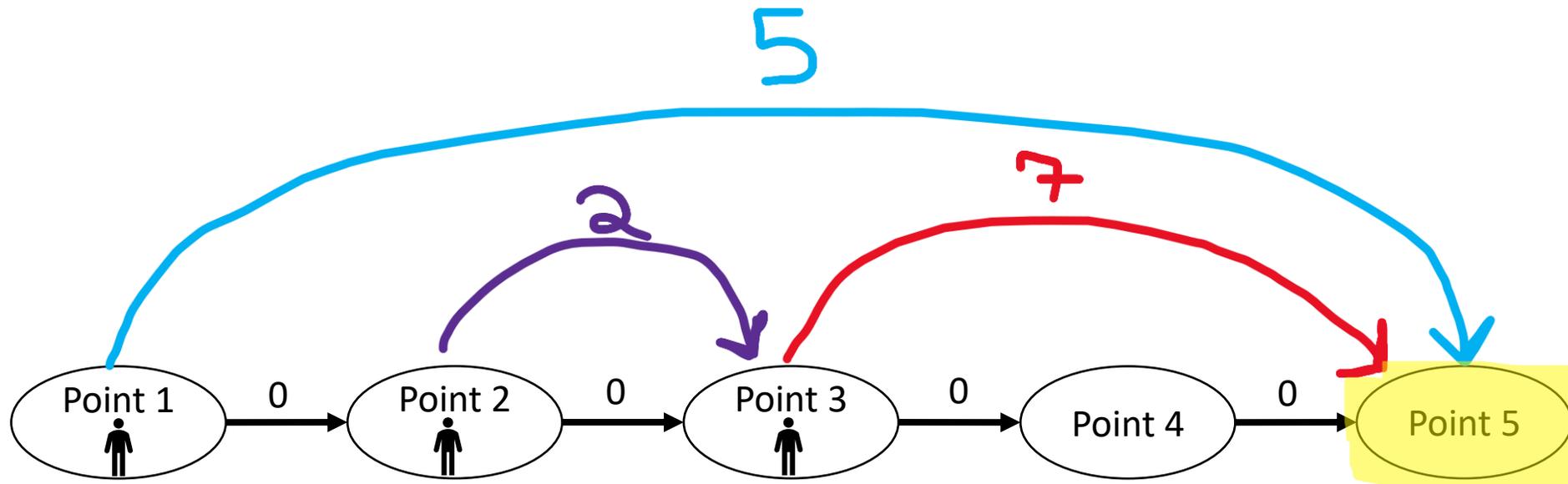
What is the **maximum profit** that the taxi can make when going from point 1 to point N?



Taxi Profit (Graph Representation)

What is the **maximum profit** that the taxi can make when going from point 1 to point N?

Let $B(x)$ be the maximum cost to go from point 1 to point x

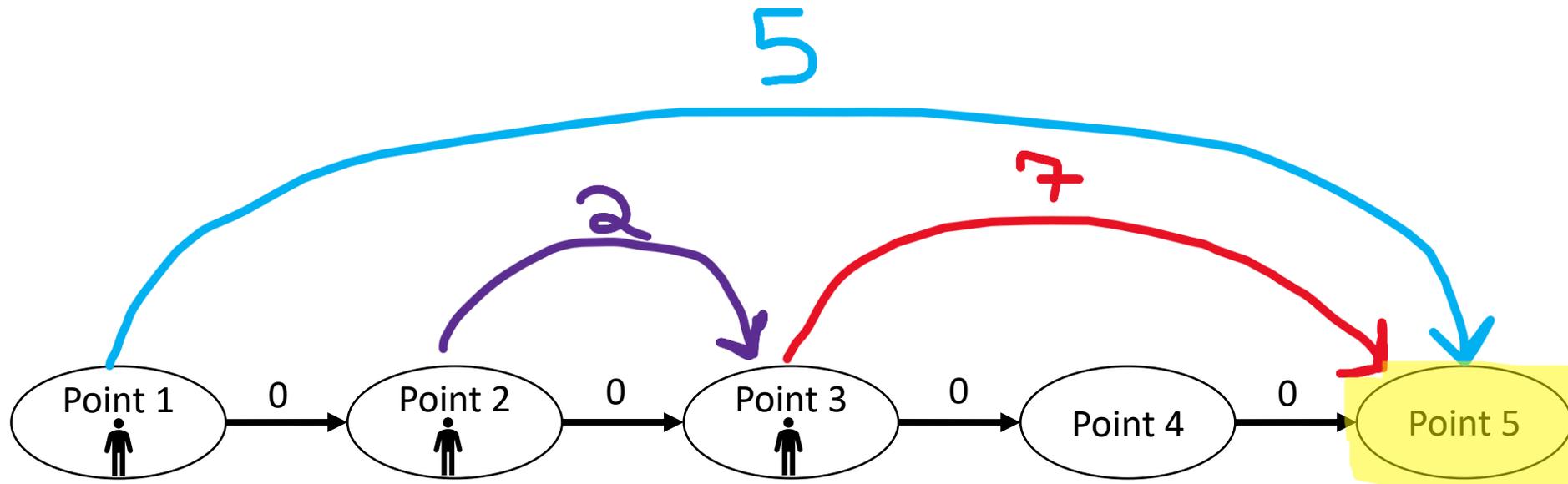


Taxi Profit (Graph Representation)

What is the **maximum profit** that the taxi can make when going from point 1 to point N?

Let $B(x)$ be the maximum cost to go from point 1 to point x

Let $P(i, x)$ be the profit made for a ride that goes from point 1 to point x



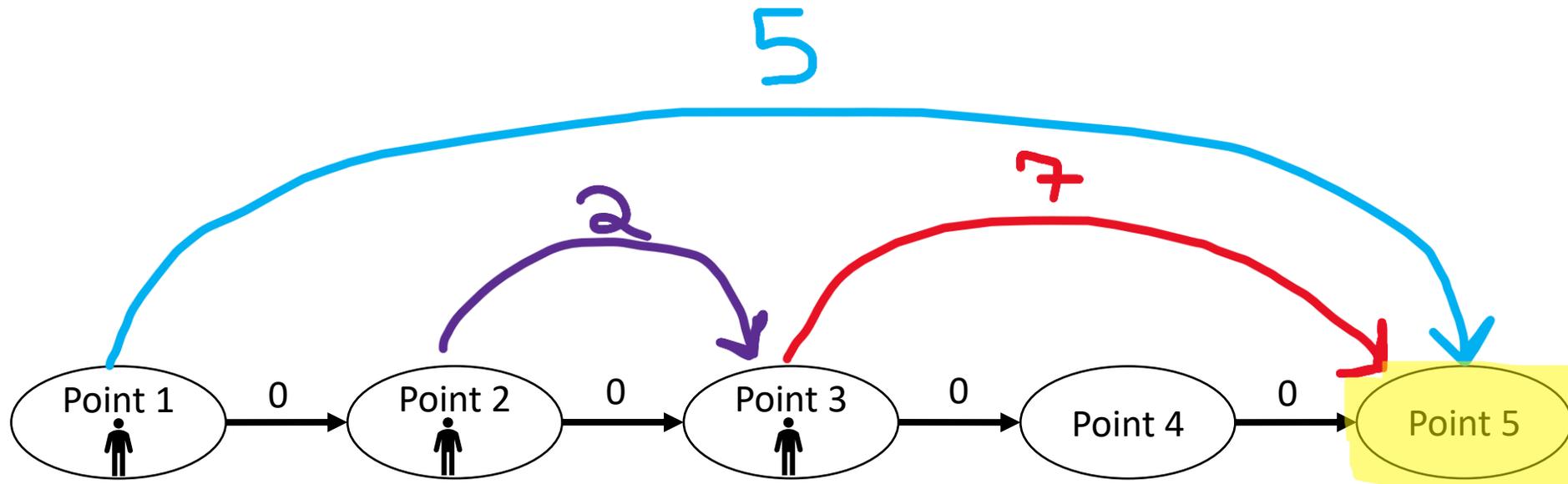
Taxi Profit (Graph Representation)

