NFA/DFA Equivalence CSCI 338



0,1

S₃

 S_2

0(



Claim: Every NFA has an equivalent DFA.

Proof Approach: How did we keep track of our location in a DFA?







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 \mathbf{q}_1

 \mathbf{q}_2

 q_1

 \mathbf{q}_2

8

 \mathbf{q}_3

 q_4

0,1

 q_1

Claim: Every NFA has an equivalent DFA.

Proof Approach: How did we keep track of our location in an NFA? Set of all states we could possibly be in.

1*,ɛ*

 \mathbf{q}_2

 \mathbf{q}_1

 \mathbf{q}_2

q₂

101

 \mathbf{q}_2

2

8

 \mathbf{q}_3

 q_4

 q_4

 \mathbf{q}_1

 q_1

q₂

 q_1

ω

0,1

 \mathbf{q}_4

0

q₃

Claim: Every NFA has an equivalent DFA.

Proof Approach: How did we keep track of our location in an NFA? Set of all states we could possibly be in.

What is the set of all possible locations? Power set of states! (set of all subsets)







DFA states = Power set of NFA states.



DFA states = Power set of NFA states. Start state = ?



DFA states = Power set of NFA states. Start state = NFA's start state.



DFA states = Power set of NFA states. Start state = NFA's start state. Accept states = ?



DFA states = Power set of NFA states. Start state = NFA's start state. Accept states = Any state that includes accept state from NFA.



Where should transition out of {S₁} with character 1 go?



Where should transition out of $\{S_1\}$ with character 1 go? Wherever S_1 goes with 1 in the NFA.



Where should transition out of $\{S_1\}$ with character 1 go? Wherever S_1 goes with 1 in the NFA.



Where should transition out of {S₁} with character 0 go?



Where should transition out of $\{S_1\}$ with character 0 go? If transition is not handled by NFA, send it to \emptyset (junk state).



Where should transition out of {S₂} with character 0 go?



Where should transition out of $\{S_2\}$ with character 0 go? NFA could stay in S_2 or go to S_3 , so $\{S_2, S_3\}$



Where should transition out of $\{S_1, S_3\}$ with character 1 go?



Where should transition out of $\{S_1, S_3\}$ with character 1 go?



Where should transition out of $\{S_1, S_3\}$ with character 0 go?



Where should transition out of $\{S_1, S_3\}$ with character 0 go?



Rule?

DFA state transitions to DFA state consisting of all states it's NFA states transition to.



Rule? For each DFA state R and $e \in \Sigma$, transition $(R, e) = \{q \in NFA: q \in transition(r, e) \text{ for some } r \in R\}$

DFA vs NFA





What about ε -transitions? Define extension of DFA state R: $E(R) = \{q \in NFA: q \text{ reachable from } r \in R \text{ with } \ge 0 \varepsilon$ -transitions}



 $E(R) = \{q \in NFA: q \text{ reachable from } r \in R \text{ with } \ge 0 \epsilon \text{-transitions} \}$ Example: $E(\{S_2, S_3\}) = ?$



 $E(R) = \{q \in \mathsf{NFA}: q \text{ reachable from } r \in R \text{ with } \ge 0 \text{ } \varepsilon \text{-transitions} \}$ Example: $E(\{S_2, S_3\}) = \{S_2, S_3\}$ $E(\{S_1, S_2\}) = ?$



 $E(R) = \{q \in \mathsf{NFA}: q \text{ reachable from } r \in R \text{ with } \ge 0 \text{ } \varepsilon \text{-transitions} \}$ Example: $E(\{S_2, S_3\}) = \{S_2, S_3\}$ $E(\{S_1, S_2\}) = \{S_1, S_2, S_3\}$



Make start state = ?



Make start state = $E(\{S_1\}) = \{S_1, S_3\}$











transition(R, e) = { $q \in NFA$: $q \in \frac{\text{transition}(r, e)}{\text{for some } r \in R}$





DFA vs NFA





Suppose *w* accepted by NFA.



Suppose *w* accepted by NFA. At each step of its processing, DFA will be in state that corresponds to all possible NFA states.



Equivalent?

Suppose *w* accepted by NFA. At each step of its processing, DFA will be in state that corresponds to all possible NFA states. If NFA ends on accept state, corresponding DFA state will accept too.



Suppose *w* accepted by DFA.



Suppose *w* accepted by DFA. At each step of its processing, DFA will be in state that corresponds to all possible NFA states.



Equivalent?

Suppose *w* accepted by DFA. At each step of its processing, DFA will be in state that corresponds to all possible NFA states. If DFA ends on accept state, it includes an NFA accept state.





A language is called a <u>regular language</u> if some DFA **or NFA** recognizes it.

How do you prove a language is regular? Make a DFA **or NFA** that recognizes it.