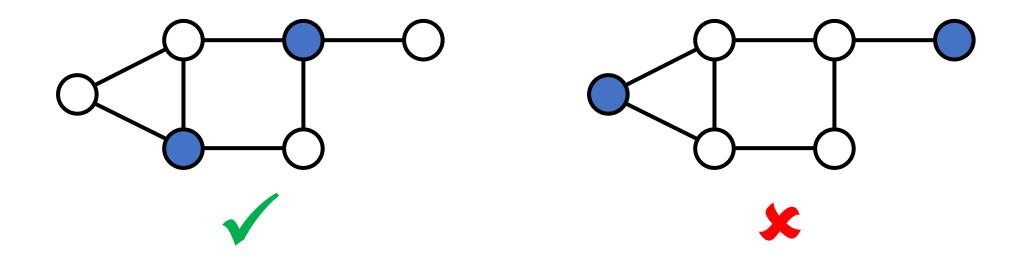
Dominating Set CSCI 338

Dominating Set: Given a graph G = (V, E) and integer  $k \leq |V|$ , is there a subset V' of size  $\leq k$ , such that every vertex  $\in V \setminus V'$  shares an edge with a vertex  $\in V'$ ?



Claim: Dominating Set  $\in NP - C$ 

Proof:

Dominating Set: For G = (V, E) and  $k \le |V|, \exists V' \subseteq V, |V'| \le k$ , s.t. each  $v \in V \setminus V'$  shares an edge with a  $u \in V'$ ?

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Proof:

1. Show Dominating Set  $\in NP$ .

2. Show  $A \leq_P$  Dominating Set, for some  $A \in NP - C$ .

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Claim: Dominating Set  $\in NP - C$ 

Proof:

1. Show Dominating Set  $\in NP$ .

Given G = (V, E), k and a subset V' of V, confirm that  $|V| \ge k$  and that for each vertex  $v \in V$ ,  $v \in V'$  or there is some edge  $(v, u) \in E$  such that  $u \in V'$ .

 $O(n^3)$  running time  $\Rightarrow$  Dominating Set  $\in NP$ .

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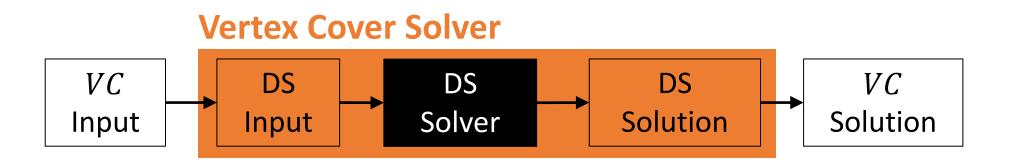
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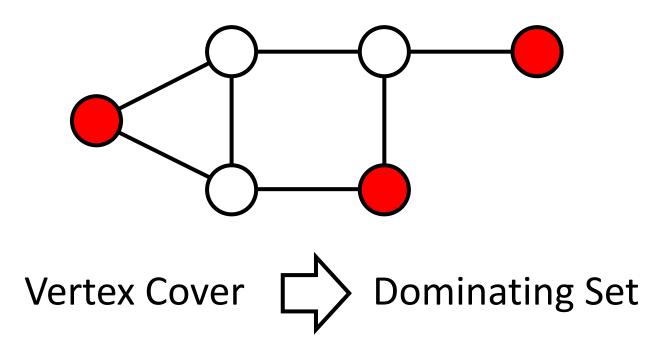
SAT? 3SAT? Clique? Vertex Cover?

Claim: Vertex Cover  $\leq_P$  Dominating Set Proof: Dominating Set: For G = (V, E) and  $k \le |V|, \exists V' \subseteq V, |V'| \le k$ , s.t. each  $v \in V \setminus V'$  shares an edge with a  $u \in V'$ ?