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$$B(5) = ?$$



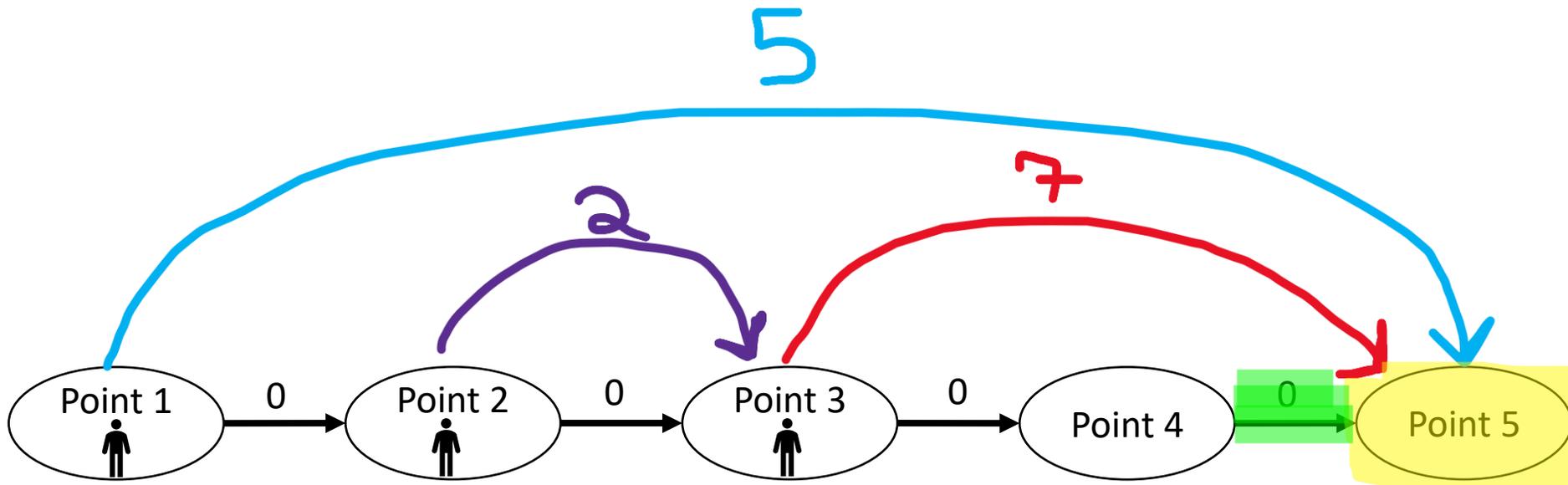
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$$B(5) = B(4) + 0$$



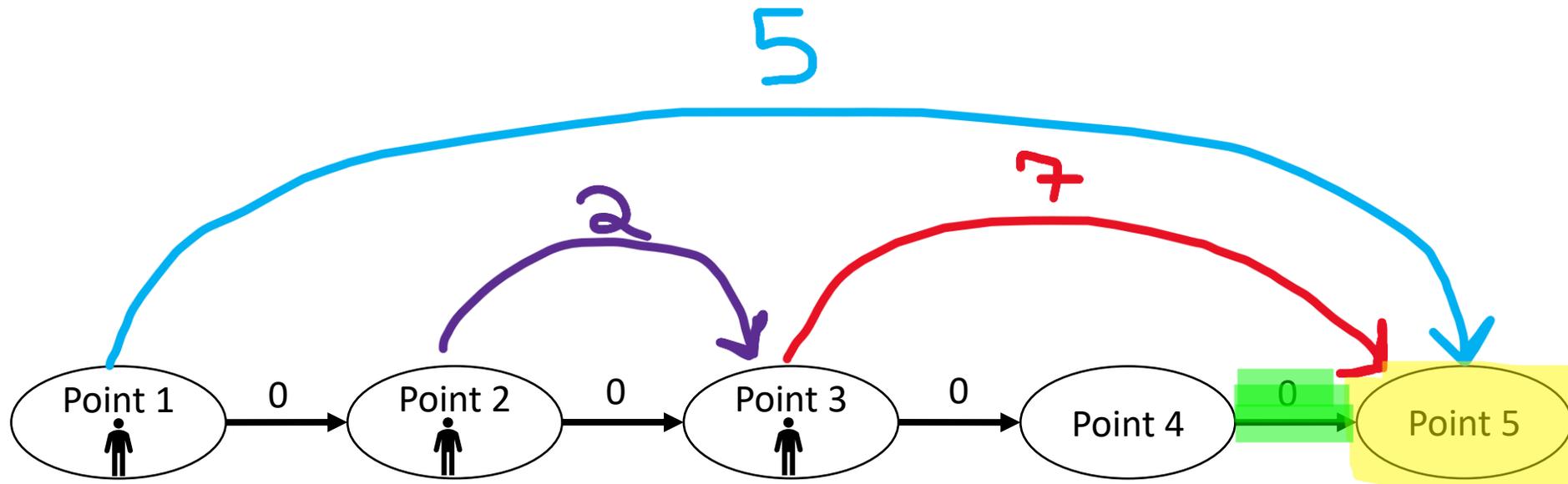
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$$B(5) = B(4) + 0 \text{ or } B(3) + 7 \text{ or } B(1) + 5$$



