Pumping Lemma
Given a regular language $L$, $\exists$ a number $p$ such that any string $s \in L$, with $|s| \geq p$, can be divided into three pieces, $s = xyz$ satisfying:
1. $xy^iz \in L, \ \forall i \geq 0$.
2. $|y| > 0$.
3. $|xy| \leq p$.

Proof Blueprint
Claim: The language $L = \langle$some language$\rangle$ is not regular.

Proof: Suppose $L$ is regular. Let $p$ be the number from the pumping lemma.
Consider $s = <$TODO: Select $s$ that will work with $s \in L$ and $|s| \geq p$$.>
Since $s \in L$ and $|s| \geq p$, the conditions of the pumping lemma must hold for $s = xyz$.

$\langle$TODO: Find conditions on what $y$ must equal$\rangle$

Consider the string $s' = xy^iz = <$TODO: Select $i$$.>
$\langle$TODO: Show what $s'$ equals$\rangle$
$\langle$TODO: Show $s'$ is not in $L$$\rangle$

$\Rightarrow s' \notin L$, which is a contradiction of the pumping lemma.
Therefore, the language is not regular.