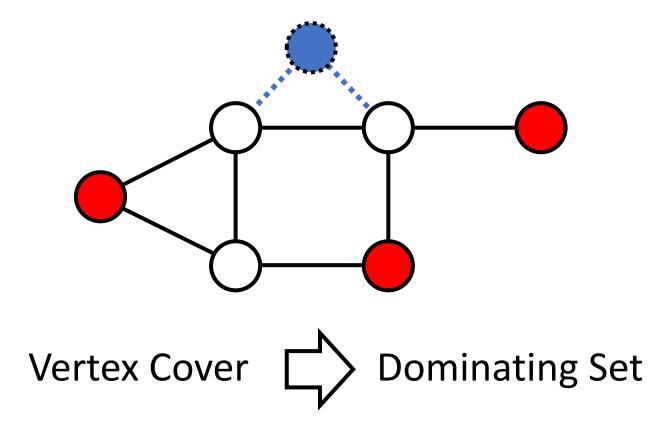


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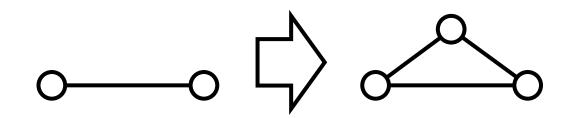
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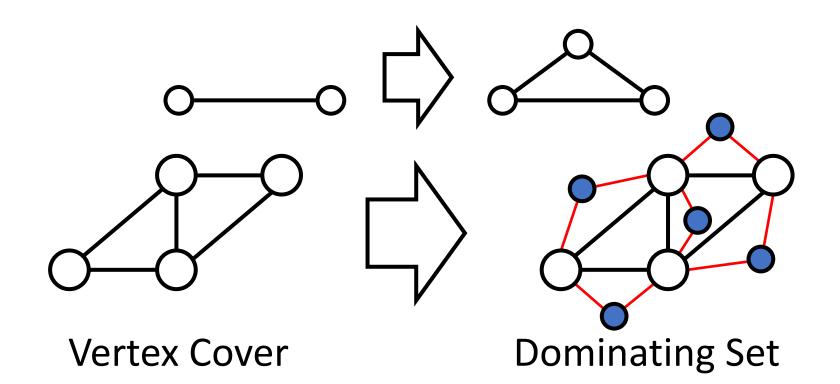
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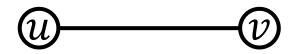
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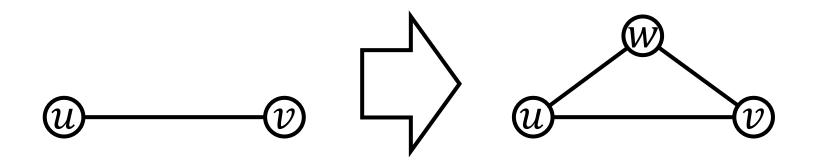
Proof: Let G = (V, E), k be input to the vertex cover problem, where |V| = n. Create G' = (V', E') as follows: For each  $e = (u, v) \in E$ , add u and v to V' and add e to E'.



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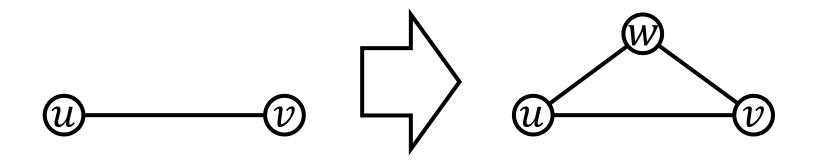
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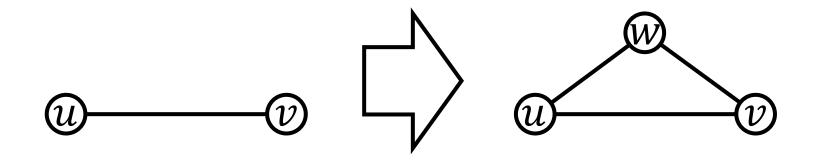
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*G* has a k-VC  $\Leftrightarrow$  *G*' has a k-DS.