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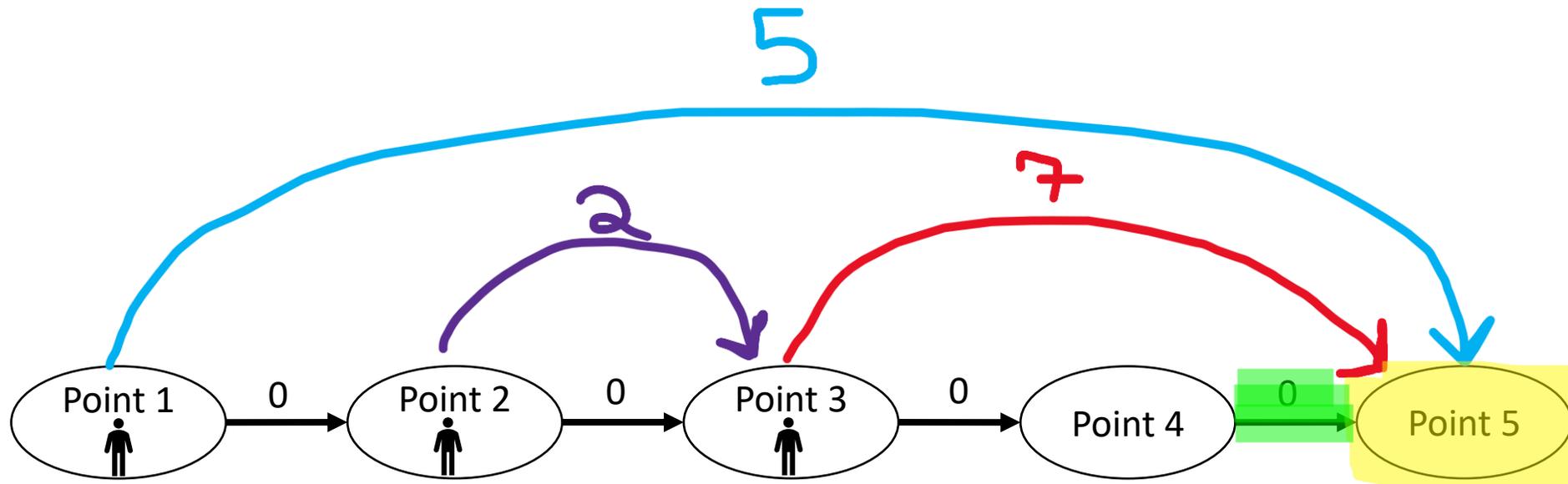
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Optimal Substructure



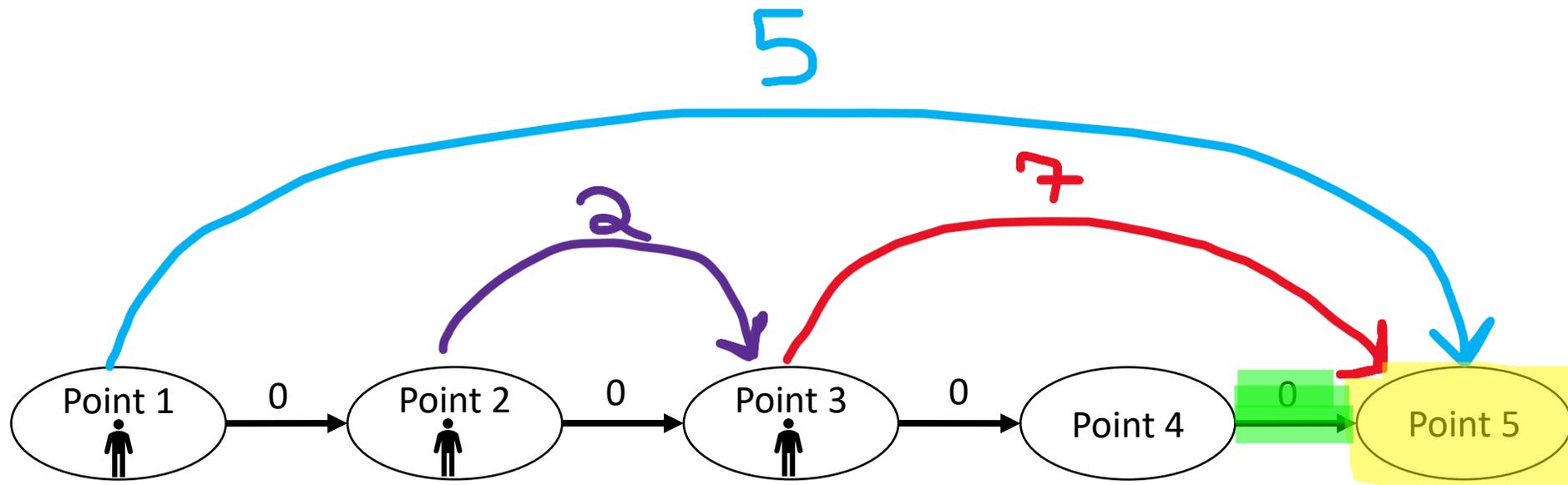
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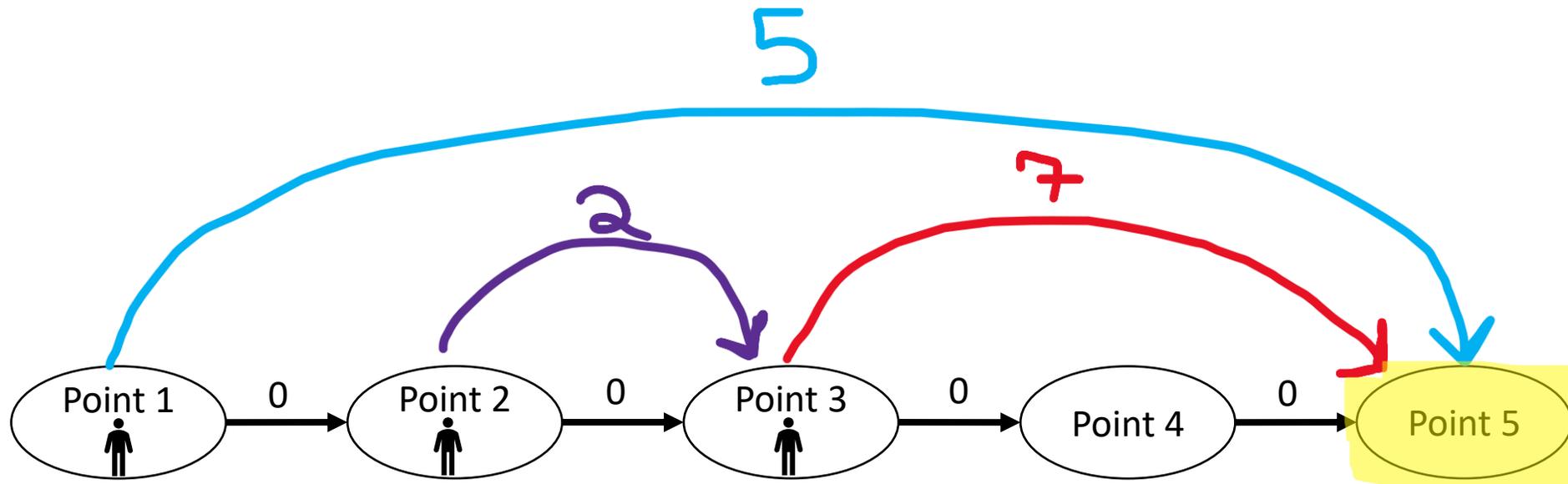
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$$B(x) = \max \begin{cases} B(x - 1) + 0 \\ B(x - i) + P(i, x) \\ \text{(for each edge that leads to point } x) \end{cases}$$



