# Finite Automata CSCI 338



DFAs either accept or reject strings.

# Deterministic Finite Automaton (DFA)

DFA string processing:

- 1. Start at start state.
- 2. Select first character in string.
- 3. Update state by following transition that corresponds to character.
- 4. Select next character in string.
- 5. Repeat step 3 and 4 until no characters remain.
- If final state is accept state, <u>accept</u>.
   Otherwise, <u>reject</u>.

DFAs either accept or reject strings.

Given string  $\omega = 01101$ , does this DFA accept or reject?



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- 2. Finite alphabet,  $\Sigma$ .

 $\Sigma$  consists of the transition characters (i.e. characters in the strings the DFA processes).



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- 3. Transition function,  $\delta: Q \times \Sigma \to Q$ .
- 4. Start state,  $q_0 \in Q$ .

Exactly one start state needed.



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- 2. Finite alphabet,  $\Sigma$ .
- 3. Transition function,  $\delta: Q \times \Sigma \to Q$ .
- 4. Start state,  $q_0 \in Q$ .
- 5. Set of accept states,  $\mathbf{F} \subseteq Q$ .

*F* is allowed to equal *Q* or be empty.



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- 4. Start state,  $q_0 \in Q$ .

5. Set of accept states,  $F \subseteq Q$ .

 $\begin{array}{c}
Q = \{q_1, q_2, q_3\} \\
\Sigma = \{0, 1\} \\
\delta: & 0 \quad 1 \\
\hline q_1 & q_1 & q_2 \\
q_2 & q_3 & q_2 \\
q_3 & q_2 & q_2
\end{array}$ 



Start state = 
$$q_1$$
  
 $F = \{q_2\}$ 

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The set of all strings A that a DFA M accepts is called its <u>language</u>, L(M) = A.

 $M \operatorname{recognizes} A.$ 

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 $L(M) = \begin{cases} \omega: \omega \text{ contains at least one 1 and an} \\ \text{even number of 0s following the final 1} \end{cases}$ 

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How do you prove a language is regular? Make a DFA that recognizes it.

Set of regular languages are "things we can do" with DFAs.

**Prove** that the following languages are regular:

1. Set of all strings over {0,1}.

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DFA Language Rules:

- 1. If the DFA accepts it, it is in the language.
- 2. If it is in the language, the DFA must accept it.

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 $\mathbf{q}_2$ 

3. Set of all strings that contain the substring: 10.



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0

 $\mathbf{q}_2$ 

 $\mathbf{q}_1$ 

 $\mathbf{q}_2$ 

0,1

**q**<sub>3</sub>

3. Set of all strings that contain the substring: 10.