## **Pumping Lemma**

Given a regular language L,  $\exists$  a number p such that any string  $s \in L$ , with  $|s| \ge p$ , can be divided into three pieces, s = xyz satisfying:

- 1.  $xy^i z \in L, \forall i \ge 0.$
- 2. |y| > 0.
- 3.  $|xy| \leq p$ .

## **Proof Blueprint**

Claim: The language L = <some language> is not regular.

Proof: Suppose L is regular. Let p be the number from the pumping lemma.

Consider  $s = \langle \text{TODO} \rangle$ : Select s that will work with  $s \in L$  and  $|s| \ge p >$ . Since  $s \in L$  and  $|s| \ge p$ , the conditions of the pumping lemma must hold for s = xyz.

<TODO: Find conditions on what y must equal>

Consider the string  $s' = xy^{<\text{TODO: Select }i>}z = <\text{TODO: Show what }s' \text{ equals>}$ 

<TODO: Show s' is not in L>

 $\Rightarrow$  s' ∉ L, which is a contradiction of the pumping lemma. Therefore, the language is not regular.