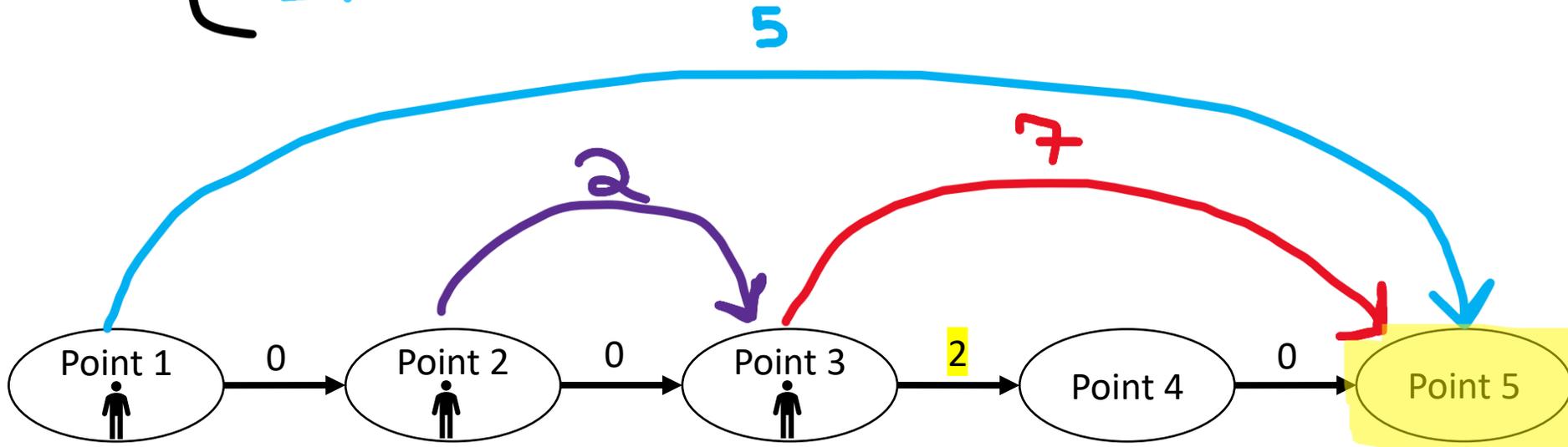
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$$B(x) = \max \begin{cases} B(4) + 0 = 2 \\ B(3) + 7 = 2 + 7 = 9 \\ B(1) + 5 = 5 \end{cases}$$



Taxi Profit (Graph Representation)

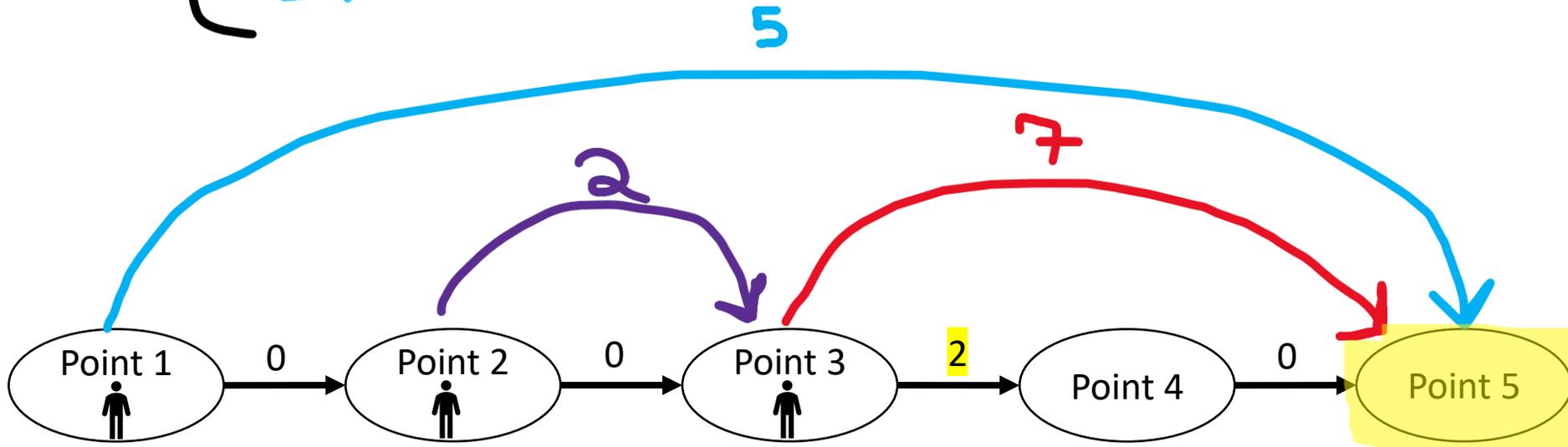
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Maximum profit = \$9



Taxi Profit (Graph Representation)

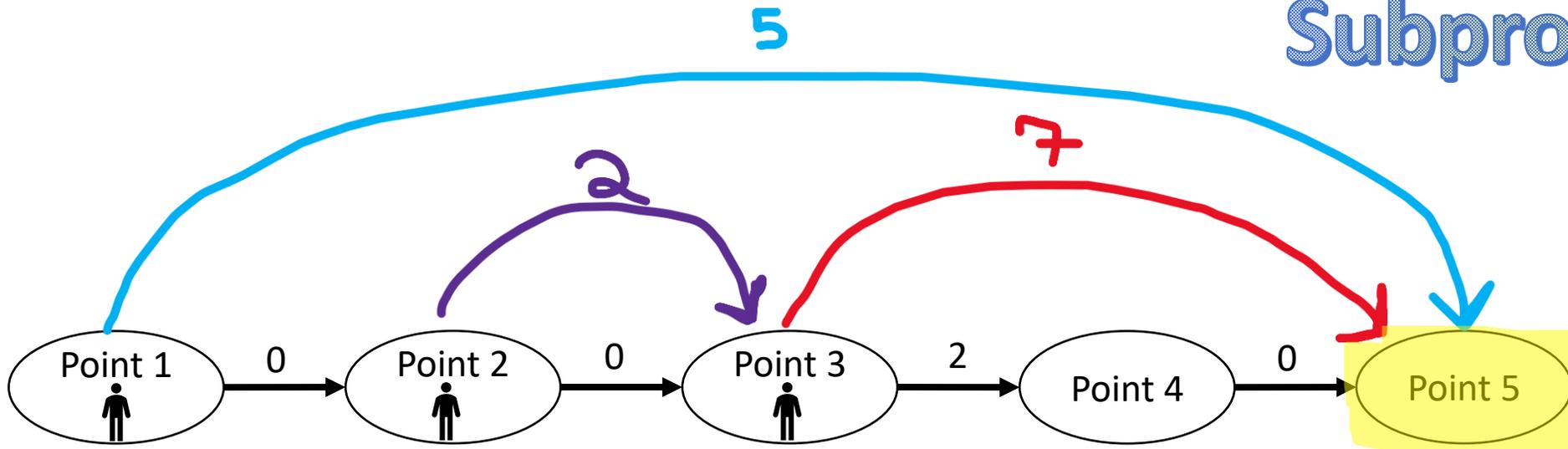
What is the **maximum profit** that the taxi can make when going from point 1 to point N?

To solve maximum profit from 1 to 5, we will solve :

- 1 to 4
- 1 to 3 (This may include *several paths* by picking up customers along the way ie 2 to 3)
- 1 to 2

2 to 3 may be a subproblem we solve several times

Overlapping Subproblems



Taxi Profit (G

ake when going from point 1 to point N?

To
•
•
•

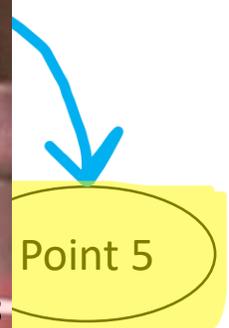
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verlapping
problems

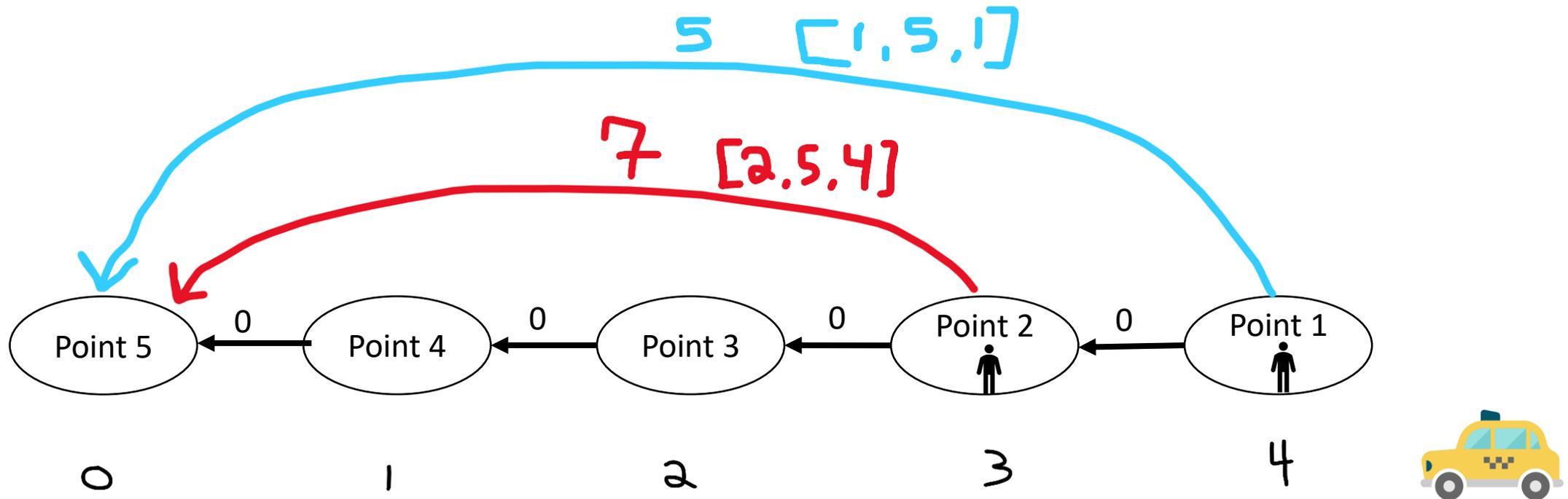


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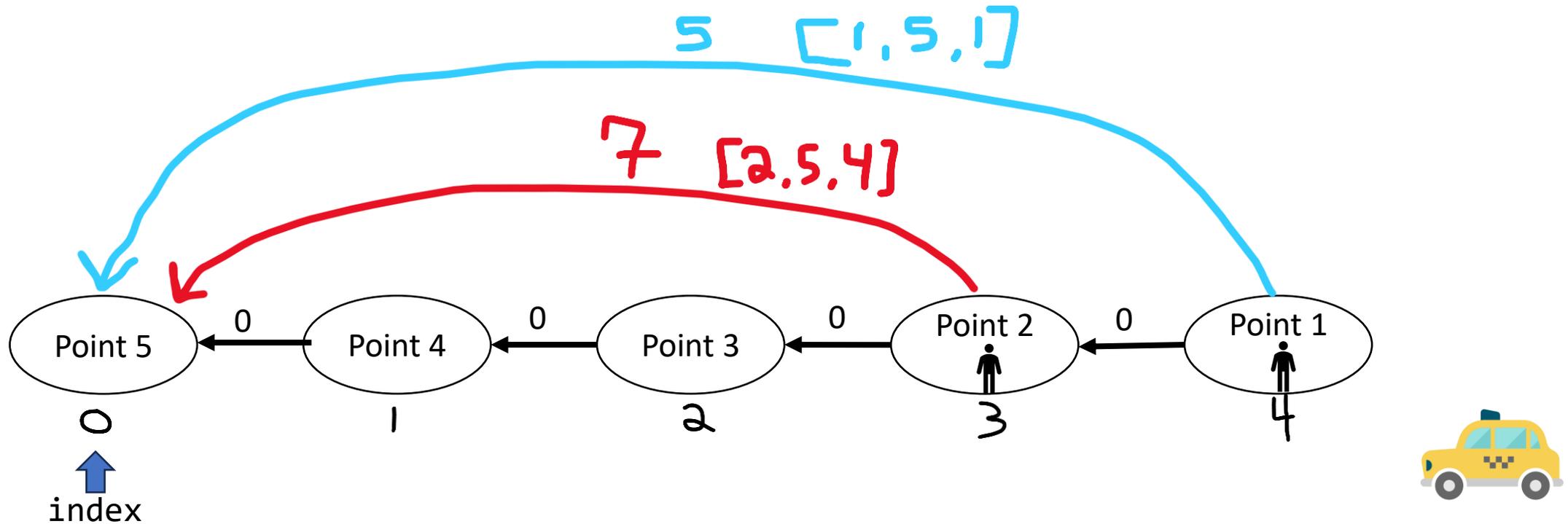


Taxi Profit

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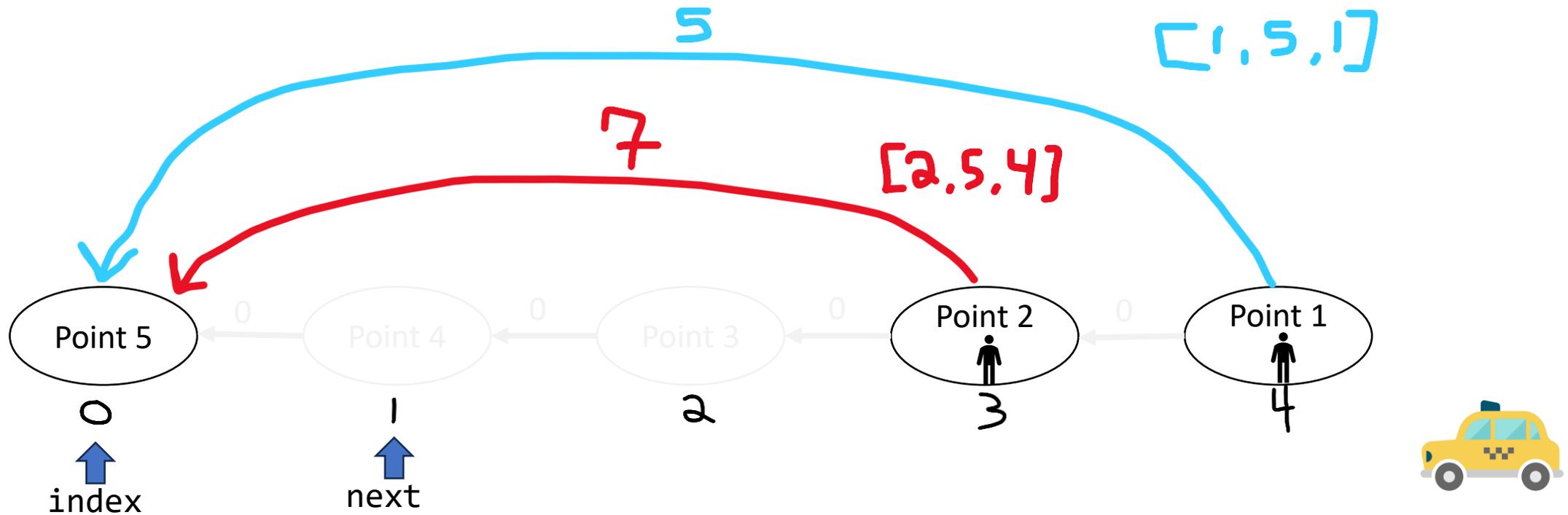
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To solve B(5) we must solve B(4). There isn't an entry that starts at point 4, so must find the next closest customer

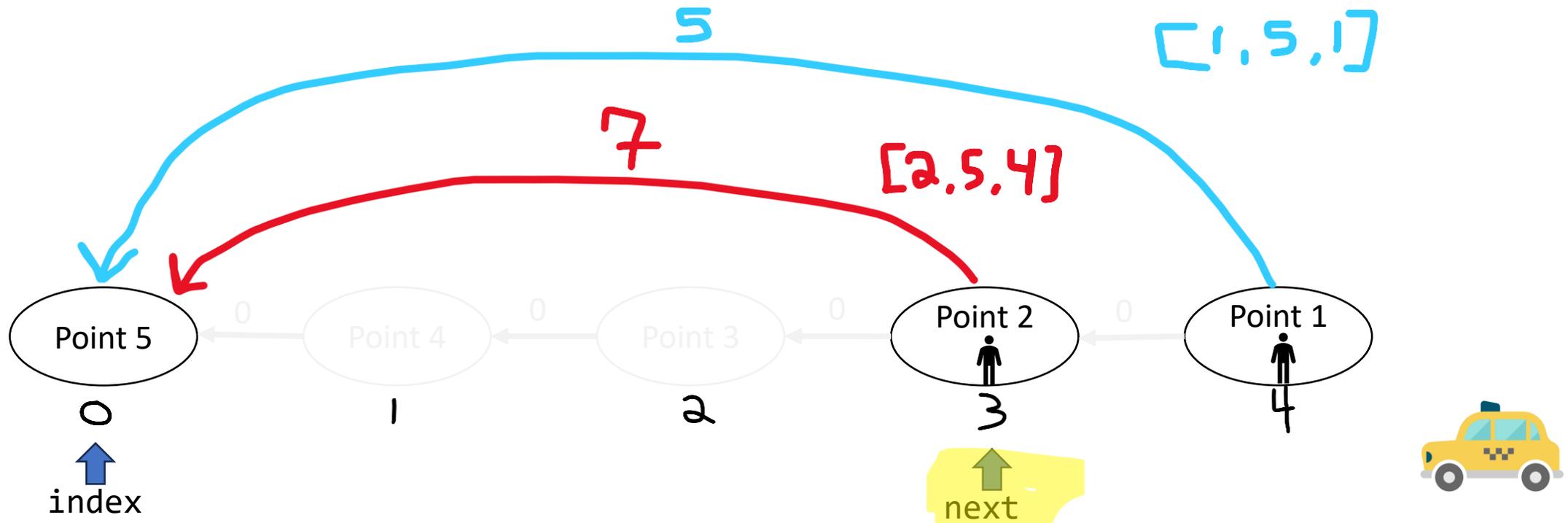
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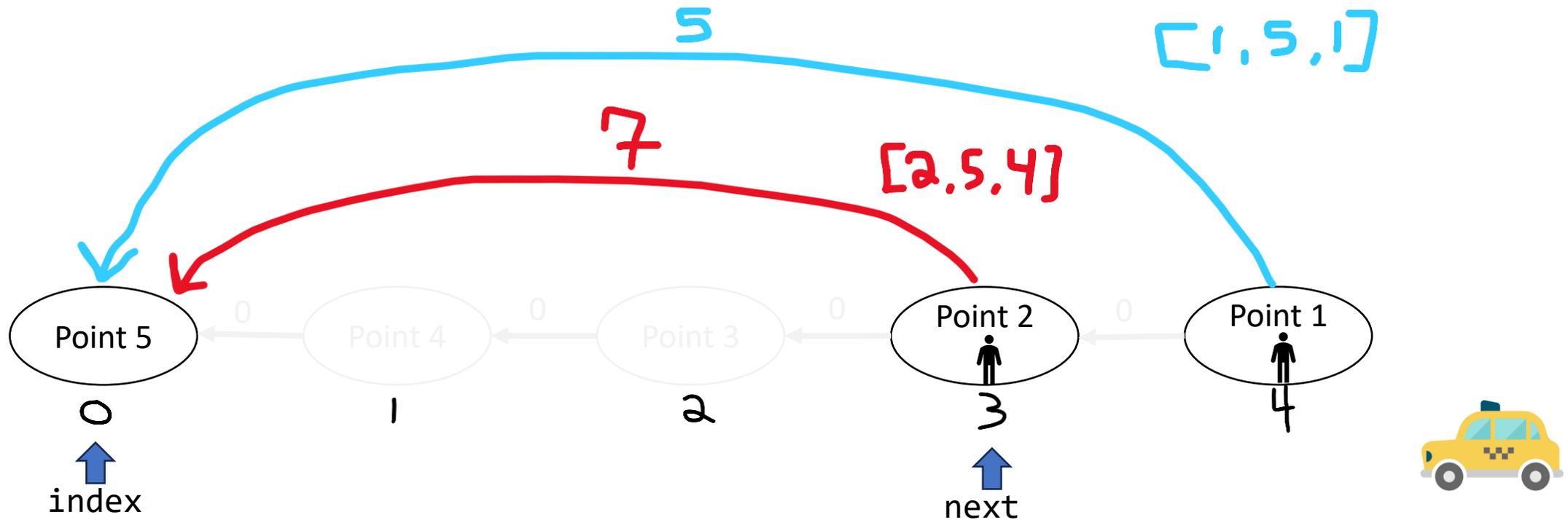
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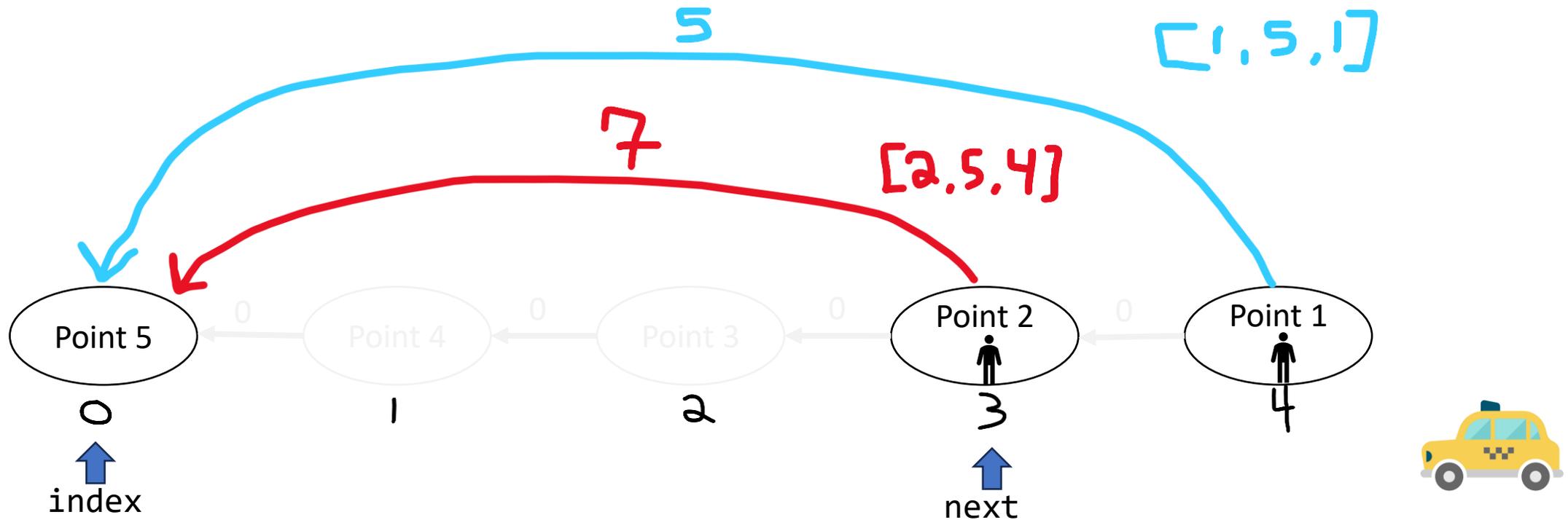
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Either we **take** this customer, or we **skip**

take: $(\text{ride.end} - \text{ride.start} + \text{ride.tip}) + B(2)$ ↖ Recursive call!

“We pick up the customer, and add their payment to our profit so far”



Taxi Profit

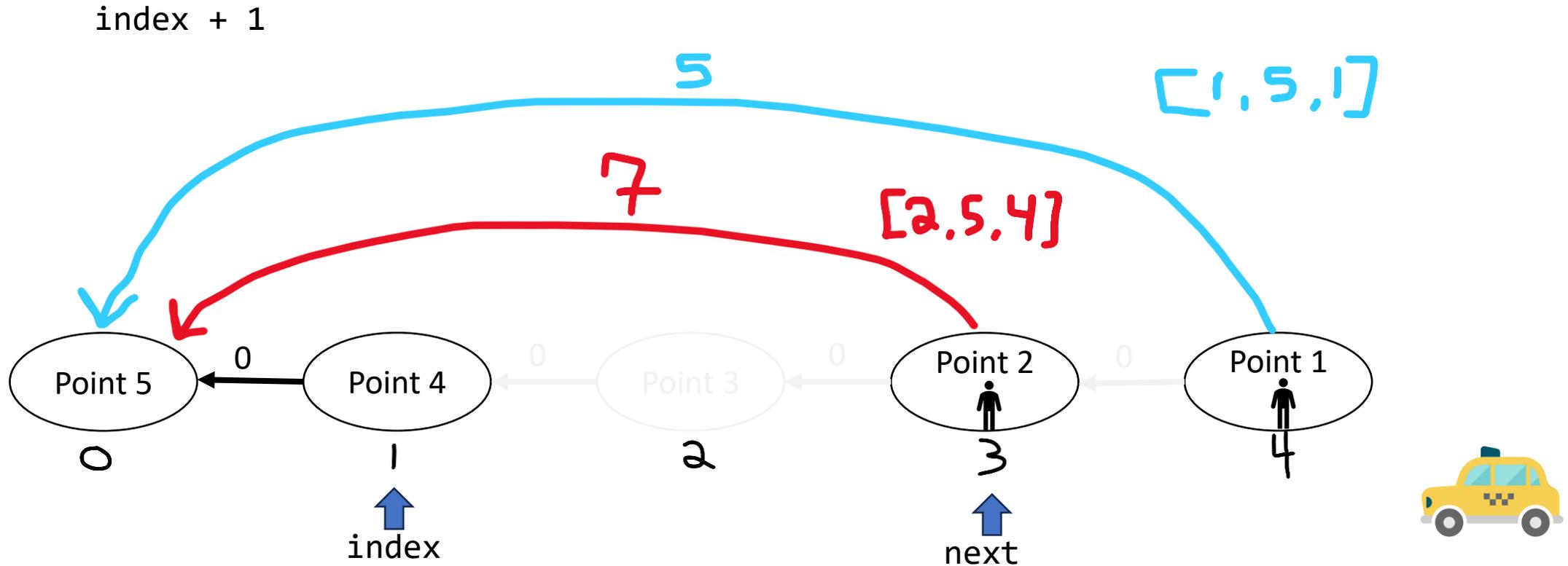
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Taxi Profit

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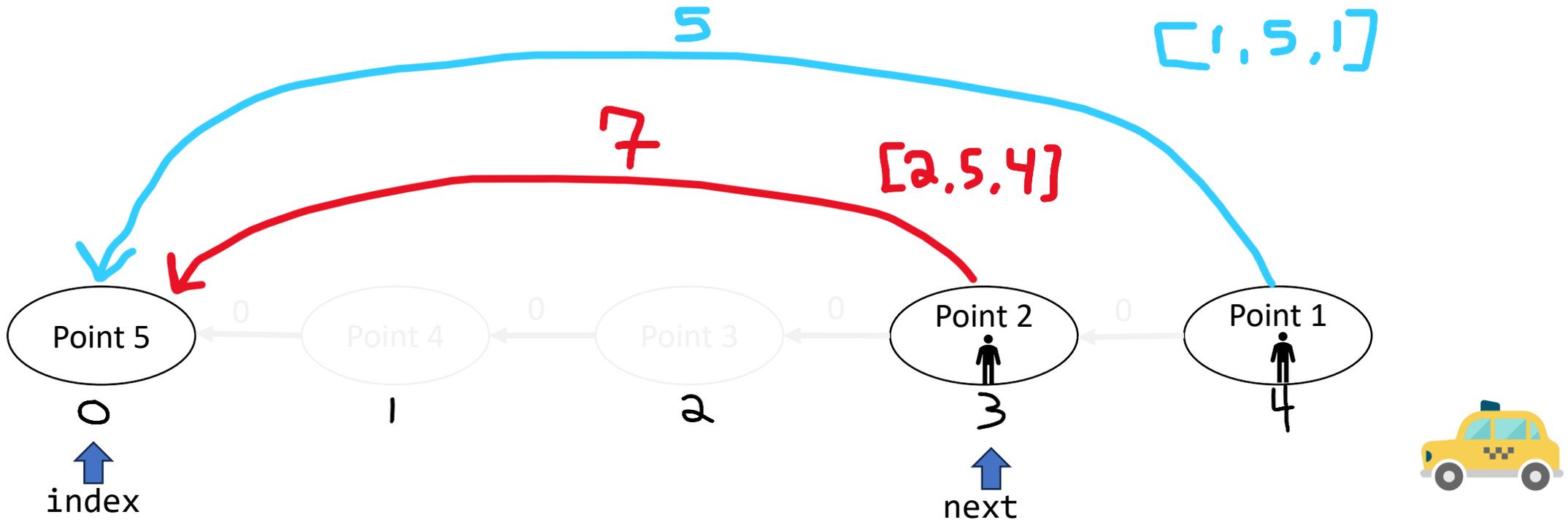
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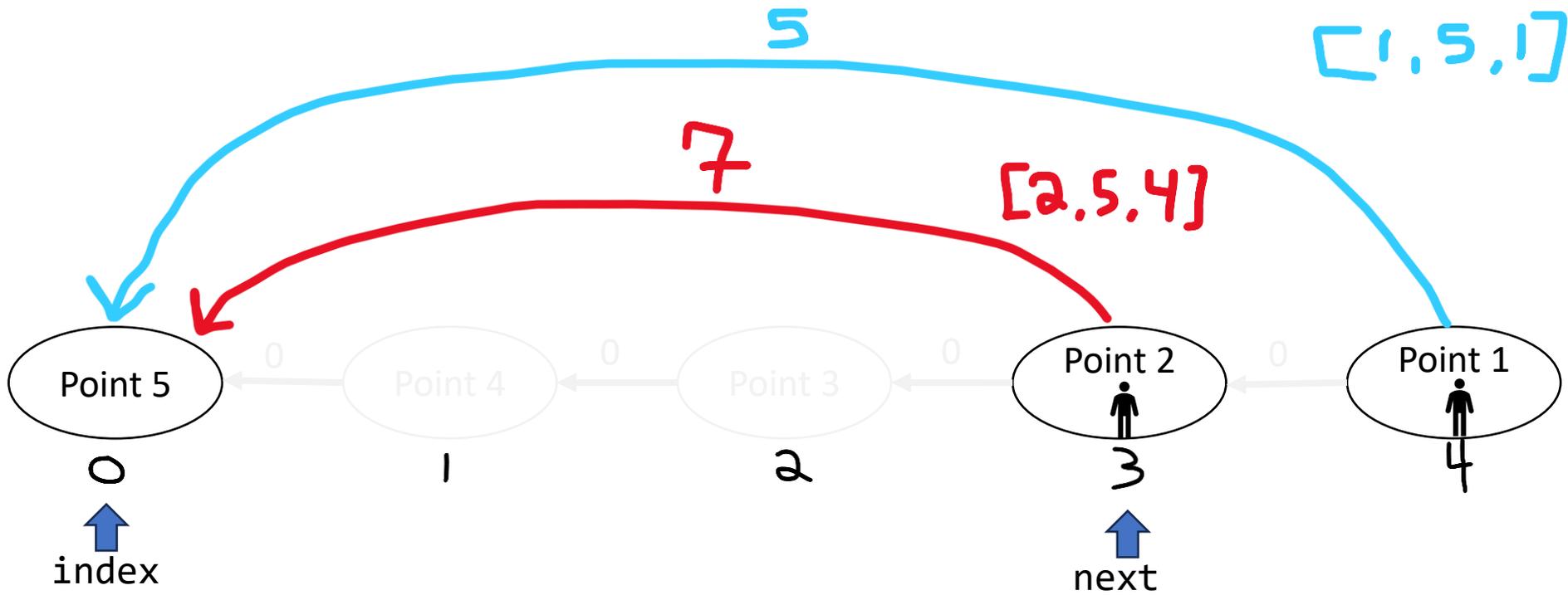
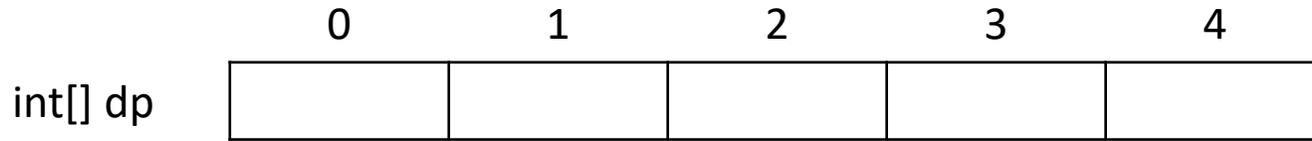
Take the Max choice, and store answer in memorization table



Taxi Profit

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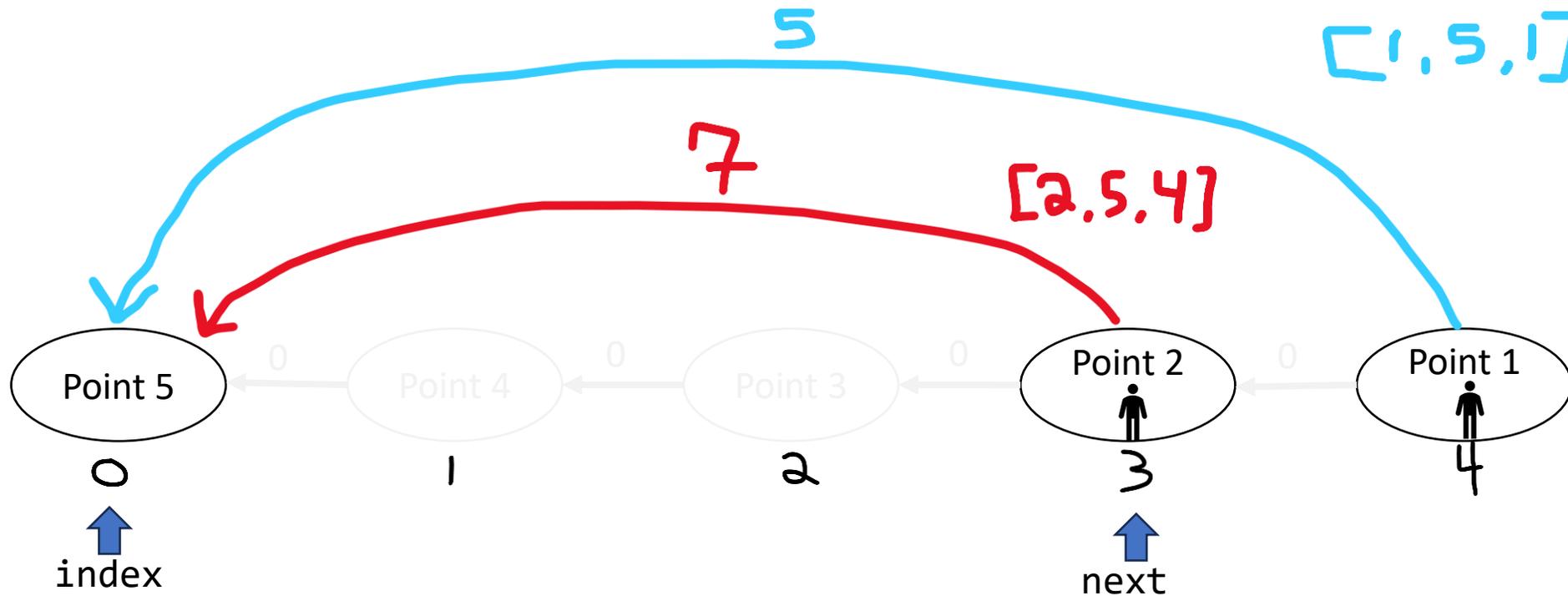
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| | | | | | |
|----------|---|---|---|---|---|
| | 0 | 1 | 2 | 3 | 4 |
| int[] dp | 7 | 0 | 0 | 0 | 0 |

↑ “The maximum profit to go from point 1 to point 5 (index 0) is \$7”



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